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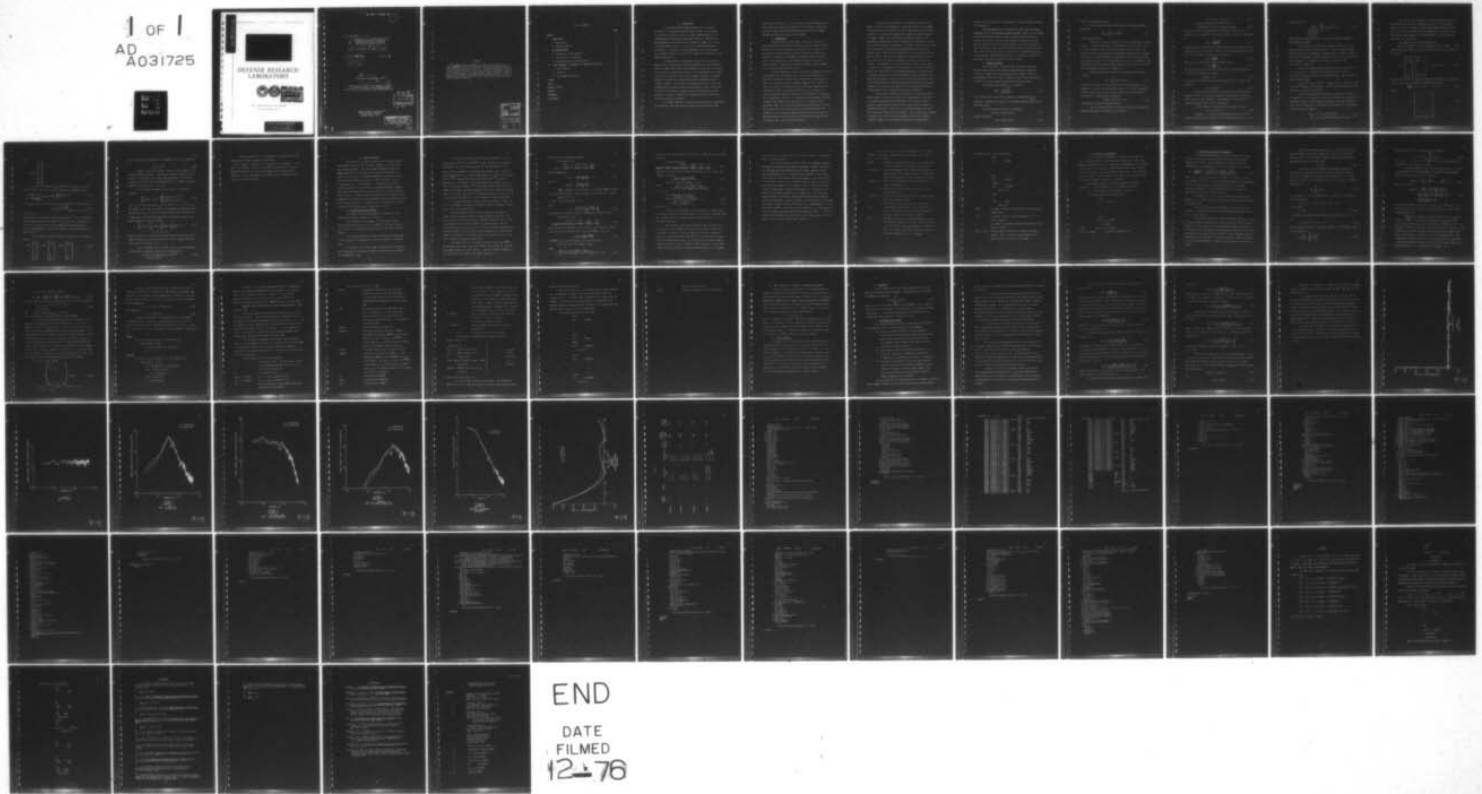
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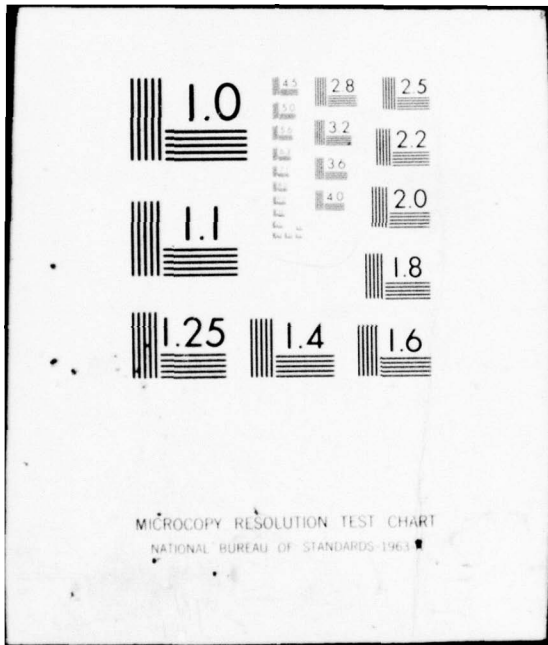
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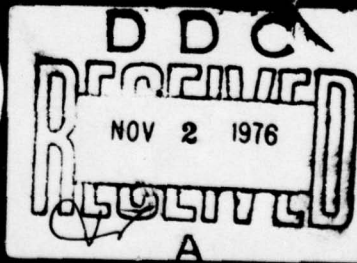
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6 GENERATION OF A STATIONARY GAUSSIAN
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POWER SPECTRAL DENSITY FUNCTION.

10 by
W. A. Matuska, Jr. and G. S. Innis

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ABSTRACT

This paper, which is based on the technique published by Joel N. Franklin, gives a method for generating digitally a time series with a given power spectral density function. A computer program to carry out this method was written for the Control Data Corporation 3200 computer, and the program was tried on some specific example problems. Then for each of these examples, the power spectral density function of the generated time series was compared with the specified, theoretical power spectral density function.

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I. INTRODUCTION

A stationary Gaussian random process, $x(t)$, is called a time series if t represents discrete values of time. The collection of functions of time $x(t) = x_r(t)$ is called a random process if r is a random variable ranging over some measure space, R . The process is called Gaussian if, for every finite collection of times $t_1 < \dots < t_n$, the random variables $x(t_1), \dots, x(t_n)$ have a multivariate Gaussian distribution. The process is called stationary if, for any increment Δt , the random variables $x(t_i + \Delta t)$ have the same joint distribution as the random variables $x(t_i)$.¹

A common method for obtaining a time series is to record on magnetic tape the amplitude of a physical quantity, such as the pressure waves in a large body of water, the atmosphere, or the earth or the waves on the surface of a large body of water. This analog data can be thought of as a time series which can be either analyzed by means of an analog computer or converted into digital data and then analyzed by means of a digital computer. This method for obtaining a time series can be time consuming and the conditions which are necessary to obtain a time series with specified statistical properties are often difficult to achieve; therefore, a method is desired for computing numerically a time series with given statistical properties.

This paper summarizes a technique, proposed by Joel N. Franklin², for computing a stationary Gaussian random process with a given power

spectral density function. It also contains a computer program for the Control Data 3200 computer which applies this technique in computing a specified stationary Gaussian random process along with four example problems which have been tried on this program.

A. Applications

The method given in this paper can be used to generate a stationary Gaussian statistical process with a specified power spectral density function which satisfies the conditions (1.8). Forms for the power spectral density function of many statistical processes can be found in the literature; for example, the general form for the power spectral density function of low frequency atmospheric turbulence is given by Bendat.³ Problem four in Chapter III gives a specific example of such a power spectral density function.

As stated previously, a time series can be obtained by recording the amplitude of a statistical process, such as a pressure wave, at regular time intervals. However, when a time series is obtained in the field, the time series which is recorded on magnetic tape, called the total time series, is a combination of the time series obtained from the signal and the time series obtained from the noise. A signal is a detectable physical quantity or impulse by which messages or information can be transmitted, and noise is an unwanted detectable physical quantity or impulse which exists either by itself or in interference with a specified signal. Now if the power spectral density function of either of the three time series mentioned above is known and satisfies the conditions (1.8), a stationary Gaussian random process, which has the same power spectral density function as the time series obtained from the physical wave, can be generated digitally.

A problem of much interest is the study of detecting a known signal in various noise backgrounds. By the method given in this paper a time series can be generated which has the same power spectral density function as the time series obtained from a specified noise background. For example, suppose the time series of a noise-free signal of the infrasound (pressure waves having a frequency lower than about 16 cycles per second) produced by a thunderstorm is obtained. Then this time series can be viewed under any specified background conditions if the power spectral density function of the time series of the specified noise background is known. This can be accomplished by combining these two time series with the use of either the digital or analog facilities of a computer. The magnitude of the time series of the signal can be varied in relation to the magnitude of the time series of the noise, and by this method, the time series of the infrasound produced by a thunderstorm at a specified distance with a specified background can be simulated. The power spectral density functions of these total time series can be used in developing methods for detecting and tracking thunderstorms by giving examples of the power spectral densities of the infrasound produced by thunderstorms at various distances with various backgrounds.

The ability to obtain a time series for a given power spectral density function could also be useful in the advancement of the theory of optimum linear prediction and filtering. This theory is used in designing engineering systems which either project into the future by information obtained in the past or recover desired signals which have been distorted by random noise disturbances. These systems can be applied to communication, meteorological forecasting, and economic analysis. A specific example is the technique of combining two independently

obtained sets of related perturbed messages to obtain an improved estimate of the message.⁴

The last application to be mentioned is the study of smoothing techniques such as the hanning and hamming windows.⁵ Here it is of interest to learn how the minor lobes of a power spectral density function are affected by the major lobe when these smoothing techniques are applied. This can be done by constructing time series for which the distances between the major and minor lobes of their power spectral density functions vary. The relative size of these lobes can also be varied along with the distance between them; therefore the effect of a smoothing technique on a given power spectral density function can be studied.

B. Summary of Method

The method under consideration for computing a stationary Gaussian random process is the method proposed by Joel N. Franklin.⁶ In order to compute such a random process by this method, either the autocorrelation function, which is defined in terms of the expected value operator, or the power spectral density function must be known.

The expected value operator E is defined as

$$E[x_i] = \lim_{N \rightarrow \infty} \frac{1}{N} \sum_{i=1}^N x_i \quad (1.1)$$

if the limit exists. For a finite sequence x_1 , the operation $\lim_{N \rightarrow \infty}$ is neglected. Throughout this paper it is assumed that $E[x(t_i)] = 0$ for each time series $x(t)$.

The autocorrelation function is defined as

$$R(t_1, t_2) = E[x(t_1)x(t_2)] \quad (1.2)$$

which is written as

$$R(\tau) = E[x(t)x(t-\tau)] \quad (1.3)$$

if $x(t)$ is a stationary process.

The power spectral density function for a stationary process is defined as

$$V(\omega) = \int_{-\infty}^{\infty} R(\tau) e^{-i\omega\tau} d\tau \quad (1.4)$$

if the integral exists.

The Fourier transform indicated by (1.4) can generally be inverted; therefore, if $V(\omega)$ is known and can be inverted, $R(\tau)$ can be found. Since it is more natural to work with the power spectral density function and since $V(\omega)$ and $R(\tau)$ are generally interchangeable, this method is designed to construct a stationary Gaussian random process from a given power spectral density function. Even though $V(\omega)$ and $R(\tau)$ are interchangeable, small perturbations in $V(\omega)$ may cause great changes in $R(\tau)$. This case is illustrated by example problem four in Chapter III.

Franklin attacks the problem by considering the linear transformation

$$x(t) = \int_{-\infty}^t g(t-s)v(s)ds \quad (1.5)$$

He then sets v equal to the Dirac delta function; therefore $g(t)$ is the response of a filter to a delta function input. It is shown by Davenport and Root⁷ that the power spectral density functions $V_x(\omega)$ and $V_v(\omega)$ of $x(t)$ and $v(t)$, respectively, are related by

$$V_x(\omega) = |G(\omega)|^2 V_v(\omega) \quad (1.6)$$

where $G(\omega)$ is the Fourier transform of $g(t)$.

Now if a random process $v(t)$, known as white noise which has the property that $V_v(\omega)=1$, is obtained, (1.6) becomes

$$V(\omega) = V_x(\omega) = |G(\omega)|^2 \quad (1.7)$$

It is shown by Davenport and Root⁸ that if

$$0 < V(\omega) < \infty, V(\omega) = V(-\omega), \text{ and } V(\omega) \rightarrow 0 \text{ as } \omega \rightarrow \pm\infty, \quad (1.8)$$

and that if $V(\omega)$ can be represented as a rational function (i.e., the quotient of two polynomials) in ω , then for real ω , $V(\omega)$ can be represented by

$$V(\omega) = \left| \frac{P(i\omega)}{Q(i\omega)} \right|^2, \quad (1.9)$$

where P and Q are polynomials with real coefficients and of degree m and n , respectively, $m < n$, and the zeros of $Q(\xi)$ lie in the half-plane $\text{Re } \xi < 0$. Now by (1.7) and (1.9), an obvious choice for $G(\omega)$ is

$$G(\omega) = \frac{P(i\omega)}{Q(i\omega)}, \quad (1.10)$$

and if D is the differential operator, d/dt , we have

$$x(t) = \frac{P(D)}{Q(D)} v(t) \quad (1.11)$$

The specified time series $x(t)$ is found by first finding a steady-state solution to the differential equation

$$Q(D)\varphi(t) = v(t), \quad (1.12)$$

for all real t and then combining the derivatives of $\varphi(t)$ of order less than n by

$$x(t) = P(D)\varphi(t) \quad (1.13)$$

The first step in finding a solution $\varphi(t)$ to the differential equation (1.12) is to generate a completely equidistributed sequence v_n which is also a "white sequence." This sequence is to be used as the digital representation of a white noise source.

A "white sequence" is defined to be a sequence for which

$$R(\tau) = 0 \text{ for } \tau \neq 0 \quad (1.14)$$

A sequence v_n is said to be equidistributed by k 's if k is a positive integer and if for every set of k intervals (a_i, b_i) , $i=1, \dots, k$,

where $0 \leq a_i < b_i \leq 1$,

$$\frac{1}{N} \sum_{\substack{a_i \leq v_{n+i} < b_i \\ (i=1, \dots, k) \\ n=1, \dots, N}} 1 \rightarrow \prod_{i=1}^k (b_i - a_i) \text{ as } N \rightarrow \infty, \quad (1.15)$$

and a sequence equidistributed by k 's for all k is called a completely equidistributed sequence.

Let the notation $\{\varphi\}$ denote the fractional part of φ . It is shown by Franklin that for almost every $\theta > 1$, $\{\theta^n\}$ is a completely equidistributed sequence⁹; furthermore he suggests that the sequence $v_n = \{\theta^n\}$, where θ is a transcendental number greater than one, be used as the required "white sequence." For example, if $\theta \approx 2.72$, $v_1 \approx \{2.72\} = .72$, $v_2 \approx \{2.72^2\} = .3984$, etc.

Because of computational difficulties which limit the length of v_n for a given $\theta > 1$, several different sequences, $v_n^{(1)}, \dots, v_n^{(s)}$, ($n = 0, 1, 2, \dots$), may be needed for a single application. These sequences should be interlaced by

$$v_{ns+1} = v_n^{(1)}, v_{ns+2} = v_n^{(2)}, \dots, v_{ns+s} = v_n^{(s)}, n = 0, 1, 2, \dots \quad (1.16)$$

to form the required "white sequence." A few of these difficulties will be pointed out in Chapter II, Section A.

Now that an equidistributed sequence v_n has been constructed on the range $0 \leq v < 1$, a sequence w_n , which will be called the w -sequence, of independent samples from the Gaussian distribution with mean 0 and variance 1 can be constructed by the method of Box and Muller.¹⁰ This method is given by the two formulas,

$$\begin{aligned} w_{2n-1} &= (-2 \ln v_{2n-1})^{\frac{1}{2}} \cos 2\pi v_{2n} \\ w_{2n} &= (-2 \ln v_{2n-1})^{\frac{1}{2}} \sin 2\pi v_{2n} \text{ for } n=1, 2, 3, \dots \end{aligned} \quad (1.17)$$

Now that the "white sequence," v_n , and the w-sequence, w_n , have been defined, the next step in finding a solution to the differential equation (1.12) is to consider the given power spectral density function $V(\omega)$, which must be a rational function satisfying the conditions (1.8), and the given time increment Δt at which samples of the specified time series are to be calculated.

The equation (1.12) can be represented as

$$a_1 \varphi^{(n)}(t) + a_2 \varphi^{(n-1)}(t) + \dots + a_{n+1} \varphi(t) = v(t) \text{ for } -\infty < t < \infty \quad (1.18)$$

where $a_1=1$ and a_i , $1 \leq i \leq n+1$, are the real coefficients of the polynomial Q .

The vectors

$$z(t) = \begin{bmatrix} z_1(t) \\ z_2(t) \\ \cdot \\ \cdot \\ \cdot \\ z_n(t) \end{bmatrix} = \begin{bmatrix} \varphi(t) \\ \varphi'(t) \\ \cdot \\ \cdot \\ \cdot \\ \varphi^{(n-1)}(t) \end{bmatrix}, \quad t = 0, \Delta t, 2\Delta t, \dots \quad (1.19)$$

must be found in order to compute the specified time series by the equation (1.13). The vector z satisfies the differential equation

$$\frac{dz(t)}{dt} = A z(t) + f(t), \quad -\infty < t < \infty \quad (1.20)$$

where

$$A = \begin{bmatrix} 0 & 1 & 0 & \dots & 0 \\ 0 & 0 & 1 & \dots & 0 \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ 0 & 0 & 0 & \dots & 1 \\ -a_{n+1} & -a_n & -a_{n-1} & \dots & -a_2 \end{bmatrix} \quad (1.21)$$

and

$$f(t) = \begin{bmatrix} 0 \\ 0 \\ \cdot \\ \cdot \\ 0 \\ v(t) \end{bmatrix} \quad (1.22)$$

In order to find $z(0)$, the solution vector (m_1, m_2, \dots, m_n) must first be found to the set of n linear equations in n unknowns given by

$$(-1)^{k-1} \sum_{\substack{k \\ 2} + 1 \leq q \leq \frac{(n+k-1)}{2} + 1} (-1)^{q-1} a_{n-2(q-1)+k} m_q = \begin{cases} 0, & k=1, \dots, n-1 \\ 1/2, & k=n \end{cases} \quad (1.23)$$

The elements of this solution vector are then placed into the matrix M by

$$m_{ij} = \begin{cases} 0, & i+j \text{ odd} \\ (-1)^{(j-i)/2} m_{(j+i)/2}, & i+j \text{ even} \end{cases} \quad (1.24)$$

for $1 \leq i, j \leq n$.

The formulas (1.23) and (1.24) are proved by Franklin.¹¹ M is the positive definite moment matrix for the vector $z(t)$ for all time t , including $t = 0$. Since M is a positive definite, real, symmetric matrix, there exists a lower-triangular matrix T with positive elements on the main diagonal such that $M = T T^*$.¹²

Now it becomes necessary to define the sequence of n -dimensional vectors,

$$w^{(0)} = \begin{bmatrix} w_1 \\ \cdot \\ \cdot \\ \cdot \\ w_n \end{bmatrix}, \quad w^{(1)} = \begin{bmatrix} w_{n+1} \\ \cdot \\ \cdot \\ \cdot \\ w_{2n} \end{bmatrix}, \quad w^{(2)} = \begin{bmatrix} w_{2n+1} \\ \cdot \\ \cdot \\ \cdot \\ w_{3n} \end{bmatrix}, \quad \dots \quad (1.25)$$

whose elements are the elements of the w -sequence, and $z(0)$ is determined by

$$z(0) = T w^{(0)} \quad (1.26)$$

Franklin shows that if $z(k\Delta t)$ is known, the following steps are needed to compute $z((k+1)\Delta t)$. $z(0)$ has been computed in equation (1.26); therefore, it is used as the starting value and $z((k+1)\Delta t)$ is found by induction. The first step in finding $z((k+1)\Delta t)$ is to compute the matrix $\exp(\Delta t A)$ with elements e_{ij} , $1 \leq i, j \leq n$.

It then becomes necessary to solve the set of simultaneous equations

$$\sum_{k=1}^n (a_{ik} \mu_{kj} + a_{jk} \mu_{ik}) = \begin{cases} \hat{b}_{ij} = e_{in} e_{jn}, & i < n \text{ or } j < n \\ b_{nn} = e_{nn}^2 - 1, & i = n \text{ and } j = n \end{cases} \quad (1.27)$$

where μ_{ij} are the unknowns and a_{ij} are the elements of A for $1 \leq i, j \leq n$. The elements μ_{ij} are also the elements of the symmetric moment matrix M_r ; therefore, $\mu_{ij} = \mu_{ji}$. It also can be seen that $\hat{b}_{ij} = \hat{b}_{ji}$; therefore, the n^2 equations in n^2 unknowns represented by (1.27) can be reduced to the $\frac{1}{2}n(n+1)$ equations in $\frac{1}{2}n(n+1)$ unknowns,

$$\sum_{k \leq j} a_{ik} \mu_{jk} + \sum_{k > j} a_{ik} \mu_{kj} + \sum_{k \leq i} a_{jk} \mu_{ik} + \sum_{k > i} a_{jk} \mu_{ki} = \hat{b}_{ij} \quad (1.28)$$

for $1 \leq j \leq i \leq n$.

When the solution to (1.28) has been found, the positive definite, real, symmetric matrix M_r is known, and therefore a lower triangular matrix T_r can be found such that $M_r = T_r T_r^*$.

T_r , Δt , $\exp(\Delta t A)$, $z(0)$, and $w^{(k)}$, $k=1, 2, \dots$, are now known, and the last step in evaluating $z((k+1)\Delta t)$ is carried out by

$$z((k+1)\Delta t) = \exp(\Delta t A) z(k\Delta t) + T_r w^{(k+1)} \quad (1.29)$$

for $k = 0, 1, 2, \dots$

Now the solution $\varphi(t)$ to the equation (1.12) and its first $(n-1)$ derivatives are known. Equation (1.13) becomes

$$x(t) = b_1 z_{m+1}(t) + b_2 z_m(t) + \dots + b_{m+1} z_1(t), \quad t = 0, \Delta t, 2\Delta t, \dots \quad (1.30)$$

where $b_1 \neq 0$ and b_i , $1 \leq i \leq m+1$, are the real coefficients of the polynomial P and $z_i(t)$, $1 \leq i \leq n$, are the elements of the vector $z(t)$ as defined by (1.19). This completes the process since the specified time series $x(t)$ has been obtained.

II. COMPUTER PROGRAMS

The method for calculating a time series with specified power spectral density function, as presented in Chapter I, is divided into three separate programs. The first program, ISLASHO with subroutine RANDOM, generates a "white sequence" for a given θ and buffers this sequence out onto tape. The program CLLCNVRT, pronounced Call Convert, with subroutine CONVERT then converts the "white sequence" into the w-sequence by (1.17). This w-sequence is then buffered out onto a second tape which is the input tape to the program GAUSSIAN. GAUSSIAN then shapes the w-sequence into a time series with a given power spectral density function. The power spectral density function is defined in this program by the coefficients of the polynomials P and Q as defined by (1.9). This specified time series is buffered out onto a third tape. The same w-sequence can be shaped into any specified time series.

A. Generating a "White Sequence"

The purpose of the program ISLASHO is to generate a "white sequence" as described in Chapter I, Section B.

For the purpose of this paper, the transcendental number greater than one was chosen to be $e = 2.71828183$. The "white sequence", $v_n = \{e^n\}$ for $1 \leq n \leq 4266$, was then generated on the Control Data 3200 computer.

The following approach was taken to the problem of generating a "white sequence." The approach leads to a description of the method used.

For every $n > 1$, e^n can no longer be stored to full accuracy as an eleven digit, floating point number; therefore, this is obviously not the approach to take.

It can be seen that about 1740 digits are necessary to store the integer part of e^{4000} since $\log(e^{4000}) \approx 1740$. Initially e is rounded to nine digits; then for the assurance that the fractional part of e^n , v_n , will not start repeating, e^n is saved with no round-off error. Since the present approximation of e contains eight digits to the right of the decimal point, $\{e^{4000}\}$ contains 32,000 digits, and therefore e^{4000} consists of approximately 33,740 digits. Now if these digits were stored in groups of threes as fixed point numbers, e^{4000} would require about 12,000 fixed point numbers for storage. The storage capacity of the Control Data 3200 is about 32,000 24-bit words; therefore, if two 12,000-word arrays are needed for the multiplication, the machine still has the capacity to handle the program.

The largest fixed point number that can be stored in a 24-bit word is $2^{23}-1 \approx 8,388,000$; therefore, all six-digit and most seven-digit integers can be stored in one 24-bit word.

Now the properties that the Control Data 3200 can handle two 12,000, 24-bit word arrays and all six-digit integers can be stored in a 24-bit word will be used in the program ISLASHO to raise any seven-, eight-, or nine-digit number to powers and save the fractional part. The two 12,000-word arrays, IX and IY, are used to store the products as three-digit, fixed point numbers. These products should be thought of as numbers represented in the base 1000. The notation used in equations (2.1), (2.4), and (2.7) is that of "long hand" multiplication.

The number which is to be raised to powers is read into ISLASHO as the three, three-digit, fixed point numbers, IX3, IX2, and IX1, with IX1 containing the three right-most digits. In the method used in ISLASHO, we first set $IX(1) = IX1$, $IX(2) = IX2$, $IX(3) = IX3$, and $L = 3$

and then start the following procedure.

In step one we have

$$IY(L+1) \frac{IX(L) \cdots IX(3) \quad IX(2) \quad IX(1)}{IX1} \quad IY(L) \cdots IY(3) \quad IY(2) \quad IY(1)} \quad (2.1)$$

$IY(1)$ is computed by

$$IY(1) = R \left(\frac{IX(1) \times IX1}{1000} \right) ; \quad (2.2)$$

also let

$$ICARRY_1 = I \left(\frac{IX(1) \times IX1}{1000} \right) .$$

$R \left(\frac{a}{b} \right)$ is defined to be the remainder of a/b , and $I \left(\frac{a}{b} \right)$ is defined to be the integral number of times b divides a where a and b are integers.

Then for $2 \leq i \leq L$,

let $IY(i) = R(T_i)$ and $ICARRY_i = I(T_i)$

where

$$T_i = \left(\frac{IX(i) \times IX1 + ICARRY_{i-1}}{1000} \right) . \quad (2.3)$$

To terminate this step, set $IY(L+1) = ICARRY_L$ and increase L by one.

In step two we have

$$IY(L+1) \frac{IX(L-1) \cdots IX(3) \quad IX(2) \quad IX(1)}{IX2} \quad IY(L) \quad IY(L-1) \cdots IY(3) \quad IY(2) \quad IY(1)} \quad (2.4)$$

where $IY(1)$ has the same value as in step one, $L-1$ has the same integer value as L in step one, and $IY(2) = R(T_2)$ and $ICARRY_2 = I(T_2)$ for

$$T_2 = \frac{IX(1) \times IX2 + \overline{IY(2)}}{1000} . \quad (2.5)$$

Let $\overline{IY(i)}$ denote $IY(i)$ as computed in the previous step.

Let $IY(i) = R(T_i)$ and $ICARRY_i = I(T_i)$,

where

$$T_i = \left(\frac{IX(i-1) \times IX2 + \overline{IY(i)} + ICARRY_{i-1}}{1000} \right) \text{ for } 3 \leq i \leq L . \quad (2.6)$$

Step two is also terminated by setting $IY(L+1) = ICARRY_L$ and then increasing L by one.

In step three we have

$$\begin{array}{ccccccc} & & & IX(L-2) & \cdots & IX(3) & IX(2) & IX(1) \\ & & & & & IX3 & & \\ \hline IY(L+1) & IY(L) & IY(L-1) & IY(L-2) & \cdots & IY(3) & IY(2) & IY(1) \end{array} \quad (2.7)$$

where $IY(1)$ and $IY(2)$ remain invariant from step two, $IY(3) = R(T_3)$, and $ICARRY_3 = I(T_3)$ where

$$T_3 = \left(\frac{IX(1) \times IX3 + \overline{IY(3)}}{1000} \right) . \quad (2.8)$$

Let $IY(i) = R(T_i)$ and $ICARRY_i = I(T_i)$

where

$$T_i = \left(\frac{IX(i-2) \times IX3 + \overline{IY(i)} + ICARRY_{i-1}}{1000} \right) \quad (2.9)$$

for $4 \leq i \leq L$.

To terminate step three we set

$$IY(L+1) = \begin{cases} ICARRY_L, & \text{if } ICARRY_L \neq 0 \\ \text{not defined,} & \text{if } ICARRY_L = 0 \end{cases} , \quad (2.10)$$

$$\hat{L} = \begin{cases} L+1, & \text{if } ICARRY_L \neq 0 \\ L, & \text{if } ICARRY_L = 0 \end{cases} , \quad (2.11)$$

and then $L = \hat{L}$.

It can be seen from equations (2.1) through (2.9) that no fixed point number calculated by the above method will be greater than $999^2 + 999 + 999 = 999,999$ which is a six-digit number and can therefore be handled by the Control Data 3200.

The IY array is then placed into the IX array and eight, nine, or ten correct digits (depending on where the theoretical decimal point is located in relation to the three digits of $IX(i)$ which contains the theoretical decimal point) after the theoretical decimal point are placed into the floating point X array. The X array therefore contains numbers between zero and one; this is the desired "white sequence." The process, starting with step one and continuing through the placing of a new number

into the X array, is repeated until L is even and $ISTOP < L \leq 12,000$ where ISTOP is an input parameter.

The entire IX array is placed on tape before ISLASHO terminates; therefore, this calculation can later be continued. However, in order to calculate many more than 4266 points in this "white sequence" by this method, ISLASHO would have to be modified such that part of the IX and IY arrays could be stored elsewhere than in the memory of the computer during the calculation of each point of X because these two arrays would soon outgrow the storage capacity of the computer. Also the longer these arrays become, the time needed to calculate each point of X increases until this method would no longer be practical. For example, v_n , $1 \leq n \leq 500$, was computed in about 45 seconds as compared with thirteen minutes needed to compute the 500 points, v_n , $3501 \leq n \leq 4000$. It was noted that if t minutes were needed to compute a block of 500 points of v_n , approximately $(t + 1.75)$ minutes were needed to compute the next 500 points of v_n . It therefore may become necessary to generate several "white sequences," $v_n^{(1)}, \dots, v_n^{(s)}$, for different transcendental numbers and then interlace these sequences by the method (1.16) to form the desired "white sequence."

The following is the list of input parameters to the program ISLASHO. The numbers in parentheses were the numbers used to find (e^n) for $1 \leq n \leq 4266$.

NO DEC the number of digits to the right of the decimal point of the number which is to be raised to power. (8).

IX3, IX2, IX1 these three parameters each contain three digits of the number which is to be raised to powers. (2.71828183 was used; therefore, IX3 = 271, IX2 = 828, and IX1 = 183).

IF PRNT SQ = 1, the X array is buffered out onto tape and printed in blocks of IBUF points,
 = 0, the X array is only buffered out on tape in blocks of IBUF points. (1).

IF PRNT ME = 1, buffer out onto tape and print the entire IX array just before the program terminates,
 = 0, only buffer out onto tape the entire IX array (0).

ISTOP the program terminates when L is even and $ISTOP < L \leq 12,000$. ISTOP should never exceed six less than the dimension on the IX and IY arrays. (11990).

IBUF the number of points of the "white sequence," X array, that should be handled in one block. IBUF should be an even number and not exceed the dimension of X which is 500 in the listing. Also $IBUF \leq 588$ if the subroutine SOLUTION in GAUSSIAN is not to be changed.

The output tape has the following form:

```

IBUF
X(1) . . . . X(IBUF)
.
.
.
IBUF
X(1) . . . . X(IBUF)
ICOUNT
X(1) . . . . X(ICOUNT)
      end-of-file
I DEC
L
IX(1) . . . . IX(L)
      end-of-file .

```

ICOUNT the number of points in the last block of the X array.
 $ICOUNT \leq IBUF$.

I DEC the number of digits in the decimal part of the last
 number in memory.

L the number of three-digit numbers which compose the last
 number in memory.

IX(1) — IX(L) the last number in memory put on tape as three-digit
 numbers. IX(1) is the right most and IX(L) is the
 the left most three digits of this number.

B. Constructing the w-sequence

The only purpose of CLLCNVRT is to convert a "white sequence" into its w-sequence by (1.17) and place this w-sequence onto tape. This program is almost trivial, but it simplifies the programming of the problem. It can be seen by (1.17) why it is desired that an even number of points of the "white sequence" be placed on tape in each block. If the number of points in a block is odd, the last point in that block will not be used by the subroutine CONVERT in forming the w-sequence.

The program CLLCNVRT has one input parameter, IF PRINT.

IF PRINT = 1, print and buffer out the w-sequence onto tape,
 = 0, only buffer out the w-sequence onto tape.

The output tape has the following form:

```

      I BUF
      W(1) ..... W(I BUF)
      .
      .
      .
      I BUF
      W(1) ... W(I BUF)
      end-of-file.

```

I BUF same as on input tape.
 W(1) ... W(I BUF) a block of w-sequence points.

C. Constructing a Specified Time Series

The program GAUSSIAN shapes the input w-sequence into a time series $x(t)$ with a given power spectral density function and a given sampling time interval. The power spectral density function must be defined by the real numbers a_i , $1 \leq i \leq n+1$, and b_i , $1 \leq i \leq m+1$ for

$$v(\omega) = \left| \frac{P(i\omega)}{Q(i\omega)} \right|^2 = \left| \frac{b_1(i\omega)^m + \dots + b_m(i\omega) + b_{m+1}}{a_1(i\omega)^n + \dots + a_n(i\omega) + a_{n+1}} \right|^2 \quad (2.12)$$

where $m < n$, ω is real, and the zeros of $a_1\xi^n + \dots + a_n\xi + a_{n+1}$ lie in the halfplane $\text{Re}(\xi) < 0$. It should be noted that each of the n -dimensional vectors in (1.25) is used to construct one point of the specified time series; therefore, if the w-sequence consists of J points, the specified time series will contain J/n points.

The numbers a_i , $1 \leq i \leq n+1$, are read into the vector SA; the numbers b_i , $1 \leq i \leq m+1$, are read into the vector SB; and the time interval Δt is read into DT.

Subroutine FIND B places the numbers, a_i , $1 \leq i \leq n+1$, into a constant, $n \times n$ matrix B and constructs the n -dimensional constant vector $[0, \dots, 0, \frac{1}{2}]$ by (1.23).

The subroutine GAUSS P then finds the solution vector $[m_1, \dots, m_n]$ by means of Gaussian elimination with partial pivoting.¹³ Since the determinant of the input matrix can be found by simply computing the product of the elements on the main diagonal and multiplying this product by the proper sign after the input matrix has been reduced to upper triangular form, the determinant of the input matrix is always calculated by GAUSS P and printed out by the program GAUSSIAN as a guide to the reliability of the solution vector.

The subroutine FIND CM places m_i , $1 \leq i \leq n$, into the positive definite moment matrix M whose elements, m_{ij} , $1 \leq i, j \leq n$, have been found by (1.24). Then as stated in Chapter I, Section B, a lower triangular matrix T must be found such that

$$M = T T^* \quad (2.13)$$

Since M is a positive definite, real, symmetric matrix, the desired lower triangular matrix T can be found by Crout factorization. Let t_{ij} be the elements of T and m_{ij} be the elements of M . It is known that $t_{ij} = 0$ for $j > i$; therefore, only t_{ij} for $j \leq i$ need to be computed. Then by (2.13)

$$m_{ij} = \sum_{k=1}^j t_{ik} t_{jk} \quad (2.14)$$

can be obtained. It is then suggested that t_{ij} , for $j \leq i$, be computed in the following order:

$$ij = 11, 21, \dots, n1; 22, 32, \dots, n2; \dots; nn.$$

By this order,

$$t_{11} = m_{11}^{\frac{1}{2}} \quad (2.15)$$

is computed first, and the other elements in the first column are given by

$$t_{i1} = t_{11}^{-1} m_{i1} \text{ for } i = 2, \dots, n \quad (2.16)$$

If the preceding columns $k < j$ have been computed, the j^{th} diagonal element can be computed by

$$t_{ij} = \left(m_{ij} - \sum_{k=1}^{j-1} t_{jk}^2 \right)^{\frac{1}{2}} \quad (2.17)$$

Now the elements below the diagonal can be computed by

$$t_{ij} = t_{jj}^{-1} m_{ij} - \sum_{k=1}^{j-1} t_{ik} t_{jk} \quad (2.18)$$

for $i = j+1, \dots, n$ where $j < n$.¹⁴ This method of Crout factorization is carried out in the subroutine T T STAR.

The matrix $\exp(\Delta t A)$, where A has the form as shown in (1.21) and is constructed by the subroutine MAKE A, is computed by the subroutine EXP ADT. This exponential matrix is defined by

$$\exp(\Delta t A) = I + \frac{\Delta t}{1!} A + \frac{\Delta t^2}{2!} A^2 + \dots + \frac{\Delta t^k}{k!} A^k + \dots \quad (2.19)$$

which is the same as

$$\begin{aligned} \exp(\Delta t A) = & I + \left(\frac{\Delta t}{1}\right) A \left[I \right] + \left(\frac{\Delta t}{2}\right) A \left[\frac{\Delta t}{1!} A \right] \\ & + \left(\frac{\Delta t}{3}\right) A \left[\frac{\Delta t^2}{2!} A^2 \right] + \left(\frac{\Delta t}{4}\right) A \left[\frac{\Delta t^3}{3!} A^3 \right] + \dots \quad (2.20) \\ & \dots + \left(\frac{\Delta t}{k}\right) A \left[\frac{\Delta t^{k-1}}{(k-1)!} A^{k-1} \right] + \dots \end{aligned}$$

The method of summation given by (2.20), where the $(k-1)^{\text{th}}$ term is multiplied by $(\Delta t/k) A$ and then added to the previous sum, is used to calculate $\exp(\Delta t A)$ in EXP A DT.

The summation terminates when the condition that $b_{ij} < \epsilon_{chk} = \max \{ |a_{ni}| \times 10^{-12} \mid 1 \leq i \leq n \text{ and } a_{ni} \text{ is an element of } A \}$ holds for every b_{ij} , $1 \leq i, j \leq n$, where b_{ij} are the elements of $B = \frac{\Delta t^k}{k} A^k$ for $k \geq 1$. The maximum of the above set was taken instead of the minimum so that the computer would not set ϵ_{chk} equal to some very small floating point numbers which should actually be zero. However, if all a_{ni} are approximately the same in magnitude, no problem should arise. If a_{ni} vary greatly in magnitude, it may become necessary to define a floating point zero and change the subroutine to check for the minimum value of $\{ |a_{ni}| \mid 1 \leq i \leq n \}$.

Let \hat{a}_{ij} be the elements of \hat{A} where

$$\hat{A} = A \times B = A \left(\frac{\Delta t^k}{k!} A^k \right) = \left(\frac{\Delta t^k}{k!} A A^k \right) = \left(\frac{\Delta t^k}{k!} A^k \right) A = B \times A \quad (2.21)$$

Now since $A \times B = B \times A$, the matrix multiplication for (2.21) can be defined by

$$\begin{aligned} \hat{a}_{ij} &= b_{i+1, j}, \quad 1 \leq i \leq n-1, \quad 1 \leq j \leq n \\ \hat{a}_{n1} &= b_{nn} a_{n1}, \end{aligned} \quad (2.22)$$

and $\hat{a}_{n, j+1} = b_{nj} + b_{nn} a_{n, j+1}, \quad 1 \leq j \leq n-1$.

This is the method of matrix multiplication used in EXPADT.

Now subroutine BIGSET constructs the $\frac{1}{2} n(n+1)$ equations in the same number of unknowns by (1.28) where the \hat{b}_{ij} for $1 \leq j \leq i \leq n$ are defined by (1.27). The unknowns μ_{ij} , $1 \leq j \leq i \leq n$, can be written in the form of an $n \times n$ lower triangular matrix L , and a one-to-one mapping f can be defined which maps the elements on and below the main diagonal of L onto an \hat{n} -dimensional vector v where $\hat{n} = \frac{1}{2} n(n+1)$. This vector is then the vector of unknowns associated with the $\hat{n} \times \hat{n}$ constant matrix defined by (1.28). After the value of v has been found by GAUSS P, an inverse mapping can be defined to take the elements of v back into L . L can then be made into the desired real symmetric matrix M_r by replacing the zeros above the main diagonal with the corresponding elements below the diagonal.

Let $f:L \rightarrow v$ be the mapping

$$f: \begin{pmatrix} \mu_{11} & 0 & \dots & 0 \\ \mu_{21} & \mu_{22} & \dots & \cdot \\ \cdot & \cdot & \dots & \cdot \\ \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & 0 \\ \mu_{n1} & \mu_{n2} & \dots & \mu_{nn} \end{pmatrix} \longrightarrow \quad (2.23)$$

$$= (\mu_{11}, \mu_{21}, \dots, \mu_{n1}; \mu_{22}, \dots, \mu_{n2}; \mu_{33}, \dots; \mu_{nn})$$

$$= (v_1, v_2, \dots, v_{\hat{n}})$$

When μ_{ij} is under consideration, all n elements in the first $j-1$ columns plus the upper $i-1$ elements in the j^{th} column less the upper triangular set of zero elements have been mapped into v ; therefore, it can intuitively be seen that for μ_{ij} , $1 \leq j \leq i \leq n$, and v_k , $1 \leq k \leq \hat{n}$, the mapping f can be defined by

$$k = i + n(j-1) - [j^2 - \frac{1}{2}j(j+1)] \quad (2.24)$$

which reduces to

$$k = i + (j-1)(2n-j)/2 \quad (2.25)$$

It must be shown that f , as defined by (2.24), is the mapping indicated by (2.23). Let $f(\mu_{ij}) = (v_k)$ be denoted by $f(i, j) = k$ for $1 \leq j \leq i \leq n$ and $1 \leq k \leq \hat{n}$. Therefore show that if $1 \leq j = \bar{j} \leq n-1$ and $\bar{i} - i = 1$, or if $i = n$, $1 \leq j \leq n-1$, and $\bar{i} = \bar{j} = j+1$, then $\bar{k} - k = f(\bar{i}, \bar{j}) - f(i, j) = 1$.

Case I:

Let $1 \leq j = \bar{j} \leq n-1$ and $\bar{i} - i = 1$; this implies that

$$\begin{aligned} \bar{k} - k &= f(\bar{i}, \bar{j}) - f(i, j) \\ &= \bar{i} + (\bar{j}-1)(2n-\bar{j})/2 - i - (j-1)(2n-j)/2 \\ &= \bar{i} - i = 1 \end{aligned}$$

Case II:

Let $i=n$, $1 \leq j \leq n-1$, and $\bar{i} = \bar{j} = j+1$; this implies that

$$\begin{aligned} \bar{k} - k &= f(\bar{i}, \bar{j}) - f(i, j) \\ &= \bar{i} + (\bar{j}-1)(2n-\bar{j})/2 - i - (j-1)(2n-j)/2 \\ &= j+1 + (j+1-1)(2n-j-1)/2 \\ &\quad - n - (j-1)(2n-j)/2 \\ &= j+1 + (2nj - j^2 - j)/2 \\ &\quad - n - (2nj - j^2 - 2n + j)/2 \\ &= 1 \end{aligned}$$

Therefore f is the mapping indicated by (2.23). This mapping is defined in the subroutine INDEX which is used by BIGSET. It is obvious that f is one-to-one, and therefore f^{-1} exists. f^{-1} is defined by taking the elements of the solution vector v in order from left to right and placing them into the lower triangular matrix L in column order.

TTSTAR is then called to find the lower triangular matrix T_r such that $M_r = T_r T_r^*$.

Now that T_r , $\exp(\Delta t A)$, and T have been computed and the vectors $w^{(k)}$ for $k = 0, 1, 2, \dots$ can be constructed from the w -sequence, $z(0)$ can be computed by (1.26) and the sequence of vectors $z(k\Delta t)$ for $k = 1, 2, \dots$ can be computed by (1.29). The input vector SB is also known; therefore, the specified time series, $x(t)$ for $t = 0, \Delta t, 2\Delta t, \dots$, can be computed by (1.30). This process is carried out in the subroutine SOLUTION. This completes the construction of the specified time series.

The card input to GAUSSIAN is of two types, problem data and flags. The sole purpose of the flags is to control the flow of the program. The flag NOSKIP is read once and only once each time the program is compiled; all other input parameters must be read anew for each time series that is to be generated.

The following is the list of problem data:

N	n in equation (2.12). N is also used as a flag in that GAUSSIAN terminates if $N = 0$; this is to be used for normal exit.
M	m in equation (2.12).
SA(1), ..., SA(N+1)	a_1, \dots, a_{n+1} in equation (2.12).
SB(1), ..., SB(M+1)	b_1, \dots, b_{m+1} in equation (2.12).
DT	Δt , time interval at which samples of the time series are to be calculated.

The following is the list of flags:

NOSKIP	the number of end-of-files (ie., previously computed time series) which must be skipped on the output tape before the first time series is computed and then buffered out onto this output tape.
ISTOP	the number of points of the time series which are to be buffered out onto the output tape in each block (except for the last block which contains the remaining points). It is necessary that $ISTOP \leq 5604$.
IPOWER T	used when factoring the matrix M.
IPOWER TR	used when factoring the matrix M_r . IPOWER T and IPOWER TR are called IPOWER in the subroutine TTSTAR. The check, "is $ m_{ii} < 10^{-IPOWER} ?$," is made before m_{ii} is used as a divisor in TTSTAR where m_{ii} are diagonal elements of M or M_r .
IPOWERB	used when solving the set of equations (1.23).
IPOWERC	used when solving the set of equations (1.28). IPOWERB and IPOWERC are called IPOWER in the subroutine GAUSS P. The check, "is $ a_{ii} < 10^{-POWER} ?$," is made before a_{ii} is used as a divisor in GAUSSP where a_{ii} are the diagonal elements of the constant matrix under consideration.
IFT	associated with IPOWER T.
IFTR	associated with IPOWER TR.
IFB	associated with IPOWERB.

IFC associated with IPOWERC. Let IF denote any of the above four flags. If IF = 1 and a zero divisor has been detected by its IPOWER, a diagnostic will be given, but the calculation of this time series will continue. However, if IF = 0 and a zero divisor has been detected, a diagnostic will be given, but the program will go to the next time series.

IF N PRINT = 1 print the number of points which are in the following block of time series points.

 = 0, do not print the number of points.

IF X PRINT = 1 print the computed time series in blocks of ISTOP points.

 = 0, do not print the time series points.

The data deck to GAUSSIAN must be stacked as follows:

NO SKIP (Format I4)

N, M (Format 2I4)

SA(1), ..., SA(N+1) (Format 4F20.6)

SB(1), ..., SB(M+1) (Format 4F20.6)

DT (Format F20.6)

ISTOP, IPOWER, IPOWERTR, IPOWERB, IPOWERC

(Format 5I4)

IFNPRINT, IFXPRINT, IFT, IFTR, IFB, IFC

(Format 6I1)

Repeated
for each
time series
to be
generated.

(Blank card; flags normal exit.)

The problem data are printed and labelled. The intermediate matrices and vectors, which are found by the various subroutines are printed

unlabelled as they are calculated.

More than one time series can be shaped from the same w-sequence each time GAUSSIAN is compiled. This input tape which contains the w-sequence must be of the same format as defined for the output tape from CLLCNVRT. Each block of the w-sequence should not exceed 588 points. GAUSSIAN terminates when and only when either $N = 0$ or a parity error has been detected on either the input or output tape.

The output tape has the following form:

ISTOP

X(1) X(ISTOP)

.

ISTOP

X(1) X(ISTOP)

IPRINT

X(1) X(IPRINT)

end-of-file

ISTOP

X(1) X(ISTOP)

.

IPRINT

X(1) X(IPRINT)

end-of-file.

X

array of time series points.

IPRINT

number time series points in the last block.

III. "WHITE SEQUENCE," w-SEQUENCE, AND EXAMPLE TIME SERIES

A spectral analysis program similar to the program SPECT, as found in a report by Ellis and Boston,¹⁵ is used to find the autocorrelation function and power spectral density function of the "white sequence" generated by ISLASHO and four example time series generated by GAUSSIAN. The autocorrelation function is defined by (1.3) and the power spectral density function is defined by (1.4). In SPECT the time τ in (1.3) and (1.4) is taken at regular intervals $\Delta\tau$, and therefore

$$\tau = l\Delta\tau, l=0,1,2, \dots, \quad (3.1)$$

where l is called the lag number. For practical applications on the computer, τ can only take on finite values and the sequences considered will only have a finite number of points; therefore a finite number of lags will be taken in SPECT. The maximum lag number will generally be about ten percent of the number of points in the given sequence.

A. "White Sequence"

The condition (1.14) states that the autocorrelation function of a "white sequence" must be zero for all but the zeroth lag. The "white sequence" under consideration is $\{e^n\}$, $1 \leq n \leq 4266$. It is also stated in Chapter I, Section B, that the power spectral density function of v_n must be constant. The plots of the autocorrelation function and the power spectral density function of v_n as computed by SPECT for time interval $\Delta\tau = .1$ sec. and $l = 200$ lags are shown in Figures 1 and 2, respectively. It can be seen from these two plots that the sequence $\{e^n\}$ is a good numerical approximation to a "white sequence."

B. w-Sequence

A short program, which is not listed, was written to find the mean and the variance of the w-sequence. Let the mean be $m = E(w_i)$ and the variance be defined as

$$\sigma^2 = \frac{1}{N} \sum_{i=1}^N (w_i - m)^2, \quad (3.2)$$

where w_i , $1 \leq i \leq N$, are the N points of the w-sequence. It is stated in Chapter I that the w-sequence, as constructed by (1.17), should have mean 0 and variance 1. The numerical w-sequence constructed from $\{e^n\}$ by CLLCNVRT was found to have $m \approx -.029$ and $\sigma^2 \approx .9577$.

C. Four Example Time Series

Four example problems are now to be considered. The following three steps are taken in the consideration of these examples:

- (1) Given the constant coefficients of the polynomials $P(i\omega)$ and $Q(i\omega)$ of the theoretical power spectral density function $V(\omega)$ as given by (2.12) and a time interval $\Delta\tau$, GAUSSIAN is used to shape a time series with this power spectral density function.
- (2) A log-log plot of the theoretical power spectral density function as a function of frequency f , is made where $f = \frac{\omega}{2\pi}$ and ω is the angular velocity given in radians/sec.
- (3) The actual autocorrelation function and power spectral density function are then calculated by SPECT; the autocorrelation function is plotted linearly against time while the power spectral density function is plotted against frequency on a log-log plot.

Only the power spectral density function is considered in the first three examples; however, both the autocorrelation and the power spectral

density functions for example four are given by Richie¹⁶ and will therefore be considered.

In all cases the theoretical and actual log-log plots of the power spectral density functions were found to have approximately the same shape; however, for each example, the two curves were found to differ by some arbitrary constant. If the magnitude of the points of the power spectral density function is specified and not just the shape, the specified time series can be obtained by multiplying each point of the original time series by the square root of the constant whose logarithm is the difference between the theoretical and the actual curves.

It can be seen by (1.3) that if each point of the time series $x(t)$ is multiplied by a scale factor s , then for every τ , the autocorrelation function is multiplied by s^2 . Since the constant s^2 can be removed from the integral in (1.4), the power spectral density function is also multiplied by s^2 .

Let t_i and a_i be corresponding points on the theoretical and actual power spectral density curves, respectively, and if there exists an $s > 0$ such that for every i where t_i and a_i are defined, $\log s^2 = \log t_i - \log a_i$, then $t_i = s^2 a_i$. The converse is also true. Therefore the specified magnitude of the power spectral density function can be obtained by multiplying each point of the time series by the scale factor s . This scale factor may have physical units.

In Figures 3 through 6 the crooked continuous curve is the actual power spectral density function as computed by SPECT and the smooth dotted curve is a multiple of the theoretical curve which gives an approximate fit to the points of the actual power spectral density function.

The theoretical power spectral density function of example one is given to be

$$V_1(\omega) = \frac{9\omega^2 + 1}{\omega^4 - 6\omega^2 + 25}, \quad (3.3)$$

which has zeros at the points $(0, 1/3)$ and $(0, -1/3)$ and poles at the points $r, \bar{r}, -r,$ and $-\bar{r}$ where $r = (2, 1)$. Therefore for all real ω , the condition $0 < V_1(\omega) < \infty$ holds. Also for real ω , the conditions $V_1(\omega) = V_1(-\omega)$ and $V_1(\omega) \rightarrow \pm \infty$ hold; therefore the three conditions, as stated in (1.8), exist which insure that $V_1(\omega)$ can be represented in the form of (2.12). Specifically, $V_1(\omega)$ can be represented as

$$V_1(\omega) = \left| \frac{3(i\omega) + 1}{(i\omega)^2 + 2(i\omega) + 5} \right|^2 \quad (3.4)$$

where $\xi^2 + 2\xi + 5$ has zeros at $(-1, 2)$ and $(-1, -2)$. Both of these zeros lie in the halfplane $\text{Re } \xi < 0$.

The theoretical power spectral density function of example two is given to be

$$V_2(\omega) = \frac{\omega^4 - 70\omega^2 + 1369}{\omega^6 + 14\omega^4 + 49\omega^2 + 36}. \quad (3.5)$$

The numerator of $V_2(\omega)$ has zeros at the points $r, \bar{r}, -r,$ and $-\bar{r}$ where $r = (6, 1)$, and the denominator of $V_2(\omega)$ is greater than zero for all real ω ; therefore, for all real ω , the condition $0 < V_2(\omega) < \infty$ holds. Also the other two conditions of (1.8) hold; therefore $V_2(\omega)$ can be represented, as indicated in (2.12), by

$$V_2(\omega) = \left| \frac{(i\omega)^2 + 2(i\omega) + 37}{(i\omega)^3 + 6(i\omega)^2 + 11(i\omega) + 6} \right|^2 \quad (3.6)$$

where $\xi^3 + 6\xi^2 + 11\xi + 6$ has zeros at $(-1, 0), (-2, 0),$ and $(-3, 0)$. All these zeros lie in the halfplane $\text{Re } \xi < 0$.

The theoretical power spectral density function of example three

is given to be

$$V_3(\omega) = \frac{400\omega^2 + 1}{\omega^6 + 14\omega^4 + 49\omega^2 + 36} \quad (3.7)$$

$V_3(\omega)$ has zeros at the points $(0, \frac{1}{20})$ and $(0, -\frac{1}{20})$. The denominator of this example is the same as in example two. It can be seen, as in examples one and two, that $V_3(\omega)$ can be represented by

$$V_3(\omega) = \left| \frac{20(i\omega) + 1}{(i\omega)^3 + 6(i\omega)^2 + 11(i\omega) + 6} \right|^2 \quad (3.8)$$

Example four is an example problem which appears as channel 4 in a report by Richie.¹⁷ The theoretical power spectral density function is given to be

$$V_4(\omega) = 4Ak \frac{\omega^2 + (k^2 + c^2)}{\omega^4 + 2(k^2 - c^2)\omega^2 + (k^2 + c^2)^2} \quad (3.9)$$

where $A = 2149.2$, $k = .658$, and $c = 2.71$. Since A , c , and k are positive and $k > c$, the conditions (1.8) hold for real ω . Therefore, $V_4(\omega)$ can be represented as

$$V_4(\omega) = \left| \frac{b_1(i\omega) + b_2}{a_1(i\omega)^2 + a_2(i\omega) + a_3} \right|^2 \quad (3.10)$$

where $b_1 \approx 75.210999197$, $b_2 \approx 53.521755647$, $a_1 = 1.0$, $a_2 = 1.316$, and $a_3 = .506405$.

Both zeros of $a_1\xi^2 + a_2\xi + a_3$ lie in the halfplane $\text{Re}\xi < 0$ since $a_1 > 0$, $a_2 > 0$, and $a_3 > 0$. This fact can be shown by letting $x = \text{Re}\xi$ and $y = \text{Im}\xi$ and then assuming that there exists a ξ with $x \geq 0$ such that $a_1\xi^2 + a_2\xi + a_3 = 0$. This complex equation is equivalent to the following two real equations:

$$a_1(x^2 - y^2) + a_2x + a_3 = 0 \quad (3.11)$$

and

$$2a_1xy + a_2y = 0 \quad (3.12)$$

By (3.12), $y = 0$ since $x \geq 0$. Then by (3.11), $a_1(x^2 - y^2) + a_2x + a_3 \neq 0$ which is a contradiction; therefore both zeros of $a_1\xi^2 + a_2\xi + a_3$ lie in the halfplane $\text{Re}\xi < 0$.

The actual power spectral density function of the time series generated by GAUSSIAN for this example was found to differ from the theoretical power spectral density function only by a multiplicative constant; however, the theoretical autocorrelation function given by Richie¹⁸ and the actual autocorrelation function differed greatly. This illustrates the fact, as pointed out previously, that small perturbations in $V(\omega)$ may result in great changes of $R(\tau)$. These two autocorrelation functions are shown in Figure 7, and the power spectral density functions are shown in Figure 6.

The degree and coefficients of the polynomials P and Q , which are used as input parameters to GAUSSIAN, can be taken from equations (3.4), (3.6), (3.8), and (3.10). These parameters, along with the time intervals used in generating a time series for each of these examples, are also given in Table 1.



FIGURE 1
WHITE NOISE

DRL - UT
DWG AS-66-688
WM - ORS
7 - 11 - 66

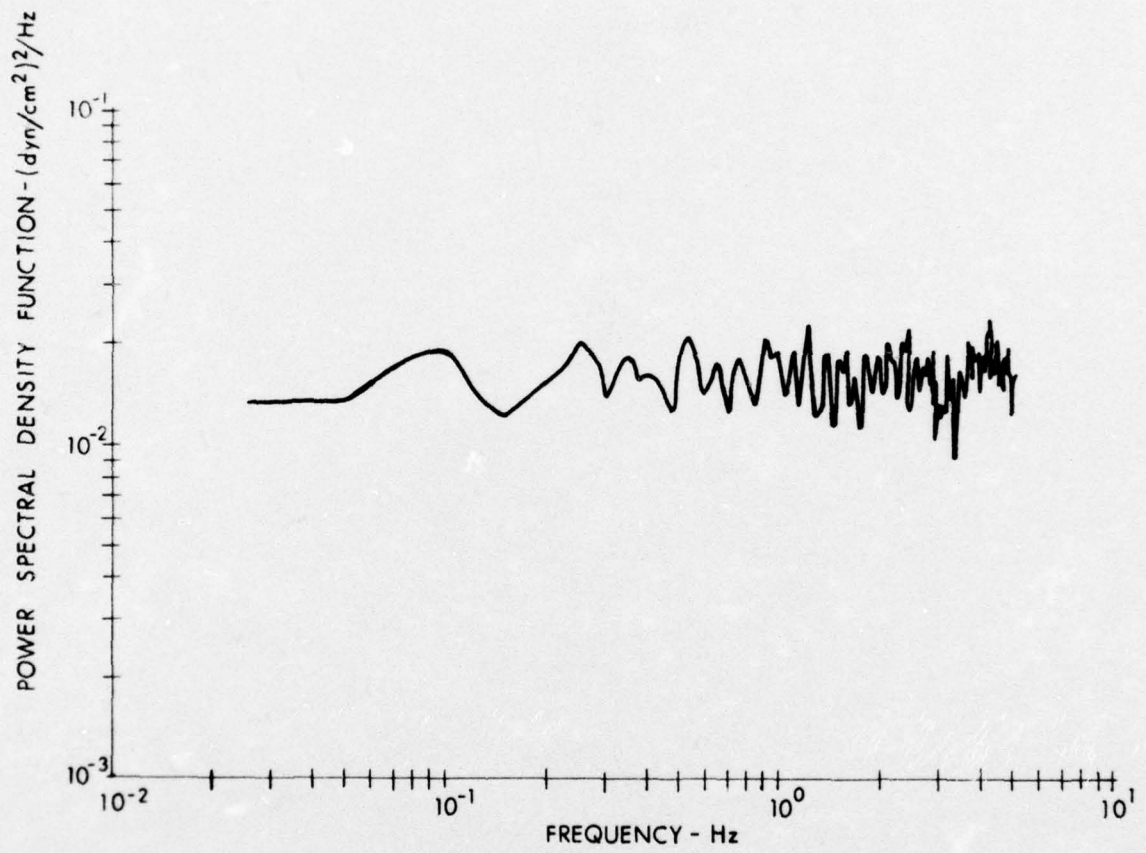


FIGURE 2
WHITE NOISE

DRL - UT
DWG AS-66-689
WM - ORS
7 - 11 - 66

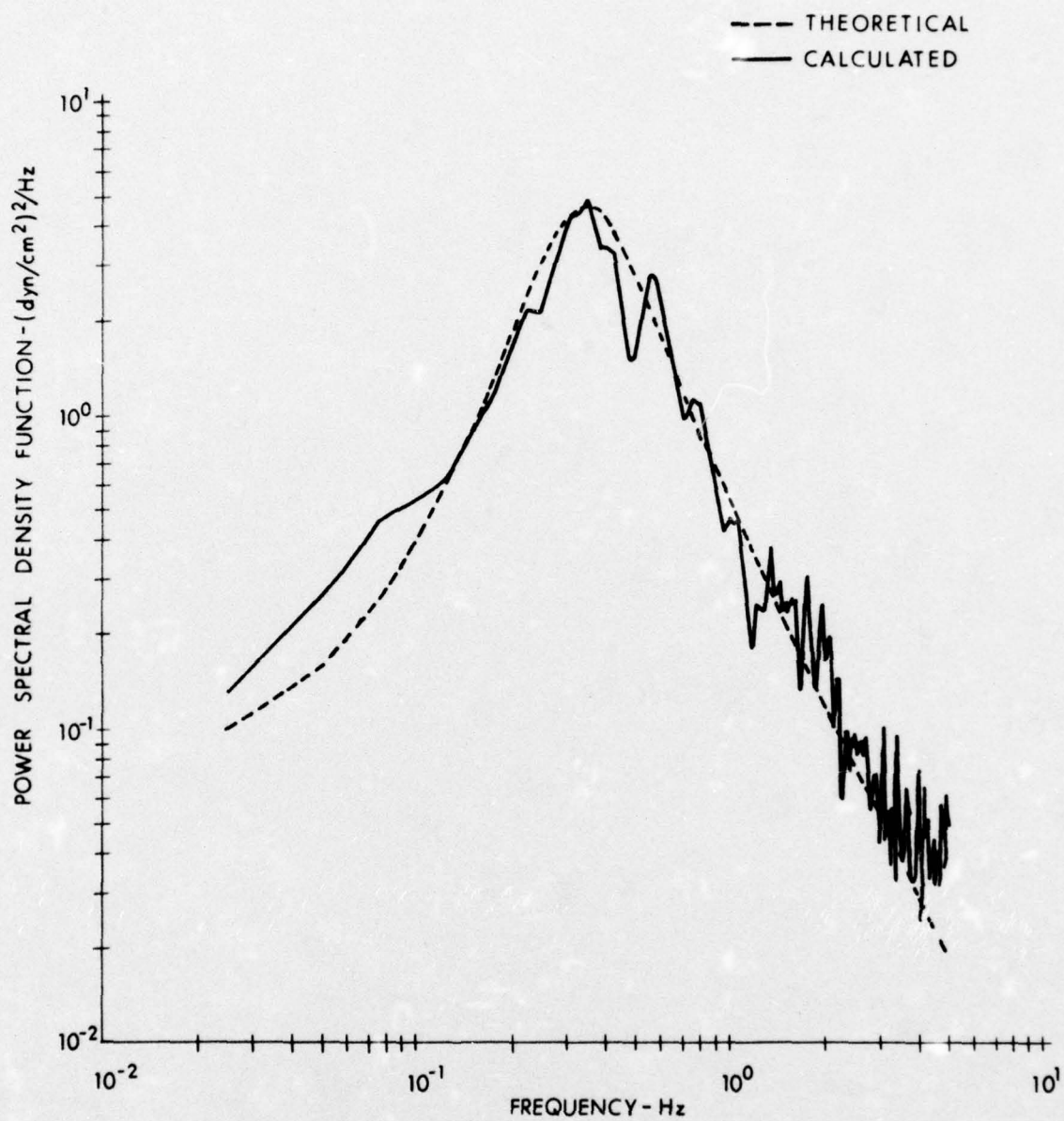


FIGURE 3
EXAMPLE 1

$$V_1(\omega) = \frac{9\omega^2 + 1}{\omega^4 - 6\omega^2 + 25}$$

DRL - UT
 DWG AS - 66 - 690
 WM - ORS
 7 - 12 - 66

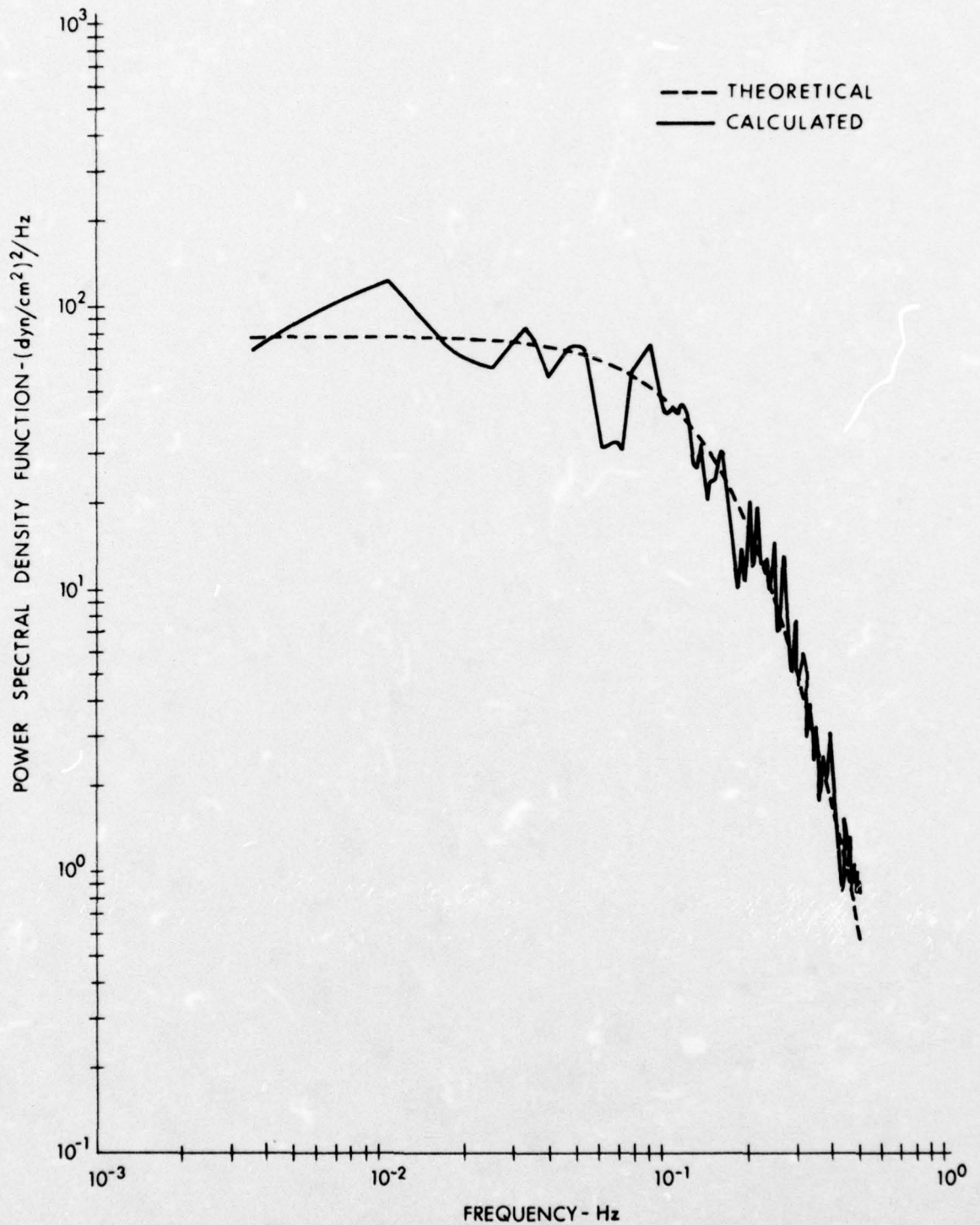


FIGURE 4
EXAMPLE 2

$$V_2(\omega) = \frac{\omega^4 - 70\omega^2 + 1369}{\omega^6 + 14\omega^4 + 49\omega^2 + 36}$$

DRL - UT
DWG AS-66-691
WM - ORS
7 - 12 - 66

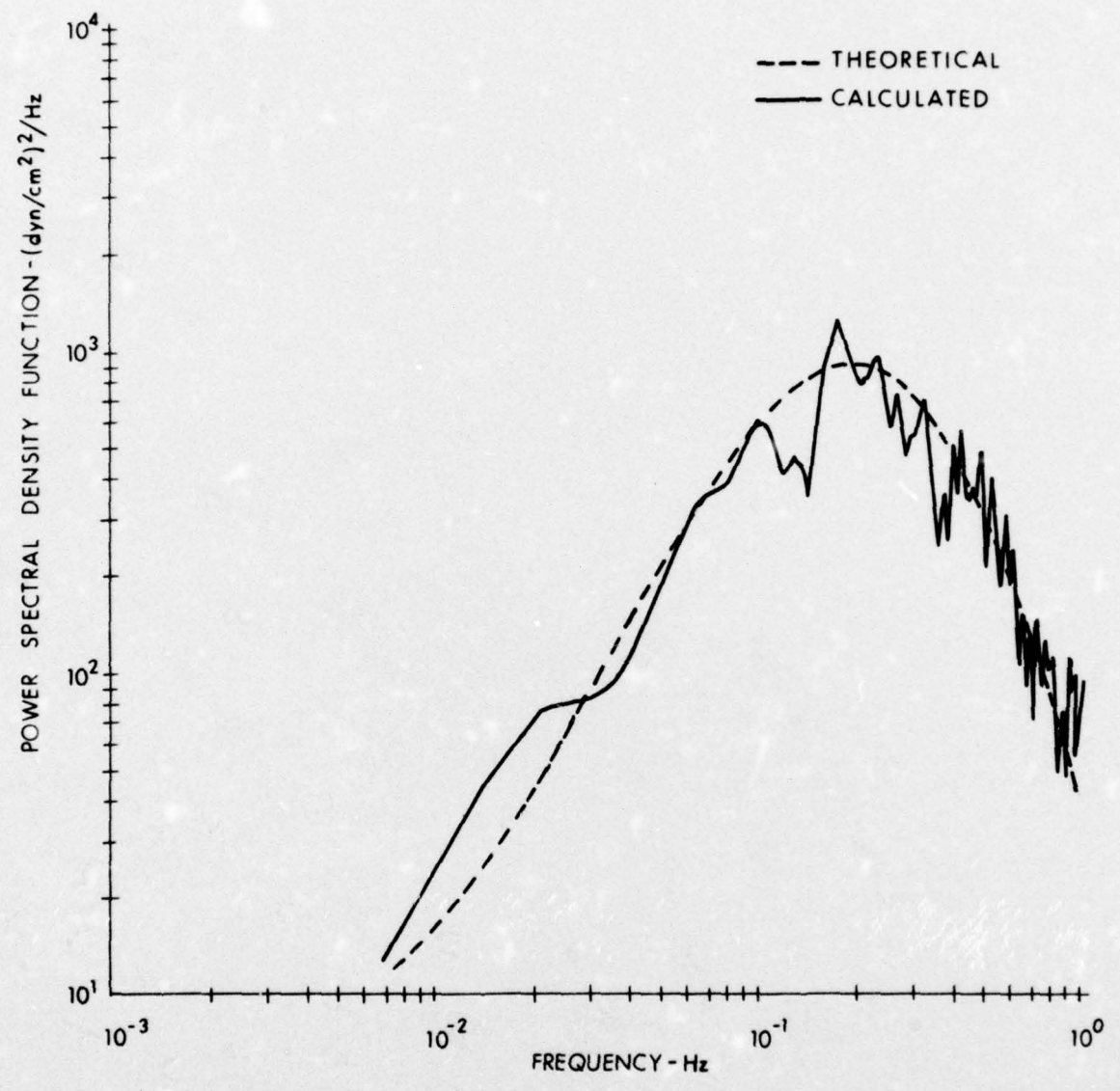


FIGURE 5
EXAMPLE 3

$$V_3(\omega) = \frac{400\omega^2 + 1}{\omega^6 + 14\omega + 49\omega^2 + 36}$$

DRL - UT
DWG AS - 66-692
WM - ORS
7 - 13 - 66

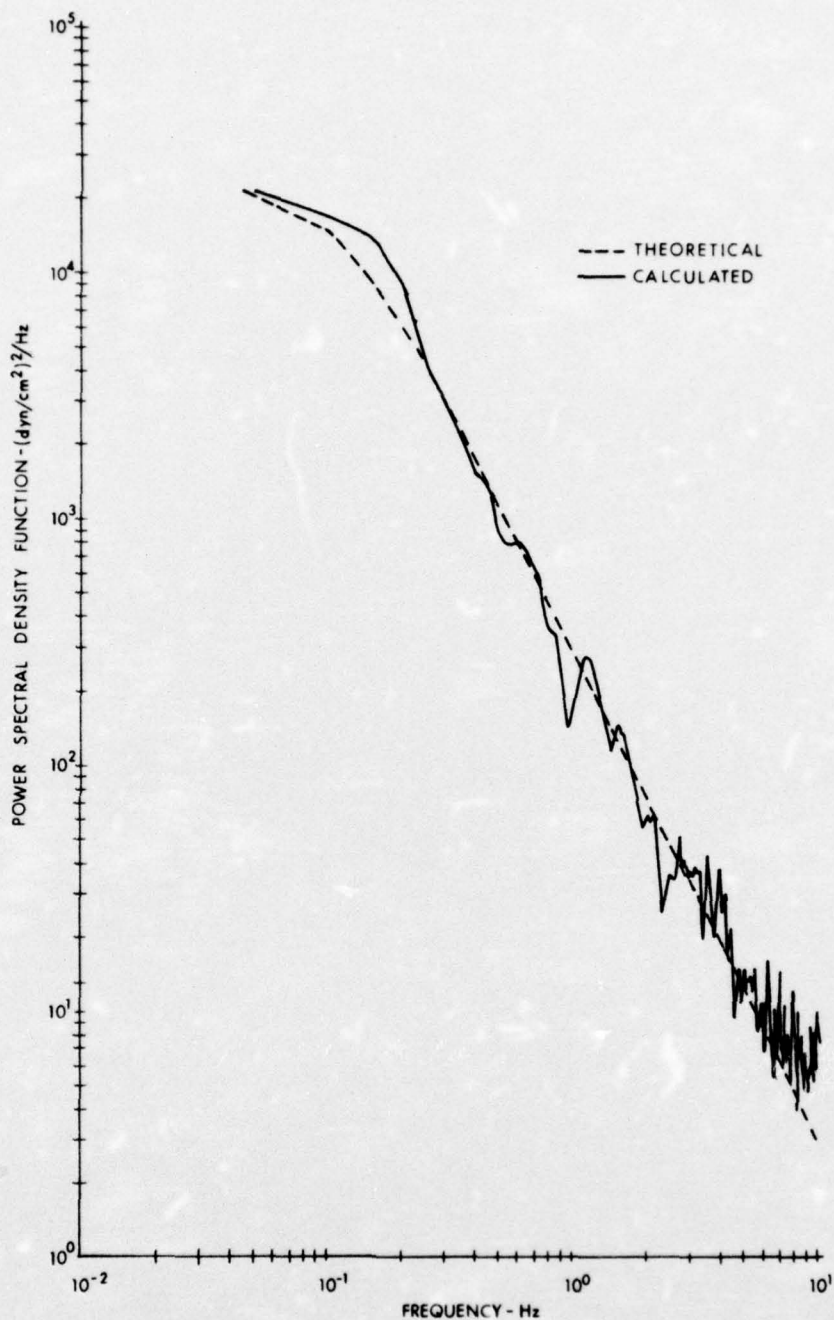


FIGURE 6
EXAMPLE 4
RICHIE'S SEQUENCE
CHANNEL 4

DRL - UT
DWG 85-66-693
WM - ORS
7 - 13 - 66

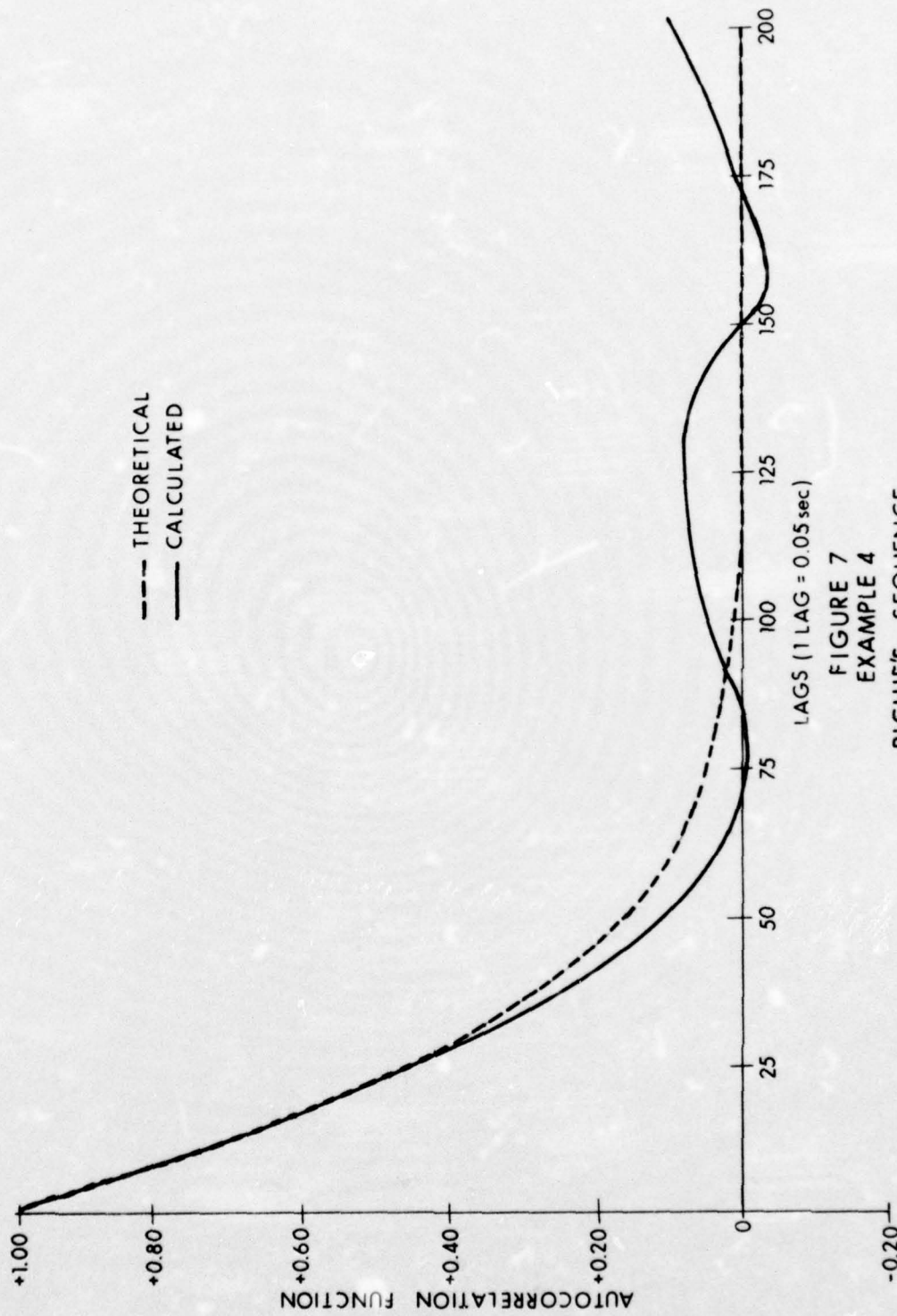


FIGURE 7
EXAMPLE 4
RICHIE'S SEQUENCE
CHANNEL 4

DRL - UT
DWG AS-66-694
WM - ORS
7 - 13 - 66

Table 1

Example	degree of P m	degree of Q n	coefficients of P	coefficients of Q	time interval $\Delta\tau$ (sec)	number of points in time series	number of lags in SPECT
Example 1	1	2	$b_1 = 3$	$a_1 = 1$.1	2133	200
			$b_2 = 1$	$a_2 = 2$			
				$a_3 = 5$			
Example 2	2	3	$b_1 = 1$	$a_1 = 1$	1.0	1420	140
			$b_2 = 2$	$a_2 = 6$			
			$b_3 = 37$	$a_3 = 11$			
			$a_4 = 6$				
Example 3	1	3	$b_1 = 20$	$a_1 = 1$.5	1420	140
			$b_2 = 1$	$a_2 = 6$			
				$a_3 = 11$			
				$a_4 = 6$			
Example 4	1	2	$b_1 = 75.210999197$	$a_1 = 1.0$.05	2133	200
			$b_2 = 53.521755647$	$a_2 = 1.316$			
				$a_3 = .506405$			

3200 FORTRAN (2.1)

02/03/66

```

PROGRAM I SLASH 0
DIMENSION X(500)
COMMON/DATA/NO DEC, I DEC, IX1, IX2, IX3, L, ITHOUS, ICARRY
COMMON IX(12000), IY(12000)
ITHOUS = 1000
REWIND 5
1 FORMAT (16I5)
2 FORMAT (//)
3 FORMAT (1H1)
98 FORMAT (2I1)
16 FORMAT (6F20.11)
17 FORMAT (2I10)
TEN1=1.0E-1
TEN2=1.0E-2
TEN3=1.0E-3
TEN4=1.0E-4
TEN5=1.0E-5
TEN6=1.0E-6
TEN7=1.0E-7
TEN8=1.0E-8
TEN9=1.0E-9
TEN10=1.0E-10
TEN11=1.0E-11
TEN12=1.0E-12
PRINT 3
READ 1, NO DEC
READ 1, IX3, IX2, IX1
READ 98, IF PRNT SQ, IF PRNT ME
READ 17, ISTOP, IRUF
IX(1) = IX1
IX(2) = IX2
IX(3) = IX3
I DEC = NO DEC
L = 3
I COUNT = 1
I = I DEC/3
IR = I DEC - I*3
IF (IR) 5, 4
4 X(1) = IX2*TEN3 + IX1*TEN6
GO TO 6
5 XX=IX3*10.0**(-IR) - IX2*10.0**(-3-IR) + IX1*10.0**(-6-IR)
I=XX
X(1)=XX-I
6 I COUNT = I COUNT + 1
7 CALL RANDOM
I = I DEC/3
IR = I DEC - I*3
IF (IR-1) 8, 9, 10
8 X(I COUNT) = IX(1)*TEN3 + IX(I-1)*TEN6 + IX(I-2)*TEN9 + IX(I-3)*TEN12
GO TO 12
9 XX=IX(I+1)*TEN1 + IX(I)*TEN4 + IX(I-1)*TEN7 + IX(I-2)*TEN10
GO TO 11
10 XX= IX(I+1)*TEN2 + IX(I)*TEN5 + IX(I-1)*TEN8 + IX(I-2)*TEN11
11 I = XX
X(I COUNT) = XX-I
12 IF (L .GT. ISTOP) 60, 14
60 JJ = I COUNT/2
JJ = JJ*2
IF (I COUNT .EQ. JJ) 13, 6
14 IF (I COUNT - IRUF) 6, 15

```

```
15 PRINT 1, I BUF
   IF (IF PRNT SQ) 30, 31
30 PRINT 16, (X(1), I=1, I BUF)
   PRINT 2
31 BUFFER OUT (5,1) (I BUF, I BUF)
32 GO TO (32, 33, 33, 77), UNITSTF(5)
33 BUFFER OUT(5,1)(X(1), X(I BUF))
34 GO TO (34, 35, 35, 77), UNITSTF(5)
35 I COUNT = 1
   GO TO 7
13 PRINT 1, I COUNT
   IF (IF PRNT SQ) 37, 38
37 PRINT 16, (X(1), I=1, I COUNT)
   PRINT 2
38 BUFFER OUT(5,1)(I COUNT, I COUNT)
39 GO TO (39, 40, 40, 77), UNITSTF(5)
40 BUFFER OUT(5,1)(X(1), X(I COUNT))
41 GO TO (41, 42, 42, 77), UNITSTF(5)
42 END FILE 5
   IY(1)=IX1
   IY(2)=IX2
   IY(3)=IX3
   PRINT 1, NO DEC
   PRINT 1, (IY(1), I=1, 3)
   PRINT 2
   PRINT 17, I DEC, L
   IF (IF PRNT ME) 43, 44
43 PRINT 1, (IX(1), I=1, L)
   PRINT 2
44 BUFFER OUT (5,1) (I DEC, I DEC)
45 GO TO (45, 46, 46, 77), UNITSTF(5)
46 BUFFER OUT( 5,1)(L,L)
47 GO TO (47, 48, 48, 77), UNITSTF(5)
48 BUFFER OUT(5,1)(IX(1), IX(L))
49 GO TO (49, 50, 50, 77), UNITSTF(5)
50 END FILE 5
   GO TO 99
77 PRINT 78
78 FORMAT (25H PARITY ERROR ON MAG TAPE)
99 CONTINUE
   END
```

3200 FORTRAN DIAGNOSTIC RESULTS - FOR ISLASHD

NO ERRORS
COMPASS.L.X

COMPASS-32

(2.1)

RANDOM

					ENTRY	RANDOM	02/03/66	PAGE	2
00000	01077777	01	0	77777	0	RANDOM	UJP	**	
00001	14600000	14	1	00000	2		ENA	0	
00002	40000007	40	0	000007	0		STA	CARRY	
00003	14177776	14	0	77776	1		ENT	-1.1	
00004	20000005	20	0	000005	0		LDA	L	
00005	15477776	15	1	77776	0		INA.S	-1	
00006	44000007	44	0	000007	0		SWA	**1	
00007	10177777	10	0	77777	1	INDEX1	IST	**1	
00010	01000012	01	0	000012	0		UJP	**2	
00011	01000022	01	0	000022	0		UJP	FINISH1	
00012	20100000	20	0	000000	1		LDA	IX.1	
00013	50000002	50	0	000002	0		MUA	IX1	
00014	30000007	30	0	000007	0		ADA	CARRY	
00015	13077747	13	0	77747	0		SHAQ	-24	
00016	51000006	51	0	000006	0		DVA	THOUSAND	
00017	41127340	41	0	C27340	1		STO	IY.1	
00020	40000007	40	0	000007	0		STA	CARRY	
00021	01000007	01	0	000007	0		UJP	INDEX1	
00022	20000005	20	0	000005	0	FINISH1	LDA	L	
00023	53500000	53	1	00000	1		IAT	1	
00024	20000007	20	0	000007	0		LDA	CARRY	
00025	40127340	40	0	C27340	1		STA	IY.1	
00026	14600000	14	1	00000	2		ENA	0	
00027	40000007	40	0	000007	0		STA	CARRY	
00030	14177776	14	0	77776	1		ENT	-1.1	
00031	20000005	20	0	000005	0		LDA	L	
00032	15477776	15	1	77776	0		INA.S	-1	
00033	44000034	44	0	000034	0		SWA	**1	
00034	10177777	10	0	77777	1	INDEX2	IST	**1	
00035	01000037	01	0	000037	0		UJP	**2	
00036	01000050	01	0	000050	0		UJP	FINISH2	
00037	20100000	20	0	000000	1		LDA	IX.1	
00040	50000003	50	0	000003	0		MUA	IX2	
00041	30000007	30	0	000007	0		ADA	CARRY	
00042	30127341	30	0	C27341	1		ADA	IY+1.1	
00043	13077747	13	0	77747	0		SHAQ	-24	
00044	51000006	51	0	000006	0		DVA	THOUSAND	
00045	41127341	41	0	C27341	1		STO	IY+1.1	
00046	40000007	40	0	000007	0		STA	CARRY	
00047	01000034	01	0	000034	0		UJP	INDEX2	
00050	20000005	20	0	000005	0	FINISH2	LDA	L	
00051	15600001	15	1	00001	2		INA	!	
00052	53500000	53	1	00000	1		IAT	1	
00053	20000007	20	0	000007	0		LDA	CARRY	
00054	40127340	40	0	C27340	1		STA	IY.1	
00055	14600000	14	1	00000	2		ENA	0	
00056	40000007	40	0	000007	0		STA	CARRY	
00057	14177776	14	0	77776	1		ENT	-1.1	
00060	20000005	20	0	000005	0		LDA	L	
00061	15477776	15	1	77776	0		INA.S	-1	
00062	44000063	44	0	000063	0		SWA	**1	
00063	10177777	10	0	77777	1	INDEX3	IST	**1	
00064	01000066	01	0	000066	0		UJP	**2	

COMPASS-32	(2.1)	RANDOM			
00065	01000077	01 0 P00077	0	UJP	FINISH3
00066	20100000	20 0 C00000	1	LDA	IX.1
00067	50000004	50 0 D00004	0	MVA	IX3
00070	30000007	30 0 000007	0	ADA	CARRY
00071	30127342	30 0 C27342	1	ADA	IY+2.1
00072	13077747	13 0 77747	0	SHAQ	-24
00073	51000006	51 0 000006	0	DVA	THOUSAND
00074	41127342	41 0 C27342	1	STO	IY+2.1
00075	40000007	40 0 D00007	0	STA	CARRY
00076	01000003	01 0 P00003	0	UJP	INDEX3
00077	20000007	20 0 000007	0	FINISH3 LDA	CARRY
00100	03100105	03 0 P00105	1	AZLINE	ZERO
00101	20000005	20 0 000005	0	LDA	L
00102	15600002	15 1 00002	2	INA	2
00103	40000005	40 0 D00005	0	STA	L
00104	01000114	01 0 P00114	0	UJP	CHANGE
00105	20000005	20 0 000005	0	ZERO LDA	L
00106	15600002	15 1 00002	2	INA	2
00107	53500000	53 1 00000	1	TAT	1
00110	15600001	15 1 00001	2	INA	1
00111	40000005	40 0 000005	0	STA	L
00112	20000007	20 0 000007	0	LDA	CARRY
00113	40127340	40 0 C27340	1	STA	IY.1
00114	14177776	14 0 77776	1	CHANGE FNT	-1.1
00115	20000005	20 0 000005	0	LDA	L
00116	15477776	15 1 77776	0	INA.S	-1
00117	44000120	44 0 P00120	0	SWA	*+1
00120	10177777	10 0 77777	1	INDEX4 IST	**+1
00121	01000123	01 0 P00123	0	UJP	*+2
00122	01000126	01 0 P00126	0	UJP	FINISH4
00123	20127340	20 0 C27340	1	LDA	IY.1
00124	40100000	40 0 C00000	1	STA	IX.1
00125	01000120	01 0 P00120	0	UJP	INDEX4
00126	20000001	20 0 000001	0	FINISH4 LDA	IDFC
00127	30000000	30 0 000000	0	ADA	NODEC
00130	40000001	40 0 000001	0	STA	IDFC
00131	01000000	01 0 P00000	0	UJP	RANDOM
00000				DATA	
00001				NODEC	BSS 1
00002				IDFC	BSS 1
00003				IX1	BSS 1
00004				IX2	BSS 1
00005				IX3	BSS 1
00006				L	BSS 1
00007				THOUSAND	BSS 1
00008				CARRY	BSS 1
00009				COMMON	
00010				IX	BSS 12000
27340				IY	BSS 12000
				END	

NUMBER OF LINES WITH DIAGNOSTICS

3200 FORTRAN (2.1)

03/03/66

```
PROGRAM CLL CVRT
1 FORMAT (11)
2 FORMAT (20H PARITY ERROR ON OUTPUT TAPE)
3 FORMAT (20H PARITY ERROR ON INPUT TAPE)
4 FORMAT (7H////)
READ 1, IF PRINT
CALL CONVERT (IF PRINT, I PARITY 1, I PARITY 0)
IF (I PARITY 1 .EQ. 1) GO TO 11
10 PRINT 3
PRINT 4
GO TO 99
11 IF (I PARITY 0 .EQ. 1) GO TO 99
12 PRINT 2
PRINT 4
99 CONTINUE
END
```

3200 FORTRAN DIAGNOSTIC RESULTS - FOR CLLCVRT

NO ERRORS

3200 FORTRAN (2.1)

03/03/66

```

SUBROUTINE CONVERT( IF PRINT, I PARITY 1, I PARITY 0)
DIMENSION X(500)
REAL LOGF
1 FORMAT (110)
2 FORMAT (5E20.10)
3 FORMAT (//)
5 FORMAT (1H1)
REWIND 1
REWIND 5
I PARITY 1 = 0
I PARITY 0 = 0
PRINT 5
20 BUFFER IN (5,1)(I BUF, I BUF)
21 GO TO (21, 22, 23, 77), UNITSTF(5)
23 END FILE 1
PRINT 5
RETURN
77 I PARITY 1 = 1
RETURN
22 I BUF = 1 BUF
BUFFER IN (5,1)(X(1),X(I BUF))
24 GO TO (24, 25, 23, 77), UNITSTF(5)
25 I BUF 2 = I BUF/2
PI = 3.1415926536
TPI = 2.0* PI
DO 28 I=1, I BUF 2
I2=I*2
I2M1= I2-1
X1= SQRT(-2.0*LOGF(X(I2M1)))
X2= TPI*X(I2)
X(I2M1)=X1*COSE(X2)
28 X(I2)=X1*SINF(X2)
BUFFER OUT (1,1)( I BUF, I BUF)
31 GO TO (31, 32, 32, 78), UNITSTF(1)
78 I PARITY 0 = 1
RETURN
32 I BUF = I BUF
BUFFER OUT (1,1)(X(1), X(I BUF))
33 GO TO (33, 34, 34, 78), UNITSTF(1)
34 IF( IF PRINT .EQ. 1) 40, 20
40 PRINT 1, I BUF
PRINT 2, (X(I), I=1, I BUF)
PRINT 3
GO TO 20
END

```

3200 FORTRAN DIAGNOSTIC RESULTS - FOR CONVERT

```

NO ERRORS
EQUIP.1=MT
LOAD.56
RUN

```

```

PROGRAM GAUSSIAN
COMMON SA(13), SB(12), C(78, 79), X(78), T(12, 12), F(12, 12),
1      A(12, 12), B(12, 12)
1 FORMAT (5I4)
2 FORMAT (4F20.6)
3 FORMAT (///, 25H THE H MATRIX IS SINGULAR)
4 FORMAT (///)
5 FORMAT (6F20.10)
6 FORMAT (///, 26H BIGSET MATRIX IS SINGULAR)
7 FORMAT (///, 27H PARITY ERROR ON INPUT TAPE)
8 FORMAT (///, 28H PARITY ERROR ON OUTPUT TAPE)
9 FORMAT (///, 23H T TSTAR DOES NOT EXIST)
10 FORMAT (///, 25H TR TRSTAR DOES NOT EXIST)
11 FORMAT (6I1)
12 FORMAT (14H DETERMINANT = ,E20.10,/)
100 FORMAT (1H1)
66 FORMAT (36H CONSTANTS OF DENOMINATOR POLYNOMIAL )
67 FORMAT (34H CONSTANTS OF NUMERATOR POLYNOMIAL )
68 FORMAT (8H N M )
69 FORMAT (5X, 13H TIME INTERVAL )
70 FORMAT (10X, 5H (SEC) )
REWIND 2
READ 1, NO SKIP
DO 61 I=1, NO SKIP
63 BUFFER IN (2,1) (ISKIP, ISKIP)
62 GO TO (62, 63, 64, 64), UNITSTF(2)
64 PRINT 65
65 FORMAT (40H PARITY ERROR ON MAG TAPE WHILE SKIPPING )
GO TO 63
61 CONTINUE
60 READ 1, N, M
IF (N.EQ.0) 99, 71
71 NP1 = N+1
MP1 = M+1
PRINT 100
PRINT 58
PRINT 1, N, M
READ 2, (SA(I), I=1, NP1)
READ 2, (SB(I), I=1, MP1)
READ 2, DT
READ 1, ISTOP, IPOWER, IPOWERTR, IPOWERRB, IPOWERC
READ 11, IF N PRINT, IF X PRINT, IFT, IFTB, IFB, IFC
PRINT 4
PRINT 66
PRINT 5, (SA(I), I=1, NP1)
PRINT 4
PRINT 67
PRINT 5, (SB(I), I=1, MP1)
PRINT 4
PRINT 69
PRINT 70
PRINT 5, DT
PRINT 4
CALL FIND B(N)
DO 50 I=1, N
50 PRINT 5, (C(I,J), J=1, NP1)
PRINT 4
CALL GAUSS P (N, KK, IPOWERRB, DET)
PRINT 12, DET

```

```

IF(KK) 21, 20
21 PRINT 3
IF (IFH .EQ. 1) 20, 60
20 CALL FIND CM(N)
PRINT 5, (X(I), I=1, N)
PRINT 4
DO 51 I=1, N
51 PRINT 5, (C(I,J), J=1, N)
CALL T T STAR ( N, KK, IPOWERT)
IF(KK) 41, 40
41 PRINT 9
IF (IFT .EQ. 1) 40, 60
40 CALL MAKE A(N)
PRINT 4
DO 52 I=1, N
52 PRINT 5, (T(I,J), J=1, N)
PRINT 4
DO 53 I=1, N
53 PRINT 5, (A(I,J), J=1, N)
CALL EXP A DT (N, DT)
PRINT 4
DO 54 I=1, N
54 PRINT 5, (E(I,J), J=1, N)
CALL HIGSET(N)
NN= (N*(N+1))/2
NNP1= NN+1
PRINT 4
DO 55 I=1, NN
55 PRINT 5, (C(I,J), J=1, NNP1)
CALL GAUSS P (NN, KK, IPOWERC, DET )
PRINT 4
PRINT 12, DET
IF(KK) 22, 23
22 PRINT 6
IF (IFC .EQ. 1) 23, 60
23 L=0
PRINT 4
PRINT 5, (X(I), I=1, NN)
DO 24 I=1, N
DO 24 J=1, N
L=L+1
C(I,J)= X(L)
24 C(J,I)= X(L)
PRINT 4
DO 56 I=1, N
56 PRINT 5, (C(I,J), J=1, N)
DO 25 I=1, N
DO 25 J=1, N
25 A(I,J)= T(I,J)
CALL T T STAR(N, KK, IPOWERT)
IF(KK) 43, 42
43 PRINT 10
IF (IFIP .EQ. 1) 42, 60
42 REWIND 1
PRINT 4
DO 57 I=1, N
57 PRINT 5, (T(I, J), J=1, N)
PRINT 100
CALL SOLUTION(N,M,IPARITYI,IPARITYO,ISTOP,IFNPRINT,IFXPRINT)
IF(IPARITYI) 26, 98
26 PRINT ?
GO TO 99

```

```
98 IF (IPARITY) 27, 60  
27 PRINT R  
99 PRINT 100  
END
```

3200 FORTRAN DIAGNOSTIC RESULTS - FOR GAUSSIAN

UNDEFINED SIMPLE VARIABLES
ISKIP

3200 FORTRAN (2.1)

31/03/66

```
SUBROUTINE FINDB(N)
COMMON SA(13),SH(12),H(78,79)
INTEGER Q, QM1
NM1= N-1 $ NP1= N+1
DO 1 I=1, NM1
1 H(I, NP1)= 0.0
  H(N, NP1)= 0.5
DO 2 I=1, N
DO 2 J=1, N
2 H(I, J)= 0.0
DO 3 K=1, N
KM1= K-1 $ KL=KM1/2.0 +.51 +1
KU= (N+KM1)/2+1 $ A3= (-1.0)**KM1
DO 4 Q=KL, KU
QM1= Q-1 $ K4= N-2*QM1+K
3 H(K,Q)=A3*(-1.0)**QM1*SA(K4)
  END
```

3200 FORTRAN DIAGNOSTIC RESULTS - FOR FINDB

NO ERRORS

3200 FORTRAN (2.1)

31/03/66

```

SUBROUTINE FINDCM(N)
COMMON SA(13),SR(12),CM(78,79),SM(78)
DO 1 I=1, N
DO 1 J=1, N
1 CM(I,J)= 0.0
DO 2 I=1, N
A= -1.0
DO 2 J=1, N, 2
A= -A $ K= (I+J)/2
CM(I,J)= A*SM(K)
2 CM(J,I)= CM(I,J)
END
```

3200 FORTRAN DIAGNOSTIC RESULTS - FOR FINDCM

NO ERRORS

3200 FORTRAN (2.1)

5/03/66

```

SUBROUTINE T T STAR (N, KK, IPOWER)
COMMON SA(13),SH(12),CM(78,79),SM(78),T(12,12)
C*****
C      THIS SUBROUTINE IS DESIGNED TO FACTOR A POSITIVE DEFINITE,
C      REAL, SYMMETRIC MATRIX, CM, INTO A LOWER-TRIANGULAR MATRIX, T,
C      WITH POSITIVE ELEMENTS ON THE MAIN DIAGONAL, SUCH THAT CM=T T*.
C*****
C      KK=0 IF THE TRIANGULAR MATRIX SOLUTION EXISTS, IF ALL THE
C      ELEMENTS ON THE MAIN DIAGONAL OF THE CM MATRIX ARE POSITIVE.
C      KK=1 IF THE SOLUTION DOES NOT EXIST.
C*****
      KK=0
      TEN POWER = 10.0**(-IPOWER)
      DO 6 I=1, N
        IF (CM(I,I) .LT. TEN POWER) 7, 6
      7 KK=1
      GO TO 4
      6 CONTINUE
      8 NM1 = N-1
      DO 5 I=1, NM1
        IP1= I+1
        DO 5 J=IP1, N
          5 T(I,J)= 0.0
          T(1,1)= SQRTF(CM(1,1))
          DO 1 I=2, N
            1 T(I,1)= CM(I,1)/T
            DO 2 J=2, N
              SUM= 0.0 $ JM1= J-1
              DO 3 K=1, JM1
                3 SUM= SUM+ T(J,K)*T(J,K)
              T(J,J)=SQRTF(CM(J,J)-SUM)
              JP1= J+1
              DO 2 I=JP1, N
                SUM=0.0
                DO 4 K=1, JM1
                  4 SUM= SUM+T(I,K)*T(I,K)
                2 T(I,J)=(CM(I,J)-SUM)/T(J,J)
              END

```

3200 FORTRAN DIAGNOSTIC RESULTS - FOR TTSTAR

NO ERRORS

3200 FORTRAN (2.1)

31/03/66

```
SUBROUTINE MAKE A(N)
COMMON SA(13),SH(12),C(78,79),X(78),T(12,12),F(12,12),A(12,12)
NM1= N-1
NP2= N+2
DO 1 J=1, NM1
DO 1 J=1, N
1 A(I,J)= 0.0
DO 3 I=1, NM1
IP1= I+1
3 A(I,IP1)=1.0
DO 4 I=1, N
NI= NP2-I
4 A(N,I)= -SA(NI)
END
```

3200 FORTRAN DIAGNOSTIC RESULTS - FOR MAKEA

NO ERRORS

```

SUBROUTINE EXP A DT(N,DT)
COMMON SA(13),SH(12),AN(12,12),B(12,12),DUMMY(5874),X(78),T(12,12)
1  ,E(12,12), A(12,12)
  NM1 = N-1
  FACN= 1.0
  DO 1 I=1, N
  DO 1 J=1, N
  AN(I,J)= A(I,J)
1  E(I,J)= 0.0
  DO 2 I=1, N
2  E(I,1)= 1.0
  EPSCHK= 0.0
  DO 7 I=1, N
  CHECK= ABSF(A(N,I))
  IF(CHECK .GT. EPSCHK) 8,7
8  EPSCHK= CHECK
7  CONTINUE
  EPSCHK = EPSCHK*1.0E-12
5  TFN= DT/FACN
  DO 12 I=1, N
  DO 12 J=1, N
12  B(I,J)= AN(I,J)*TFN
  DO 28 J=1, N
  DO 28 I=1, N
  IF(ABSF(B(I,J)) .LT. EPSCHK) 28, 9
28  CONTINUE
  GO TO 99
9  DO 6 I=1, N
  DO 6 J=1, N
6  E(I,J)= E(I,J)+B(I,J)
  DO 11 I=1, NM1
  IP1 = I+1
  DO 11 J=1, N
11  AN(I,J) = B(IP1,J)
  BNN = B(N,N)
  AN(N,1) = BNN*A(N,1)
  DO 10 I=1, NM1
  IP1 = I + 1
10  AN(N,IP1) = B(N,I)+BNN*A(N,IP1)
  FACN= FACN*1.0
  GO TO 5
99  CONTINUE
  END

```

3200 FORTRAN DIAGNOSTIC RESULTS - FOR EXPADT

NO ERRORS
LOAD,56
RUN

3200 FORTRAN (2.1)

31/03/66

```

SUBROUTINE GAUSS P (N, IF SINGLR, IPOWER, DET)
COMMON SA(13), SR(12), A(78, 79), X(79)
REAL M
INTEGER R, RP1, RMAX
ISIGN = 1
NM1 = N-1
NP1 = N+1
IF SINGLR = 0
TEN POWER = 10.0**(-IPOWER)
DO 1 R=1, NM1
AMAX = ABSF(A(R,R))
RMAX = R
RP1 = R+1
DO 3 I=RP1, N
ABSFA = ABSF(A(I,R))
IF (ABSFA .GT. AMAX) 2, 3
2 AMAX = ABSFA
RMAX = I
3 CONTINUE
IF (AMAX .LT. TEN POWER) 5, 4
5 IF SINGLR = 1
4 IF (R .EQ. RMAX) 10, 9
9 DO 6 I=R, N
ATEMP = A(R,I)
A(R,I) = A(RMAX,I)
6 A(RMAX,I) = ATEMP
ATEMP = A(R,NP1)
A(R,NP1) = A(RMAX,NP1)
A(RMAX,NP1) = ATEMP
ISIGN = -ISIGN
10 DO 11 I=RP1, N
M = -A(I,R)/A(R,R)
DO 7 J=RP1, N
7 A(I,J) = A(I,J) + M*A(R,J)
1 A(I,NP1) = A(I,NP1) + M*A(R,NP1)
IF (ABSF(A(N,N)) .LT. TENPOWER) 11, 8
11 IF SINGLR = 1
8 CONTINUE
DO 21 I=1, N
IM1 = I - 1
NN = N-IM1
XX = A(NN, NP1)
DO 22 J=1, IM1
NP1J = NP1 - J
22 XX = XX - A(NN, NP1J)*X(NP1J)
21 X(NN) = XX/A(NN, NN)
DET = 1.0
DO 31 I=1, N
31 DET = DET*A(I,I)
DET = DET*ISIGN
END

```

3200 FORTRAN DIAGNOSTIC RESULTS - FOR GAUSSP

NO ERRORS

3200 FORTRAN (2.1)

31/03/66

```
SUBROUTINE INDEX(JJ,I,J,N2)
JJ= J*((I-1)*(N2-1))/2
END
```

3200 FORTRAN DIAGNOSTIC RESULTS - FOR INDEX

NO ERRORS

3200 FORTRAN (2.1)

31/03/66

```
SUBROUTINE HIGSET (N)
COMMON SA(13), SR(12), C(78, 79), X(78), T(12, 12), E(12, 12),
1   A(12, 12), H(12, 12)
NN=(N*(N+1))/2 $ NNP1= NN+1
DO 7 I=1, N
DO 7 J=1, N
7 B(I,J)= E(I,N)*E(J,N)
H(N,N)= H(N,N)-1.0
DO 1 I=1, NN
DO 1 J=1, NNP1
1 C(I,J)= 0.0
N2= 2*N $ L=0
DO 2 J=1, N
JP1= J+1
DO 2 I=J, N
L=L+1
IP1= 1+1
DO 3 K=1, J
CALL INDEX(JJ,K,J,N2)
3 C(L,JJ)=C(L,JJ)+A(I,K)
DO 4 K=JP1, N
CALL INDEX(JJ,J,K,N2)
4 C(L,JJ)=C(L,JJ)+A(I,K)
DO 5 K=1, I
CALL INDEX(JJ,K,T,N2)
5 C(L,JJ)= C(L,JJ)+A(J,K)
DO 6 K=IP1, N
CALL INDEX(JJ,I,K,N2)
6 C(L,JJ)=C(L,JJ)+A(J,K)
2 C(L,NNP1)= H(I,J)
END
```

3200 FORTRAN DIAGNOSTIC RESULTS - FOR HIGSET

NO ERRORS

3200 FORTRAN (2.1)

31/03/66

```

SUBROUTINE SOLUTION(N,M,IPARITY1,IPARITY0,ISTOP,IFNPRINT,IFXPRINT)
COMMON SA(12), SR(12), Z(12), ZP(12), WP(12), W(600), X(5604),
1 TW(12, 12), F(12, 12), T(12, 12)
IPARITY1=0
IPARITY0=0
BUFFER IN (1,1)(MAXW, MAXW)
30 GO TO (30, 31, 22, 3), UNITSTF(1)
31 IWMAX= MAX W
BUFFER IN(1,1)(W(1),W(IWMAX))
1 GO TO (1, 2, 22, 3), UNITSTF(1)
3 IPARITY1=1
RETURN
2 MP1= M+1
MP2= M+2
II=N
ISTART=2
DO 4 I=1, N
Z(I)=0.0
DO 4 J=1, I
4 Z(I)= Z(I)+T(I,J)*W(J)
X(I)= 0.0
DO 5 I=1, MP1
J= MP2-I
5 X(I)=X(I)+SR(I)*Z(J)
7 DO 8 I=ISTART, ISTOP
IPRINT = I-1
9 II= II+N
IF(IWMAX .LT. II) 10, 11
10 IF(II-N .EQ. IWMAX) 12, 13
12 JJ= 0
GO TO 14
13 JJ= II-IWMAX
KL= II-N
DO 15 KJ=1, JJ
KL= KL+1
15 W(KJ)= W(KL)
14 BUFFER IN (1,1)(MAX W, MAX W)
32 GO TO (32, 33, 34, 3), UNITSTF(1)
34 IF (IPRINT) 35, 22
35 IF (IF N PRINT .EQ. 1) 70, 71
70 PRINT 72, IPRINT
72 FORMAT(////,I10, 30H POINTS ARE IN THIS DATA BLOCK, /)
71 IF (IF X PRINT .EQ. 1) 73, 74
73 PRINT 28, (X(J), J=1, IPRINT)
74 BUFFER OUT (2,1)(IPRINT, IPRINT)
62 GO TO (62, 63, 63, 25), UNITSTF(2)
63 BUFFER OUT(2,1)(X(1), X(IPRINT))
64 GO TO (64, 22, 22, 25), UNITSTF(2)
33 IWMAX= MAX W + JJ
JJ=JJ+1
BUFFER IN(1,1)(W(JJ),W(IWMAX))
23 GO TO(23, 24, 22, 3), UNITSTF(1)
24 II=0
GO TO 1
11 JJ= II-N
DO 16 J=1, N
JJ=JJ+1
16 WP(J)= W(JJ)
DO 17 K=1, N
ZZP=0.0

```

```
DO 18 J=1, N
18 ZP= ZP+F(K,J)*7(J)+TR(K,J)*WP(J)
17 ZP(K)= ZP
  XX=0.0
DO 19 K=1, MP1
  J= MP2-K
19 XX=XX+SB(K)*ZP(J)
  X(I)= XX
DO 3 J=1, N
  Z(J)= ZP(J)
  IF(IF N PRINT .EQ. 1) 75, 76
75 PRINT 72, ISTOP
76 IF (IF X PRINT .EQ. 1) 77, 78
77 PRINT 28, (X(I), I=1, ISTOP)
28 FORMAT (6E20.10)
78 BUFFER OUT(2,1)(ISTOP, ISTOP)
60 GO TO (60, 61, 61, 25), UNITSTF(2)
61 BUFFER OUT (2,1)(X(1), X(ISTOP))
21 GO TO (21, 6, 6, 25), UNITSTF(2)
25 IPARITY=1
  RETURN
  6 ISTART=1
  GO TO 7
22 END FILE 2
  CONTINUE
  ENDD
```

3200 FORTRAN DIAGNOSTIC RESULTS - FOR SOLUTION

UNDEFINED SIMPLE VARIABLES
MAXV
FQUIP, 2=MT
LOAD.56
RUN

APPENDIX

In addition to the "white sequence," $\{e^i\}$, $1 \leq i \leq 4266$, discussed in the text of this paper, seven other "white sequences" have been generated by ISLASHO. These eight "white sequences" have been permanently stored on magnetic tape; this tape is labeled Tape B.

The eight "white sequences" (in the order in which they are stored on Tape B) are:

$$v_i^{(1)} = \{\theta^i\}, 1 \leq i \leq 4004; \theta = 9.63778431 \approx \log_{10} N$$

$$v_i^{(2)} = \{\theta^i\}, 1 \leq i \leq 4010; \theta = 1.27323954 \approx 4/\pi$$

$$v_i^{(3)} = \{\theta^i\}, 1 \leq i \leq 4006; \theta = 3.04800610 \approx 0.2(\text{number of cm/ft})$$

$$v_i^{(4)} = \{\theta^i\}, 1 \leq i \leq 4006; \theta = 1.46459189 \approx \sqrt[3]{\pi}$$

$$v_i^{(5)} = \{\theta^i\}, 1 \leq i \leq 4006; \theta = 4.97149872 \approx \log_{10} \pi$$

$$v_i^{(6)} = \{\theta^i\}, 1 \leq i \leq 4266; \theta = 2.71828183 \approx e$$

$$v_i^{(7)} = \{\theta^i\}, 1 \leq i \leq 4008; \theta = 2.30258509 \approx 1/N$$

$$v_i^{(8)} = \{\theta^i\}, 1 \leq i \leq 4006; \theta = 5.72957795 \approx 0.1(180/\pi).$$

In $v_i^{(1)}$ and $v_i^{(7)}$, define $N = \log_{10} e$.

```

IDEC(8)
L(8)
IX(8)(1) . . . . IX(8)(L(8))
      end-of-file
      end-of-file.

```

Here IDEC, L, and IX(I) are defined the same as in the section of the text on ISLASHO.

The "white sequence" $v_i^{(1)}$ on Tape B can be used as direct input to GAUSSIAN. In order to use $v_i^{(s)}$, $2 \leq i \leq 8$, as input to GAUSSIAN, either the specified "white sequence" must first be placed on an auxiliary tape or GAUSSIAN must be modified to skip over previous "white sequences."

The first 4000 points of each of the eight above "white sequences" have been interlaced by the scheme

$$v_{8(i-1)+1} = v_i^{(1)}, v_{8(i-1)+2} = v_i^{(2)}, \dots, v_{8(i-1)+8} = v_i^{(8)},$$

for $1 \leq i \leq 4000$. This 32,000 point "white sequence" has also been stored permanently on a second magnetic tape which is labeled Tape C.

Tape C has the format:

```

500
v1 . . . . v500
.
.
.
500
v31,501 . . v32,000
      end-of-file
      end-of-file.

```

Tape C can be used directly as input to GAUSSIAN.

Tape B has the following format:

500

$v_1^{(1)} \dots v_{500}^{(1)}$

.

.

.

500

$v_{3501}^{(1)} \dots v_{4000}^{(1)}$

4

$v_{4001}^{(1)} \dots v_{4004}^{(1)}$

end-of-file

IDEC⁽¹⁾

L⁽¹⁾

IX⁽¹⁾₍₁₎ . . . IX⁽¹⁾_(L)

end-of-file

.

.

.

500

$v_1^{(8)} \dots v_{500}^{(8)}$

.

.

.

500

$v_{3501}^{(8)} \dots v_{4000}^{(8)}$

6

$v_{4001}^{(8)} \dots v_{4006}^{(8)}$

end-of-file

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