

AD-786 066

MULTI-VALUED CROSS-CORRELATION  
FUNCTIONS BETWEEN TWO MAXIMAL LINEAR  
RECURSIVE SEQUENCES

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Prepared for:

Army Research Office

January 1972

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## DOCUMENT CONTROL DATA - R &amp; D

(Security classification of title, body of abstract and indexing annotation must be entered when the overall report is classified)

|                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                     |                                                                             |                                                                                                      |  |
|-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|-----------------------------------------------------------------------------|------------------------------------------------------------------------------------------------------|--|
| 1. ORIGINATING ACTIVITY (Corporate author)<br>Electronic Sciences Laboratory<br>University of Southern California<br>Los Angeles, California 90007                                                                                                                                                                                                                                                                                                                                                                                                                                  |                                                                             | 2a. REPORT SECURITY CLASSIFICATION<br>UNCLASSIFIED                                                   |  |
|                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                     |                                                                             | 2b. GROUP                                                                                            |  |
| 3. REPORT TITLE<br>MULTI-VALUED CROSS-CORRELATION FUNCTIONS BETWEEN TWO MAXIMAL LINEAR RECURSIVE SEQUENCES                                                                                                                                                                                                                                                                                                                                                                                                                                                                          |                                                                             |                                                                                                      |  |
| 4. DESCRIPTIVE NOTES (Type of report and inclusive dates)<br>Scientific Interim                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                     |                                                                             |                                                                                                      |  |
| 5. AUTHOR(S) (First name, middle initial, last name)<br>Yoji Niho                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                   |                                                                             |                                                                                                      |  |
| 6. REPORT DATE<br>January 1972                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                      | 7a. TOTAL NO. OF PAGES<br>115                                               | 7b. NO. OF REFS<br>22                                                                                |  |
| 8a. CONTRACT OR GRANT NO.<br>DA ARO-D-31-124-70-G104<br>DA ARO-D-31-124-70-G930                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                     | 8b. ORIGINATOR'S REPORT NUMBER(S)<br>USCEE Report 409                       |                                                                                                      |  |
| 8c. PROJECT NO.<br>7193-RT, 7198-RT                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                 | 8d. OTHER REPORT NO(S) (Any other numbers that may be assigned this report) |                                                                                                      |  |
| 10. DISTRIBUTION STATEMENT<br>This document has been approved for public release and sale; its distribution is unlimited.                                                                                                                                                                                                                                                                                                                                                                                                                                                           |                                                                             |                                                                                                      |  |
| 11. SUPPLEMENTARY NOTES                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                             |                                                                             | 12. SPONSORING MILITARY ACTIVITY<br>The U.S. Army Research Office-Durham,<br>Durham, North Carolina. |  |
| 13. ABSTRACT<br><br>The cross-correlation function, $\Delta_r(y)$ , between two maximal linear recursive sequences is defined by<br>$\Delta_r(y) = \sum_{x \in GF(2^n)} (-1)^{\text{Tr}(xy + x^r)}$<br>for some $r$ , $\text{GCD}(r, 2^n - 1) = 1$ . $\Delta_r(y)$ is analyzed and evaluated for two types of decimation $r$ . For the first type, $r \equiv 2^k \pmod{2^{n/2} - 1}$ . It is shown that $\Delta_r(y)$ is restricted to the form<br>$\Delta_r(y) = 2^{n/2}(j-1), \quad 0 \leq j \leq J,$<br>where $j$ is the number of distinct solutions to the system of two equa- |                                                                             |                                                                                                      |  |

DD FORM 1473  
NOV 66

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Item 13(Cont)

tions over  $GF(2^n)$  and  $J$  is the degree of one of the two equations. For the second type,  $r = (2^{mk}+1)/(2^k+1)$  and for this case  $\Delta_r(y)$  is restricted to the form

$$\Delta_r(y) = 0, \quad \pm 2^{(n+de)/2}, \quad 0 \leq d \leq m-2, \quad d \text{ odd}$$

where  $e = \text{GCD}(n,k)$  and  $d$  depends on the rank of a quadratic form over  $GF(2^e)$ .

The explicit evaluation of  $\Delta_r(y)$  is equivalent to the explicit evaluation of weight distribution of the  $(2^n-1, 2n)$  cyclic code whose dual code is generated by  $f_1(x)f_r(x)$ , the product of two primitive polynomials of degree  $n$ .

MULTI-VALUED CROSS-CORRELATION FUNCTIONS BETWEEN TWO  
MAXIMAL LINEAR RECURSIVE SEQUENCES

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January 1972

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A dissertation presented to the Graduate School, University of Southern California in partial fulfillment of the requirements for the degree of Doctor of Philosophy (Electrical Engineering).

This work was supported in part by the U.S. Army Research Office-Durham under Grant DA-ARO-D-31-124-70-G930 (G1044, G51, G104).

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## ACKNOWLEDGEMENTS

The author wishes to express his most profound gratitude and appreciation to Dr. Lloyd R. Welch, the Chairman of his Dissertation Committee, for his continual guidance, assistance and encouragement. Numerous ideas on evaluation of cross-correlation functions were suggested by Dr. Welch. The author also wishes to express his gratitude and appreciation to Dr. Solomon W. Golomb and Dr. Albert L. Whiteman who have served on the Dissertation Committee.

# ABSTRACT

The cross-correlation function,  $\Delta_r(y)$ , between two maximal linear recursive sequences is defined by

$$\Delta_r(y) = \sum_{x \in GF(2^n)} (-1)^{\text{Tr}(xy + x^r)}$$

for some  $r$ ,  $\text{GCD}(r, 2^n - 1) = 1$ .  $\Delta_r(y)$  is analyzed and evaluated for two types of decimation  $r$ . For the first type,  $r \equiv 2^k \pmod{2^n - 1}$ . It is shown that  $\Delta_r(y)$  is restricted to the form

$$\Delta_r(y) = 2^{n/2(j-1)}, \quad 0 \leq j \leq J,$$

where  $j$  is the number of distinct solutions to the system of two equations over  $GF(2^n)$  and  $J$  is the degree of one of the two equations. For the second type,  $r = (2^{mk} + 1)/(2^k + 1)$  and for this case  $\Delta_r(y)$  is restricted to the form

$$\Delta_r(y) = 0, \quad \pm 2^{(n+de)/2}, \quad 0 \leq d \leq m-2, \quad d \text{ odd}$$

where  $e = \text{GCD}(n, k)$  and  $d$  depends on the rank of a quadratic form over  $GF(2^e)$ .

The explicit evaluation of  $\Delta_r(y)$  is equivalent to the explicit evaluation of weight distribution of the  $(2^n - 1, 2n)$  cyclic code whose dual code is generated by  $f_1(x)f_r(x)$ , the product of two primitive polynomials of degree  $n$ .

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**MULTI-VALUED CROSS-CORRELATION FUNCTIONS BETWEEN TWO  
MAXIMAL LINEAR RECURSIVE SEQUENCES**

# CHAPTER I

## INTRODUCTION

### 1. INTRODUCTION

In recent years, maximal linear recursive sequences have been studied extensively, notably by Zierler [1], Gold [2] [3] [4] [5], Kasami [6] [7] [8], Solomon [9], Golomb [10] [11], Welch [12] and Trachtenberg. [13] Maximal linear recursive sequences have ideal auto-correlation values:  $C_{aa}(\tau) = -1$  for all  $\tau \neq 0$  and  $C_{aa}(0) = \text{period of the sequence}$ . This ideal auto-correlation property is useful in synchronization technique in communication systems. Some maximal linear recursive sequences have uniformly low cross-correlation values which is responsible for extensive application in spread spectrum communication systems for multiplexing operations. [4]

What is involved in study of the cross-correlation functions between two maximal linear recursive sequences is to evaluate explicitly a quantity  $\Delta_r(y)$  where

$$\Delta_r(y) = \sum_{x \in GF(2^n)} (-1)^{\text{Tr}(xy + x^r)}.$$

Generally it is difficult to evaluate  $\Delta_r(y)$  algebraically. However, when the decimation  $r$  takes on some special values, the explicit evaluation of  $\Delta_r(y)$  is possible. Gold, Kasami and Solomon evaluated  $\Delta_r(y)$  for the case  $r = 2^k + 1$ ,  $n/\text{GCD}(n,k)$  odd. Kasami and Solomon found

the weight distribution formula of the  $(2^n-1, 2n)$  cyclic code generated by the polynomial  $f_1(x)f_r(x)$  via linear recursion. The evaluation of weights of the above cyclic code is equivalent to the evaluation of  $\Delta_r(y)$ . Welch evaluated  $\Delta_r(y)$  for the case  $r = 2^{2k} - 2^k + 1$ ,  $n/\text{GCD}(n, k)$  odd. Trachtenberg extended the above two results to non-binary cases. He introduced a quantity

$$\Delta'_r(y) = \sum_{x \in \text{GF}(p^n)} \rho^{\text{Tr}(xy - x^r)}$$

where  $p$  is an odd prime and  $\rho$  is a complex  $p$ -th root of unity. He evaluated  $\Delta'_r(y)$  for the cases  $r = (p^{2k} + 1)/2$  and  $r = (p^{2k} - p^k + 1)$ ,  $n$  odd,  $r \not\equiv p^j \pmod{p^n-1}$  for any  $j$ . In this paper only the binary cases are considered.

In CHAPTER I fundamental concepts of linear recursive sequences are introduced. All materials are discussed extensively in [10] and [14]. Also given are known results on 3-valued  $\Delta_r(y)$  and a relationship between weights of the  $(2^n-1, 2n)$  cyclic code and cross-correlation values.

In CHAPTER II some properties are presented on  $\Delta_r(y)$  including  $\Delta_{-1}(y)$ . Also presented are useful lemmas on greatest common divisors of two integers and the complete results on 3-valued, 4-valued and 5-valued  $\Delta_r(y)$ .

In CHAPTER III  $\Delta_r(y)$  for the case  $r \equiv 2^k \pmod{2^{n/2}-1}$ ,  $\text{GCD}(r, 2^n-1) = 1$ , is considered.

In CHAPTER IV  $\Delta_r(y)$  for the case  $r = (2^{mk}+1)/(2^k+1)$ ,  $n/\text{GCD}(n, k)$  and  $m$  odd, is considered. Also given are some

conjectures on 3-valued and 5-valued  $\Delta_r(y)$ .

In CHAPTER V summary and comments are given.

## 2. LINEAR RECURSIVE SEQUENCES

A linear recursive sequence  $a_0, a_1, a_2, \dots$  is a sequence  $\{a_i\}$  which satisfies a recursion relation of the form:

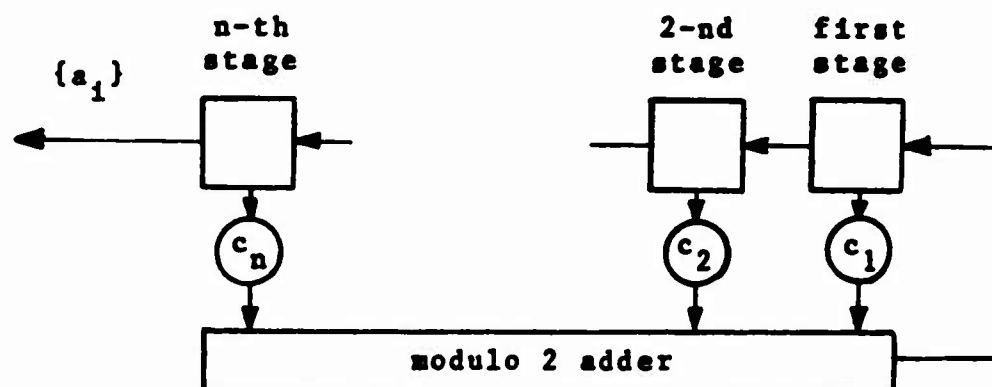
$$a_{n+1} = \sum_{j=1}^n c_j a_{n+1-j} \quad i = 0, 1, 2, \dots \quad (1.1)$$

where

$$a_i, c_i \in GF(2)$$

and  $a_0, a_1, \dots, a_{n-1}$  are pre-assigned.

Such a sequence  $\{a_i\}$  can be generated by the  $n$ -stage shift register and a modulo 2 adder as shown below.



Constants  $c_1, c_2, \dots, c_n$  are feedback coefficients;  
 $c_i = 1$  if there is a tap connected to the  $i$ -th stage and  
 $c_i = 0$  if there is no tap connected to the  $i$ -th stage.  
 The linear recursive sequence  $\{a_i\}$  of (1.1) can be thought  
 as the sequence of outputs observed from the  $n$ -th stage.

Initial conditions  $a_0, a_1, \dots, a_{n-1}$  are pre-assigned with the  $k$ -th stage containing  $a_{n-k}$ ,  $1 \leq k \leq n$ . Contents of stages are added modulo 2 according to the recursion relation (1.1) and the sum is fed back to the first stage as the content of the  $i$ -th stage is shifted into the  $(i+1)$ -st stage,  $1 \leq i \leq n-1$ , and the output  $a_j$  is shifted out from the  $n$ -th stage.

It is said that the sequence  $\{a_i\}$  is generated by the polynomial  $f(x)$  via linear recursion where

$$f(x) = x^n + c_1 x^{n-1} + c_2 x^{n-2} + \dots + c_n. \quad (1.2)$$

In this paper only sequences considered are those whose recursion polynomial  $f(x)$  is irreducible.

In order to describe the sequence  $\{a_i\}$  mathematically, consider the sequence  $\{b_j\}$  defined by  $b_j = a_{n+j}$ . Then, the sequence  $\{b_j\}$  can be characterized by its Z-transform in terms of the recursion polynomial  $f(x)$  of the sequence  $\{a_i\}$  and initial conditions  $a_0, a_1, \dots, a_{n-1}$ . Let  $B(z)$  be the Z-transform of the sequence  $\{b_j\}$ :

$$B(z) = b_0 + b_1 z + b_2 z^2 + b_3 z^3 + \dots$$

Then,  $B(z)$  can be expressed as follows:

$$B(z) = \frac{\sum_{j=1}^n c_j z^j \left\{ \sum_{i=0}^{j-1} a_{n-j+i} z^{1-i} \right\}}{z^n f(1/z)} \quad (1.3)$$

Many interesting properties of the sequence  $\{b_j\}$  can be derived by investigating  $B(z)$ .<sup>[10]</sup> However, it is necessary to expand (1.3) if one is to determine  $\{b_j\}$ .

The period  $p$  of the sequence  $\{a_i\}$  is the smallest positive integer  $p$  such that  $a_{i+p} = a_i$  for all  $i$ . In terms of the recursion polynomial  $f(x)$ , the period  $p$  is the smallest positive integer  $p$  such that  $f(x)$  divides  $(x^p + 1)$ . When  $p = 2^n - 1$ , the polynomial  $f(x)$  is said to be primitive. Since  $f(x)$  is irreducible, the period  $p$  of the sequence  $\{a_i\}$  is some divisor of  $(2^n - 1)$ . When  $p = 2^n - 1$ , the sequence  $\{a_i\}$  is said to be a maximal linear recursive sequence or a maximal-length linear shift register sequence. The sequence  $\{a_i\}$  is a maximal linear recursive sequence if and only if its recursion polynomial  $f(x)$  is primitive. For the remainder of this paper, only the maximal linear recursive sequence  $\{a_i\}$  is considered and the analysis is made on  $\{a_i\}_{i=0}^{2^n-2}$ . A maximal linear recursive sequence is henceforth simply referred to as a maximal sequence.

A more algebraic representation of a maximal sequence  $\{a_i\}_{i=0}^{2^n-2}$  generated by (1.1), which gives more insight to the structure of such a sequence, is the following: [14]

$$a_i = \text{Tr}_1^n(\beta \alpha^i) \quad (1.4)$$

where

$\alpha$  is the root of  $f(x)$  of (1.2),

$\text{Tr}_e^n(x)$  is the trace linear functional  $\text{Tr}_e^n$  on  $\text{GF}(2^n)$  defined by

$$\text{Tr}_e^n(x) = x + x^{2^e} + x^{2^{2e}} + \dots + x^{2^{(n-1)e}} \quad (1.5)$$

and  $\beta$  is some element of  $\text{GF}(2^n)$ .

The element  $\beta$  is determined uniquely by initial conditions  $a_0, a_1, \dots, a_{n-1}$  from the following matrix-vector equation:

$$\underline{D} \underline{x}^t = \underline{y}^t$$

where

$$\underline{D} = (d_{i,j}), \quad 1 \leq i, j \leq n,$$

$$\underline{x} = (x_i), \quad \underline{y} = (y_i), \quad 1 \leq i \leq n,$$

$$d_{i,j} = \alpha^{(i-1)2^{j-1}},$$

$$x_i = \beta^{2^{i-1}},$$

$$y_i = a_{i-1},$$

$$a_i, \alpha \text{ and } \beta \text{ are as defined above}$$

and  $t$  indicates the transpose.

Since the determinant  $|D|$  is a van der Monde determinant, a vector  $\underline{x}$  and hence  $\beta$  can be determined uniquely from  $\underline{y}$ .

Since for any  $x, y \in GF(2^n)$  and for all  $i$ ,  $(x + y)^{2^i} = x^{2^i} + y^{2^i}$  and  $x^{2^n} = x$ , the following important properties of  $\text{Tr}_e^n(x)$  are true:

$$\text{Tr}_e^n(x) \in GF(2^e)$$

$$\text{Tr}_e^n(x + y) = \text{Tr}_e^n(x) + \text{Tr}_e^n(y)$$

$$\text{Tr}_e^n(x^{2^i}) = \text{Tr}_e^n(x)$$

A maximal sequence  $\{a_i\}_{i=0}^{2^n-2}$  is said to be in natural orientation if  $a_i = a_{2i}$  for all  $i$ , indices reduced modulo  $2^n-1$ . Such a sequence does exist and it is obtained when initial conditions  $a_0, a_1, \dots, a_{n-1}$  are given by:

$$a_i = \text{Tr}_1^n(\alpha^i) \quad (1.6)$$

That is,  $\beta = 1$  in (1.4). Henceforth, when we consider a maximal sequence  $\{a_i\}_{i=0}^{2^n-2}$  in natural orientation, it is understood that  $a_i$  is given by (1.6) for all  $i$ .

Throughout this paper a symbol "n" is reserved to represent the degree of the recursion polynomial  $f(x)$  in (1.2) or equivalently the number of stages used to generate a sequence  $\{a_i\}$ . A notation  $\text{Tr}$  is often used in place of  $\text{Tr}_1^n$ . When a confusion can arise,  $\text{Tr}_1^n$  will be used explicitly.

Given two integers  $r$  and  $q$ ,  $\text{GCD}(r, q)$  is the greatest positive integer that divides both  $r$  and  $q$ , and  $\text{LCM}(r, q)$  is the least positive integer that is divisible by both  $r$  and  $q$ . Given two polynomials  $s(x)$  and  $t(x)$ ,  $\text{GCD}(s(x), t(x))$  is the monic polynomial of greatest degree that divides both  $s(x)$  and  $t(x)$ .

A sequence  $\{b_i\}$  is said to be obtained by decimation  $r$  from a sequence  $\{a_i\}$  when  $b_i = a_{ri}$ , indices reduced modulo  $2^n-1$ . Let  $\{b_i\}$  be the sequence obtained from  $\{a_i\}$  by decimation  $r$ . When  $\{a_i\}$  is maximal,  $\{b_i\}$  is also maximal provided that  $\text{GCD}(r, 2^n-1) = 1$ . If  $\{a_i\}$  is in natural orientation and  $\text{GCD}(r, 2^n-1) = 1$ , then  $\{b_i\}$  is also in natural orientation. If  $f(x)$  is a recursion polynomial for  $\{a_i\}$ , then  $h(x)$  is a recursion polynomial for  $\{b_i\}$  provided that  $h(x)$  is the minimal polynomial of  $\alpha^r$  and  $\alpha$  is the root of  $f(x)$ .

There are  $\phi(p)$  integers from 1 to  $p$  which are relatively prime to  $p$ , where  $\phi(p)$  is the Euler  $\phi$ -function. These  $\phi(p)$  integers form a group  $G$  under multiplication modulo  $p$ . Let  $p = 2^n - 1$ . Then, the set  $H = \{1, 2, 2^2, \dots, 2^{n-1}\}$  forms a multiplicative subgroup of  $G$ . Proper cyclotomic cosets,  $C_i$ , of  $H$  can be obtained by multiplying elements of  $H$  by an element of  $G$ . That is,

$$C_1 = g_1 H, \quad g_1 = 1$$

$$C_2 = g_2 H, \quad g_2 \in G \text{ but } g_2 \notin C_1$$

.....

$$C_i = g_i H, \quad g_i \in G \text{ but } g_i \notin C_j \text{ for all } j, 1 \leq j < i$$

.....

There are exactly  $\phi(2^n - 1)/n$  proper cyclotomic cosets. A cyclotomic coset  $C = rH$  is called improper if  $\text{GCD}(r, 2^n - 1) \neq 1$ ,  $1 \leq r \leq 2^n - 1$ . Proper cyclotomic cosets and improper cyclotomic cosets constitute cyclotomic cosets modulo  $2^n - 1$ . There are  $\{-1 + \frac{1}{n} \sum_{d|n} \phi(d) 2^{n/d}\}$  cyclotomic cosets in all.  $q$  and  $r$  belong to the same cyclotomic coset if and only if there exists some integer  $j$ ,  $0 \leq j \leq n-1$ , such that  $q \equiv 2^j r \pmod{2^n - 1}$ . The smallest member of the cyclotomic coset is called the cyclotomic coset leader.

Let  $C_q$  and  $C_r$  be cyclotomic cosets containing  $q$  and  $r$  respectively. Then,  $C_q$  is called the inverse cyclotomic coset of  $C_r$  if there exists some integer  $j$ ,  $0 \leq j \leq n-1$ , such that  $qr \equiv 2^j \pmod{2^n - 1}$ .

Two maximal sequences  $\{b_i\}$  and  $\{b'_i\}$  obtained by

decimations  $r$  and  $r'$  respectively from the same maximal sequence  $\{a_i\}$  in natural orientation are identical if  $r$  and  $r'$  belong to the same proper cyclotomic coset.

There are  $\phi(2^n-1)/n$  distinct maximal sequences of period  $2^n - 1$ . They are generated by  $\phi(2^n-1)/n$  distinct primitive polynomials of degree  $n$ . Given one maximal sequence  $\{a_i\}$ , remaining  $\{\phi(2^n-1)/n - 1\}$  maximal sequences can be obtained from  $\{a_i\}$  by decimations  $r_2, r_3, \dots, r_{\phi(2^n-1)/n}$ , where  $r_1$  is any member of the  $i$ -th proper cyclotomic coset  $C_i$ .

The following lemma is a well known result of finite fields. It is often used in evaluation of cross-correlation functions between two maximal sequences.

LEMMA 1-1:

$$\sum_{x \in \text{GF}(2^n)} (-1)^{\text{Tr}(\sigma x)} = \begin{cases} 2^n & \text{if } \sigma = 0 \\ 0 & \text{if } \sigma \neq 0, \sigma \in \text{GF}(2^n) \end{cases}$$

proof:

If  $\sigma = 0$ , the result is trivial. If  $\sigma \neq 0$ ,  $\sigma x$  is equal to each element of  $\text{GF}(2^n)$  exactly once as  $x$  ranges over  $\text{GF}(2^n)$ . Hence,

$$\sum_{x \in \text{GF}(2^n)} (-1)^{\text{Tr}(\sigma x)} = \sum_{x \in \text{GF}(2^n)} (-1)^{\text{Tr}(x)}$$

Since  $x \in \text{GF}(2^n)$  if and only if  $x + x^{2^n} = 0$  and

$$\begin{aligned} x + x^{2^n} &= (x + x^2 + \dots + x^{2^{n-1}})(1 + x + x^2 + x^{2^2} + \dots + x^{2^{n-1}}) \\ &= \text{Tr}(x)\{1 + \text{Tr}(x)\}, \end{aligned}$$

exactly half of the elements have trace 1 and the others have trace 0. Hence

$$\sum_{x \in GF(2^n)} (-1)^{\text{Tr}(x)} = 0 \quad \text{QED}$$

### 3. CROSS-CORRELATION FUNCTIONS BETWEEN 2 MAXIMAL SEQUENCES

An auto-correlation function  $C_{aa}(\tau)$  of a maximal sequence  $\{a_i\}_{i=0}^{2^n-2}$  is defined as the number of bit-by-bit agreements minus the number of bit-by-bit disagreements between  $\{a_{i+\tau}\}_{i=0}^{2^n-2}$  and  $\{a_i\}_{i=0}^{2^n-2}$ , indices reduced modulo  $2^n-1$ . Then,  $C_{aa}(\tau)$  can be expressed as

$$C_{aa}(\tau) = \sum_{i=0}^{2^n-2} (-1)^{a_{i+\tau}} (-1)^{a_i}, \quad 0 \leq \tau \leq 2^n-2.$$

The auto-correlation function  $C_{aa}(\tau)$  of a maximal sequence is a two-valued function as shown in the following well-known theorem.

#### THEOREM 1-2: [1] [10]

$$C_{aa}(\tau) = \begin{cases} -1 & \text{for } \tau \neq 0 \\ 2^n-1 & \text{for } \tau = 0 \end{cases}$$

proof:

Using (1.6),  $C_{aa}(\tau)$  can be expressed as

$$\begin{aligned} C_{aa}(\tau) &= \sum_{i=0}^{2^n-2} (-1)^{\text{Tr}(\alpha^{i+\tau})} (-1)^{\text{Tr}(\alpha^i)} \\ &= \sum_{i=0}^{2^n-2} (-1)^{\text{Tr}(\alpha^{i+\tau} + \alpha^i)}. \end{aligned}$$

As  $i$  ranges from 0 to  $2^n-2$ ,  $\alpha^i$  takes on all non-zero elements of  $GF(2^n)$  once. Then, by letting  $\alpha^i = y \in GF(2^n)$ ,

$$\begin{aligned}
C_{aa}(\tau) &= \sum_{\substack{x \in GF(2^n) \\ x \neq 0}} (-1)^{\text{Tr}(xy + x)} \\
&= \sum_{\substack{x \in GF(2^n) \\ x \neq 0}} (-1)^{\text{Tr}\{x(y+1)\}} \\
&= -1 + \sum_{x \in GF(2^n)} (-1)^{\text{Tr}\{x(y+1)\}}
\end{aligned}$$

Using Lemma 1-1,

$$\begin{aligned}
&= \begin{cases} -1 & \text{if } y \neq 1 \\ -1+2^n & \text{if } y = 1 \end{cases} \\
&= \begin{cases} -1 & \text{if } \tau \neq 0 \\ -1+2^n & \text{if } \tau = 0 \end{cases}
\end{aligned}$$

QED

A cross-correlation function  $C_{ab}(\tau)$  between two maximal sequences  $\{a_i\}_{i=0}^{2^n-2}$  and  $\{b_i\}_{i=0}^{2^n-2}$  is defined similarly.

$$C_{ab}(\tau) = \sum_{i=0}^{2^n-2} (-1)^{a_{i+\tau}} (-1)^{b_i}$$

Let  $\{b_i\}_{i=0}^{2^n-2}$  be a maximal sequence obtained from  $\{a_i\}_{i=0}^{2^n-2}$  by proper decimation  $r$ . That is,  $b_i = a_{ri}$  and  $\text{GCD}(r, 2^n-1) = 1$ . Then,

$$\begin{aligned}
C_{ab}(\tau) &= \sum_{i=0}^{2^n-2} (-1)^{a_{i+\tau}} (-1)^{a_{ri}} \\
&= \sum_{i=0}^{2^n-2} (-1)^{\text{Tr}(a^{i+\tau} + a^{ri})} \\
&= \sum_{\substack{x \in GF(2^n) \\ x \neq 0}} (-1)^{\text{Tr}(xy + x^r)}
\end{aligned}$$

where  $y = a^r$ . Then,

$$C_{ab}(\tau) = -1 + \sum_{x \in GF(2^n)} (-1)^{\text{Tr}(xy + x^r)}.$$

The following two theorems are known.

THEOREM 1-3: Gold [5], Kasami [6] [7] [8], Solomon [9]

If  $r = 2^k + 1$  and  $n/e$  is odd where  $e = \text{GCD}(n, k)$ ,  $C_{ab}(\tau)$  takes on only three values:  $-1$ ,  $-1 \pm 2^{(n+e)/2}$ .

$$C_{ab}(\tau) = \begin{cases} -1 & 2^{n-e-1} - 1 \text{ times} \\ -1 + 2^{(n+e)/2} & 2^{n-e-1} + 2^{(n-e-2)/2} \text{ times} \\ -1 - 2^{(n+e)/2} & 2^{n-e-1} - 2^{(n-e-2)/2} \text{ times} \end{cases} \quad (1.7)$$

THEOREM 1-4: Golomb [11], Welch [12]

If  $r = 2^{2k} - 2^k + 1$  and  $n/e$  is odd where  $e = \text{GCD}(n, k)$ ,  $C_{ab}(\tau)$  takes on only three values. The distribution of cross-correlation values are again given by (1.7).

The conjecture [11] that for  $n$  odd  $r = 2^{(n-1)/2} + 2^d + 1$  where  $d$  is a divisor of  $(n-1)$  leads to the three-valued  $C_{ab}(\tau)$  was shown to be false. Some of the counter-examples are given below.

| $n$ | $d$ | $r$              | $C_{ab}(\tau)$ |
|-----|-----|------------------|----------------|
| 7   | 6   | 73               | 7-valued       |
| 11  | 2   | 37               | 5-valued       |
| 11  | 10  | $1057 \equiv 67$ | 9-valued       |
| 13  | 2   | 69               | 10-valued      |

|    |    |                   |           |
|----|----|-------------------|-----------|
| 13 | 3  | 73                | 18-valued |
| 13 | 4  | 81                | 9-valued  |
| 13 | 12 | $4161 \equiv 131$ | 19-valued |

However, the Welch's conjecture [1] that  $r = 2^{(n-1)/2} + 3$  leads to the 3-valued  $C_{ab}(\tau)$  has been verified for  $n \leq 17$ .

Define  $\Delta_r(y)$  by

$$\begin{aligned}\Delta_r(y) &= \sum_{x \in GF(2^n)} (-1)^{\text{Tr}(xy + x^r)} \\ &= C_{ab}(\tau) + 1\end{aligned}$$

where

$$y = \alpha^r \text{ and } \{b_i\}_{i=0}^{2^n-2} = \{a_{ri}\}_{i=0}^{2^n-2}.$$

For the remainder of this paper, by the cross-correlation function, we shall mean  $\Delta_r(y)$  rather than  $C_{ab}(\tau)$ .

Let  $n(\Delta_r)$  denote the number of times  $\Delta_r(y)$  takes on the value  $\Delta_r$  as  $y$  ranges over  $GF(2^n)$ . In terms of  $\Delta_r(y)$  Theorem 1-3 and Theorem 1-4 become:

**THEOREM 1-5:**

If  $r = 2^k + 1$  and  $n/e$  is odd where  $e = \text{GCD}(n, k)$ ,  $\Delta_r(y)$  takes on three values:  $0, \pm 2^{(n+e)/2}$

| $\Delta_r$     | $n(\Delta_r)$               |       |
|----------------|-----------------------------|-------|
| $2^{(n+e)/2}$  | $2^{n-e-1} + 2^{(n-e-2)/2}$ |       |
| 0              | $2^n - 2^{n-e}$             | (1.8) |
| $-2^{(n+e)/2}$ | $2^{n-e-1} - 2^{(n-e-2)/2}$ |       |

**THEOREM 1-6:**

If  $r = 2^{2k} - 2^k + 1$  and  $n/e$  is odd where  $e = \text{GCD}(n, k)$ ,  $\Delta_r(y)$  takes on three values.  $\Delta_r$  and  $n(\Delta_r)$  are given by (1.8).

For the remainder of this paper, an analysis is made on  $\Delta_r(y)$  instead of  $C_{ab}(\tau)$ .

The complete evaluation of  $\Delta_r(y)$  for  $n = 3$  through  $n = 12$  has been carried out on IBM 370 at the University Computing Center. For  $n = 13$  through  $n = 15$ , only those proper decimations that lead to 7 or less distinct cross-correlation values are determined. For  $n = 16$ , only those that lead to 5 or less cross-correlation values are determined. In obtaining these results, the author used the Fourier transform algorithm developed by Welch.

Table 1-1 lists coset leaders of the cyclotomic cosets containing  $r$  and  $r^{-1}$  where  $r$  is given by Theorems 1-5 and 1-6. Table 1-2 lists decimations that result in 3-valued  $\Delta_r(y)$  for  $n = 3$  through  $n = 16$  that are not covered by either Theorem 1-5 or Theorem 1-6. Two decimations  $r$  and  $r'$  are given in pair as the cyclotomic cosets containing  $r$  and  $r'$  are inverse of each other. The complete results appear in APPENDIX A.

#### 4. $(2^n - 1, 2n)$ CYCLIC CODES

Theorems on cross-correlation values  $\Delta_r$  and their

## CYCLOTOMIC COSET LEADERS GIVEN BY THEOREMS 1-5 &amp; 1-6

|     |    | THEOREM 1-5 |   |    |          | THEOREM 1-6 |          |
|-----|----|-------------|---|----|----------|-------------|----------|
|     |    | k           | e | r  | $r^{-1}$ | r           | $r^{-1}$ |
| n = | 3  | 1           | 1 | 3  | 3        | 3           | 3        |
| n = | 5  | 1           | 1 | 3  | 11       | 3           | 11       |
|     |    | 2           | 1 | 5  | 7        | 11          | 3        |
| n = | 6  | 2           | 2 | 5  | 13       | 13          | 5        |
| n = | 7  | 1           | 1 | 3  | 43       | 3           | 43       |
|     |    | 2           | 1 | 5  | 27       | 13          | 11       |
|     |    | 3           | 1 | 9  | 15       | 23          | 29       |
| n = | 9  | 1           | 1 | 3  | 171      | 3           | 171      |
|     |    | 2           | 1 | 5  | 103      | 13          | 59       |
|     |    | 3           | 3 | 9  | 57       | 57          | 9        |
|     |    | 4           | 1 | 17 | 31       | 47          | 87       |
| n = | 10 | 2           | 2 | 5  | 205      | 13          | 79       |
|     |    | 4           | 2 | 17 | 181      | 79          | 13       |
| n = | 11 | 1           | 1 | 3  | 683      | 3           | 683      |
|     |    | 2           | 1 | 5  | 411      | 13          | 315      |
|     |    | 3           | 1 | 9  | 231      | 57          | 413      |
|     |    | 4           | 1 | 17 | 365      | 143         | 43       |
|     |    | 5           | 1 | 33 | 63       | 95          | 151      |

|        |   |   |     |       |     |       |
|--------|---|---|-----|-------|-----|-------|
| n = 12 | 4 | 4 | 17  | 241   | 241 | 17    |
| n = 13 | 1 | 1 | 3   | 2731  | 3   | 2731  |
|        | 2 | 1 | 5   | 1639  | 13  | 635   |
|        | 3 | 1 | 9   | 911   | 57  | 723   |
|        | 4 | 1 | 17  | 1453  | 241 | 171   |
|        | 5 | 1 | 33  | 1243  | 287 | 1691  |
|        | 6 | 1 | 65  | 127   | 191 | 1245  |
| n = 14 | 2 | 2 | 5   | 3277  | 13  | 1339  |
|        | 4 | 2 | 17  | 2893  | 241 | 205   |
|        | 6 | 2 | 65  | 2773  | 319 | 979   |
| n = 15 | 1 | 1 | 3   | 10923 | 3   | 10923 |
|        | 2 | 1 | 5   | 6555  | 13  | 2523  |
|        | 3 | 3 | 9   | 3641  | 57  | 575   |
|        | 4 | 1 | 17  | 1935  | 241 | 3671  |
|        | 5 | 5 | 33  | 993   | 993 | 33    |
|        | 6 | 3 | 65  | 3529  | 575 | 57    |
|        | 7 | 1 | 129 | 255   | 383 | 4791  |

TABLE 1-1

INVERSE PAIRS OF COSET LEADERS GIVING 3-VALUED  $\Delta_r(y)$ 

|        |            |             |
|--------|------------|-------------|
| n = 9  | 19 - 27    |             |
| n = 10 | 25 - 41    | 49 - 107    |
| n = 11 | 35 - 117   | 107 - 249   |
| n = 13 | 67 - 367   | 71 - 347    |
| n = 14 | 113 - 145  | 193 - 1613  |
| n = 15 | 131 - 4815 | 1371 - 2033 |

TABLE 1-2

distribution  $\eta(\Delta_r)$  as  $y$  ranges over  $GF(2^n)$  can be translated to theorems on weight  $w$  and weight distribution  $\eta(w)$  of a  $(2^n-1, 2n)$  cyclic code.

Let  $C$  be the  $(2^n-1, 2n)$  cyclic code generated by the polynomial  $f_1(x)f_r(x)$  via linear recursion. That is,  $C$  is the  $(2^n-1, 2n)$  cyclic code whose generator polynomial [15] is given by:

$$(x^{2^n-1} + 1)/x^{2n}f_1(1/x)f_r(1/x)$$

where

$f_1(x)$  is the minimal polynomial of  $\alpha^1$  and  $\alpha$  is a primitive element of  $GF(2^n)$ .

Then, given a codeword  $\underline{a} = (a_0, a_1, a_2, \dots, a_{2^n-2})$  in  $C$ , there exist two elements  $c$  and  $d$  in  $GF(2^n)$  such that the Mattson-Solomon polynomial associated with  $\underline{a}$  is given by  $g_{\underline{a}}(x) = \text{Tr}(cx) + \text{Tr}(dx^r)$  and  $a_i = g_{\underline{a}}(\alpha^i)$ . Note that  $c = 0 = d$  corresponds to the case  $w(\underline{a}) = 0$  and  $c = 0, d \neq 0$  or  $c \neq 0, d = 0$  corresponds to the case  $w(\underline{a}) = 2^{n-1}$ .

The non-zero weight  $w$  and the weight distribution  $\eta(w)$  of the  $(2^n-1, 2n)$  cyclic code  $C$  are related to  $\Delta_r$  and  $\eta(\Delta_r)$  as follows:

$$\begin{aligned} w &= (2^n - \Delta_r)/2 \\ \eta(w) &= (2^n - 1)\eta(\Delta_r) && \text{for } \Delta_r \neq 0 \\ \eta(w) &= (2^n - 1)\{\eta(\Delta_r) + 1\} && \text{for } \Delta_r = 0 \end{aligned} \quad (1.9)$$

The minimum weight  $w_{\min}$  of  $C$  is given by:

$$w_{\min} = \{2^n - \max_y \Delta_r(y)\}/2.$$

In terms of cyclic codes, Theorem 1-5 and Theorem 1-6 become :

The  $(2^n-1, 2n)$  cyclic code generated by the polynomial  $f_1(x)f_r(x)$  via linear recursion is a tri-weight code with the following weight distribution for non-zero weight.

| $w$                       | $n(w)$                                 |
|---------------------------|----------------------------------------|
| $2^{n-1} - 2^{(n+e-2)/2}$ | $(2^n - 1)(2^{n-e-1} + 2^{(n-e-2)/2})$ |
| $2^{n-1}$                 | $(2^n - 1)(2^n - 2^{n-e} + 1)$         |
| $2^{n-1} + 2^{(n+e-2)/2}$ | $(2^n - 1)(2^{n-e-1} - 2^{(n-e-2)/2})$ |

where

$$r = 2^k + 1 \text{ or } r = 2^{2k} - 2^k + 1 ,$$

$$e = \text{GCD}(n, k)$$

and  $n/e$  is odd.

## CHAPTER II

### PROPERTIES OF CROSS-CORRELATION FUNCTIONS

#### 1. FURTHER RESULTS ON $\Delta_r(y)$

In this section, we derive some results on  $\Delta_r(y)$ . The first three lemmas are the direct results of the structures of finite fields. Proofs are omitted.

##### LEMMA 2-1:

$$\Delta_r(y) = \Delta_r(y^{2^k}) \text{ for all } k.$$

##### LEMMA 2-2:

If  $r$  and  $r'$  belong to the same proper cyclotomic coset,  $\Delta_r(y) = \Delta_{r'}(y)$  for all  $y$ .

##### LEMMA 2-3:

If two proper decimations  $r$  and  $r'$  are such that  $r \cdot r' \equiv 2^k \pmod{2^n - 1}$ ,  $\Delta_r(y) = \Delta_{r'}(y^{-r})$ .

##### LEMMA 2-4:

$$\Delta_r(y) \equiv 0 \pmod{4} \text{ for all } y \text{ and } r, \text{ GCD}(r, 2^n - 1) = 1.$$

##### proof:

Let  $N(i, j)$  be the number of times the ordered pair  $\{\text{Tr}(xy) = i \text{ and } \text{Tr}(x^r) = j\}$  occurs as  $x$  ranges over  $\text{GF}(2^n)$ .

Both mappings  $x \rightarrow xy$ ,  $y \neq 0$ , and  $x \rightarrow x^r$ ,  $\text{GCD}(r, 2^n - 1)$

$= 1$ , permute elements of  $GF(2^n)$ . Since exactly half of the elements have trace 1 and the others have trace 0, the following 4 equalities hold.

$$N(0,0) + N(0,1) = 2^{n-1}$$

$$N(1,0) + N(1,1) = 2^{n-1}$$

$$N(0,0) + N(1,0) = 2^{n-1}$$

$$N(0,1) + N(1,1) = 2^{n-1}$$

which imply

$$N(0,0) = N(1,1)$$

$$N(1,0) = N(0,1)$$

$$N(0,0) \text{ and } N(1,0) \text{ are both even or both odd.}$$

From the definition of the cross-correlation function,

$$\begin{aligned} \Delta_r(y) &= \{N(0,0) + N(1,1)\} - \{N(1,0) + N(0,1)\} \\ &= 2\{N(0,0) - N(1,0)\}. \end{aligned}$$

Since  $N(0,0)$  and  $N(1,0)$  are both even or both odd,

$\{N(0,0) - N(1,0)\}$  is even. Hence,  $\Delta_r(y) \equiv 0 \pmod{4}$ .

Clearly  $\Delta_r(y) = 0$  when  $y = 0$ .

QED

In view of Lemma 2-4,  $\Delta_r(y)$  exhibits interesting characteristics for  $r = 2^{n-1} - 1$ . Since  $2r \equiv -1 \pmod{2^{n-1}}$ ,  $\Delta_r(y)$  for  $r = 2^{n-1} - 1$  is a cross-correlation function between a maximal sequence and its reverse sequence. In [16] Dowling and McEliece gave the bound on  $\Delta_{-1}(y)$ . They applied the result on exponential sums in finite fields given by Carlitz and Uchiyama. [17]

$$\left| \sum_{\substack{x \in GF(p^n) \\ x \neq 0}} \exp\left\{\frac{2\pi i}{p} \text{Tr}(xy + x^{-1})\right\} \right| \leq 2 \cdot p^{n/2}$$

where

$$y \in GF(p^n).$$

Letting  $p = 2$ ,

$$\begin{aligned} & \left| \sum_{\substack{x \in GF(2^n) \\ x \neq 0}} (-1)^{\text{Tr}(xy + x^{-1})} \right| \\ &= |\Delta_{-1}(y) - 1| \leq 2^{(n+2)/2} \end{aligned}$$

Hence,

$$-2^{(n+2)/2} + 1 \leq \Delta_{-1}(y) \leq 2^{(n+2)/2} + 1 \quad (2.1)$$

Since  $2^{(n+2)/2} \equiv 0 \pmod{4}$  for  $n$  even, we have

$$-2^{(n+2)/2} + 4 \leq \Delta_{-1}(y) \leq 2^{(n+2)/2}, \quad n \text{ even} \quad (2.2)$$

Let  $L_n$  denote the number of distinct values that  $\Delta_{-1}(y)$  takes on as  $y$  ranges over  $GF(2^n)$ .

It has been verified that the bounds on  $\Delta_{-1}(y)$  given by (2.1) and (2.2) are extremely tight for  $n \leq 18$  as shown in Table 2-1. For  $n$  odd, the largest and the smallest values between the bounds of (2.1) which are of the form  $4K$  are attained by  $\Delta_{-1}(y)$  for some  $y$ . For  $n$  even, both the upper and the lower bounds of (2.2) are attained. Furthermore, since  $L_n = \{\max\{\Delta_{-1}(y)\} - \min\{\Delta_{-1}(y)\}\}/4 + 1$  for all  $n \leq 18$ ,  $\Delta_{-1}(y)$  takes on all values of the form  $4K$  between the bounds of (2.1).

## CROSS-CORRELATION VALUES OF REVERSE SEQUENCES

| $n$ | $2^{(n+2)/2}$ | $\max_y \{\Delta_{-1}(y)\}$ | $\min_y \{\Delta_{-1}(y)\}$ | $L_n$ |
|-----|---------------|-----------------------------|-----------------------------|-------|
| 3   | 5.7           | 4                           | -4                          | 3     |
| 4   | 8             | 8                           | -4                          | 4     |
| 5   | 11.3          | 12                          | -8                          | 6     |
| 6   | 16            | 16                          | -12                         | 8     |
| 7   | 22.6          | 20                          | -20                         | 11    |
| 8   | 32            | 32                          | -28                         | 16    |
| 9   | 45.3          | 44                          | -44                         | 23    |
| 10  | 64            | 64                          | -60                         | 32    |
| 11  | 90.5          | 88                          | -88                         | 45    |
| 12  | 128           | 128                         | -124                        | 64    |
| 13  | 181.02        | 180                         | -180                        | 91    |
| 14  | 256           | 256                         | -252                        | 128   |
| 15  | 362.0         | 360                         | -360                        | 181   |
| 16  | 512           | 512                         | -508                        | 256   |
| 17  | 724.1         | 724                         | -720                        | 362   |
| 18  | 1024          | 1024                        | -1020                       | 512   |

TABLE 2-1

LEMMA 2-5:

$$\sum_{y \in GF(2^n)} \{\Delta_r(y)\} = 2^n$$

$$\sum_{y \in GF(2^n)} \{\Delta_r(y)\}^2 = 2^{2n}$$

proof:

$$\begin{aligned} & \sum_{y \in GF(2^n)} \{\Delta_r(y)\} \\ &= \sum_{y \in GF(2^n)} \sum_{x \in GF(2^n)} (-1)^{\text{Tr}(xy + x^r)} \\ &= \sum_{x \in GF(2^n)} (-1)^{\text{Tr}(x^r)} \sum_{y \in GF(2^n)} (-1)^{\text{Tr}(xy)} = 2^n \end{aligned}$$

The last step follows from Lemma 1-1 since the second sum is equal to 0 for all  $x \neq 0$  and  $2^n$  for  $x = 0$ .

$$\begin{aligned} & \sum_{y \in GF(2^n)} \{\Delta_r(y)\}^2 \\ &= \sum_{y \in GF(2^n)} \sum_{x \in GF(2^n)} (-1)^{\text{Tr}(xy + x^r)} \sum_{z \in GF(2^n)} (-1)^{\text{Tr}(zy + z^r)} \\ &= \sum_{x, z \in GF(2^n)} (-1)^{\text{Tr}(x^r + z^r)} \sum_{y \in GF(2^n)} (-1)^{\text{Tr}\{(x + z)y\}} \end{aligned}$$

Again from Lemma 1-1, the second sum is equal to  $2^n$  if  $x + z = 0$  or  $x = z$  and the sum is equal to zero if  $x + z \neq 0$ . Hence,

$$\begin{aligned} &= 2^n \sum_{\substack{x, z \in GF(2^n) \\ x = z}} (-1)^{\text{Tr}(x^r + z^r)} = 2^n \sum_{\substack{x, z \in GF(2^n) \\ x = z}} 1 \\ &= 2^{2n} \end{aligned}$$

QED

Lemma 2-5 is useful in evaluating the distribution of  $\Delta_r(y)$ . In order to obtain the higher moment identities, it

is necessary to evaluate the term

$$\sum_{x_1} \sum_{x_2} \dots \sum_{x_k} (-1)^{\text{Tr}(x_1^r + x_2^r + \dots + x_k^r)} \\ x_1 + x_2 + \dots + x_k = 0$$

However, it is not trivial to evaluate the above sum.

Let  $C$  and  $C'$  be the  $(2^n-1, 2n)$  and the  $(2^n-1, 2^n-2n-1)$  cyclic codes whose generator polynomials are  $(x^{2^n-1} + 1)/x^{2n}f_1(1/x)f_r(1/x)$  and  $f_1(x)f_r(x)$  respectively.

Let  $a_1$  and  $b_1$  be the number of codewords of weight 1 in  $C$  and  $C'$  respectively. Since  $C'$  is a subcode of a cyclic Hamming code,  $b_1 = 0 = b_2$ . Therefore, following power moment identities of Pless [18] must hold.

$$\sum_j j a_j = 2^{2n-1}(2^n-1)$$

$$\sum_j j^2 a_j = 2^{2n-2}(2^n-1)2^n = 2^{3n-2}(2^n-1)$$

These two power moment identities and two identities of Lemma 2-5 become identical with the substitution of  $j$  for  $w$  and  $a_j$  for  $\eta(w)$ . The two quantities  $w$  and  $\eta(w)$  are as given in (1.9)

## 2. GREATEST COMMON DIVISORS

For  $\Delta_r(y)$  to be a legitimate cross-correlation function between two maximal sequences, it is necessary and sufficient that  $\text{GCD}(r, 2^n-1) = 1$ . Hence, in the analysis of  $\Delta_r(y)$  for some  $r$ , it must first be shown that  $\text{GCD}(r, 2^n-1) = 1$ . However, for the remainder of this paper, when  $\Delta_r(y)$

is mentioned and no specific form is assumed for  $r$ , it is understood that  $\text{GCD}(r, 2^n - 1) = 1$ . The following three lemmas are well known results on greatest common divisors.

LEMMA 2-6:

$$\text{GCD}(2^m - 1, 2^n - 1) = 2^{\text{GCD}(m, n)} - 1$$

LEMMA 2-7:

$$\text{GCD}(2^m + 1, 2^n - 1) = 1 \text{ if and only if } n/\text{GCD}(m, n) \text{ is odd.}$$

LEMMA 2-8:

$$\text{If } m \text{ and } n/\text{GCD}(n, k) \text{ are both odd, } \text{GCD}\left(\frac{2^{mk} + 1}{2^k + 1}, 2^n - 1\right) = 1.$$

Lemma 2-7 was assumed in Theorems 1-3 and 1-5.

Lemma 2-8 with  $m = 3$  was assumed in Theorems 1-4 and 1-6 since  $(2^{3k} + 1)/(2^k + 1) = 2^{2k} - 2^k + 1$ . Using Lemmas 2-6 and 2-7, it is easy to show that the following decimations are all proper. That is,  $\text{GCD}(r_i, 2^n - 1) = 1$ .

$$r_1 = 2^{n/2+1} - 1 \quad n \equiv 0 \pmod{4} \quad (2.3)$$

$$r_2 = (2^{n/4} - 1)(2^{n/2} + 1) + 2 \quad n \equiv 0 \pmod{4} \quad (2.4)$$

$$r_3 = 2^{n/2} + 3 \quad n \equiv 0 \pmod{2} \quad (2.5)$$

$$r_4 = 2^{n/2} + 2^{n/2-1} - 1 \quad n \equiv 2 \pmod{4} \quad (2.6)$$

$$r_5 = 2^{n/2+2} - 3 \quad n \equiv 0 \pmod{2} \quad (2.7)$$

$$r_6 = 2^{n/2+2} + 2^{n/2} - 3 \quad n \equiv 0 \pmod{2} \quad (2.8)$$

In CHAPTER III  $\Delta_{r_i}(y)$  for these 6 decimations are analyzed.

In CHAPTER IV decimations of the form  $(2^{mk}+1)/(2^k+1)$  for  $m$  and  $n/\text{GCD}(n,k)$  odd are considered. However, due to Lemma 2-3, some results can be obtained on 3-valued  $\Delta_r(y)$  for decimations of this type. For  $m = n$  odd,  $\Delta_r(y)$  is a 3-valued function as given in the following lemma.

LEMMA 2-9:

For  $n$  odd and  $r = (2^{nk}+1)/(2^k+1)$ ,  $\Delta_r(y)$  takes on 3 values. Both  $\Delta_r$  and  $n(\Delta_r)$  are given by (1.8).

proof:

From Lemma 2-8,  $\text{GCD}(r, 2^n-1) = 1$ . Let  $r' = 2^k + 1$ . From Lemma 2-7,  $\text{GCD}(r', 2^n-1) = 1$ .  $rr' = 2^{nk} + 1 \equiv 2 \pmod{2^n-1}$ . The result follows immediately from Theorem 1-5 and Lemma 2-3. QED

In order to find the multiplicative inverse of  $r = 2^{2k} - 2^k + 1 = (2^{3k}+1)/(2^k+1)$  modulo  $2^n-1$  when  $\text{GCD}(n,3) = 1$ , consider the following.

Let  $m'$  be the multiplicative inverse of 3 mod  $n$ .  $m'$  exists if and only if  $\text{GCD}(n,3) = 1$ .  $m'$  is odd if  $n$  is even. If  $n$  is odd,  $m'$  can be even or odd.

Define  $m$  so that

$$\begin{aligned} m &= m' \text{ when } m' \text{ is odd and} \\ m &= n - m' \text{ when } m' \text{ is even.} \end{aligned} \tag{2.9}$$

Note that  $m$  is always odd.

Now consider  $r' = (2^{3km}+1)/(2^{3k}+1)$  where  $m$  is given

by (2.9). Then,

$$r \cdot r' = \frac{2^{3k+1}}{2^k + 1} \frac{2^{3km+1}}{2^{3k} + 1} = \frac{2^{3km+1}}{2^k + 1}.$$

When  $m'$  is odd,  $3m = 3m' = 1 + qn$  for some  $q$ .

$$r \cdot r' = \frac{2^{k+qnk+1}}{2^k + 1} = \frac{2^k (2^n)^{qk+1}}{2^k + 1} \equiv 1 \pmod{2^n - 1}$$

When  $m'$  is even,  $3m = 3(n-m') = 3n - (1+qn) = -1 + (3-q)n$  for some  $q$ .

$$\begin{aligned} r \cdot r' &= \frac{2^{-k+(3-q)nk+1}}{2^k + 1} = \frac{2^{-k} (2^n)^{(3-q)k+1}}{2^k + 1} \\ &\equiv \frac{2^{-k} + 1}{2^k + 1} \pmod{2^n - 1} \\ &\equiv 2^{n-k} \pmod{2^n - 1} \end{aligned}$$

Let  $3k \equiv j \pmod{n}$  and define  $t$  so that

$$\begin{aligned} t &= n-j && \text{when } j > (n-1)/2 \text{ and} \\ t &= j && \text{when } j \leq (n-1)/2 \end{aligned} \tag{2.10}$$

For this choice of  $t$ ,  $r' = (2^{3km+1})/(2^{3k}+1)$  and  $r'' = (2^{mt}+1)/(2^t+1)$  belong to the same cyclotomic coset modulo  $2^n-1$  and  $r \cdot r'' \equiv 2^i \pmod{2^n-1}$  for some  $i$ . Since  $\text{GCD}(n,3) = 1$ ,  $\text{GCD}(n,k) = \text{GCD}(n,3k) = \text{GCD}(n,t)$ . Hence, the following lemma follows directly from Theorem 1-6 and Lemma 2-3.

**LEMMA 2-10:**

Suppose  $\text{GCD}(3,n) = 1$ ,  $\text{GCD}(t,n) = e$  and  $n/e$  is odd. Let  $r'' = (2^{mt}+1)/(2^t+1)$ . Then,  $\Delta_{r''}(y)$  takes on 3 values.  $\Delta_{r''}$  and  $n(\Delta_{r''})$  are given by (1.8),

where

$$m = \begin{cases} m' & \text{when } m' \text{ is odd} \\ n - m' & \text{when } m' \text{ is even} \end{cases}$$

$m'$  is the multiplicative inverse of 3 mod  $n$ .

### 3. COMPUTED RESULTS FOR 3-, 4- and 5-VALUED $\Delta_r(y)$

Before proceeding to CHAPTER III, we give the complete results on proper cyclotomic coset leaders  $r$  that lead to 3-valued, 4-valued and 5-valued cross-correlation functions  $\Delta_r(y)$  for  $n \leq 16$ . They are listed in Tables 2-2, 2-3 and 2-4 respectively. Two coset leaders  $r$  and  $q$  are given in pair if  $rq \equiv 2^i \pmod{2^n-1}$  for some  $i$ ,  $0 \leq i \leq n-1$ . If  $r^2 \equiv 2^i \pmod{2^n-1}$  for some  $i$ ,  $0 \leq i \leq n-1$ , the coset leader  $r$  is given by itself. The letters following coset leaders give theorems, lemmas and conjectures that explain the corresponding  $\Delta_r(y)$ . The following notations are used:

|                   |                          |
|-------------------|--------------------------|
| A - Theorem 1-5   | B - Theorem 1-6          |
| C - Theorem 2-9   | D - Theorem 2-10         |
| E - Conj. 4-5 (1) | F - Conj. 4-5 (2) or (3) |
| G - Conj. 4-5 (4) | H - Conj. 4-5 (5)        |
| J - Theorem 3-6   | K - Theorem 3-7          |
| L - Theorem 3-5   | M - Theorem 3-8          |
| N - Conj. 4-6 (5) | $\emptyset$ - Lemma 4-1  |
| P - Conj. 4-3     | Q - Conj. 4-4            |
| R - Conj. 4-6 (1) | S - Conj. 4-6 (2)        |
| T - Conj. 4-6 (3) | U - Conj. 4-6 (4)        |

3-VALUED  $\Delta_T(y)$ 

|        |              |             |              |  |
|--------|--------------|-------------|--------------|--|
| n = 3  | 3ABCDEF      |             |              |  |
| n = 5  | 3ABD- 11BCD  | 5AF-        | 7CE          |  |
| n = 6  | 5A - 13BGH   |             |              |  |
| n = 7  | 3AB- 43CD    | 5A - 27C    | 9A - 15C     |  |
|        | 11DE- 13B    | 23B - 29DF  |              |  |
| n = 9  | 3AB- 171C    | 5A - 103C   | 9A - 57BC    |  |
|        | 13B - 59     | 17A - 31C   | 19EF- 27     |  |
|        | 47B - 87     |             |              |  |
| n = 10 | 5A - 205     | 13BD- 79BD  | 17A - 181    |  |
|        | 25 - 41H     | 49G - 107   |              |  |
| n = 11 | 3AB- 683CD   | 5A - 411C   | 9A - 231C    |  |
|        | 13B - 315D   | 17A - 365C  | 33A - 65C    |  |
|        | 35E - 117    | 43D - 143B  | 57B - 413D   |  |
|        | 95B - 151D   | 107 - 249F  |              |  |
| n = 12 | 17A - 241B   |             |              |  |
| n = 13 | 3AB- 2731CD  | 5A - 1639C  | 9A - 911C    |  |
|        | 13B - 635D   | 17A - 1453C | 33A - 1243C  |  |
|        | 57B - 723D   | 65A - 127C  | 67E - 367    |  |
|        | 71F - 347    | 171D - 241B | 191B - 1245D |  |
|        | 287B - 1691D |             |              |  |

|        |              |             |             |
|--------|--------------|-------------|-------------|
| n = 14 | 5A - 3277    | 13B - 1339D | 17A - 2893  |
|        | 65A - 2773   | 113 - 145H  | 193G - 1613 |
|        | 205D - 241B  | 319B - 979D |             |
| n = 15 | 3AB-10923C   | 5A - 6555C  | 9A - 3641C  |
|        | 13B - 2523   | 17A - 1935C | 33A - 993BC |
|        | 57B - 575B   | 65A - 3529C | 129A - 255C |
|        | 131E - 4815  | 241B - 3671 | 383B - 4791 |
|        | 1371 - 2033F |             |             |

TABLE 2-2

4-VALUED  $\Delta_T(y)$ 

|        |                |       |               |
|--------|----------------|-------|---------------|
| n = 4  | 7JK            |       |               |
| n = 8  | 31J - 91L      | 53K   |               |
| n = 12 | 127J -1387L    | 457K  |               |
| n = 16 | 511J -21931L   | 3857K | 7399L - 9947L |
|        | 11093L -13133L |       |               |

TABLE 2-3

5-VALUED  $\Delta_r(y)$ 

|        |              |              |              |  |
|--------|--------------|--------------|--------------|--|
| n = 6  | 11M - 23L    |              |              |  |
| n = 8  | 19M - 47L    | 23L - 61N    |              |  |
| n = 9  | 110PS - 93P  | 23R - 25RT   | 43P - 1070P  |  |
|        | 1090P        |              |              |  |
| n = 10 | 35M - 95L    | 101L - 157L  |              |  |
| n = 11 | 110Q- 189Q   | 25 - 87      | 37 - 83      |  |
|        | 47R - 49R    | 81 - 139     | 1210Q- 423Q  |  |
|        | 141 - 363    | 171Q - 2050Q | 187 - 427    |  |
|        | 2210Q- 343Q  | 229 - 295    | 2350Q- 429Q  |  |
|        | 311          |              |              |  |
| n = 12 | 67M - 191L   | 73 - 731     | 253N - 599L  |  |
| n = 13 | 110Q- 7450Q  | 19S - 485    | 43Q - 381Q   |  |
|        | 95R - 97R    | 113T - 363   | 147 - 949    |  |
|        | 161 - 973    | 2050Q- 9190Q | 225 - 405    |  |
|        | 4450Q- 4970Q | 483Q - 1645Q | 631 - 1707   |  |
|        | 683Q - 1643Q | 749Q - 1367Q | 869Q - 1461Q |  |
|        | 939Q - 953Q  |              |              |  |
| n = 14 | 131M - 383L  |              |              |  |
| n = 15 | 110P- 2979   | 43P - 765P   | 113U - 1451  |  |
|        | 171P - 5557  | 191R - 193R  | 2050P- 4965  |  |
|        | 683P - 5805P | 8930P- 3229  | 995 - 19430P |  |

|              |               |               |
|--------------|---------------|---------------|
| 1119 - 3543  | 1275P - 6605P | 1913P - 2895P |
| 2295P - 3755 | 2521 - 6827   | 2731P - 3277P |
| 2981 - 5463  | 3643 - 6557P  | 5783P - 6567P |
| 5813P        |               |               |

|        |             |               |               |
|--------|-------------|---------------|---------------|
| n = 16 | 259M - 767L | 383L - 13261L | 1021N - 9949L |
|--------|-------------|---------------|---------------|

TABLE 2-4

## CHAPTER III

### MULTI-VALUED CROSS-CORRELATION FUNCTIONS I

#### 1. PRELIMINARY

As mentioned in CHAPTER I, generally it is difficult to evaluate  $\Delta_r(y)$  algebraically. The case  $r = 2^k + 1$ ,  $n/\text{GCD}(n,k)$  odd, has been solved by Gold, Kasami and Solomon. The case  $r = 2^{2k} - 2^k + 1$ ,  $n/\text{GCD}(n,k)$  odd, has been solved by Welch. In this chapter we consider the case:

$$n = 2m \text{ and } r \equiv 2^k \pmod{2^m - 1}.$$

It can be shown that  $\Delta_r(y)$  is restricted to the form:

$$\Delta_r(y) = 2^m(j-1), \quad 0 \leq j \leq J \text{ for some } J.$$

$j$  is the number of distinct solutions to two equations over  $\text{GF}(2^n)$ . Hence, the upper bound  $J$  for  $\Delta_r(y)$  can be determined immediately by finding the degrees of equations. In general, however, equations are of high degree and the explicit evaluation of  $\Delta_r(y)$  is not trivial.

As mentioned in CHAPTER II, the following six decimations are analyzed in Section 2 of this chapter.

$$r_1 = 2^{n/2+1} - 1 \quad n \equiv 0 \pmod{4} \quad (3.1)$$

$$r_2 = (2^{n/4} - 1)(2^{n/2} + 1) + 2 \quad n \equiv 0 \pmod{4} \quad (3.2)$$

$$r_3 = 2^{n/2} + 3 \quad n \equiv 0 \pmod{2} \quad (3.3)$$

$$r_4 = 2^{n/2} + 2^{n/2-1} - 1 \quad n \equiv 2 \pmod{4} \quad (3.4)$$

$$r_5 = 2^{n/2+2} - 3 \quad n \equiv 0 \pmod{2} \quad (3.5)$$

$$r_6 = 2^{n/2+2} + 2^{n/2} - 3 \quad n \equiv 0 \pmod{2} \quad (3.6)$$

They are analyzed in Theorems 3-6, 3-7, ... , 3-11 respectively.

In Section 3 computed results are listed. Before the main theorem of this chapter, Theorem 3-5, is introduced, we consider the two preliminary lemmas and two results in solving equations over  $GF(2^n)$ .

LEMMA 3-1:

When  $n = 2m$ , non-zero element  $x$  of  $GF(2^n)$  can be represented as  $x = \alpha\beta$ , where  $\alpha \in GF(2^m)$ ,  $\alpha \neq 0$ , and  $\beta$  is  $(2^m+1)$ -st root of unity in  $GF(2^{2m})$ .

proof:

Let  $x \in GF(2^{2m})$ ,  $x \neq 0$ .

$$2^{2m} = 2^{m-1}(2^m + 1) + 2^{m-1}(2^m - 1)$$

$$x = x^{2^{2m}} = x^{2^{m-1}(2^m+1)} \cdot x^{2^{m-1}(2^m-1)}$$

Letting  $\alpha = x^{2^{m-1}(2^m+1)}$  and  $\beta = x^{2^{m-1}(2^m-1)}$ , it is easy to see that  $\alpha^{2^m-1} = 1$  and hence  $\alpha \in GF(2^m)$  and that  $\beta^{2^m+1} = 1$ .

To show the uniqueness of the representation, let  $\omega$  be a primitive element of  $GF(2^{2m})$ . Then, elements  $\alpha$  and  $\beta$  can be expressed as

$$\alpha = \omega^{(2^m+1)i}, \quad i = 1, 2, \dots, 2^m-1$$

$$\text{and } \beta = \omega^{(2^m-1)j}, \quad j = 1, 2, \dots, 2^m+1.$$

Assume that there exist elements  $\alpha, \beta, \alpha'$  and  $\beta'$  such that  $\alpha\beta = x = \alpha'\beta'$ . This says that given some  $i$  and  $j$ ,

there exist  $(2^m+1)s = \alpha'$  and  $(2^m-1)t = \beta'$  such that

$$\omega(2^m+1)i \omega(2^m-1)j = \omega(2^m+1)s \omega(2^m-1)t,$$

which implies

$$\begin{aligned} (2^m+1)i + (2^m-1)j &\equiv (2^m+1)s + (2^m-1)t \pmod{2^{2m}-1} \\ (2^m+1)(i-s) &\equiv (2^m-1)(t-j) \pmod{2^{2m}-1} \end{aligned} \quad (3.7)$$

Let  $i \geq s$  without loss of generality. Note that

$0 \leq (i-s) \leq 2^m-2$  and  $|t-j| \leq 2^m$ . From (3.7) we have

$$\begin{aligned} (2^m+1)(i-s) &= (2^m-1)(t-j) + k(2^{2m}-1) \\ &= (2^m-1)\{k(2^m+1) + (t-j)\} \end{aligned} \quad (3.8)$$

for some  $k$ ,  $k = 0$  or  $k = 1$ . (3.8) implies that

$$\text{LCM}(2^m+1, 2^m-1) \leq (2^m+1)(2^m-2) = 2^{2m} - 2^m - 2.$$

But since  $\text{GCD}(2^m+1, 2^m-1) = 1$ , we must have

$$\text{LCM}(2^m+1, 2^m-1) = (2^m+1)(2^m-1) = 2^{2m} - 1.$$

This is a contradiction. The only way in which (3.8) holds without contradiction is  $i = s$ ,  $t = j$  and  $k = 0$ . Or  $\alpha = \alpha'$  and  $\beta = \beta'$ . QED

### LEMMA 3-2:

Let  $n = 2m$ ,  $\alpha \in \text{GF}(2^m)$  and  $\beta \in \text{GF}(2^{2m})$ . Then

$$\text{Tr}_1^{2m}(\alpha\beta) = \text{Tr}_1^m\{\alpha(\beta + \beta^{2^m})\}.$$

proof:

$$\begin{aligned} \text{Tr}_1^{2m}(\alpha\beta) &= \sum_{j=0}^{2m-1} (\alpha\beta)^{2^j} \\ &= \sum_{j=0}^{m-1} (\alpha\beta)^{2^j} + \sum_{j=m}^{2m-1} (\alpha\beta)^{2^j} \end{aligned}$$

$$\begin{aligned}
&= \sum_{j=0}^{m-1} (\alpha\beta)^{2^j} + \sum_{j=0}^{m-1} (\alpha\beta)^{2^j} 2^{2^m} \\
&= \sum_{j=0}^{m-1} (\alpha\beta)^{2^j} + \sum_{j=0}^{m-1} (\alpha\beta^{2^m})^{2^j} \\
&= \sum_{j=0}^{m-1} \{\alpha(\beta + \beta^{2^m})\}^{2^j} \\
&= \text{Tr}_1^m \{\alpha(\beta + \beta^{2^m})\}
\end{aligned}$$

QED

In order to derive an explicit formula for  $\Delta_r(y)$ , it becomes necessary to find the number of distinct solutions to the equations over  $\text{GF}(2^n)$ . Following two lemmas are useful in finding the number of solutions.

LEMMA 3-3: McEliece [19]

The equation  $x^{2^k} + x = y$ ,  $y \in \text{GF}(2^n)$ ,  $k < n$ , has either no roots or  $2^e$  roots in  $\text{GF}(2^n)$ . It has  $2^e$  roots if and only if  $\text{Tr}_e^n(y) = 0$ , where  $e = \text{GCD}(n, k)$ .

LEMMA 3-4:

A polynomial  $g(x)$  is a repeated factor of a polynomial  $f(x)$  if and only if  $g(x)$  divides  $\text{GCD}(f(x), f'(x))$ , where  $f'(x)$  is the formal derivative of  $f(x)$ .

2.  $\Delta_r(y)$  FOR  $r \equiv 2^k \pmod{2^{n/2}-1}$

The crux for analyzing this type of  $\Delta_r(y)$  is the following theorem, which was suggested by Welch.

THEOREM 3-5:

If  $n = 2m$ ,  $\text{GCD}(r, 2^n - 1) = 1$  and  $r \equiv 2^k \pmod{2^m - 1}$  for some  $k$ ,  $0 \leq k \leq m-1$ , then  $\Delta_r(y) = 2^m(j-1)$ ,  $0 \leq j \leq J$ , for some  $j$  and  $J$ .

proof:

$$\begin{aligned}\Delta_r(y) &= \sum_{x \in \text{GF}(2^n)} (-1)^{\text{Tr}(xy + x^r)} \\ &= 1 + \sum_{\substack{x \in \text{GF}(2^n) \\ x \neq 0}} (-1)^{\text{Tr}(xy + x^r)}\end{aligned}$$

Using Lemma 3-1, let  $x = \alpha\beta$  where  $\alpha \in \text{GF}(2^m)$ ,  $\alpha \neq 0$ , and  $\beta^{2^m+1} = 1$ ,  $\beta \in \text{GF}(2^{2m})$ . Then,

$$\Delta_r(y) = 1 + \sum_{\substack{\beta \in \text{GF}(2^{2m}) \\ \beta^{2^m+1} = 1}} \sum_{\substack{\alpha \in \text{GF}(2^m) \\ \alpha \neq 0}} (-1)^{\text{Tr}_1^{2m}\{(y\alpha\beta + \alpha^r\beta^r)\}}$$

Furthermore, let  $r \equiv s \pmod{2^m+1}$ .

$$\begin{aligned}\text{Tr}_1^{2m}(y\alpha\beta + \alpha^r\beta^r) &= \text{Tr}_1^{2m}(y\alpha\beta + \alpha^{2^k}\beta^s) \\ &= \text{Tr}_1^{2m}\{y\alpha\beta + (\alpha^{2^k}\beta^s)2^{n-k}\} \\ &= \text{Tr}_1^{2m}\{y\alpha\beta + \alpha\beta s^{2^{-k}}\}\end{aligned}$$

Substituting this into the above,

$$\Delta_r(y) = 1 + \sum_{\substack{\beta \in \text{GF}(2^{2m}) \\ \beta^{2^m+1} = 1}} \sum_{\substack{\alpha \in \text{GF}(2^m) \\ \alpha \neq 0}} (-1)^{\text{Tr}_1^{2m}\{\alpha(y\beta + \beta s^{2^{-k}})\}}$$

Since there are  $2^m+1$   $(2^m+1)$ -st roots of unity in  $\text{GF}(2^{2m})$ , the above reduces to:

$$\Delta_r(y) = 1 + \sum_{\beta \in GF(2^{2m})} \sum_{\alpha \in GF(2^m)} (-1)^{\text{Tr}_1^{2m} \{ \alpha(y\beta + \beta s^{2^{-k}}) \}} \beta^{2^m+1} = 1 - (2^m + 1).$$

From Lemma 3-2,

$$\begin{aligned} & \text{Tr}_1^{2m} \{ \alpha(y\beta + \beta s^{2^{-k}}) \} \\ &= \text{Tr}_1^m \{ \alpha(y\beta + \beta s^{2^{-k}} + y^{2^m} \beta^{2^m} + \beta^{2^m} s^{2^{-k}}) \} \\ &= \text{Tr}_1^m \{ \alpha(y\beta + \beta s^{2^{-k}} + y^{2^m} \beta^{-1} + \beta^{-s^{2^{-k}}}) \}. \end{aligned}$$

Note that  $(y\beta + \beta s^{2^{-k}} + y^{2^m} \beta^{-1} + \beta^{-s^{2^{-k}}}) \in GF(2^m)$ . Let  $N_\beta(y)$  denote the number of distinct solutions  $\beta$  in  $GF(2^{2m})$  to (3.9) and (3.10).

$$y\beta + \beta s^{2^{-k}} + y^{2^m} \beta^{-1} + \beta^{-s^{2^{-k}}} = 0 \quad (3.9)$$

$$\beta^{2^m+1} = 1 \quad (3.10)$$

where

$$y \in GF(2^{2m}).$$

$$\begin{aligned} \Delta_r(y) &= -2^m + \sum_{\beta \in GF(2^{2m})} \sum_{\alpha \in GF(2^m)} (-1)^{\text{Tr}_1^m \{ \alpha(y\beta + \beta s^{2^{-k}} + y^{2^m} \beta^{-1} + \beta^{-s^{2^{-k}}}) \}} \beta^{2^m+1} = 1 \end{aligned}$$

Then using Lemma 1-1,

$$\begin{aligned} \Delta_r(y) &= -2^m + 2^m N_\beta(y) \\ &= 2^m \{ N_\beta(y) - 1 \} \end{aligned} \quad (3.11)$$

QED

Raising (3.9) to the power  $2^k$ , we obtain

$$y^{2^k} \beta^{2^k} + \beta^s + y^{2^{m+k}} \beta^{-2^k} + \beta^{-s} = 0.$$

Multiplying by  $\beta^j$  and raising to the power  $2^i$  for some  $j$  and  $i$ , the above equation can be transformed to the form:

$$z_1 \beta^{e_1} + z_2 \beta^{e_2} + z_3 \beta^{e_3} + z_4 = 0 \quad (3.12)$$

where

$e_1, e_2$  and  $e_3$  are non-negative integers with at least one odd  $e_i$  and  $z_1$  is some power of  $y$ .

Let  $e_{\max} = \max\{e_1, e_2, e_3\}$ . Note that  $e_{\max} \leq 2 \cdot \max\{|s|, 2^k\}$ .

Then  $\Delta_r(y) \leq 2^{e_{\max}}(e_{\max} - 1)$  since (3.12) can have at most  $e_{\max}$  roots for  $\beta$  in  $GF(2^{2m})$ . This upper bound on  $\Delta_r(y)$  gives the bound on the minimum weight of the corresponding  $(2^{2m}-1, 4m)$  cyclic code.

$$w_{\min} \geq (2^{2m} + 2^m - e_{\max} \cdot 2^m) / 2 \quad (3.13)$$

From (3.11),  $\min_y \Delta_r(y) = -2^m$ . This gives the upper bound on the weight of the cyclic code.

$$w_{\max} \leq (2^{2m} + 2^m) / 2 \quad (3.14)$$

From (3.11),  $\Delta_r(y)$  is negative if and only if  $N_\beta(y) = 0$ .

In view of Lemma 2-5, there must exist at least one  $y$  in  $GF(2^{2m})$  such that  $N_\beta(y) = 0$ . This implies that the upper bound in (3.14) is always achieved.

Therefore, if  $n = 2m$ ,  $\text{GCD}(r, 2^n - 1) = 1$ ,  $r \equiv 2^k \pmod{2^m - 1}$  and  $r \equiv s \pmod{2^m + 1}$ , the following bounds must hold for  $\Delta_r(y)$  and weight  $w$  of the corresponding  $(2^{2m}-1, 4m)$

cyclic code. Furthermore, the lower bound for  $\Delta_r(y)$  and the upper bound for  $w$  are always attained.

$$-2^m \leq \Delta_r(y) \leq 2^m(e_{\max} - 1) \quad (3.15)$$

$$2^{2m-1} + 2^{m-1} - e_{\max} \cdot 2^{m-1} \leq w \leq 2^{2m-1} + 2^{m-1} \quad (3.16)$$

$e_{\max}$  is itself bounded by:

$$e_{\max} \leq 2 \cdot \max\{|s|, 2^k\}.$$

We are now in a position to analyze decimations (3.1) through (3.6). They are considered in Theorems 3-6 through 3-11 respectively.

#### THEOREM 3-6:

If  $n \equiv 0 \pmod{4}$ ,  $n = 2m$  and  $r = 2^{m+1} - 1$ ,  $\Delta_r(y)$  is a 4-valued function.  $\Delta_r$  and  $\eta(\Delta_r)$  are given by:

| $\Delta_r$ | $\eta(\Delta_r)$         |
|------------|--------------------------|
| $2^{m+1}$  | $(2^{2m-1} - 2^{m-1})/3$ |
| $2^m$      | $2^m$                    |
| 0          | $2^{2m-1} - 2^{m-1}$     |
| $-2^m$     | $(2^{2m} - 2^m)/3$       |

#### proof:

From (2.3),  $\text{GCD}(r, 2^n - 1) = 1$ .

$$r = 2(2^m - 1) + 1 \equiv 1 \pmod{2^m - 1}$$

$$r = 2(2^m + 1) - 3 \equiv -3 \pmod{2^m + 1}$$

With  $k = 0$ ,  $s = -3$  and  $n = 2m$ , (3.9) becomes:

$$y\beta + \beta^{-3} + y^{2^m}\beta^{-1} + \beta^3 = 0 \quad (3.17)$$

Multiplying (3.17) by  $\beta^3$ ,

$$\beta^6 + y\beta^4 + y^{2^m}\beta^2 + 1 = 0 \quad (3.18)$$

Raising (3.18) to the power  $2^{-1}$  or equivalently  $2^{n-1}$ ,  $N_\beta(y)$  becomes the number of distinct solutions  $\beta$  in  $GF(2^{2m})$  to (3.19) and (3.20):

$$\beta^3 + y^{1/2}\beta^2 + y^{2^{m-1}}\beta + 1 = 0 \quad (3.19)$$

$$\beta^{2^m+1} = 1 \quad (3.20)$$

Since (3.19) is of degree 3, it can have no root, one root or three roots. Suppose (3.19) has three distinct roots:  $\beta_1$ ,  $\beta_2$  and  $\beta_3$ . Furthermore, suppose  $\beta_1$  and  $\beta_2$  satisfy (3.20). Then, since  $\beta_1\beta_2\beta_3 = 1$ ,

$$\begin{aligned} \beta_3^{2^m+1} &= ((\beta_1\beta_2)^{-1})^{2^m+1} = (\beta_1^{2^m}\beta_2^{2^m})^{2^m+1} \\ &= (\beta_1^{2^m+1})^{2^m}(\beta_2^{2^m+1})^{2^m} = 1. \end{aligned}$$

This says that  $\beta_3$  also satisfies (3.20). Therefore,  $N_\beta(y) = 2$  if and only if there exist one repeated root of multiplicity 2 and another root.

Let  $N_1$  be the number of times  $N_\beta(y) = 1$  as  $y$  ranges over  $GF(2^{2m})$ .

Let  $f(\beta) = \beta^3 + y^{1/2}\beta^2 + y^{2^{m-1}}\beta + 1$  and  $f'(\beta)$  be its formal derivative. From Lemma 3-4,  $g(\beta)$  is a repeated factor of  $f(\beta)$  if and only if  $g(\beta)$  divides  $\text{GCD}(f(\beta), f'(\beta))$ .

$$f'(\beta) = \beta^2 + y^{2^m-1}$$

By the Euclid's algorithm, it is easy to show that

$$\text{GCD}(f(\beta), f'(\beta)) = \text{constant if } y^{2^m+1} \neq 1 \text{ and}$$

$$\text{GCD}(f(\beta), f'(\beta)) = \beta^2 + y^{2^m-1} \text{ if } y^{2^m+1} = 1.$$

Hence,  $f(\beta)$  has a repeated factor if and only if  $y^{2^m+1} = 1$ .

Now, let  $y^{2^m+1} = 1$  and factor  $f(\beta)$ .

$$f(\beta) = (\beta + y^{\frac{1}{2}})(\beta + y^{2^m-2})^2$$

Note that  $f(\beta) = 0$  has a repeated root of multiplicity 3

when  $y^{\frac{1}{2}} = y^{2^m-2}$  or

$$y = y^{2^m-1} \quad (3.21)$$

(3.21) says  $y \in \text{GF}(2^{m-1})$ . Since  $y \in \text{GF}(2^{2m})$  and

$\text{GCD}(2m, m-1) = 1$ ,  $y \in \text{GF}(2^{2m}) \cap \text{GF}(2^{m-1}) = \text{GF}(2)$ . Hence,

$f(\beta) = 0$  has a repeated root of multiplicity 3 if and only if  $y = 1$ .

Therefore,  $f(\beta) = 0$  has one repeated root of multiplicity 2

$$\beta_1 = y^{2^m-2} \quad (3.22)$$

and another root

$$\beta_2 = y^{\frac{1}{2}} \quad (3.23)$$

if and only if  $y^{2^m+1} = 1$ ,  $y \neq 1$ .

The root given by (3.22) satisfies (3.20) since

$y^{2^m+1} = 1$ . The root given by (3.23) satisfies (3.20) since

$y^{2^m+1} = 1$  and 2 does not divide  $2^m+1$ .

Therefore,  $N_\beta(y) = 2$  and  $\Delta_r(y) = 2^m$  if and only if  $y^{2^m+1} = 1$ ,  $y \neq 1$ . Since there exist  $2^m+1$   $(2^m+1)$ -st roots of unity including a unit element in  $GF(2^{2m})$ ,  $N_2 = 2^m$ .

Using Lemma 2-5, the following equalities must hold.

$$2^{m+1} \cdot N_3 + 2^m 2^m - 2^m \cdot N_0 = 2^{2m} \quad (3.24)$$

$$2^{2m+2} \cdot N_3 + 2^{2m} 2^m + 2^{2m} \cdot N_0 = 2^{4m} \quad (3.25)$$

Dividing (3.24) by  $2^m$  and (3.25) by  $2^{2m}$ , we obtain

$$2 \cdot N_3 - N_0 = 0 \quad (3.26)$$

$$4 \cdot N_3 + N_0 = 2^{2m} - 2^m \quad (3.27)$$

Solving for  $N_0$  and  $N_3$  in (3.26) and (3.27),

$$N_0 = (2^{2m} - 2^m)/3$$

$$N_3 = (2^{2m-1} - 2^{m-1})/3$$

Since  $N_0 + N_1 + N_2 + N_3 = 2^{2m}$ ,

$$\begin{aligned} N_1 &= 2^{2m} - (2^{2m} - 2^m)/3 - 2^m - (2^{2m-1} - 2^{m-1})/3 \\ &= 2^{2m-1} - 2^{m-1} \end{aligned} \quad \text{QED}$$

### THEOREM 3-7:

If  $n = 4m$  and  $r = (2^m-1)(2^{2m}+1) + 2$ ,  $\Delta_r(y)$  is a 4-valued function.  $\Delta_r$  and  $\eta(\Delta_r)$  are given by:

| $\Delta_r$ | $\eta(\Delta_r)$      |
|------------|-----------------------|
| $2^{3m}$   | $2^m$                 |
| $2^{2m}$   | $2^{4m-1} - 2^{3m-1}$ |
| 0          | $2^{3m} - 2^m$        |
| $-2^{2m}$  | $2^{4m-1} - 2^{3m-1}$ |

proof:

From (2.4),  $\text{GCD}(r, 2^n - 1) = 1$ .

$$r \equiv 2 \pmod{2^{2m} + 1}$$

$$\begin{aligned} r &= (2^m - 1)\{(2^m + 1)(2^m - 1) + 2\} + 2 \\ &= (2^{2m} - 1)(2^m - 1) + 2^{m+1} \\ &\equiv 2^{m+1} \pmod{2^{2m} - 1} \end{aligned}$$

With  $k = m+1$ ,  $s = 2$  and  $n = 4m$ , (3.9) becomes:

$$y\beta + \beta^{2 \cdot 2^{-m-1}} + y^{2^{2m}} \beta^{-1} + \beta^{-2 \cdot 2^{-m-1}} = 0.$$

$$\text{Or } y\beta + \beta^{2^{-m}} + y^{2^m} \beta^{-1} + \beta^{-2^{-m}} = 0 \quad (3.28)$$

(3.10) becomes:

$$\beta^{2^{2m}+1} = 1$$

which implies

$$\beta^{-2^{-m}} = \beta^{2^m} \text{ and } \beta^{2^{-m}} = \beta^{-2^m} \quad (3.29)$$

Using (3.29),  $N_\beta(y)$  is the number of distinct solutions  $\beta$  in  $\text{GF}(2^{4m})$  to (3.30) and (3.31)

$$y\beta + y^{2^{2m}} \beta^{-1} + (\beta + \beta^{-1})^{2^m} = 0 \quad (3.30)$$

$$\beta^{2^{2m}+1} = 1 \quad (3.31)$$

First, consider the case  $y \in \text{GF}(2^{2m})$ . (3.30) becomes:

$$y(\beta + \beta^{-1}) = (\beta + \beta^{-1})^{2^m} \quad (3.32)$$

Clearly  $\beta = 1$  satisfies (3.32) and (3.31). Hence, we now consider the equation:

$$y = (\beta + \beta^{-1})^{2^m - 1}, \quad y \in \text{GF}(2^{2m}) \quad (3.33)$$

We will show that

1. There exists no solution to (3.33) and (3.31) if  $y = 1$  or  $y^{2^m+1} \neq 1$ .

2. There exist  $2^m$  distinct solutions to (3.33) and (3.31) if  $y^{2^m+1} = 1$ ,  $y \neq 1$ .

Raising (3.33) to the power  $2^m+1$ ,

$$y^{2^m+1} = (\beta + \beta^{-1})^{2^{2m}-1} = 1.$$

Hence, there is no solution to (3.33) if  $y^{2^m+1} \neq 1$ .

Next, let  $y = 1$  in (3.33).

$$1 = (\beta + \beta^{-1})^{2^m-1} \quad (3.34)$$

This says that  $(\beta + \beta^{-1}) = \delta$  for some  $\delta \in GF(2^m)$ ,  $\delta \neq 0$ .

Multiplying  $\beta$  on both sides,

$$\beta^2 + \delta\beta + 1 = 0. \quad (3.35)$$

Transform (3.35) by introducing a new variable  $\omega$ :

$$\omega = \delta^{-1}\beta$$

Then  $\beta = \delta\omega$  and (3.35) becomes:

$$\delta^2\omega^2 + \delta^2\omega + 1 = 0.$$

$$\text{Or } \omega^2 + \omega = \delta^{-2} \quad (3.36)$$

Raising both sides of (3.36) to the power  $2^i$  and adding

$i = 0, 1, \dots, 2m-1$ , we obtain:

$$\sum_{i=0}^{2m-1} (\omega^2 + \omega)^{2^i} = \sum_{i=0}^{2m-1} (\delta^{-2})^{2^i}. \quad (3.37)$$

(3.37) reduces to:

$$\omega^{2^{2m}} + \omega = \text{Tr}_1^{2m}(\delta^{-2}).$$

Note that

$$\text{Tr}_1^{2m}(\delta^{-2}) = \text{Tr}_1^{2m}(\delta^{-1}) = 2 \cdot \text{Tr}_1^m(\delta^{-1}) = 0$$

since  $\delta^{-1} \in GF(2^m)$ .

Hence,  $\omega \in GF(2^{2m})$ . Since  $\beta = \delta\omega$ , this implies that

$\beta \in GF(2^{2m})$ . However,  $\beta = 1$  is the only solution to (3.31) in  $GF(2^{2m})$ . Clearly  $\beta = 1$  is not a solution to (3.34).

Hence, there is no solution to (3.34) and (3.31).

Now, suppose  $y^{2^m+1} = 1$ ,  $y \neq 1$ .

Let  $y = \omega^{2^m-1}$ ,  $\omega \neq 0$ ,  $\omega \in GF(2^{2m})$ . Furthermore, suppose  $y = \omega^{2^m-1}$  satisfies (3.33). Then,  $y = \delta\omega^{2^m-1}$  satisfies (3.33) for  $\delta \in GF(2^m)$ ,  $\delta \neq 0$ . Now consider the equation:

$$\beta + \beta^{-1} = \delta\omega. \quad (3.38)$$

Multiplying (3.38) by  $\beta$  on both sides,

$$\beta^2 + \delta\omega\beta + 1 = 0. \quad (3.39)$$

Setting  $\beta = (\delta\omega)\lambda$ , transform (3.39) to:

$$(\delta\omega)^2\lambda^2 + (\delta\omega)^2\lambda + 1 = 0.$$

$$\text{Or } \lambda^2 + \lambda = (\delta\omega)^{-2} \quad (3.40)$$

From Lemma 3-3, (3.40) has 2 solutions for  $\lambda$  and therefore 2 solutions for  $\beta$  to (3.39) and (3.31) if and only if

$$\text{Tr}_1^{2m}\{(\delta\omega)^{-2}\} = \text{Tr}_1^{2m}\{(\delta\omega)^{-1}\} = 1$$

$$\text{and } \text{Tr}_1^{4m}\{(\delta\omega)^{-2}\} = 0.$$

Clearly  $\text{Tr}_1^{4m}\{(\delta\omega)^{-2}\} = 0$  for all  $\delta \in GF(2^m)$  and for all  $\omega \in GF(2^{2m})$ .

Note that

$$\text{Tr}_1^{2m}\{(\delta^{-1}\omega^{-1})\} = \begin{cases} 0 & \text{for } 2^{m-1}-1 \text{ choices of } \delta \\ 1 & \text{for } 2^{m-1} \text{ choices of } \delta \end{cases}$$

$$\text{or } \text{Tr}_1^{2m}\{(\delta^{-1}\omega^{-1})\} = 0 \text{ identically.}$$

Since  $\text{Tr}_1^{2m}\{(\delta^{-1}\omega^{-1})\} = \text{Tr}_1^m\{\delta^{-1}(\omega^{-1} + \omega^{-2^m})\}$ ,  $\text{Tr}_1^{2m}\{(\delta^{-1}\omega^{-1})\}$

= 0 identically if and only if  $\omega^{-1} = \omega^{-2^m}$  or  $\omega = \omega^{2^m}$ . But

$\omega^{2^m} = \omega$  implies  $y = \omega^{2^m-1} = 1$ . However, the case  $y = 1$  is being excluded. Therefore, there exist  $2 \cdot 2^{m-1} = 2^m$  solutions to (3.33) if  $y^{2^m+1} = 1$ ,  $y \neq 1$ .

To show that solutions of (3.39) satisfy (3.31), suppose  $\text{Tr}_1^{2^m}\{(\delta\omega)^{-1}\} = 1$ . Repeatedly raising (3.40) to the power  $2^i$  and adding  $i = 0, 1, \dots, 2^m-1$ , we obtain:

$$\sum_{i=0}^{2^m-1} (\lambda^{2^i} + \lambda)^{2^i} = \sum_{i=0}^{2^m-1} \{(\delta\omega)^{-2^i}\}^{2^i} \quad (3.41)$$

(3.41) reduces to:

$$\lambda^{2^{2^m}} + \lambda = \text{Tr}_1^{2^m}\{(\delta\omega)^{-1}\} = 1 \quad (3.42)$$

Since  $\beta = (\delta\omega)\lambda$ ,  $\beta^{2^{2^m}} = (\delta\omega)^{2^{2^m}} \lambda^{2^{2^m}} = (\delta\omega)\lambda^{2^{2^m}}$ . Using (3.40) and (3.42),

$$\beta^{2^{2^m+1}} = (\delta\omega)^2 \lambda(\lambda + 1) = (\lambda^2 + \lambda)^{-1}(\lambda^2 + \lambda) = 1.$$

Hence, for  $y \in \text{GF}(2^{2^m})$

$N_\beta(y) = 2^{m+1}$  and  $\Delta_r(y) = 2^{3m}$  if  $y^{2^m+1} = 1$ ,  $y \neq 1$  and  $N_\beta(y) = 1$  and  $\Delta_r(y) = 0$  otherwise.

In order to finish the proof of Theorem 3-7, we use the result due to Welch.

The system of two equations

$$\begin{aligned} y\beta + y^{2^{2^m}}\beta^{-1} + (\beta + \beta^{-1})^{2^m} &= 0 \\ \beta^{2^{2^m+1}} &= 1 \end{aligned}$$

where

$$y \in \text{GF}(2^{4m}) - \text{GF}(2^{2m})$$

has at most 2 solutions for  $\beta$  in  $\text{GF}(2^{4m})$ .

Again let  $N_1$  denote the number of times  $N_\beta(y) = 1$  as  $y$  ranges over  $GF(2^{4m})$ . Then, from Lemma 2-5, we must have:

$$2^{3m} \cdot 2^m + 2^{2m} \cdot N_2 - 2^{2m} \cdot N_0 = 2^{4m}$$

$$2^{6m} \cdot 2^m + 2^{4m} \cdot N_2 + 2^{4m} \cdot N_0 = 2^{8m}$$

Solving for  $N_2$  and  $N_0$ ,

$$N_2 = N_0 = 2^{4m-1} - 2^{3m-1}.$$

Hence,

$$N_1 = 2^{4m} - (2^m + N_2 + N_0) = 2^{3m} - 2^m. \quad \text{QED}$$

### THEOREM 3-8:

If  $n = 2m$  and  $r = 2^m + 3$ ,  $\Delta_r(y)$  is at most a 5-valued function.

$$\Delta_r(y) = 2^m(j-1), \quad 0 \leq j \leq 4.$$

proof:

From (2.5),  $\text{GCD}(r, 2^n - 1) = 1$ .

$$r \equiv 2^2 \pmod{2^m - 1}$$

$$r \equiv 2 \pmod{2^m + 1}$$

With  $k = 2$ ,  $s = 2$  and  $n = 2m$ , (3.9) becomes:

$$y\beta + \beta^{2 \cdot 2^{-2}} + y^{2^m} \beta^{-1} + \beta^{-2 \cdot 2^{-2}} = 0.$$

Raising to the power  $2^1$ ,

$$y^2 \beta^2 + \beta + y^{2^{m+1}} \beta^{-2} + \beta^{-1} = 0.$$

Multiplying by  $\beta^2$ ,  $N_\beta(y)$  becomes the number of distinct solutions  $\beta$  in  $GF(2^{2m})$  to (3.43) and (3.44):

$$y^2 \beta^4 + \beta^3 + \beta + y^{2^{m+1}} = 0 \quad (3.43)$$

$$\beta^{2^m+1} = 1 \quad (3.44)$$

Hence,  $\Delta_r(y) = 2^m \{N_\beta(y) - 1\}$  where  $N_\beta(y) = 0, 1, 2, 3$  or  $4$ .

QED

A further analysis can be made for this case.

First, consider the case  $y \in GF(2^m)$ . Since  $y^{2^m} = y$ ,

(3.43) becomes:

$$y^2 \beta^4 + \beta^3 + \beta + y^2 = 0. \quad (3.45)$$

(3.45) can be factored to become:

$$y^2(\beta + 1)^2(\beta^2 + y^{-2}\beta + 1) = 0. \quad (3.46)$$

Clearly  $\beta = 1$  is a solution to (3.46) and (3.44). Now

consider the second factor of (3.46):

$$\beta^2 + y^{-2}\beta + 1 = 0 \quad (3.47)$$

(3.47) has either no root or two distinct roots. Transform

(3.47) by introducing a new variable  $\omega$ :

$$\omega = y^2 \beta \quad (3.48)$$

Then  $\beta = y^{-2}\omega$  and (3.47) becomes:

$$(y^{-2})^2 \omega^2 + (y^{-2})^2 \omega + 1 = 0. \quad (3.49)$$

Multiplying (3.49) by  $y^4$ , we obtain:

$$\omega^2 + \omega = y^4 \quad (3.50)$$

In order to have solutions for  $\omega$  in (3.50) and hence

solutions for  $\beta$  in (3.47) which also satisfy (3.44), we

must have:

$$\text{Tr}_1^m(y^4) = \text{Tr}_1^m(y) = 1$$

$$\text{and } \text{Tr}_1^{2m}(y^4) = \text{Tr}_1^{2m}(y) = 0.$$

This follows again from Lemma 3-3, for if  $\text{Tr}_1^m(y) = 0$ , two roots for  $\omega$  in (3.50) belong to  $\text{GF}(2^m)$  and hence two roots for  $\beta$  in (3.47) belong to  $\text{GF}(2^m)$ . But (3.44) has only one solution in  $\text{GF}(2^m)$ , namely  $\beta = 1$ . Clearly  $\beta = 1$  is not a root of (3.47).

Since  $\text{Tr}_1^{2^m}(y) = 0$  for all  $y \in \text{GF}(2^m)$ , (3.47) has two distinct roots for  $\beta$  if and only if  $\text{Tr}_1^m(y) = 1$ .

To show that two roots of (3.47) satisfy (3.44), suppose  $\text{Tr}_1^m(y) = 1$ . Repeatedly raising (3.50) to the power  $2^i$  and adding  $i = 0, 1, \dots, m-1$ , we obtain:

$$\sum_{i=0}^{m-1} (\omega^2 + \omega)^{2^i} = \sum_{i=0}^{m-1} (y^4)^{2^i} \quad (3.51)$$

(3.51) reduces to:

$$\omega^{2^m} + \omega = \text{Tr}_1^m(y) = 1 \quad (3.52)$$

Since  $\beta = y^{-2}\omega$ ,  $\beta^{2^m} = (y^{-2})^{2^m}\omega^{2^m} = y^{-2}\omega^{2^m}$ . Using (3.50) and (3.52),

$$\beta^{2^m+1} = (y^{-2})^2 \omega \omega^{2^m} = y^{-4} \omega (\omega + 1) = 1.$$

Therefore, for  $y \in \text{GF}(2^m)$

$$N_\beta(y) = 1 \text{ when } \text{Tr}_1^m(y) = 0 \quad (3.53)$$

$$\text{and } N_\beta(y) = 3 \text{ when } \text{Tr}_1^m(y) = 1.$$

Next, consider the case  $y \in \text{GF}(2^{2m}) - \text{GF}(2^m)$ .

Multiplying (3.43) by  $y^{-2}$ , we obtain:

$$\beta^4 + y^{-2}\beta^3 + y^{-2}\beta + y^{2(2^m-1)} = 0 \quad (3.54)$$

Suppose (3.54) has 4 distinct roots:  $\beta_1, \beta_2, \beta_3$  and  $\beta_4$ .

From (3.54) we have

$$\beta_1 \beta_2 \beta_3 \beta_4 = y^{2(2^m-1)}$$

$$\text{or } \beta_4 = (\beta_1 \beta_2 \beta_3)^{-1} \cdot y^{2(2^m-1)}$$

Furthermore, suppose  $\beta_1, \beta_2$  and  $\beta_3$  satisfy (3.44). Then,

$$(\beta_1 \beta_2 \beta_3)^{-1} = (\beta_1 \beta_2 \beta_3)^{2^m}$$

$$\text{and } \beta_4^{2^m+1} = (\beta_1^{2^m+1} \cdot \beta_2^{2^m+1} \cdot \beta_3^{2^m+1})^{2^m} \cdot y^{2(2^{2m}-1)} = 1$$

Hence,  $\beta_4$  also satisfies (3.44). This implies that if  $N_\beta(y) = 3$ , (3.54) must have 2 distinct roots and one repeated root of multiplicity 2.

Assume (3.54) has a repeated root  $\beta_1$ . Dividing (3.54) by  $(\beta + \beta_1)^2$ , we obtain:

$$\begin{aligned} & (\beta^2 + y^{-2}\beta + \beta_1^2)(\beta + \beta_1)^2 + (\beta_1^2 + 1)y^{-2}\beta \\ & + y^{2(2^m-1)} + \beta_1^4 = 0 \end{aligned}$$

Hence, we must have

$$(\beta_1^2 + 1) = 0 \tag{3.55}$$

$$y^{2(2^m-1)} = \beta_1^4 \tag{3.56}$$

(3.55) says  $\beta_1 = 1$ , which in turn implies  $y^{2^m-1} = 1$  from (3.56). This says  $y \in \text{GF}(2^m)$ .

Hence, if  $y \in \text{GF}(2^{2m}) - \text{GF}(2^m)$ ,

$$N_\beta(y) \neq 3. \tag{3.57}$$

Finally combining (3.53) and (3.57), we obtain:

$$N_{\beta}(y) = 3 \text{ and } \Delta_r(y) = 2^{m+1}$$

if and only if  $y \in \text{GF}(2^m)$  and  $\text{Tr}_1^m(y) = 1$ .

For the case  $y \in \text{GF}(2^{2m}) - \text{GF}(2^m)$ , (3.54) can be transformed to a more standard form by introducing a new variable  $\sigma$ :

$$\sigma = (\beta + 1)^{-1} \quad (3.58)$$

Then, the number of distinct solutions  $\beta$  to (3.54) and (3.44) is equal to the number of distinct solutions  $\sigma$  in  $\text{GF}(2^{2m})$  to the system of two equations:

$$\sigma^4 + \lambda \sigma^2 + \lambda \sigma + \lambda y^2 = 0$$

$$\sigma^{2^m} + \sigma = 1$$

where

$$\lambda = (y + y^{2^m})^{-2}, \quad y \in \text{GF}(2^{2m}) - \text{GF}(2^m).$$

Furthermore, it is conjectured that the following distribution holds for  $m > 2$ .

| $\Delta_r$    | $n(\Delta_r)$ for $m$ odd      | $n(\Delta_r)$ for $m$ even     |
|---------------|--------------------------------|--------------------------------|
| $3 \cdot 2^m$ | $(2^{2m-3} - 2^{m-2})/3$       | $(2^{2m-3} - 2^{m-1})/3$       |
| $2^{m+1}$     | $2^{m-1}$                      | $2^{m-1}$                      |
| $2^m$         | $2^{2m-2} - 2^{m-1}$           | $2^{2m-2}$                     |
| 0             | $(2^{2m} + 5 \cdot 2^{m-1})/3$ | $(2^{2m} + 2^{m-1})/3$         |
| $-2^m$        | $3(2^{2m-3} - 2^{m-2})$        | $(3 \cdot 2^{2m-3} - 2^{m-1})$ |

REMARK: If  $n = 2m$  and  $r = 2^{2m-1} - 2^{m-1} + 1$ ,  $\Delta_r(y)$  is at most a 5-valued function. For the proof, let  $r' = 2^m + 3$ . Then,

$$\begin{aligned} 2 \cdot r \cdot r' &= (2^{2m} - 2^m + 2)(2^m + 3) \\ &\equiv (3 - 2^m)(3 + 2^m) \pmod{2^{2m}-1} \\ &\equiv 2^3 \pmod{2^{2m}-1}. \end{aligned}$$

The desired result follows directly from Lemma 2-3 and Theorem 3-8.

As an alternative proof, consider the following. For this decimation  $r$ , it is easy to show that  $r \equiv 1 \pmod{2^m-1}$ ,  $r \equiv 2 \pmod{2^m+1}$  and hence  $\text{GCD}(r, 2^{2m}-1) = 1$ . With  $k = 0$ ,  $s = 2$  and  $n = 2m$ , (3.9) becomes:

$$y\beta + \beta^2 + y^{2^m}\beta^{-1} + \beta^{-2} = 0$$

Multiplying the above by  $\beta^2$ , we get:

$$\beta^4 + y\beta^3 + y^{2^m}\beta + 1 = 0$$

Therefore,  $\Delta_r(y) = 2^m(j-1)$ ,  $0 \leq j \leq 4$ .

#### THEOREM 3-9;

If  $n \equiv 2 \pmod{4}$ ,  $n = 2m$  and  $r = 2^m + 2^{m-1} - 1$ ,  $\Delta_r(y)$  is at most a 6-valued function.

$$\Delta_r(y) = 2^m(j-1), \quad 0 \leq j \leq 5.$$

#### proof:

From (2.6),  $\text{GCD}(r, 2^n-1) = 1$ .

$$r \equiv 2^{m-1} \pmod{2^m-1}$$

$$r \equiv 2^{m-1} - 2 \pmod{2^m+1}$$

With  $k = m-1$ ,  $s = 2^{m-1}-2$  and  $n = 2m$ , (3.9) becomes:

$$\begin{aligned} & y\beta + \beta^{(2^{m-1}-2)2^{-m+1}} + y^{2^m}\beta^{-1} + \beta^{-(2^{m-1}-2)2^{-m+1}} \\ &= y\beta + \beta^{1-2^{-m+2}} + y^{2^m}\beta^{-1} + \beta^{-(1-2^{-m+2})} = 0 \end{aligned}$$

Since  $\beta^{2^m} = \beta^{-1}$  and  $\beta = \beta^{-2^m}$ ,

$$\begin{aligned} & y\beta + \beta(\beta^{2^m})^{2^{-m+2}} + y^{2^m}\beta^{-1} + \beta^{-1}(\beta^{-2^m})^{2^{-m+2}} \\ &= y\beta + \beta \cdot \beta^{2^2} + y^{2^m}\beta^{-1} + \beta^{-1} \cdot \beta^{-2^2} \\ &= y\beta + \beta^5 + y^{2^m}\beta^{-1} + \beta^{-5} = 0 \end{aligned}$$

Multiplying by  $\beta^5$ ,

$$\beta^{10} + y\beta^6 + y^{2^m}\beta^4 + 1 = 0$$

Raising to the power  $2^{-1}$ ,

$$\beta^5 + y^{1/2}\beta^3 + y^{2^{m-1}}\beta^2 + 1 = 0$$

QED

### THEOREM 3-10:

If  $n = 2m$  and  $r = 2^{m+2} - 3$ ,  $\Delta_r(y)$  is at most an 8-valued function.

$$\Delta_r(y) = 2^m(j-1), \quad 0 \leq j \leq 7.$$

proof:

From (2.7),  $\text{GCD}(r, 2^n-1) = 1$ .

$$r = 4(2^m-1) + 1 \equiv 1 \pmod{2^m-1}$$

$$r = 4(2^m+1) - 7 \equiv -7 \pmod{2^m+1}$$

With  $k = 0$ ,  $s = -7$  and  $n = 2m$ , (3.9) becomes:

$$y\beta + \beta^{-7} + y^{2^m}\beta^{-1} + \beta^7 = 0$$

Multiplying by  $\beta^7$ ,

$$\beta^{14} + y\beta^8 + y^{2^m}\beta^6 + 1 = 0$$

Raising to the power  $2^{-1}$ ,

$$\beta^7 + y^{\frac{1}{2}}\beta^4 + y^{2^{m-1}}\beta^3 + 1 = 0$$

QED

REMARK: Computed results indicate that for  $n = 2m = 4m'$  and  $r = 2^{m+2} - 3$ ,  $\Delta_r(y)$  is at most 5-valued.

$$\Delta_r(y) = 2^m(j-1), \quad j = 0, 1, 2, 3 \text{ or } 5.$$

REMARK: When  $n = 8$ , a further improvement can be made on Theorem 3-10.

$$r = 2^{4+2} - 3 = 61$$

$$r^{-1} = 23 \cdot 2$$

Let  $r' = 23$ . Then  $r' \equiv 2^3 \pmod{15}$  and  $r' \equiv 6 \pmod{17}$ .

With  $k = 3$ ,  $s = 6$  and  $n = 2 \cdot 4$ , (3.9) becomes:

$$y\beta + \beta^{6 \cdot 2^{-3}} + y^{2^4}\beta^{-1} + \beta^{-6 \cdot 2^{-3}} = 0$$

$$y\beta + \beta^{3 \cdot 2^{-2}} + y^{2^4}\beta^{-1} + \beta^{-3 \cdot 2^{-2}} = 0$$

Since  $\beta^{17} = 1$ ,  $\beta^{-1} = \beta^{2^4}$  and  $\beta = \beta^{-2^4}$ .

$$y\beta + (\beta^{-2^4})^{3 \cdot 2^{-2}} + y^{2^4}\beta^{-1} + (\beta^{2^4})^{3 \cdot 2^{-2}} = 0$$

$$y\beta + \beta^{-12} + y^{16}\beta^{-1} + \beta^{12} = 0$$

$$y\beta + \beta^5 + y^{16}\beta^{-1} + \beta^{-5} = 0$$

Multiplying by  $\beta^5$  and then raising to the power  $2^{-1}$ ,

$$\beta^5 + y^{\frac{1}{2}}\beta^3 + y^8\beta^2 + 1 = 0$$

Hence,

$$\Delta_{23}(y^{-61}) = \Delta_{61}(y) = 2^4(j-1), \quad 0 \leq j \leq 5.$$

THEOREM 3-11:

If  $n = 2m$  and  $r = 2^{m+2} + 2^m - 3$ ,  $\Delta_r(y)$  is at most a 9-valued function.

$$\Delta_r(y) = 2^m(j-1), \quad 0 \leq j \leq 8.$$

proof:

From (2.8),  $\text{GCD}(r, 2^n - 1) = 1$ .

$$r = 5(2^m - 1) + 2 \equiv 2 \pmod{2^m - 1}$$

$$r = 5(2^m + 1) - 8 \equiv -2^3 \pmod{2^m + 1}$$

With  $k = 1$ ,  $s = -2^3$  and  $n = 2m$ , (3.9) becomes:

$$y\beta + \beta^{-2^3 \cdot 2^{-1}} + y^{2^m} \beta^{-1} + \beta^{2^3 \cdot 2^{-1}} = 0$$

$$y\beta + \beta^{-4} + y^{2^m} \beta^{-1} + \beta^4 = 0$$

Multiplying by  $\beta^4$ ,

$$\beta^8 + y\beta^5 + y^{2^m} \beta^3 + 1 = 0$$

QED

## 3. COMPUTED RESULTS

Table 3-1 gives the cyclotomic coset leaders given by (3.1) through (3.6), which are considered in Theorems 3-6 through 3-11 respectively. Table 3-2 gives the actual values of  $N_\beta(y) = \{\Delta_r(y)/2^m + 1\}$ . It is observed that the bounds given by Theorems 3-8 and 3-9 are tight whereas the bounds given by Theorems 3-10 and 3-11 are not.

For  $n = 12$  and  $n = 14$ , there exist many other decimations which are not covered by 6 theorems of the previous section. Degrees of (3.9) that result from those decimations are found to be high. However, the computed

results indicate that values which  $\Delta_r(y)$  takes on are restricted. With 2 exceptions,  $r = 331$  and  $r = 631$  for  $n = 12$ ,  $\Delta_r(y)$  for which  $r \equiv 2^k \pmod{2^{n/2}-1}$  are seen to be 7-valued or less. Table 3-3 lists all decimations of this type for  $n = 4$  through  $n = 14$ . The number in parentheses following decimation  $r$  is the number of distinct values that  $\Delta_r(y)$  takes on.

## CYCLOTOMIC COSET LEADERS GIVEN BY (3.1) THROUGH (3.6)

| n     | 4 | 6  | 8  | 10  | 12  | 14  | 16   |
|-------|---|----|----|-----|-----|-----|------|
| $r_1$ | 7 |    | 31 |     | 127 |     | 511  |
| $r_2$ | 7 |    | 53 |     | 457 |     | 3857 |
| $r_3$ | 7 | 11 | 19 | 35  | 67  | 131 | 259  |
| $r_4$ |   | 11 |    | 47  |     | 191 |      |
| $r_5$ | 7 | 23 | 61 | 125 | 253 | 509 | 1021 |
| $r_6$ |   | 11 | 53 | 157 | 317 | 637 | 1277 |

TABLE 3-1

NUMBER OF SOLUTIONS TO (3.9) and (3.10),  $N_\beta(y)$ 

|       |   |   |   |             |   |                |                     |                  |
|-------|---|---|---|-------------|---|----------------|---------------------|------------------|
| $r_1$ | 0 | 1 | 2 | 3           |   |                |                     |                  |
| $r_2$ | 0 | 1 | 2 | $2^{n/4}+1$ |   |                |                     |                  |
| $r_3$ | 0 | 1 | 2 | 3           | 4 | for $n \geq 6$ |                     |                  |
| $r_4$ | 0 | 1 | 2 | 3           | 4 | for $n = 6$    |                     |                  |
|       | 0 | 1 | 2 | 3           | 4 | 5              | for $n = 10, 14$    |                  |
| $r_5$ | 0 | 1 | 2 | 3           |   |                | for $n = 4$         |                  |
|       | 0 | 1 | 2 | 3           | 4 |                | for $n = 6$         |                  |
|       | 0 | 1 | 2 | 3           | 5 |                | for $n = 8, 12, 16$ |                  |
|       | 0 | 1 | 2 | 3           | 4 | 5              | for $n = 10, 14$    |                  |
| $r_6$ | 0 | 1 | 2 | 3           | 4 |                | for $n = 6, 10$     |                  |
|       | 0 | 1 | 2 | 5           |   |                | for $n = 8$         |                  |
|       | 0 | 1 | 2 | 3           | 4 | 5              | 6                   | for $n = 12, 14$ |

TABLE 3-2

DECIMATIONS OF TYPE  $2^k \text{ MOD } 2^{n/2}-1$ 

|        |         |         |         |         |         |
|--------|---------|---------|---------|---------|---------|
| n = 4  | 7(4)    |         |         |         |         |
| n = 6  | 11(5)   | 23(5)   |         |         |         |
| n = 8  | 19(5)   | 23(5)   | 31(4)   | 47(5)   | 53(4)   |
|        | 61(5)   | 91(4)   |         |         |         |
| n = 10 | 35(5)   | 47(6)   | 95(5)   | 101(5)  | 109(6)  |
|        | 125(6)  | 157(5)  | 221(6)  | 343(6)  |         |
| n = 12 | 67(5)   | 71(6)   | 79(7)   | 127(4)  | 191(5)  |
|        | 197(7)  | 253(5)  | 317(7)  | 319(7)  | 331(8)  |
|        | 347(7)  | 379(7)  | 443(7)  | 457(4)  | 473(7)  |
|        | 599(5)  | 631(8)  | 701(6)  | 757(6)  | 821(7)  |
|        | 823(7)  | 827(7)  | 1387(4) |         |         |
| n = 14 | 131(5)  | 143(7)  | 191(6)  | 383(5)  | 389(7)  |
|        | 397(6)  | 413(6)  | 445(6)  | 509(6)  | 637(7)  |
|        | 667(7)  | 763(6)  | 893(7)  | 905(7)  | 953(6)  |
|        | 1145(7) | 1147(7) | 1151(6) | 1175(6) | 1207(6) |
|        | 1271(6) | 1399(6) | 1405(6) | 1429(6) | 1525(7) |
|        | 1655(7) | 1715(6) | 1907(6) | 1909(6) | 1913(7) |
|        | 2429(6) | 2477(7) | 2669(6) | 2671(7) | 2675(7) |
|        | 2683(6) | 2731(6) | 2923(7) | 3431(6) | 3437(7) |
|        | 3445(6) |         |         |         |         |

TABLE 3-3

## CHAPTER IV

### MULTI-VALUED CROSS-CORRELATION FUNCTIONS II

1.  $\Delta_r(y)$  FOR  $r = (2^{mk}+1)/(2^k+1)$

In this chapter we consider  $\Delta_r(y)$  for the case

$$r = (2^{mk}+1)/(2^k+1), \text{ m and } n/\text{GCD}(n,k) \text{ both odd.}$$

From Lemma 2-8,  $\text{GCD}(r, 2^n-1) = 1$ . It can be shown that

$\Delta_r(y)$  is restricted to the form:

$$\Delta_r(y) = 0, \pm 2^{(n+de)/2}$$

where

$$e = \text{GCD}(n,k)$$

and  $d$  is a some positive odd integer.

$$\text{Given } r = (2^{mk}+1)/(2^k+1)$$

$$= 2^{(m-1)k} - 2^{(m-2)k} + \dots - 2^k + 1,$$

$$\text{consider } q = \{2^{m(n-k)}+1\}/\{2^{(n-k)}+1\}$$

$$= 2^{-(m-1)k} \{2^{(m-1)n} - 2^{k_2(m-2)n} + 2^{2k_2(m-3)n} \\ - \dots - 2^{(m-2)k_2n} + 2^{(m-1)k}\}$$

$$\equiv 2^{-(m-1)k} \{1 - 2^k + 2^{2k} - \dots + 2^{(m-1)k}\}$$

$$\text{mod } 2^n-1.$$

This says that  $r$  and  $q$  belong to the same proper cyclotomic coset. Hence, to determine  $\Delta_r(y)$  for some  $k$ , it suffices to consider only those  $k$  such that  $2k \leq n$  for  $n$  odd and  $2k \leq (n-2)$  for  $n$  even. Furthermore, note that

$$2^{m+jn} \equiv 2^m \text{ mod } 2^n-1$$

$$\{2^{(n-m)k}+1\}/(2^k+1) \equiv 2^{-mk}(1+2^{mk})/(2^k+1) \text{ mod } 2^n-1$$

Therefore, in determining  $\Delta_r(y)$ , it suffices to consider

$$\begin{aligned}
 & 3 \leq m \leq n/2, \quad 2k \leq n-2 \quad \text{for } n \text{ even} \\
 & \text{and } 3 \leq m \leq n, \quad 2k \leq n \quad \text{for } n \text{ odd.}
 \end{aligned}
 \tag{4.1}$$

The case  $m = 3$  reduces to the Welch's case, Theorem 1-6. The case  $m = n$  odd is considered in Lemma 2-9. The case in which  $3m \equiv \pm 1 \pmod n$  and  $\text{GCD}(3, n) = 1$  is considered in Lemma 2-10.

In CHAPTER II it was shown that if  $\text{GCD}(3, n) = 1$ , the multiplicative inverse of  $r = (2^{3k}+1)/(2^k+1)$  can be given by  $r'' = 2^s(2^{m't}+1)/(2^t+1)$  for some  $s$  where  $m$  and  $t$  are given by (2.9) and (2.10) respectively. Similarly, if  $\text{GCD}(m, n) = 1$ , the multiplicative inverse of  $r = (2^{mk}+1)/(2^k+1) \pmod{2^n-1}$  can be given by

$$r'' = 2^s(2^{m''k''}+1)/(2^{k''}+1) \text{ for some } s$$

where

$$m \cdot m' \equiv 1 \pmod n$$

$$j \equiv m \cdot k \pmod n$$

$$m'' = \begin{cases} n - m' & \text{when } m' \text{ is even} \\ m' & \text{when } m' \text{ is odd} \end{cases}$$

$$\text{and } k'' = \begin{cases} n - j & \text{when } j > (n-1)/2 \\ j & \text{when } j \leq (n-1)/2 \end{cases}$$

For the analysis of  $\Delta_r(y)$  of this type, we follow the same arguments used by Welch when he proved Theorem 1-6.

$$\Delta_r(y) = \sum_{x \in \text{GF}(2^n)} (-1)^{\text{Tr}_1^n \{xy + x^{(2^{mk}+1)/(2^k+1)}\}}$$

From Lemma 2-7,  $\text{GCD}(2^k+1, 2^n-1) = 1$ . Hence, a mapping

$x \rightarrow x^{2^{k+1}}$  permutes elements of  $GF(2^n)$ . Then,

$$\Delta_r(y) = \sum_{x \in GF(2^n)} (-1)^{\text{Tr}_1^n \{ yx^{2^{k+1}} + x^{2^{mk+1}} \}}$$

Let  $e = \text{GCD}(n, k)$ . Then, it is easy to show that

$$\text{Tr}_1^n(x) = \text{Tr}_1^e \{ \text{Tr}_e^n(x) \}$$

Let  $x \in GF(2^n)$ . Then, the element  $x$  can be expressed in the form

$$x = \sum_{i=1}^{n/e} x_i \sigma_i$$

where

$$x_i \in GF(2^e)$$

$$\sigma_i \in GF(2^n)$$

and  $\{\sigma_1, \sigma_2, \dots, \sigma_{n/e}\}$  is the basis of  $GF(2^n)$  over  $GF(2^e)$ .

Then,

$$\begin{aligned} \text{Tr}_1^n(x) &= \text{Tr}_1^n \left\{ \sum_{i=1}^{n/e} x_i \sigma_i \right\} \\ &= \text{Tr}_1^e \left\{ \text{Tr}_e^n \left\{ \sum_{i=1}^{n/e} x_i \sigma_i \right\} \right\} \\ &= \text{Tr}_1^e \left\{ \sum_{i=1}^{n/e} x_i \text{Tr}_e^n(\sigma_i) \right\} \end{aligned}$$

Expanding  $x^{2^{mk+1}}$ ,

$$\begin{aligned} x^{2^{mk+1}} &= \left\{ \sum_{i=1}^{n/e} x_i \sigma_i \right\}^{2^{mk}} \left\{ \sum_{j=1}^{n/e} x_j \sigma_j \right\} \\ &= \left\{ \sum_{i=1}^{n/e} x_i^{2^{mk}} \sigma_i^{2^{mk}} \right\} \sum_{j=1}^{n/e} x_j \sigma_j \\ &= \left\{ \sum_{i=1}^{n/e} x_i \sigma_i^{2^{mk}} \right\} \sum_{j=1}^{n/e} x_j \sigma_j \end{aligned}$$

$$= \sum_{i,j} x_i x_j \sigma_i^{2^{mk}} \sigma_j$$

Define  $Q(x)$  by:

$$Q(x) = \text{Tr}_e^n(yx^{2^k+1} + x^{2^{mk}+1})$$

where

$$y \in \text{GF}(2^n).$$

$$\begin{aligned} Q(x) &= \text{Tr}_e^n\{y \cdot \sum_{i,j} x_i x_j \sigma_i^{2^k} \sigma_j + \sum_{i,j} x_i x_j \sigma_i^{2^{mk}} \sigma_j\} \\ &= \text{Tr}_e^n\{ \sum_{i,j} x_i x_j (y \sigma_i^{2^k} \sigma_j + \sigma_i^{2^{mk}} \sigma_j) \} \\ &= \sum_{i,j} x_i x_j \text{Tr}_e^n(y \sigma_i^{2^k} \sigma_j + \sigma_i^{2^{mk}} \sigma_j) \\ &= \sum_{i,j} x_i x_j \cdot \delta_{ij} \end{aligned}$$

where

$$\delta_{ij} = \text{Tr}_e^n(y \sigma_i^{2^k} \sigma_j + \sigma_i^{2^{mk}} \sigma_j) \in \text{GF}(2^e).$$

Hence,  $Q(x)$  is a quadratic form over  $\text{GF}(2^e)$ .

A quadratic form with coefficients in  $\text{GF}(2^e)$  can be reduced to one of the following 2 canonical forms: [20] [21]

$$\text{Type I: } QF = x_1 x_2 + \dots + x_{2s-1} x_{2s} + x_{2s+1}^2$$

$$\begin{aligned} \text{Type II: } QF_\lambda &= x_1 x_2 + \dots + x_{2s-1} x_{2s} \\ &\quad + \lambda(x_{2s-1}^2 + x_{2s}^2) \end{aligned}$$

where

$$\lambda = 0$$

$$\text{or } \lambda \in \text{GF}(2^e) \text{ and } \text{Tr}_1^e(\lambda) = 1.$$

First, consider  $\Delta_r(y)$  for Type I.

$$\begin{aligned}\Delta_r(y) &= \sum_{x \in GF(2^n)} (-1)^{\text{Tr}_1^e(Q(x))} \\ &= \sum_{x_1 \dots x_{n/e}} \dots \sum_{x_{2s+1}} (-1)^{\text{Tr}_1^e(x_1 x_2 + x_3 x_4 + \dots + x_{2s+1}^2)}\end{aligned}$$

where

$x_i$ 's range over  $GF(2^e)$ .

$$\begin{aligned}\Delta_r(y) &= \sum_{x_1} \sum_{x_2} (-1)^{\text{Tr}_1^e(x_1 x_2)} \sum_{x_3} \sum_{x_4} (-1)^{\text{Tr}_1^e(x_3 x_4)} \dots \\ &\quad \sum_{x_{2s-1}} \sum_{x_{2s}} (-1)^{\text{Tr}_1^e(x_{2s-1} x_{2s})} \sum_{x_{2s+1}} (-1)^{\text{Tr}_1^e(x_{2s+1}^2)} \\ &\quad \sum_{x_{2s+2}} \dots \sum_{x_{n/e}} (1)\end{aligned}$$

From Lemma 1-1,

$$\sum_{x_{2s+1}} (-1)^{\text{Tr}_1^e(x_{2s+1}^2)} = \sum_{x_{2s+1}} (-1)^{\text{Tr}_1^e(x_{2s+1})} = 0$$

Hence, if  $Q(x)$  is of Type I,  $\Delta_r(y) = 0$ .

Next, consider  $\Delta_r(y)$  for Type II.

$$\begin{aligned}\Delta_r(y) &= \sum_{x_1} \sum_{x_2} (-1)^{\text{Tr}_1^e(x_1 x_2)} \sum_{x_3} \sum_{x_4} (-1)^{\text{Tr}_1^e(x_3 x_4)} \dots \\ &\quad \sum_{x_{2s-1}} \sum_{x_{2s}} (-1)^{\text{Tr}_1^e\{x_{2s-1} x_{2s} + \lambda(x_{2s-1}^2 + x_{2s}^2)\}} \\ &\quad \sum_{x_{2s+1}} \dots \sum_{x_{n/e}} (1)\end{aligned}$$

Consider the  $s$ -th sum. By the substitution of

$x_{2s-1} \rightarrow (x_{2s-1} + \lambda^{\frac{1}{2}})$  and  $x_{2s} \rightarrow (x_{2s} + \lambda^{\frac{1}{2}})$ , the exponent of the  $s$ -th sum becomes:

$$\begin{aligned} & \text{Tr}_1^e \{ (x_{2s-1} + \lambda^{\frac{1}{2}})(x_{2s} + \lambda^{\frac{1}{2}}) + \lambda(x_{2s-1}^2 + \lambda + x_{2s}^2 + \lambda) \} \\ &= \text{Tr}_1^e \{ x_{2s-1}x_{2s} + \lambda + \lambda^{\frac{1}{2}}(x_{2s-1} + x_{2s}) + \lambda(x_{2s-1}^2 + x_{2s}^2) \} \\ &= \text{Tr}_1^e (x_{2s-1}x_{2s} + \lambda) \end{aligned}$$

Then the  $s$ -th sum becomes:

$$(-1)^{\text{Tr}_1^e(\lambda)} \sum_{x_{2s-1}} \sum_{x_{2s}} (-1)^{\text{Tr}_1^e(x_{2s-1}x_{2s})}$$

Hence,

$$\begin{aligned} \Delta_r(y) &= (-1)^{\text{Tr}_1^e(\lambda)} \left\{ \sum_{x_1} \sum_{x_2} (-1)^{\text{Tr}_1^e(x_1x_2)} \right\}^s \\ &\quad \sum_{x_{2s+1}} \cdots \sum_{x_{n/e}} (1) \end{aligned}$$

Since

$$\sum_{x_1} \sum_{x_2} (-1)^{\text{Tr}_1^e(x_1x_2)} = 2^e$$

$$\text{and } \sum_{x_{2s+1}} \cdots \sum_{x_{n/e}} (1) = (2^e)^{n/e - 2s} = 2^{n-2es},$$

we have

$$\Delta_r(y) = (-1)^{\text{Tr}_1^e(\lambda)} (2^e)^s \cdot 2^{n-2es} = \pm 2^{n-es}$$

Thus we have:

$$\text{If } Q(x) \text{ is of Type I, } \Delta_r(y) = 0$$

$$\text{If } Q(x) \text{ is of Type II, } \Delta_r(y) = \pm 2^{n-es}$$

(4.2)

where

$2s$  is the rank of  $QF_\lambda$ .

In order to evaluate  $\Delta_r(y)$  explicitly, it suffices to find the rank of the quadratic form over  $GF(2^e)$  of the Type II,  $QF_\lambda$ .

Suppose an element  $z$  in  $GF(2^n)$  has its first  $2s$  coordinates equal to zero in that coordinates which produced the canonical form. Then

$$Q(x + z) = Q(x) + Q(z), \quad x \in GF(2^n).$$

Furthermore, if the quadratic form is of Type II,  $Q(z) = 0$ .

Hence, to find the rank of the quadratic form of Type II, we must find the number of  $z$  in  $GF(2^n)$  such that

$$\begin{aligned} 0 &= Q(x + z) + Q(x) \\ &= \text{Tr}_e^n \{ y(x+z)^{2^k+1} + (x+z)^{2^{mk}+1} \} \\ &\quad + \text{Tr}_e^n \{ yx^{2^k+1} + x^{2^{mk}+1} \} \\ &= \text{Tr}_e^n \{ yx^{2^k+1} + yx^{2^k}z + yxz^{2^k} + yz^{2^k+1} + x^{2^{mk}+1} \\ &\quad + x^{2^{mk}}z + xz^{2^{mk}} + z^{2^{mk}+1} + yx^{2^k+1} + x^{2^{mk}+1} \} \\ &= \text{Tr}_e^n \{ yx^{2^k}z + yxz^{2^k} + x^{2^{mk}}z + xz^{2^{mk}} \} \\ &\quad + \text{Tr}_e^n \{ yz^{2^k+1} + z^{2^{mk}+1} \} \\ &= \text{Tr}_e^n \{ y^{2^{(m-1)k}} x^{2^{mk}} z^{2^{(m-1)k}} + y^{2^{mk}} x^{2^{mk}} z^{2^{(m+1)k}} \\ &\quad + x^{2^{mk}}z + x^{2^{mk}}z^{2^{2mk}} \} + Q(z) \end{aligned}$$

$$= \text{Tr}_e^n \{ x^{2^{mk}} \{ z^{2^{2mk}} + z + y^{2^{mk}} z^{2^{(m+1)k}} + y^{2^{(m-1)k}} z^{2^{(m-1)k}} \} \}$$

Since the above must hold as  $x$  ranges over  $\text{GF}(2^n)$ , we must have

$$z^{2^{2mk}} + z + y^{2^{mk}} z^{2^{(m+1)k}} + y^{2^{(m-1)k}} z^{2^{(m-1)k}} = 0 \quad (4.3)$$

Suppose  $z_1$  is a solution to (4.3). Then, for  $\alpha_1 \in \text{GF}(2^e)$ ,  $z = \alpha_1 z_1$  is also a solution to (4.3) since

$$z^{2^{jk}} = \alpha_1^{2^{jk}} z_1^{2^{jk}} = \alpha_1 \cdot z_1^{2^{jk}} \quad \text{for any } j.$$

If  $z_2$  is another solution to (4.3), then for  $\alpha_2 \in \text{GF}(2^e)$   $z = (\alpha_1 z_1 + \alpha_2 z_2)$  is also a solution to (4.3). This says that the set of solutions to (4.3) is a linear space over  $\text{GF}(2^e)$ . Hence, the number of solutions  $z$  in  $\text{GF}(2^n)$  to (4.3) is of the form  $(2^e)^d$ , where  $d$  is the dimension of the solution space over  $\text{GF}(2^e)$ .

The rank of the quadratic form of Type II over  $\text{GF}(2^e)$ ,  $2s$ , is equal to:

$$2s = n/e - d \quad (4.4)$$

Since  $n/e$  is odd,  $d$  must also be odd.

It can be shown that if  $\{z_1, z_2, \dots, z_{n/e}\}$  is the basis for  $\text{GF}(2^n)$  over  $\text{GF}(2^e)$ , then  $\{z_1, z_2, \dots, z_{n/e}\}$  is also the basis of  $\text{GF}(2^{kn/e})$  over  $\text{GF}(2^{2k})$ .

The set of solutions to (4.3) is also a linear space over  $\text{GF}(2^{2k})$ . In view of the fact that (4.3) is of

degree  $2^{mk} = (2^{2k})^m$ , the dimension of the solution space is at most  $m$ . Hence,

$$1 \leq d \leq m, \quad d \text{ odd} \quad (4.5)$$

Substituting (4.4) into (4.2), we have

$$\Delta_r(y) = \begin{cases} 0 \\ \pm 2^{(n+de)/2}, \end{cases} \quad 1 \leq d \leq m, \quad d \text{ odd} \quad (4.6)$$

In [22] Kasami evaluates the weight distribution of the dual of the cyclic code whose generator polynomial is

$$\prod_{i=0}^{u-1} f_{1+2^{k(2i+1)}}(x).$$

The weight of this code is restricted to the form:

$$2^{n-1} \pm 2^{(n-e)/2 + ie - 1} \quad 1 \leq i \leq u-1$$

where

$$e = \text{GCD}(n, k).$$

The  $(2^n-1, 2n)$  cyclic code generated by

$$\frac{x^{2^n-1} + 1}{x^{2n \cdot f_{1+2^k(1/x) \cdot f_{1+2^{mk}(1/x)}}}}$$

is a subcode of the Kasami's code for  $u = (m+1)/2$ . Hence the weight restriction implies that (4.6) can be improved to:

$$\Delta_r(y) = \begin{cases} 0 \\ \pm 2^{(n+de)/2}, \end{cases} \quad 1 \leq d \leq m-2, \quad d \text{ odd} \quad (4.7)$$

(4.7) immediately implies that

**LEMMA 4-1:**

If  $r = (2^{5k}+1)/(2^k+1)$ ,  $e = \text{GCD}(n, k)$  and  $n/e$  is odd,

$\Delta_r(y)$  is at most 5-valued.

$$\Delta_r(y) = 0, \pm 2^{(n+e)/2}, \pm 2^{(n+3e)/2}.$$

## 2. COMPUTED RESULTS AND CONJECTURES

Table 4-1 lists the cyclotomic coset leaders given by  $r = (2^{mk}+1)/(2^k+1)$  where both  $m$  and  $n/\text{GCD}(n,k)$  are odd.  $m$  and  $k$  are restricted to (4.1). The number in parentheses following decimation  $r$  is the number of the distinct values that  $\Delta_r(y)$  takes on. For  $n=15$ , the coset leaders given by  $m=15$  are omitted from Table 4-1.

These results exhibit definite patterns and they can be summarized by the following 3 conjectures. As before,  $e = \text{GCD}(n,k)$  and  $r = (2^{mk}+1)/(2^k+1) \not\equiv 2^j \pmod{2^n-1}$  for any  $j$ ,  $0 \leq j \leq n-1$ , with  $m$  and  $n/e$  both odd.

### CONJECTURE 4-2:

If  $e > 1$ , then  $\Delta_r(y)$  is a 3-valued function.  $\Delta_r$  and  $n(\Delta_r)$  are given by (1.8).

### CONJECTURE 4-3:

If  $e = 1$  but  $n$  is not a prime, then  $\Delta_r(y)$  is at most a 5-valued function.  $\Delta_r(y)$  is of the form (4.7).

### CONJECTURE 4-4:

If  $n$  is a prime,  $\Delta_r(y)$  is at most a 5-valued function.

$$\Delta_r(y) = 0, \pm 2^{(n+1)/2}, \pm 2^{(n+3)/2}.$$

**REMARK:** It is observed that  $\Delta_r(y)$  given in Conjecture 4-4 is a 5-valued function if  $m$  is restricted to

$$5 \leq m \leq n-2$$

$$\text{and } 3m \not\equiv \pm 1 \pmod{n}.$$

Next, we give conjectures on decimations of the types not considered in this chapter. Among the decimations that lead to the 3-valued  $\Delta_r(y)$ , there remain only a few cases that are not covered by Theorems 1-5, 1-6 or Lemmas 2-9, 2-10. They are listed in Table 1-2. It is seen that they are still not covered by Conjecture 4-2. The following conjecture covers all of these remaining cases.

**CONJECTURE 4-5:**

The following 5 decimations lead to the 3-valued  $\Delta_r(y)$ .

- (1)  $r = 2^{(n-1)/2} + 3 \quad n \equiv 1 \pmod{2}$
- (2)  $r = 2^{(n-1)/2} + 2^{(n-1)/4} - 1 \quad n \equiv 1 \pmod{4}$
- (3)  $r = 2^{(n-1)/2} + 2^{(3n-1)/4} - 1 \quad n \equiv 3 \pmod{4}$
- (4)  $r = 2^{(n+2)/2} + 3 \quad n \equiv 2 \pmod{4}$
- (5)  $r = 2^{n/2} + 2^{(n+2)/4} + 1 \quad n \equiv 2 \pmod{4}$

$\Delta_r$  and  $n(\Delta_r)$  for (1), (2) and (3) are given by (1.8) with  $e = 1$ .  $\Delta_r$  and  $n(\Delta_r)$  for (4) and (5) are given by (1.8) with  $e = 2$ .

**REMARK:** The case (1) has been known for some time by Welch.

It is of interest to note that all 3 decimations of Conjecture 4-5 for  $n$  odd are of the form  $2^{(n-1)/2} + 2^j - 1$  for some  $j$ . The next conjecture is on the 5-valued  $\Delta_r(y)$ .

CONJECTURE 4-6:

The following 5 decimations lead to at most the 5-valued  $\Delta_r(y)$ .

- (1)  $r = 2^{(n+1)/2} - 2^{(n-3)/2} \pm 1 \quad n \equiv 1 \pmod{2}$
- (2)  $r = 2^{(n+3)/4} + 3 \quad n \equiv 1 \pmod{4}$
- (3)  $r = 2^{(n+1)/2} - 2^{(n+3)/4} + 1 \quad n \equiv 1 \pmod{4}$
- (4)  $r = 2^{(n-1)/2} - 2^{(n+1)/4} + 1 \quad n \equiv 3 \pmod{4}$
- (5)  $r = 2^{n/2+2} - 3 \quad n \equiv 0 \pmod{4}$

Conjectures 4-2 through 4-6 have been verified for  $n \leq 16$ . Conjecture 4-4 with  $k = 1$  and  $m = 5, 7$  &  $9$ , Conjecture 4-5 (1) & (2) and Conjecture 4-6 (2) still hold for  $n = 17$ .

CYCLOTOMIC COSET LEADERS GIVEN BY  $(2^{mk}+1)/(2^k+1)$ 

| n  | k | e | m      |        |        |        |        |
|----|---|---|--------|--------|--------|--------|--------|
|    |   |   | 3      | 5      | 7      | 9      | 11     |
| 3  | 1 | 1 | 3(3)   |        |        |        |        |
| 5  | 1 | 1 | 3(3)   | 11(3)  |        |        |        |
|    | 2 | 1 | 11(3)  | 7(3)   |        |        |        |
| 6  | 2 | 2 | 13(3)  |        |        |        |        |
| 7  | 1 | 1 | 3(3)   | 11(3)  | 43(3)  |        |        |
|    | 2 | 1 | 13(3)  | 29(3)  | 27(3)  |        |        |
|    | 3 | 1 | 23(3)  | 43(3)  | 15(3)  |        |        |
| 9  | 1 | 1 | 3(3)   | 11(5)  | 43(5)  | 171(3) |        |
|    | 2 | 1 | 13(3)  | 107(5) | 109(5) | 103(3) |        |
|    | 3 | 3 | 57(3)  | 1(2)   | 1(2)   | 57(3)  |        |
|    | 4 | 1 | 47(3)  | 109(5) | 93(5)  | 31(3)  |        |
| 10 | 2 | 2 | 13(3)  | 205(3) |        |        |        |
|    | 4 | 2 | 79(3)  | 181(3) |        |        |        |
| 11 | 1 | 1 | 3(3)   | 11(5)  | 43(3)  | 171(5) | 683(3) |
|    | 2 | 1 | 13(3)  | 205(5) | 413(3) | 423(5) | 411(3) |
|    | 3 | 1 | 57(3)  | 235(5) | 683(3) | 343(5) | 231(3) |
|    | 4 | 1 | 143(3) | 121(5) | 151(3) | 429(5) | 365(3) |
|    | 5 | 1 | 95(3)  | 221(5) | 315(3) | 189(5) | 63(3)  |
| 12 | 4 | 4 | 241(3) | 1(2)   |        |        |        |

| n  | k | e | m      |         |         |         |         |         |
|----|---|---|--------|---------|---------|---------|---------|---------|
|    |   |   | 3      | 5       | 7       | 9       | 11      | 13      |
| 13 | 1 | 1 | 3(3)   | 11(5)   | 43(5)   | 171(3)  | 683(5)  | 2731(3) |
|    | 2 | 1 | 13(3)  | 205(5)  | 1643(5) | 1691(3) | 1645(5) | 1639(3) |
|    | 3 | 1 | 57(3)  | 919(5)  | 1367(5) | 2731(3) | 939(5)  | 911(3)  |
|    | 4 | 1 | 241(3) | 497(5)  | 483(5)  | 723(3)  | 1461(5) | 1453(3) |
|    | 5 | 1 | 287(3) | 745(5)  | 869(5)  | 1245(3) | 749(5)  | 1243(3) |
|    | 6 | 1 | 191(3) | 445(5)  | 953(5)  | 635(3)  | 381(5)  | 127(3)  |
| 14 | 2 | 2 | 13(3)  | 205(3)  | 3277(3) |         |         |         |
|    | 4 | 2 | 241(3) | 979(3)  | 2893(3) |         |         |         |
|    | 6 | 2 | 319(3) | 1339(3) | 2773(3) |         |         |         |
| 15 | 1 | 1 | 3(3)   | 11(5)   | 43(5)   | 171(5)  | 683(5)  | 2731(5) |
|    | 2 | 1 | 13(3)  | 205(5)  | 3277(5) | 6557(5) | 6605(5) | 6567(5) |
|    | 3 | 3 | 57(3)  | 3641(3) | 57(3)   | 1(2)    | 1(2)    | 57(3)   |
|    | 4 | 1 | 241(3) | 1943(5) | 5783(5) | 5813(5) | 5805(5) | 2895(5) |
|    | 5 | 5 | 993(3) | 1(2)    | 1(2)    | 993(3)  | 1(2)    | 1(2)    |
|    | 6 | 3 | 575(3) | 3529(3) | 575(3)  | 1(2)    | 1(2)    | 575(3)  |
|    | 7 | 1 | 383(3) | 893(5)  | 1913(5) | 2295(5) | 1275(5) | 765(5)  |

TABLE 4-1

## CHAPTER V

### SUMMARY AND COMMENTS

We have considered the cross-correlation function between two maximal linear recursive sequences:  $\{a_i\}_{i=0}^{2^n-2}$  and  $\{a_{ri}\}_{i=0}^{2^n-2}$ . We have defined the cross-correlation function  $\Delta_r(y)$  by:

$$\Delta_r(y) = \sum_{x \in GF(2^n)} (-1)^{\text{Tr}(xy + x^r)}$$

and analyzed  $\Delta_r(y)$  for two types of decimation  $r$ :

$$(1) \quad r \equiv 2^k \pmod{2^n-1}$$

$$\text{and } (2) \quad r = (2^{mk}+1)/(2^k+1).$$

For the type (1)  $\Delta_r(y)$  is of the form:

$$\Delta_r(y) = 2^{n/2}(j-1), \quad 0 \leq j \leq J, \text{ for some } J.$$

The value  $j$  is the number of distinct solutions to the two equations over  $GF(2^n)$ . For the type (2)  $\Delta_r(y)$  is of the form:

$$\Delta_r(y) = 0, \pm 2^{(n+de)/2}, \quad 1 \leq d \leq m-2, \quad d \text{ odd and} \\ e = \text{GCD}(n, k).$$

The value  $d$  depends on the rank of the quadratic form over  $GF(2^e)$ .

For the both types the complete evaluation of  $n(\Delta_r)$  depends on one's ability to determine the number of solutions to equations in  $GF(2^n)$ . However, using Lemma 2-5,  $n(\Delta_r)$  can be determined completely even though

$n(\Delta_r)$  may be known only for certain  $\Delta_r$  as in the case of Theorems 3-6 and 3-7. For  $r = 2^m + 3$ ,  $n = 2m$  (Theorem 3-8),  $n(\Delta_r)$  can be determined completely if it is possible to

1. evaluate  $\Sigma\{\Delta_r(y)\}^3$
- or 2. find any one of  $N_0$ ,  $N_1$ ,  $N_2$  and  $N_4$  where  
 $N_1$  is the number of times  $N_8(y) = 1$ .

From the observation of results obtained, it is seen that many decimations that lead to 3-valued, 4-valued and 5-valued  $\Delta_r(y)$  are one of the above two types. We have also presented some conjectures on 3-valued and 5-valued  $\Delta_r(y)$  that are not covered by the known theorems.

## APPENDIX A

### CROSS-CORRELATION VALUES

In APPENDIX A  $\Delta_r$  and  $\eta(\Delta_r)$  for all  $r$  are tabulated for  $3 \leq n \leq 12$ . For  $n = 13, 14$  and  $15$ ,  $\Delta_r$  and  $\eta(\Delta_r)$  are given if  $\Delta_r(y)$  is 7-valued or less. For  $n = 16$ ,  $\Delta_r$  and  $\eta(\Delta_r)$  are given if  $\Delta_r(y)$  is 5-valued or less.  $\Delta_r$  and  $\eta(\Delta_r)$  for  $r = 2^{n-1} - 1$  are also given for  $13 \leq n \leq 16$ .

Suppose we have 2 proper cyclotomic coset leaders  $r$  and  $q$  such that  $r \cdot q \equiv 2^k \pmod{2^n - 1}$ . Then, in view of Lemma 2-3,  $\Delta_r$  and  $\eta(\Delta_r)$  are given provided  $r \leq q$ . See APPENDIX B for the inverse pair relation of cyclotomic coset leaders.

In APPENDIX A the first column is the cyclotomic coset leader  $r$ . The second column gives the number of distinct values that  $\Delta_r(y)$  takes on. The numbers to the right give the distribution.  $\eta(\Delta_r)$  and  $\Delta_r$  are given in pair: the first number is  $\eta(\Delta_r)$  and the second number in parentheses is  $\Delta_r$ .

EXAMPLE: For  $n = 7$ ,  $\Delta_9(y)$  is 3-valued.  $\Delta_9(y) = 16$  36 times,  $\Delta_9(y) = 0$  64 times and  $\Delta_9(y) = -16$  28 times as  $y$  ranges over  $GF(2^7)$ . Note that  $9 \cdot 15 = 135 \equiv 2^3 \pmod{127}$ . The cyclotomic cosets containing 9 and 15 are inverse of each other, and  $\Delta_9(y)$  is given but  $\Delta_{15}(y)$  is omitted.



## DEGREE 8

| DEGREE 8 |    |    |    |    |    |    |    |    |    |    |    |
|----------|----|----|----|----|----|----|----|----|----|----|----|
|          | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 | 21 |
| 60       | 11 | 14 | 16 | 18 | 20 | 22 | 24 | 26 | 28 | 30 | 32 |
| 70       | 12 | 15 | 17 | 19 | 21 | 23 | 25 | 27 | 29 | 31 | 33 |
| 80       | 13 | 16 | 18 | 20 | 22 | 24 | 26 | 28 | 30 | 32 | 34 |
| 90       | 14 | 17 | 19 | 21 | 23 | 25 | 27 | 29 | 31 | 33 | 35 |
| 100      | 15 | 18 | 20 | 22 | 24 | 26 | 28 | 30 | 32 | 34 | 36 |
| 110      | 16 | 19 | 21 | 23 | 25 | 27 | 29 | 31 | 33 | 35 | 37 |
| 120      | 17 | 20 | 22 | 24 | 26 | 28 | 30 | 32 | 34 | 36 | 38 |
| 130      | 18 | 21 | 23 | 25 | 27 | 29 | 31 | 33 | 35 | 37 | 39 |
| 140      | 19 | 22 | 24 | 26 | 28 | 30 | 32 | 34 | 36 | 38 | 40 |
| 150      | 20 | 23 | 25 | 27 | 29 | 31 | 33 | 35 | 37 | 39 | 41 |
| 160      | 21 | 24 | 26 | 28 | 30 | 32 | 34 | 36 | 38 | 40 | 42 |
| 170      | 22 | 25 | 27 | 29 | 31 | 33 | 35 | 37 | 39 | 41 | 43 |
| 180      | 23 | 26 | 28 | 30 | 32 | 34 | 36 | 38 | 40 | 42 | 44 |
| 190      | 24 | 27 | 29 | 31 | 33 | 35 | 37 | 39 | 41 | 43 | 45 |
| 200      | 25 | 28 | 30 | 32 | 34 | 36 | 38 | 40 | 42 | 44 | 46 |
| 210      | 26 | 29 | 31 | 33 | 35 | 37 | 39 | 41 | 43 | 45 | 47 |
| 220      | 27 | 30 | 32 | 34 | 36 | 38 | 40 | 42 | 44 | 46 | 48 |
| 230      | 28 | 31 | 33 | 35 | 37 | 39 | 41 | 43 | 45 | 47 | 49 |
| 240      | 29 | 32 | 34 | 36 | 38 | 40 | 42 | 44 | 46 | 48 | 50 |
| 250      | 30 | 33 | 35 | 37 | 39 | 41 | 43 | 45 | 47 | 49 | 51 |
| 260      | 31 | 34 | 36 | 38 | 40 | 42 | 44 | 46 | 48 | 50 | 52 |
| 270      | 32 | 35 | 37 | 39 | 41 | 43 | 45 | 47 | 49 | 51 | 53 |
| 280      | 33 | 36 | 38 | 40 | 42 | 44 | 46 | 48 | 50 | 52 | 54 |
| 290      | 34 | 37 | 39 | 41 | 43 | 45 | 47 | 49 | 51 | 53 | 55 |
| 300      | 35 | 38 | 40 | 42 | 44 | 46 | 48 | 50 | 52 | 54 | 56 |
| 310      | 36 | 39 | 41 | 43 | 45 | 47 | 49 | 51 | 53 | 55 | 57 |
| 320      | 37 | 40 | 42 | 44 | 46 | 48 | 50 | 52 | 54 | 56 | 58 |
| 330      | 38 | 41 | 43 | 45 | 47 | 49 | 51 | 53 | 55 | 57 | 59 |
| 340      | 39 | 42 | 44 | 46 | 48 | 50 | 52 | 54 | 56 | 58 | 60 |
| 350      | 40 | 43 | 45 | 47 | 49 | 51 | 53 | 55 | 57 | 59 | 61 |
| 360      | 41 | 44 | 46 | 48 | 50 | 52 | 54 | 56 | 58 | 60 | 62 |

**DEGREE 9**

[illegible]

|           |    |            |            |             |             |             |             |             |             |             |
|-----------|----|------------|------------|-------------|-------------|-------------|-------------|-------------|-------------|-------------|
| 41        | 9  | 91 641     | 91 441     | 451 321     | 1181 161    | 1191 01     | 991 (-161)  | 631 (-321)  | 91 (-481)   | 11(-1121)   |
| 43        | 5  | 91 641     | 1291 121   | 2861 01     | 1781 (-321) | 11(-641)    |             |             |             |             |
| 45        | 7  | 271 641    | 371 121    | 1381 161    | 1541 01     | 631 (-161)  | 811 (-321)  | 121 (-481)  |             |             |
| 47        | 3  | 1361 321   | 2561 01    | 1271 (-321) |             |             |             |             |             |             |
| 51        | 13 | 371 401    | 271 401    | 91 321      | 271 241     | 161 161     | 821 81      | 821 01      | 811 (-81)   | 541 (-161)  |
|           |    | 441 (-121) | 91 (-401)  | 11(-721)    |             |             |             |             |             | 271 (-261)  |
| 55        | 14 | 11 1001    | 11 441     | 211 481     | 271 401     | 91 321      | 271 241     | 341 161     | 721 41      | 1001 01     |
|           |    | 631 (-161) | 161 (-241) | 271 (-321)  | 191 (-401)  |             |             |             |             | 721 (-81)   |
| 51        | 8  | 271 401    | 541 321    | 1021 161    | 1481 01     | 1171 (-161) | 541 (-321)  | 91 (-481)   | 11(-1121)   |             |
| 74        | 4  | 11 901     | 31 641     | 191 481     | 551 321     | 1081 161    | 1001 01     | 1621 (-161) | 541 (-321)  |             |
| 79        | 8  | 11 501     | 271 441    | 541 321     | 991 161     | 1481 01     | 1171 (-161) | 541 (-321)  | 121 (-481)  |             |
| 91        | 9  | 31 641     | 91 401     | 811 321     | 1081 161    | 1091 01     | 1351 (-161) | 631 (-321)  | 31 (-481)   | 11(-1121)   |
| 85        | 12 | 141 441    | 91 441     | 271 321     | 361 241     | 771 161     | 641 81      | 581 01      | 901 (-81)   | 451 (-161)  |
|           |    | 271 (-121) | 181 (-481) |             |             |             |             |             |             | 481 (-261)  |
| 179       | 5  | 91 641     | 191 121    | 2861 01     | 1041 (-121) | 11(-641)    |             |             |             |             |
| 255       | 23 | 91 641     | 181 401    | 91 361      | 91 321      | 451 281     | 211 241     | 181 201     | 451 161     | 181 121     |
|           |    | 451 41     | 191 01     | 191 (-41)   | 361 (-81)   | 271 (-121)  | 271 (-161)  | 271 (-201)  | 91 (-241)   | 181 (-281)  |
|           |    | 181 (-361) | 91 (-401)  | 31 (-641)   |             |             |             |             |             | 361 (-321)  |
| DEGREE 10 |    |            |            |             |             |             |             |             |             |             |
| 5         | 3  | 1361 641   | 7481 01    | 1201 (-641) |             |             |             |             |             |             |
| 7         | 13 | 51 961     | 21 801     | 401 641     | 701 481     | 1151 321    | 2001 161    | 1961 01     | 1601 (-161) | 1351 (-321) |
|           |    | 201 (-641) | 101 (-801) | 11(-961)    |             |             |             |             |             | 701 (-481)  |
| 13        | 3  | 1361 641   | 7681 01    | 1201 (-641) |             |             |             |             |             |             |
| 17        | 3  | 1361 641   | 7681 01    | 1201 (-641) |             |             |             |             |             |             |
| 19        | 6  | 101 961    | 461 641    | 7601 321    | 4281 01     | 2101 (-321) | 701 (-641)  |             |             |             |
| 23        | 11 | 71 801     | 551 641    | 501 481     | 1101 321    | 1751 161    | 1711 01     | 2441 (-161) | 1001 (-321) | 401 (-481)  |
|           |    | 101 (-901) |            |             |             |             |             |             |             | 401 (-641)  |
| 25        | 3  | 1361 641   | 7641 01    | 1201 (-641) |             |             |             |             |             |             |
| 29        | 11 | 11 1441    | 151 641    | 1251 481    | 1501 321    | 1151 161    | 1731 01     | 1551 (-161) | 1801 (-321) | 701 (-481)  |
|           |    | 101 (-121) |            |             |             |             |             |             |             | 101 (-441)  |
| 35        | 5  | 401 641    | 161 641    | 2401 321    | 1681 01     | 1601 (-321) |             |             |             |             |









|      |    |           |           |            |           |           |            |            |           |           |           |           |           |
|------|----|-----------|-----------|------------|-----------|-----------|------------|------------|-----------|-----------|-----------|-----------|-----------|
| 307  | 29 | 111 1281  | 111 1201  | 111 1121   | 111 1041  | 111 961   | 111 881    | 111 801    | 111 721   | 111 641   | 111 561   | 111 481   | 111 401   |
|      |    | 1451 471  | 1101 321  | 1031 241   | 1431 161  | 1431 81   | 1101 01    | 1441 01    | 1761 01   | 1431 01   | 1101 01   | 1101 01   | 1101 01   |
|      |    | 771 -471  | 841 -481  | 441 -561   | 331 -641  | 331 -721  | 331 -801   | 331 -881   | 441 -961  | 221 -1041 | 221 -1121 | 221 -1201 | 221 -1281 |
| 311  | 9  | 221 1241  | 4401 641  | 11581 01   | 4081 -641 | 221 -1281 |            |            |           |           |           |           |           |
| 317  | 8  | 771 961   | 2531 641  | 4291 321   | 5291 01   | 4951 -321 | 1991 -641  | 551 -961   | 111 -1281 |           |           |           |           |
| 319  | 9  | 111 1281  | 641 941   | 2201 641   | 4621 321  | 5621 01   | 4621 -321  | 1881 -641  | 111 -1281 |           |           |           |           |
| 323  | 8  | 111 1281  | 641 961   | 2091 641   | 4841 321  | 5741 01   | 4101 -321  | 1901 -641  | 881 -961  |           |           |           |           |
| 341  | 29 | 111 1281  | 111 1121  | 111 1041   | 331 881   | 551 801   | 331 721    | 331 641    | 441 561   | 991 481   | 1321 401  |           |           |
|      |    | 841 321   | 1331 241  | 1541 161   | 1321 81   | 1551 01   | 1871 -81   | 941 -161   | 1101 -241 | 1211 -321 | 771 -401  |           |           |
|      |    | 641 -441  | 771 -541  | 441 -641   | 111 -721  | 331 -801  | 331 -881   | 331 -961   | 111 -1041 | 111 -1121 |           |           |           |
| 347  | 9  | 111 1281  | 641 941   | 2201 641   | 4621 321  | 5621 01   | 4621 -321  | 1881 -641  | 111 -1281 |           |           |           |           |
| 357  | 9  | 111 1281  | 551 941   | 2861 641   | 3301 321  | 6611 01   | 4511 -321  | 1881 -641  | 441 -961  | 221 -1281 |           |           |           |
| 371  | 16 | 111 1281  | 111 1121  | 111 1041   | 331 881   | 551 801   | 331 721    | 331 641    | 441 561   | 991 481   | 1321 401  |           |           |
|      |    | 2311 -321 | 1541 -481 | 771 -641   | 661 -801  | 441 -961  | 221 -1121  | 2311 321   | 2641 161  | 3201 01   | 2311 -161 |           |           |
| 415  | 9  | 221 1241  | 331 941   | 2311 641   | 5171 321  | 5071 01   | 4511 -321  | 2211 -641  | 551 -961  | 111 -1281 |           |           |           |
| 443  | 16 | 111 1281  | 121 1121  | 221 941    | 551 801   | 991 641   | 1321 481   | 2421 321   | 3301 161  | 2981 01   | 1981 -161 |           |           |
|      |    | 2211 -321 | 1981 -481 | 771 -641   | 661 -801  | 441 -961  | 111 -1281  |            |           |           |           |           |           |
| 463  | 15 | 111 1281  | 121 1121  | 221 941    | 551 801   | 991 641   | 1321 481   | 1651 321   | 3301 161  | 3531 01   | 2201 -161 |           |           |
|      |    | 1981 -321 | 1321 -481 | 881 -641   | 991 -801  | 551 -961  |            |            |           |           |           |           |           |
| 477  | 16 | 111 1281  | 111 1121  | 331 941    | 661 801   | 1211 641  | 1871 481   | 2201 321   | 2971 161  | 2431 01   | 2531 -161 |           |           |
|      |    | 2311 -321 | 1541 -481 | 1211 -641  | 551 -801  | 441 -961  | 111 -1121  |            |           |           |           |           |           |
| 491  | 19 | 221 1121  | 221 941   | 661 801    | 1431 441  | 1651 481  | 2201 321   | 2311 161   | 2981 01   | 3411 -161 | 1541 -321 |           |           |
|      |    | 1321 -481 | 1431 -641 | 551 -801   | 441 -961  | 111 -1121 |            |            |           |           |           |           |           |
| 497  | 16 | 111 1281  | 121 1121  | 331 941    | 551 801   | 991 641   | 2091 481   | 2201 321   | 2421 161  | 2871 01   | 3081 -161 |           |           |
|      |    | 2311 -321 | 1211 -481 | 991 -641   | 551 -801  | 441 -961  | 221 -1121  |            |           |           |           |           |           |
| 1223 | 45 | 221 841   | 111 841   | 331 801    | 441 761   | 221 721   | 231 681    | 441 641    | 441 601   | 331 541   | 991 521   |           |           |
|      |    | 641 441   | 331 441   | 881 401    | 331 361   | 441 321   | 661 281    | 441 241    | 661 201   | 441 161   | 661 121   |           |           |
|      |    | 221 41    | 771 41    | 551 01     | 331 -41   | 1211 -81  | 221 -121   | 441 -161   | 1101 -241 | 441 -321  | 331 -401  |           |           |
|      |    | 771 -121  | 331 -541  | 441 -601   | 661 -641  | 331 -681  | 221 -721   | 551 -761   | 441 -801  | 331 -841  | 661 -881  |           |           |
|      |    | 221 -721  | 221 -761  | 441 -801   | 111 -841  | 111 -881  |            |            |           |           |           |           |           |
| 11   | 7  | 11 1841   | 241 1921  | 1931 1781  | 11091 641 | 14901 01  | 10521 -641 | 2281 -1281 |           |           |           |           |           |
| 17   | 3  | 1361 2561 | 34401 01  | 1201 -2541 |           |           |            |            |           |           |           |           |           |
| 19   | 8  | 11 2561   | 241 1921  | 2821 1321  | 8641 641  | 16891 01  | 10441 -641 | 1801 -1281 | 131 -1921 |           |           |           |           |

DEGREE 12

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|     |    |                                  |                                  |                                  |                                   |                                    |                        |                       |                       |                       |                       |
|-----|----|----------------------------------|----------------------------------|----------------------------------|-----------------------------------|------------------------------------|------------------------|-----------------------|-----------------------|-----------------------|-----------------------|
| 23  | 13 | 71 2501<br>2441 -961             | 121 2201<br>1111-1201            | 481 1601<br>121-1721             | 1411 1201                         | 2241 961                           | 4961 641               | 6721 321              | 8771 01               | 8041 -321             | 4441 -641             |
| 29  | 13 | 671 1021<br>771-1201             | 671 1601<br>401-1601             | 711 1201<br>171-2201             | 2161 961                          | 4401 641                           | 7241 321               | 8781 01               | 8781 -321             | 4801 -641             | 1801 -961             |
| 31  | 23 | 41 1021<br>2471 321<br>241-1201  | 241 1701<br>4671 161<br>361-1601 | 241 1601<br>4611 01<br>361-1601  | 361 1441<br>4021 -161<br>361-1601 | 691 1201<br>3881 -321<br>2921 -501 | 1201 1121              | 1241 961<br>2641 -641 | 1861 801<br>2521 -801 | 2161 641<br>1201 -961 | 2541 481<br>481-1121  |
| 37  | 16 | 31 2501<br>2521 -961             | 241 1021<br>361-1201             | 241 1401<br>491-1601             | 1741 1201<br>241-1921             | 2881 961                           | 3721 641               | 7121 321              | 9391 01               | 7201 -721             | 5161 -641             |
| 41  | 7  | 91 2501                          | 721 1021                         | 1761 1201                        | 8921 641                          | 16511 01                           | 11561 -641             | 1381-1201             |                       |                       |                       |
| 43  | 7  | 11 3861                          | 241 1021                         | 1931 1201                        | 11091 641                         | 14901 01                           | 10521 -641             | 2281-1201             |                       |                       |                       |
| 47  | 23 | 171 2501<br>331 641<br>541-1121  | 31 1021<br>3101 321<br>621-1121  | 121 1701<br>4621 161<br>241-1601 | 241 1601<br>4851 01<br>241-1601   | 521 1441<br>3421 -161<br>11-2241   | 961 1201<br>3391 -321  | 641 1121<br>3401 -481 | 1021 961<br>2401 -641 | 1381 801<br>2041 -801 | 2161 641<br>1201 -961 |
| 53  | 22 | 21 3861<br>3471 -961             | 71 2501<br>1041-1201             | 481 1601                         | 1081 1201                         | 2681 961                           | 5291 641               | 6841 321              | 7911 01               | 7081 -321             | 5041 -641             |
| 59  | 23 | 11 2841<br>3941 321<br>341-1201  | 41 2501<br>3421 161<br>481-1501  | 201 1701<br>4471 01<br>41-301    | 341 1441<br>4861 -161             | 741 1201<br>2761 -321              | 1141 1121<br>3401 -481 | 1231 961<br>2221 -641 | 1671 801<br>1681 -801 | 2311 641<br>1321 -961 | 2881 481<br>1201-1121 |
| 61  | 25 | 21 3861<br>2441 641<br>1471 -961 | 41 2201<br>2441 641<br>341-1121  | 151 1921<br>3721 321<br>361-1201 | 61 1701<br>4541 161<br>121-1441   | 271 1601<br>4011 01<br>241-1601    | 361 1441<br>2871 -161  | 181 1201<br>2951 -321 | 721 1121<br>3981 -481 | 1201 961<br>2581 -641 | 2641 801<br>1921 -801 |
| 67  | 5  | 1671 1021                        | 321 1201                         | 10241 641                        | 13761 01                          | 15041 -641                         |                        |                       |                       |                       |                       |
| 71  | 6  | 211 2501                         | 441 1021                         | 2771 1201                        | 7721 641                          | 15091 01                           | 14761 -641             |                       |                       |                       |                       |
| 73  | 5  | 3341 1201                        | 1691 641                         | 19201 01                         | 7681 -641                         | 3041-1201                          |                        |                       |                       |                       |                       |
| 77  | 7  | 121 3271                         | 71 2501                          | 641 1021                         | 2341 1201                         | 7441 641                           | 16071 01               | 14281 -641            |                       |                       |                       |
| 83  | 24 | 41 1021<br>1721 321<br>121-1201  | 121 1701<br>4371 161<br>241-1601 | 441 1601<br>4131 01<br>241-1601  | 121 1441<br>3961 -151<br>121-1741 | 871 1201<br>2741 -321              | 1321 1121<br>3401 -481 | 1101 961<br>3451 -641 | 1681 801<br>2041 -801 | 1931 641<br>1441 -961 | 2561 481<br>601-1121  |
| 89  | 15 | 21 3271<br>2441 -961             | 301 1021<br>701-1201             | 361 1601<br>241-1601             | 1051 1201<br>121-1921             | 1541 961<br>121-2241               | 5491 641               | 8041 321              | 8121 01               | 7081 -321             | 4841 -641             |
| 97  | 7  | 121 2501                         | 2721 1201                        | 9321 641                         | 18961 01                          | 8321 -641                          | 2401-1201              | 121-2561              |                       |                       |                       |
| 101 | 15 | 21 1021<br>5441 -641             | 61 2501<br>3121 -961             | 41 1021<br>911-1201              | 121 1601<br>121-1601              | 1171 1201<br>241-1921              | 2481 961               | 5401 641              | 7541 321              | 8181 01               | 6241 -321             |
| 107 | 13 | 71 2501<br>3121 -961             | 121 2201<br>471-1201             | 621 1601<br>361-1601             | 1021 1201                         | 2481 961                           | 4481 641               | 8841 321              | 8951 01               | 8941 -321             | 5161 -641             |
| 113 | 6  | 421 1021                         | 2251 1201                        | 9841 641                         | 15181 01                          | 11581 -641                         | 1691-1201              |                       |                       |                       |                       |

|     |    |            |           |            |            |           |            |            |           |           |           |
|-----|----|------------|-----------|------------|------------|-----------|------------|------------|-----------|-----------|-----------|
| 121 | 14 | 11 2551    | 121 2241  | 131 1921   | 601 1601   | 541 1281  | 2721 961   | 5641 641   | 6241 321  | 8051 01   | 7321 -321 |
|     |    | 5541 -641  | 1241 -941 | 601 -1241  | 241 -1601  |           |            |            |           |           |           |
| 127 | 4  | 6721 1741  | 641 641   | 20151      | 13441 -641 |           |            |            |           |           |           |
| 137 | 14 | 21 3441    | 41 2441   | 71 2541    | 41 1921    | 721 1601  | 1081 1241  | 2881 941   | 3961 641  | 5041 321  | 10031 01  |
|     |    | 4641 -321  | 4541 -641 | 3121 -961  | 721 -1241  |           |            |            |           |           |           |
| 137 | 14 | 11 2551    | 221 1921  | 341 1401   | 1171 1241  | 2681 961  | 4861 641   | 6841 321   | 8551 01   | 6721 -321 | 5281 -641 |
|     |    | 3121 -941  | 971 -1241 | 121 -141   | 41 -1921   |           |            |            |           |           |           |
| 149 | 8  | 11 3441    | 301 1921  | 2041 1281  | 10241 641  | 15841 01  | 10581 -641 | 1681 -1281 | 241 -1921 |           |           |
| 151 | 13 | 11 2541    | 221 1921  | 341 1401   | 1531 1281  | 2881 961  | 4441 641   | 5741 321   | 8741 01   | 7441 -321 | 6061 -641 |
|     |    | 2421 -941  | 751 -1241 | 241 -1401  |            |           |            |            |           |           |           |
| 157 | 24 | 41 2441    | 41 2441   | 31 1921    | 121 1601   | 61 1441   | 781 1281   | 1381 1121  | 1711 961  | 1801 831  | 1801 641  |
|     |    | 941 41     | 241 321   | 3121 161   | 4511 01    | 4901 -161 | 4421 -321  | 1421 -641  | 2101 -641 | 2521 -801 | 1211 -941 |
|     |    | 441 -1121  | 241 -1241 | 481 -1441  | 241 -1921  | 21 -2041  |            |            |           |           |           |
| 163 | 13 | 121 1921   | 361 1601  | 1351 1241  | 3001 961   | 4601 641  | 6361 321   | 9421 01    | 5841 -321 | 5441 -641 | 3361 -941 |
|     |    | 511 -1241  | 241 -1601 | 121 -1921  |            |           |            |            |           |           |           |
| 167 | 13 | 241 1921   | 441 1601  | 1531 1281  | 2721 961   | 2881 641  | 8041 321   | 8341 01    | 8521 -321 | 4041 -641 | 3161 -941 |
|     |    | 511 -1241  | 121 -1601 | 241 -1921  |            |           |            |            |           |           |           |
| 173 | 23 | 151 2241   | 41 1921   | 131 1601   | 521 1441   | 121 1281  | 1321 1121  | 1861 941   | 1441 801  | 2191 641  | 3081 481  |
|     |    | 4021 321   | 3701 141  | 3311 01    | 3941 -161  | 3301 -321 | 3341 -641  | 2881 -641  | 2741 -801 | 1721 -941 | 841 -1121 |
|     |    | 241 -1241  | 341 -1441 | 121 -1601  |            |           |            |            |           |           |           |
| 181 | 24 | 41 2441    | 41 2441   | 151 1921   | 241 1761   | 11 1601   | 61 1441    | 241 1281   | 1401 1121 | 1291 941  | 1501 801  |
|     |    | 2441 641   | 3201 481  | 3721 321   | 1601 161   | 4151 01   | 4561 -161  | 4321 -321  | 2161 -641 | 1321 -641 | 3121 -801 |
|     |    | 1021 -941  | 441 -1121 | 1201 -1281 | 181 -1441  |           |            |            |           |           |           |
| 187 | 22 | 11 2441    | 121 1921  | 301 1441   | 641 1281   | 1141 1121 | 1501 961   | 2501 801   | 2161 641  | 3041 481  | 3301 321  |
|     |    | 3721 161   | 4371 01   | 3141 -161  | 1471 -321  | 3441 -641 | 2191 -641  | 2041 -801  | 1461 -941 | 841 -1121 | 281 -1281 |
|     |    | 401 -1441  | 121 -1921 |            |            |           |            |            |           |           |           |
| 197 | 7  | 121 321    | 161 2541  | 401 1921   | 2401 1281  | 7801 641  | 15681 01   | 14401 -641 |           |           |           |
| 199 | 23 | 41 2541    | 111 1921  | 121 1761   | 241 1601   | 361 1441  | 271 1281   | 1201 1121  | 1561 961  | 1201 801  | 2111 641  |
|     |    | 4441 431   | 2641 321  | 1491 161   | 3971 01    | 5421 -161 | 2741 -321  | 2641 -641  | 2221 -641 | 2421 -801 | 1441 -941 |
|     |    | 941 -1121  | 541 -1241 | 471 -1441  |            |           |            |            |           |           |           |
| 209 | 24 | 121 2441   | 31 1921   | 121 1761   | 41 1601    | 341 1441  | 511 1241   | 881 1121   | 841 941   | 1841 831  | 3241 641  |
|     |    | 4041 431   | 3721 321  | 3121 161   | 4251 01    | 1401 -161 | 1401 -161  | 3121 -641  | 2141 -641 | 1801 -801 | 1381 -841 |
|     |    | 1541 -1121 | 941 -1241 | 341 -1441  | 21 -1601   | 121 -1761 | 11 -1921   |            |           |           |           |
| 211 | 25 | 41 321     | 71 1921   | 241 1761   | 181 1601   | 241 1441  | 751 1241   | 601 1121   | 1161 961  | 2141 801  | 1811 641  |
|     |    | 3121 431   | 4141 321  | 3841 161   | 4691 01    | 3441 -161 | 3001 -321  | 3241 -641  | 2701 -641 | 2641 -801 | 1321 -941 |
|     |    | 601 -1121  | 441 -1241 | 161 -1441  | 121 -1601  | 121 -1761 |            |            |           |           |           |
| 209 | 44 | 11 4441    | 41 2441   | 21 2221    | 41 2441    | 41 2441   | 41 2441    | 151 1921   | 121 1441  | 121 1441  | 181 1281  |
|     |    | 401 1241   | 601 1121  | 841 1041   | 1321 961   | 721 941   | 781 801    | 841 721    | 1141 641  | 1201 561  | 1261 481  |
|     |    | 1021 431   | 1141 321  | 1481 241   | 1331 161   | 1441 41   | 2411 01    | 1441 -801  | 2641 -161 | 1801 -241 | 1801 -321 |
|     |    | 1441 -641  | 1441 -641 | 1401 -561  | 1931 -641  | 601 -721  | 721 -801   | 1201 -841  | 311 -941  | 341 -1041 | 481 -1121 |
|     |    | 441 -1241  | 121 -1241 | 121 -1241  | 361 -1441  |           |            |            |           |           |           |

|     |    |           |           |           |           |           |           |           |           |           |           |
|-----|----|-----------|-----------|-----------|-----------|-----------|-----------|-----------|-----------|-----------|-----------|
| 253 | 27 | 11 2801   | 61 2621   | 31 2241   | 21 1761   | 121 1601  | 61 1441   | 1321 1281 | 841 1121  | 781 961   | 2501 801  |
|     |    | 2151 641  | 2821 441  | 3421 321  | 4441 161  | 4571 01   | 3721 161  | 3601 321  | 3281 481  | 1951 641  | 2041 801  |
|     |    | 1201 961  | 601 1121  | 741 1281  | 301 1441  | 121 1601  | 121 1761  | 121 1921  |           |           |           |
| 255 | 5  | 331 2561  | 3421 1291 | 7241 641  | 15211 01  | 14761 641 |           |           |           |           |           |
| 277 | 9  | 11 3001   | 61 3201   | 41 2561   | 221 1921  | 2401 121  | 9041 641  | 16481 01  | 10681 641 | 1711 1281 |           |
| 283 | 24 | 61 2561   | 241 1761  | 141 1601  | 441 1441  | 461 1781  | 481 1121  | 1051 961  | 1261 801  | 2281 641  | 4541 481  |
|     |    | 4201 321  | 2821 161  | 4451 01   | 4781 161  | 3751 321  | 2741 481  | 1351 641  | 2281 801  | 1091 961  | 481 1121  |
|     |    | 901 1281  | 601 1441  | 121 1601  | 121 1921  |           |           |           |           |           |           |
| 293 | 12 | 41 441    | 31 5121   | 1621 1281 | 2481 961  | 2441 641  | 8641 321  | 10791 01  | 5141 321  | 4701 641  | 2801 961  |
|     |    | 441 1281  | 341 1601  |           |           |           |           |           |           |           |           |
| 307 | 26 | 31 2241   | 41 1721   | 121 1761  | 241 1601  | 581 1441  | 421 1281  | 341 1121  | 1201 961  | 2381 801  | 1481 641  |
|     |    | 441 641   | 3241 321  | 4541 161  | 3211 01   | 4561 161  | 1241 321  | 4621 481  | 2761 641  | 2161 801  | 721 961   |
|     |    | 1341 1121 | 511 1291  | 121 1601  | 121 1761  | 121 1921  | 11 2881   |           |           |           |           |
| 319 | 7  | 121 3201  | 71 2441   | 641 1921  | 2341 1281 | 7441 641  | 16071 01  | 14281 641 |           |           |           |
| 331 | 8  | 21 3441   | 121 3201  | 71 2561   | 361 1921  | 2761 1281 | 7441 641  | 15791 01  | 14401 641 |           |           |
| 341 | 14 | 71 1721   | 241 1761  | 401 1441  | 661 1281  | 721 1121  | 3601 801  | 2641 641  | 5941 481  | 6161 161  | 3821 01   |
|     |    | 4401 161  | 4201 441  | 2571 641  | 3161 901  | 1841 1121 | 441 1281  | 241 1441  | 121 1761  |           |           |
| 349 | 24 | 191 1021  | 121 1761  | 341 1601  | 161 1441  | 601 1281  | 721 1121  | 1401 961  | 1321 801  | 2881 641  | 2521 481  |
|     |    | 4341 321  | 3441 161  | 3131 01   | 4081 161  | 3541 321  | 3641 481  | 2851 641  | 2281 801  | 1001 961  | 841 1121  |
|     |    | 441 1281  | 121 1441  | 121 1601  | 121 1921  |           |           |           |           |           |           |
| 341 | 13 | 11 2561   | 41 1921   | 241 1601  | 1291 1781 | 2401 961  | 4201 641  | 6961 321  | 8071 01   | 7681 321  | 5641 641  |
|     |    | 2541 961  | 471 1231  | 241 1401  |           |           |           |           |           |           |           |
| 347 | 13 | 41 2561   | 441 1601  | 871 1281  | 2161 961  | 5701 641  | 8201 321  | 7441 01   | 6161 321  | 4281 641  | 3001 961  |
|     |    | 961 1281  | 121 1601  | 221 1921  |           |           |           |           |           |           |           |
| 377 | 7  | 121 3201  | 101 2541  | 521 1921  | 2521 1281 | 7321 641  | 16101 01  | 14281 641 |           |           |           |
| 397 | 14 | 11 3441   | 61 2541   | 121 2241  | 121 1921  | 721 1601  | 1311 1281 | 1921 961  | 3961 641  | 8081 321  | 6521 01   |
|     |    | 3141 321  | 5281 641  | 3041 961  | 641 1281  |           |           |           |           |           |           |
| 409 | 22 | 41 2561   | 241 1761  | 441 1601  | 541 1441  | 391 1281  | 601 1121  | 961 961   | 1921 801  | 1441 641  | 4681 481  |
|     |    | 2741 321  | 5141 161  | 3431 01   | 4081 161  | 1961 321  | 4181 481  | 2941 641  | 1921 801  | 641 961   | 1561 1121 |
|     |    | 341 1231  | 431 1441  |           |           |           |           |           |           |           |           |
| 431 | 23 | 21 2241   | 141 2081  | 121 1601  | 441 1441  | 781 1281  | 841 1121  | 1081 961  | 1321 801  | 3151 641  | 3241 481  |
|     |    | 3341 321  | 4341 161  | 3981 01   | 3761 161  | 3541 321  | 3241 481  | 2041 641  | 1801 801  | 1801 961  | 721 1121  |
|     |    | 441 1281  | 501 1441  | 11 1921   |           |           |           |           |           |           |           |
| 433 | 23 | 11 3441   | 31 5441   | 31 5441   | 121 1761  | 161 1441  | 721 1281  | 241 1121  | 1681 961  | 1861 801  | 1621 641  |
|     |    | 4501 441  | 4741 321  | 2701 161  | 4091 01   | 4261 161  | 2701 321  | 4081 481  | 2941 641  | 1021 801  | 2121 961  |
|     |    | 641 1121  | 121 1281  | 341 1441  |           |           |           |           |           |           |           |
| 457 | 4  | 41 5121   | 1721 641  | 5041 01   | 17921 641 |           |           |           |           |           |           |
| 441 | 9  | 61 2541   | 41 1921   | 2641 1281 | 9161 641  | 17601 01  | 8721 641  | 2401 1281 | 121 1921  | 21 2541   |           |

|      |    |                                                                              |                                                                  |                                                                  |                                                                   |                                                       |                                                        |                                                        |                                                        |                                                         |                                                                    |
|------|----|------------------------------------------------------------------------------|------------------------------------------------------------------|------------------------------------------------------------------|-------------------------------------------------------------------|-------------------------------------------------------|--------------------------------------------------------|--------------------------------------------------------|--------------------------------------------------------|---------------------------------------------------------|--------------------------------------------------------------------|
| 463  | 9  | 51 5121                                                                      | 11 1941                                                          | 21 3201                                                          | 14141 641                                                         | 15861 01                                              | 8201 -641                                              | 2191-1281                                              | 361-1921                                               | 121-2561                                                |                                                                    |
| 467  | 11 | 41 2561<br>1521 -961                                                         | 141 1021<br>611-1281                                             | 441 1601<br>361-1601                                             | 911 1281                                                          | 3401 961                                              | 4321 641                                               | 6241 321                                               | 9801 01                                                | 6481 -321                                               | 4741 -641                                                          |
| 479  | 13 | 181 1921<br>941-1281                                                         | 471 1601<br>441-1601                                             | 1201 1281<br>21-2561                                             | 2201 961                                                          | 5041 641                                              | 7481 321                                               | 8701 01                                                | 6841 -321                                              | 4381 -641                                               | 3001 -961                                                          |
| 481  | 11 | 191 1921<br>441-1231                                                         | 441 1601<br>241-1601                                             | 1641 1281<br>241-1921                                            | 1401 961                                                          | 4491 641                                              | 7321 321                                               | 9001 01                                                | 7081 -321                                              | 5341 -641                                               | 2281 -961                                                          |
| 497  | 24 | 21 4431<br>1481 641<br>1081 -961                                             | 41 2561<br>3041 481<br>441-1121                                  | 121 1921<br>2501 321<br>421-1281                                 | 241 1741<br>3121 161<br>121-1601                                  | 781 1601<br>4911 01                                   | 244 1441<br>4641 -161                                  | 781 1281<br>5701 -321                                  | 1081 1121<br>4121 -681                                 | 601 961<br>1801 -641                                    | 1321 801<br>2161 -801                                              |
| 509  | 24 | 31 5431<br>3641 441<br>941-1121                                              | 41 5151<br>441 121<br>361-1281                                   | 11 4481<br>3001 161<br>171-1441                                  | 121 1741<br>5491 01<br>121-1601                                   | 401 1441<br>4561 -161                                 | 961 1281<br>2461 -321                                  | 361 1121<br>4081 -481                                  | 1381 961<br>2581 -641                                  | 1921 801<br>1441 -801                                   | 901 641<br>1321 -961                                               |
| 619  | 24 | 31 3201<br>2441 321<br>241-1281                                              | 121 2241<br>3121 161<br>121-1441                                 | 421 1601<br>4021 01<br>121-1601                                  | 241 1441<br>5641 -161<br>241-2081                                 | 361 1281<br>3781 -321                                 | 1201 1121<br>1401 -481                                 | 1321 961<br>2821 -641                                  | 1921 801<br>1561 -801                                  | 2771 641<br>1201 -961                                   | 2841 481<br>841-1121                                               |
| 671  | 13 | 11 3841<br>941-1281                                                          | 241 1601<br>401-1601                                             | 1311 1281<br>171-1321                                            | 3081 961                                                          | 3781 641                                              | 6481 321                                               | 9161 01                                                | 7481 -321                                              | 4501 -641                                               | 2641 -961                                                          |
| 683  | 6  | 111451                                                                       | 211341                                                           | 12181 641                                                        | 18041 01                                                          | 9241 -641                                             | 1471-1281                                              |                                                        |                                                        |                                                         |                                                                    |
| 701  | 6  | 121 2561                                                                     | 441 1021                                                         | 2961 1281                                                        | 4881 641                                                          | 15401 01                                              | 14561 -641                                             |                                                        |                                                        |                                                         |                                                                    |
| 921  | 7  | 41 9321                                                                      | 11 7631                                                          | 241 1721                                                         | 2421 1241                                                         | 7681 641                                              | 17311 01                                               | 13241 -641                                             |                                                        |                                                         |                                                                    |
| 927  | 7  | 111241                                                                       | 41 7431                                                          | 121 2561                                                         | 2641 1241                                                         | 7201 641                                              | 17091 01                                               | 12961 -641                                             |                                                        |                                                         |                                                                    |
| 929  | 13 | 71 2561<br>2271 -961                                                         | 91 1921<br>1381-1281                                             | 601 1401<br>241-1401                                             | 1141 1281                                                         | 2281 961                                              | 4861 641                                               | 7201 321                                               | 8371 01                                                | 7321 -321                                               | 5221 -641                                                          |
| 977  | 12 | 41 2741<br>3121 -961                                                         | 121 1721<br>1711-1281                                            | 241 1601                                                         | 1911 1281                                                         | 1921 961                                              | 4081 641                                               | 8641 321                                               | 8461 01                                                | 6521 -321                                               | 4201 -641                                                          |
| 997  | 13 | 11 2561<br>2271 -961                                                         | 241 1921<br>1041-1281                                            | 721 1601<br>481-1601                                             | 1081 1281                                                         | 1941 961                                              | 5461 641                                               | 5641 321                                               | 9911 01                                                | 7681 -321                                               | 4621 -641                                                          |
| 2947 | 44 | 121 1241<br>421 961<br>1171 481<br>1241 41<br>641-121<br>481-721<br>481-1121 | 121 1241<br>1091 481<br>1271 441<br>721 41<br>961-361<br>481-761 | 301 1201<br>811 801<br>441 401<br>731 01<br>781 -401<br>241 -801 | 521 1161<br>481 761<br>721 361<br>961 -41<br>601 -441<br>961 -861 | 241 1121<br>721 721<br>441 321<br>481 -81<br>721 -481 | 491 1091<br>491 641<br>941 241<br>961 -121<br>841 -521 | 601 1041<br>481 641<br>481 201<br>961 -161<br>481 -561 | 161 1001<br>481 601<br>441 201<br>361 -201<br>481 -601 | 721 961<br>961 561<br>961 161<br>1441 -241<br>1201 -641 | 481 921<br>481 521<br>601 121<br>1081 -281<br>601 -681<br>241-1081 |

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DEGREE 13

|     |   |            |            |             |             |           |  |  |  |
|-----|---|------------|------------|-------------|-------------|-----------|--|--|--|
| 3   | 3 | 2080( 128) | 4096( 0)   | 2016( -128) |             |           |  |  |  |
| 5   | 3 | 2080( 128) | 4096( 0)   | 2016( -128) |             |           |  |  |  |
| 9   | 3 | 2080( 128) | 4096( 0)   | 2016( -128) |             |           |  |  |  |
| 11  | 5 | 91( 256)   | 1716( 128) | 4642( 0)    | 1652( -128) | 91( -256) |  |  |  |
| 13  | 3 | 2080( 128) | 4096( 0)   | 2016( -128) |             |           |  |  |  |
| 17  | 3 | 2080( 128) | 4096( 0)   | 2016( -128) |             |           |  |  |  |
| 19  | 5 | 91( 256)   | 1716( 128) | 4642( 0)    | 1652( -128) | 91( -256) |  |  |  |
| 33  | 3 | 2080( 128) | 4096( 0)   | 2016( -128) |             |           |  |  |  |
| 43  | 5 | 91( 256)   | 1716( 128) | 4642( 0)    | 1652( -128) | 91( -256) |  |  |  |
| 57  | 3 | 2080( 128) | 4096( 0)   | 2016( -128) |             |           |  |  |  |
| 65  | 3 | 2080( 128) | 4096( 0)   | 2016( -128) |             |           |  |  |  |
| 67  | 3 | 2080( 128) | 4096( 0)   | 2016( -128) |             |           |  |  |  |
| 71  | 3 | 2080( 128) | 4096( 0)   | 2016( -128) |             |           |  |  |  |
| 95  | 5 | 91( 256)   | 1716( 128) | 4642( 0)    | 1652( -128) | 91( -256) |  |  |  |
| 113 | 5 | 91( 256)   | 1716( 128) | 4642( 0)    | 1652( -128) | 91( -256) |  |  |  |
| 147 | 5 | 65( 256)   | 1827( 128) | 4466( 0)    | 1756( -128) | 65( -256) |  |  |  |
| 149 | 5 | 91( 256)   | 1716( 128) | 4642( 0)    | 1652( -128) | 91( -256) |  |  |  |
| 171 | 3 | 2080( 128) | 4096( 0)   | 2016( -128) |             |           |  |  |  |
| 179 | 3 | 2080( 128) | 4096( 0)   | 2016( -128) |             |           |  |  |  |
| 205 | 5 | 91( 256)   | 1716( 128) | 4642( 0)    | 1652( -128) | 91( -256) |  |  |  |
| 225 | 5 | 65( 256)   | 1827( 128) | 4466( 0)    | 1756( -128) | 65( -256) |  |  |  |
| 287 | 3 | 2080( 128) | 4096( 0)   | 2016( -128) |             |           |  |  |  |
| 445 | 5 | 91( 256)   | 1716( 128) | 4642( 0)    | 1652( -128) | 91( -256) |  |  |  |
| 483 | 5 | 91( 256)   | 1716( 128) | 4642( 0)    | 1652( -128) | 91( -256) |  |  |  |
| 631 | 5 | 78( 256)   | 1763( 128) | 4564( 0)    | 1704( -128) | 78( -256) |  |  |  |
| 683 | 5 | 91( 256)   | 1716( 128) | 4642( 0)    | 1652( -128) | 91( -256) |  |  |  |







|       |     |      |       |       |       |        |       |       |       |      |        |      |       |
|-------|-----|------|-------|-------|-------|--------|-------|-------|-------|------|--------|------|-------|
| 1913  | 5   | 3801 | 5121  | 68001 | 2561  | 184721 | 01    | 68001 | -2561 | 3161 | -5121  | 1051 | 3321  |
| 2295  | 5   | 331  | 10241 | 79201 | 2561  | 168941 | 01    | 79201 | -2561 | 11   | -10241 | 901  | 3001  |
| 2521  | 5   | 3951 | 5121  | 67401 | 2561  | 185621 | 01    | 67401 | -2561 | 3311 | -5121  | 2101 | 3001  |
| 2731  | 5   | 3801 | 5121  | 68001 | 2561  | 184721 | 01    | 68001 | -2561 | 3161 | -5121  | 901  | 2001  |
| 2981  | 5   | 3801 | 5121  | 67361 | 2561  | 184641 | 01    | 66081 | -2561 | 3801 | -5121  | 901  | 1721  |
| 3653  | 5   | 331  | 10241 | 79201 | 2561  | 168941 | 01    | 79201 | -2561 | 11   | -10241 | 2851 | 1401  |
| 5783  | 5   | 3801 | 5121  | 68001 | 2561  | 184721 | 01    | 68001 | -2561 | 3161 | -5121  | 1201 | 1701  |
| 5813  | 5   | 331  | 10241 | 79201 | 2561  | 168941 | 01    | 79201 | -2561 | 11   | -10241 | 3001 | 761   |
| 16303 | 181 | 351  | 3631  | 351   | 3561  | 1201   | 3521  | 451   | 3481  | 301  | 3441   | 1201 | 3401  |
|       |     | 901  | 3261  | 901   | 3241  | 1051   | 3201  | 1651  | 3161  | 751  | 3121   | 901  | 3781  |
|       |     | 901  | 2801  | 2801  | 2801  | 901    | 2801  | 1901  | 2841  | 3751 | 2801   | 2701 | 2441  |
|       |     | 3751 | 2321  | 1201  | 2281  | 1051   | 2241  | 3601  | 2201  | 1201 | 2151   | 1801 | 2121  |
|       |     | 1401 | 2701  | 1651  | 1941  | 2101   | 1921  | 1201  | 1841  | 1151 | 1841   | 1501 | 1401  |
|       |     | 1351 | 1401  | 1501  | 1641  | 4201   | 1601  | 1501  | 1561  | 1651 | 1521   | 2401 | 1541  |
|       |     | 2451 | 1361  | 2101  | 1321  | 1801   | 1281  | 2131  | 1241  | 1651 | 1701   | 901  | 1161  |
|       |     | 1501 | 1041  | 5251  | 1001  | 1551   | 961   | 1801  | 921   | 3151 | 881    | 1651 | 941   |
|       |     | 1951 | 721   | 1351  | 691   | 2251   | 641   | 2251  | 601   | 1651 | 561    | 1501 | 441   |
|       |     | 2701 | 401   | 1251  | 361   | 1501   | 321   | 2851  | 281   | 1401 | 241    | 2401 | 161   |
|       |     | 1751 | 81    | 2251  | 41    | 2841   | 71    | 1201  | -41   | 3401 | -81    | 1501 | -121  |
|       |     | 1251 | -241  | 3051  | -281  | 3601   | -321  | 1351  | -361  | 1801 | -401   | 3001 | -441  |
|       |     | 3701 | -561  | 2701  | -601  | 1501   | -641  | 5551  | -681  | 1051 | -721   | 1201 | -761  |
|       |     | 1801 | -881  | 2101  | -921  | 1201   | -961  | 1651  | -1001 | 1951 | -1041  | 2251 | -1081 |
|       |     | 3701 | -1201 | 1401  | -1241 | 3301   | -1281 | 1351  | -1321 | 1351 | -1361  | 3751 | -1401 |
|       |     | 3401 | -1521 | 901   | -1561 | 1501   | -1601 | 2401  | -1641 | 2401 | -1681  | 1201 | -1721 |
|       |     | 1051 | -1841 | 3301  | -1881 | 901    | -1921 | 1501  | -1961 | 3901 | -2001  | 3001 | -2041 |
|       |     | 1051 | -2141 | 1651  | -2201 | 2551   | -2241 | 1201  | -2281 | 1051 | -2321  | 2551 | -2361 |
|       |     | 2551 | -2441 | 1501  | -2521 | 1351   | -2561 | 2501  | -2601 | 1201 | -2641  | 1051 | -2681 |
|       |     | 1231 | -2801 | 1801  | -2841 | 1501   | -2881 | 1201  | -2921 | 1651 | -2961  | 1051 | -3001 |
|       |     | 601  | -3121 | 601   | -3161 | 2251   | -3201 | 451   | -3241 | 951  | -3281  | 751  | -3321 |
|       |     | 901  | -3441 | 601   | -3481 | 301    | -3521 | 451   | -3561 | 151  | -3601  | 601  | -3641 |

DEGREE 16

|      |   |             |             |             |              |              |
|------|---|-------------|-------------|-------------|--------------|--------------|
| 259  | 5 | 2698( 740)  | 128( 512)   | 16384( 256) | 21898( 0)    | 24448( -256) |
| 303  | 5 | 548( 1024)  | 5407( 512)  | 11216( 256) | 24620( 0)    | 23952( -256) |
| 511  | 4 | 13880( 512) | 256( 256)   | 32640( 0)   | 21760( -256) |              |
| 1021 | 5 | 540( 1024)  | 5480( 512)  | 11056( 256) | 24540( 0)    | 23920( -256) |
| 3857 | 4 | 161 4096)   | 30720( 256) | 4080( 0)    | 30720( -256) |              |
| 7399 | 4 | 10880( 512) | 256( 256)   | 32640( 0)   | 21760( -256) |              |

[illegible]

## APPENDIX B

### INVERSE PAIR RELATION OF CYCLOTOMIC COSET LEADERS

In APPENDIX B the inverse pair relation of cyclotomic coset leaders is given for  $3 \leq n \leq 14$ . If two cyclotomic coset leaders  $r$  and  $q$  are such that  $r \cdot q \equiv 2^k \pmod{2^n - 1}$  and  $r \leq q$ , then  $r$  and  $q$  are given in pair with  $q$  in parentheses.

EXAMPLE: For  $n = 7$ , two cyclotomic cosets containing 9 and 15 are inverse of each other as mentioned above. For  $n = 9$  the cyclotomic coset containing 75 is a self-inverse since  $(75)^2 = 5625 \equiv 2^2 \pmod{511}$ .

INVERSE PAIR RELATION OF CYCLOTOMIC COSET LEADERS

|           |          | DEGREE 3 |          | 2 PAIRS   |           |
|-----------|----------|----------|----------|-----------|-----------|
| 1( 1)     | 3( 3)    |          |          |           |           |
|           |          |          |          |           |           |
|           |          | DEGREE 4 |          | 2 PAIRS   |           |
| 1( 1)     | 7( 7)    |          |          |           |           |
|           |          |          |          |           |           |
|           |          | DEGREE 5 |          | 4 PAIRS   |           |
| 1( 1)     | 3( 11)   | 5( 7)    | 15( 15)  |           |           |
|           |          |          |          |           |           |
|           |          | DEGREE 6 |          | 4 PAIRS   |           |
| 1( 1)     | 5( 13)   | 11( 23)  | 31( 31)  |           |           |
|           |          |          |          |           |           |
|           |          | DEGREE 7 |          | 10 PAIRS  |           |
| 1( 1)     | 3( 43)   | 5( 27)   | 7( 55)   | 9( 15)    | 11( 13)   |
| 63( 63)   |          |          |          | 19( 47)   | 21( 31)   |
|           |          |          |          |           |           |
|           |          | DEGREE 8 |          | 12 PAIRS  |           |
| 1( 1)     | 7( 57)   | 11( 29)  | 13( 59)  | 19( 47)   | 23( 61)   |
| 127( 127) |          |          |          | 31( 91)   | 43( 43)   |
|           |          |          |          |           |           |
|           |          | DEGREE 9 |          | 24 PAIRS  |           |
| 1( 1)     | 3( 171)  | 5( 103)  | 9( 57)   | 11( 93)   | 13( 59)   |
| 23( 23)   | 29( 43)  | 37( 181) | 39( 95)  | 41( 107)  | 43( 107)  |
| 59( 223)  | 61( 111) | 75( 75)  | 79( 123) | 83( 117)  | 85( 127)  |
|           |          |          |          |           |           |
|           |          |          |          | 15( 239)  | 17( 31)   |
|           |          |          |          | 45( 123)  | 47( 47)   |
|           |          |          |          | 109( 109) | 255( 255) |
|           |          |          |          |           | 19( 27)   |
|           |          |          |          |           | 51( 191)  |

## DEGREE 10 32 PAIRS

|      |      |      |      |      |      |      |      |      |      |      |      |      |      |      |      |
|------|------|------|------|------|------|------|------|------|------|------|------|------|------|------|------|
| 11   | 11   | 51   | 2051 | 71   | 4301 | 131  | 791  | 171  | 1011 | 191  | 1751 | 251  | 411  | 291  | 2471 |
| 351  | 951  | 371  | 431  | 431  | 1101 | 471  | 1091 | 491  | 1071 | 531  | 2511 | 511  | 1511 | 711  | 2451 |
| 741  | 1771 | 451  | 451  | 911  | 2151 | 1011 | 1571 | 1031 | 1491 | 1151 | 4471 | 1671 | 2301 | 1731 | 4791 |
| 1791 | 3031 | 2231 | 3571 | 2351 | 3791 | 3431 | 1631 | 5111 | 5111 |      |      |      |      |      |      |

## DEGREE 11 90 PAIRS

|      |      |      |      |      |      |      |      |      |      |      |      |      |      |      |       |
|------|------|------|------|------|------|------|------|------|------|------|------|------|------|------|-------|
| 11   | 11   | 31   | 4631 | 51   | 4111 | 71   | 2931 | 91   | 2311 | 111  | 1391 | 151  | 1371 | 171  | 1651  |
| 191  | 4711 | 211  | 991  | 251  | 871  | 271  | 771  | 291  | 751  | 311  | 7311 | 331  | 1171 | 371  | 831   |
| 191  | 1051 | 511  | 4791 | 411  | 1431 | 431  | 911  | 471  | 471  | 511  | 2011 | 531  | 1151 | 571  | 4131  |
| 591  | 2631 | 511  | 2371 | 671  | 1431 | 711  | 1711 | 731  | 351  | 751  | 1431 | 771  | 5291 | 791  | 1511  |
| 1011 | 5271 | 1071 | 1531 | 1071 | 6951 | 1111 | 2031 | 1131 | 6711 | 1191 | 5711 | 1191 | 4231 | 1211 | 2731  |
| 1251 | 4791 | 1271 | 8471 | 1411 | 3631 | 1471 | 3491 | 1491 | 7351 | 1591 | 1021 | 1571 | 3391 | 1591 | 2111  |
| 1471 | 3311 | 1501 | 3731 | 1711 | 2051 | 1751 | 5031 | 1791 | 2231 | 1911 | 4751 | 1851 | 2131 | 1911 | 4931  |
| 1491 | 4951 | 2151 | 2191 | 2171 | 3431 | 2291 | 2931 | 2351 | 6291 | 2291 | 4291 | 2471 | 4731 | 2511 | 3671  |
| 2551 | 7311 | 3011 | 9591 | 3071 | 7671 | 3111 | 3111 | 3171 | 3011 | 4131 | 3791 | 3331 | 3751 | 3411 | 3511  |
| 3501 | 7271 | 3711 | 4691 | 4151 | 4391 | 4431 | 4951 | 4631 | 7031 | 4771 | 7511 | 4911 | 7631 | 4071 | 10231 |

## DEGREE 12 76 PAIRS

|      |       |      |       |      |       |       |       |      |       |      |       |      |       |      |      |
|------|-------|------|-------|------|-------|-------|-------|------|-------|------|-------|------|-------|------|------|
| 11   | 11    | 111  | 1731  | 171  | 2411  | 191   | 2291  | 231  | 8931  | 231  | 7891  | 371  | 10231 | 411  | 1031 |
| 431  | 6671  | 671  | 4371  | 531  | 2731  | 591   | 4291  | 611  | 10071 | 671  | 1911  | 731  | 1511  | 791  | 4731 |
| 631  | 15311 | 891  | 5791  | 971  | 7191  | 1011  | 2231  | 1071 | 4211  | 1091 | 1131  | 1271 | 13971 | 1371 | 7511 |
| 1391 | 3431  | 1491 | 6591  | 1511 | 4071  | 1571  | 4791  | 1631 | 1791  | 1671 | 8591  | 1741 | 1411  | 1871 | 5431 |
| 1971 | 3171  | 1991 | 2271  | 2091 | 17831 | 2111  | 9911  | 2191 | 19191 | 2511 | 4111  | 2541 | 2811  | 2431 | 9411 |
| 2071 | 5471  | 3791 | 3771  | 3131 | 3571  | 3311  | 6111  | 3411 | 6531  | 3411 | 13211 | 3611 | 7431  | 3791 | 4431 |
| 3071 | 7331  | 4791 | 6911  | 4311 | 10141 | 4391  | 4791  | 4571 | 4571  | 4611 | 4591  | 4631 | 14731 | 4791 | 6271 |
| 4011 | 6131  | 4931 | 14711 | 5371 | 17501 | 6191  | 10111 | 4711 | 7271  | 6031 | 13671 | 7011 | 8231  | 4271 | 9271 |
| 4201 | 6011  | 6771 | 14991 | 5471 | 14931 | 20471 | 20471 |      |       |      |       |      |       |      |      |

## DEGREE 13 316 PAIRS

|      |       |      |       |      |       |      |       |      |       |      |       |      |       |      |       |
|------|-------|------|-------|------|-------|------|-------|------|-------|------|-------|------|-------|------|-------|
| 11   | 11    | 11   | 27411 | 51   | 14311 | 71   | 35111 | 91   | 9111  | 111  | 7451  | 151  | 34231 | 171  | 14531 |
| 191  | 4451  | 211  | 17911 | 231  | 17911 | 251  | 7011  | 271  | 15171 | 291  | 2341  | 311  | 12631 | 331  | 7031  |
| 371  | 6051  | 391  | 1711  | 511  | 9911  | 431  | 3911  | 451  | 12771 | 471  | 1771  | 491  | 10031 | 511  | 13911 |
| 551  | 1431  | 571  | 7231  | 591  | 1411  | 611  | 9411  | 631  | 47311 | 651  | 1271  | 671  | 9311  | 691  | 1471  |
| 731  | 14111 | 751  | 12141 | 771  | 15411 | 791  | 9451  | 811  | 8291  | 831  | 6911  | 851  | 35671 | 871  | 17491 |
| 911  | 14311 | 931  | 9491  | 951  | 971   | 971  | 10031 | 1011 | 6111  | 1041 | 15111 | 1051 | 17191 | 1071 | 20291 |
| 1111 | 3451  | 1131 | 3431  | 1151 | 30631 | 1171 | 27391 | 1191 | 4131  | 1211 | 4331  | 1231 | 7291  | 1251 | 20191 |
| 1311 | 1451  | 1331 | 4751  | 1351 | 2931  | 1391 | 9251  | 1411 | 1791  | 1431 | 1731  | 1451 | 18991 | 1471 | 5891  |
| 1511 | 5471  | 1531 | 12111 | 1551 | 9791  | 1611 | 9731  | 1631 | 17591 | 1651 | 6951  | 1671 | 19191 | 1691 | 2411  |
| 1711 | 14511 | 1731 | 11451 | 1751 | 7631  | 1811 | 11431 | 1831 | 20371 | 1911 | 12451 | 1931 | 3951  | 1971 | 25431 |
| 2111 | 27971 | 2311 | 4771  | 2041 | 9191  | 2071 | 2771  | 2091 | 9211  | 2111 | 13591 | 2131 | 14731 | 2171 | 4411  |

|      |      |     |      |     |      |     |      |     |      |     |      |     |      |     |      |
|------|------|-----|------|-----|------|-----|------|-----|------|-----|------|-----|------|-----|------|
| 21-1 | 6371 | 221 | 1001 | 221 | 551  | 221 | 401  | 221 | 1261 | 221 | 461  | 231 | 871  | 231 | 1701 |
| 237  | 1811 | 237 | 1811 | 237 | 1811 | 237 | 1811 | 237 | 1811 | 237 | 1811 | 237 | 1811 | 237 | 1811 |
| 251  | 1751 | 251 | 1751 | 251 | 1751 | 251 | 1751 | 251 | 1751 | 251 | 1751 | 251 | 1751 | 251 | 1751 |
| 261  | 1721 | 261 | 1721 | 261 | 1721 | 261 | 1721 | 261 | 1721 | 261 | 1721 | 261 | 1721 | 261 | 1721 |
| 271  | 1671 | 271 | 1671 | 271 | 1671 | 271 | 1671 | 271 | 1671 | 271 | 1671 | 271 | 1671 | 271 | 1671 |
| 281  | 1621 | 281 | 1621 | 281 | 1621 | 281 | 1621 | 281 | 1621 | 281 | 1621 | 281 | 1621 | 281 | 1621 |
| 291  | 1571 | 291 | 1571 | 291 | 1571 | 291 | 1571 | 291 | 1571 | 291 | 1571 | 291 | 1571 | 291 | 1571 |
| 301  | 1521 | 301 | 1521 | 301 | 1521 | 301 | 1521 | 301 | 1521 | 301 | 1521 | 301 | 1521 | 301 | 1521 |
| 311  | 1471 | 311 | 1471 | 311 | 1471 | 311 | 1471 | 311 | 1471 | 311 | 1471 | 311 | 1471 | 311 | 1471 |
| 321  | 1421 | 321 | 1421 | 321 | 1421 | 321 | 1421 | 321 | 1421 | 321 | 1421 | 321 | 1421 | 321 | 1421 |
| 331  | 1371 | 331 | 1371 | 331 | 1371 | 331 | 1371 | 331 | 1371 | 331 | 1371 | 331 | 1371 | 331 | 1371 |
| 341  | 1321 | 341 | 1321 | 341 | 1321 | 341 | 1321 | 341 | 1321 | 341 | 1321 | 341 | 1321 | 341 | 1321 |
| 351  | 1271 | 351 | 1271 | 351 | 1271 | 351 | 1271 | 351 | 1271 | 351 | 1271 | 351 | 1271 | 351 | 1271 |
| 361  | 1221 | 361 | 1221 | 361 | 1221 | 361 | 1221 | 361 | 1221 | 361 | 1221 | 361 | 1221 | 361 | 1221 |
| 371  | 1171 | 371 | 1171 | 371 | 1171 | 371 | 1171 | 371 | 1171 | 371 | 1171 | 371 | 1171 | 371 | 1171 |
| 381  | 1121 | 381 | 1121 | 381 | 1121 | 381 | 1121 | 381 | 1121 | 381 | 1121 | 381 | 1121 | 381 | 1121 |
| 391  | 1071 | 391 | 1071 | 391 | 1071 | 391 | 1071 | 391 | 1071 | 391 | 1071 | 391 | 1071 | 391 | 1071 |
| 401  | 1021 | 401 | 1021 | 401 | 1021 | 401 | 1021 | 401 | 1021 | 401 | 1021 | 401 | 1021 | 401 | 1021 |
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| 421  | 921  | 421 | 921  | 421 | 921  | 421 | 921  | 421 | 921  | 421 | 921  | 421 | 921  | 421 | 921  |
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| 471  | 671  | 471 | 671  | 471 | 671  | 471 | 671  | 471 | 671  | 471 | 671  | 471 | 671  | 471 | 671  |
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| 491  | 571  | 491 | 571  | 491 | 571  | 491 | 571  | 491 | 571  | 491 | 571  | 491 | 571  | 491 | 571  |
| 501  | 521  | 501 | 521  | 501 | 521  | 501 | 521  | 501 | 521  | 501 | 521  | 501 | 521  | 501 | 521  |
| 511  | 471  | 511 | 471  | 511 | 471  | 511 | 471  | 511 | 471  | 511 | 471  | 511 | 471  | 511 | 471  |
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| 581  | 121  | 581 | 121  | 581 | 121  | 581 | 121  | 581 | 121  | 581 | 121  | 581 | 121  | 581 | 121  |
| 591  | 71   | 591 | 71   | 591 | 71   | 591 | 71   | 591 | 71   | 591 | 71   | 591 | 71   | 591 | 71   |
| 601  | 21   | 601 | 21   | 601 | 21   | 601 | 21   | 601 | 21   | 601 | 21   | 601 | 21   | 601 | 21   |

## DEGREE 14 380 PAIRS

|      |        |      |       |       |       |       |       |       |       |       |       |       |       |       |       |       |       |
|------|--------|------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|
| 11   | 11     | 51   | 3271  | 71    | 2341  | 111   | 1501  | 131   | 1331  | 171   | 2091  | 191   | 8691  | 231   | 7131  | 251   | 9831  |
| 21   | 441    | 31   | 5291  | 31    | 4491  | 371   | 441   | 411   | 1201  | 471   | 3931  | 491   | 3151  | 531   | 9171  | 551   | 2381  |
| 31   | 3951   | 41   | 1841  | 61    | 2771  | 471   | 241   | 711   | 9231  | 731   | 11231 | 771   | 1911  | 791   | 2091  | 831   | 33751 |
| 41   | 5501   | 91   | 3521  | 91    | 1621  | 951   | 1871  | 971   | 35471 | 1011  | 14071 | 1031  | 1691  | 1071  | 29291 | 1091  | 1551  |
| 1131 | 1451   | 1151 | 461   | 1191  | 1521  | 1211  | 6771  | 1251  | 42631 | 1311  | 1931  | 1331  | 1351  | 1371  | 25131 | 1391  | 5731  |
| 1431 | 14371  | 1491 | 5351  | 1511  | 2171  | 1571  | 1781  | 1611  | 9191  | 1631  | 11171 | 1671  | 18671 | 1711  | 9471  | 1751  | 42771 |
| 1791 | 17391  | 1811 | 24871 | 1831  | 1511  | 1871  | 2631  | 1911  | 9431  | 1931  | 16131 | 1971  | 40751 | 1991  | 2471  | 2031  | 35511 |
| 2051 | 2411   | 2111 | 2471  | 2211  | 27431 | 2231  | 5511  | 2271  | 3611  | 2291  | 17171 | 2331  | 26411 | 2351  | 7671  | 2391  | 6171  |
| 2511 | 7191   | 2531 | 16191 | 2551  | 4331  | 2601  | 12791 | 2711  | 6651  | 2751  | 12131 | 2771  | 24291 | 2811  | 14911 | 2831  | 71331 |
| 2871 | 37991  | 2931 | 25511 | 2951  | 7271  | 2951  | 6111  | 2991  | 15891 | 3051  | 12211 | 3071  | 30491 | 3111  | 7911  | 3131  | 13111 |
| 3171 | 1651   | 3191 | 9791  | 3231  | 37591 | 3251  | 54931 | 3291  | 56271 | 3311  | 27671 | 3371  | 25491 | 3411  | 60971 | 3431  | 34391 |
| 3471 | 1341   | 3531 | 16431 | 3551  | 1721  | 3551  | 40151 | 3591  | 12351 | 3671  | 6711  | 3711  | 19431 | 3731  | 57131 | 3771  | 26511 |
| 3751 | 33351  | 3831 | 6371  | 3911  | 29751 | 3951  | 9711  | 3971  | 26931 | 4011  | 33911 | 4071  | 14231 | 4071  | 24631 | 4091  | 6011  |
| 4131 | 11511  | 4151 | 7551  | 4191  | 19991 | 4211  | 13291 | 4251  | 11951 | 4271  | 37931 | 4311  | 13111 | 4371  | 19331 | 4391  | 8631  |
| 4451 | 117991 | 4511 | 14511 | 4551  | 19251 | 4571  | 4751  | 4611  | 50151 | 4671  | 23511 | 4701  | 15511 | 4711  | 31991 | 4731  | 19931 |
| 4871 | 29271  | 4911 | 15511 | 4931  | 13631 | 4971  | 13791 | 4991  | 8731  | 5031  | 6451  | 5051  | 19791 | 5091  | 12071 | 5111  | 76631 |
| 5331 | 16311  | 5351 | 24231 | 5391  | 13071 | 5411  | 14991 | 5471  | 7491  | 5531  | 29331 | 5571  | 7791  | 5631  | 27731 | 5691  | 7811  |
| 5751 | 20231  | 5811 | 14191 | 5831  | 14931 | 5871  | 13231 | 5911  | 9771  | 5951  | 15971 | 5991  | 14531 | 6031  | 4531  | 6071  | 37551 |
| 6131 | 13511  | 6191 | 30671 | 6211  | 9731  | 6251  | 30671 | 6291  | 16931 | 6311  | 5571  | 6471  | 37731 | 6491  | 20031 | 6531  | 11291 |
| 6551 | 35831  | 6591 | 7111  | 6611  | 11451 | 6671  | 29531 | 6701  | 9611  | 6731  | 18231 | 6801  | 26871 | 6811  | 4651  | 6831  | 37011 |
| 7071 | 16411  | 7111 | 17531 | 7131  | 64791 | 7171  | 18771 | 7191  | 19671 | 7251  | 9911  | 7251  | 95391 | 7311  | 17211 | 7331  | 39471 |
| 7531 | 11711  | 7571 | 17351 | 7591  | 23591 | 7611  | 11411 | 7631  | 24271 | 7651  | 15071 | 7671  | 10951 | 7731  | 9711  | 7751  | 17891 |
| 7931 | 3411   | 7971 | 17751 | 8011  | 14451 | 8091  | 20051 | 8111  | 14551 | 8151  | 9851  | 8211  | 13371 | 8231  | 27471 | 8271  | 14471 |
| 8491 | 13611  | 8551 | 19271 | 8571  | 15591 | 8591  | 14451 | 8611  | 15071 | 8631  | 27471 | 8711  | 14391 | 8731  | 11391 | 8751  | 17671 |
| 8971 | 1971   | 8991 | 24771 | 9011  | 55511 | 9051  | 14451 | 9071  | 14271 | 9111  | 12971 | 9131  | 35351 | 9171  | 17491 | 9251  | 25331 |
| 9311 | 12671  | 9351 | 34671 | 9411  | 11031 | 9491  | 14331 | 9511  | 14291 | 9551  | 13211 | 9591  | 10591 | 9651  | 14351 | 9671  | 70311 |
| 9951 | 14631  | 9971 | 31071 | 10031 | 30391 | 10071 | 15631 | 10091 | 29391 | 10151 | 24051 | 10191 | 30071 | 10931 | 54711 | 10991 | 14251 |

## DEGREE 15 904 PAIRS

|             |             |             |             |             |             |             |             |             |
|-------------|-------------|-------------|-------------|-------------|-------------|-------------|-------------|-------------|
| 11091 53151 | 11091 53831 | 11111 24111 | 11151 24851 | 11171 40931 | 11271 13311 | 11331 19911 | 11351 34211 | 11471 26711 |
| 11631 16311 | 11751 17151 | 11771 19671 | 11811 15191 | 11831 17591 | 11871 27191 | 11891 13091 | 11931 13871 | 11991 32931 |
| 12111 16371 | 12141 30591 | 12211 17911 | 12251 14451 | 12291 14911 | 12311 37611 | 12371 60131 | 12411 17031 | 12431 16631 |
| 12531 39651 | 12551 13241 | 12591 30711 | 12611 32971 | 12711 14051 | 12731 61391 | 12771 20131 | 13011 55031 | 13151 16571 |
| 13131 19751 | 13271 19531 | 13571 38131 | 13671 39211 | 13691 25491 | 13731 20951 | 13751 40871 | 13791 19611 | 13951 98791 |
| 13911 24151 | 14131 15331 | 14211 17411 | 14351 71591 | 14391 20271 | 14591 14711 | 14631 23631 | 14591 17111 | 14591 38031 |
| 14531 58591 | 14531 17651 | 15171 27331 | 15251 26751 | 15291 17051 | 15351 23471 | 15551 24491 | 15591 19531 | 16151 20391 |
| 14571 57491 | 14991 55111 | 15551 19131 | 15791 67791 | 15851 59991 | 15871 23971 | 15951 23871 | 17271 26471 | 17151 23891 |
| 17271 24731 | 17511 29111 | 17571 17811 | 17711 29031 | 17831 30231 | 18311 54571 | 18351 24911 | 18471 38231 | 18531 30371 |
| 14551 23411 | 18711 36451 | 19011 26631 | 19071 34451 | 19091 26691 | 19141 71931 | 19491 37431 | 19571 61071 | 19571 61071 |
| 19731 24531 | 20031 23751 | 20111 37071 | 20291 26751 | 20351 55571 | 20451 35171 | 20471 54511 | 21571 35811 | 21571 35811 |
| 21131 25331 | 21051 25251 | 24191 34151 | 24311 25551 | 24531 60711 | 24551 29831 | 24591 24451 | 24711 33251 | 24791 40811 |
| 24311 54331 | 24571 57571 | 24911 24631 | 24971 71511 | 25571 34151 | 26771 59631 | 27171 34991 | 27311 27311 | 27311 27311 |
| 27351 37311 | 27611 60231 | 27711 55091 | 27831 37971 | 27911 47211 | 27971 58871 | 28611 76791 | 29091 35031 | 29091 35031 |
| 29351 37311 | 29711 63111 | 29931 56111 | 30051 39311 | 32451 35651 | 33831 38211 | 34131 34131 | 37031 70391 | 37571 61271 |
| 40611 40111 | 61311 81211 |             |             |             |             |             |             |             |

|            |             |            |            |             |             |             |            |            |
|------------|-------------|------------|------------|-------------|-------------|-------------|------------|------------|
| 11 11      | 3110231     | 51 45531   | 91 36411   | 111 29791   | 131 25231   | 151 21851   | 171 10351  | 191 17231  |
| 31 14251   | 251 13111   | 271 15171  | 291 11311  | 311 9931    | 371 11071   | 391 4611    | 411 9991   | 431 7651   |
| 451 7311   | 471 35731   | 511 6451   | 531 24731  | 551 6331    | 571 5751    | 591 27771   | 611 32231  | 631 35291  |
| 671 14711  | 691 4751    | 711 32311  | 731 46411  | 751 4371    | 791 37331   | 811 28571   | 831 19771  | 851 3071   |
| 471 33771  | 491 121511  | 511 3451   | 531 47431  | 551 4311    | 1011 22711  | 1031 9451   | 1071 9191  | 1091 9771  |
| 1111 3691  | 1131 14511  | 1151 2451  | 1171 70191 | 1211 111031 | 1231 3331   | 1251 18471  | 1271 16251 | 1291 25531 |
| 1311 44151 | 1331 55831  | 1371 7251  | 1391 32651 | 1411 11911  | 1431 6911   | 1451 6851   | 1491 11031 | 1531 2151  |
| 1571 14011 | 1591 22671  | 1631 70391 | 1651 2111  | 1671 13751  | 1691 17451  | 1711 55571  | 1731 24631 | 1751 14511 |
| 1791 9211  | 1811 53311  | 1831 19711 | 1851 19491 | 1871 75351  | 1911 1931   | 1931 78991  | 1971 4991  | 1991 62931 |
| 2011 91511 | 2031 49831  | 2071 75191 | 2091 74791 | 2131 10771  | 2191 14671  | 2211 81571  | 2231 23511 | 2251 27411 |
| 2271 10111 | 2291 47231  | 2331 23911 | 2351 18171 | 2371 31811  | 2391 50771  | 2411 34711  | 2431 17531 | 2451 49231 |
| 2491 4591  | 2511 45511  | 2531 34071 | 2551 13811 | 2591 44591  | 2651 9811   | 2671 60331  | 2691 32471 | 2711 16031 |
| 2751 4351  | 2771 7571   | 2811 13411 | 2831 38211 | 2891 23811  | 2911 65671  | 2931 6711   | 2951 23331 | 2971 18791 |
| 3011 23591 | 3031 7571   | 3051 9671  | 3071 7031  | 3091 50371  | 3111 116951 | 3131 9431   | 3171 38411 | 3191 27031 |
| 3211 12551 | 3231 35511  | 3251 37431 | 3271 7031  | 3291 28411  | 3371 6811   | 3391 46731  | 3431 12431 | 3451 84531 |
| 3491 54711 | 3511 57411  | 3531 58151 | 3551 4391  | 3611 39031  | 3631 37011  | 3651 19751  | 3671 24111 | 3731 28991 |
| 3751 49731 | 3791 23151  | 3811 91811 | 3831 47911 | 3891 39591  | 3911 49871  | 3931 16051  | 3951 10811 | 3971 4231  |
| 4011 54851 | 4051 18411  | 4071 39451 | 4091 38851 | 4111 27911  | 4131 37111  | 4171 14931  | 4191 13641 | 4211 19471 |
| 4251 62531 | 4291 34191  | 4311 29651 | 4331 15271 | 4371 36211  | 4391 27111  | 4411 79191  | 4451 30191 | 4471 8431  |
| 4491 4711  | 4511 59851  | 4531 60991 | 4551 47231 | 4591 26311  | 4631 29371  | 4651 18951  | 4711 4471  | 4731 32571 |
| 4771 58491 | 4791 10471  | 4811 33831 | 4831 14231 | 4891 81751  | 4911 22091  | 4931 9971   | 4951 70431 | 5011 40551 |
| 5031 35531 | 5051 47911  | 5071 7191  | 5091 92631 | 5111 47171  | 5131 15691  | 5191 30391  | 5211 14271 | 5231 78631 |
| 5251 11731 | 5271 5471   | 5291 33191 | 5311 43511 | 5331 6571   | 5351 25771  | 5391 28731  | 5411 54711 | 5431 75231 |
| 5451 34171 | 5471 81131  | 5491 27331 | 5511 79791 | 5531 17791  | 5551 55471  | 5591 54931  | 5611 77251 | 5631 24331 |
| 5671 34231 | 5691 112551 | 5711 17131 | 5731 20311 | 5791 46191  | 5811 26171  | 5831 27171  | 5851 23291 | 5871 32931 |
| 5911 13771 | 5931 6751   | 5951 16551 | 5971 59391 | 6011 16411  | 6031 40631  | 6051 80231  | 6071 74131 | 6091 60931 |
| 6131 40551 | 6151 72511  | 6171 50451 | 6191 31911 | 6231 12631  | 6251 19571  | 6271 96431  | 6311 51091 | 6331 27191 |
| 6371 15231 | 6391 98531  | 6411 53971 | 6431 30251 | 6471 17911  | 6491 59571  | 6511 140711 | 6531 29011 | 6571 20451 |
| 6591 2971  | 6611 7951   | 6631 55951 | 6651 40631 | 6671 76591  | 6711 35191  | 6731 40551  | 6751 27491 | 6771 75291 |
| 6811 55091 | 6831 22211  | 6851 34491 | 6871 42111 | 6891 34951  | 6931 33971  | 6951 16231  | 6971 67631 | 7011 70931 |
| 7031 15111 | 7051 29711  | 7071 33751 | 7091 42151 | 7131 33751  | 7151 7851   | 7171 28731  | 7191 20331 | 7211 70931 |
| 7231 47371 | 7251 117751 | 7271 29711 | 7291 8571  | 7311 66191  | 7331 40911  | 7351 27371  | 7371 46731 | 7391 47031 |
| 7411 26771 | 7431 1751   | 7451 64311 | 7471 47191 | 7491 55991  | 7511 75111  | 7531 76631  | 7551 40471 | 7571 32731 |
| 7591 11591 | 7611 11591  | 7631 41251 | 7651 12151 | 7671 15911  | 7691 39151  | 7711 18311  | 7731 49171 | 7751 49171 |
| 7771 13431 | 7791 26471  | 7811 54551 | 7831 29251 | 7851 54551  | 7871 39151  | 7891 32291  | 7911 80711 | 7931 12731 |
| 7951 29161 | 7971 110671 | 8011 14951 | 8031 14951 | 8051 81111  | 8071 57611  | 8091 57611  | 8111 57611 | 8131 14751 |

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