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²⁰ components, where $M=1$ neither implies nor is implied by ergodicity of \underline{S} , such that conditional state distributions within each component geometrically approach an initial-value-independent process in the manner described above, and one or more equivalent components eventually dominate the others. A method for drift-free finite-memory approximation (or realization) of this process is also introduced.

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STATE-ESTIMATION OF PARTIALLY-OBSERVED MARKOV CHAINS:
DECOMPOSITION, CONVERGENCE, AND COMPONENT IDENTIFICATION***

by

Loren K. Platzman

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STATE-ESTIMATION OF PARTIALLY-OBSERVED MARKOV CHAINS:
DECOMPOSITION, CONVERGENCE, AND COMPONENT IDENTIFICATION

BY

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SUMMARY

A partially-observed Markov chain $(\underline{S}, \underline{Y})$ consists of an N -state Markov chain \underline{S} , along with a process \underline{Y} of noisy observations of the transitions of \underline{S} . A metric on stochastic N -vectors and a generalized ergodic coefficient on the transition probability matrices of $(\underline{S}, \underline{Y})$ are defined, resulting in a notion (similar to weak ergodicity) of deteriorating dependence on initial value in a process of distributions of the state (of \underline{S}) conditioned on past observations. If $(\underline{S}, \underline{Y})$ is stationary, then \underline{S} may be decomposed into $M \leq N$ components, where $M=1$ neither implies nor is implied by ergodicity of \underline{S} , such that conditional state distributions within each component geometrically approach an initial-value-independent process in the manner described above, and one or more equivalent components eventually dominate the others. A method for drift-free finite-memory approximation (or realization) of this process is also introduced.

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1. Introduction

Consider a probability space $(\Omega, \underline{F}, \underline{P})$. For ease of notation, random variables (i.e., symbols that implicitly denote measurable functions of (Ω, \underline{F})) will be shown in boldface.

Let $\underline{S} = \{\underline{s}(k) \in S : k \in I\}$ and $\underline{Y} = \{\underline{y}(k) \in Y : k \in I\}$ be stochastic processes where $S = \{1, \dots, N\}$, Y is a finite set, and $I = \{\dots, -1, 0, 1, \dots\}$. Also let the following "splitting" condition be satisfied:

For any $k \in I$, the "past" $\{\underline{s}(k'), \underline{y}(k') : k' < k\}$
and the "future" $\{\underline{y}(k'-1), \underline{s}(k') : k' > k\}$
are (conditionally) independent given $\underline{s}(k)$.

Then \underline{S} is a finite-state Markov chain (FSMC) (the state process), and \underline{Y} is a process of noisy observations of the transitions of \underline{S} (the observation process). The pair $(\underline{S}, \underline{Y}) = \{(\underline{s}(k), \underline{y}(k)) : k \in I\}$ is called a partially-observed Markov chain (POMC).

Since $(\underline{S}, \underline{Y})$ is itself a FSMC, the probability measure \underline{P} may be completely specified by a finite array of time-indexed transition probabilities. We will adopt a notation associated with probabilistic automata, $\{P^k(\underline{y}) : \underline{y} \in Y, k \in I\}$, where $P^k(\underline{y}) = [P_{ij}^k(\underline{y})]$ is an $N \times N$ matrix having entries

$$(1.1) \quad P_{ij}^k(\underline{y}) = \Pr(\underline{s}(k+1)=j, \underline{y}(k)=\underline{y} \mid \underline{s}(k)=i) \quad i, j \in S, \underline{y} \in Y, k \in I.$$

Thus $P^k(y)$ is a substochastic matrix and $\sum_{y \in Y} P^k(y)$ is a stochastic matrix, the transition probability matrix of \underline{S} .

Consider an observer who wishes to characterize the "future" $\{\underline{s}(k'), \underline{y}(k') : k' \geq k\}$, given a finite string of most recent observations

$$(1.2) \quad \underline{y}(k_0, k) = \underline{y}(k_0) \underline{y}(k_0+1) \dots \underline{y}(k-1).$$

This is accomplished by computing the horizontal N-vector $\underline{n}(k_0, k)$ having entries

$$(1.3) \quad \underline{n}_i(k_0, k) = \Pr\{\underline{s}(k)=i \mid \underline{y}(k_0, k)\} \quad i \in S$$

and lying in the simplex

$$(1.4) \quad \Pi = \{\tau : \tau_i \geq 0 \forall i \in S, \sum_{i \in S} \tau_i = 1\}.$$

Clearly, $\underline{n}(k_0, k)$ is the a posteriori probability distribution of $\underline{s}(k)$ given observations $\underline{y}(k_0, k)$. Moreover, in the sense that each state in S is associated with an extreme point of Π , $\underline{n}(k_0, k)$ is a conditional state expectation, and hence a least-squares state estimate, at time k , given observations $\underline{y}(k_0, k)$.

Let $v = (1, 1, \dots, 1)^t$ and define

$$(1.5) \quad T(\tau, P) = \tau P / \tau P v \quad \tau P v > 0.$$

The state estimate may now be recursively updated according to the identity:

$$(1.6) \quad \underline{\eta}(k_0, k+1) = T(\underline{\eta}(k_0, k), P^k(\underline{y}(k))).$$

We note that $\{\underline{\eta}(k_0, k) : k=k_0, k_0+1, \dots\}$ is therefore a Markov process and may, itself, be viewed as the "state" in an alternate realization of \underline{X} .

Multiple-step transitions are similarly expressed.

Let $z = y_1, \dots, y_n$ and define the matrix product:

$$(1.7) \quad P^k(z) = P^k(y_1) \cdot P^{k+1}(y_2) \cdot \dots \cdot P^{k+n-1}(y_n).$$

It is easily seen that

$$(1.8) \quad \underline{\eta}(k_0, k+n) = T(\underline{\eta}(k_0, k), P^k(\underline{y}(k, k+n))).$$

We also define a process of state estimate approximations constructed as in (1.8), but from an initial value \underline{v} :

$$(1.9) \quad \underline{\eta}(\underline{v}; k_0, k) = T(\underline{v}, P^k(\underline{y}(k_0, k))).$$

For any $k \in I$, the process $\{\underline{\eta}(k-n, k) : n = 0, 1, \dots\}$ is a uniformly bounded martingale. Hence, by Proposition V-2-6 of Neveu [9], there exists a vector $\underline{\eta}(k)$ satisfying

$$(1.10) \quad \underline{\eta}(k) = \lim_{k_0 \rightarrow -\infty} \underline{\eta}(k_0, k) \quad \text{a.s.}$$

Since $T(\cdot, P)$ is continuous, (1.8) and (1.9) yield

$$(1.11) \quad \underline{\eta}(k) = T(\underline{\eta}(k_0), P^{k_0}(\underline{y}(k_0, k))) = \underline{\eta}(\underline{\eta}(k_0); k_0, k).$$

This paper examines the evolution of the state estimate approximation process $\{\underline{\eta}(\underline{x}; k_0, k) : k=k_0, k_0+1, \dots\}$. Specifically, we establish conditions under which $\underline{\eta}(\underline{x}; k_0, k) \rightarrow \underline{\eta}(k)$, in some sense, as $k \rightarrow \infty$. To this end, a metric Δ on \mathbb{R}^n and a generalized ergodic coefficient α on substochastic matrices are defined, in Sections 3 and 4, respectively, so that

$$(1.12) \quad \frac{1}{2} \sum_{i \in S} |v_i - v'_i| \leq \Delta[v, v'] \leq 1,$$

$$(1.13) \quad \alpha[PQ] \leq \alpha[P] \alpha[Q] \leq 1,$$

$$(1.14) \quad \Delta[T(v, P), T(v', P)] \leq \alpha[P] \Delta[v, v'].$$

It follows that $\Delta[\underline{\eta}(\underline{x}; k_0, k), \underline{\eta}(k)] \leq \sum_{k'=k_0}^{k-1} \alpha[P^{k'}(\underline{y}(k'))]$, and that both sides of the inequality are nonincreasing in k , a.s. Sufficient conditions for this bound to vanish as $k \rightarrow \infty$ are given by Theorems (5.3) and (5.4). Theorem (6.10) shows that if $(\underline{S}, \underline{Y})$ is stationary and certain conditions (concerning future-dependence of \underline{x} and state-minimality of $(\underline{S}, \underline{Y})$) are satisfied, then

$$(1.15) \quad \lim_{k \rightarrow \infty} \sum_{i \in S} |\eta_i(\underline{x}; k_0, k) - \eta_i(k)| = 0 \quad \text{a.s.}$$

although it may simultaneously hold that

$$(1.16) \quad \Delta[\underline{\eta}(\underline{x}; k_0, k), \underline{\eta}(k)] = 1 \quad \text{a.s., } k=k_0, k_0+1, \dots$$

Section 6 also introduces a stochastic process $\{\underline{\Sigma}(k) : k \in I\}$ which, under these conditions, displays the following properties (made explicit by (6.1), (6.2) and (6.7) respectively):

- (a) Decomposition: $\underline{I}(k)$ is a collection of disjoint subsets of S and $\underline{g}(k)$ is contained by some element $\underline{\sigma}(k)$ of $\underline{I}(k)$ a.s.
- (b) Recursive Calculability: $\underline{\sigma}(k)$ may be expressed as a (deterministic) function of $\underline{\sigma}(k_0)$ and $\underline{y}(k_0, k)$.
- (c) Convergence: $\Delta[\underline{\eta}(\underline{x}; k_0, k), \underline{\eta}(k)] \rightarrow 0$, geometrically in k , as $k \rightarrow \infty$, provided that $\{i : \underline{x}_i > 0\} \subseteq \underline{\sigma}(k_0)$.

Thus $\underline{\eta}(\underline{x}; k_0, k)$ may be decomposed into $M = \# \underline{I}(k_0)$ components, one of which may converge (in Δ) to $\underline{\eta}(k)$ as $k \rightarrow \infty$. The convergent component is identified by maximum likelihood on the basis of past or future observations. A necessary and sufficient condition for $\{\underline{\eta}(k)\}$ itself to be a FSMC is given by Theorem (6.14).

In analogy to (finite-dimensional) linear system theory, the ergodic classes of \underline{S} in S may be likened to controllable subspaces, whereas the subsets $\underline{I}(k_0)$ of S may be likened to observable subspaces. It is consistent with this analogy that $M = \# \underline{I}(k) = 1$ neither implies, nor is implied by, ergodicity of \underline{S} .

POMCs have been studied by many authors, notably Blackwell [2], Drake [5] Astron [1], Sondik [13], (also see [12]), and Rudemo [11]. Using a result of Furstenberg and Kesten [6], Kaijser [8] has obtained a weaker version of Theorem (5.3), and conjectures Theorem (6.10). A similar result has been

derived by Devore [4]. Ergodic coefficients are discussed at length in Chapter 5 of Isaacson and Madsen [7]. A POMC is also a free probabilistic automaton; see Paz [10].

The analysis given here may be generalized to processes having denumerable state-output sets, to certain continuous-time processes (following Rudemo [11]), and to semi-Markov processes (following White [14]).

2. Notations

Sets: If A and B are sets, then $A-B$ is the set of elements in A that are not contained in B. $\#A$ is the number of elements in A. \emptyset is the null set.

Strings: $Y^k(n)$ is the set of strings $z = y_1 y_2 \dots y_n$ satisfying $p_k(z) \neq 0$. $Y^k(*) = \bigcup_{n=1}^{\infty} Y^k(n)$. If $z = y_1 \dots y_n$ and $\tilde{z} = \tilde{y}_1 \dots \tilde{y}_m$ then $z\tilde{z}$ denotes the concatenation $y_1 \dots y_n \tilde{y}_1 \dots \tilde{y}_m$.

Scalars: $a \dagger b$ denotes integer quotient rounded down; i.e., $q = a \dagger b$ is the integer of largest magnitude satisfying $|qb| \leq |a|$ and $\text{sgn}(q) = \text{sgn}(a/b)$. Also $(a)^+ = \max(0, a)$.

Vectors: R_N is the Euclidean space of vertical N-vectors.

For $\sigma \in S$, e^σ is a row N-vector with entries $e_i^\sigma = 1$ if $i = \sigma$, $e_i^\sigma = 0$ otherwise. For $i \in S$, e^i denotes the transpose of $e^{i\}$, the i^{th} unit stochastic vector in Π .

Faces of Π : For any subset σ of S , $\Pi(\sigma)$ denotes the set of stochastic vectors τ satisfying:

$$\tau_i > 0 \iff i \in \sigma, \quad \forall i \in S.$$

Bayes' Operator: For $\tau \in \Pi$, $w \in \mathbb{R}_M^+$, with $w_i > 0 \forall i \in S$ and $\tau w > 0$, let $\tau \circ w$ denote the vector in Π having entries

$$(\tau \circ w)_i = \tau_i w_i / \tau w.$$

This may be interpreted as follows: Let τ represent a priori probabilities of some random variable, \underline{s} , on sample space S . Consider an event A occurring with conditional probability w_i provided that $\underline{s} = i$. Then $\tau \circ w$ is the vector of a posteriori probabilities of random variable \underline{s} , given A .

Stochastic Vector Restriction: For $\tau \in \Pi$ and $\sigma \subseteq S$ with $\sum_{i \in \sigma} \tau_i > 0$, define

$$[\tau | \sigma] = \tau \circ e^\sigma.$$

Subrectangularity: An $N \times N$ substochastic matrix P is M-subrectangular if there exist collections $\{I_1, \dots, I_M\}$ and $\{J_1, \dots, J_M\}$, of mutually disjoint, nonempty subsets of S such that

$$\{(i, j) : P_{ij} > 0\} = (I_1 \times J_1) \cup \dots \cup (I_M \times J_M).$$

A 1-subrectangular matrix is also called subrectangular.

State Estimate Approximations: $H(k)$ denotes the set of random variables \underline{r} having values in \mathbb{R} such that \underline{r} is (conditionally) independent of the "future" $\{\underline{y}(k'-1), \underline{z}(k') : k' > k\}$ given $\underline{z}(k)$, and $\underline{r}_S(k) > 0$, a.s. $H(k)$ contains each of the following: v/N , $\underline{n}(k_0, k)$, $\underline{n}(k)$, $e^{\underline{z}(k)}$, and $\underline{n}(\underline{r}; k_0, k)$ if $\underline{r} \in H(k_0)$.

Stopping Times: Let $\underline{Y}(k_0, k)$ denote the sigma-algebra generated by $\{\underline{y}(k_0, k)\}$. Then $\theta_f(k_0)$ is the set of random variables \underline{n} having positive integer values and satisfying $\{\underline{n} = k - k_0\} \in \underline{Y}(k_0, k) \forall k > k_0$; and $\theta_b(k)$ is the set of random variables \underline{n} having positive integer values and satisfying $\{\underline{n} = k - k_0\} \in \underline{Y}(k_0, k) \forall k_0 < k$.

3. A Metric on \mathbb{R}

This section introduces a metric that is used to measure the "closeness" of state estimate approximations.

(3.1) Definition: For $\tau, \tau' \in \mathbb{R}$, define

$$(a) \quad |\tau - \tau'| = \sum_{i \in S} |\tau_i - \tau'_i|;$$

$$(b) \quad \delta(\tau, \tau') = \sum_{i \in S} (\tau_i - \tau'_i)^+;$$

$$(c) \quad \Delta(\tau, \tau') = \sup\{\delta(\tau_0 w, \tau'_0 w) : w \in \mathbb{R}_N, w_i \geq 0 \forall i \in S, \tau w > 0, \tau' w > 0\}.$$

(3.2) Lemma: $|\cdot - \cdot|$, δ and Δ are metrics on \mathbb{H} , and

$$0 \leq \frac{1}{2}|\pi - \pi'| = \delta[\pi, \pi'] \leq \Delta[\pi, \pi'] \leq 1.$$

Proof: Trivial.

(3.3) Proposition: (Evaluation of Δ). For $\pi, \pi' \in \mathbb{H}$, define:

$$c_1 = \min\{\pi'_1/\pi_1 : \pi_1 > 0\};$$

$$c_2 = \min\{\pi_1/\pi'_1 : \pi'_1 > 0\}.$$

Then

$$\Delta[\pi, \pi'] = \frac{1 - \sqrt{c_1 c_2}}{1 + \sqrt{c_1 c_2}}.$$

Proof: The proof is given in Section 3'.

The metric δ , also known as the Hajnal measure, has many applications in the theory of ergodic Markov chains [7]. Informally, $\delta[\pi, \pi']$ is the (minimal) "quantity of probability" that would have to be "reassigned" in order to transform probability distribution π into probability distribution π' . Similarly, $\Delta[\pi, \pi']$ is the least upper bound on the quantity of conditional probability by which π and π' might differ if they were conditioned on identical observations.

The distinction between δ and Δ is also illuminated by an examination of the topologies they induce on \mathbb{I} : the topology induced by δ is connected, but Δ causes \mathbb{I} to be separated into the 2^N-1 faces $\{\mathbb{I}(\sigma), \emptyset \subset \sigma \subseteq S\}$.

In the computation of state estimates, this separation can be significant. Consider a POMC in which $\underline{\eta}(0) = (1-\epsilon, \epsilon)$, $\epsilon \ll 1$, but it is desired to approximate $\underline{\eta}(0)$ by e^1 . In a δ sense, $\underline{\eta}(0)$ is "near" the approximation e^1 ; the unconditional expectation of a bounded function of the initial state will be little affected by this approximation. Suppose, however, that every output that subsequently evolves corresponds to

transition probabilities $Q = \begin{bmatrix} .1 & 0 \\ 0 & .9 \end{bmatrix}$. Now $\underline{\eta}(\underline{\eta}(0); 0, k) = T(\underline{\eta}(0), [Q]^k)$

tends to e^2 ; yet if the approximation $\underline{\eta}(0) = e^1$ is used, then $\underline{\eta}(e^1; 0, k) = T(e^1, [Q]^k) = e^1$ is obtained.

Thus an initial error, of δ -sense magnitude $\epsilon \ll 1$, may lead to an eventual error of δ -sense magnitude close to 1, albeit after conditioning on an event of probability less than ϵ . Proposition (4.3) will show that such increase in Δ -sense approximation error cannot occur.

3'. Proof of Proposition (3.3)

Clearly $\Delta(v, v) = 0$. If $\{i : v_i > 0\} \neq \{i : v'_i > 0\}$ then $\Delta(v, v') = 1$.

To see this, assume without loss of generality that $i \in S$ is such that $v_i > 0$ and $v'_i = 0$, and define $w^m = (1 - 1/m)e^i + (1/m)v/N$.

Then $\{w^m\}$ is a sequence in \mathbb{R}^N for which $\lim_{m \rightarrow \infty} \delta(v \cdot w^m, v' \cdot w^m) = 1$; since $(v \cdot w^m)_i \rightarrow 1$ as m increases, but $(v' \cdot w^m)_i = 0$.

By (3.2), the sequence $\{w^m\}$ is supremal.

The case $\{i : v_i > 0\} = \{i : v'_i > 0\}$, $v \neq v'$, remain.

Assume without loss of generality that $v_i > 0$ and $v'_i > 0$, $\forall i \in S$.

Now $0 < c_1 < 1$ and $0 < c_2 < 1$; hence $0 < c_1 < c_2^{-1} < \infty$. Define:

$$\begin{aligned} \Delta_\zeta(v, v') &= \sup\{\delta(v \cdot w, v' \cdot w) : w \in \mathbb{R}_+^N, w_i > 0 \forall i \in S, v' \cdot w / v \cdot w = \zeta\} \\ &= \sup\{ \sum_{i \in S} v_i w_i - \frac{v'_i w_i}{\zeta} : w \in \mathbb{R}_+^N, w_i > 0 \forall i \in S, v \cdot w = 1, v' \cdot w = \zeta \} \end{aligned}$$

which exists for all $c_1 \leq \zeta \leq c_2^{-1}$. Clearly

$$\Delta(v, v') = \sup\{\Delta_\zeta(v, v') : c_1 \leq \zeta \leq c_2^{-1}\}.$$

But $\Delta_\zeta(v, v')$ may be expressed as the solution of a linear program

$$\Delta_\zeta(v, v') = \left[\begin{array}{ll} \text{max:} & v \cdot w \\ \text{subject to:} & v \cdot w = 1 \\ & \tilde{v}' \cdot w = 1 \\ & w \geq 0 \end{array} \right]$$

where $a_1 = (v_i - \tilde{v}'_i)^+$, $\tilde{v}'_i = v'_i / \zeta$. Any optimal basic w that solves this linear program has at most two non-zero entries. Let their indices be

denoted (i,j) and assume without loss of generality that

$$(i,j) \in \Lambda = \{(i,j) : (\bar{v}_i/\bar{v}_1) < (\bar{v}'_j/\bar{v}_j)\}.$$

The optimal index pair (i,j) satisfies $a_i > 0$ and $a_j = 0$;

for otherwise one of the following contradictions is implied

$$(i) \quad a_i = 0, a_j = 0 \longrightarrow \Delta_{\zeta}[\bar{v}, \bar{v}'] = 0;$$

$$(ii) \quad a_i > 0, a_j > 0 \longrightarrow \Delta_{\zeta}[\bar{v}, \bar{v}'] = a_i \bar{v}_1 + a_j \bar{v}_j = (\bar{v}_1 - \bar{v}'_1) \bar{v}_1 + (\bar{v}_j - \bar{v}'_j) \bar{v}_j = 1 - 1 = 0; \text{ or}$$

$$(iii) \quad a_i = 0, a_j > 0 \longrightarrow (i,j) \notin \Lambda.$$

Hence ζ is such that $(\bar{v}'_i/\bar{v}_i) \leq \zeta \leq (\bar{v}'_j/\bar{v}_j)$. The basic feasible solution with indices (i,j) is now seen to take the form:

$$w_i^* = \frac{\bar{v}'_j - \bar{v}_j}{\bar{v}_1 \bar{v}'_j - \bar{v}_j \bar{v}'_1} \geq 0, \quad w_j^* = \frac{\bar{v}_1 - \bar{v}'_1}{\bar{v}_1 \bar{v}'_j - \bar{v}_j \bar{v}'_1} \geq 0.$$

Now

$$\Delta_{\zeta}[\bar{v}, \bar{v}'] = \max \{ \Delta_{\zeta, i, j}[\bar{v}, \bar{v}'] : (\bar{v}'_i/\bar{v}_i) \leq \zeta \leq (\bar{v}'_j/\bar{v}_j) \}$$

where

$$\Delta_{\zeta, i, j}[\bar{v}, \bar{v}'] = a_i w_i^* = \frac{\bar{v}_1 \bar{v}'_j + \bar{v}_j \bar{v}'_1 - \zeta \bar{v}_1 \bar{v}_j - \zeta^{-1} \bar{v}'_1 \bar{v}'_j}{\bar{v}_1 \bar{v}'_j - \bar{v}_j \bar{v}'_1}.$$

Consequently

$$\begin{aligned} \Delta(w, w') &= \sup\{\Delta_{\zeta}(w, w') : c_1 \leq \zeta \leq c_2^{-1}\} \\ &= \sup\{\Delta_{\zeta, 1, j}(w, w') : (1, j) \in A, (w'_1/w_1) \leq \zeta \leq (w'_j/w_j)\} \\ &= \max_{(1, j) \in A} \sup\{\Delta_{\zeta, 1, j}(w, w') : (w'_1/w_1) \leq \zeta \leq (w'_j/w_j)\}. \end{aligned}$$

Since $\Delta_{\zeta, 1, j}(w, w')$ is concave in ζ , it achieves a unique maximum at

$$\zeta = \zeta^* = \sqrt{\frac{w'_1 w'_j}{w_1 w_j}}. \quad \text{Thus}$$

$$\Delta(w, w') = \max_{(1, j) \in A} \{\Delta_{\zeta^*, 1, j}(w, w')\}$$

$$= \max_{(1, j) \in A} \left\{ \frac{1 - \sqrt{\frac{w'_1 w'_j}{w_1 w_j}}}{1 + \sqrt{\frac{w'_1 w'_j}{w_1 w_j}}} \right\}$$

$$= \frac{1 - \sqrt{c_1 c_2}}{1 + \sqrt{c_1 c_2}}$$

4. The Generalized Ergodic Coefficient

It is well known that if P is a stochastic matrix and

$$\alpha[P] = \max_{i,j \in S} \delta[e^i P, e^j P] < 1$$

then, for any $v, v' \in \mathbb{R}$,

$$\delta[vP, v'P] \leq \alpha[P] \delta[v, v'].$$

i.e., the transformation $f[v] = vP$ is a contraction in \mathbb{R} .

One consequence of this property of P is that $\{v(P)^n\}$ approaches a unique limit as $n \rightarrow \infty$. The rate of convergence $\alpha[P]$ is called the ergodic coefficient of the stochastic matrix P . This section generalizes the concept of an ergodic coefficient to substochastic matrices; a corresponding contraction property of $T(\cdot, P)$ is also established.

(4.1) Definition: If P is a nonzero substochastic matrix, then define

$$\alpha[P] = \max\{\Delta[T(e^i, P), T(e^j, P)] : e^i P \neq 0, e^j P \neq 0\}$$

Remark: The evaluation of $\alpha[P]$ by (3.3) requires N^3 operations. This is comparable to the effort expended when multiplying two $N \times N$ matrices.

The generalized ergodic coefficient $\alpha[P]$ has the following properties:

- (4.2) Lemma: (a) $0 \leq \alpha[P] \leq 1$ for all substochastic matrices $P \neq 0$.
(b) $\alpha[P] < 1 \iff P$ is subrectangular.
(c) $\alpha[P] = 0 \iff \text{rank}[P] = 1$.

Proof: Trivial.

- (4.3) Proposition: (Contraction Property of T)

$$\Delta[T(n,P), T(n',P)] \leq \alpha[P] \Delta[n, n'], \quad nP \neq 0, n'P \neq 0.$$

Proof: The proof is given in Section 4'.

- (4.4) Corollary: $\alpha[P] = \sup \{ \Delta[T(n,P), T(n',P)] : nP \neq 0, n'P \neq 0 \}$.

- (4.5) Corollary: $\alpha[PQ] \leq \alpha[P] \alpha[Q]$.

This analysis may be modified to obtain a generalized ergodic coefficient for M-subrectangular matrices.

(4.6) Definition: If P is a nonzero substochastic matrix, define:

$$\alpha_M(P) = \left\{ \begin{array}{ll} \max \{ \Delta(T(e^i, P), T(e^j, P)) : P_{jk} > 0 \text{ and} \\ P_{jk} > 0, \text{ some } k \in S \}, & \text{if } P \text{ is } M\text{-subrectangular} \\ 1, & \text{otherwise} \end{array} \right.$$

Remark: $\alpha(P) = \alpha_1(P)$.

(4.7) Lemma: If P^1, \dots, P^M are subrectangular and $P = \sum_{i=1}^M P^i$ is M -subrectangular, then $\alpha_M(P) = \max_i \{\alpha(P^i)\}$.

(4.8) Corollary:

- (a) $0 \leq \alpha_M(P) \leq 1$
- (b) $\alpha_M(P) < 1 \iff P$ is M -subrectangular.
- (c) $\alpha_M(P) = 0 \iff \text{rank}[P] = M$
and P is M -subrectangular.

Proof: Trivial; compare (4.2).

(4.9) Corollary: If PQ is M -subrectangular then

$$\alpha_M(PQ) \leq \alpha_M(P) \alpha_M(Q).$$

Proof: Trivial; compare (4.5)

4'. Proof of Proposition (4.3)

We require the following well-known Lemma:

(4'.1) Lemma: If $w \in \mathbb{R}_n^+$ and $v, v' \in \mathbb{R}^n$ then $|vw - v'w|$
 $\leq \delta(v, v') \cdot \max_{1,1' \in S} (w_1 - w_{1'})$.

Proof: By (3.2), $\sum_{j \in S} (v_j - v'_j)^+ = \sum_{j \in S} (v'_j - v_j)^+$.

Now $wv - v'w = \sum_{j \in S} (v_j - v'_j) w_j = \sum_{j \in S} [(v_j - v'_j)^+ - (v'_j - v_j)^+] w_j$
 $\leq \{[\sum_{j \in S} (v_j - v'_j)^+] \cdot [\max_{1 \in S} w_1]\} - \{[\sum_{j \in S} (v'_j - v_j)^+] \cdot [\min_{1' \in S} w_{1'}]\}$
 $= \delta(v, v') \cdot [\max_{1 \in S} w_1 - \min_{1' \in S} w_{1'}] = \delta(v, v') \cdot [\max_{1,1' \in S} w_1 - w_{1'}]$.

By a similar argument $wv - v'w \geq \delta(v, v') \cdot [\min_{1 \in S} w_1 - \max_{1' \in S} w_{1'}]$
 $= \delta(v, v') \cdot [-\max_{1,1' \in S} w_1 - w_{1'}]$.

In order to prove (4.3) define:

$$r^1 = T(e^1, P), \quad e^1 P > 0;$$

$$W = \{w \in \mathbb{R}_n^+ : w_1 > 0 \forall 1 \in S, \eta P w > 0, \eta' P w > 0\};$$

$$I(w) = \{1 : e^1 P w > 0\}.$$

Since it contains v , the N -vector of one's, W is nonempty.

Also, if $w \in W$, then $Pw \neq 0$ and $I(w)$ is nonempty. Now

$$\begin{aligned} & \Delta[T(\eta, P), T(\eta', P)] \\ &= \sup_{w \in W} \left\{ \sum_{j \in S} \left(\frac{\sum_{i \in I(w)} \eta_i P_{1j} w_j}{\eta P w} - \frac{\sum_{i \in I(w)} \eta'_i P_{1j} w_j}{\eta' P w} \right) + \right\} \\ &= \sup_{w \in W} \max_{J \subseteq S} \left\{ \sum_{j \in J} \left(\frac{\sum_{i \in I(w)} \eta_i P_{1j} w_j}{\eta P w} - \frac{\sum_{i \in I(w)} \eta'_i P_{1j} w_j}{\eta' P w} \right) \right\} \\ &= \sup_{w \in W} \max_{J \subseteq S} \left\{ \sum_{i \in I(w)} \left(\frac{\sum_{j \in J} \eta_i P_{1j} w_j}{\eta P w} - \frac{\sum_{j \in J} \eta'_i P_{1j} w_j}{\eta' P w} \right) \right. \\ & \quad \left. \cdot \left(\frac{\sum_{j \in S} P_{1j} w_j}{\sum_{j \in S} P_{1j} w_j} \right) \right\} \\ &= \sup_{w \in W} \max_{J \subseteq S} \left\{ \sum_{i \in I(w)} \left(\frac{\sum_{j \in S} \eta_i P_{1j} w_j}{\eta P w} - \frac{\sum_{j \in S} \eta'_i P_{1j} w_j}{\eta' P w} \right) \right. \\ & \quad \left. \cdot \left(\frac{\sum_{j \in J} P_{1j} w_j}{\sum_{j \in S} P_{1j} w_j} \right) \right\}. \end{aligned}$$

Application of (4'.1) now yields

$$\begin{aligned}
 & \Delta[I(n,P), I(n',P)] \\
 & \leq \sup_{w \in W} \max_{J \subseteq S} \left\{ \left[\sum_{i \in I(w)} \left(\frac{\sum_{j \in S} \eta_i P_{ij} w_j}{n w} - \frac{\sum_{j \in S} \eta_i P_{ij} w_j}{n' w} \right) \right] \right. \\
 & \quad \cdot \left. \left[\max_{i, i' \in I(w)} \left(\frac{\sum_{j \in J} \eta_i^1 w_j}{\eta_i^1 w} - \frac{\sum_{j \in J} \eta_{i'}^1 w_j}{\eta_{i'}^1 w} \right) \right] \right\} \\
 & \leq \Delta(n, n') \cdot \alpha(P).
 \end{aligned}$$

5. Convergence of Information Vector Approximations

(5.1) Definition: The approximation closeness is defined by

$$\xi(\underline{v}; k_0, k) = \Delta[\eta(\underline{z}; k_0, k), \eta(k)], \quad \underline{v} \in H(k_0).$$

(5.2) Lemma:

$$\frac{1}{2} |\eta(\underline{v}; k_0, k) - \eta(k)| \leq \xi(\underline{v}; k_0, k) \leq \xi(\underline{v}; k_0, k_1) \cdot \alpha[P(\underline{y}(k_1, k))], \\ \underline{v} \in H(k_0), \quad k_0 \leq k_1 \leq k \in I.$$

Proof: Apply (1.11), (3.2) and (4.3).

The following theorems give easily verified bounds on $\xi(\cdot; \cdot, \cdot)$.

(5.3) Theorem: (Strong ~~approximate~~ convergence). Define

$$\alpha(n) = \sup_{k \in I} \{ \max_{z \in Y^k(n)} \{ \alpha[P^k(z)] \} \}.$$

Then

$$\alpha(n \cdot m) \leq \alpha(n) \alpha(m)$$

and

$$\xi(\underline{v}; k_0, k) \leq \alpha(k - k_0) \leq [\alpha(n)]^{(k - k_0) \div n}, \quad \text{a.s. } \forall \underline{v} \in H(k_0).$$

Moreover, if $\alpha(n^*) < 1$, for some positive integer n^* , then

$$\lim_{k \rightarrow \infty} \xi(\underline{v}; k_0, k) = 0 \quad \text{a.s. } \forall \underline{v} \in H(k_0)$$

Proof: Immediate, using (4.5) and (5.2).

(5.4) Theorem (Weak ^{GEOMETRIC} convergence). Define

$$\bar{u}(n) = \sup_{k \in I} \{ \max_{j \in S} E\{\alpha [P^k(\underline{y}(k, k+n))] \mid \underline{a}(k)=j\} \}.$$

Then

$$\bar{u}(n+m) \leq \bar{u}(n)\bar{u}(m)$$

and

$$E\{\xi(\underline{v}; k_0, k) \mid \underline{a}(k_0)\} \leq \bar{u}(k-k_0) \leq [\bar{u}(n)]^{(k-k_0)/n} \text{ a.s. } \forall \underline{v} \in H(k_0).$$

Moreover, if $\bar{u}(n^*) < 1$, for some positive integer n^* , then

$$\lim_{k \rightarrow \infty} \xi(\underline{v}; k_0, k) = 0 \text{ a.s. } \forall \underline{v} \in H(k_0)$$

Proof:

$$\begin{aligned} \bar{u}(n+m) &\leq \sup_{k \in I} \{ \max_{j \in S} E\{\alpha [P^k(\underline{y}(k, k+n))] \\ &\quad \cdot \alpha [P^{k+n}(\underline{y}(k+n, k+n+m))] \mid \underline{a}(k)=j\} \} \\ &\leq \sup_{k \in I} \{ \max_{j \in S} E\{\alpha [P^k(\underline{y}(k, k+n))] \\ &\quad \cdot E\{\alpha [P^{k+n}(\underline{y}(k+n, k+n+m))] \mid \underline{a}(k)=j, \underline{y}(k, k+n)\} \mid \underline{a}(k)=j\} \} \\ &= \sup_{k \in I} \{ \max_{j \in S} E\{\alpha [P^k(\underline{y}(k, k+n))] \\ &\quad \cdot [E_{j \in S} E\{\alpha [P^{k+n}(\underline{y}(k+n, k+n+m))] \mid \underline{a}(k+n)=j\} \\ &\quad \Pr\{\underline{a}(k+n)=j \mid \underline{a}(k)=j, \underline{y}(k, k+n)\}] \mid \underline{a}(k)=j\} \} \\ &\leq \sup_{k \in I} \{ \max_{j \in S} E\{\alpha [P^k(\underline{y}(k, k+n))] \} \cdot \bar{u}(m) \mid \underline{a}(k)=j\} \\ &\leq \bar{u}(n) \bar{u}(m). \end{aligned}$$

The remainder is immediate, using (4.5) and (5.2).

Theorems (5.3) and (5.4) give conditions under which $\underline{n}(k)$ may be arbitrarily closely approximated by a finite-state automaton. This automaton is of course time-varying if $(\underline{S}, \underline{Y})$ is nonstationary. In (5.3), assume $\alpha(n) < 1$, for some n . Then for any $\epsilon > 0$ there is a $n(\epsilon)$ such that $\alpha(n(\epsilon)) \leq \epsilon$. Construct an automaton whose state at time k is $\underline{x}(k) = \underline{y}(k-n(\epsilon), k)$ and associate with $\underline{x}(k)$ any element $\hat{\underline{n}}(k)$ in the range of $T(\cdot, P^{k-n(\epsilon)}(\underline{x}(k)))$. Now $\frac{1}{2} |\hat{\underline{n}}(k) - \underline{n}(k)| \leq \epsilon \quad \forall k \in I$. Defining $t(n) = \log \alpha(n) / \log \#Y(n)$, (where the logarithm may be taken to any desired base,) the automaton having Q states achieves an error bounded above by $Q^{t(n)} / \alpha(n)$. Similarly, defining $\bar{t}(n) = \log \bar{\alpha}(n) / \log \#Y(n)$, the automaton having Q states achieves an average error bounded above in the sense of (5.4) by $Q^{\bar{t}(n)} / \bar{\alpha}(n)$. If $t(n)$ or $\bar{t}(n) \ll -1$, then relatively small values of Q can be made to yield extremely close approximations of the information vector. This type of approximation enjoys the advantage, over grid techniques, that it is not subject to drift.

If $\alpha(n) = 0$ for some n , then $\underline{n}(k)$ may be exactly determined by a finite-state automaton; for $\{\underline{y}(k-n, k) : k \in I\}$ is now a FSMC; and hence $\{\underline{n}(k)\}$ is FSMC. A sufficient condition for $\{\underline{n}(k)\}$ to be a FSMC is given by Theorem (6.14) in Section 6.

6. Analysis of Stationary POMCs

In this section, we study POMCs that satisfy the following assumptions:

(A1) Time-Invariance: $p^k(y) = p^{k'}(y) \forall k, k' \in I, y \in Y$.

In view of (A1), the superscript "k" will be omitted throughout this section.

(A2) State Nontransitivity: If $P_{ij}(z) > 0$, then there is a $\tilde{z} \in Y(*)$ such that $P_{ji}(\tilde{z}) > 0$.

(A3) Output Recurrence: For any $i \in S, z \in Y(*)$, there is a $\tilde{z} \in Y(*)$ such that $e^{-1}P(\tilde{z}) \neq 0$. This may be satisfied by considering, one at a time, appropriate combinations of recurrent state classes in S.

The elimination of redundant states (following Paz[10]), although desirable in certain applications, will not be required.

Let F denote a matrix whose columns span the linear space generated by $\{P(z)v : z \in Y(*)\}$ or, equivalently, by $\{P(z)v : z \in Y(N)\}$; the equivalence follows from Theorem I.B.2.1 of Paz [10]. Two stochastic vectors satisfying $vF = v'F$ are said to be equivalent since $vP(z)v = v'P(z)v \forall z \in Y(*)$, and thus they convey the same information about future observations.

Central to the analysis of stationary POMCs is the concept of a recursively computable decomposition of the state set into components whose transition probabilities are asymptotically subrectangular in the sense of (5.4).

(6.1) **Theorem:** Assume (A1), (A2), (A3) and let $\underline{I}(k_0, k)$ denote the collection of subsets of S defined by

$$\underline{I}(k_0, k) = \{ \{j : P_{1j}(y(k_0, k)) > 0\} : i \in S \} - \emptyset.$$

Then there exist a stochastic process $\{\underline{I}(k) : k \in I\}$ and an integer $M \leq N$ such that

- (a) $\# \underline{I}(k) = M$, a.s. $\forall k \in I$.
- (b) The elements of $\underline{I}(k)$ are disjoint, a.s. $\forall k \in I$.
- (c) $\underline{g}(k)$ is contained by a single element $\underline{g}(k)$ of $\underline{I}(k)$, a.s. $\forall k \in I$.
- (d) For any $k \in I$, there is a random variable $\underline{n}(k) \in \theta_p(k)$, $E\{\underline{n}(k)\} < \infty$, such that $\underline{I}(k_0, k) = \underline{I}(k) \iff k_0 \leq k - \underline{n}(k)$, a.s.
- (e) For any $k_0 \in I$, there is a random variable $\underline{m}(k_0) \in \theta_f(k_0)$, $E\{\underline{m}(k_0)\} < \infty$, such that $\underline{I}(k_0, k) = \underline{I}(k) \iff k \geq k_0 + \underline{m}(k_0)$, a.s.
- (f) $\underline{I}(k_0, k) = \underline{I}(k) \iff P(\underline{g}(k_0, k))$ is M -subrectangular, a.s.

Proof: The proof is given in Section 6'.

If $\underline{g}(k_0)$ is known to the observer, then subsequent values of $\underline{g}(k)$ may be computed according to the identity

$$(6.2) \quad \underline{g}(k) = \tau(\underline{g}(k_0), \underline{y}(k_0, k)), \quad \tau(\sigma, z) = \{j : P_{1j}(z) > 0, i \in \sigma\}.$$

Similarly, (6.2) provides a means for determining $\underline{I}(k)$ from $\underline{I}(k_0)$ and $\underline{y}(k_0, k)$.

We now show that $\underline{\eta}(\underline{y}; k_0, k)$ may be decomposed into components $\{\underline{\eta}(\underline{y}; k_0, k) | \sigma\}$, $\sigma \in \underline{I}(k)$, each of which converges in a manner similar to (5.3) or (5.4). This decomposition is preserved by (1.9) in the sense that

$$(6.3) \quad \underline{\eta}(\underline{v}|\sigma; k_0, k) = \underline{\eta}(\underline{v}; k_0, k) | \tau(\sigma, \underline{y}(k_0, k))$$

a.s. $\forall \underline{v} \in H(k_0), \sigma \in \Sigma(k_0), \underline{v}e^\sigma > 0.$

The convergence properties are stated as follows. In analogy to (5.1) define

$$\begin{aligned} & \underline{\xi}_M(\underline{v}; k_0, k) \\ & = \max\{\Delta[\underline{\eta}(\underline{v}|\sigma; k_0, k), \underline{\eta}(\underline{v}(k_0)|\sigma; k_0, k)] : \sigma \in \Sigma(k_0), \underline{v}e^\sigma > 0, \underline{\eta}(k_0)e^\sigma > 0\} \\ (6.4) \quad & = \max\{\Delta[\underline{\eta}(\underline{v}; k_0, k)|\sigma, \underline{\eta}(k)|\sigma] : \sigma \in \Sigma(k), \underline{\eta}(\underline{v}; k_0, k)e^\sigma > 0, \underline{\eta}(k)e^\sigma > 0\}, \end{aligned}$$

$\underline{v} \in H$

so that, in analogy to (5.2),

$$(6.5) \quad \frac{1}{2} |\underline{\eta}(\underline{v}; k_0, k) - \underline{\eta}(k)| \leq \underline{\xi}_M(\underline{v}; k_0, k) \leq \underline{\xi}_M(\underline{v}; k_0, k_1) \cdot \alpha_M[P(\underline{y}(k_1, k))],$$

$\underline{v} \in H(k_0), k_0 \leq k_1 \leq k \in I.$

The immediate consequences of (5.3) and (5.4) are now:

(6.6) Theorem: ^{GEOMETRIC} (Strong ~~exponential~~ convergence within a component.)

Assume (A1), (A2), (A3), and define

$$\alpha_M(n) = \max_{z \in \Upsilon(n)} \{\alpha_M[P(z)]\}.$$

Then

$$\alpha_M(n+m) \leq \alpha_M(n) \alpha_M(m)$$

and

$$\underline{\xi}_M(\underline{v}; k_0, k) \leq \alpha_M(k-k_0) \leq [\alpha_M(n)]^{(k-k_0)/n}, \quad \text{a.s. } \forall \underline{v} \in H(k_0).$$

Moreover, if $\alpha_M(n^*) < 1$, for some positive integer n^* , then

$$\lim_{k \rightarrow \infty} \underline{\xi}_M(\underline{v}; k_0, k) = 0, \quad \text{a.s. } \forall \underline{v} \in H(k_0)$$

(6.7) Theorem: ^{GEOMETRIC} (Weak ~~convergence~~ convergence within a component.)

Assume (A1), (A2), (A3), and define

$$\bar{\alpha}_M(n) = \max_{\underline{1} \in S} E\{\alpha_M[P(\underline{Y}(0,n))] \mid \underline{s}(0)=1\}.$$

Then

$$\bar{\alpha}_M(n^*) < 1 \text{ for some positive integer } n^*;$$

$$\bar{\alpha}_M(n+m) \leq \bar{\alpha}_M(n) \bar{\alpha}_M(m);$$

$$E\{\underline{1}_M(\underline{1}; k_0, k) \mid \underline{s}(k_0)\} \leq \bar{\alpha}_M(k-k_0) \leq [\bar{\alpha}_M(n)]^{(k-k_0)/n}, \quad \text{a.s. } \forall \underline{1} \in H(k),$$

$$\lim_{k \rightarrow \infty} \underline{1}_M(\underline{1}; k_0, k) = 0, \quad \text{a.s. } \forall \underline{1} \in H(k_0).$$

Finally, we show that a weighted sum of equivalent components of $\underline{n}(\underline{1}; k_0, k)$ converges to $\underline{n}(k)$. Define:

$$\tau_1(\hat{\underline{1}}) = \Pr\{\underline{s}(0)=1 \mid \underline{1}(0) = \hat{\underline{1}}\}$$

$$f(\hat{\sigma}, \underline{1}(k)) = \bigcup \{\sigma \in \underline{1}(k) : [\tau(\underline{1}(k)) \mid \sigma] F = [\tau(\underline{1}(k)) \mid \hat{\sigma}] F\}$$

(6.8) $\underline{1}^*(k) = \{f(\sigma, \underline{1}(k)) : \sigma \in \underline{1}(k)\}$

$$\underline{g}^*(k) = f(\underline{g}(k), \underline{1}(k))$$

$$\psi(\hat{\sigma}, \hat{\sigma}^*, \hat{\underline{1}}) = [\tau(\hat{\underline{1}}) \mid \hat{\sigma}^*] \sigma \hat{\sigma}$$

We may interpret $f(\hat{\sigma}, \underline{1}(k))$ as the union of elements in $\underline{1}(k)$ that are indistinguishable from $\hat{\sigma}$; likewise $\underline{1}^*(k)$ is the collection of indistinguishable unions in $\underline{1}(k)$ and $\underline{g}^*(k)$ is the element of $\underline{1}^*(k)$ that contains $\underline{g}(k)$.

(6.9) Lemma: (Component Identification) Assume (A1), (A2), (A3)

and consider the events

$$A = \{I(k_0) = \hat{I}\}$$

$$B = A \cap \{\sigma(k_0) = \hat{\sigma}\}$$

where $\Pr(A) \geq \Pr(B) > 0$. Then

$$\lim_{\substack{k' \rightarrow \infty \\ k \rightarrow -\infty}} \Pr(B | A, \underline{y}(k, k')) = \psi(\hat{\sigma}, \hat{\sigma}^h(k_0), \hat{I}) \quad \text{a.s.}$$

Proof: Throughout this proof, consider all random variables and events to be conditioned on A. For $n \in \mathbb{I}$ define.

$$\underline{k}(n) = \begin{cases} \min\{k > k(n-1) : I(k) = \hat{I}, \tau(\sigma, \underline{y}(k_0, k)) = \sigma \quad \forall \sigma \in \hat{I}, & n > 0 \\ k_0, & n = 0 \\ \max\{k < k(n+1) : I(k) = \hat{I}, \tau(\sigma, \underline{y}(k, k_0)) = \sigma \quad \forall \sigma \in \hat{I}, & n < 0 \end{cases}$$

$$\underline{g}'(n) = \underline{g}(\underline{k}(n))$$

$$\underline{y}'(n, n') = \underline{y}(\underline{k}(n), \underline{k}(n'))$$

For $\sigma \in \hat{I}$, also define

$$\underline{\lambda}_\sigma(n, n') = \ln(\Pr\{\underline{y}'(n, n') \mid B\} / \Pr\{\underline{y}'(n, n') \mid \underline{\sigma}(0) = \sigma\}).$$

Now $\underline{\lambda}_\sigma(n, n')$ and $\underline{\lambda}_\sigma(0, n'-n)$ are identically distributed, so

$$E(\underline{\lambda}_\sigma(n, n') \mid B) = \bar{\lambda}_\sigma(n'-n).$$

By convexity of the $-\ln$ function, $\bar{\lambda}_\sigma(n)$ is nondecreasing in n .

Moreover, $\underline{g}'(n)$ is a FSMC whose transition probability matrix P' satisfies $P'_{1j} > 0 \quad \forall 1, j \in \sigma \in \hat{I}$; hence by a standard argument of convergence of powers of positive stochastic matrices, there is, for every $\epsilon > 0$,

$$\begin{aligned} & \text{a function } d^\epsilon \text{ satisfying } \ln(\Pr\{\underline{y}'(d^\epsilon(n)-n, d^\epsilon(n)) \mid \underline{y}'(0, n), B\} \\ & / \Pr\{\underline{y}'(d^\epsilon(n)-n, d^\epsilon(n)) \mid \underline{y}'(0, n), \underline{\sigma}(0) = \sigma\}) \geq (1-\epsilon) \underline{\lambda}_\sigma(d^\epsilon(n)-n, d^\epsilon(n)) \quad \text{a.s.} \end{aligned}$$

Now consider $n_1 = 1$, $n_{1+1} = d^\epsilon(n_1)$, and $\underline{\phi}_\sigma(1) = \ln(\Pr\{\underline{y}'(0, n_1),$

$$\underline{y}'(n_{1+1}-n_1, n_{1+1}) \mid B\} / \Pr\{\underline{y}'(0, n_1), \underline{y}'(n_{1+1}-n_1, n_{1+1}) \mid \underline{\sigma}(0) = \sigma\}).$$

By convexity of the $-\ln$ function,

$$\begin{aligned} & \lambda_{\sigma}(0, n_1) + (1-\epsilon) \bar{\lambda}_{\sigma}(n_1) \\ & \leq E(\phi_{\sigma}(1) \mid \lambda_{\sigma}(0, n_1), B) \\ & \leq E(\lambda_{\sigma}(0, n_1+1) \mid \lambda_{\sigma}(0, n_1), B) \end{aligned}$$

Hence, either $\lambda_{\sigma}(n, n') = 0$ a.s. in $B \iff \underline{Y}$ and $\underline{\sigma}(k_0)$ are conditionally independent given $(\underline{\sigma}(k_0) = \sigma \text{ or } \hat{\sigma}) \iff \sigma \subseteq f(\hat{\sigma}, \underline{Y})$; or

else $\lim_{n' \rightarrow \infty} \lambda_{\sigma}(n, n') = 0$ a.s. in $B \iff \lim_{\substack{k' \rightarrow \infty \\ k' \leq k}} \Pr(\underline{\sigma}(k_0) \in \sigma \mid A, \underline{y}(k, k')) = 0$ a.s. in $B \iff \sigma \subseteq g-f(\hat{\sigma}, \underline{Y})$. The desired result follows trivially.

(6.10) Theorem: (Overall convergence) Assume (A1), (A2), (A3).

Then

- (a) $\underline{\eta}(k) = [\underline{\eta}(k) \mid \underline{\sigma}^*(k)]$
- (b) $\underline{\eta}(k) e^{\sigma} = \psi(\sigma, \underline{\sigma}^*(k), \underline{I}(k)) \quad \forall \sigma \in \underline{I}(k)$
- (c) $[\underline{\eta}(k) \mid \sigma] F = \underline{\eta}(k) F \quad \forall \sigma \in \underline{I}(k), \sigma \subseteq \underline{\sigma}^*(k)$

and for all $\underline{x} \in H(k_0)$:

- (a') $\lim_{k \rightarrow \infty} \underline{\eta}(\underline{x}; k_0, k) e^{\underline{\sigma}^*(k)} = 1 \quad \text{a.s.}$
- (b') $\lim_{k \rightarrow \infty} | \sum_{\sigma \in \underline{I}(k)} (\sum_{\sigma \in \underline{I}^*(k)} \psi(\sigma, \underline{\sigma}^*(k), \underline{I}(k)) \underline{\eta}(\underline{x}; k_0, k) e^{\sigma}) \cdot [\underline{\eta}(\underline{x}; k_0, k) \mid \sigma] - \underline{\eta}(k) | = 0 \quad \text{a.s.}$
- (c') $\lim_{k \rightarrow \infty} | \underline{\eta}(\underline{x}; k_0, k) F - \underline{\eta}(k) F | = 0 \quad \text{a.s.}$

Remark: If $\underline{I}^*(k) = \underline{I}(k)$ a.s., e.g., if F has rank N , then (c') becomes

$$\lim_{k \rightarrow \infty} | \underline{n}(\underline{y}; k_0, k) - \underline{n}(k) | = 0 \quad \text{a.s.}$$

which establishes the conjecture of Kaijser [8]. This convergence can occur under numerous conditions readily derived from (b'); for example, it is implied by $\underline{y} = v(\underline{I}(k_0))$.

Proof: Parts (a), (b), (c) are the consequence of (6.9); parts (a'), (b'), (c') then follow from (1.11) and (6.7).

Hence, the state estimates may be approximated arbitrarily closely (in a time average or mean $|\cdot|$ sense) by a finite-state automaton. If $\underline{g}^*(k_0)$ is known initially, the automaton is constructed as discussed in Section 5. Memory-efficient approximations to $\underline{n}(\underline{y}; k_0, k) e^{\sigma^*}$, $\sigma^* \in \underline{I}^*(k)$, may be obtained by generalized hypothesis testing procedures; see Section 2.2 of [3].

The state estimates for a POMC satisfying (A1), (A2), and (A3) can be exactly computed by a (properly initialized) finite-state automaton if $\{\underline{n}(k)\}$ is a FSMC, i.e., if there is a finite set $G \subset I$ such that $\underline{n}(k) \in G \quad \forall k \in I$. In this case $\{(\underline{n}(k), \underline{y}(k)) : k \in I\}$ is a state-calculable POMC. This condition may be verified in the manner described below.

Let Γ denote the partition of S into ergodic classes; hence for $\gamma \in \Gamma$ there holds $\gamma = \{j : P_{1j}(z) > 0, z \in Y(\gamma)\}$. For each $\gamma \in \Gamma$, select $\hat{\sigma}_\gamma \in \hat{I}_\gamma$, $\hat{z}_\gamma \in Y(\gamma)$ with $\hat{\sigma}_\gamma^* = f(\hat{\sigma}_\gamma, \hat{z}_\gamma)$, $M_\gamma = \#\{\sigma \in \hat{I}_\gamma : \sigma \subseteq \hat{\sigma}_\gamma^*\}$ such that

$$\begin{aligned}
 & \hat{\sigma}_\gamma \subseteq \gamma; \\
 & \Pr\{\underline{\sigma}(0) = \hat{\sigma}_\gamma, \underline{z}(0) = \hat{z}_\gamma\} > 0; \\
 & \sigma = \tau(\sigma, \hat{z}_\gamma) \quad \forall \sigma \in \hat{I}_\gamma, \sigma \subseteq \hat{\sigma}_\gamma^*;
 \end{aligned}$$

(6.11)

$$\hat{P}_{1j}^\gamma = \begin{cases} P_{1j}(\hat{z}_\gamma), & \text{if } 1 \in \hat{\sigma}_\gamma^* \\ 0, & \text{otherwise} \end{cases} \quad \text{is } M_\gamma\text{-subrectangular.}$$

Now $\alpha_{M_\gamma}[\hat{P}^\gamma] < 1$; hence by (4.9) there exists a unique limit

$$(6.12) \quad \hat{\nu}^\gamma = \lim_{n \rightarrow \infty} T([\nu(\hat{z}_\gamma) | \hat{\sigma}_\gamma^*], (\hat{P}^\gamma)^n).$$

Finally, construct

$$\begin{aligned}
 (6.13) \quad G(0) &= \{\hat{\nu}^\gamma : \gamma \in \Gamma\}; \\
 G(m+1) &= G(m) \cup \{T(\nu, P(y)) : \nu \in G(m), y \in Y, \nu P(y) \neq 0\}; \\
 \hat{G}(m) &= \{\nu F : \nu \in G(m)\}.
 \end{aligned}$$

We may now demonstrate:

(6.14) Theorem: Assume (A1), (A2), (A3). Then

(a) $\{\underline{n}(k)\}$ is a FSMC iff $G(m^*) = G(m^*+1)$ for some positive integer m^* . Moreover, if such an m^* exists, then $G = G(m^*)$ is the minimal subset of Π satisfying $\underline{n}(k) \in G$ a.s. $\forall k \in I$.

(b) $\{\underline{n}(k)\}$ is equivalent to a FSMC iff $\hat{G}(m^*) = \hat{G}(m^*+1)$ for some positive integer m^* . Moreover, if such an m^* exists, then $\hat{G} = \hat{G}(m^*)$ is the minimal set satisfying $\underline{n}(k)F \in \hat{G}$ a.s. $\forall k \in I$; the equivalent FSMC must assume at least $\#\hat{G}$ distinct values.

Remark: This proposition provides the basis for an algorithm that decides whether $\{\eta(k)\}$ is (or is equivalent to) a FSMC having, at most, a specified number of states. It cannot, however, be used to decide whether $\{\eta(k)\}$ is (or is equivalent to) a FSMC when the size of G (resp. \hat{G}) is unconstrained.

Proof: Only the proof of (a) is given. The proof of (b) is similar.

(Necessity) Suppose $G(m^*) = G(m^*+1)$. Then for every $\epsilon > 0$ there is a $\tilde{v}^1 \in G(m^*)$ such that $\tilde{v}_1^1 > 0$; now $\tilde{\eta}(0) = \tilde{v}_1^1(0) \in H(0)$. Define $\tilde{\eta}(k) = \eta(\tilde{\eta}(0); 0, k)$. By (6.7) $\lim_{k \rightarrow \infty} |\tilde{\eta}(k) - \eta(k)| = 0$. But $\tilde{\eta}(k) \in G(m^*+k) = G(m^*)$, and $\{\eta(k)\}$ is stationary. Hence $\eta(k) \in G(m^*)$ a.s. $\forall k \in I$.

(Sufficiency) Assume that $\{\eta(k)\}$ is a FSMC and let G be the minimal set satisfying $\eta(k) \in G$ a.s. $\forall k \in I$. Clearly $G \cap H(\hat{\sigma}_\gamma^*)$ is nonempty, $\forall \gamma \in \Gamma$. Thus $G(0) \subseteq G$; for otherwise some $\gamma \in \Gamma$ satisfies $\hat{v}^\gamma \notin G \Rightarrow \min_{v \in G} \frac{1}{2} |\hat{v}^\gamma - v| = \epsilon > 0$; but $\alpha_{M_\gamma} [\hat{P}^\gamma]^n < \epsilon$ (some n); now the concatenation z^* of \hat{z}^γ with itself n times and $v^* = T(v, P(z^*))$ (some $v \in G \cap H(\hat{\sigma}_\gamma^*)$) satisfy $v^* \in G$ and $|\hat{v}^\gamma - v^*| < \epsilon$, which is a contradiction. Since $\{T(v, P(y)) : v \in G, y \in Y, vP(y) \neq 0\} = G$, it follows that $G(m) \subseteq G \Rightarrow G(m+1) \subseteq G$. Now $G(0) \subseteq G(1) \subseteq \dots \subseteq G$. Since G is a finite set, m^* exists and $G(m^*) = G(m^*+1) \subseteq G$. But now $G(m^*) \subseteq G$ by the argument given above; thus $G(m^*) = G$.

6' Proof of Theorem (6.1)

Before giving the proof, we introduce some new terminology.

Let S^* denote the power set (collection of subsets) of S ; S^{**}

in turn will be the power set of S^* . Now define $\Delta : S^{**} \rightarrow S^{**}$,

$\Pi : S^{**} \rightarrow S^{**}$, $R : S \times Y(*) \rightarrow S^*$, $I : Y(*) \rightarrow S^{**}$, and $U : Y(*) \rightarrow S^{**}$,

as follows:

$$\Delta[I] = \{\sigma \in I : \text{no element of } I \text{ is properly contained in } \sigma\} \subseteq I$$

$$\Pi[I] = \{\sigma_1 \cup \sigma_2 \cup \dots \cup \sigma_n : \sigma_1, \dots, \sigma_n \in I, n > 0\}$$

$$B_{1j}(z) = 1 \text{ if } P_{1j}(z) > 0, 0 \text{ if } P_{1j}(z) = 0$$

$$R(i, z) = \{j : P_{1j}(z) > 0\}$$

$$I(z) = \{R(i, z) : i \in S\} - \{\emptyset\}$$

$$U(z) = \Pi[I(z)]$$

$$M = \min\{\#I(z) : z \in Y(*)\}$$

$$M^* = \min\{\#U(z) : z \in Y(*)\}.$$

Note that $\underline{I}(k_0, k) = I(\underline{y}(k_0, k))$. Also define the partial order " $<$ " on $Y(*)$:

$$\tilde{z} < z \text{ if } \tilde{z} = z'z \text{ for some } z' \in Y(*) .$$

Finally, define

$$E = \{z \in Y(*) : U(\tilde{z}) = U(z) \forall \tilde{z} < z\}.$$

(6'.1) Lemma: (a) $U(z'z) \subseteq U(z)$

(b) $\#U(z'z) \leq \#U(z')$

Proof: By definition, $U(z) = \{\bigcup_{i \in C} R(i, z) : C \in S^*\} - \{\emptyset\}$; and

$R(i, z'z) = \bigcup_{j \in R(i, z')} R(j, z)$; so $U(z'z) = \{\bigcup_{j \in C} R(j, z) : C \in U(z')\} - \{\emptyset\}$.

(6'.2) Corollary: $E = \{z : \#U(z) = M^*\}$.

Proof: By (6'.1)(a), $\#U(z) = M^*$ implies $z \in E$. Now suppose $\#U(z) > M^*$.

By definition of M^* , we may select z^* such that $\#U(z^*) = M^*$, and by (A3), we may select \tilde{z} such that $z^* \tilde{z} z \in Y^*$. Now $\#U(z^* \tilde{z} z) = M^*$ by (6'.1) (b); so $z \notin E$.

(6'.3) Corollary: For any $k \in I$, there is a random variable $\underline{n}(k) \in \theta_p(k)$, $E\{\underline{n}(k)\} < \infty$, such that $\underline{y}(k, k) \in E \iff k_0 \leq k - \underline{n}(k)$

Proof: Select $\underline{n}(k) = \min\{\underline{n} : \#U(\underline{y}(k-\underline{n}, k)) = M^*\}$. $E\{\underline{n}(k)\} < \infty$ because $\{[\underline{s}(k-n), B(\underline{y}(k-n, k)), U(\underline{y}(k-n, k))]\} : n=0, \dots\}$ is a FSMC whose recurrent states $[s, B, U]$ have the property $\#U = M^*$, by (A2), (A3) and (6'.1) (b), and since $\underline{n}(k)$ is therefore bounded above by the number of transitions until this FSMC enters a recurrent stage, a random variable having finite expectation.

(6'.4) Lemma: $I(z) = \Delta[U(z)]$, $\forall z \in E$.

Proof: (By contradiction): Clearly $I(z) \supseteq \Delta[U(z)]$.

Now assume that $z \in E$, $\sigma \in I(z)$, $\sigma' \in U(z)$ and $\sigma' \not\subseteq \sigma$. We may select $\sigma'' \in \Delta[U(z)]$ so that $\sigma'' \subseteq \sigma'$. Let i be any element of σ'' and pick j so that $R(j, z) = \sigma$; now $P_{ji}(z) > 0$. By (A?), there exists $\tilde{z} \in Y^*$ satisfying $P_{ji}(\tilde{z}) > 0$. Now construct $z_0 = z$ and $z_n = z_{n-1} \tilde{z} z_{n-1}$, $n \in S$. Trivially $P_{ji}(z_n) > 0$ so $z_n \in Y^*$, and hence $z_n \in E$. But $z \in E$, so $\sigma'' \in \Delta[U(z)] = \Delta[U(z_n)] \subseteq I(z_n)$, and there is a k_n such that $R(k_n, z_n) = \sigma''$. Moreover, by construction of z_n , $R(k_1, z_n) \supseteq \sigma > \sigma''$, $i=1, \dots, n-1$, so $k_n \neq k_{n'}$ when $n \neq n'$. But now $R(i, z_n) \supseteq \sigma > \sigma''$, $\forall i \in S$ and $\sigma'' \notin U(z_n) = U(z) \supseteq \Delta[U(z)]$, which contradicts the original assumption that $\sigma'' \in \Delta[U(z)]$.

(6'.5) Corollary: $\hat{z} < z \in E \implies I(\hat{z}) = I(z)$

(6'.6) Lemma: The elements of $I(z)$ are disjoint, $\forall z \in E$.

Proof: Suppose (i) $\hat{z} < z \in E$ and (ii) $B(\hat{z}\hat{z}) = B(\hat{z})$.

Then (i) implies $\hat{z} \in E$ and hence, by (6'.5), $I(\hat{z}) = I(z)$.

Now, by (ii), $R(j, \hat{z}) = \bigcup_{1 \in R(j, \hat{z})} R(1, z)$.

So $k \in R(1, \hat{z}) \cap R(j, \hat{z})$ implies $R(k, \hat{z}) \subseteq R(1, \hat{z})$ and $R(k, \hat{z}) \subseteq$

$R(j, \hat{z})$. Hence, by (6'.4), $R(1, \hat{z}) = R(j, \hat{z})$, and the elements of $I(\hat{z})$

are disjoint. It only remains to show that there exists a \hat{z} satisfying (i) and

Pick i, j so $P_{ij}(z) > 0$ and pick \tilde{z} so $P_{ij}(\tilde{z}) > 0$;

\tilde{z} exists by (A2). Define $z_1 = \tilde{z} z$ and $z_{n+1} = z_n \tilde{z} z$.

By construction $z_n z_m = z_{n+m}$. Moreover $P_{ij}(z'z) > 0$

so $P(z_n) \neq 0$ and $z_n \in E$. Since $\{B(z_n) : n=1, 2, \dots\}$ is a finite

set, there exist $n, m > 0$ such that $B(z_n) = B(z_{n+m})$. Hence,

$B(z_n) = B(z_{n+m}) = B(z_m) B(z_n) = B(z_m) B(z_{n+m}) = B(z_{n+2m}) = \dots = B(z_{n+nm})$

and similarly $B(z_{(m-1)n} z_n) = B(z_{(m-1)n} z_{n+nm})$; hence $B(z_{nm}) = B(z_{nm} z_{nm})$.

Identify $\hat{z} = z_{nm}$ to complete the proof.

(6'.7) Corollary: $E = \{z : \#I(z) = M\}$.

Proof: By (6'.2) and (6'.6), $M^* = \#U(z) = 2^{\#I(z)} - 1$, $z \in E$, and

$M^* < \#U(z) \leq 2^{\#I(z)} - 1$, $z \notin E$.

Proof of (6.1): Define $\underline{n}(k)$ as in (6'.3), and let $\underline{I}(k) = \underline{I}(k-\underline{n}(k), k)$.

Now $\underline{I}(k) = \underline{I}(\underline{y}(k-\underline{n}(k)), k)$, where $\underline{y}(k-\underline{n}(k), k) \in \mathbb{E}$.

(a) Follows from (6'.7).

(b) Follows from (6'.6).

(c) If j is contained by no element of $\underline{I}(k) = \underline{I}(k-\underline{n}(k), k)$, then $P_{1j}(\underline{y}(k-\underline{n}(k), k)) = 0 \forall i \in S$, by definition of $\underline{I}(\cdot, \cdot)$, and hence $\Pr\{\underline{g}(k)=j \mid \underline{y}(k-\underline{n}(k), k)\} = 0$.

Thus $\underline{g}(k)$ is contained by an element $\underline{g}(k)$ of $\underline{I}(k)$ a.s.

Uniqueness of $\underline{g}(k)$ follows from (b) above.

(d) Follows from (6'.3) and (6'.5).

(e) Define $\tilde{P}_{1j}(y) = \Pr\{\underline{g}(k-1)=j, \underline{y}(k)=y \mid \underline{g}(k)=i\}$.

and $\tilde{P}(y_1 \dots y_n) = \tilde{P}(y_n) \dots \tilde{P}(y_1)$.

Then $\{\tilde{P}(y)\}$ defines a POMC with state process

$\{\tilde{g}(k)=\underline{g}(-k) : k \in I\}$ and observation process $\{\tilde{y}(k)=\underline{y}(-k) : k \in I\}$.

Also $\tilde{P}_{1j}(z) > 0 \iff P_{1j}(z) > 0$. Hence by (f), below, and (6'.3),

there is an $\underline{m}(k) \in \theta_f(k)$ such that $k \geq k_0 + \underline{m}(k_0)$

$\iff \tilde{P}(\tilde{y}(-k, -k_0))$ is M -subrectangular $\iff P(\underline{y}(k_0, k))$ is

M -subrectangular.

(f) Follows from (6'.6) and (6'.7).

7. Examples

We now illustrate by means of simple stationary POMCs the concepts introduced in this paper.

(7.1) Example: Suppose that \underline{S} is an ergodic Markov chain with

$$P_{ij}(y) = \begin{cases} P_{ij}, & \text{if } y = j \\ 0, & \text{otherwise} \end{cases}$$

Now $y(k) = \underline{g}(k) \forall k \in I$, i.e. there is perfect state observation.

$\alpha(1) = 0$, and $\underline{\eta}(\underline{y}; k_0, k) = \underline{\eta}(k) = \underline{a}^{\underline{y}}(k)$ whenever $k > k_0$, and $\underline{y} \in H(k_0)$.

Here $\{\underline{\eta}(k)\}$ is a FSMC.

(7.2) Example: Suppose that \underline{S} is a Markov chain with strictly positive transition probability matrix. Let $Y = S$ and

$$P_{ij}(y) = \begin{cases} P_{ij}, & \text{if } y = i-j \text{ or } y = N+1-j \\ 0, & \text{otherwise} \end{cases}$$

Now $\underline{g}(k) = \{\underline{g}(k_0) + \sum_{k'=k_0+1}^k \underline{y}(k')\} \bmod N$,

so $(\underline{S}, \underline{Y})$ is state-calculable. This result is expressed by (6.1),

where $\underline{M} = N$ and $\underline{I}(k) = \{\{1\}, \dots, \{N\}\}$, the set of singletons in S .

Clearly $\underline{g}(k) = \{\underline{g}(k)\}$.

Hence, if no two rows of F coincide, and if $\underline{y}^{\text{CH}}(k_0)$, then

$$\underline{\eta}(k) = \underline{a}^{\text{CH}}(k);$$

$$\Delta[\underline{\eta}(\underline{y}; k_0, k), \underline{\eta}(k)] = 1, \text{ whenever } \underline{y} = \underline{a}^{\text{CH}}(k_0);$$

$$\lim_{k \rightarrow \infty} |\underline{\eta}(\underline{y}; k_0, k) - \underline{\eta}(k)| = 0 \text{ a.s.}$$

Note that $|\underline{\eta}(\underline{y}; k_0, k) - \underline{\eta}(k)|$ is not monotone decreasing in k .

Again $\{\underline{\eta}(k)\}$ is a FSMC.

(7.3) Example: Suppose

$$P(1) = \begin{bmatrix} 1/4 & 1/4 \\ 0 & 1/2 \end{bmatrix}, P(2) = \begin{bmatrix} 1/2 & 0 \\ 1/4 & 1/4 \end{bmatrix}.$$

Now $\underline{\eta}(k+1) = 1/2[(1,0) + \underline{\eta}(k)]$ if $y(k+1) = 1$, and $\underline{\eta}(k+1) = 1/2[(0,1) + \underline{\eta}(k)]$

if $y(k+1) = 2$. Thus it happens that $\frac{1}{2}|\underline{\eta}(\underline{y}; k_0, k) - \underline{\eta}(k)| \leq (.5)^{k-k_0}$.

This convergence is predicted by Theorem (5.4) which states that

$E\{\Delta[\underline{\eta}(\underline{y}; k_0, k), \underline{\eta}(k)]\} \leq \bar{\alpha}(k-k_0)$. Since $\bar{\alpha}(2) = .634$, we have

$\bar{\alpha}(n) \leq (.634)^{n+2}$. Moreover $\Delta[\underline{\eta}(\underline{y}; k_0, k), \underline{\eta}(k)]$ is monotone

decreasing in k a.s., by (4.3). Clearly, $\{\underline{\eta}(k)\}$ is not a FSMC. But \underline{Y} is a

process of independent identically distributed Bernoulli trials; hence past

observations convey no information about future observations. Indeed $F = \nu$ and

all information vectors are equivalent. Since \hat{G} contains only the

element 1, $\{\underline{\eta}(k)\}$ is equivalent to a one-state FSMC.

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