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VALIDATION OF THE
PROTECTED DMS SPECIFICATIONS



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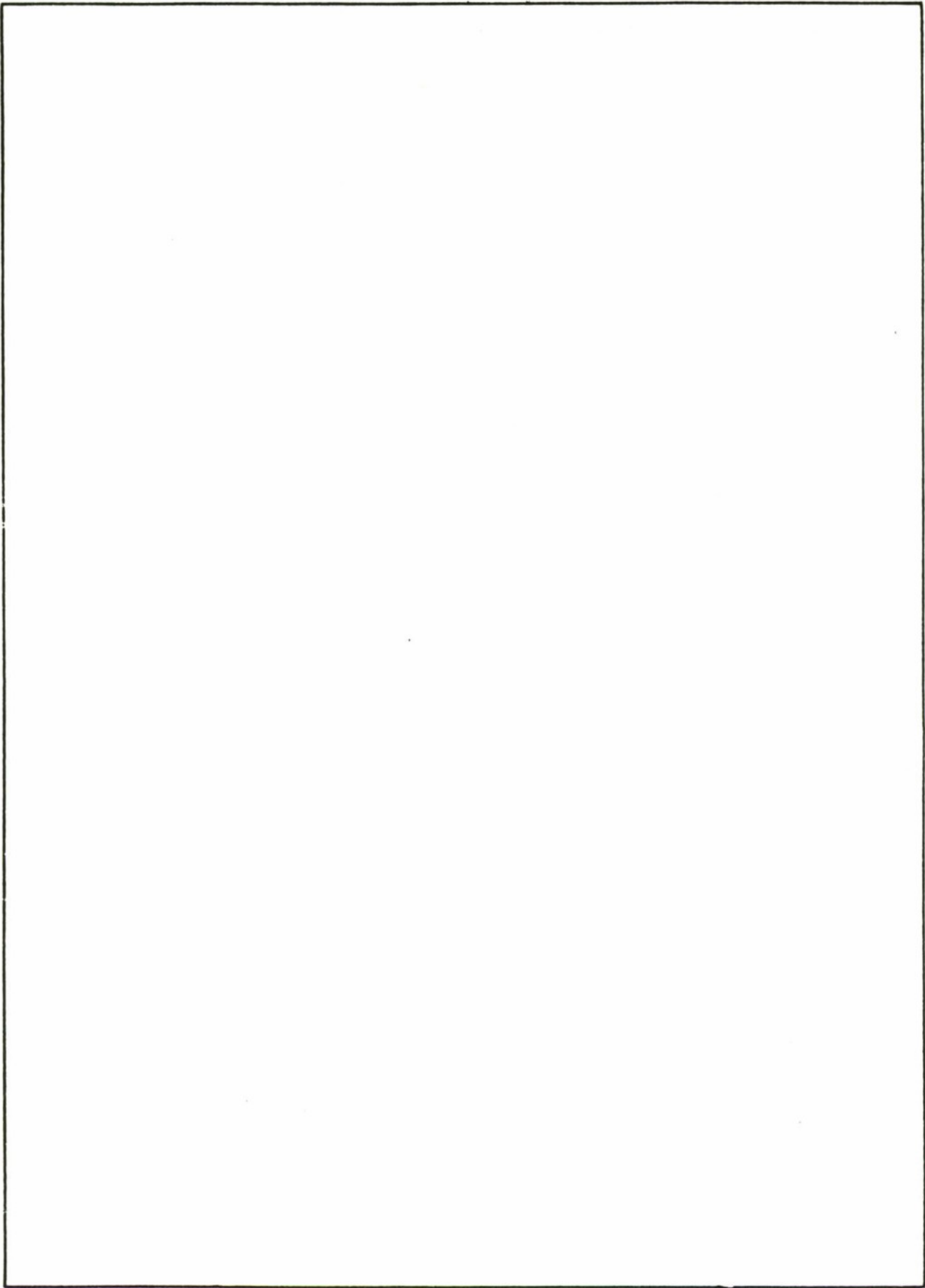
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I INTRODUCTION

The purpose of this report is to prove mathematically that the specifications of the primitives of the protected DMS [1] correctly embody the protection principles designated by the mathematical model [2].

The validation technique is described, and the detailed validation of a representative sample of O-functions is included as an appendix.

II VALIDATION TECHNIQUE

2.1 Objective

The objective of the validation is to prove that the accesses of objects by subjects (in the secure DMS), which are formally specified in appendix IV of reference [1], do in fact conform to the axioms of the mathematical model.

These axioms are taken from section VI of reference [2], and are expressed using the notation and variables of the specifications. In addition, " \succ_s " means "security level dominance", and " \succ_i " means "integrity level dominance". Let S_V and I_V be the security and integrity levels, respectively, of variable V . M_V is the permission matrix associated with variable V . CUR_SEC and CUR_INT are the current security and integrity levels, respectively, of the subject of interest. Then the model axioms are:

- (i) Direct Disclosure (Simple Security)
Subject observes $V \Rightarrow CUR_SEC \succ_s S_V$
- (ii) Indirect Disclosure
Subject modifies $V \Rightarrow CUR_SEC \succ_s S_V$
- (iii) Direct Modification
Subject modifies $V \Rightarrow CUR_INT \succ_i I_V$
- (iv) Indirect "Contamination"
Subject observes $V \Rightarrow CUR_INT \succ_i I_V$
- (v) Tranquility Principle
 - A) CUR_SEC and CUR_INT are constant for any subject; and

B) S_V and I_V are constant for any variable V .

(vi) Discretionary Access Control

Subject accesses $V \Rightarrow (\text{Subject}, \text{access}) \in M_V$.

Essentially, the validation consists of proofs based on the inspection of specification statements, invariants, and case table results. Arguments of reasonableness are used in the determination of an O-function's cases, and to justify discretionary access control.

The mathematical notation and semantics used in the validation are those of the formal specifications, and are described in §B of appendix III of reference [1]. The flow of the validation is maintained by the use of clear English statements.

2.2 Invariants

Invariants are general "relations", or conditions involving variables of the specification which hold true between O-function invocations (i.e. between system states). They are proved inductively by showing that each O-function preserves each invariant [3]. The proofs involve inspection of the effects section of the relevant O-functions.

2.3 Case Tables

The purpose of the case tables is to identify the object accesses in all cases of the specified O-functions. Each case is defined by a distinct relationship among function inputs, and presents the protection levels of all variables observed and modified.

An inspection of a case table will prove that the protection level of each modified variable dominates the levels of all observed variables. Certain lemmas may be required to clearly justify the dominance relationship, and these are proved by invariants or the TRUE/FALSE value of exception conditions. In subsection 5.2 it is explained that this relationship ensures the satisfaction of non-discretionary requirements.

The set of all security-related O-functions is partitioned in such a way that each class of the partition consists of O-functions which possess analogous case tables. That is, input parameters, the number and type of accesses (observe or modify), and their levels are the same. Then the case tables of only certain representatives from each class are included in this report, in appendix I.

III DESIGN FUNDAMENTALS

The specification language was structured to facilitate the validation of the design. The semantics of the symbols used in the specifications most relevant to validation are summarized in this section.

3.1 Objects

Variables in the specifications represent the hidden kernel entities, directories, sign-on lists and the component entities of data base and user working area objects. Variables correspond to model objects since each possesses an identifier, a protection level, and a value.

The identifier of a variable is constructed in such a way as to indicate the characteristics of the design entity it represents, such as:

- (i) user, kernel or data base environment
(W, K or D);
- (ii) the data type of the entity;
- (iii) its owner (creator);
- (iv) its name (for use in W);
- (v) the kind of compound object of which
it is a part (e.g. string, program,
relation); and
- (vi) its protection level.

The identifiers of variables isolated on a user basis (working area and kernel) do not include owner or level

information, since they are "owned" by the current user, and their levels are stored in other variables. Their name is a mnemonic suggesting their role in the design.

3.2 Protection Level Assignment

Table 3.1 lists the protection levels which are assigned to the variables in the specifications.

The mathematical model explicitly gives the levels of the data base variables. Directories and sign-on lists are assigned the level of their contents. The level of a component entity of a data base object is taken to be the level component of its identifier.

Variables in the user's working area will assume the user's current (SIGNON) level. This is because they may be both observed and modified by the user.

In the hidden kernel area, certain variables contain data describing the current user, or his activities. Since the user set this data (by parameters) it will be assigned his current level. These variables are indicated in table 3.1 by level "K_CUR_LEVEL". The reserve table reflects successful reservation requests. The space quota gives the current session space resources of the user. The levels of the accumulator and temporaries reflect past "level" parameters. The current time is included to guarantee its correctness.

The levels of the accumulator and temporaries are set equal to the contents of K_LACC and K_LX (Y and Z too) respectively, since that is their purpose.

The open table is a five dimensional boolean array, which is partitioned according to the level dimension. That is, the level of each part of the open table is the level of the objects identified in that part, unless that level is strictly dominated by K_CUR_LEVEL. In that case the level of such a part is taken to be K_CUR_LEVEL [c.f. § 4.1].

Variable	Identifier	Level
<u>Database (multi-level)</u>		
directory	D_D(level)	level
sign-on list	D_Q(level)	level
exact size	D_E(o,n,t,l)	1
format	D_F(o,n,t,l)	1
history	D_H(o,n,t,l)	1
permission matrix	D_M(o,n,t,l)	1
open list	D_O(o,n,t,l)	1
reserve queue	D_R(o,n,t,l)	1
values	D_V(o,n,t,l)	1
maximum size	D_Z(o,n,t,l)	1
<u>Kernel (isolated)</u>		
current user level	K_CUR_LEVEL	K_CUR_LEVEL
current user identifier	K_CUR_ID	K_CUR_LEVEL
current time	K_CUR_TIME	K_CUR_LEVEL
session space quota	K_CUR_QTA	K_CUR_LEVEL
reserve table	K_RESERVE	K_CUR_LEVEL
open table	K_OPEN	K_CUR_LEVEL OR K_OPEN [level] ¹
accumulator level	K_LACC	K_CUR_LEVEL
accumulator contents	K_IACC	K_LACC
accumulator format	K_FACC	K_LACC
accumulator values	K_VACC	K_LACC
level of temporary	K_LX ₂	K_CUR_LEVEL
temporary contents	K_IX ₂	K_LX
temporary format	K_FX ₂	K_LX
temporary values	K_VX	K_LX
<u>Working area (isolated)</u>		
return code	W_CODE	K_CUR_LEVEL
relation value table	W_Vname	K_CUR_LEVEL
relation format	W_Fname	K_CUR_LEVEL

Table 3.1 Protection Level Assignment

¹ The dominating level is assigned to each part of the open table.

² Y and Z temporaries as well.

3.3 Object Access

The accessing of objects is represented in the specifications by usage of a variable's identifier. An observe access is represented by a variable's appearance in a V-function derivation, an exception condition, or on the right-hand side of the "arrow" (+) in an assignment statement in an effects section. A modification access is represented by the appearance of a variable on the left-hand side of the "arrow" in an assignment statement.

3.4 Subjects

A subject is a process associated with each O-function invocation, whose sole purpose is to have the capability (authorization) to perform the observations and modifications required by the effects of the O-function.

When an O-function is invoked, one subject sets the return code (W_CODE) and the levels of the accumulator (K_LACC) and temporaries (K_LX, K_LY and K_LZ). The level of this subject is K_CUR_LEVEL [c.f. table 3.1]. Another subject performs all other modifications. An inspection of the case tables in appendix I will reveal that these modifications are all either at level "lv" (a parameter) or level "LA" (level of the accumulator), except for SIGNON, SIGNOFF and MOVE. Therefore the level of the second subject is established as that level.

Since the three exceptions are functions which perform multi-level observations and modifications, they are restricted to be performed by "trusted" processes such as the user control process (UCP) and the data base administrator (DBA).

IV PROOF OF INVARIANTS

4.1 Minimum K-level

4.1.1 Invariant

$$K_LACC \succ K_CUR_LEVEL$$

and

$$K_Lx \succ K_CUR_LEVEL, \text{ for } x = X, Y, Z.$$

The levels of the kernel accumulator and temporaries always dominate the current user's level (K_CUR_LEVEL).

4.1.2 Requirement

This invariant ensures that parameter data (which is at level K_CUR_LEVEL) is never found in a kernel entity at a strictly dominated level (prohibits "write-downs").

4.1.3 Proof

An inspection of the case tables in appendix I will reveal that K_LACC is set by the following primitives:

(i) $DKD, DKQ, LIST_DOWN,$ and $WKB;$

An inspection of their specifications will reveal that the level K_LACC effect is:

$$K_LACC \leftarrow K_CUR_LEVEL$$

(ii) $DKE, DKH, DKM, DKR, DKV,$ and $DKZ;$

An inspection of their specifications will reveal that the level K_LACC effect is:

$K_LACC \leftarrow LEV \text{ IF } LEV \succ K_CUR_LEVEL \text{ ELSE } K_CUR_LEVEL$

where LEV is the level parameter

Clearly, these assignments result in the relation:

$K_LACC \succ K_CUR_LEVEL$

(iii) ASSIGN

The level effect statement is:

$K_LACC \leftarrow K_Lsource \text{ IF } source \in \{ACC, X, Y, Z\}.$

An inspection of the case tables will reveal that only the ASSIGN primitive modifies the temporaries X, Y and Z. The function of ASSIGN is such that data can be assigned to temporaries from two sources only:

- the kernel accumulator; or
- a working area (W) value.

In the first case:

$K_Lx \succ K_CUR_LEVEL, x \in \{X, Y, Z\},$

by the conclusions of (i) and (ii), since K_Lx was in K_LACC once.

In the second case, K_Lx is not modified, so the invariant is not affected. Therefore, $K_LACC \succ K_CUR_LEVEL$ if $source \in \{\bar{X}, \bar{Y}, \bar{Z}\}.$

If $source = ACC$, there is no change in K_LACC and the invariant is unaffected.

If $source \notin \{ACC, X, Y, Z\}$, then K_LACC is not changed, so the invariant is unaffected.

Note that the assignment of a working area value will occur only when the data conforms to the format data in the kernel, prohibiting assignment of values only to null kernel entities. Therefore, the levels of the kernel entities will dominate the values by the above arguments.

(iv) SIGNOFF

The effect statements:

$K_LACC, K_Lx = \emptyset$, $x \in \{X, Y, Z\}$
 $K_CUR_LEVEL = \emptyset$

trivially satisfy the invariant.

Q.E.D.

4.2 Open Table

4.2.1 Invariant

Subscript a variable by t to indicate its value in the system state at time t . Then:

$$\begin{aligned} K_OPEN_t(\text{owner}, \text{name}, \text{type}, \text{level}, \text{access}) = \text{TRUE} \\ \implies \\ K_OPEN_{t-1}(\text{owner}, \text{name}, \text{type}, \text{level}, \text{access}) = \text{TRUE} \\ \text{OR} \\ (K_CUR_ID, \text{access}) \in D_M_{t-1}(\text{owner}, \text{name}, \text{type}, \text{level}) \\ \text{OR} \\ K_CUR_ID = \text{owner} \end{aligned}$$

If an object has some access granted to it in a user's open table, then it was there in the preceding system state, or it was authorized in the object's permission matrix in the preceding state (when OPEN was requested).

4.2.2 Requirement

This invariant assures that the discretionary authorization mechanisms function correctly.

4.2.3 Proof

An inspection of the case tables in appendix I reveals that only O_APPEND, O_DELETE and SIGNOFF modify K_OPEN. Effects statements in O_DELETE and SIGNOFF set the K_OPEN entry of interest to FALSE, so only O_APPEND is relevant to this invariant.

There are two possibilities:

- (i) O_APPEND was invoked at time (t-1), producing the system state at time t;
- or (ii) O_APPEND was not invoked at time (t-1).

Suppose O_APPEND was not invoked at time (t-1). If $K_OPEN(owner, name, type, level, access) = TRUE$ at time t, then it must be TRUE at time (t-1) as well, since K_OPEN is not modified in this case.

Otherwise, suppose O_APPEND was invoked at time (t-1). If the open table entry of interest is TRUE at time (t-1), then O_APPEND will return exception code "DO", and the entry will remain unchanged to time t.

If the entry is FALSE at time (t-1) and TRUE at time t, then exception ND = FALSE. That is, by the derivation of the ND exception:

$$\begin{aligned} & \sim(K_CUR_ID \neq owner) \wedge (\exists x(K_CUR_ID, x) \in D_M_{t-1}(id)) \\ & \quad \equiv \\ & (K_CUR_ID = owner) \vee (\exists x(K_CUR_ID, x) \in D_M_{t-1}(id)) \end{aligned}$$

Effect [1] of O_APPEND states:

$$K_OPEN(Access_set_O(id)) = TRUE$$

Access_set_0 returns the set of tuples S:

$$S = \{(id,x) : (OWN = K_CUR_ID) \vee (K_CUR_ID,x) \in D_M_{t-1}(id)\}$$

by the derivation of the Access_set_0 and Aith_0 V-functions. The "access" in the invariant is one of these x's, since these are the accesses authorized.

Q.E.D.

4.3 Identification

4.3.1 Invariant

$$\begin{aligned} D_V(id) \in K_Vx \quad \text{and} \quad D_F(id) \in K_Fx \\ \implies \\ K_Ix = id, \text{ for } x = ACC, X, Y, Z. \end{aligned}$$

At all times the identifiers (i.e. K_IACC , K_IX , K_IY , K_IZ) of the contents of the kernel accumulator and temporaries are a correct indication of the identity of their contents.

4.3.2 Requirement

This invariant is required for discretionary authorization, to ensure that data cannot masquerade when being presented to a user (KWA) or copied to the data base (KD...).

4.3.3 Proof

An inspection of the case tables in appendix I will reveal that K_IACC is modified by the O-functions in table 4.1. An inspection of the K_IACC effect in each function specification (included in table 4.1) will reveal that it is appropriate for the function.

Since only ASSIGN affects the temporaries, they contain only data previously in the accumulator. Assigning a user working area value to the values component of the accumulator or a temporary does not affect its identity.

<u>Primitive</u>	<u>Identification Effect</u>
DKD	$K_IACC \leftarrow (K_CUR_ID, 'D', 'R', lv, 'V')$
DKQ	$K_IACC \leftarrow (K_CUR_ID, 'Q', 'R', lv, 'V')$
LIST_DOWN	$K_IACC \leftarrow (K_CUR_ID, DEF_NAME, 'R', K_CUR_LEVEL, 'V')$
DKC	$K_IACC \leftarrow (id, C) , C \in \{E, H, M, R, V, Z\}$
ASSIGN	$K_Itarget \leftarrow K_Isource \text{ IF } source \in \{ACC, X, Y, Z\}$
WKB	$K_IACC \leftarrow (K_CUR_ID, n, 'S', K_CUR_LEVEL, 'V')$
SIGNOFF	$K_IACC \leftarrow \emptyset$

Table 4.1 Table of Identification Data

V JUSTIFICATION OF PROTECTION

5.1 Non-discretionary Policy

The secure DMS has been designed in such a way that there is no modification access without an observation access, and vice versa.

An O-function execution involves two subjects: One executes at the single level of the data being modified, and performs all effects except the setting of W_CODE, K_LACC, K_LX, K_LY and K_LZ. (This is "Subject 2" in appendix I.) The other subject executes at the user's current level (K_CUR_LEVEL), and sets them ("Subject 1").

Therefore, the first four (non-discretionary) model axioms [c.f. § 2.1] are maintained by the specifications if it can be proved that for each O-function, the level of the modified variables dominates the levels of all observed variables, for each subject. This follows from the definition of protection levels and protection dominance [2].

The representative sample of case tables in appendix I do indeed prove that this is true for every O-function except SIGNON, SIGNOFF and MOVE.

Since the three exceptions violate the axioms, each must be executed by a special process, "trusted" to perform its function correctly. The user control process (UCP) executes SIGNON and SIGNOFF in response to a human user's request. MOVE may be invoked only by the data base administrator's process.

These "trusted" processes are required to execute at the system-high protection level, and therefore maintain at least the simple security [c.f. § 2.1] and indirect "contamination" axioms.

5.2 Tranquility Principle

It is assumed that the protection levels of the subjects performing the O-function effects remain constant.

Variables in the user working area are tranquil for the following reasons:

- (i) they are completely isolated from other users; and
- (ii) their level, `K_CUR_LEVEL` (by definition), is modified only by `SIGNON` (initialized) and `SIGNOFF` (purged).

In the data base, the protection level of an object is a parameter of the access (by the structure of identifiers [c.f. § 3.1]). Each different level parameter indicates a different object is to be accessed. Additionally, the design allows no movement of identifiers from one directory to another.

Isolation of the kernel working area, and the constancy of `K_CUR_LEVEL` ensure that the following kernel entities maintain a constant protection level:

```
K_CUR_LEVEL, K_CUR_ID, K_CUR_QTA,  
K_CUR_TIME, K_OPEN, K_RESERVE,  
K_LACC, K_LX, K_LY, K_LZ.
```

However, the level of the contents of the accumulator or a temporary may be changed by various O-functions. This apparent violation of tranquility is acceptable for the following reasons:

- (i) an inspection of the case tables in appendix I will reveal that every modification of K_LACC, K_LX, K_LY or K_LZ has an associated modification of K_VACC, K_VX, K_VY or K_VZ, respectively; and
- (ii) this modification is specified by means of an assignment statement, which requires the previous contents of the accumulator to be purged upon re-assignment (by definition).

In conclusion, it is evident that the tranquility principle of the model is maintained.

5.3 Discretionary Authorization

Discretionary authorization policy requires explicit permission to have been extended before a user may access a data base object.

The primitive function which establishes such permission is KDM. An inspection of its specification [1] will prove its correctness. That is, the "NO" exception ensures that there is discretionary authorization to extend discretionary authorization. This is made possible by defining ownership of a variable (indicated in its identifier) to imply complete discretionary authorization. The "IC" and "IV" exceptions ensure that the accumulator contains the appropriate permission matrix.

The only means of transferring data base objects to the user is by means of the KWA primitive, and the Discretionary_kwa V-function checks for appropriate authorization. The Open Table [c.f. § 4.2] and Identification [c.f. § 4.3] invariants ensure that the mechanisms involved in this check function correctly.

The only means of storing data in the data base is by means of the WDV, KDM and KDV O-functions. The NO exception in each of these functions checks for discretionary authorization.

The Open Table and Identification invariants ensure that all data in the hidden kernel area is correctly subject to discretionary authorization.

Table 5.1 summarizes the discretionary authorization mechanisms in the O-functions, leading to the conclusion that the discretionary authorization model axiom is maintained.

Primitive	Discretionary Authorization	Primitive	Discretionary Authorization
1. APP_DIR	directory is exempt	DKM	NO exception
2. DEL_DIR	directory is exempt	DKQ	signon list is exempt
3. REP_DIR	directory is exempt	DKR	NO exception
4. INIT	ownership	DKV	NO exception
5. DESTROY	ownership	DKZ	NO exception
6. RES	NO exception	WKB	ownership
7. REQ	NO exception	KDM	NO exception
8. REL	RS exception	KDV	NO exception
9. SIGNON	trusted process	KDZ	ownership
10. SIGNOFF	trusted process	WDV	NO exception
11. O_APPEND	ND exception	KWA	Discretionary_kwa
12. O_DELETE	NO exception	PROJECTW	ownership
13. APPEND	hidden	SELECTW	ownership
14. ASSIGN	hidden	APPENDW	ownership
15. CONCAT	hidden	CROSS	ownership
16. EXTRACT	hidden	ARITH	ownership
17. SELECT	hidden	ASSIGNW	ownership
18. PROJECT	hidden	SIZE	ownership
19. LIST_DOWN	exemption	APFOR	ownership
20. DKD	directory is exempt	DIFF	ownership
21. DKE	NO exception	MOVE	directory is exempt
22. DKH	NO exception		

Table 5.1 Discretionary Authorization Relevant to Each O-function

APPENDIX I
O-FUNCTION CASES

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A.1.1 Introduction

This appendix contains detailed case tables [c.f. § 2.3] for certain O-function specifications. For each such O-function the following is included:

- (i) its formal specification from appendix IV of reference [1];
- (ii) flow diagrams illustrating the cases for each of the two subjects [c.f. § 3.4] performing the O-function; and
- (iii) its set of case tables. The purpose of each case table is to prove that the level of each modified variable dominates the levels of all variables observed in that case.

The security related primitives are the only ones requiring certification, and these may be categorized according to analogous case tables. The case tables for a representative sample of O-functions is included in this appendix. The sample consists of at least one O-function from each category (subsection A.1.3).

Three primitive functions break some of the rules of the "strict" protection policy, but they are essential for computerized data management. These are SIGNON, SIGNOFF and MOVE, and their case tables are found in subsection A.1.4.

A.1.2 Notation Used in Case Tables

The symbols and mnemonics used in the case tables are taken from the formal specifications, and are described in appendices III and IV of reference [1].

Additionally, the following abbreviations are used:

<u>abbreviation</u>	<u>meaning</u>
W	- User's current protection level (K_CUR_LEVEL)
LA	- Level of the kernel accumulator
LX	- Level of a kernel temporary
$L_1 \lceil L_2$	- The dominant level of L_1 and L_2
$L_1 \lfloor L_2$	- The dominated level of L_1 and L_2
id[LEV]	- The level component of the identifier
LEV	- An abbreviation for id[LEV]
ACC	- Kernel accumulator
UCP	- User Controller Process

A.1.3 The Sample Validations

The set of all security-related O-functions are categorized below according to analogous case tables. Those whose case tables are found in this section are so indicated.

<u>Category</u>	<u>Case Tables Included</u>	<u>Not Included</u>
1. Directory manipulation	APP_DIR	DEL_DIR, REP_DIR
2. Object existence	INIT	DESTROY
3. Object reservation	REQ	RES, REL
4. Access authorization	O_APPEND	O_DELETE
5. Accumulator manipulation	APPEND, CONCAT	ASSIGN, EXTRACT, SELECT, PROJECT, LIST_DOWN
6. Transfer to accumulator	DKD, DKE	DKH, DKM, DKQ, DKR, DKV, DKZ
7. Transfer to data base	KDM, KDV, KDZ	WDV
8. Working area	KWA	WKB
9. Access DMS	SIGNON, SIGNOFF	
10. Data base administrator	MOVE	

Table A.1.1 Table of O-function Categories

A.1.3.1 APP_DIR

(A) O-function APP_DIR(lv,n,t,lz)

```
* Append a tuple to a data base directory, where:
* (i) lz = 0 if it's an object's defining entry;
* (ii) lz ≠ 0 if it's an object registration at "lv".
```

parameter types

level lv, name n, type t, level_zero lz

exception * Check legality of parameter lz first. *

```
IR: 1 FALSE IF lz = 0 ELSE (lz = lv) v (lz ≠ lv)
IL: 2 lv ≠ K_CUR_LEVEL
*DD: (K_CUR_ID, n, t, *) ∈ D_D(lv) †
```

* Zero or strictly dominating
* Append must be a write-up
* Entry is there already

effect * Note that effect [2] gives the semantics of "*DD". *

```
[1] D D(lv) ← † ∪ (K_CUR_ID, n, t, lz)
*[2] W_CODE ← DN IF K_CUR_LEVEL = lv
* Append the entry
* Avoid a "write-down" *
```

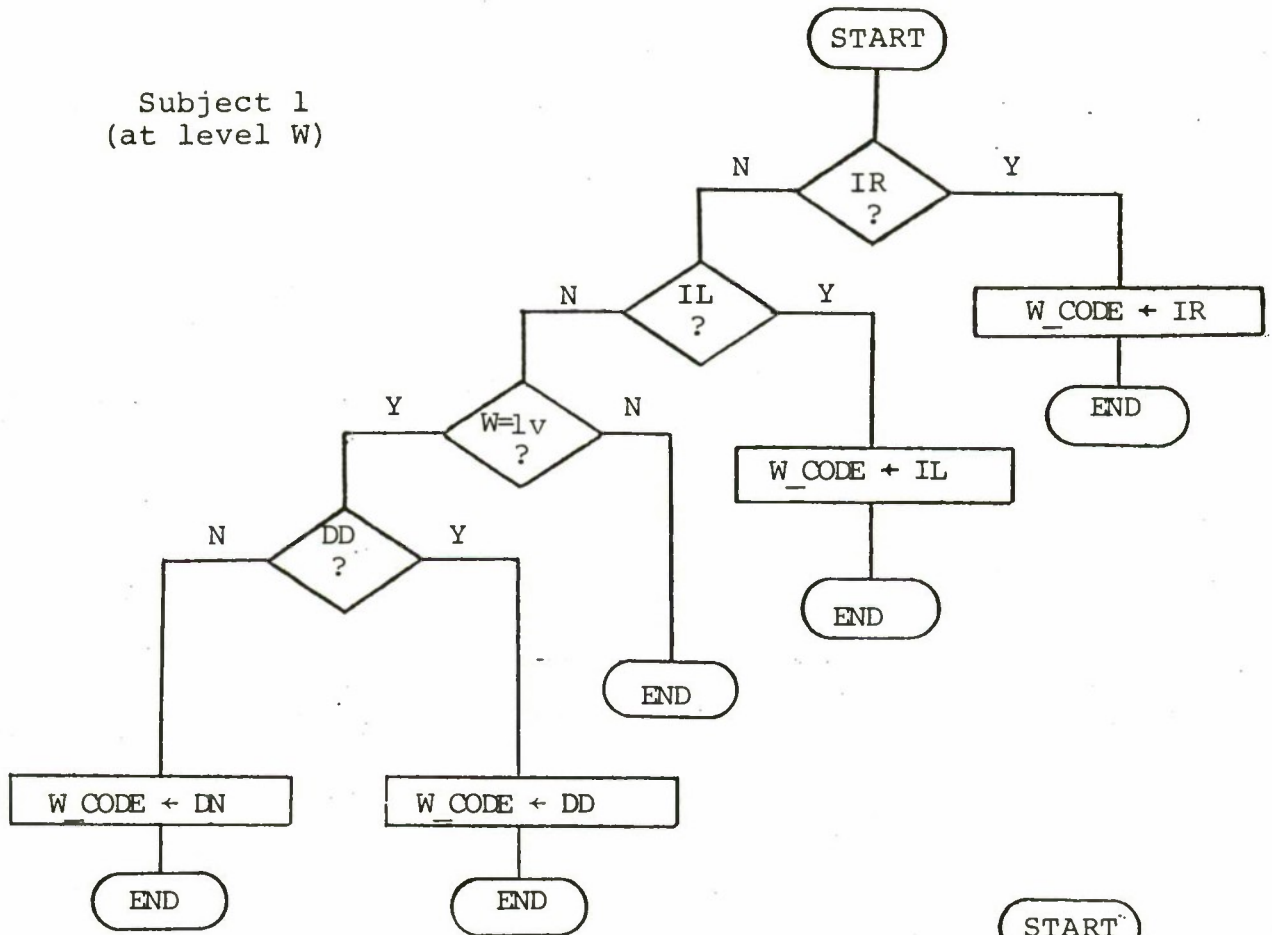
(B) Access table

Variables Observed	Variables Modified	Variables Observed and Modified
K_CUR_ID	W_CODE	D_D(lv)
K_CUR_LEVEL		

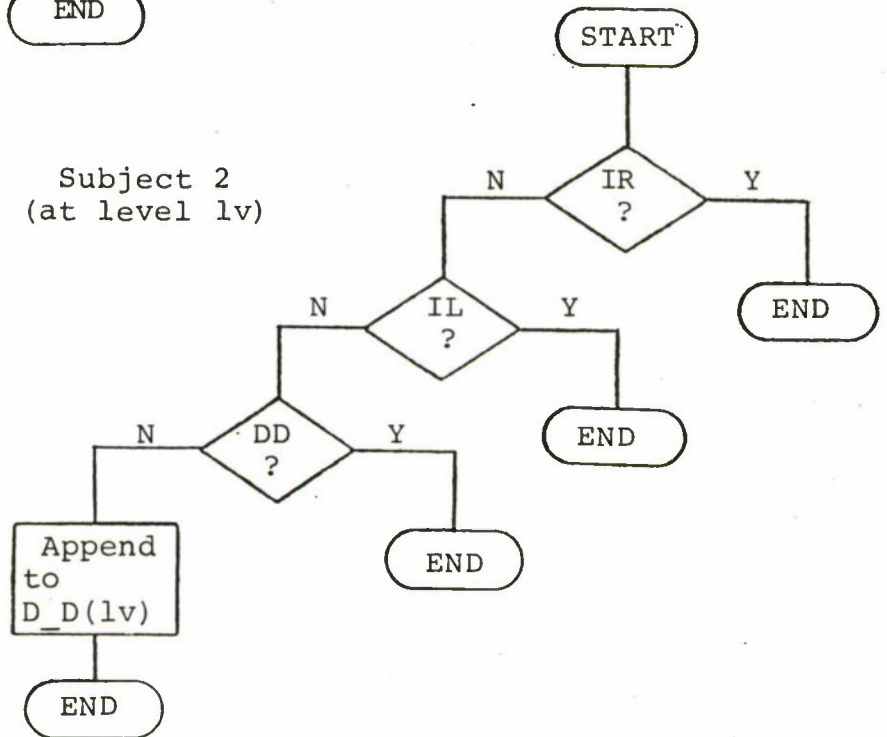
¹Exception statements containing "IF" are of the form: E₁ IF E₂ ELSE E₃, where E_i are logical expressions with TRUE/FALSE values. The semantics² of such a statement is: exception is value of E₁ if E₂ evaluates to TRUE, else exception is value of E₃.

² The asterisk by DD indicates that code DD is returned ONLY if K_CUR_LEVEL = lv. The asterisk in the directory tuple indicates that any value found in the LEVEL domain will satisfy the membership condition.

Subject 1
(at level W)



Subject 2
(at level lv)



PRIMITIVE: APP_DIR

CASE: 1

SUBJECT: 1

CONDITIONS: IR

Register attempted at level not strictly dominated by that of the definition entry.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> lv, n, t, lz	W	W_CODE	W
<u>CONSTANTS:</u> O, IR	Unclass		
<u>VARIABLES:</u>			
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: APP_DIR

CASE: 2

SUBJECT: 1

CONDITIONS: (~IR) ^ IL

Directory level does not strictly dominate the user's current level.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> lv, n, t, lz	W	W_CODE	W
<u>CONSTANTS:</u> O, IL	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL			
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: APP_DIR

CASE: 3

SUBJECT: 1

CONDITIONS: (\sim IR) \wedge (\sim IL) \wedge (W=lv) \wedge DD
The directory entry already exists.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> lv, n, t, lz	W	W_CODE	W
<u>CONSTANTS:</u> O, DD	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL D_D(lv) K_CUR_ID	W lv=W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: APP_DIR

CASE: 4

SUBJECT: 1

CONDITIONS: (\sim IR) \wedge (\sim IL) \wedge (W=lv) \wedge (\sim DD)
No exceptions.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> lv, n, t, lz	W	W_CODE	W
<u>CONSTANTS:</u> O, DN	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL D_D(lv) K_CUR_ID	W lv=W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: APP_DIR

CASE 1

SUBJECT: 2

CONDITIONS: (~IR) ^ (~IL) ^ (~DD)
No exceptions

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> lv, n, t, lz	W	D_D (lv)	lv
<u>CONSTANTS:</u> O, DN	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_CUR_ID D_D (lv)	W W lv		
HIGHEST LEVEL OBSERVED:	lv	LOWEST LEVEL MODIFIED:	lv

LEMMA: lv \neq W

PROOF: ~IL

A.1.3.2 INIT

(A) O-function INIT(n,t,lv,s)

* Initialize all component entities of an object. *
 * If parameter "size" is negative, and the session *
 * quota is less than the absolute value of the *
 * size, then the maximum size of the object is set *
 * equal to the current session quota. *

parameter types

name n, type t, level lv, size s

abbreviation

id = K_CUR_ID,n,t,lv

* Identifier of new Object *

exception

IL: lv ≠ K_CUR_LEVEL *
 *DE: (K_CUR_ID,n,t,ZERO) ≠ D_D(lv) *
 *DD: |D_E(id)| ≠ ∅ *
 *SZ: s > K_CUR_QTA *
 * Note: negative size case is included here *

* Illegal level for modification *
 * Directory entry is required *
 * It exists already *
 * Size is too large *
 * Size is too large *

effect

[1] D E(id) ← ZERO *
 [2] D_H(id) ← (K_CUR_TIME,K_CUR_ID,K_CUR_TIME) * Initialize current size *
 [3] D_Z(id) ← Size_init(s) IF K_CUR_LEVEL = lv ELSE ZERO * Initialize object history *
 [4] K_CUR_QTA ← |H - Size_init(s) IF K_CUR_LEVEL = lv * Use QUOTA only at *
 *[5] W_CODE ← DN IF K_CUR_LEVEL = lv * user's current level *

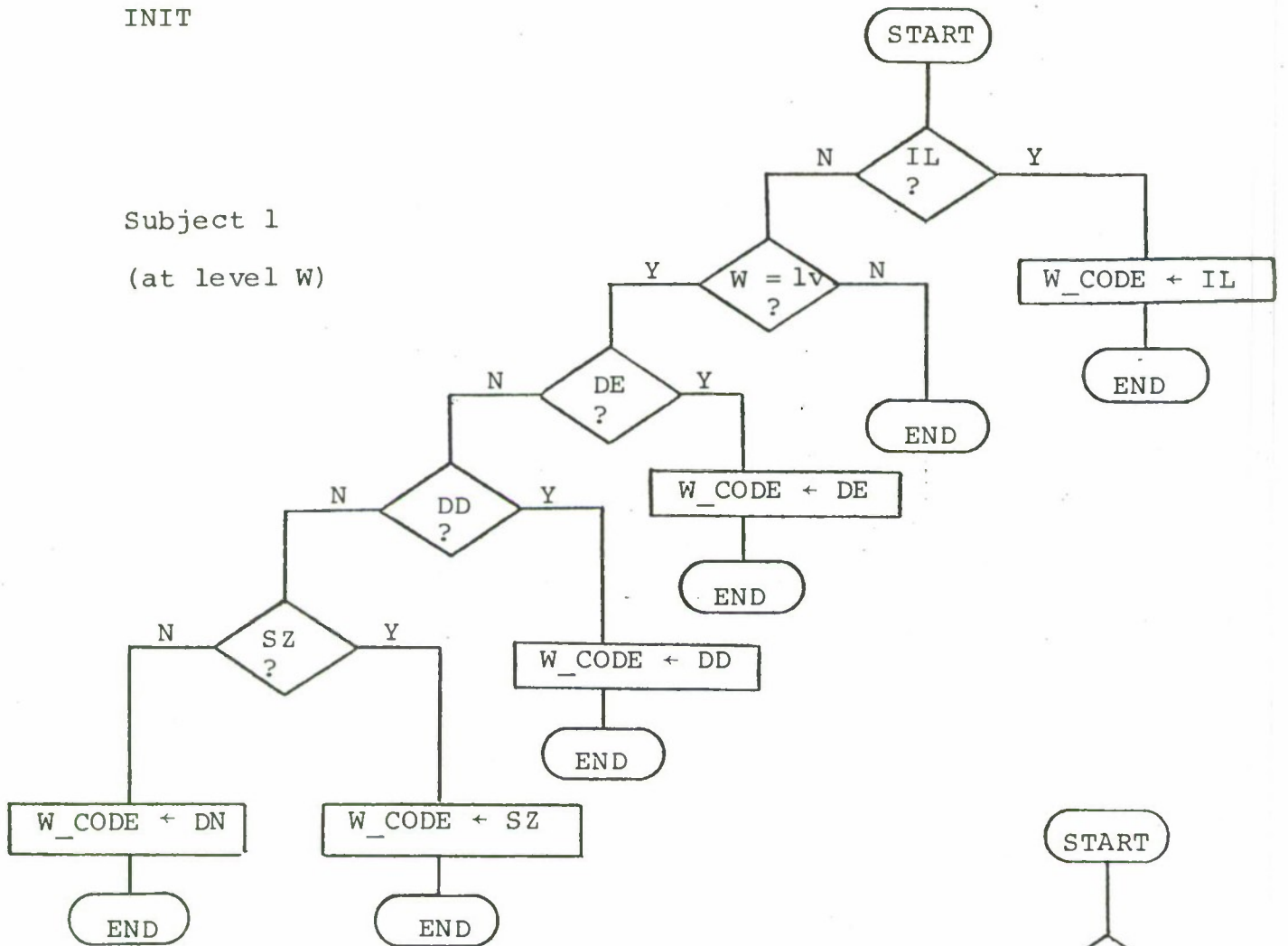
(B) Access table

Variables Observed	Variables Modified	Variables Observed and Modified
K_CUR_LEVEL		
D_D(lv)		
K_CUR_ID		
K_CUR_TIME		
D_Z(id)		
D_H(id)		
W_CODE		
		D E(id)
		K_CUR_QTA

INIT

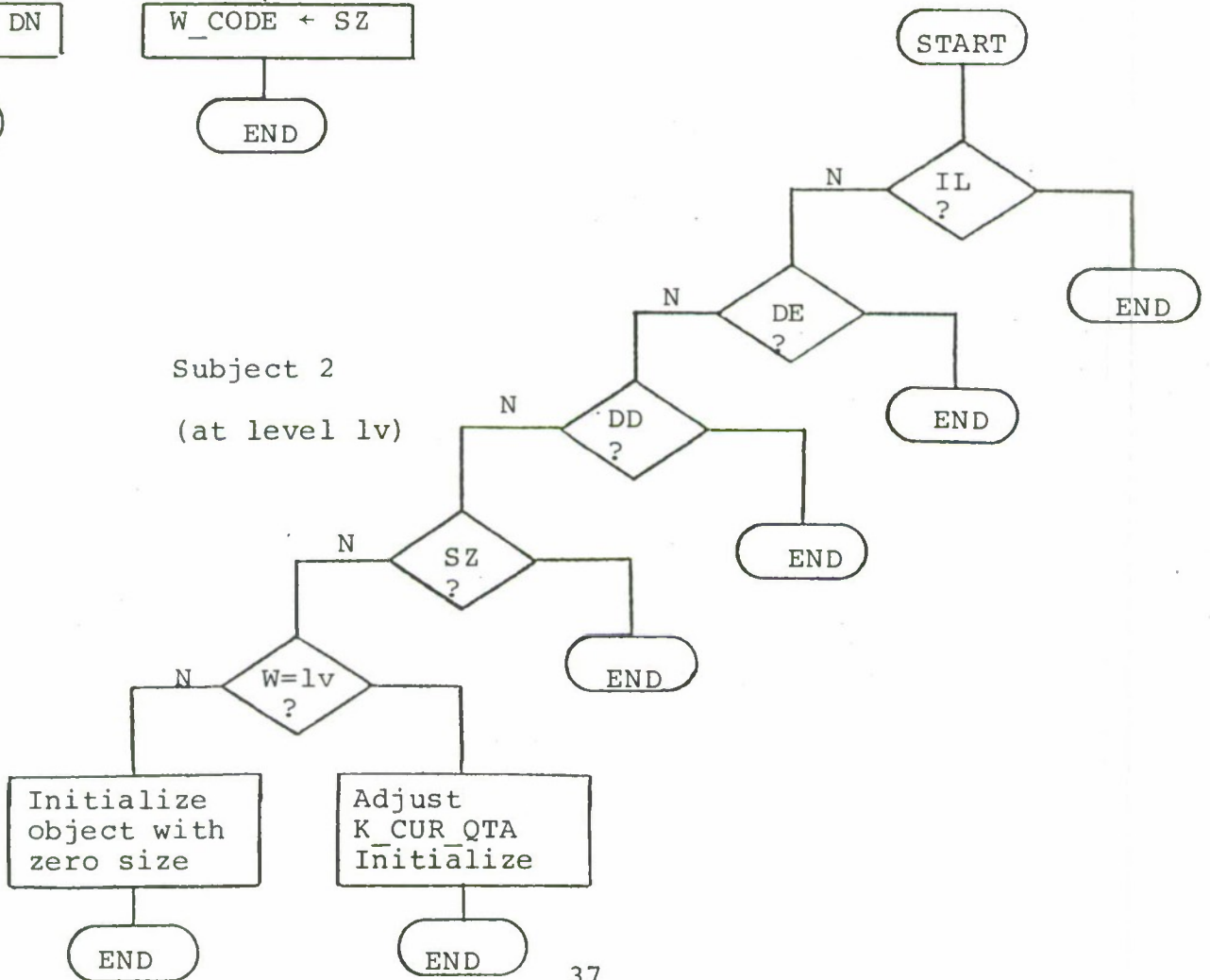
Subject 1

(at level W)



Subject 2

(at level lv)



PRIMITIVE: INIT

CASE: 1

SUBJECT: 1

CONDITIONS: IL

Level of object to be initialized does not dominate user's current level.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> n,t,lv,s	W	W_CODE	W
<u>CONSTANTS:</u> IL	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL	W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: INIT

CASE: 2

SUBJECT: 1

CONDITIONS: (~IL) ^ (W=lv) ^ DE

Object is not defined in the directory, and the object level equals the user's current level.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> n,t,lv,s	W	W_CODE	W
<u>CONSTANTS:</u> ZERO, DE	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_CUR_ID D_D(lv)	W W lv=W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: INIT

CASE: 3

SUBJECT: 1

CONDITIONS: (\sim IL) \wedge (W=lv) \wedge (\sim DE) \wedge DD
The object has been previously initialized

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> n, t, lv, s	W	W_CODE	W
<u>CONSTANTS:</u> ZERO, DD	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_CUR_ID D_D(lv) D-E(id)	W W lv=W lv=W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: INIT

CASE: 4

SUBJECT: 1

CONDITIONS: (\sim IL) \wedge (W=lv) \wedge (\sim DE) \wedge (\sim DD) \wedge SZ
The requested size exceeds user's current quota

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> n, t, lv, s	W	W_CODE	W
<u>CONSTANTS:</u> ZERO, SZ	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_CUR_ID D_D(lv) D_E(id) K_CUR_QTA	W W lv=W lv=W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: INIT

CASE: 5

SUBJECT: 1

CONDITIONS: (~IL) ^ (W=lv) ^ (~DE) ^ (~DD) ^ (~SZ)
No exceptions.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> n,t,lv,s	W	W_CODE	W
<u>CONSTANTS:</u> ZERO, Ø, DN	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_CUR_ID D_D(lv) D_E(id) K_CUR_QTA K_CUR_TIME	W W lv=W lv=W W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: INIT

CASE: 6

SUBJECT: 1

CONDITIONS: (~IL) ^ (W≠lv)
The object is initialized at a strictly dominating level.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> n,t,lv,s	W		W ³
<u>CONSTANTS:</u>	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL	W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

³Note that this is the "null" modification, that is, W_CODE is set in such a way that it signals that "lv strictly dominates K_CUR_LEVEL", by ~IL.

PRIMITIVE: INIT

CASE: 1

SUBJECT: 2

CONDITIONS: (~IL) ^ (~DE) ^ (~DD) ^ (~SZ) ^ (W=lv)
No exceptions. Initialize object at user's current level.

	OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u>	n, t, lv, s	W	D_E(id) D_H(id) D_Z(id)	lv=W lv=W lv=W
<u>CONSTANTS:</u>	ZERO, Ø, DN	Unclass	K_CUR_QTA	W
<u>VARIABLES:</u>	K_CUR_LEVEL K_CUR_ID D_D(lv) D_E(id) K_CUR_QTA K_CUR_TIME	W W lv=W lv=W W W		
<u>HIGHEST LEVEL OBSERVED:</u>		W	<u>LOWEST LEVEL MODIFIED:</u>	W

PRIMITIVE: INIT

CASE: 2

SUBJECT: 2

CONDITIONS: (~IL) ^ (~DE) ^ (~DD) ^ (~SZ) ^ (W≠lv)
No exceptions. Initialize object at level strictly dominating user's current level, giving it zero size.

	OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u>	n, t, lv, s	W	D_E(id) D_H(id) D_Z(id)	lv lv lv
<u>CONSTANTS:</u>	ZERO, Ø, DN	Unclass		
<u>VARIABLES:</u>	K_CUR_LEVEL K_CUR_ID D_D(lv) D_E(lv) K_CUR_QTA K_CUR_TIME	W W lv lv W W		
<u>HIGHEST LEVEL OBSERVED:</u>		lv	<u>LOWEST LEVEL MODIFIED:</u>	lv

Lemma: lv ∞ W
Proof: ~IL

A.1.3.3

(A) O-function REQ(o,n,t)

* Reserve an object (at the current level) if it's *
 * available. If not, wait until it is available *

parameter types

user o, name n, type t

abbreviation

id = o,n,t,K_CUR_LEVEL

exception

NO: K OPEN(id,RSRV)⁴ = FALSE
 RS: id ∈ {K RESERVE-}
 DL: (¬D_R(id) ∨ ≠ ∅) ∧ (¬K RESERVE- ≠ ∅)
 * If current process must wait while holding objects, dead-lock could result. *

* Identifier of object to reserve *

* The object has not been opened *
 * I've already reserved it *
 * Dead-lock hazard *

* Wait until object is not reserved. *
 * This specification statement is to be performed AFTER *
 * all exceptions have been tested for, and BEFORE the *
 * effects take place. *

effect

[1] D_R(id) ← K_CUR_ID
 [2] K_RESERVE ← { } ∪ (id)
 [3] W_CODE ← DN

* Reserve the object *
 * Append an entry to reserve table *

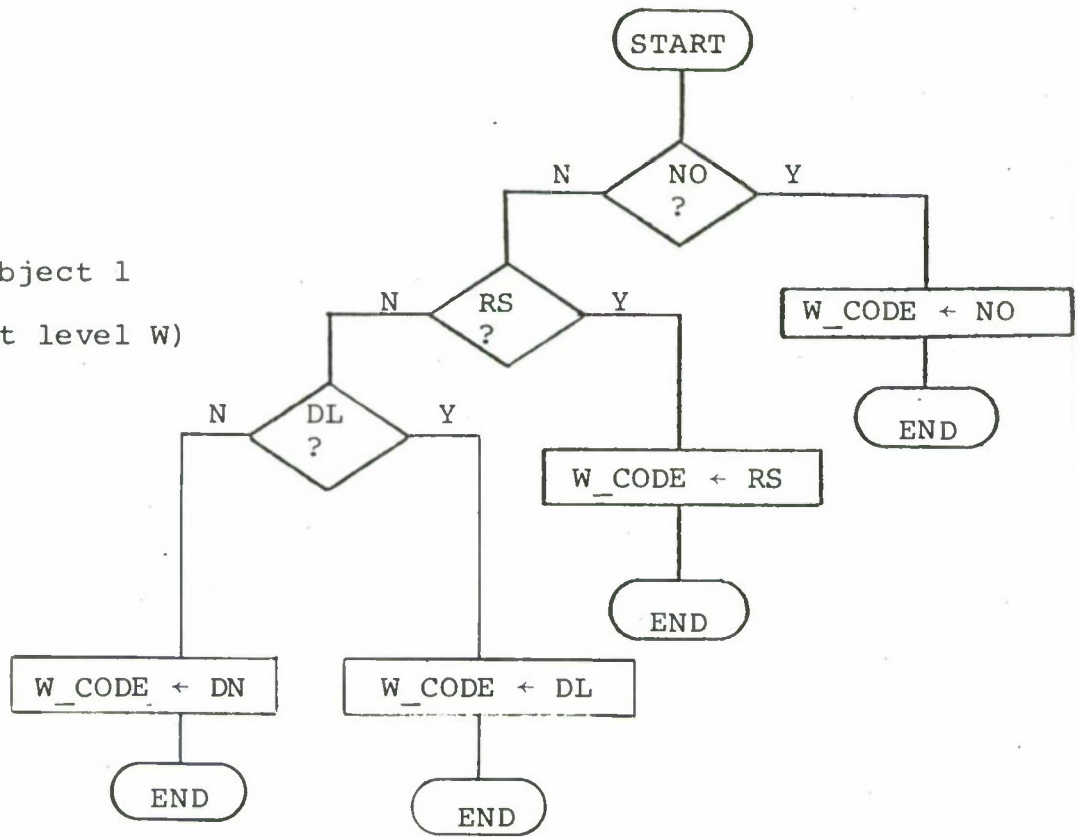
(B) Access table

Variables Observed	Variables Modified	Variables Observed and Modified
K_OPEN	W_CODE	K_RESERVE
K_CUR_LEVEL		D_R(id)
K_CUR_ID		

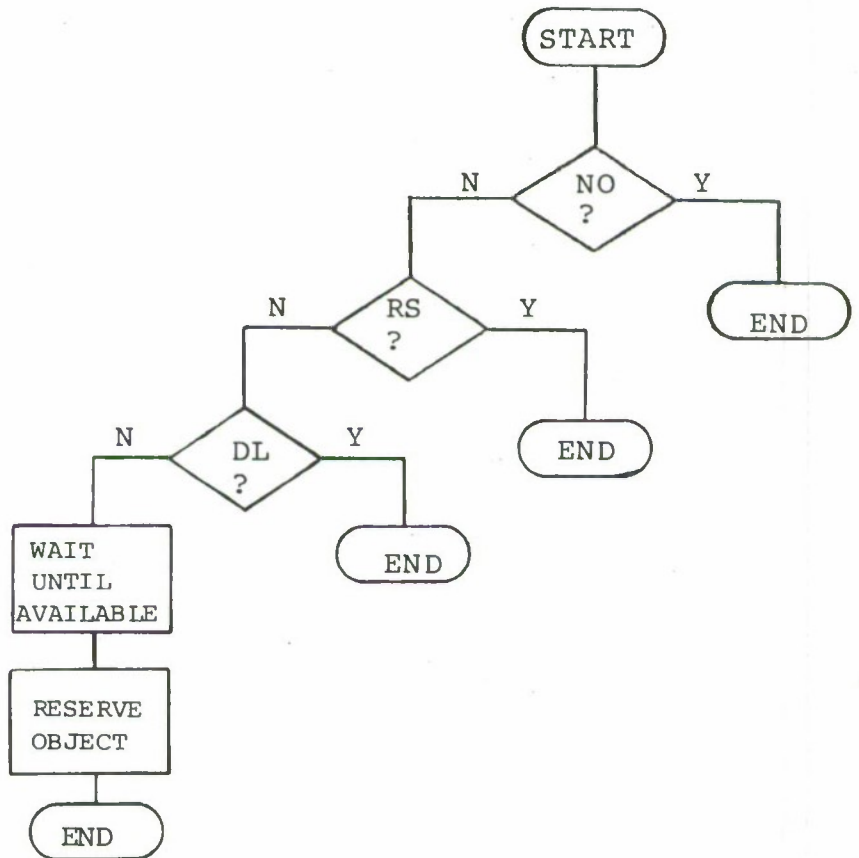
⁴The open table is a five-dimensional array, with the "level" dimension assuming a lattice structure.

REQ

Subject 1
(at level W)



Subject 2
(at level W)



PRIMITIVE: REQ

CASE: 1

SUBJECT: 1

CONDITIONS: NO

Object is not open.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS</u> : o,n,t	W	W_CODE	W
<u>CONSTANTS</u> : RSRV, NO	Unclass		
<u>VARIABLES</u> : K_CUR_LEVEL K_OPEN	W W ⁵		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: REQ

CASE: 2

SUBJECT: 1

CONDITIONS: (~NO) ^ RS

User already has object reserved.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS</u> : o,n,t	W	W_CODE	W
<u>CONSTANTS</u> : RSRV, RS	Unclass		
<u>VARIABLES</u> : K_CUR_LEVEL K_OPEN K_RESERVE	W ⁵ W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

⁵ Note that reservation is restricted to a user's current level.

PRIMITIVE: REQ

CASE: 3

SUBJECT: 1

CONDITIONS: (\sim NO) \wedge (\sim RS) \wedge DL

Object is reserved by another user, and to queue for it would risk deadlock.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> o, n, t	W	W_CODE	W
<u>CONSTANTS:</u> RSRV, \emptyset , DL	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_OPEN K_RESERVE D_R(id)	W W W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: REQ

CASE: 4

SUBJECT: 1

CONDITIONS: (\sim NO) \wedge (\sim RS) \wedge (\sim DL)
No exceptions

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> o, n, t	W	W_CODE	W
<u>CONSTANTS:</u> RSRV, \emptyset , DN	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_OPEN K_RESERVE D_R(id)	W W W W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: REQ

CASE 1

SUBJECT: 2

CONDITIONS: (~NO) ^ (~RS) ^ (~DL)

No exceptions. Reserve object on behalf of user.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> o, n, t	W	K_RESERVE	W
<u>CONSTANTS:</u> RSRV, Ø, DN	Unclass	D_R(id)	W
<u>VARIABLES:</u> K_CUR_LEVEL K_OPEN K_RESERVE D_R(id) K_CUR_ID	W W W W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

A.1.3.4 O_APPEND

(A) O-function O_APPEND(id)

- * This primitive opens an object for all possible access by W, *
- * by appending a tuple to K_OPEN for every authorized access. *
- * An open is required to: (i) allow object access; and *
- * (ii) prohibit the purging of an accessed object. *

parameter types

identifier id (OWN,NAM,TYP,LEV)⁶

exception

IL: (K CUR LEVEL \neq LEV) \wedge (LEV \neq K_CUR_LEVEL) * Incomparable level *
 DO: K_OPEN(id,) = TRUE * Object is open already *
 *NE: D_E(id) = \emptyset * Object does not exist *
 ND: (K_CUR_ID \neq OWN) \wedge ((K_CUR_ID,) \notin D_M(id)) *
 * There is no discretionary authorization whatsoever. *

effect

- [1] K_OPEN(Access_set O(id)) = TRUE * Set each authorized access to TRUE *
- [2] D_O(id) \leftarrow \vdash \vdash K_CUR_ID IF Opened O(id) \wedge (LEV \neq K_CUR_LEVEL)
- *[3] W_CODE \leftarrow DN IF Opened O(id) \wedge (K_CUR_LEVEL \neq LEV) ELSE ND IF K_CUR_LEVEL \neq LEV
- * Note that the open list, D_O(id), is NOT modified for strictly dominated objects. *

(B) Access table

Variables Observed	Variables Modified	Variables Observed and Modified
K_CUR_LEVEL	W_CODE	K_OPEN
K_CUR_ID		D_O(id)
D_E(id)		
D_M(id)		

⁶ Type "id" has previously been defined, so "(OWN,NAM,TYP,LEV)" should appear in a comment, instead of in the parameter declaration.

⁷ "LEV" is an abbreviation for "id[LEV]".

(C) O_APPEND V-functions

(i) V-function Access_set_0(id) : open

* Return a set of open table tuples; one for each authorized access mode. *
range

ZERO to NUM_MODES of (identifier,access)

parameters

identifier id (OWN,NAM,TYP,LEV)

derivation

(id,APCY) IF Auth_0(LEV,K_CUR_LEVEL,id,APCY) * Authorize each of the *
u(id,EXPM) IF Auth_0(LEV,K_CUR_LEVEL,id,EXPM) * eight modes of access *
u(id,RDHS) IF Auth_0(K_CUR_LEVEL,LEV,id,RDHS)
u(id,RDPM) IF Auth_0(K_CUR_LEVEL,LEV,id,RDPM)
u(id,RDSZ) IF Auth_0(K_CUR_LEVEL,LEV,id,RDSZ)
u(id,RETR) IF Auth_0(K_CUR_LEVEL,LEV,id,RETR)
u(id,RSRV) IF Auth_0(LEV,K_CUR_LEVEL,id,RSRV) ^ (LEV = K_CUR_LEVEL)
u(id,STOR) IF Auth_0(LEV,K_CUR_LEVEL,id,STOR)

(ii) V-function Auth_0(lv₁,lv₂,id,acc) : boolean

* Check non-discretionary and discretionary access authorization. *

range

TRUE,FALSE

parameter types

level lv₁, level lv₂, identifier id(OWN,NAM,TYP,LEV), access acc

derivation

(lv₁ > lv₂) ^ ((OWN = K_CUR ID) v ((K_CUR ID,acc) ∈ D M(id)))
* non-discretionary AND_ownership OR_discretionary authorization. *

(iii) V-function Opened_O(id) : boolean

* Return TRUE if any access mode is authorized. *

range

TRUE, FALSE

parameter types

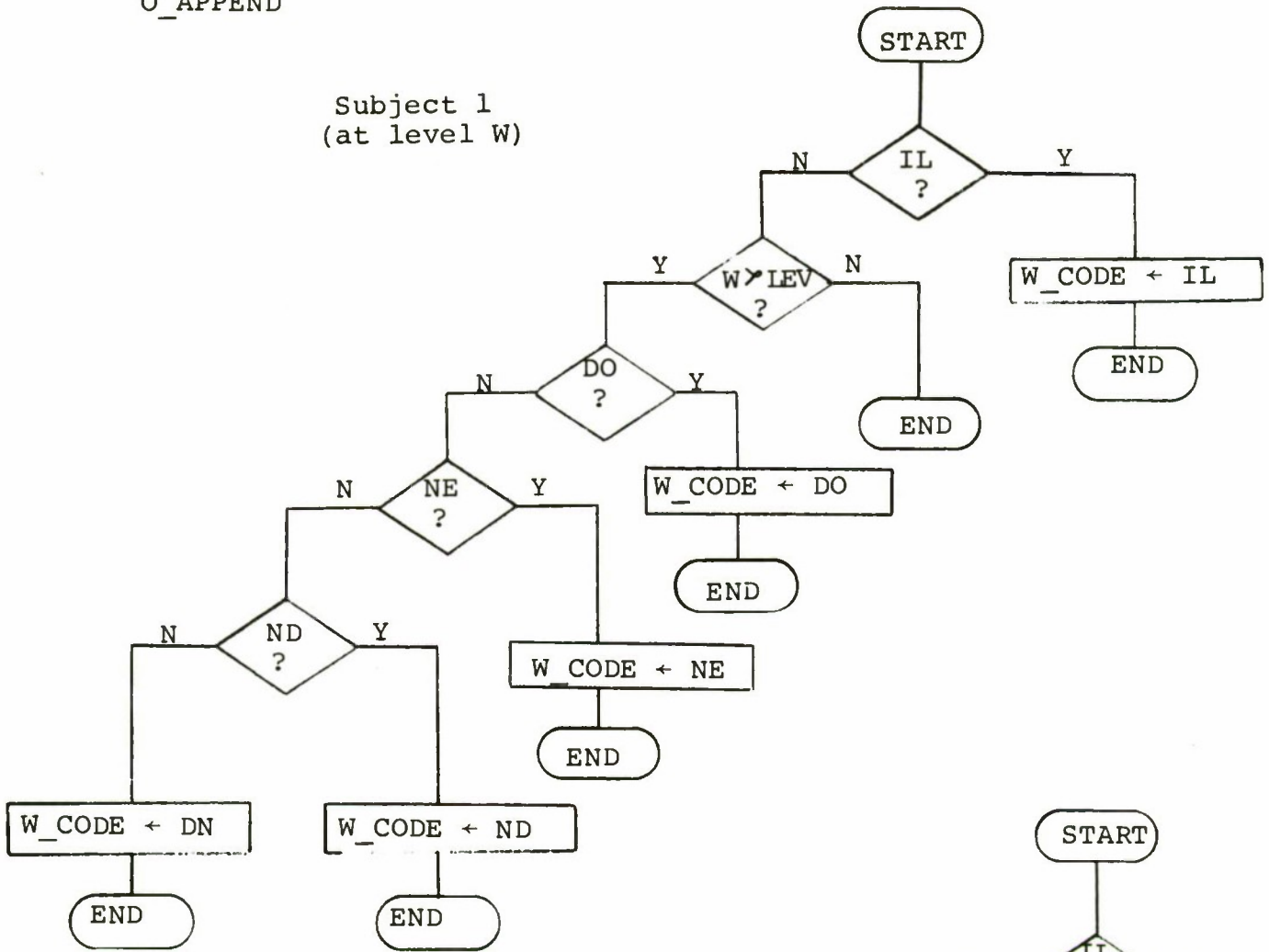
identififier id (OWN, NAM, TYP, LEV)

derivation

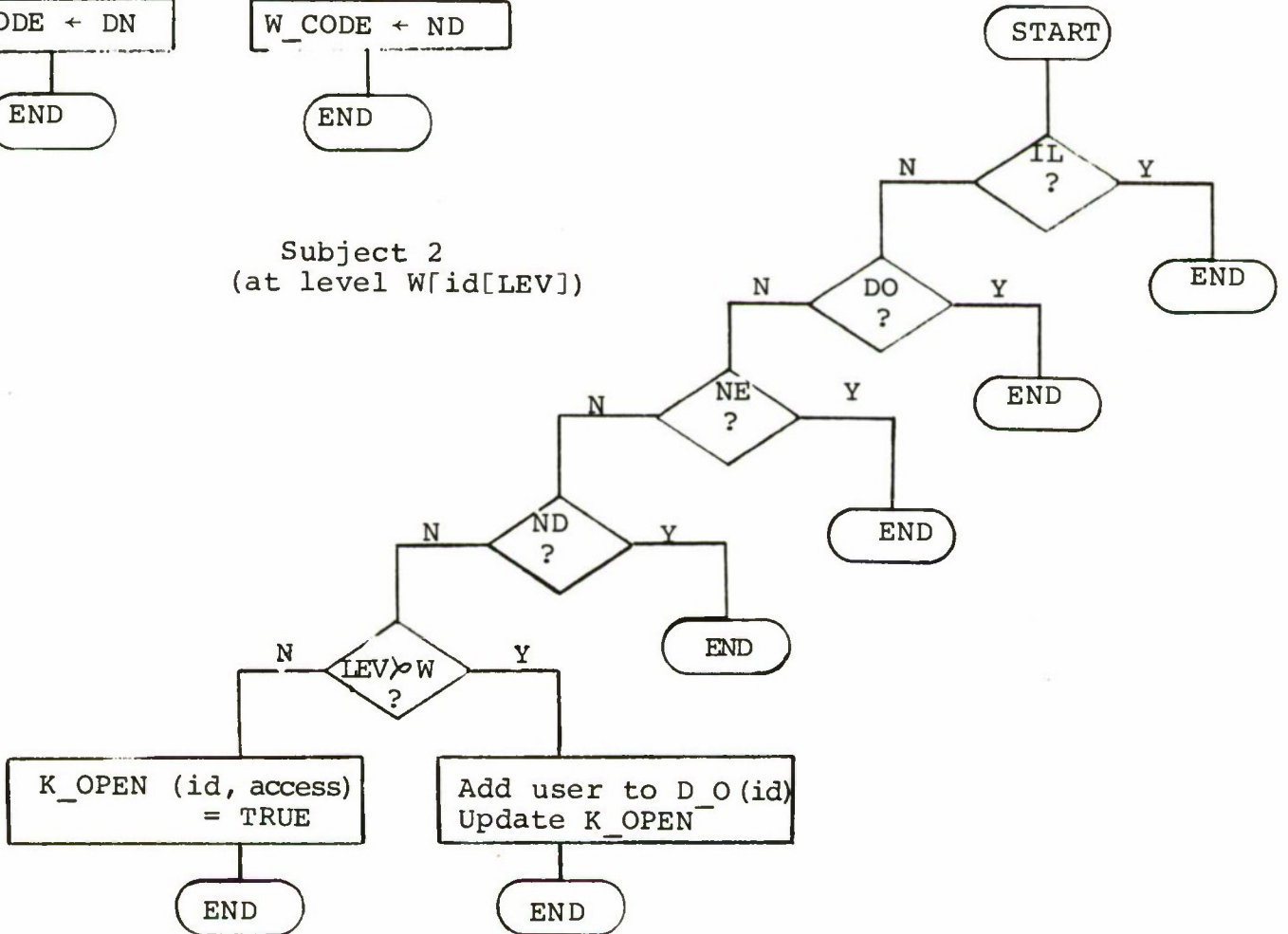
Auth_O(LEV, K CUR LEVEL, id, APCY)
vAuth_O(LEV, K_CUR_LEVEL, id, EXPM)
vAuth_O(K CUR_LEVEL, LEV, id, RDHS)
vAuth_O(K_CUR_LEVEL, LEV, id, RDPM)
vAuth_O(K_CUR_LEVEL, LEV, id, RDSZ)
vAuth_O(K_CUR_LEVEL, LEV, id, RETR)
vAuth_O(LEV, K_CUR_LEVEL, id, RSRV) ^ (LEV = K_CUR_LEVEL)
vAuth_O(LEV, K_CUR_LEVEL, id, STOR)

O_APPEND

Subject 1
(at level W)



Subject 2
(at level W[id[LEV]])



PRIMITIVE: O_APPEND

CASE: 1

SUBJECT: 1

CONDITIONS: IL

User's current level and object level are incomparable.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS</u> : id	W	W_CODE	W
<u>CONSTANTS</u> : IL	Unclass		
<u>VARIABLES</u> : K_CUR_LEVEL	W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: O_APPEND

CASE: 2

SUBJECT: 1

CONDITIONS: (~IL) ^ (W ≠ LEV) ^ DO
Object is open already.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS</u> : id	W	W_CODE	W
<u>CONSTANTS</u> : DO	Unclass		
<u>VARIABLES</u> : K_CUR_LEVEL K_OPEN[LEV]	W LEV		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: O_APPEND

CASE: 3

SUBJECT: 1

CONDITIONS: (~IL) ^ (W > LEV) ^ (~DO) ^ NE
Non-existent object

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	W_CODE	W
<u>CONSTANTS:</u> NE	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_OPEN[LEV] D_E(id)	W LEV LEV		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: O_APPEND

CASE: 4

SUBJECT: 1

CONDITIONS: (~IL) ^ (W > LEV) ^ (~DO) ^ (~NE) ^ ND
User has no discretionary authorization.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	W_CODE	W
<u>CONSTANTS:</u> ND	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_OPEN[LEV] D_E(id) K_CUR_ID D_M(id)	W LEV LEV W LEV		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: O_APPEND

CASE: 5

SUBJECT: 1

CONDITIONS: (~IL) ^ (W > LEV) ^ (~DO) ^ (~NE) ^ (~ND)
No exceptions.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	W_CODE	W
<u>CONSTANTS:</u> APCY, EXPM, RDHS, RDPM, RDSZ, RETR, RSRV, STOR, DN	Unclass		
<u>VARIABLES:</u> K_CUR LEVEL K_OPEN[LEV] D_E(id) K_CUR ID D_M(id)	W LEV LEV W LEV		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: O_APPEND

CASE: 1

SUBJECT: 2

CONDITIONS: (~IL) ^ (~DO) ^ (~NE) ^ (~ND) ^ (LEV < W)
No exceptions for a strictly dominated object.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	K_OPEN	W
<u>CONSTANTS:</u> APCY, EXPM, RDHS, RDPM, RDSZ, RETR, RSRV, STOR, DN	Unclass		
<u>VARIABLES:</u> K_CUR LEVEL K_OPEN[LEV] D_E(id) K_CUR ID D_M(id)	W LEV LEV W LEV		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: O_APPEND

CASE 2

SUBJECT: 2

CONDITIONS: (~IL) ^ (~DO) ^ (~NE) ^ (~ND) ^ (LEV > W)
No exceptions. Level of object being opened dominates user's current level.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	D_O(id)	LEV
<u>CONSTANTS:</u> APCY, EXPM, RDHS, RDPM, RDSZ, RETR, RSRV, STOR, DN	Unclass	K_OPEN[LEV]	LEV
<u>VARIABLES:</u> K_CUR LEVEL K_OPEN[LEV] D_E(id) K_CUR ID D_M(id)	W LEV LEV W LEV		
HIGHEST LEVEL OBSERVED:	LEV	LOWEST LEVEL MODIFIED:	LEV

A.1.3.5 APPEND

(A) O-function APPEND(xt)

```
* The contents of "xt" are appended to the accumulator contents, where xt is: *
* (i) a kernel temporary: X; Y; or Z; *
* or (ii) a tuple of values from the working area. *
* The APPEND fails if the primary key uniqueness property of the relation in *
* the accumulator would be destroyed. This function differs from WW3 mainly *
* in that it is hidden, and only the accumulator is modified. *
```

parameter types

temp_tup xt, contents (O,N,T,L,C) * Data in contents regs *

abbreviation

x = xt ∈ {X,Y,Z} * x assumes TRUE/FALSE values *

exception

```
*IL: K_LACC < K_Lxt IF x * Accumulator must dominate *
*IT: K_IACC[T] = 'S' * Don't do this for strings *
*IC: K_FACC[DTYPE] ≠ K_Fxt[DTYPE] IF x ELSE ~(xt CONFORMS TO K_FACC) *
*IV: ~Unique_keys(K_FACC, K_Vxt) IF x *
      ~Unique_keys(K_VACC, K_Vxt) IF x *
* Check that the primary key uniqueness property is maintained. *
```

effects

```
[1] K_VACC ← ⊥ ∪ K_Vxt IF x ELSE ⊥ ∪ xt
* Append the temporary register or a tuple of parameters. *
*[2] W_CODE ← DN IF K_CUR_LEVEL ≠ K_LACC * Return code if possible *
```

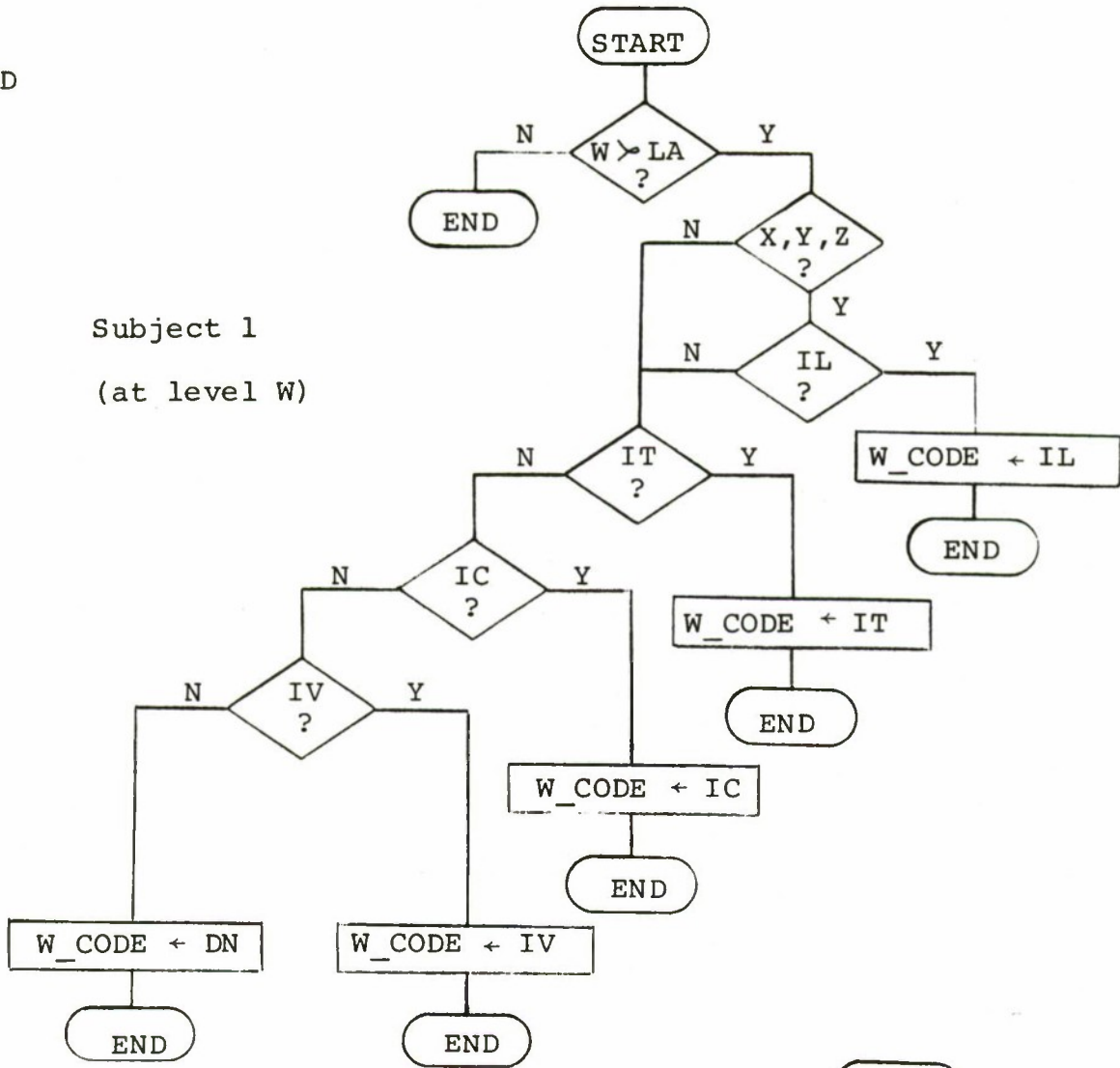
(B) Access table

Variables Observed	Variables Modified	Variables Observed and Modified
K_IACC, K_Ixt	W_CODE	K_VACC
K_FACC, K_Fxt		
K_Vxt, K_LACC, K_Lx		

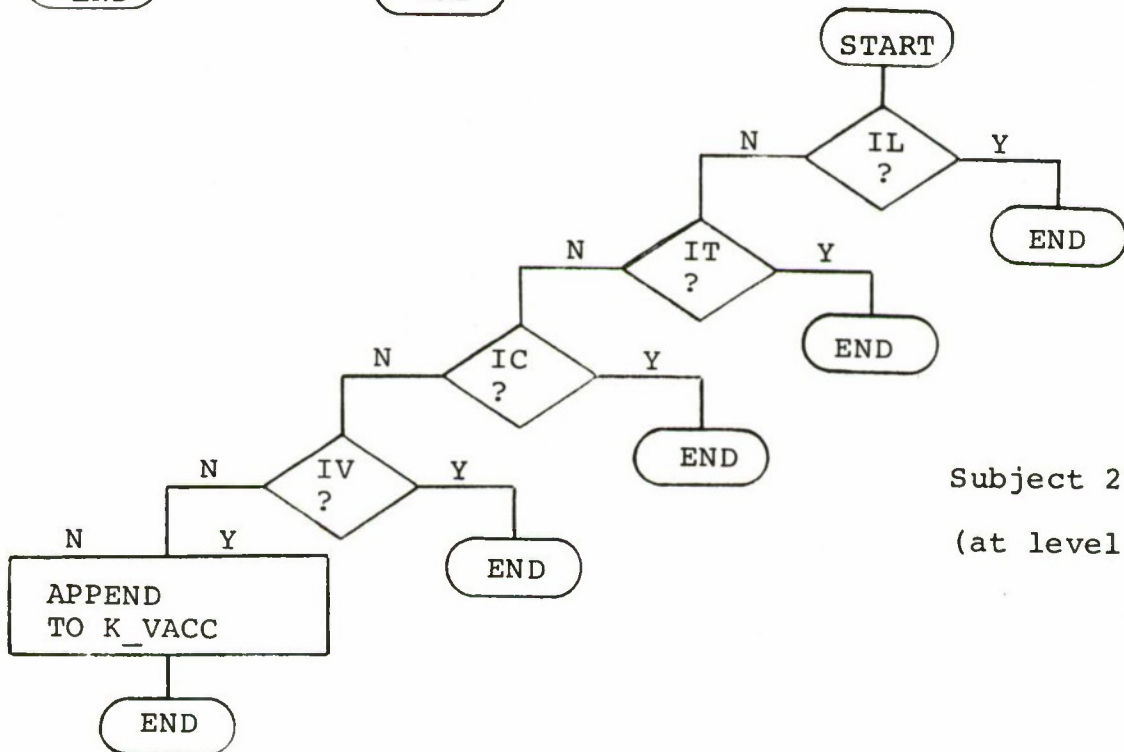
8 The level of data in these boolean expressions is the level of the accumulator contents (not a parameter). This is required for setting W_CODE (for '*').

APPEND

Subject 1
(at level W)



Subject 2
(at level LA)



PRIMITIVE: APPEND

CASE: 1

SUBJECT: 1

CONDITIONS: $(xt) \in \{X,Y,Z\} \wedge (W \succ LA) \wedge IL$
xt is a kernel temporary. However, its level is not dominated by the accumulator level.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS</u> : xt(name of temporary)	W	W_CODE	W
<u>CONSTANTS</u> : IL	Unclass		
<u>VARIABLES</u> : K_CUR_LEVEL K_LACC K_Lxt	W W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: LA = W

PROOF: $W \succ LA$ and Minimum K-level Invariant [c.f. § 4.1].

PRIMITIVE: APPEND

CASE: 2A

SUBJECT: 1

CONDITIONS: $(xt \in \{X,Y,Z\}) \wedge (W = LA) \wedge (\sim IL) \wedge IT$
APPEND cannot be used with strings.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS</u> : xt(temporary)	W	W_CODE	W
<u>CONSTANTS</u> : 'S',IT	Unclass		
<u>VARIABLES</u> : K_CUR_LEVEL K_LACC K_Lxt K_IACC	W W W LA = W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: APPEND

CASE: 2B

SUBJECT: 1

CONDITIONS: $(xt \notin \{X,Y,Z\}) \wedge (W = LA) \wedge (\sim IL) \wedge IT$
xt is a value in the user's working area.
Attempted to append this to a string in the accumulator.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> xt(value)	W	W_CODE	W
<u>CONSTANTS:</u> 'S',IT	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC K_IACC	W W LA=W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	

PRIMITIVE: APPEND

CASE: 3A

SUBJECT: 1

CONDITIONS: $(xt \in X,Y,Z) \wedge (W = LA) \wedge (\sim IL) \wedge (\sim IT) \wedge IC$
xt is a kernel temporary. The domains of
xt do not conform to those in the accumulator.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> xt(temporary)	W	W_CODE	W
<u>CONSTANTS:</u> 'DTYPE','S',IC	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC K_Lxt K_IACC K_Fxt K_FACC	W W W LA=W LX X		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: $LX = W$

PROOF: $(\sim IL) \Rightarrow LA \supset LX$

$W = LA \Rightarrow W \supset LX$

PRIMITIVE: APPEND

CASE: 3B

SUBJECT: 1

CONDITIONS: (xt \notin {X,Y,Z}) \wedge (W = LA) \wedge (~IL) \wedge (~IT) \wedge IC
xt is a user working area value which does not conform to the accumulator contents.

	OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u>	xt(value	W	W_CODE	W
<u>CONSTANTS:</u>	'DTYPE','S',IC	Unclass		
<u>VARIABLES:</u>	K_CUR_LEVEL K_LACC K_IACC K_FACC K_Fxt	W W LA=W W W		
<u>HIGHEST LEVEL OBSERVED:</u>		W	<u>LOWEST LEVEL MODIFIED:</u>	W

PRIMITIVE: APPEND

CASE: 4A

SUBJECT: 1

CONDITIONS: (xt \in {X,Y,Z}) \wedge (W = LA) \wedge (~IL) \wedge (~IT) \wedge (~IC) \wedge IV
xt is a temporary, and appending its tuples to the accumulator would result in duplicate keys.

	OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u>	xt(temporary)	W	W_CODE	W
<u>CONSTANTS:</u>	'DNAME','DTYPE' 'WIDTH','ROLE','S',IV	Unclass		
<u>VARIABLES:</u>	K_CUR_LEVEL K_LACC K_Lxt K_IACC K_Fxt K_FACC K_Vxt K_VACC	W W W LA=W LA=W LA=W LX=W LA=W		
<u>HIGHEST LEVEL OBSERVED:</u>		W	<u>LOWEST LEVEL MODIFIED:</u>	W

See Case 3A for Subject 1

PRIMITIVE: APPEND

CASE: 4B

SUBJECT: 1

CONDITIONS: (xt ∈ {X,Y,Z}) ∧ (W = LA) ∧ (~IL) ∧ (~IC) ∧ IV
xt is a user working area value, but appending it to the accumulator would produce duplicate keys.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> xt(value)	W	W_CODE	W
<u>CONSTANTS:</u> 'DNAME', 'DTYPE' 'WIDTH', 'ROLE', 'S', IV	Unclass		
<u>VARIABLES:</u> K_CUR LEVEL K_LACC K_IACC K_Fxt K_FACC K_Vxt K_VACC	W W LA=W W LA=W W LA=W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: APPEND

CASE: 5A

SUBJECT: 1

CONDITIONS: (xt ∈ {X,Y,Z}) ∧ (W = LA) ∧ (~IL) ∧ (~IT) ∧ (~IC) ∧ (~IV)
xt is a temporary, and there are no exceptions.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> xt	W	W_CODE	W
<u>CONSTANTS:</u> 'DNAME', 'DTYPE' 'WIDTH', 'ROLE', 'S', DN	Unclass		
<u>VARIABLES:</u> K_CUR LEVEL K_LACC K_IACC K_Lxt K_Fxt K_FACC K_Vxt K_VACC	W W LA=W W LA=W LA=W LX=W LA=W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: APPEND

CASE: 5B

SUBJECT: 1

CONDITIONS: (xt \notin X,Y,Z) \wedge (W = LA) \wedge (\sim IL) \wedge (\sim IT) \wedge (\sim IC) \wedge (\sim IV)
xt is a user working area value, and there are no exceptions.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> xt(value)	W	W_CODE	W
<u>CONSTANTS:</u> 'DNAME', 'DTYPE', 'WIDTH', 'ROLE', 'S', DN	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC K_IACC K_Fxt K_FACC K_Vxt K_VACC	W W LA=W W LA=W W LA=W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: APPEND

CASE: 6

SUBJECT: 1

CONDITIONS: W \neq LA

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> xt	W		W ⁹
<u>CONSTANTS:</u>	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC	W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

⁹The "null" return code is at level W.

PRIMITIVE: APPEND

CASE: 1A

SUBJECT: 2

CONDITIONS: $(xt \in \{X,Y,Z\}) \wedge (\sim IL) \wedge (\sim IT) \wedge (\sim IC) \wedge (\sim IV)$
xt is a temporary.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> xt(temporary)	W	K_VACC	LA
<u>CONSTANTS:</u> 'DNAME', 'DTYPE', 'WIDTH', 'ROLE', 'S', DN	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC K_Lxt K_IACC K_Fxt K_FACC K_Vxt K_VACC	W W W LA LX LA LA LA		
HIGHEST LEVEL OBSERVED:	LA	LOWEST LEVEL MODIFIED:	LA

LEMMA: LA \succ W
PROOF: Minimum K-level Invariant
[c.f. § 4.1]

LEMMA: LA \succ LX
PROOF: $\sim IL$

PRIMITIVE: APPEND

CASE: 1B

SUBJECT: 2

CONDITIONS: $(xt \notin X,Y,Z) \wedge (\sim IL) \wedge (\sim IT) \wedge (\sim IC) \wedge (\sim IV)$
xt is a user working area value.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> xt(value)	W	K_VACC	LA
<u>CONSTANTS:</u> 'DNAME', 'DTYPE', 'WIDTH', 'ROLE', 'S', DN	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC K_IACC W_Fxt K_FACC W_Vxt K_VACC	W W LA W LA W LA		
HIGHEST LEVEL OBSERVED:	LA	LOWEST LEVEL MODIFIED:	LA

LEMMA: LA \succ W
PROOF: Minimum K-level Invariant

A.1.3.6 CONCAT

(A) O-function CONCAT(x)

* Concatenate a string in a temporary variable to the string in the *
 * accumulator. Only fields with unique names will be concatenated. *

parameter types

temp x, format(DNAME, DTYPE, WIDTH, ROLE), contents(O, N, T, L, C)

abbreviation

common = |K_FACC[DNAME]| n K_Fx[DNAME] * Field names in common *

exceptions

- *IL: K_LACC > K_Lx * Append must be "up"
- *IV: (K_IACC[T;C] ≠ ('S', 'V')) ∨ (K_Ix[T;C] ≠ ('S', 'V'))
- * Check that these are really strings. *
- *ND: K_Fx[DNAME] ≤ |K_FACC[DNAME]|- * There are no new fields *

effects

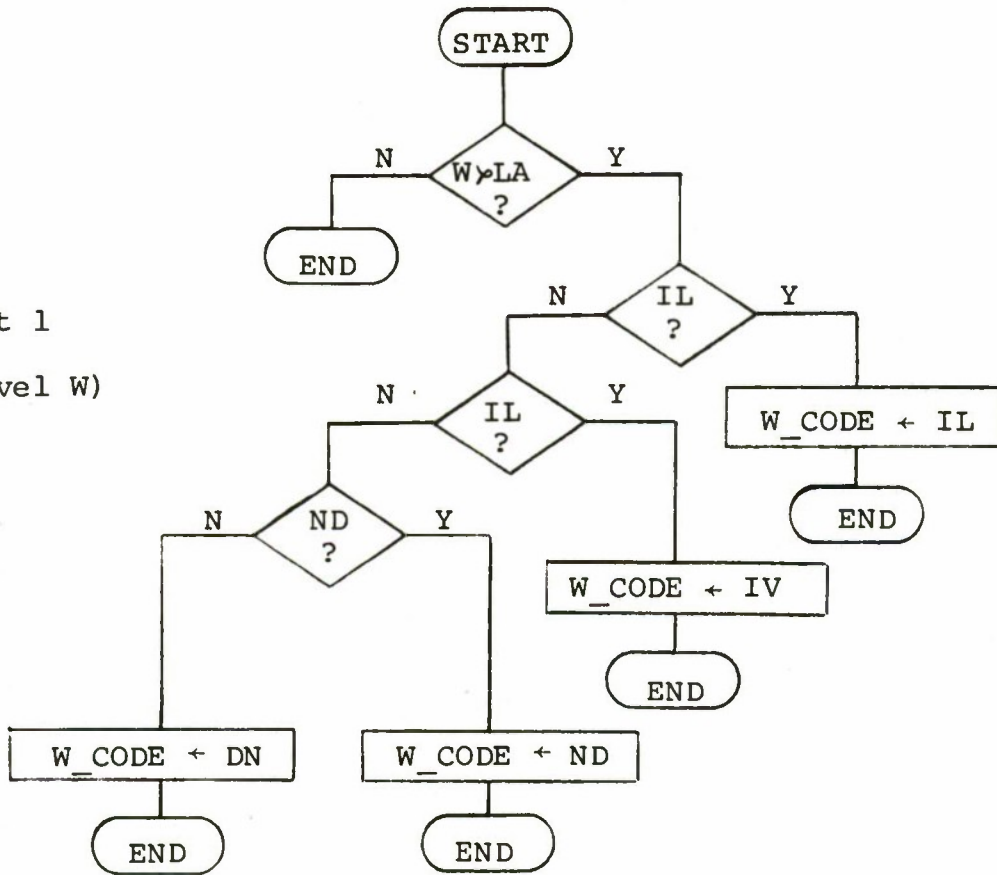
- [1] K_FACC ← | - | ∪ (K_Fx - K_Fx{DNAME ε common})
- * Append format tuples for fields with unique names. *
- [2] K_VACC ← | - | ∪ (K_Vx - K_Vx[common])
- * Append value tuples for fields with unique names. *
- *[3] W_CODE ← DN IF K_CUR_LEVEL > K_LACC

(B) Access table

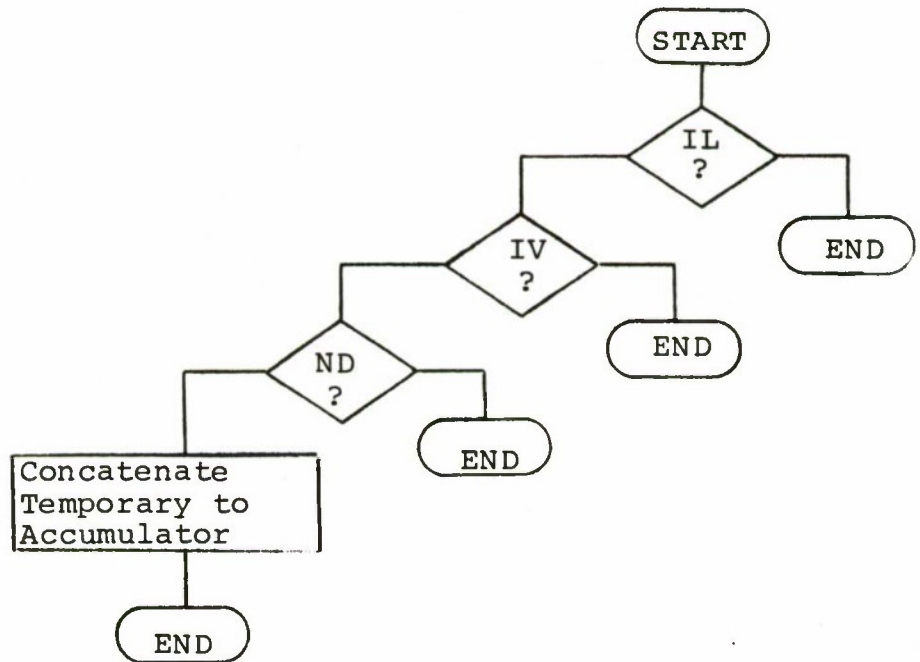
Variables Observed	Variables Modified	Variables Observed and Modified
K_IACC, K_Ix	W_CODE	
K_Fx, K_Vx		K_VACC
K_LACC, K_Lx		K_FACC

CONCAT

Subject 1
(at level W)



Subject 2
(at level LA)



PRIMITIVE: CONCAT

CASE: 1

SUBJECT: 1

CONDITIONS: (W > LA) ^ IL
Level of temporary is not dominated
by the level of the accumulator.

	OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS</u> :	x(temporary)	W	W_CODE	W
<u>CONSTANTS</u> :	IL	Unclass		
<u>VARIABLES</u> :	K_CUR_LEVEL K_LACC K_Lx	W W W		
HIGHEST LEVEL OBSERVED:		W	LOWEST LEVEL MODIFIED:	W

LEMMA: LA=W

PROOF: W > LA

LA > W by Minimum K-level Invariant [c.f. § 4.1]

PRIMITIVE: CONCAT

CASE: 2

SUBJECT: 1

CONDITIONS: (W = LA) ^ (~IL) ^ IV
Accumulator and temporary do not
bot contain strings.

	OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS</u> :	x	W	W_CODE	W
<u>CONSTANTS</u> :	IV	Unclass		
<u>VARIABLES</u> :	K_CUR_LEVEL K_LACC K_Lx K_IACC K_Ix	W W W LA=W LX		
HIGHEST LEVEL OBSERVED:		W	LOWEST LEVEL MODIFIED:	W

LEMMA: LX = W

PROOF: LA > LX By ~IL => W > LX

LX > W By Minimum K-level Invariant

PRIMITIVE: CONCAT

CASE: 3

SUBJECT: 1

CONDITIONS: $(W = LA) \wedge (\sim IL) \wedge (\sim IV) \wedge ND$
There are no new fields in the string to be concatenated.

	OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u>	x	W	W_CODE	W
<u>CONSTANTS:</u>	'DNAME', ND	Unclass		
<u>VARIABLES:</u>	K_CUR_LEVEL K_LACC K_Lx K_IACC K_Ix K_FACC K_Fx	W W W LA=W LX LA=W LX		
HIGHEST LEVEL OBSERVED:		W	LOWEST LEVEL MODIFIED:	W

LEMMA: LX = W

PROOF: $W = LA$ and $LA \not\sim LX$ By $\sim IL \Rightarrow W \not\sim LX$

$LX \not\sim W$ By Minimum K-level Invariant [c.f. § 4.1]

PRIMITIVE: CONCAT

CASE: 4

SUBJECT: 1

CONDITIONS: $(W = LA) \wedge (\sim IL) \wedge (\sim IV) \wedge (\sim ND)$
No exceptions.

	OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u>	x	W	W_CODE	W
<u>CONSTANTS:</u>	'DNAME', DN	Unclass		
<u>VARIABLES:</u>	K_CUR_LEVEL K_LACC K_Lx K_IACC K_Ix K_FACC K_Fx	W W W LA=W LX LA=W LX		
HIGHEST LEVEL OBSERVED:		W	LOWEST LEVEL MODIFIED:	W

LX = W By LEMMA in Case 3.

PRIMITIVE: CONCAT

CASE: 5

SUBJECT: 1

CONDITIONS:

(W $\not\propto$ LA)

User's current level does not dominate accumulator.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> x	W		W
<u>CONSTANTS:</u>	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC	W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: CONCAT

CASE: 1

SUBJECT: 2

CONDITIONS:

(~IL) \wedge (~IV) \wedge (~ND)

No exceptions.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> x	W	K_FACC K_VACC	LA LA
<u>CONSTANTS:</u> 'DNAME',DN	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC K_Lx K_IACC K_Ix K_FACC K_Fx	W W W LA LX LA LX		
HIGHEST LEVEL OBSERVED:	LA	LOWEST LEVEL MODIFIED:	LA

LEMMA: LX $\not\propto$ W

PROOF: Minimum K-level Invariant

LEMMA: LA $\not\propto$ LX

PROOF: ~IL

A.1.3.7 DKD

(A) O-function DKD(lv)

* Copy a data base directory to the kernel accumulator. *

parameter types

level lv

exception

IL: K CUR LEVEL \neq lv
 NF: D_D(lv) = \emptyset

* Current level must dominate *
 * The directory is not found *

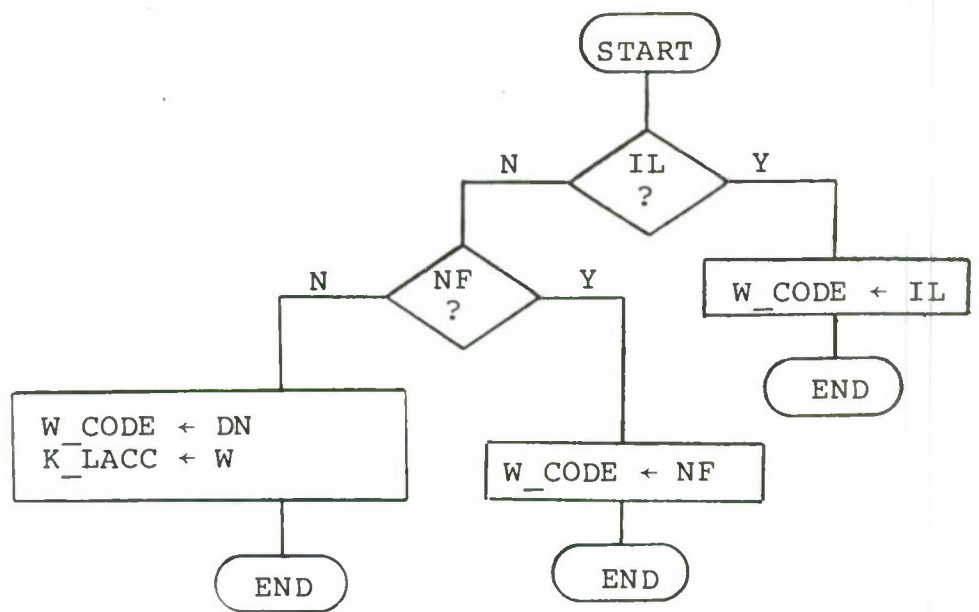
effect

- [1] K LACC ← K CUR LEVEL
- [2] K_FACC ← ('OWNER', 'I', USE_WIDTH, 1) ∪ ('NAME', 'C', MAX_NAME, 2) * This is "minimum" level *
 ∪ ('TYPE', 'C', 1, 3) ∪ ('LEVEL', 'I', LEV_WIDTH, 4) * continued *
- [3] K IACC ← (K CUR ID, 'D', 'R', lv, 'V') * Identify the data *
- [4] K_VACC ← D_D(lv) * Now copy the directory data *
- [5] W_CODE ← DN

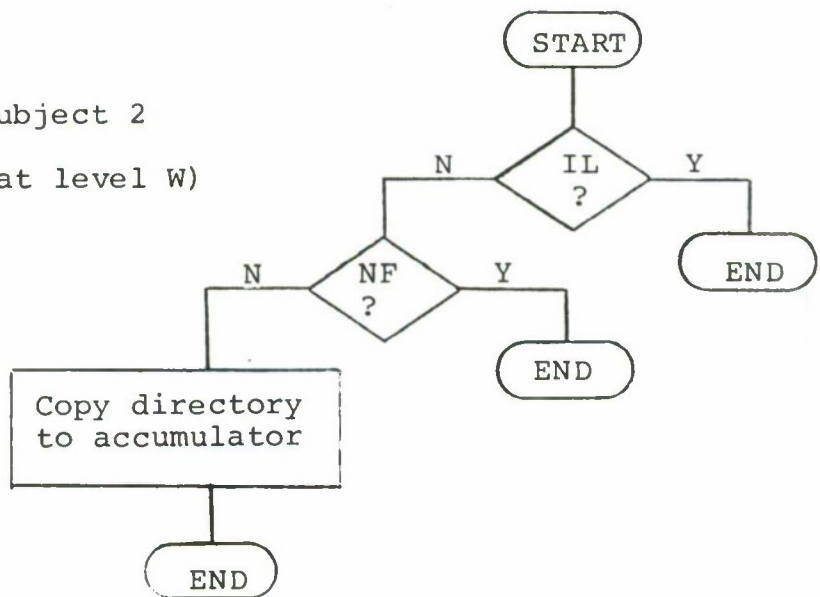
(B) Access table

Variables Observed	Variables Modified	Variables Observed and Modified
K CUR LEVEL		K_FACC
D_D(level)		K_IACC
K_CUR_ID		K_VACC
		K_LACC
		W_CODE

Subject 1
(at level W)



Subject 2
(at level W)



PRIMITIVE: DKD

CASE: 1

SUBJECT: 1

CONDITIONS: IL
User's current level does
dominate the directory level.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> lv	W	W_CODE	W
<u>CONSTANTS:</u> IL	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL	W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: DKD

CASE: 2

SUBJECT: 1

CONDITIONS: (~IL) ^ NF
No directory exists at level lv.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> lv	W	W_CODE	W
<u>CONSTANTS:</u> ∅, NF	Unclass		
<u>VARIABLES:</u> D_D(lv)	lv		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: $W \not\prec lv$

PROOF: ~IL

PRIMITIVE: DKD

CASE: 3

SUBJECT: 1

CONDITIONS: (~IL) ^ (~NF)
No exceptions.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> lv	W	W_CODE K_LACC	W W
<u>CONSTANTS:</u> ∅,1,2,3,4,'OWNER', 'NAME','TYPE','LEVEL',USE WIDTH MAX NAME,LEV WIDTH,'R','V',DN	Unclass		
<u>VARIABLES:</u> D_D(lv) K_CUR_ID	lv W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: W ∽ lv
PROOF: ~IL

PRIMITIVE: DKD

CASE: 1

SUBJECT: 2

CONDITIONS: (~IL) ^ (~NF)
No exceptions. Copy directory into accumulator.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> lv	W	K_FACC K_IACC K_VACC	LA LA LA
<u>CONSTANTS:</u> (Same as case 3)	Unclass		
<u>VARIABLES:</u> D_D(lv) K_CUR_ID	lv W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: W ∽ lv
PROOF: ~IL

LEMMA: LA = W
PROOF: Effect[1] of
specification

A.1.3.8 DKE

(A) O-function DKE(id)

* Copy the specified exact size component to the kernel accumulator. *
 * The object must be open, in order to justify this data movement. *

parameter types

identifier id (OWN,NAM,TYP,LEV)

exception

**NO: K_OPEN(id,*) = FALSE * The object is not open *

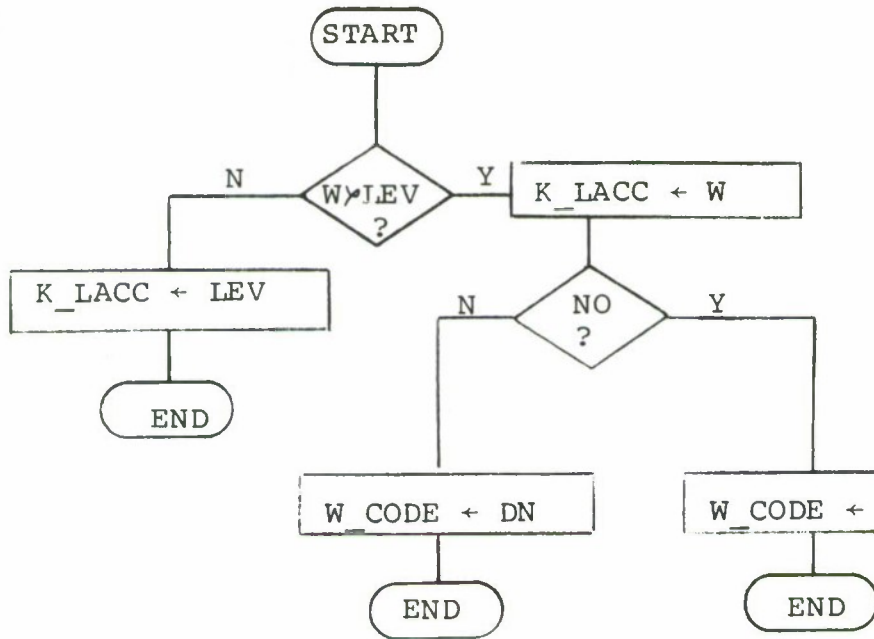
effects

```
[1] K_LACC ← LEV IF LEV ≠ K_CUR_LEVEL ELSE K_CUR_LEVEL
[2] K_FACC ← ('EXACT', 'I', SIZ_WIDTH, 1) * Set format for a single value *
[3] K_IACC ← (id, 'E') * Identify the accumulator contents *
[4] K_VACC ← D E(id) * Copy the exact size value *
*[5] W_CODE ← DN IF K_CUR_LEVEL ≠ LEV * Return code if a *
```

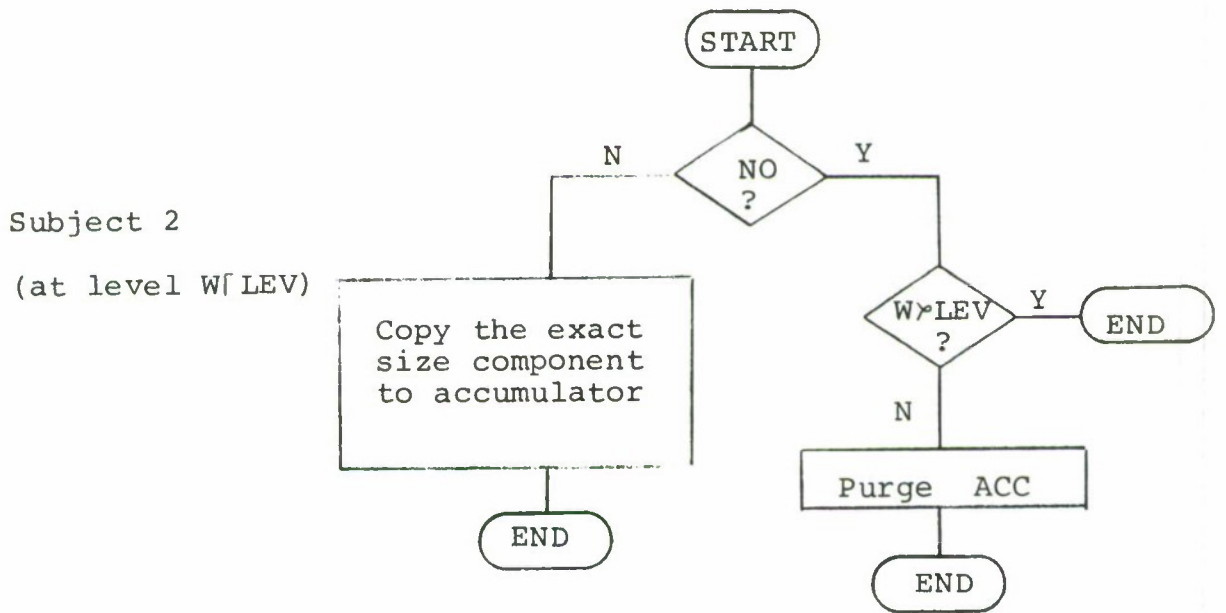
(B) Access table

Variables Observed	Variables Modified	Variables Observed and Modified
K_OPEN		K_FACC
K_CUR_LEVEL		K_IACC
D_E(id)		K_VACC
		W_CODE
		K_LACC

DKE



Subject 1
(at level W)



Subject 2
(at level W/LEV)

PRIMITIVE: DKE

CASE: 1

SUBJECT: 1

CONDITIONS: W ~~Y~~ LEV
Level of object strictly dominates

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	K_LACC	W
<u>CONSTANTS:</u>	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL	W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: DKE

CASE: 2

SUBJECT: 1

CONDITIONS: (W ~~Y~~ LEV) ^ NO
Object has not been opened.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	K_LACC W_CODE	W W
<u>CONSTANTS:</u> NO, FALSE	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_OPEN[LEV]	W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: DKE

CASE: 3

SUBJECT: 1

CONDITIONS: (W \neq LEV) \wedge (~NO)
No exceptions.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	K_LACC W_CODE	W W
<u>CONSTANTS:</u> DN, FALSE	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_OPEN[LEV]	W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: DKE

CASE: 1

SUBJECT: 2

CONDITIONS: NO \wedge (W \neq LEV)
Dominated object is not open.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W		W
<u>CONSTANTS:</u> FALSE	Unclass		
<u>VARIABLES:</u> K_OPEN[LEV] K_CUR_LEVEL	W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: DKE

CASE: 2

SUBJECT: 2

CONDITIONS: NO \wedge (W $\not\prec$ LEV)
Strictly dominating object is not open.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	K_FACC K_IACC K_VACC	LEV LEV LEV
<u>CONSTANTS:</u> \emptyset , FALSE	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_OPEN[LEV]	W LEV		
HIGHEST LEVEL OBSERVED:	LEV	LOWEST LEVEL MODIFIED:	LEV

LEMMA: LEV $\not\prec$ W
PROOF: Minimum K-level Invariant

PRIMITIVE: DKE

CASE: 3

SUBJECT: 2

CONDITIONS: ~NO
Object is open, no exception.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	K_FACC K_IACC K_VACC	W[LEV W[LEV W[LEV
<u>CONSTANTS:</u> FALSE, 'EXACT', 'I', _SIZ_WIDTH, 'E'	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_OPEN[LEV]	W W[LEV		
HIGHEST LEVEL OBSERVED:	W[LEV	LOWEST LEVEL MODIFIED:	W[LEV

A.1.3.9 KDM

(A) O-function KDM(id)

* Copy an access permission matrix from the kernel accumulator *
 * to the data base. This is required for EXTEND-PERMISSION *

parameter types

identifier id (OWN,NAM,TYP,LEV), history (CREATION,USER,MODIFICATION)

abbreviation * This is a proper format for a permission matrix. *

proper_format = ('USER', 'I', USE_WIDTH, 1) U ('VISIBLE', 'L', 1, 0)
 exceptions

- IL: K_LACC ≠ LEV * Accumulator level is wrong *
- *NO: K_OPEN(id,EXPM) = FALSE * Object is not open *
- *IC: K_IACC ≠ (id,'M') * It's not the permission matrix *
- *IV: K_FACC ≠ proper format * It's not in proper format *
- *SZ: (√(K_VACC) + √(D_F(id)) + √(D_V(id))) > D_Z(id) * Blows space limit *
- *√ is a system function to compute the size of something. *

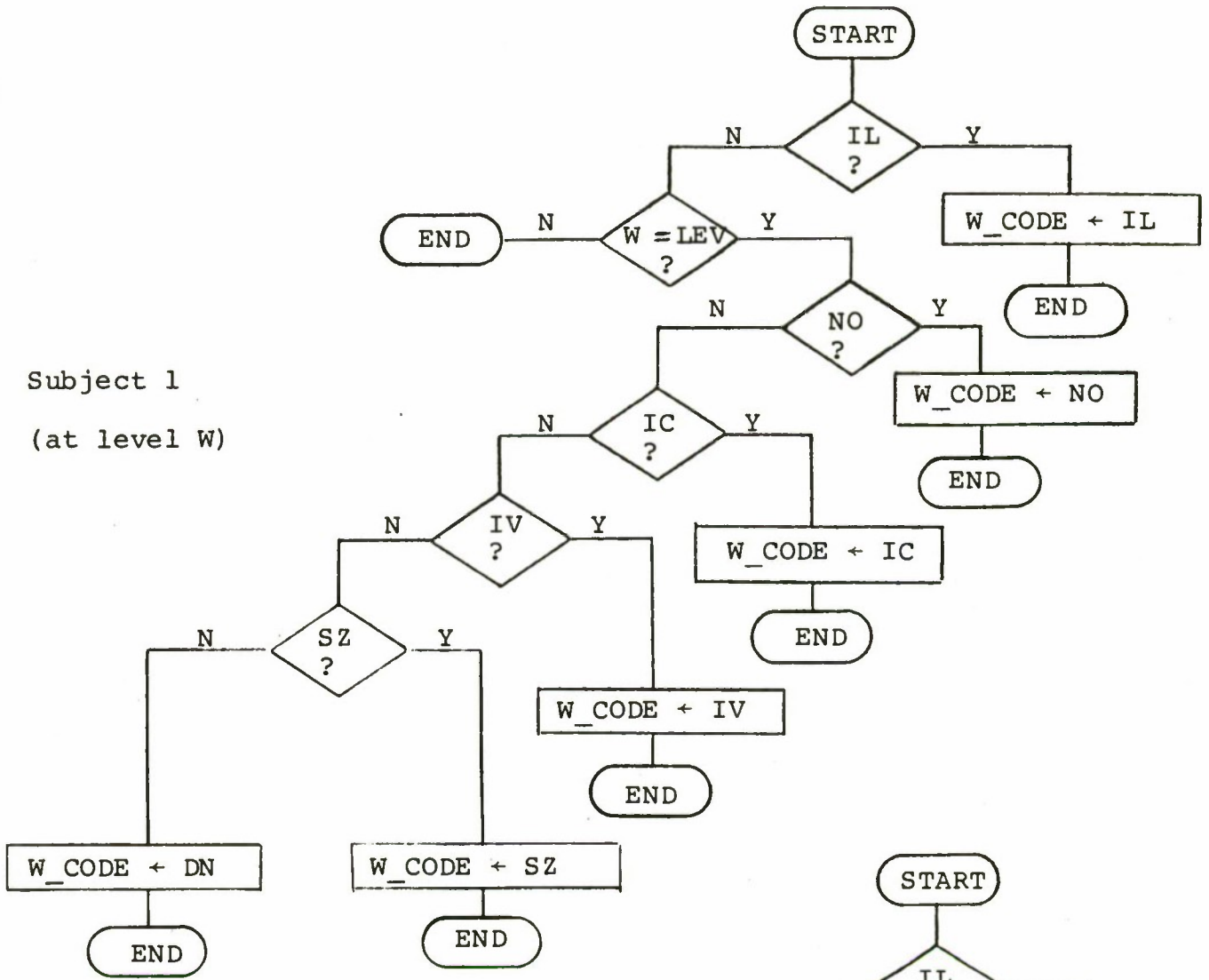
effect

- [1] D_M(id) ← K_VACC * Copy the new permission matrix *
- [2] D_E(id) ← √(K_VACC) + √(D_F(id)) + √(D_V(id)) * Compute new size *
- [3] D_H(id) ← (|-[CREATION],K_CUR_ID,K_CUR_TIME) * Copy the creation_date, and construct rest of history *
- *[4] W_CODE ← DN IF K_CUR_LEVEL = LEV

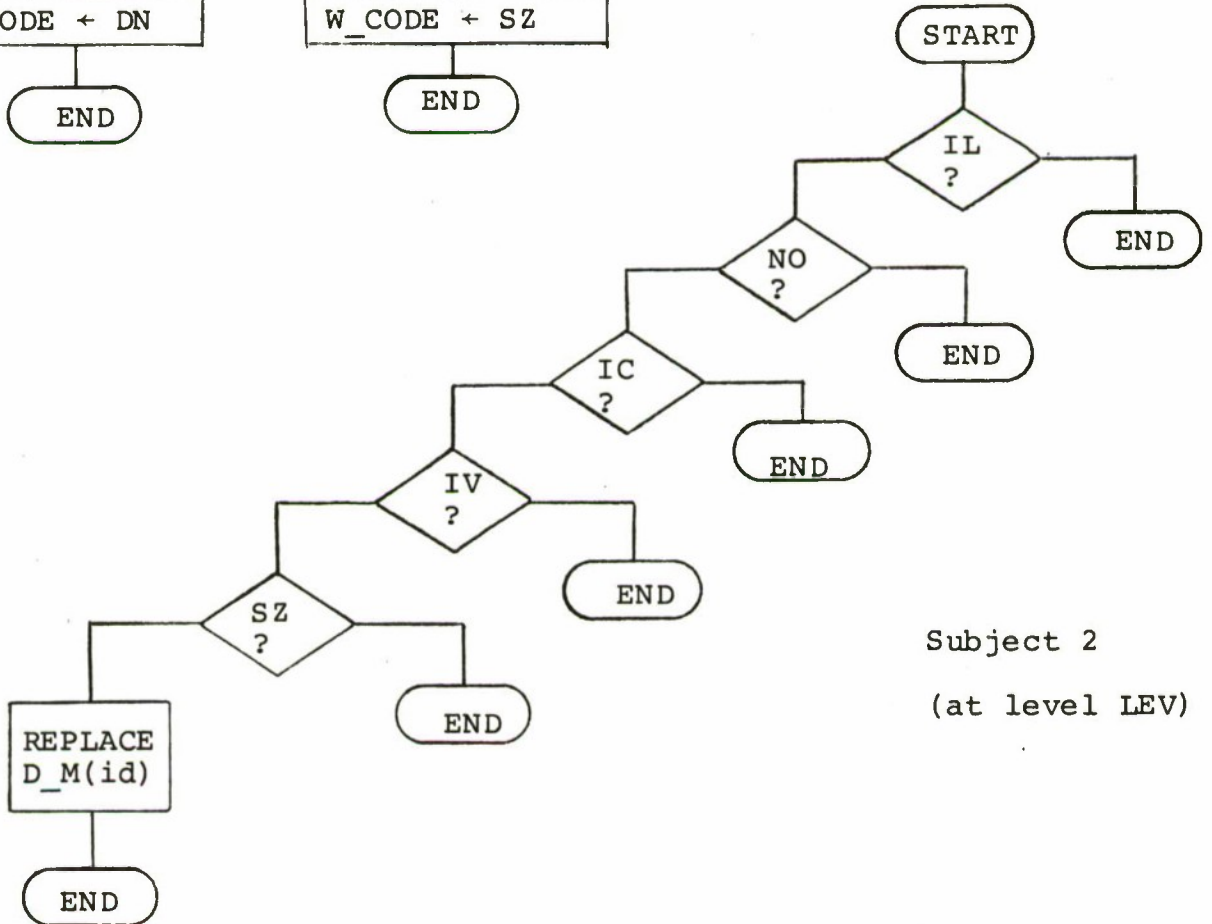
(B) Access table

Variables Observed	Variables Modified	Variables Observed and Modified
K_OPEN, K_CUR_ID	D_M(id)	D_H(id)
K_FACC, K_IACC, K_VACC	D_E(id)	
D_F(id), D_V(id), D_Z(id)	W_CODE	
K_CUR_LEVEL, K_CUR_TIME		
K_LACC		

Subject 1
(at level W)



Subject 2
(at level LEV)



PRIMITIVE: KDM

CASE: 1

SUBJECT: 1

CONDITIONS: IL
Accumulator is not at proper level (W).

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	W_CODE	W
<u>CONSTANTS:</u> IL	Unclass		
<u>VARIABLES:</u> K_CUR LEVEL K_LACC	W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: KDM

CASE: 2

SUBJECT: 1

CONDITIONS: (~IL) ^ NO ^ (W = LEV)
Dominated object is not open.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	W_CODE	W
<u>CONSTANTS:</u> NO	Unclass		
<u>VARIABLES:</u> K_CUR LEVEL K_LACC K_OPEN[LEV]	W W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: KDM

CASE: 3

SUBJECT: 1

CONDITIONS: (~IL) ^ (W = LEV) ^ (~NO) ^ IC
The accumulator does not contain
the object's permission matrix.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	W_CODE	W
<u>CONSTANTS:</u> IC	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC K_OPEN[LEV] K_IACC	W W W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: KDM

CASE: 4

SUBJECT: 1

CONDITIONS: (~IL) ^ (W = LEV) ^ (~NO) ^ (~IC) ^ IV
Format of accumulator contents is inappropriate.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	W_CODE	W
<u>CONSTANTS:</u> 0,1,'I','L','M',IV, 'USER','VISIBLE',USE_WIDTH	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC K_OPEN[LEV] K_IACC K_FACC	W W W W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: KDM

CASE 5

SUBJECT: 1

CONDITIONS: (\sim IL) \wedge (W = LEV) \wedge (\sim NO) \wedge (\sim IC) \wedge (\sim IV) \wedge SZ
This permission matrix would cause the maximum size of this object to be exceeded.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	W_CODE	W
<u>CONSTANTS:</u> o,l,'I','L','M',SZ 'USER','VISIBLE',USE_WIDTH	Unclass		
<u>VARIABLES:</u> K_CUR LEVEL K_LACC K_OPEN[LEV] K_IACC K_FACC D_F(id) D_V(id) D_Z(id)	W W LA LA LA LEV=W LEV=W LEV=W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: LA = W

PROOF: (\sim IL) \wedge (W = LEV)

PRIMITIVE: KDM

CASE 6

SUBJECT: 1

CONDITIONS: (~IL) ^ (~NO) ^ (~IC) ^ (~IV) ^ (~SZ) ^ (W = LEV)
No exceptions.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	W_CODE	W
<u>CONSTANTS:</u> 0,1,'I','L','M',DN 'USER','VISIBLE',USE_WIDTH	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC K_OPEN[LEV] K_IACC K_FACC D_F(id) D_V(id) D_Z(id) K_CUR TIME D_H(id)	W W W W W LEV=W LEV=W LEV=W W LEV=W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

A.1.3.10 KDV

(A) O-function KDV(id)

* Copy an object's format and values from *
 * the accumulator to the data base *

parameter types

identifier id (OWN,NAM,TYP,LEV), history (CREATION,USER,MODIFIED)

exceptions

IL: K LACC ≠ LEV * Accumulator level is wrong *
 *NO: K_OPEN(id,STOR) = FALSE * The object is NOT open. *
 *IC: K_IACC ≠ (id,'V') * Incorrect accumulator contents. *
 *SZ: (←(D_M(id)) + ←(K_FACC) + ←(K_VACC)) > D_Z(id) * Not enough space *

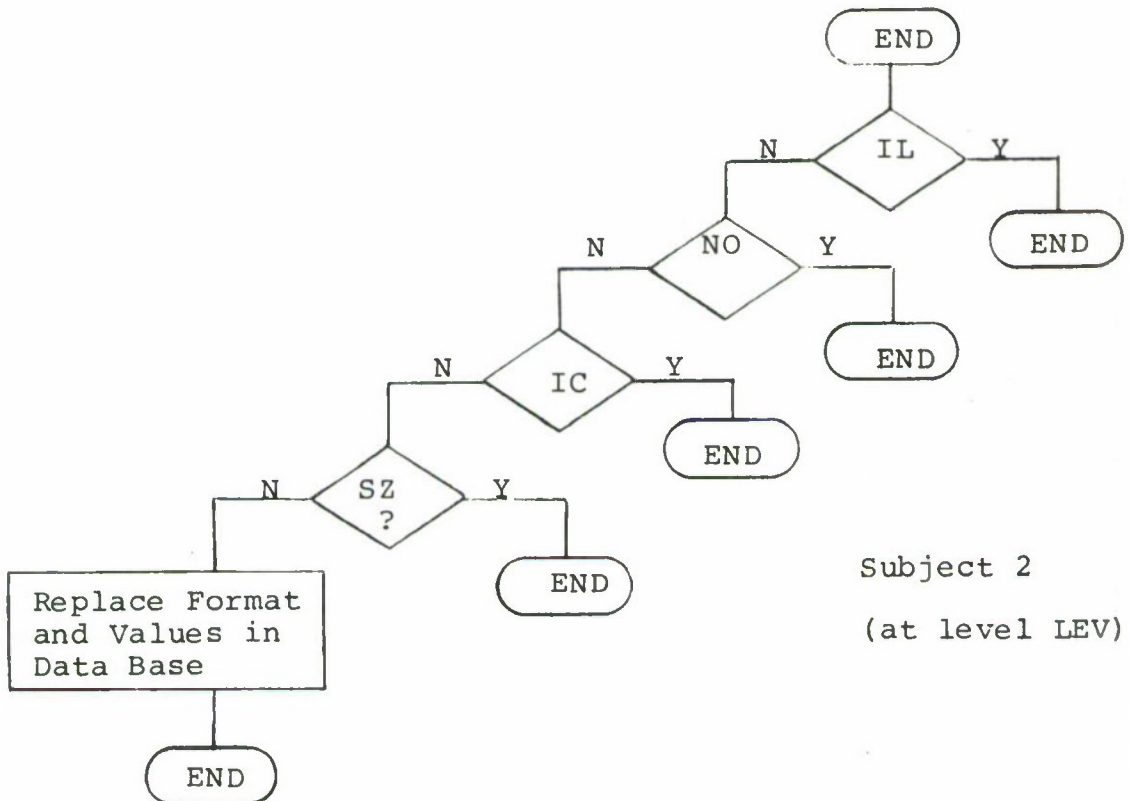
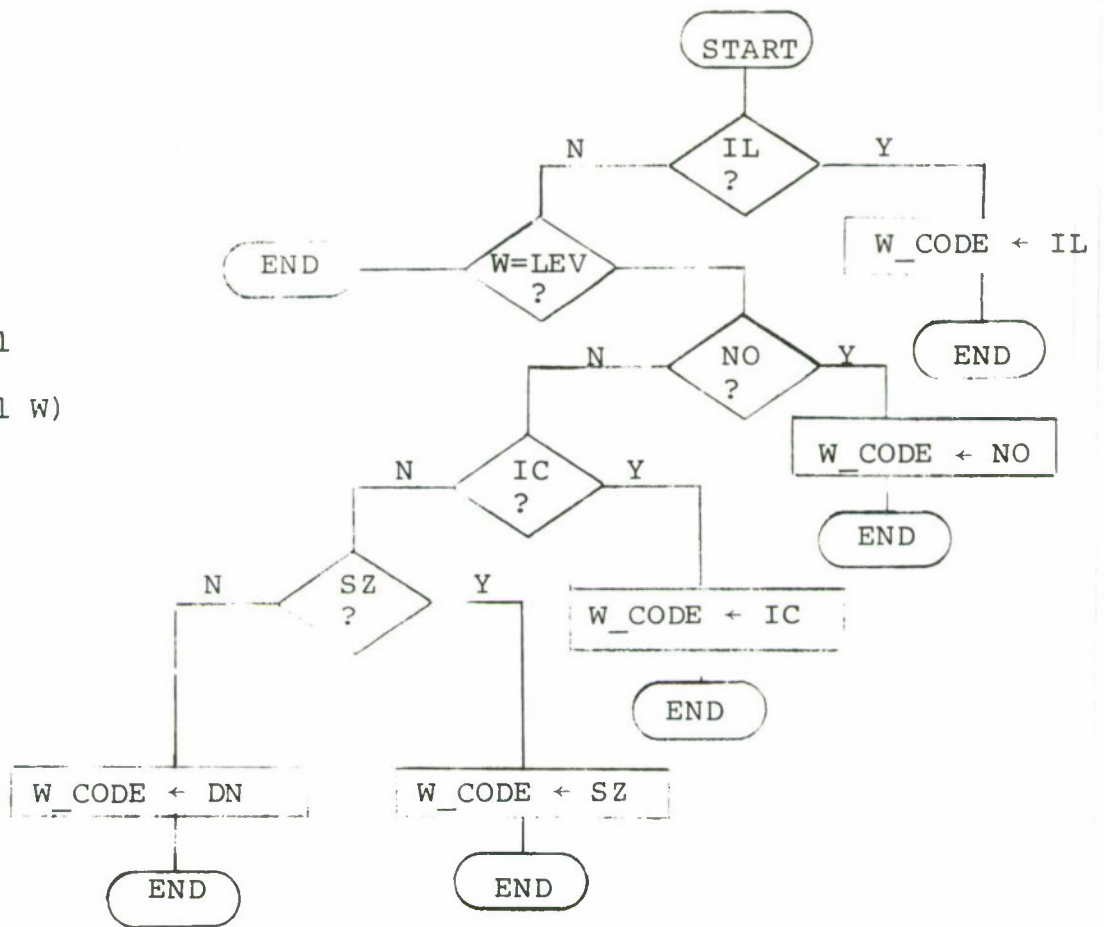
effects

[1] D_E(id) ← ←(D_M(id)) + ←(K_FACC) + ←(K_VACC) * Reset current size *
 [2] D_H(id) ← (← [CREATION],K_CUR_ID,K_CUR_TIME) * Update the history *
 [3] D_F(id) ← K_FACC * Copy the format *
 [4] D_V(id) ← K_VACC * Copy the table of values *
 *[5] W_CODE ← DN IF K_CUR_LEVEL = LEV

(B) Access table

Variables Observed	Variables Modified	Variables Observed and Modified
K_OPEN		D_H(id)
K_IACC,K_LACC		
D_M(id),D_Z(id)		
K_FACC,K_VACC		D_F(id)
K_CUR_ID		D_V(id)
K_CUR_TIME		W_CODE
K_CUR_LEVEL		

Subject 1
(at level W)



Subject 2
(at level LEV)

PRIMITIVE: KDV

CASE: 1

SUBJECT: 1

CONDITIONS: IL

Accumulator level is not equal to object level.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	W_CODE	W
<u>CONSTANTS:</u> IL	Unclass		
<u>VARIABLES:</u> K_LACC	W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: KDV

CASE: 2

SUBJECT: 1

CONDITIONS:

(~IL) ^ (W ≠ LEV)
Object being replaced is at
a strictly dominating level.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W		W
<u>CONSTANTS:</u>	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC	W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: KDV

CASE: 3

SUBJECT: 1

CONDITIONS: (~IL) ^ (W = LEV) ^ NO
Object is not open with STOR access.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	W_CODE	W
<u>CONSTANTS:</u> STOR,NO	Unclass		
<u>VARIABLES:</u> K_LACC K_CUR_LEVEL K_OPEN[LEV]	W W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: KDV

CASE: 4

SUBJECT: 1

CONDITIONS: (~IL) ^ (W = LEV) ^ (~NO) ^ IC
Accumulator does not contain
the appropriate value set.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	W_CODE	W
<u>CONSTANTS:</u> 'V',STOR,IC	Unclass		
<u>VARIABLES:</u> K_LACC K_CUR_LEVEL K_OPEN[LEV] K_IACC	W W LEV=W LA		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: LA = W

PROOF: ~IL => LA = LEV

But LEV = W

PRIMITIVE: KDV

CASE 5

SUBJECT: 1

CONDITIONS: (~IL) ^ (W = LEV) ^ (~NO) ^ (~IC) ^ SZ
User has insufficient space to store
the value set in the data base.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	W_CODE	W
<u>CONSTANTS:</u> 'V',STOR,SZ	Unclass		
<u>VARIABLES:</u> K_LACC K_OPEN[LEV] K_IACC D_M(id) K_VACC D_Z(id)	W W LA LEV=W LA LEV=W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: LA = W
PROOF: ~IL => LA = LEV
But LEV = W

PRIMITIVE: KDV

CASE 6

SUBJECT: 1

CONDITIONS: (\sim IL) \wedge (W = LEV) \wedge (\sim NO) \wedge (\sim IC) \wedge (\sim SZ)
No exceptions.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	W_CODE	W
<u>CONSTANTS:</u> 'V', STOR, DN	Unclass		
<u>VARIABLES:</u> K_LACC K_OPEN[LEV] K_IACC D_M(id) K_FACC K_VACC D_Z(id) D_H(id) K_CUR_ID K_CUR_LEVEL K_CUR_TIME	W W LA LEV=W LA LA LEV=W LEV=W W W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: LA = W
PROOF: \sim IL \Rightarrow LA = LEV
But LEV = W

PRIMITIVE: KDV

CASE 1

SUBJECT: 2

CONDITIONS: (~IL) ^ (~NO) ^ (~IC) ^ (~SZ)

No exceptions. Replace value set in data base.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id	W	D_E(id) D_V(id) D_F(id)	LEV LEV LEV
<u>CONSTANTS:</u> 'V',STOR	Unclass		
<u>VARIABLES:</u> K_LACC K_OPEN[LEV] K_IACC K_FACC K_VACC D_M(id) D_Z(id) D_H(id) K_CUR_ID K_CUR_TIME K_CUR_LEVEL	W LEV LA LA LA LEV LEV LEV W W W		
HIGHEST LEVEL OBSERVED:	LEV	LOWEST LEVEL MODIFIED:	LEV

LEMMA: LEV > W

PROOF: K_OPEN(id,EXPM) = TRUE by ~NO

The derivations of the Access set O and Auth O V-functions in O APPEND complete the proof.

LEMMA: LA = LEV

PROOF: ~IL

A.1.3.11 KDZ

(A) O-function KDZ(n,t)

- * Resize an object (maximum space) in the data base by copying
- * its maximum size value from the accumulator to the data base.
- * The session quota is updated appropriately. The object
- * needn't be open, since this can be done ONLY by an object's
- * owner, at the current level.

parameter types

name n, type t, level lv

abbreviation

id = K CUR ID,n,t,K CUR LEVEL
 change = K_VACC - D_Z(id) +

exception

IL: K CUR LEVEL ≠ K_LACC
 NE: D_Z(id) = ∅
 IC: K_IACC ≠ (id,'Z')

SZ: (K_VACC < D_E(id)) ∨ (change > K_CUR_QTA
 * new size is less than current OR change is more than session quota

effect

- [1] D_Z(id) ← K_VACC
- [2] K_CUR_QTA ← + - change
- [3] W_CODE ← DN

- * Identify the data base object
- * This is the amount of size change

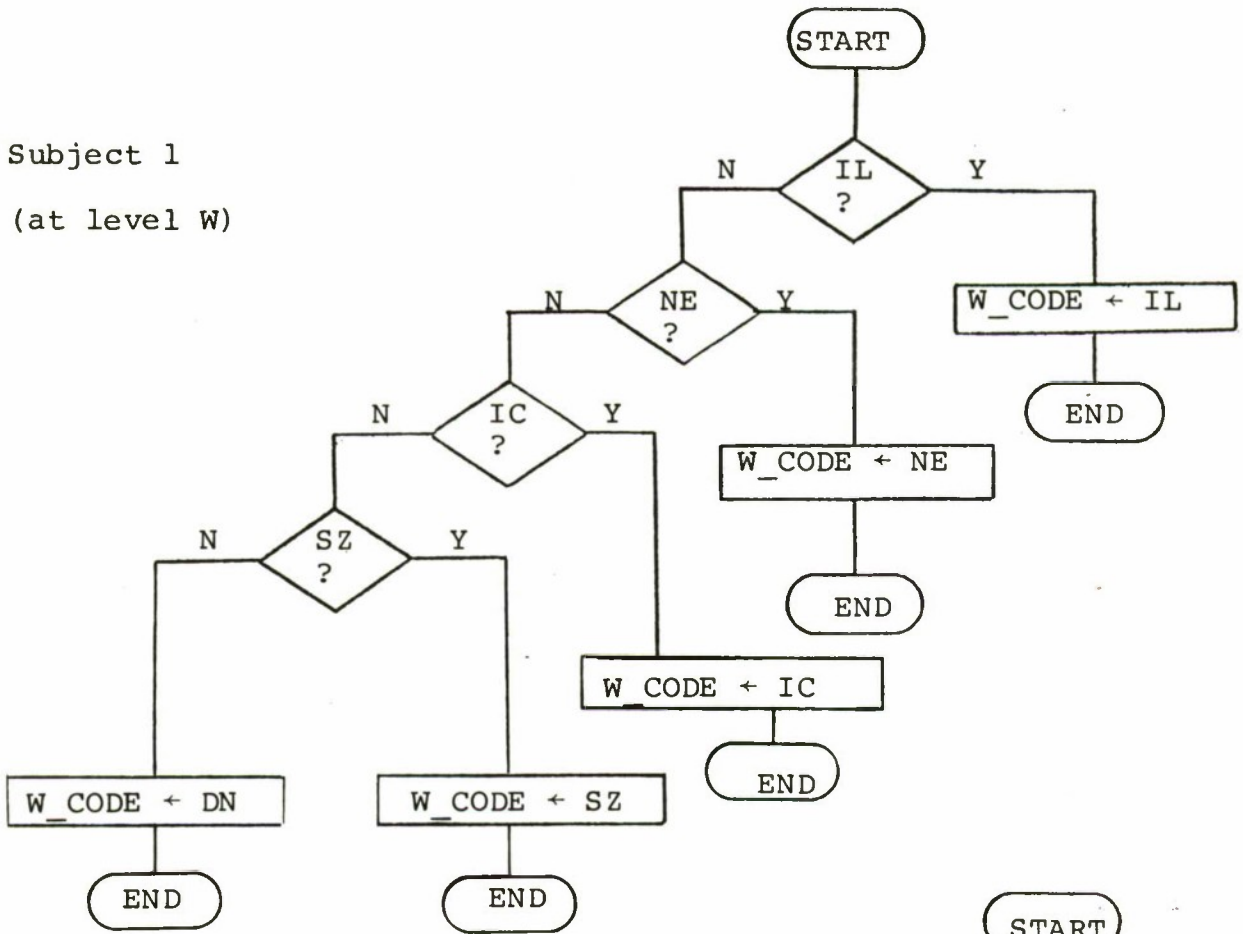
- * Accumulator level is wrong
- * Component object doesn't exist
- * It is not a maximum size

- * Set the new maximum size
- * Update session quota

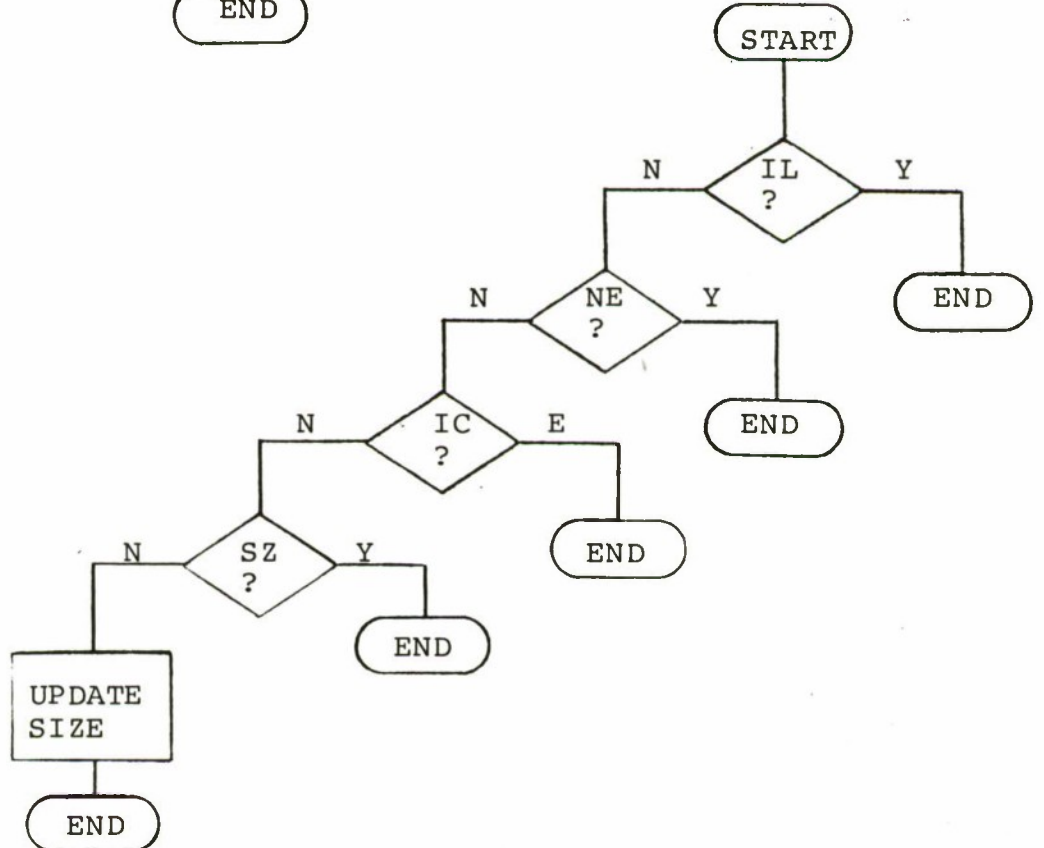
(B) Access table

Variables Observed	Variables Modified	Variables Observed and Modified
K CUR ID	W_CODE	D_Z(id)
K_VACC,K_IACC		K_CUR_QTA
K_CUR_LEVEL		
D_E(id)		
K_LACC		

Subject 1
(at level W)



Subject 2
(at level W)



PRIMITIVE: KDZ

CASE: 1

SUBJECT: 1

CONDITIONS: IL

Accumulator level is not equal to
the user's current level.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS</u> : n,t	W	W_CODE	W
<u>CONSTANTS</u> : IL	Unclass		
<u>VARIABLES</u> : K_LACC K_CUR_LEVEL	W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: KDZ

CASE: 2

SUBJECT: 1

CONDITIONS: (~IL) ^ NE
Non-existent object.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS</u> : n,t	W	W_CODE	W
<u>CONSTANTS</u> : ø,NE	Unclass		
<u>VARIABLES</u> : K_LACC K_CUR_LEVEL D_Z(i \bar{d}) K_CUR_ID	LA W W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: LA = W

PROOF: ~IL

PRIMITIVE: KDZ

CASE: 3

SUBJECT: 1

CONDITIONS: (~IL) ^ (~NE) ^ IC
The accumulator does not contain
the required exact size component.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> n,t	W	W_CODE	W
<u>CONSTANTS:</u> ∅, 'z', IC	Unclass		
<u>VARIABLES:</u> K_LACC K_CUR_LEVEL D_Z(id̄) K_IACC K_CUR_ID	LA W W LA W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: LA = W
PROOF: ~IL

PRIMITIVE: KDZ

CASE: 4

SUBJECT: 1

CONDITIONS: (~IL) ^ (~NE) ^ (~IC) ^ SZ
The current exact size exceeds
the proposed maximum.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> n,t	W	W_CODE	W
<u>CONSTANTS:</u> ∅, 'z', SZ	Unclass		
<u>VARIABLES:</u> K_LACC K_CUR_ID, K_CUR_LEVEL D_Z(id̄) K_IACC, K_VACC D_E(id) K_CUR_QTA	W W W LA W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: LA = W
PROOF: ~IL

PRIMITIVE: KDZ

CASE: 5

SUBJECT: 1

CONDITIONS: (\sim IL) \wedge (\sim NE) \wedge (\sim IC) \wedge (\sim SZ)
No exceptions.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> n, t	W	W_CODE	W
<u>CONSTANTS:</u> \emptyset , 'z', DN	Unclass		
<u>VARIABLES:</u> K_LACC K_CUR_LEVEL, K_CUR_ID D_E(id), D_Z(id) K_CUR_QTA K_IACC, K_VACC	W W W W LA		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: LA = W
PROOF: \sim IL

PRIMITIVE: KDZ

CASE: 1

SUBJECT: 2

CONDITIONS: (\sim IL) \wedge (\sim NE) \wedge (\sim IC) \wedge (\sim SZ)
No exceptions.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> n, t	W	D_Z(id) K_CUR_QTA	W W
<u>CONSTANTS:</u> \emptyset , 'z', DN	Unclass		
<u>VARIABLES:</u> K_LACC K_CUR_LEVEL, K_CUR_ID D_E(id), D_Z(id) K_CUR_QTA K_IACC, K_VACC	W W W W LA		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: LA = W
PROOF: \sim IL

A.1.3.12

(A) O-function KWA(n)

* Copy the accumulator contents to the working area *

parameter types

name n, contents (O,N,T,L,C)

exceptions * Note that returning NOTHING instead of IL *
 * transmits the same information *

IL: K CUR_LEVEL > K_LACC
 ND: ~Discretionary_kwa

* Current level must dominate *
 * No discretionary authorization *

effect

[1] W_Fn ← K_FACC
 [2] W_Vn ← K_VACC
 [3] W_CODE ← DN

* Copy the format *
 * Copy the values *

(B) Access table

Variables Observed	Variables Modified	Variables Observed and Modified
--------------------	--------------------	---------------------------------

K CUR_LEVEL		W_Fn
K_IACC,K_LACC		W_Vn
K_CUR_ID		W_CODE
K_OPEN		
K_FACC,K_VACC		

(C) KWA V-functions

(i) V_function Discretionary_kwa : boolean

* Return a boolean indication of whether or not the *
* current user has discretionary access authorization *
* to the accumulator contents. *

range

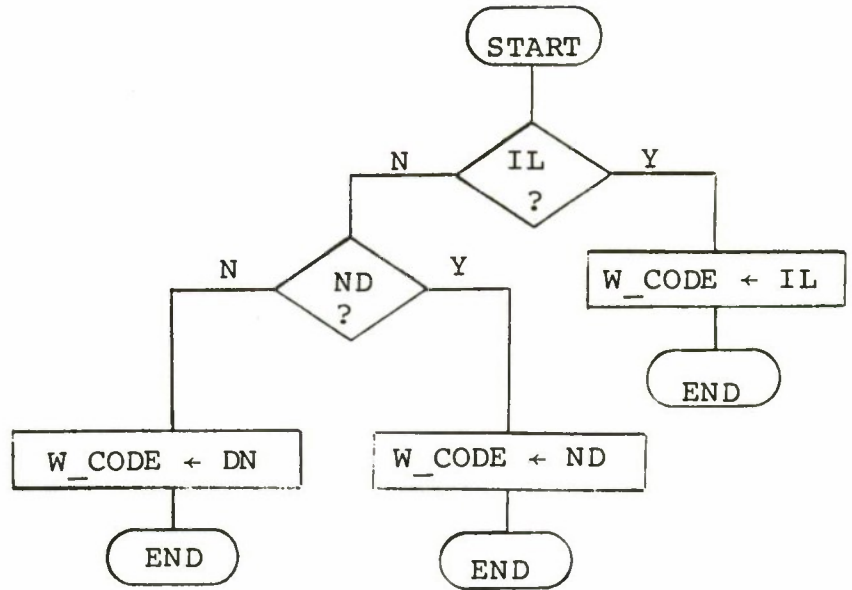
TRUE, FALSE

derivation

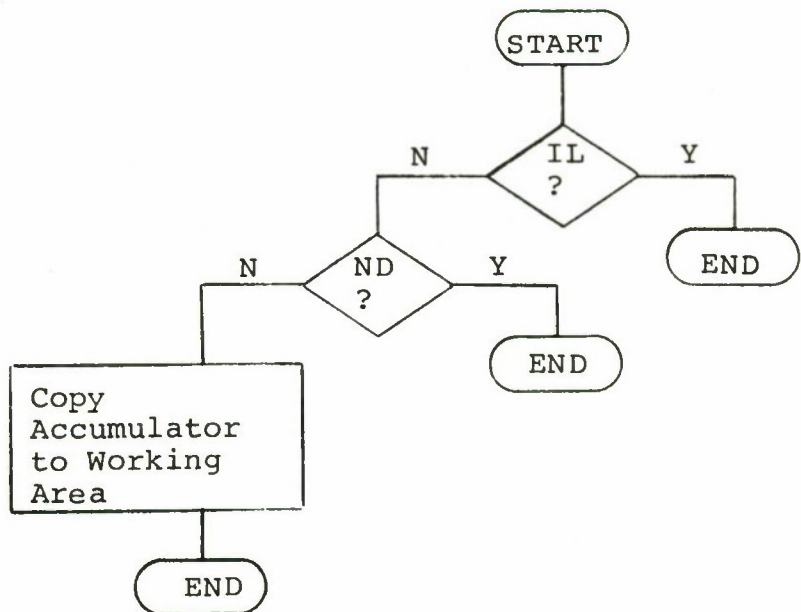
```
TRUE IF K_IACC[O] = K_CUR_ID * DKD,DKQ and LIST_DOWN DATA *  
ELSE K_OPEN(K_IACC[O;N;T;L],RDSZ) IF K_IACC[C] = 'E' * Read current size? *  
ELSE K_OPEN(K_IACC[O;N;T;L],RDHS) IF K_IACC[C] = 'H' * Read history? *  
ELSE K_OPEN(K_IACC[O;N;T;L],RDPM) IF K_IACC[C] = 'M' * Read permission matrix? *  
ELSE K_OPEN(K_IACC[O;N;T;L],RSRV) IF K_IACC[C] = 'R' * Read reservation? *  
ELSE K_OPEN(K_IACC[O;N;T;L],RETR) IF K_IACC[C] = 'V' * Read values? *  
ELSE K_OPEN(K_IACC[O;N;T;L],RDSZ) IF K_IACC[C] = 'Z' * Read max. size? *  
ELSE FALSE * Something must be wrong, fail. *
```

KWA

Subject 1
(at level W)



Subject 2
(at level W)



PRIMITIVE: KWA

CASE: 1

SUBJECT: 1

CONDITIONS: IL
User's current level does not equal accumulator level.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> n	W	W_CODE	W
<u>CONSTANTS:</u> IL	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC	W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: KWA

CASE: 2

SUBJECT: 1

CONDITIONS: (~IL) ^ ND
User has no discretionary authorization.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> n	W	W_CODE	W
<u>CONSTANTS:</u> RDSZ, RDHS, RDPM, RSRV, RETR, 'E', 'H', 'M', 'R', 'V', 'Z', ND	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC K_CUR_ID K_OPEN[LA]	W W W LA		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: LA = W

PROOF: Minimum K-level Invariant and ~IL

PRIMITIVE: KWA

CASE: 3

SUBJECT: 1

CONDITIONS: (~IL) ^ (~ND)
No exceptions.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> n	W	W_CODE	W
<u>CONSTANTS:</u> RDSZ, RDHS, RDPM, RSRV, RETR, 'E', 'H', 'M', 'R', 'V', 'Z', DN	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC K_CUR_ID K_OPEN[LA] K_FACC K_VACC	W W W LA LA LA		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: LA = W

PROOF: Minimum K-level Invariant and ~IL

PRIMITIVE: KWA

CASE: 1

SUBJECT: 2

CONDITIONS:

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> n	W	W_Fn W_Vn	W W
<u>CONSTANTS:</u> RDSZ, RDHS, RDPM, RSRV, RETR, 'E', 'H', 'M', 'R', 'V', 'Z'	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_LACC K_CUR_ID K_OPEN K_VACC K_FACC	W W W LA LA LA		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

LEMMA: LA = W

PROOF: Minimum K-level Invariant and ~IL

A.1.4 SIGNON, SIGNOFF and MOVE

The security kernel functions, SIGNON, SIGNOFF and MOVE are discussed in Section V of this report. All three functions are in violation of the *-property for a successful execution, although the simple security principle is always upheld.

In this section, validations are included for these functions, to illustrate precisely how the *-property is violated. It will be seen that nothing other than a successful execution can induce a *-property violation and that all three functions may only be invoked by trusted subjects operating at the highest protection level.

The concept of the trusted subject, known as the User Controller Process (UCP), was introduced in the functional design report , Section 4.1. In initiating a sign-on to the DMS, the UCP must be restricted, by the hardware if necessary, to respond only to a human user's request to link to the DMS and not to the request of a process acting on behalf of a user.

The Data Base Administrator (DBA) is also considered to be a trusted subject and only a process acting at system high (SYS_HI), on behalf of the DBA, is permitted to declassify database objects. This is the purpose of the primitive MOVE.

A.1.4.1 SIGNON

(A) O-function SIGNON(u,lv,s,v)

```
* This is a request by the user control (and authentication) *
* process to sign a user on to the secure DMS. Users cannot *
* request kernel functions until they have been signed on. *
* A working area and an initialized set of variables are *
* allocated to the user. The user process is spawned. *
```

parameter types

user u, level lv, size s, boolean v * v is visibility indicator *

abbreviations

usent = (D V(DBA,'DBA_ULIST','R',SYS_HI) {D_F.USERID = u} * relation type is user_ent *
 * Select the user's entry from the DBA's user relation. *

exceptions

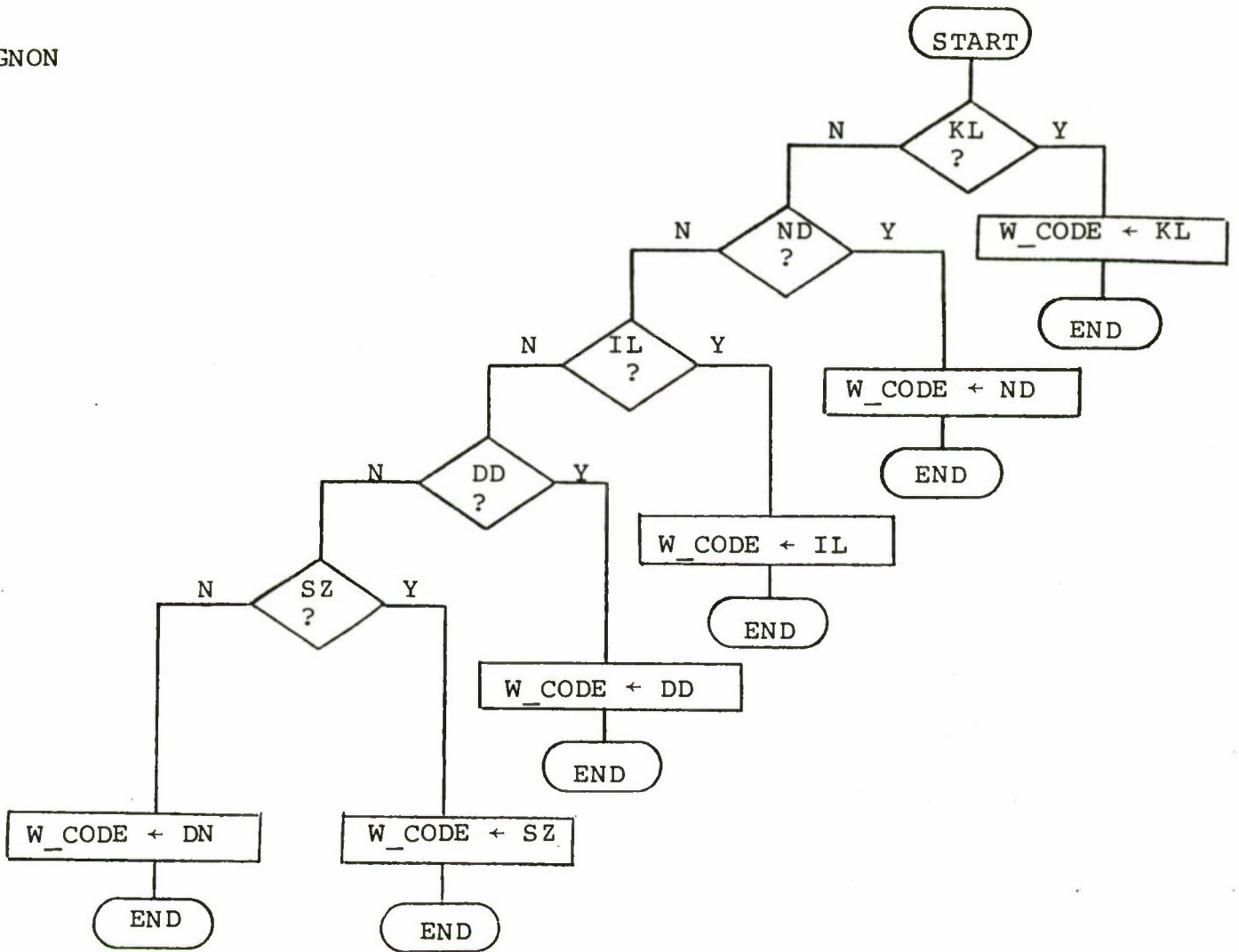
```
KL: (K_CUR_LEVEL ≠ SYS_HI) v (K_CUR_ID ≠ UCP) * It must be user control process *
ND: usent = ∅ * No such user *
IL: usent[MAX_LEVEL] - lv * Error if maximum level doesn't dominate *
DD: (u,*) ∈ D_Q(lv) - | * User is signed on there already *
SZ: s > usent[LIMIT] - usent[SUM] - | * Too much space requested *
```

effects

```
For USER u:
[1] K_CUR_ID ← u * Set identification of K-Variables *
[2] K_CUR_LEVEL ← lv * Set the current level *
[3] K_CUR_QTA ← s * Set session space quota *
[4] ACTIVATE K_CUR_TIME * Activate the timer for this user *
[5] D_Q(lv) ← usent[SUM] - usent[SUM] * Append user entry to sign on list *
[6] usent[SUM] ← usent[SUM] + s * Update user's sum of space used *
[7] K_OPEN ← FALSE * Initialize whole 5 dimensional array *
For UCP process:
[8] W_CODE ← DN * Set code for user control process *
```

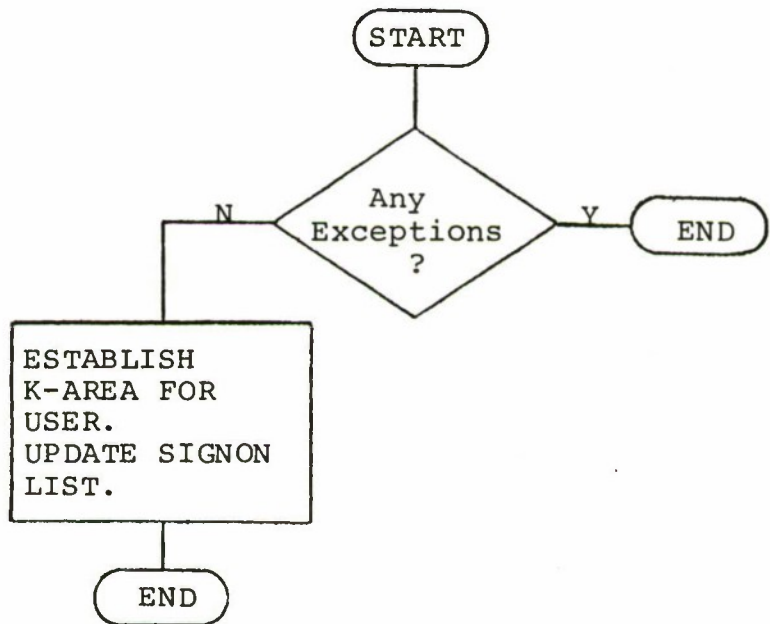
(B) Access table	Variables Observed	Variables Modified	Variables Observed and Modified
		K CUR_ID, K_CUR_LEVEL K_CUR_QTA K_CUR_TIME K_OPEN W_CODE	D Q(1v) D_V(DBA, 'DBA_ULIST')

SIGNON



Subject 2

(on behalf of UCP
at level SYS_HI)



PRIMITIVE: SIGNON CASE: 1 SUBJECT: 1

CONDITIONS: KL
 An attempt to invoke SIGNON by other than the trusted UCP at level SYS_HI.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> u,lv,s,r	W	W_CODE	W
<u>CONSTANTS:</u> UCP,SYS_HI,KL	Unclass		
<u>VARIABLES:</u> K_CUR_ID K_CUR_LEVEL	W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: SIGNON CASE: 2 SUBJECT: 1

CONDITIONS: (~KL) ^ ND
 An attempt to sign-on by a non-registered data base user.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> u,lv,s,r	W	W_CODE	W
<u>CONSTANTS:</u> DBA,'DBA_ULIST','R', UCP,SYS_HI,Ø,ND	Unclass		
<u>VARIABLES:</u> K_CUR_ID K_CUR_LEVEL usent	W W SYS_HI		
HIGHEST LEVEL OBSERVED:	SYS_HI	LOWEST LEVEL MODIFIED:	W

LEMMA: W = SYS_HI

PROOF: ~KL

PRIMITIVE: SIGNON CASE: 3 SUBJECT: 1

CONDITIONS: (~KL) ^ (~ND) ^ IL
 User's maximum level does not dominate
 requested sign-on level.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> u,lv,s,v	W	W_CODE	W
<u>CONSTANTS:</u> DBA,'DBA_ULIST','R', UCP,SYS_HI,Ø,MAX_LEVEL,IL	Unclass		
<u>VARIABLES:</u> K_CUR_ID K_CUR_LEVEL usent	W W SYS_HI		
HIGHEST LEVEL OBSERVED:	SYS_HI	LOWEST LEVEL MODIFIED:	W

LEMMA: W = SYS_HI
 PROOF: ~KL

PRIMITIVE: SIGNON CASE: 4 SUBJECT: 1

CONDITIONS: (~KL) ^ (~ND) ^ (~IL) ^ DD
 User is already signed on at this level.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> u,lv,s,v	W	W_CODE	W
<u>CONSTANTS:</u> DBA,'DBA_ULIST','R', UCP,SYS_HI,Ø,MAX_LEVEL,DD	Unclass		
<u>VARIABLES:</u> K_CUR_ID K_CUR_LEVEL usent D_Q(lv)	W W SYS_HI lv		
HIGHEST LEVEL OBSERVED:	SYS_HI	LOWEST LEVEL MODIFIED:	W

LEMMA: W = SYS_HI
 PROOF: ~KL

PRIMITIVE: SIGNON CASE: 5 SUBJECT: 1

CONDITIONS: (~KL) ^ (~ND) ^ (~IL) ^ (~DD) ^ SZ
 The user attempting to sign on
 has requested too much space.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> u,lv,s,v	W	W_CODE	W
<u>CONSTANTS:</u> DBA,'DBA ULIST','R', UCP,SYS_HI,Ø,MAX_LEVEL,SZ	Unclass		
<u>VARIABLES:</u> K_CUR_ID K_CUR_LEVEL usent D_Q(lv)	W W SYS_HI lv		
HIGHEST LEVEL OBSERVED:	SYS_HI	LOWEST LEVEL MODIFIED:	W

LEMMA: W = SYS_HI
 PROOF: ~KL

PRIMITIVE: SIGNON CASE: 6 SUBJECT: 1

CONDITIONS: (~KL) ^ (~ND) ^ (~IL) ^ (~DD) ^ (~SZ)
 No exception.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> u,lv,s,v	W	W_CODE	W
<u>CONSTANTS:</u> DBA,'DBA ULIST','R', UCP,SYS_HI,Ø,MAX_LEVEL,DN	Unclass		
<u>VARIABLES:</u> K_CUR_ID K_CUR_LEVEL usent D_Q(lv)	W W SYS_HI lv		
HIGHEST LEVEL OBSERVED:	SYS_HI	LOWEST LEVEL MODIFIED:	W

LEMMA: W = SYS_HI
 PROOF: ~KL

PRIMITIVE: SIGNON CASE 1 SUBJECT: 2

CONDITIONS: (~KL) ^ (~ND) ^ (~IL) ^ (~DD) ^ (~SZ)
 No exceptions. Establish the new DMS user.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> u,lv,s,v	W		
<u>CONSTANTS:</u> DBA, 'DBA_ULIST', 'R', SYS_HI, 0, MAX_LEVEL, SUM, DN, UCP, LIMIT	Unclass	K_CUR_ID(user) K_CUR_LEVEL(user) K_CUR_QTA(user) K_CUR_TIME(user) D_Q(lv) usent[SUM]	lv lv lv lv lv SYS_HI
<u>VARIABLES:</u> K_CUR_ID(UCP) K_CUR_LEVEL(UCP) usent D_Q(lv)	W W SYS_HI lv		
HIGHEST LEVEL OBSERVED:	SYS_HI	LOWEST LEVEL MODIFIED:	lv

The UCP must violate the *-property in order to sign users on to the system.

A.1.4.2 SIGNOFF

(A) SIGNOFF

* The user control (and authentication) process will sign the current *
 * off the secure DMS, cleaning everything up. (A possible means of *
 * requesting a SIGNOFF is to turn the terminal off.) *

abbreviation

usent = (D_V(DBA,'DBA_ULIST','R',SYS_HI) {D_F.USERID = K_CUR_ID} * type is user_ent *
 * Select the current user's entry from the DBA's user relation. *

open_obj = LOC_OP * This function locates all TRUE's in the *
 * open table (K_OPEN), and returns 5-tuples *
 * consisting of their co-ordinates. *

exceptions

KL: (K_CUR_LEVEL ≠ SYS_HI) ∨ (K_CUR_ID ≠ UCP) * It must be user control process *

effect

- For USER u:
- [1] D R(¬K RESERVE) ← ∅ * Remove all reservations *
 - [2] D O(¬open_obj) ← ∅ * Remove all open list entries *
 - [3] D Q(¬K CUR_LEVEL) ← ∅ * Remove user entries *
 - [4] usent[SUM] ← ∅ * Return unused space *
 - [5] K CUR_ID, K CUR_LEVEL, K CUR_QTA, K CUR_TIME ← ∅ * Purge everything *
 - [6] K_FACC, K_IACC, K_VACC, K_OPEN, K_RESERVE, K_LACC ← ∅
 - [7] K_FX, K_IX, K_VX, K_FY, K_IY, K_VY, K_FZ, K_IZ, K_VZ, K_LX, K_LY, K_LZ ← ∅

For UCP Process:

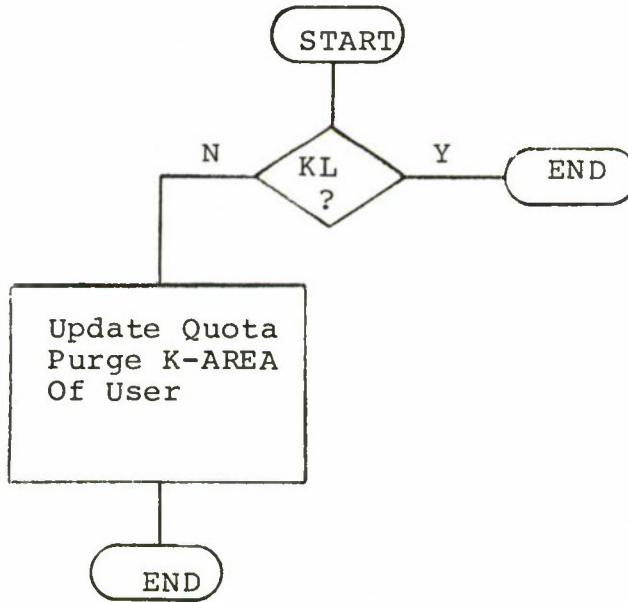
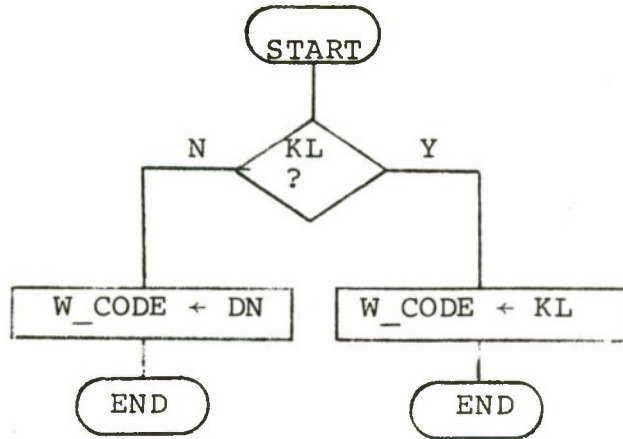
- [8] W_CODE ← DN

(B) Access table

Variables Observed	Variables Modified	Variables Observed and Modified
K CUR TIME	K CUR ID, K CUR LEVEL	K CUR ID, K CUR LEVEL
K_FACC, K_IACC	D_V(DBA, 'DBA_ULIST')	D_V(DBA, 'DBA_ULIST')
K_VACC, K_LACC	K_OPEN, K_RESERVE	K_OPEN, K_RESERVE
K_Fx, K_Ix, K_Vx	D_R(K_RESERVE)	D_R(K_RESERVE)
W_CODE	D_O(open_obj)	D_O(open_obj)
	D_Q(K_CUR_LEVEL)	D_Q(K_CUR_LEVEL)
	K_CUR_QTA	K_CUR_QTA

SIGNOFF

Subject 1
(at level W)



Subject 2
(for UCP at
level SYS_HI)

PRIMITIVE: SIGNOFF CASE: 1 SUBJECT: 1

CONDITIONS: KL
 An attempt to invoke SIGNOFF
 by other than the UCP.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u>	W	W_CODE	W
<u>CONSTANTS:</u> UCP,SYS_HI,KL	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_CUR_ID	W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: SIGNOFF CASE: 2 SUBJECT: 1

CONDITIONS: DN
 No exceptions.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u>	W	W_CODE	W
<u>CONSTANTS:</u> UCP,SYS_HI, DN	Unclass		
<u>VARIABLES:</u> K_CUR_LEVEL K_CUR_ID	W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: SIGNOFF CASE 1 SUBJECT: 2

CONDITIONS: No exceptions. The UCP executes SIGNOFF as the result of an explicit user request, or because of an implicit one, such as turning the terminal off.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u>	W	D_R(+K_RESERVE+)	W
		D_O(+open_obj-)	W
<u>CONSTANTS:</u> UCP, SYS_HI, Ø, DN	Unclass	D_Q(W)	W
		usent	SYS_HI
		K_CUR_ID	W
<u>VARIABLES:</u>		K_CUR_LEVEL	W
D_O(+open_obj-)	*9	K_CUR_QTA	W
K_CUR_ID	W	K_CUR_TIME	W
K_CUR_LEVEL	W	K_OPEN	*
D_Q(W)	W	ACC, X, Y, Z	*
usent	SYS_HI		
K_CUR_QTA	W		
K_OPEN	*		
<u>HIGHEST LEVEL OBSERVED:</u>	SYS_HI	<u>LOWEST LEVEL MODIFIED:</u>	*9

⁹ Any protection level

A.1.4.3 MOVE

(A) O-function MOVE(id,lv)

- * Move the specified object to the given protection level. *
- * Only the DBA may use this primitive, since it potentially *
- * violates the star-property (i.e. modifies "lower" data. *

parameter types

identifier id (OWN,NAM,TYP,LEV), level lv, history (CREATION,USER,MODIFIED)

exceptions

- ND: K CUR_ID ≠ DBA *
 - IL: K_CUR_LEVEL ≠ SYS_HI *
 - NE: D_E(id) = ∅ *
 - IR: (OWN,NAM,TYP,ZERO) ≠ |D_D(LEV) | *
 - DO: D_O(id) ≠ ∅ *
 - DD: (OWN,NAM,TYP,*) ∈ |D_D(lv) | *
- * It must be the DBA
 * User's max must dominate both
 * Object doesn't exist.
 * Registrations are outstanding
 * Object is open (or reserved)
 * Name has already been used

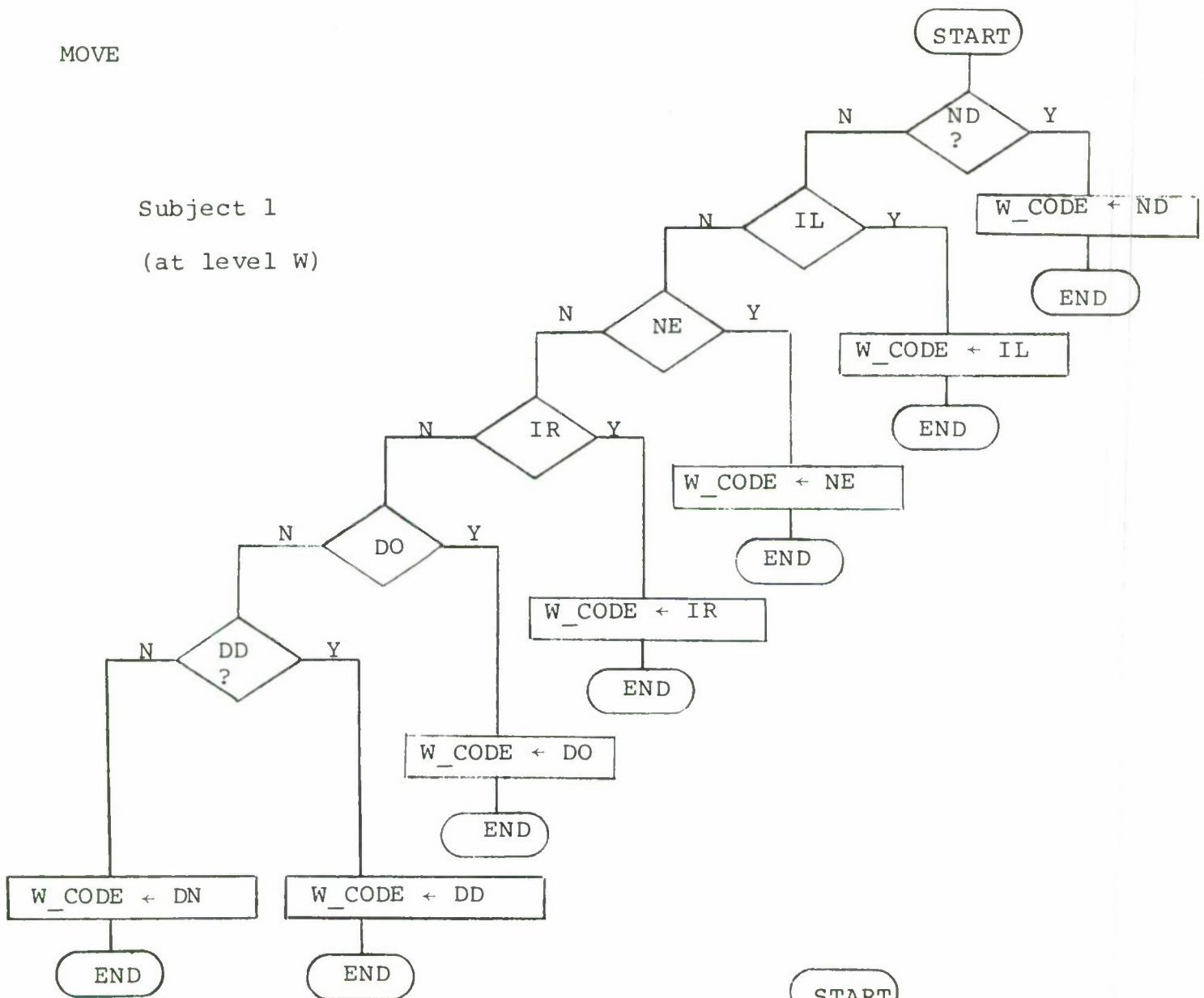
effect

- [1] D_D(lv) ← | ∪ (OWN,NAM,TYP,ZERO) *
- [2] D_E(OWN,NAM,TYP,lv) ← |D_E(id) | *
- [3] D_F(OWN,NAM,TYP,lv) ← |D_F(id) | *
- [4] D_H(OWN,NAM,TYP,lv) ← (|D_H(id)|[CREATION]) |
 * Take creation date, create rest of history. *
- [5] D_M(OWN,NAM,TYP,lv) ← |D_M(id) | *
- [6] D_V(OWN,NAM,TYP,lv) ← |D_V(id) | *
- [7] D_Z(OWN,NAM,TYP,lv) ← |D_Z(id) | *
- [8] D_E(id),D_F(id),D_H(id),D_M(id),D_V(id),D_Z(id) ← ∅ * Purge old object*
- [9] D_D(LEV) ← | - (OWN,NAM,TYP,ZERO) * Remove the directory entry *
- [10] W_CODE ← DN

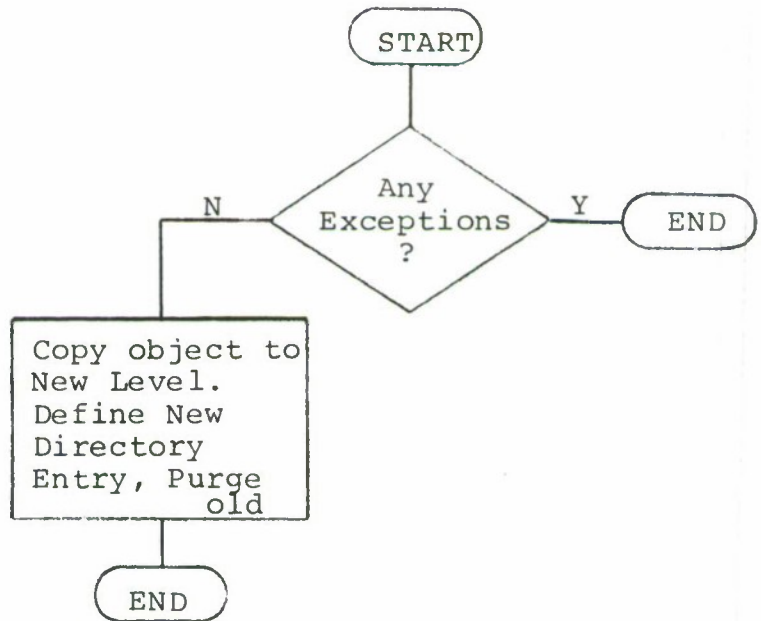
(B) Access table	Variables Observed	Variables Modified	Variables Observed and Modified
	K_CUR_ID, K_CUR_LEVEL D_V(DBA, 'DBA_ULIST')	W_CODE	D D(lv) D_D(LEV) D_E(id), D_F(id) D_H(id), D_M(id) D_V(id), D_Z(id)

MOVE

Subject 1
(at level W)



Subject 2
(at level SYS_HI)



PRIMITIVE: MOVE

CASE: 1

SUBJECT: 1

CONDITIONS: ND

User is not the DBA.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id,lv	W	W_CODE	W
<u>CONSTANTS:</u> DBA,ND	Unclass		
<u>VARIABLES:</u> K_CUR_ID	W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: MOVE

CASE: 2

SUBJECT: 1

CONDITIONS: (~ND) ^ IL

DBA is not signed on at the system high level.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id,lv	W	W_CODE	W
<u>CONSTANTS:</u> DBA,SYS_HI,IL	Unclass		
<u>VARIABLES:</u> K_CUR_ID K_CUR_LEVEL	W W		
HIGHEST LEVEL OBSERVED:	W	LOWEST LEVEL MODIFIED:	W

PRIMITIVE: MOVE CASE: 5 SUBJECT: 1

CONDITIONS: (~ND) ^ (~IL) ^ (~NE) ^ (~IR) ^ DO
Object is open at present.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id,lv	W	W_CODE	SYS_HI
<u>CONSTANTS:</u> DBA_SYS_HI,Ø,ZERO,DO	Unclass		
<u>VARIABLES:</u> K_CUR_ID K_CUR_LEVEL D_E(id) D_D(lv) D_O(id)	SYS_HI SYS_HI lv lv lv		
HIGHEST LEVEL OBSERVED:	SYS_HI	LOWEST LEVEL MODIFIED:	SYS_HI

PRIMITIVE: MOVE CASE: 6 SUBJECT: 1

CONDITIONS: (~ND) ^ (~IL) ^ (~NE) ^ (~IR) ^ (~DO) ^ DD
The object already exists at the new level.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id,lv	W	W_CODE	SYS_HI
<u>CONSTANTS:</u> DBA,SYS_HI,Ø,ZERO,DD	Unclass		
<u>VARIABLES:</u> K_CUR_ID K_CUR_LEVEL D_E(id) D_D(lv) D_O(id)	SYS_HI SYS_HI lv lv lv		
HIGHEST LEVEL OBSERVED:	SYS_HI	LOWEST LEVEL MODIFIED:	SYS_HI

PRIMITIVE: MOVE CASE 1 SUBJECT: 2

CONDITIONS: (~ND) ^ (~IL) ^ (~NE) ^ (~IR) ^ (~DO) ^ (~DD)
 No exceptions. Copy object to new level;
 append entry to new directory and delete
 entry from old.

OBSERVED	LEVEL	MODIFIED	LEVEL
<u>PARAMETERS:</u> id,lv	W	D_D(lv)	lv
		D_D(new-lv)	lv
<u>CONSTANTS:</u> DBA,SYS_HI,∅,ZERO,DN	Unclass	D_E(id)	lv
		D_F(id)	lv
		D_H(id)	lv
		D_M(id)	lv
<u>VARIABLES:</u> K_CUR_ID	SYS_HI	D_V(id)	lv
K_CUR_LEVEL	SYS_HI	D_Z(id)	lv
D_E(id)	lv	D_E(new-id)	lv
D_D(lv)	lv	D_F(new-id)	lv
D_O(id)	lv	D_H(new-id)	lv
D_F(id)	lv	D_M(new-id)	lv
D_H(id)	lv	D_V(new-id)	lv
D_M(id)	lv	D_Z(new-id)	lv
D_V(id)	lv		
D_Z(id)	lv		
K_CUR_TIME	SYS_HI		
HIGHEST LEVEL OBSERVED:	SYS_HI	LOWEST LEVEL MODIFIED:	lv

This function violates the *-property in a controlled manner since only the DBA may reclassify information. It does not violate the SS-property because the DBA must have a system high current level.

APPENDIX II
RESULTS OF THE VALIDATION

The course of the validation revealed certain problems in the specification. These problems and their solution are summarized in table A.2.1.

Table A.2.1 Problems with the Specifications

<u>Problem Description</u>	<u>Problem Type</u>	<u>Solution</u>	<u>Functions Affected</u>
1. '*'-property violations	Design error	Define new entities: K_LACC, K_LX, K_LY, K_LZ at_level = K_CUR_LEVEL	DKD, DKE, DKH, DKM, DKQ, DKR, DKV, DKZ, LIST_DOWN, ASSIGN, KWA
2. Accumulator could not be "lower" than parameter data	Design error	Set minimum kernel level = K_CUR_LEVEL	Same as 1.
3. "NO" exceptions for "higher" data provided a communication path down	Design error	Purge accumulator if exception is TRUE	DKE, DKH, DKM, DKR, DKV, DKZ
4. Using quota at "higher" levels constituted a "write-down" to K_CUR_QTA	Design error	Restrict K_CUR_QTA modification to current level.	INIT, KDZ
5. Failed to check for a string (i.e. single tuple). Also, K_CUR_ID, K_CUR_LEVEL omitted from access table	Simple omission	Include them	WKB
6. Return code condition and asterisks on exception codes were omitted	Simple omission	Include them	INIT
7. ND exception was redundant (already in Unique_keys of IV exception)	Logic error	Remove ND	APPEND, Unique_keys

(Table A.2.1 continued...)

<u>Problem Description</u>	<u>Problem Type</u>	<u>Solution</u>	<u>Functions Affected</u>
8. Exceptions IT and IL were in the wrong order	Logic error	Switch them	APPEND
9. There was no check that SIGNON and SIGNOFF were executed by UCP.	Simple omission	Include check	SIGNON, SIGNOFF
10. Superfluous format check	Logic error	Remove it	KDZ
11. Return code condition excluded a check for kernel to kernel transfer. Asterisk on W_CODE was omitted	Serious omission Simple omission	Include it Include it	ASSIGN
12. Incorrect (obsolete) OPEN ARRAY test for "NO" exception.	Oversight	Replace in specs: *NO: K_OPEN(id,EXPM)=FALSE *NO: K_OPEN(id,STOR)=FALSE	KDM KDV,WDV

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