

Mixed-Trust Scheduling in AADL

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Mixed-Trust Scheduling in AADL

Modeling Mixed-Trust Scheduling in AADL

- Modeling Mixed-Trust Processors
- Modeling Mixed-Trust Tasks

OSATE Mixed-Trust Scheduling Plugin

Modeling Mixed-Trust Processors

To execute a mixed-trust task, a processor must have

- A Guest OS
- A HyperVisor

The Guest OS and HyperVisor are best represented in AADL using `virtual processor` components:

- A virtual processor “[represents] a virtual machine such as the Java VM, a partition of a processor, or a scheduler”
[*Model-Based Engineering with AADL*, Feiler & Gluch].

Guest OS and HyperVisor as Virtual Processors

A **virtual processor** component must be **bound** to a **processor** component.

- This is an AADL language requirement.

Two ways to declare binding:

- Use the **Actual_Processor_Binding** property on the virtual processors.
 - Usually done using a contained property association.
- Declare the virtual processors as subcomponents of the processor.

Example: Binding Virtual Processors

```

system implementation Example1.impl
subcomponents
  P: processor;
  GuestOS: virtual processor;
  HyperVisor: virtual processor;

properties
  Actual_Processor_Binding =>
    (reference(P)) applies to
      GuestOS, HyperVisor;
end Example1.impl;

```

Explicit binding to processor.

```

processor implementation MTP.i
subcomponents
  GuestOS: virtual processor;
  HyperVisor: virtual processor;
end MixedTrustProcessor.i;

```

Subcomponents are *implicitly* bound to the containing processor.

```

system implementation Example2.impl
subcomponents
  P: processor MTP.i;
end Example2.impl;

```

Declaring a Mixed-Trust Processor

A processor with bound virtual processors is not necessary meant for mixed-trust scheduling.

- Virtual processors are used to represent many things.

Must **declare** that the processor is **intended for mixed-trust scheduling**.

- Use a property association to do this.
 - Indicates which virtual processor is
 - The Guest OS.
 - The HyperVisor.

Mixed-Trust Processor Property

```
property set Mixed_Trust_Properties is

  Mixed_Trust_Bindings: type record (
    GuestOS: reference (virtual processor);
    HyperVisor: reference (virtual processor);
  );

  Mixed_Trust_Processor:
    Mixed_Trust_Properties::Mixed_Trust_Bindings
    applies to (processor);

  -- . . .
end Mixed_Trust_Properties;
```

Example: Mixed-Trust Processor Property

```
system implementation Example1.impl
```

```
subcomponents
```

```
  P: processor;
```

```
  GuestOS: virtual processor;
```

```
  HyperVisor: virtual processor;
```

```
properties
```

```
  Mixed_Trust_Properties::Mixed_Trust_Processor =>
```

```
    [GuestOS => reference(GuestOS);
```

```
     HyperVisor => reference(HyperVisor);]
```

```
  applies to P;
```

```
  Actual_Processor_Binding =>
```

```
    (reference(P)) applies to GuestOS, HyperVisor;
```

```
end Example1.impl;
```

Modeling Mixed-Trust Tasks — Pairing Threads

A mixed-trust task is fundamentally a pair of threads:

- One executes on the Guest OS.
- One executes on the HyperVisor.

Use standard AADL thread binding mechanism to assign threads to virtual processors:

- The `Actual_Processor_Binding` property (again).

But, need to know **which pair of threads** corresponds to a single mixed-trust task.

- Again, **declare** this using a property association.

Modeling Mixed-Trust Tasks — Timing

Each mixed-trust task has timing information:

- **Period**
 - Applies to the mixed-trust task, not just a single thread
 - Standard AADL `period` property would need to be consistent across the thread pair.
- **Deadline**
 - Applies to the mixed-trust task, not just a single thread
 - Standard AADL `period` property would need to be consistent across the thread pair.
- **Execution times**
 - Applies to the individual threads
 - Use the standard AADL `compute_execution_time` property.

Mixed-Trust Tasks Property

```
property set Mixed_Trust_Properties is
  Mixed_Trust_Task: type record (
    Name: aadlstring;
    Period: Time;
    Deadline: Time;
    GuestTask: reference (Thread);
    HyperTask: reference (Thread);
    E: Time;
  );
```

Used in place of the
Period and Deadline
properties on the
threads themselves.

Really an output of the scheduler.
Ideally this would be automatically
inserted into the instance model.

```
Mixed_Trust_Tasks:
  list of Mixed_Trust_Properties::Mixed_Trust_Task
  applies to (System);
```

```
-- . . .
```

```
end Mixed_Trust_Properties;
```

Example: Mixed-Trust Tasks Property (1)

```
process MTT
end MixedTrustTask;
```

```
process implementation MTT.basic
  subcomponents
    GuestThread: thread;
    HyperThread: thread;
end MTT.basic;
```

```
process implementation MTT.t1
  subcomponents
    GuestThread: thread {
      Compute_Execution_Time => 4ms .. 4ms;
    };
    HyperThread: thread {
      Compute_Execution_Time => 0ms .. 0ms;
    };
end MixedTrustTask.t1;
```

```
process implementation MTT.t2
  subcomponents
    GuestThread: thread {
      Compute_Execution_Time => 2ms .. 2ms;
    };
    HyperThread: thread {
      Compute_Execution_Time => 3ms .. 3ms;
    };
end MixedTrustTask.t2;
```

Current approach **does not** require the two threads to be in the same process, but it is a convenient abstraction.

Example: Mixed-Trust Tasks Property (2)

```
system implementation Example1.impl
```

```
  subcomponents
```

```
    P: processor;
```

```
    GuestOS: virtual processor;
```

```
    HyperVisor: virtual processor;
```

```
    MTT1: process MTT.task1;
```

```
    MTT2: process MTT.task2;
```

```
  properties
```

```
    Actual_Processor_Binding => (reference(P))
```

```
    applies to GuestOS, HyperVisor;
```

```
    Mixed_Trust_Properties::Mixed_Trust_Processor =>
```

```
      [GuestOS => reference(GuestOS);
```

```
        HyperVisor => reference(HyperVisor);]
```

```
    applies to P;
```

```
    Actual_Processor_Binding => (reference(GuestOS))
```

```
    applies to MTT1.GuestThread, MTT2.GuestThread;
```

```
    Actual_Processor_Binding => (reference(HyperVisor))
```

```
    applies to MTT1.HyperThread, MTT2.HyperThread;
```

```
    Mixed_Trust_Properties::Mixed_Trust_Tasks =>
```

```
      ([Name => "MT 1";
```

```
        Period => 8 ms;
```

```
        Deadline => 8 ms;
```

```
        GuestTask => reference(MTT1.GuestThread);
```

```
        HyperTask => reference(MTT1.HyperThread);],
```

```
      [Name => "MT 2";
```

```
        Period => 14 ms;
```

```
        Deadline => 14 ms;
```

```
        GuestTask => reference(MTT2.GuestThread);
```

```
        HyperTask => reference(MTT2.HyperThread);]);
```

```
  end Example1.impl;
```

OSATE Analysis Plug-In

OSATE Plug-in builds an analysis around Dio's scheduling code.

- Based on the modeling approach just described.
- Performs consistency checking on the model.

Input is an AADL instance model.

- Checks each system operation mode for schedulability.
- Outputs results to a CSV file.
 - Schedulable?
 - E values.

Consistency Checks

The analysis checks consistency based on

- Model structure.
- Mixed-trust property associations.

Reports problems using Eclipse error markers.

For example

- A `Mixed_Trust_Processor` property association must provide values for both the `GuestOS` and `HyperVisor` fields.
 - (AADL does not require this.)
- The virtual processors referenced by the `Mixed_Trust_Processor` fields must be bound to the associated processor component.

More Consistency Checks

Checking `Mixed_Trust_Tasks` property associations:

- A `GuestTask` thread must be bound to a virtual processor identified as a `GuestOS`.
- A `HyperTask` thread must be bound to a virtual processor identified as a `HyperVisor`.
- The two virtual processors must be bound to the same processor.
 - Cannot split a task across *two different* mixed-trust domains.

There are more checks; this is to just give an idea of what is done.