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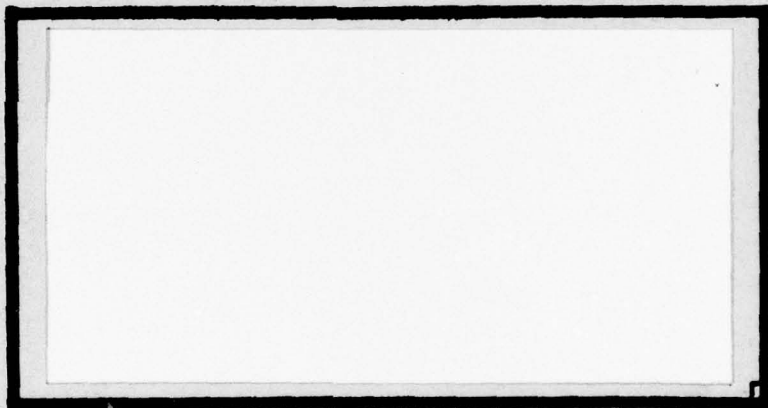
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⑨ master's THESIS

⑪ Mar 77

⑫ AFIT/GCS/MA/78-1

⑩ Alfred H. Linder III
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SIMULATION OF THE ACQUISITION MANAGEMENT
INFORMATION SYSTEM

THESIS

Presented to the Faculty of the School of Engineering
of the Air Force Institute of Technology
Air University
in Partial Fulfillment of the
Requirements for the Degree of
Master of Science

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by

Alfred H. Linder III, B.S.

2Lt USAF

Graduate Computer Science

March 1977

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PREFACE

This investigation is a result of an effort to provide AFSC, Wright-Patterson A.F.B., Ohio with a computer simulation analysis of the IBM 370/155 computer system's throughput performance and batch-interactive-mix. I hope that the results will prove useful to the Acquisition Management Information System (AMIS) department.

I wish to express my sincere thanks to Major Kenneth Melendez of AFIT/ENC for his advice and leadership as my advisor and his interest in this effort. I also want to thank my wife Linda for her emotional support throughout this undertaking.

Alfred H. Linder, III

38
40
41
43
43
46
48
49
51
66
120

CONTENTS

Page

face ii

Table of Figures v

Table of Tables vi

Preface vii

Introduction 1

 Background 1

 General Discussion 3

 Problem Statement 4

 Scope 4

 Approach 5

IBM 370 Operating Characteristics 8

 System Organization 8

 Main Storage 9

 Central Processing Unit 10

 Input/Output 10

 Byte Multiplexer Channels 11

 Block Multiplexer Channels 11

 Job Scheduler 12

 Multiprogramming 14

 System 2000 14

 Data Bank 15

 Data Base Management System 16

 Data Dictionary/Directory System 16

 Data Base Administrator 16

 User System Interface 16

System Requirements, Definitions and Performance Measures 18

 Job Processing Sequence 18

 Workload Parameters 20

 Performance Measures 22

 Development of Simulation Model 26

Simulation Steps 28

 Model Formulation 28

 Logic Flow 31

 Data Preparation 33

 Programming Techniques 33

 Validation of the Simscript Model 35

 Assumptions 35

 Steady State 36

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LIST OF FIGURES

<u>Figure</u>		<u>Page</u>
1	Computer Response Time	2
2	IBM System/370 Logical Structures	8
3	Organization of Buffer Memory Hardware	9
4	Data Base Management: Conceptual Environment	15
5.	Simulation Process Flowchart	24
6	Job Processing Sequence	29
7	Transient Warm-up Period	37
8	Nomograph for Cost Estimates	41

LIST OF TABLES

<u>Table</u>		<u>Page</u>
I	Scheduling Priorities	13

ABSTRACT

↙ The emphasis of this investigation is on the development of a simulation model that will allow an analysis to be performed on the batch-interactive-mix and throughput performance of the IBM 370/155 computer, used by the Acquisition Management Information System (AMIS) users. The model is driven from data obtained from the AMIS major production jobs and hence the validity of the results are not accurate. In the recommendation section of this effort, cluster analysis is presented as a method to collect data about all the types of AMIS jobs so that a follow on study could be made to achieve valid results from the simulation model that has been constructed.

↘ The basics of hardware and software are discussed, along with some information concerning the System 2000 data base management software package. After the operating characteristics have been discussed, the need for system requirements, definitions and performance measures are presented and the necessary simulation steps are enumerated to show how the simulation model was constructed.

The main variables of interest (CPU utilization and gain factor) are analyzed to determine the throughput performance of the system as the interactive workload is increased and as the number of interactive ports are decreased. ↗

SIMULATION OF THE ACQUISITION MANAGEMENT INFORMATION SYSTEM

I. Introduction

A time sharing system should provide the best possible throughput performance for batch and interactive jobs. Throughput is the amount of useful work completed per unit of time given a particular workload. (Ref. 30:16) Even after a computer system has been designed and implemented it may be necessary to perform a batch-interactive-mix analysis to determine if the system is being used in such a manner that high throughput performance is achieved. A computer simulation analysis of the operating system might reveal pertinent information to aid in the determination of the exact number of batch and interactive terminals to be used in achieving good throughput performance. A simulation model was developed to analyze the throughput performance for an Air Force System Command (AFSC) organization.

Background

The Air Force System Command (AFSC) at Wright Patterson Air Force Base uses the System 2000 software package to store contracts as they are developed from the conception phase through the completion phase, as well as during the purchasing phase of the acquired system. There are approximately 75 terminals, used at several locations across the United States. These are connected to the System 2000, which runs on an IBM 370/155 model computer.

The Acquisition Management Information System (AMIS) was implemented to support contract administration and disbursement activities. One of the major objectives is to implement Source Data Automation (SDA) at the buying activities, AF Plant Representative Offices (AFPROs), and the

Air Force Contract Management Department (AFCMD) - thus providing them with an interactive capability to update and query the central data base.

Throughput performance could possibly be improved if a batch-interactive-mix analysis is performed and the best ratio of batch to interactive terminals is used by AFSC. A simulation program which models the System 2000, I/O channels, controllers, memory allocation technique of the IBM 370, job scheduling, and outputting of jobs would provide results which can be used to determine the best batch-interactive-mix. Good throughput performance assures that the work being processed will be completed in time for its intended use as illustrated in Figure 1.

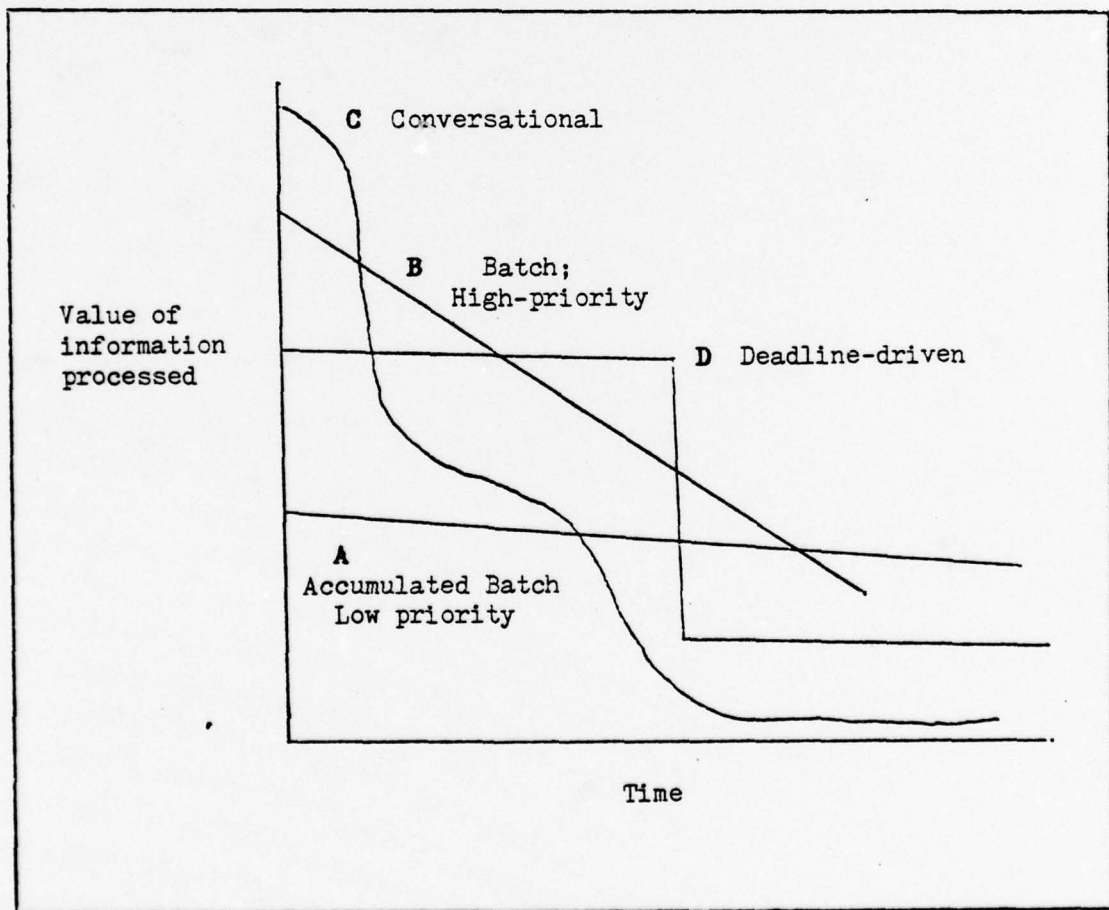


Fig 1. Computer Response Time (From Ref. 8:11)

If there is a deadline to getting a job completed and the throughput performance of the system is low then the job that needs to be executed may not be completed in time. Efficient use of computer terminals is necessary to minimize the expense of processing data and of retrieving meaningful results from the computer.

General Discussion

Batch and interactive terminals have advantages and disadvantages and a batch-interactive-mix analysis based upon the types of jobs run by an organization will provide results that can be used to efficiently use computer terminal resources.

Terminals that provide access to and from the computer are usually either batch or interactive. The term batch implies that the entire program or data is routed to the computer concurrently and the results are routed back to the terminal after the complete execution of the program. Interactive processing allows small segments of data to be transmitted to the computer with results being sent back to the terminal before more data is required.

When a batch job is submitted through the terminal, the job will be scheduled for processing according to the job size, according to the availability of the needed peripheral processors, according to the job priority, and in accordance with any other constraints the operating system designers may have employed. These scheduling constraints can cause delays, especially if large jobs require a large section of central memory. Several small jobs can tie up most of the central memory available, and when one job finishes there still may not be enough central memory available for the extremely large job that has been waiting in the job queue. This will force the job scheduler to by-pass the larger job because a smaller job in the queue is available for processing. If the

large job has a high enough priority, some of the smaller jobs might be rolled out of central memory to allow the larger job enough memory to begin execution.

Scheduling batch jobs can become complex. "Its [scheduling system] objective is to select jobs from all jobs available to provide a well-balanced mix, thus enabling the computer to use its resources efficiently, and at the same time meet all its deadlines (Ref. 3:27)."

Use of interactive terminals can minimize these scheduling problems because their demands for processing is recognized almost instantaneously due to the higher priority of interactive jobs. However, this creates a new significant problem. "With interactive processing ... system resources must always be fully prepared to handle a worst case situation (i.e., amount of memory needed)." (Ref. 8:14)

A single computer can handle hundreds of terminal connections, and the number of batch terminals versus the number of interactive terminals can significantly alter the operational costs of using the computer. The approximate ratio of batch to interactive terminals that should be used is dependent upon the computer system and the needs of the organization the computer services.

Problem Statement

AFSC is interested in the results obtained from batch-interactive-mix analysis performed on the AMIS jobs run on the IBM 370. The results might aid in making operational changes that will increase the throughput performance for AMIS users.

Scope

The objective of this investigation is to build a computer simulation model of the IBM 370 and System 2000 which will allow a batch-interactive-mix analysis to be performed and the results presented to AFSC. The model

will be driven by estimates of the workload characteristics of the AMIS jobs. A section is presented dealing with obtaining valid input data which will lead to more accurate results from the simulation analysis.

Approach

One way to evaluate and determine if an organization is achieving near optimal throughput performance by using the best batch-interactive-mix, is to construct a computer model of the system being used and to use input data that is representative of jobs being run. When large samples of workload characteristic data is needed, a computer simulation model will help the analyst arrive at meaningful benchmarks. The sheer volume of data, needed to drive the simulation model, would rule out the use of achieving the analysis by use of a mathematical probabilistic model.

Simsript II.5 is a computer language that is ideal for simulating the computer operating system because it allows the user to schedule events when a pre-determined simulation time has elapsed. Results from the Simsript computer simulation run can be used to evaluate the performance of the system. System performance actually involves two primary considerations. The first is the effectiveness of how the system handles a specific job or request; the second involves efficiency or how the system uses the resources available. Understanding both of these considerations is essential when analyzing the performance of the system. The most efficient use of system resources may not provide the best throughput performance for a specific job, but it certainly will provide good throughput performance over a given time period, with a substantially reduced operational cost.

The computer simulation model must be capable of handling certain workload parameters such as: job CPU time, job I/O request, CPU service

time, I/O service time, interarrival time, priority, memory requests, number of simultaneous users, number of jobs in the system, etc. (Ref. 29: 12-13). Without understanding these workload parameters, or using them improperly in the simulation model, the performance evaluation results will be invalid or, at best, somewhat misleading. Workload characterization of jobs and tasks should be fairly accurate if the workload parameters are representative of the jobs run by the organization.

Workload parameters can be combined, consequently producing data that is representative of the jobs and tasks. Through a clustering analysis technique the analyst produces data that is consistent, regardless of the possible extreme values of a given parameter. According to Anderberg,

"Cluster analysis has been employed as an effective tool in scientific inquiry. One of its most useful roles is to generate hypotheses about category structure. An algorithm can assemble observations into groups which prior misconceptions and ignorance would otherwise preclude. An algorithm can also apply a principle of grouping more consistently in a large problem than can a human. ... cluster analysis may be used to reveal structure and relations in the data. It is a tool of discovery. (Ref. 3:4)

Developing a simulation model which uses correct workload parameters and valid input data which allows the analyst to examine the systems performance is more than an intuitive art. Shannon [28] has suggested some criteria that are instrumental in developing a successful simulation model. The 11 suggested steps will help the simulation designer during the initial state of "system definition" through the "documentation" phase of the design. (Ref. 28:23) These steps will be discussed in detail in Chapter 3.

Understanding the system and analyzing the operation were important in building the Simscript II.5 model of the AMIS System. This enhances the development of a workload analysis methodology. Chapter 2 deals with the operating characteristics. Chapter 3 discusses the job processing

sequence and performance measures while Chapter 4 discusses the model formulation, logic flow and data preparation. The first phase of a performance improvement effort is that of understanding the computer operating system and the logical structure of the computer, which is the topic of the next chapter.

II IBM 370 Operating Characteristics

System Organization

"The IBM system 370 model 155 is a high-performance data processing system that provides the reliability, availability, and convenience demanded by business and scientific users, as well as by users with applications in communications or control." (Ref. IBM System/370 System Summary: 6-9)

The following sections: system organization, system control and System 2000 will be described using information from the IBM System/370 System Summary, IBM System/370 Principles of Operation and System Operation and Guide to Data Base Mangement.

Figure 2 shows the logical structure of the model 155 system.

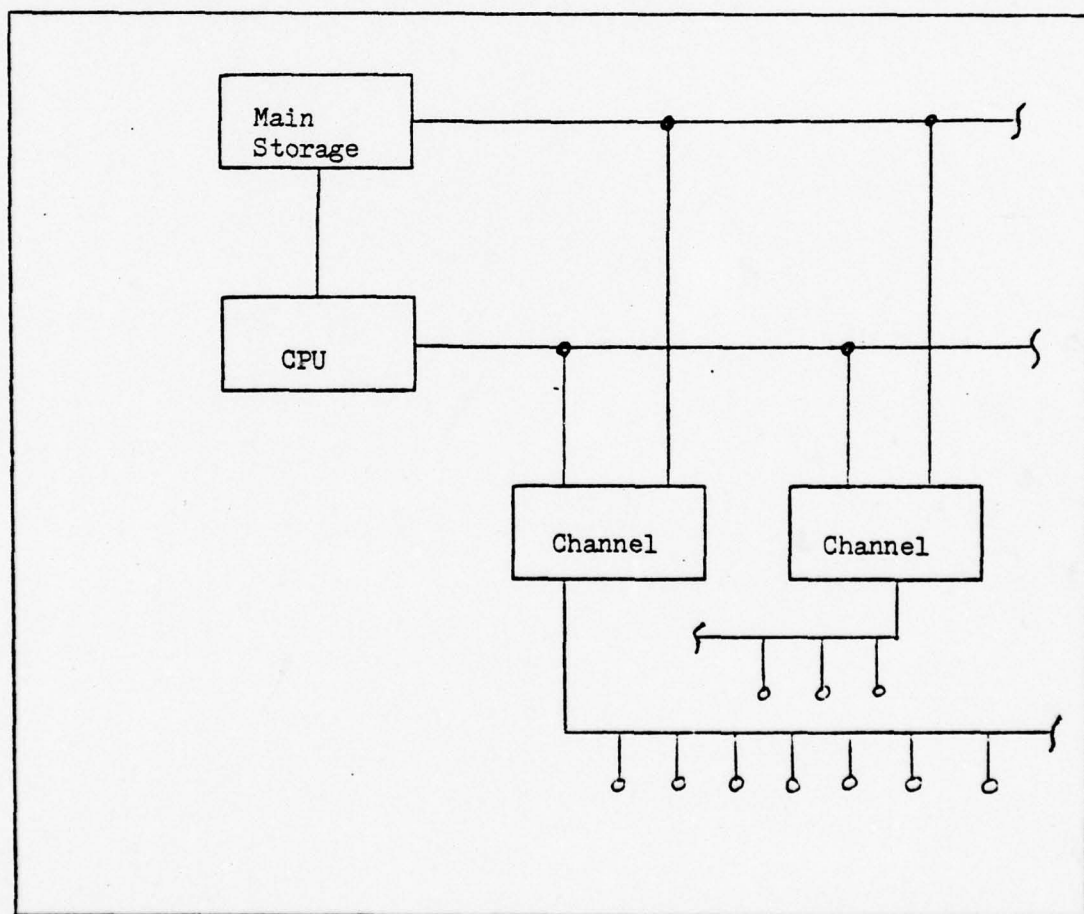


Fig 2. IBM System/370 Logical Structures (From Ref. 21:13)

The single CPU system logical structure is comprised of main storage, a central processing unit, a selector and multiplexer channels which communicate with each other by connecting paths.

Main Storage

Both data and program must be loaded into main storage before processing is allowed. The main storage is set up to provide direct addressable fast-access storage of data. The main storage is comprised of a large-volume access buffer called a cache.

In this type of a buffer storage system a Storage Control Unit (SCU) is placed between the processor and main storage unit and is illustrated in figure 3.

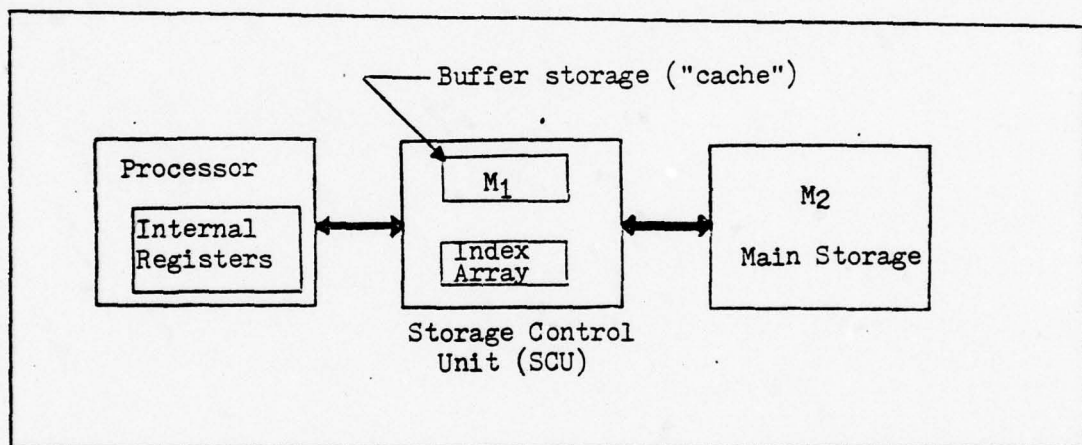


Fig 3. Organization of buffer memory hardware (From Ref. 14:192)

Exact copies of main memory (32 bytes at a time) are brought into the cache whenever a reference is made to an address that is not already existing in the cache. A block is stored over the existing block in the cache. Whenever a fetch request is generated, a check is performed to determine if there is need to bring in a different block from main memory. This algorithm is simple enough to be built into the hardware

of the system and allows a much faster fetch request from "copied" portions of main memory which actually reside in the buffer storage "cache" unit.

Central Processing Unit

The CPU is the controlling center of the computer system and handles system related functions such as execution, interrupts, initial program loading, etc. The CPU will handle 5 basic classes of instructions: system control, general decimal, floating point, and input/output instructions. The basic instruction sets include decimal, floating point, extended precision floating point, direct control, byte oriented operand, dynamic address translation, extended control mode, and store status and program reset.

Input/Output

Data is transferred between devices and main storage by attaching the I/O devices to channels, and the actual communication between control units and the specific channel takes place over a connector called the I/O interface. I/O devices such as card reader, punches, magnetic tape units, disc storage, drum storage, typewriter-keyboard devices, printers, teleprocessing devices and sensor-based equipment are handled by the associated I/O interface which provides information format and control signal sequences. I/O devices fall into several categories, most of which are used for:

Auxiliary storage.

Machine and manual (keyed) input, both local and remote.

Teleprocessing.

Reading (or output) of external documents and displays.

process control.

and data acquisition.

"An Input/Output operation transfers data between main storage and an I/O device. An I/O operation is initiated by a program instruction that generates a command to a channel. A control unit receives the command via the I/O interface, decodes it, and starts the I/O device. (Ref. 21:2-7)

Byte Multiplexer Channels

Channels are the direct controllers of I/O devices and control units, and allow the system to read, write and compute. This is accomplished while relieving the CPU of having to communicate directly with the I/O devices. "The byte multiplexer channels separate the operations of high-speed devices from those of lower-speed devices." (Ref. System Summary: 2-7) This allows channel communication to take place in one of two different modes: the byte mode using the slower data rate I/O devices and the burst mode using the higher data rate I/O devices.

When operating in the byte mode, the single data path can be used by several low speed I/O devices (such as card readers, printers and terminals) and each device takes turn in sending data over the multiplexed line. This is accomplished when the channel receives and sends data to the I/O on demand and is controlled by the channel program.

In the burst mode, I/O devices (such as magnetic tape units, discs, or data cell storage) are not under the control of the programmer and after these high-speed devices have established a logical connection with a channel, large amounts of data can be transferred in "bursts" which are not multiplexed.

Block Multiplexer Channels

The block multiplexer channels operate high-speed I/O devices on a single data path and can also operate in one of two modes: block multiplex or selector mode. In the block multiplex mode the channels permit interleaving of channel programs and allow initiation and

termination of these programs to occur sooner than is allowed by selector channels. The block multiplexing mode allows more data to be transferred during its burst mode than can be routed by the byte burst mode. Entire records and blocks can be transferred during each burst so that block multiplex channels are used with faster I/O devices than are used by the byte multiplexor channels.

When operating in the selector mode, I/O devices are attached to the selector channel which transmits data to or from a single I/O device at a time. These select channels can be operated with either the slow or high speed devices but are especially suitable with high speed I/O devices. Once a selector channel has attached a particular I/O device it will transmit data until all data has been handled and no other I/O device can interfere with the selector channel.

Information about the hardware configuration has been presented and it is hoped that the reader is a little more familiar with the computer system. General information concerning the software of the system is discussed now to further aid in the understanding of the system.

Job Scheduler

All jobs are classified into 11 membership classes. The requirements for each class is listed in table 1. where K is the region size in memory (in blocks of 1024 bytes) and T is the CPU time in minutes. The job class represents the programmer's estimate of the resources required for the job and is used to schedule and prioritize jobs. The following classes are used on the IBM 370/155 by A.S.D.

Table 1.
Scheduling Priorities

Class	Region	Time	Tape Mounts Permitted	Disk Mount Permitted
A	$K \leq 200$	$T \leq 5$	NO	NO
B	$K \leq 260$	$T \leq 60$	NO	NO
and (K 200 or T 5), I.E., JOB CANNOT BE RUN IN CLASS = A				
C	$260 < K \leq 500$	$T \leq 60$	NO	NO
D	$K \leq 200$	$T \leq 5$	YES	NO
E	$K \leq 260$	$T \leq 60$	YES	NO
and (K 200 or T 5), I.E., JOB CANNOT BE RUN IN CLASS = D				
F	$260 < K \leq 500$	$T \leq 60$	YES	NO
G	$K \leq 200$	$T \leq 5$	YES	YES
H	$K \leq 260$	$T \leq 60$	YES	YES
and (K 200 or T 5), I.E., JOB CANNOT BE RUN IN CLASS = G				
I	$260 < K \leq 500$	$T \leq 60$	YES	YES
J	$K \leq 500$	$T \leq 60$	YES	NO
L	$K \leq 500$	$T \leq 60$	YES	YES

Multiprogramming

Once a job has been assigned a class priority it is also ranked by priority within that class. The highest priority job will always have access to the CPU under this multiprogramming scheme. Once a job has accessed the CPU it will continue being processed until the job must relinquish the CPU because of an I/O or until a higher priority job has been scheduled or a higher priority job has completed its I/O. The job that was interrupted is put into a wait state and must compete once again with all jobs available to be processed in accordance with the priority established in Table 1. The amount of I/O a job performs does not affect the priority algorithm and this multiprogramming scheme works well if the higher priority jobs are I/O bound while the lower priority jobs are CPU bound.

Because there is no time slicing, the lower priority jobs must wait until all the higher priority jobs are handling I/O and this means that these jobs will be waiting for long periods of time in between opportunities to access the CPU. CPU utilization will usually be high with good turnaround time for the higher priority jobs. Degredation of turnaround time will occur for the lower priority jobs. Fortunately the AMIS jobs perform a lot of I/Os when run on the System 2000 data base management system. The large number of I/Os allow the lower priority jobs to frequently access the CPU.

System 2000

The system 2000 software package supports the AMIS requirement of storing and modifying contracts for Air Force Systems. This system simplifies the actual storing and accessing of contract related data.

"A data base is generally acknowledged to be a collection of multiple logical files containing interrelated but nonredundant data which can be accessed by one or more applications. A data base management system is a software tool -- actually a collection of routine -- used to define and maintain the data base's logical structure, and provide a means by which data can be retrieved." (Ref. 4:1)

The system 2000 data base is comprised of 5 functional elements: Data Bank, Data Dictionary/Directory System, Data Base Administrator, Data Base Management System and the User System Interface. The inter-relationship between these functional elements is shown in Figure 4.

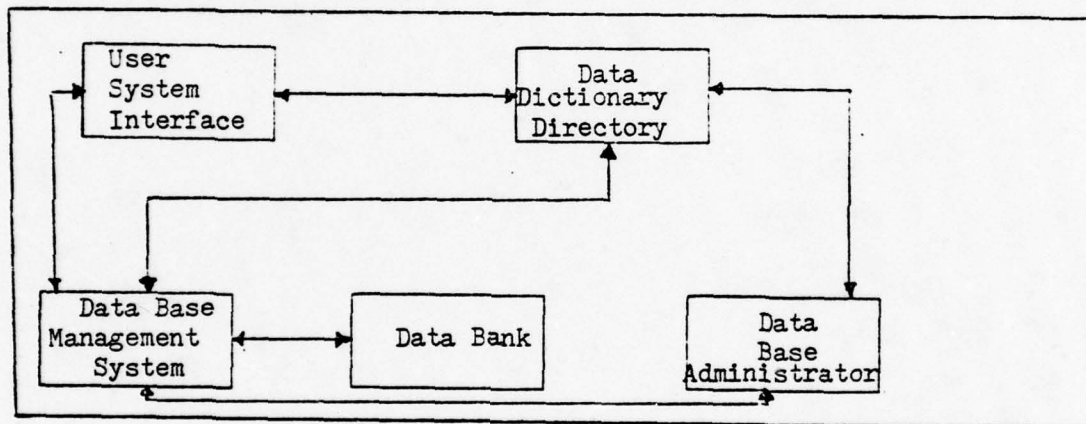


Fig 4. Data Base Management: Conceptual Environment (From Ref. 4:5)

Data Bank

The data bank is comprised of a collection of data bases organized to provide maximum performance of the system. The physical location of these data bases is not important as each is logically connected with the Data Dictionary/Directory System which is capable of coordinating centralized or decentralized data base files. Queries, updates, referencing, etc. can access the Data Bank through the Data Base

Management System.

Data Base Management System

The Data Base Management System consults the Data Dictionary/Directory System "... for information about data to provide the parameters necessary for the generalized software/hardware to execute user request." (Ref. 4:7) The Data Base Management System is the actual software package for accessing and storing data.

Data Dictionary/Directory System

The Data Dictionary/Directory System has two objectives. The first is the collection and dissemination of data which supplies the user with meta-data (data about data). The second function of the Data Dictionary/Directory System is that of establishing standards for coding conventions, data naming and usage.

Data Base Administrator

The major responsibilities of the Data Base Administrator (Ref. 4:6) are:

- Definition of the content and structure of the data base.
- Control of data access and modification rights to the data base.
- Advising data base users on efficient techniques for extracting data.
- Establishing data entry, edit, and validation standards.
- Maintenance of the Data Dictionary/Directory System.
- Maintenance of the Data Base Management System.
- and keeping track of available physical storage.

User System Interface

To adequately serve its users the data base must use interfaces that allow for 3 main areas in the User System Interface Support.

The first is the language capability. The use of a natural language will help the user formulate requests and problem definitions that are easier to use when communicating with the system.

The second, "Interactive Capability -- the ability to browse in search of solutions supports the user decision process by providing the means to develop new alternatives." (Ref. 4:7) The last is the "Auxiliary Subsystems -- use will provide subsystems to assist users in putting system output into the most humanely comprehensible form. This category includes graphic techniques, algorithmic processes, and modeling and/or simulation tools." (Ref. 4:7)

The first phase in improving the performance of the computer system is that of that of understanding the system in terms of hardware configuration and related software programs in use. The second phase involves the analyzing of operations and understanding the system requirements, definitions and performance measures.

III System Requirements, Definitions and Performance Measures

Job Processing Sequence

The system simulation program that has been constructed deals with the following basic job processing sequence as outlined by MacDougall (Ref. 26:191-192). The simscript program itself is located in Appendix B.

1. Job arrival.
2. Request for central memory (CM) - if available, allocated, if not, the job is entered into the CM queue.
3. Request for central processor - if available, the job is assigned to it and executes until an I/O request is encountered or until execution completes or a higher priority job interrupts the CPU. If the central processor is not available then the job is entered into a CPU queue.
4. Request for an I/O - the job is released from the central processor and if another job is waiting for the CPU, the job waiting will be assigned to the CPU. If the disc is free, the disc is assigned to the job to process the I/O request; if the disc is busy, the request is entered in a queue.
5. On completion of processing of an I/O request, the disc is released and the central processor requested once again. (When the disc is released, the disc queue is checked; if there is a waiting request, it is assigned to the disc.)
6. When a job completes execution, it releases the central processor and its central memory space is released. (The CM queue is checked to

determine if there is a job waiting to which space now can be assigned, and the CPU queue checked to determine if there is a waiting job which now can be assigned to the central processor.)

7. The job leaves the system and the sequence is repeated.

The above processing sequence is modified slightly for interactive jobs because the interactive job can be thought of operating in either in one of two states. (Ref. 9:172) The think state is the time between a computer's response to the terminal request and when another request is made by the user, such as the hitting of the carriage return. The system state is the time from the user's request until the computer has processed that request. When the processed request has been completed, an "interaction" has occurred and the interactive process has once again entered the think state.

Buzen [9] has verified that the average response time can be calculated by the following formula:

$$R = \frac{1}{J} \sum_{k=1}^N r(k) \quad (1)$$

where,

" R = Average response time (i.e., average amount of time in system state per interaction).

r(k) = total time that the k-th interactive process (i.e., the interactive process associated with the k-th terminal) spends in system state during the observation interval (k=1,2,...N)."

J = Number of interactions completed during the observation time interval.

Buzen also claims that the average think time can be calculated by:

$$Z = \frac{1}{J'} \sum_{k=1}^N z(k) \quad (2)$$

Where,

" Z = Average think time (i.e., average amount of time in think state per transition think state to system state).

J' = total number of transitions from think state to system state during the observation interval.

z(k) = total time that the k-th interactive process spends in think state during the observation interval (k=1,2,...N)." (Ref. 9:172-173)

Since this interactive process occurs, the simulation model must account for this response time, and during the system state points 3 through 5 of the job processing sequence justify considering the interactions as "separate jobs".

Workload Parameters

Within the basic job processing sequence the simulation program handles the following workload parameters as modified from the list compiled by Svobodova. (Ref. 29:12-13)

Job Cpu Time	Total Cpu time requested by the job or CPU time requested by a single interactive command.
Job I/O Requests	Total number of I/O requests needed by a single job.

CPU Service Time	Time required to process a single CPU operation.
I/O Service Time	Time to complete or process a single I/O task.
Interarrival Time	Time between two successive requests for any given resource.
Priority	Priority assigned to a job by the algorithm shown in Table 1.
Blocked Time	Time the job must wait for the CPU service in a wait state.
Memory Requests	Amount of core required by the individual job.
User Response Time	Time it takes the user to submit another request after a response from the CPU.
Number of Simultaneous users	The total number of interactive users concurrently logged on.
Number in the system	Total number of batch and interactive jobs operating within the system.

The Simscript programming language allows an entity to possess certain attributes which allow the simulation program to "move" a batch or interactive job "through" the system, carrying with it the necessary workload parameters to analyze the job workload characteristics. As the job moves through the system - enters queues, is assigned to the central processor, etc. - the job carries with it all the information needed throughout the simulation process. Once the "job output simulation"

occurs the individual job is destroyed in the system but the statistics about the job are kept by use of a TALLY STATEMENT. "The Tally Statement computes statistical quantities and prepares histograms for time-independent variables."(Ref. Simscript Manual) This allows the simulated job to be destroyed while specific data collection counters and routines to compute statistical quantities are used in conjunction with an ACCUMULATE statement to determine statistical analysis of the collective job workload characterizations. These statistics are useful in determining some important system performance measures, such as those listed by Svobodova. (Ref. 29:16-18)

Performance Measures

Throughput is a good measure of system responsiveness and since it is the amount of useful work completed within a given time period, given a particular workload, there are several performance measures that will give an indication of the system throughput performance. The turnaround time is the elapsed time between submitting a job and the time the results are received at the printer or terminal. The turnaround time can be a good indication of throughput performance. The elapsed time multiprogramming factor (ETMF) is a numerical value calculated by dividing the turnaround time of a particular job run under multiprogramming by the turnaround time of the job had this been the only job in the system. The ETMF is a measure of only one job in the system where gain factor is a measure of several sequential jobs.

The gain factor is determined by finding the time needed to execute a set of jobs under multiprogramming divided by the total system time needed to execute the same jobs sequentially without the capability of multiprogramming. Since types of jobs programmed may

vary considerably from day to day and during different time periods within a day, it is important that the gain factor be based upon results obtained from different time periods. Another measure of system performance is the CPU productivity or CPU utilization.

The CPU productivity is the percent of time the CPU is in use performing useful work. Multiprogramming should allow the system to perform with high CPU utilization because as one job has completed a CPU task and relinquished the CPU for an I/O there will usually be a job that has been waiting for the CPU in the CPU queue. Of course there is some overhead time, which is CPU time required by the operating system. CPU productivity could be very low if the organization jobs are consistently I/O bound.

The wait time for I/O is the time necessary to process an I/O task and therefore, if all the multiprogrammed jobs are involved in I/O tasks, the CPU will be idle until one of the jobs completes its I/O. If this occurrence of I/O bound jobs is common, there may be many times the CPU is idle and hence the CPU productivity will be low. If the CPU productivity is constantly low then the availability of the system will be high.

The availability of a system is the percentage of time the system is available to the users. Low availability will cause the external delay factor to be high (job turnaround time/the total CPU processing time required) and the throughput to be low.

A computer simulation model that uses good workload parameters and is capable of analyzing throughput performance with adequate performance measures will yield results that will help to remedy poor CPU utilization and improve the turnaround time for the organizational users. A reliable

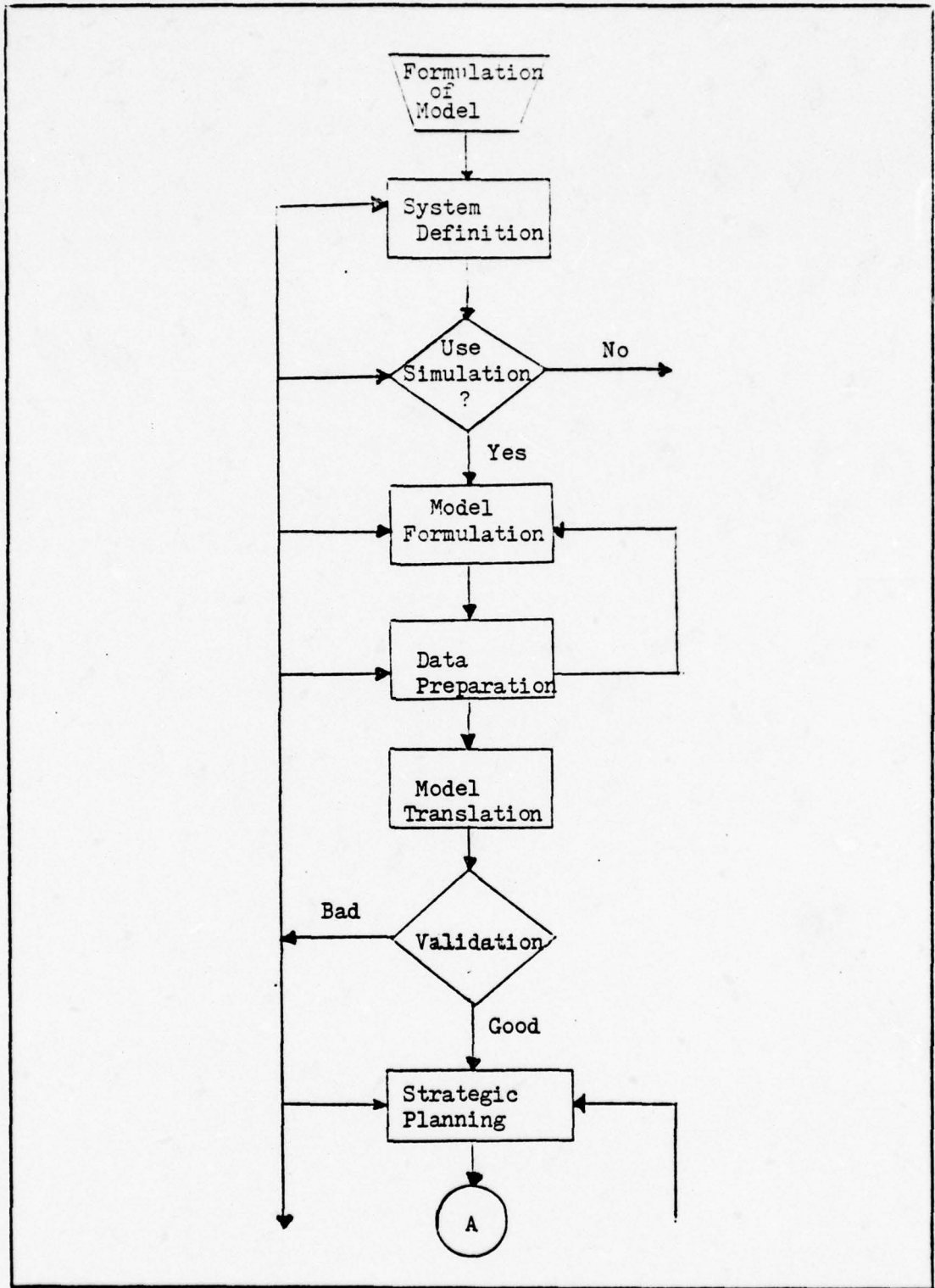


Fig 5A. Simulation Process Flowchart (From Ref. 28:24 Fig 1.3)

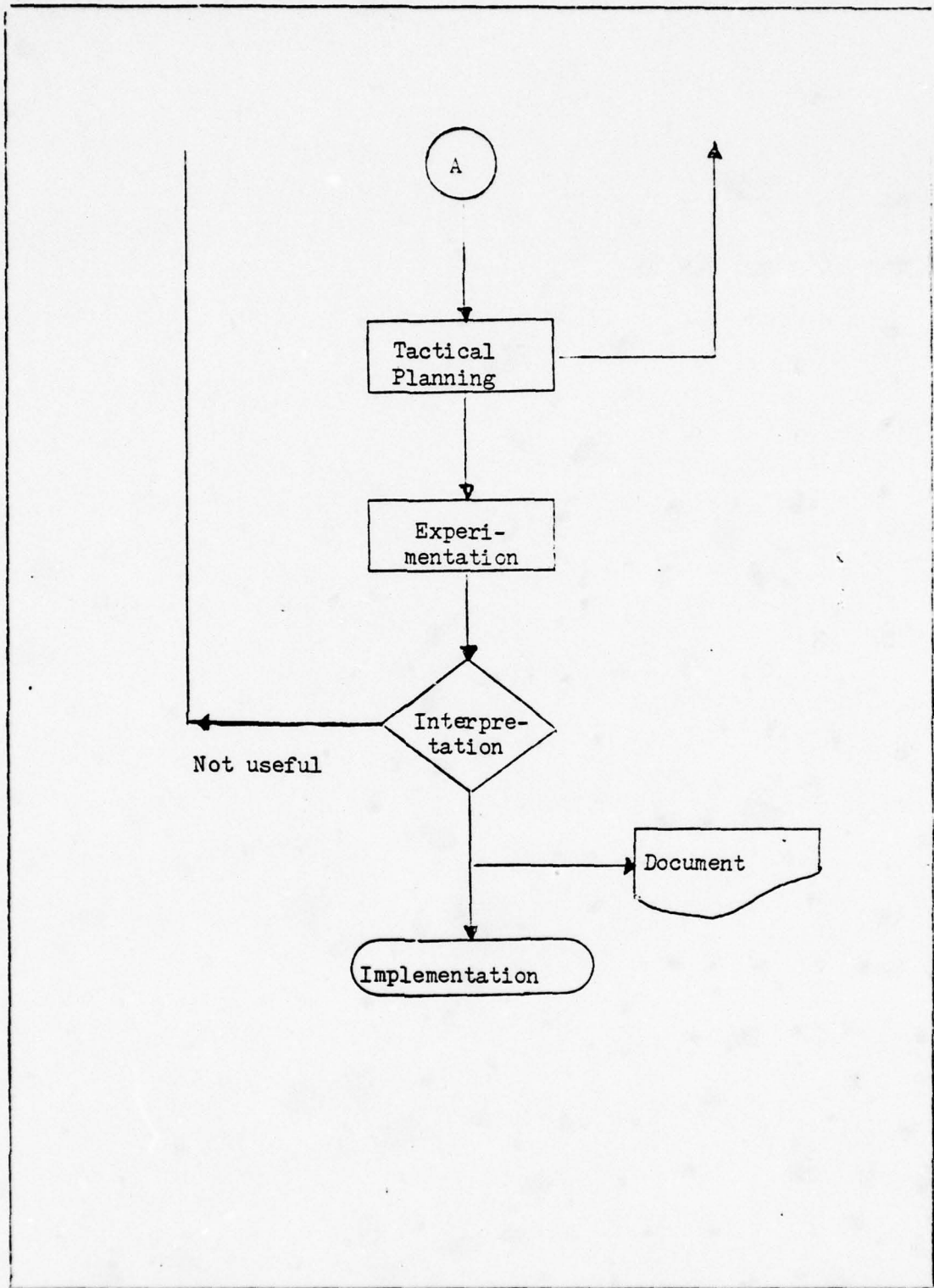


Fig 5B

and accurate model is constructed by following the steps in the flowchart of Figure 5. (Ref. 28:24)

Development of Simulation Model

The flow chart on pages 24 and 25 shows the necessary stages that must be considered and implemented in order to achieve the construction of a reliable simulation model. The 11 steps that Shannon has incorporated into this process will be described in an abbreviated format. (Ref. 28:23)

- 1) System Definition - The Simgscript II.5 model was developed within the boundary of considering the CPU time, I/O time, priority of the job and the need for disc or tapes. The measures of effectiveness of the model include: job turnaround time, CPU utilization, and gain factor.
- 2) Model Formulation - Reduction of the real system to be simulated into a logic flow diagram. This is shown in Figure 6A and 6B.
- 3) Data Preparation - Identification of data needed by the Simgscript model which is described on page 33.
- 4) Model Translation - Description of the model in the Simgscript language.
- 5) Validation - Increasing the level of confidence in the model and the results the model generates. This is discussed on page 35.
- 6) Strategic Planning - Design of an experiment that shows how the gain factor and CPU utilization varies with a change in the workload of the system.
- 7) Tactical Planning - Determining how the two experiments are to be run and tested.
- 8) Experimentation - Execution of the simulation model to perform these experiments and to perform a sensitivity analysis of the system. The sensitivity of the model is presented on page 38 and a

discussion of the experiments is located on page 40.

- 9) Interpretation - Drawing inferences from the data gathered from the experimentations. This will be done in the conclusion section of this paper.
- 10) Implementation - The results of these experimentations will be presented to AFSC.
- 11) Documentation - The results of the experimentations are listed in Appendix A.

The formulation of the problem has been presented in the Introduction while the system definition was presented in Chapter 3. It is now time to consider the use of a simulation language and to look at what is involved in the Model Formulation stage of the simulation process.

IV Simulation Steps

Model Formulation

Having discussed the problem to be investigated and showing the need for a computer simulation analysis, Shannon (Ref. 28:24) suggests that the next step concerns the reduction of the real system to a logic flow diagram that accurately depicts the process to be simulated. This flow approach to analyzing the IBM 370/155 will allow the workload parameters and performance measures to be incorporated into the simulation model. Figure 6 on pages 29 and 30 is a flow diagram of the necessary events to simulate the job processing of batch and interactive jobs.

Once the real system is represented by logic flow, the necessary data input and starting conditions are determined. Some data generation is internal to the program such as pseudorandom numbers and stochastic variates and must be generated appropriately. Once it is known how the real system should be represented, it is important that a simulation language be chosen that will allow a meaningful description to be implemented in a computer simulation program. Unfortunately, the best simulation language for a particular investigation is overlooked simply because the analyst may be familiar with a specific language already and may feel that too much time and effort would be needed in determining which language is best and then having to learn that new language.

Heidenreich and Blitt (Ref. 7:38) have listed the advantages and disadvantages of different languages by comparing 9 languages including Simscript and GPSS simulation languages. Their comparison

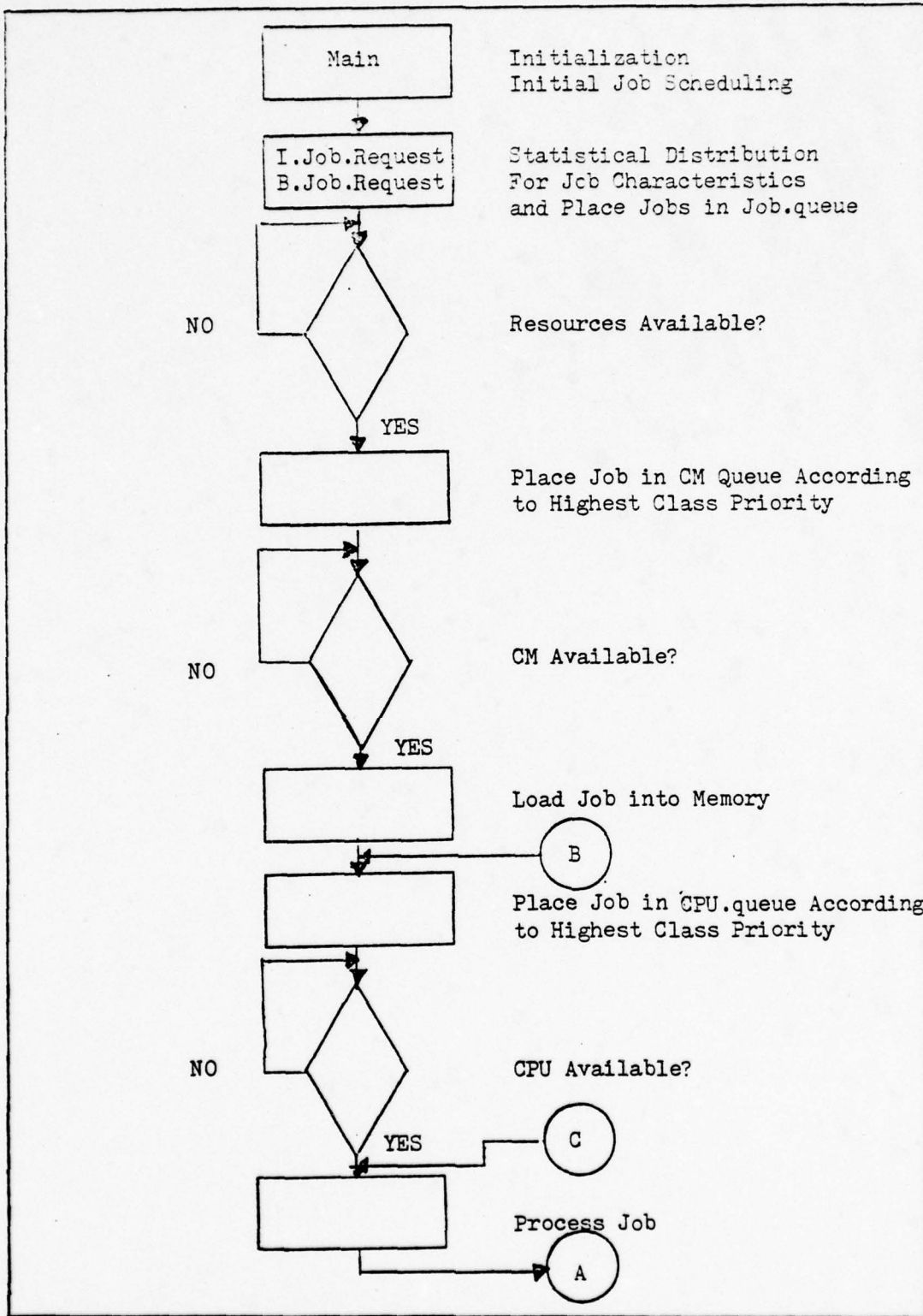


Fig 6A. Job Processing Sequence

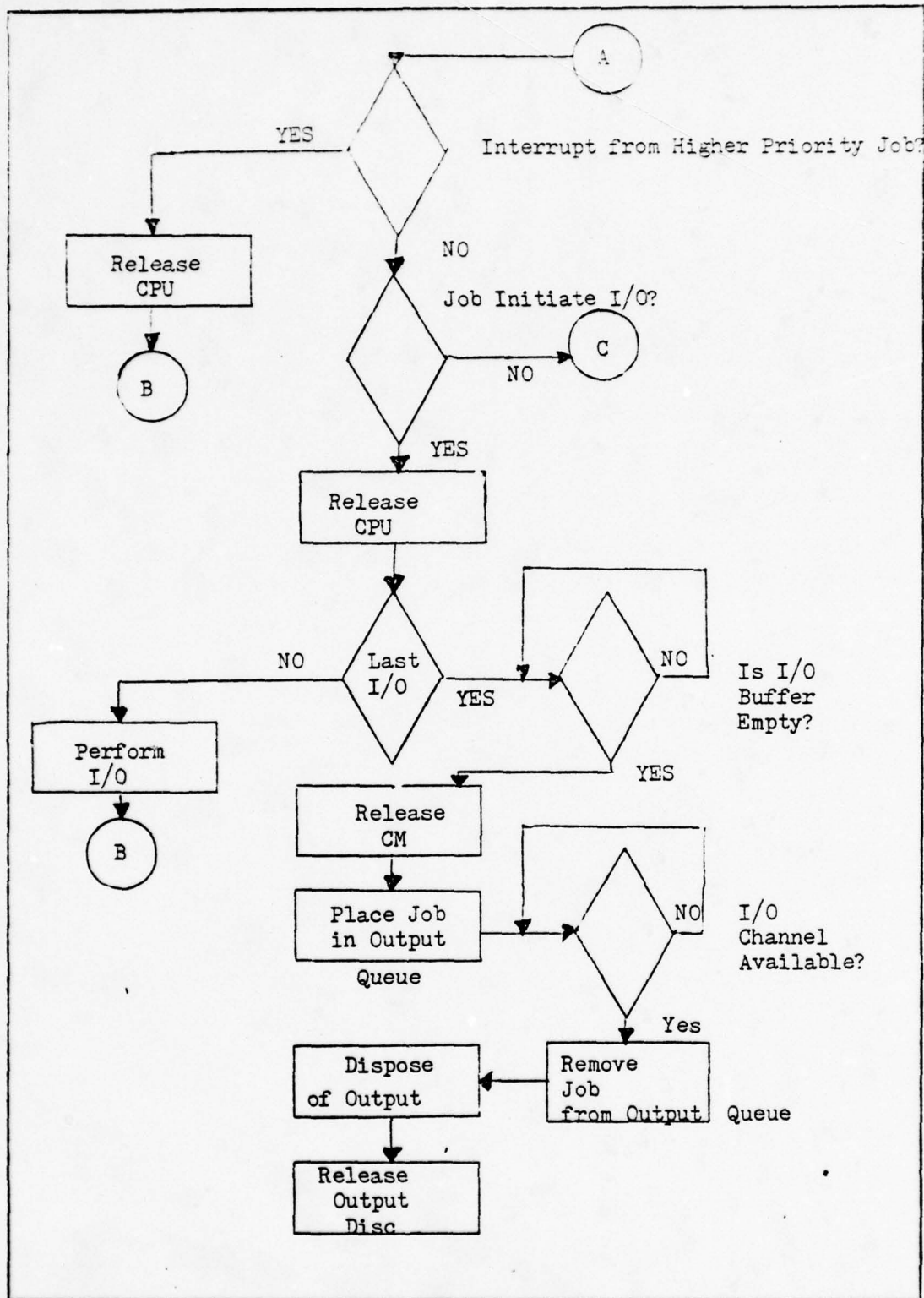


Fig 6B.

suggests that Simscript is designed for simulation because it is event oriented but is slow to execute and not well known by its users. GPSS is also a simulation language but appears to be easier to learn than Simscript. A prime consideration between the languages might be simply "Which one is available to the investigator?". Simscript was chosen because it is suited to general programming and discrete-event simulation modeling, and because of familiarity with the language.

Logic Flow

Again, figure 6. shows the logic flow for the continuous simulation model which represents the AMIS system employed on the IBN 370/155. The actual program is listed in Appendix B. and a description of the logic flow will coincide with the events in the program model.

The "Main" routine is used to establish initialization of variables and to schedule the first interactive job and the first 6 batch requests. All simulation time is based upon parts of a Day, so all units must be established here in "Main" to avoid errors that result from working with wrong units. Once the variables have been initialized and all events scheduled, the "... 'Start Simulation' statement begins the simulation by passing control to the timing routine which removes the first event notice from the event set and executes that event." (Ref. 28:210)

Next, statistical generation of job input parameters is accomplished in the I.JOB.REQUEST (Interactive job request) and B.JOB.REQUEST (Batch job request) events. The individual job CPU time, I/O time, number of I/O's per job and carriage return, core requirements and job priority are generated within these two events. The job priority is determined by the algorithm shown in Table 1. and is based solely upon the amount of core requested, CPU time, and tape mount and disc mount requests. Once

the job characteristic parameters and job priority have been determined the 'job is placed into the job.queue'. The individual "job" entity has certain job attributes (job parameter values listed in the preamble) which are "moved" with the job throughout the simulation process by use of pointers when referencing the job. Once the job has been placed into the job queue the job scheduler scans the jobs in the queue to determine which job will get access to the central memory.

"... the job scheduler chooses a small subset of the jobs submitted and lets them 'into' the system. That is, the job scheduler creates processes for these jobs and assigns the processes some resources. It must decide, for example, which two or three jobs of the 100 submitted will have any resources assigned to them. The process scheduler decides which of the processes within this subset will be assigned a processor, at what time, and for how long." (Ref. 14:211)

This processing management checks to see if necessary devices are available, creates processes for the job and then assigns memory when possible. The request for memory is based upon the job class priority (and the priority within a class), and when memory is available the job is loaded into core to wait for access to the CPU.

The job is now in the CPU.queue; this queue being formed every time the CPU is not available for job processing. When the CPU is released a job is selected from the CPU.queue according to the highest class priority and scheduled for the event "CPU.processing".

Once the job is processing, the job may be pre-empted by an interrupt that is initiated by a job which has a higher priority. At this point the CPU is released and the current job being processed is placed back into the "CPU.queue" and the higher priority job is scheduled for the event "CPU.processing". If the job being processed initiates an I/O, the job will release the CPU and the I/O task completed before the job is re-inserted into the CPU.queue (unless of course this is the job's last I/O).

If the job has completed its last I/O then as soon as the I/O buffer is empty, the core relinquished by the job may be used by another job waiting in the CM queue. When the I/O channel becomes free, the job will be placed in an output queue and appropriately disposed to the output terminal. After the job has been disposed to the output terminal the output disc is released for use by other jobs.

The event AN.DISC.RELEASE "destroys" the individual batch or interactive job but tallies up statistics about the job before removing the job's attributes from the simulation program.

Data Preparation

In order to simulate the job processing accurately there are five major parameters that must be generated for each job; namely, the amount of I/O time, the need for discs, the need for tapes, the amount of core needed and the amount of CPU time to process the job. To insure that sample jobs with these characteristics are selected and truly represent the job workload of the system would require an additional effort. This additional effort would require more time than has been allocated for this study. It was decided that the average statistics of S2K jobs (AMIS major production jobs, which account for about 25% of the workload) would be adequate to drive this simulation model. In the recommendations section of this paper the use of cluster analysis will be presented as a means for acquiring sample inputs which would reduce the amount of data that would need to be read into the program during execution.

Programming Techniques

The discrete-event capability of Simscript II.5 allows for good modularity throughout the simulation program. Modularization allowed the program to be divided into subprograms (events) which are called when the

timing routine has been scheduled. With this type of event, the need for global flags was minimized and consequently the hierarchical structure was forced to simulate the real system more realistically. Without the need for global flags the modifiability of the model is enhanced as well.

Modifiability implies that when changes are made to a portion of the program the changes in one subroutine make few, if any, changes in other subroutines throughout the program. This modifiability principle was used in constructing the two events which "parameterize" the batch and interactive jobs.

The internal generation of job attributes can easily be changed to read data as is recommended in the last chapter of this paper. This simplifies the transition process of modifying the program. Another way to make changes easy is to use the principle of understandability.

Simscrip II.5 allows variables, routines and events to be in alphanumeric form and hence the principle of understandability is strengthened throughout the program. As long as the first 5 characters differ, there is no confusion when transferring execution from one event to another event or routine. This means that the names of the labels, events, routines and variables can be spelled out in a string of short words separated by "periods" when coding the program. Understandability is not merely a property of legibility as the entire conceptual structure of the model is involved.

The principles of modularity, modifiability, and understandability were achieved by using top down design methods in building the conceptual structure of this model. After the model had been built and debugged it was necessary to validate the model.

Validation of the Simscript Model

The investigator must have confidence in inferences and results obtained from the computer simulation results. Confidence in the model is established if the results from the model compare favorably and accurately with the real system results. If the model results vary drastically from those of the real system then changes must be made in the structure of the model or in the variables and parameter estimates that were used in the data preparation stage. If the model is simple to understand and easy for the user to control and manipulate then the analyst will find this iterative process manageable.

Before inferences can be drawn from the validated model, certain assumptions must be made. The input parameters must be as accurate as possible before a high level of confidence in the simulation results is achieved.

Assumptions

The AMIS jobs account for 90% of the workload on the IBM 370/155 and during a given typical week of November 1977, this would amount to about 12-15 hundred jobs. During this period the S2K jobs (one type of AMIS major production jobs) accounted for about one quarter of the jobs run. Data about these jobs revealed that of 266 interactive jobs an "average" job initiated 2,739 I/O's needing about 44 CPU-seconds for execution of the job. The typical batch job initiated 1,215 I/O's with about 34 CPU-seconds required. Since I/O time is not needed to determine class priorities it is difficult to ascertain close figures for job I/O times. It is known that the AMIS jobs are highly I/O bound jobs because of the time needed to search the data bases, so a ratio of 20 to 1 was chosen to represent I/O vs CPU time for batch jobs, and a

ratio of 4.5 to 1 for interactive jobs was selected. Assuming that the S2K jobs are "representative" of the workload, a 4 hour simulation time period was chosen to gather data for analysis. Again, the time allocated for this study did not allow for an analysis of data beyond this single 4 hour period.

Assuming that 28 batch jobs are scheduled during this 4 hour period and that 40 interactive jobs are run during this same time, the initial base line run showed that the model reflected a CPU utilization of 87% with a gain factor of 2.3684. (See definition of gain factor on page 22) This compares favorably with the 85% CPU utilization of the real system. Appendix A contains information about the baseline run and all other subsequent runs as well. The results were calculated after a warm-up period that is explained and discussed next.

Steady State

Shannon [28] recommends three possible ways to eliminate errors from simulation results that were introduced early in the warm-up of the system simulation.

- 1) Run the computer simulation long enough and the initial warm-up errors will be absorbed into an accurate average of the collected statistics.
- 2) Choose initial starting conditions that accurately reflect a given operating time period which will reduce errors during this transient period.
- 3) Throw out or just don't calculate statistics during an appropriate warm-up period.

The method employed in this investigation merged ideas from both

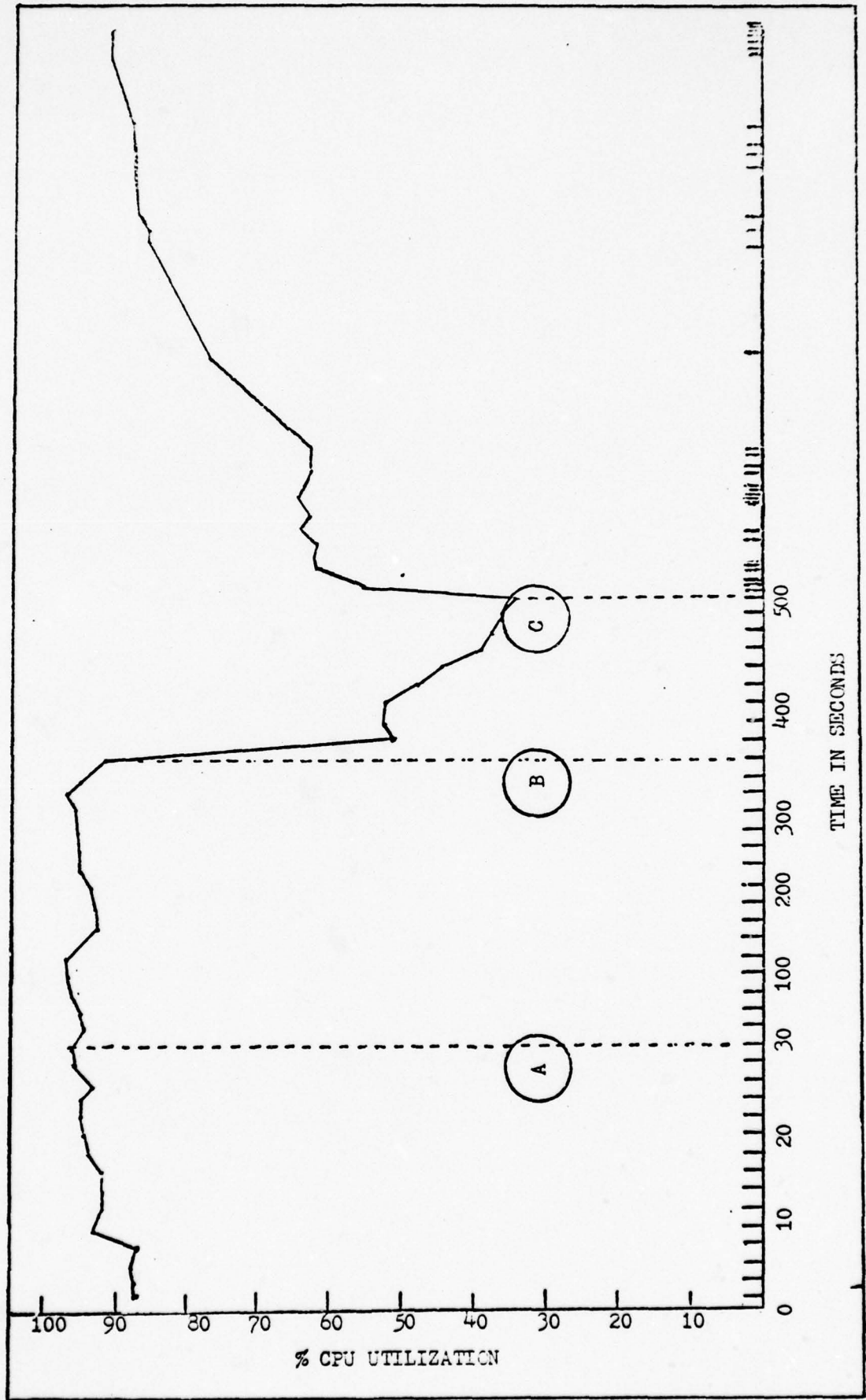


Fig 7. Transient Warm-up Period

points 2 and 3 above. The initial starting conditions provided for the scheduling of 1 interactive and 6 batch jobs to begin execution within the first minute of simulation time. After about 5 3/4 minutes a flag was set to start gathering job workload statistics (Figure 7 on page 37). shows what effect this warm-up period had on CPU utilization of the system. The time interval between "start simulation" and time A represents 30 seconds, while between point A and C each division represents a 20 second time interval. The cumulative CPU utilization remains fairly constant up to point B (approximately 5 3/4 minutes) when a flag was set to recalculate the cumulative CPU utilization. From point B to point C it can be seen that larger variations occurred. Point C represents that last regular interval where the CPU utilization was calculated and every point beyond C reveals the cumulative CPU utilization when any given job was completed. It appears that the cumulative CPU utilization stabilizes in the neighborhood of about 82 to 88 percent. By using point B as the starting reference point the results are weighted negatively by the rapid drop from point B to point C; but the immediate rise following point C will help keep this negative weight factor from drastically altering the true cumulative CPU utilization as it stabilizes near the end of the 4 hour run. Now that the transient warm-up errors have been reduced, a sensitivity analysis of the model's predictions is appropriate.

Sensitivity Analysis

A total of 29 simulation runs were made to determine what effect changes in I/O and CPU time would have on the cumulative CPU utilization and gain factor. The first graph on page 52 shows how a workload increase will affect the two variables. In this graph, as well as the rest of the graphs in Appendix A, the vertical dashed line represents the initial base

line run. It appears that the gain factor is on a steady increase and that CPU utilization falls below the baseline value of 86.67% after an increase of 10% in CPU and I/O time.

The chart on page 53 shows what effect an increase in I/O and decrease in CPU time will have on the variables of interest. The CPU utilization is on a "saw tooth" decline while the gain factor is also on a "saw tooth" rise which seems to rise when the CPU utilization decreases and falls when the CPU utilization increases. The baseline run is at a relative maximum point for CPU utilization with a slightly lower than average point calculated for the gain factor.

The chart on page 54 shows what effect changes in CPU time will have on the results when the I/O time is kept constant. The CPU utilization is on a constant rise until the CPU time exceeds the neighborhood of a plus 10% increase, at which point the CPU utilization decreases by about 5%. The baseline run reveals that a 10% increase in the CPU time of all jobs run will allow a 3% increase in CPU utilization while decreasing the gain factor by a factor of .2927.

The results on page 55 show that when keeping the CPU time (of all the jobs run) constant, an increasing change in I/O time will cause a gradual decrease in CPU utilization until going beyond a 10% increase, where this variable drops sharply. The gain factor remains on a gradual increase throughout the increase in I/O time.

The next eight charts, page 56 through page 63 show variations of changing the CPU and I/O times of all the jobs run. These graphs reveal that the baseline run is slightly below or generally above the CPU utilization figures and that the gain factor is at about the mid-point or generally below the values calculated for the other runs. These sensitivity runs

show that there are only two main areas where the results vary drastically from the baseline run. The first is where the CPU time is increased by 10% and the I/O time decreased by 10%; the other is where the CPU time is decreased by 20% and the I/O time is increased by 20%. Using the workload characteristics of the baseline run 10 more computer runs were made to help in the analysis of the batch-interactive-mix problem.

Experimentation

Two separate single factor experiments were run to determine what effect an increase in the interactive workload and a decrease in the number of ports available to run interactive jobs would have on the CPU utilization and gain factor. Five runs were made for each investigation with the intent that each separate run be considered a datum point.

The chart on page 64 shows that as the number of interactive jobs increase up to an additional 50% increase in the interactive workload, the CPU utilization climbs by about 8 1/2% while the gain factor decreases by a factor of .9550. There is every indication that the CPU could be used more efficiently but at the same time it appears that a decline in throughput and turnaround performance might result. The results on page 65 (decrease in the number of interactive terminals) are just about the same as those obtained when increasing the interactive workload.

Allowing just one factor (i.e., number of interactive terminals or interactive workload) to vary with each computer run provides a limited analysis concerning throughput performance. Ideally interactions between these two factors should be considered but due to the cost of each computer run it was decided that these extra runs would not be made. As an aid to the reader, the following discussion will help in the determination of the

important factors and the number of runs to be made during a simulation analysis.

Design Tradeoffs

Shannon [28], (Section 4.6) discusses a tradeoff study between the number of factors (input parameters, or variables), number of factor levels, number of replications and total number of computer runs required. Figure 8 shows a nomograph and the dark arrows show how to enter the nomograph and calculate the expected total computer cost and number of computer runs to be made.

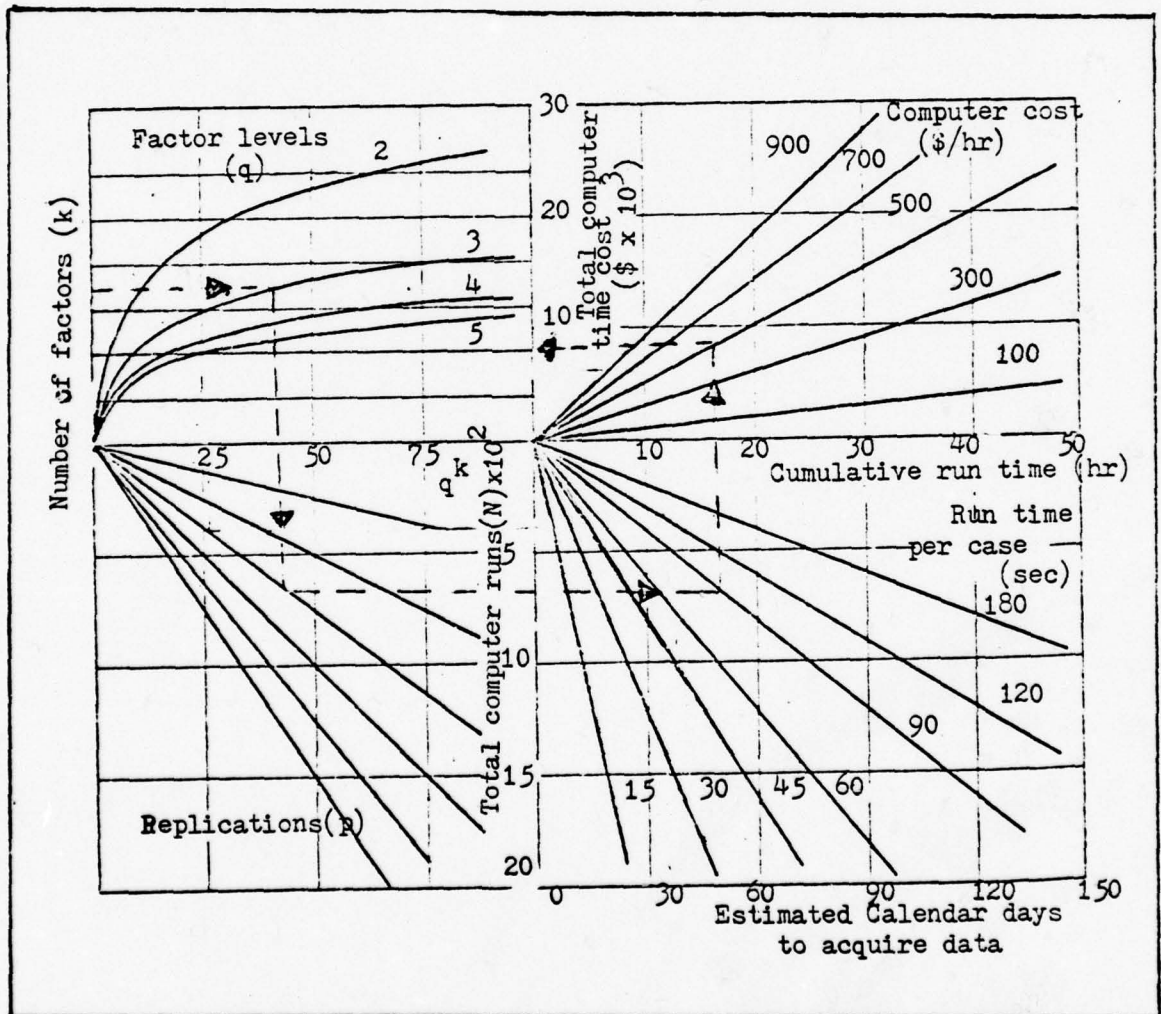


Fig 8. Nomograph for cost estimates (From Ref 28:156 Figure 4.2)

Where

k = number of factors (input parameters or variables)

q = number of factor levels

p = number of replications

N = total number of computer runs required

In section 4.6, Shannon lists 3 equations for determining which of the variables (level, factor and replication) are the most dominant. By knowing which variable is dominant the investigator can save computer costs by decreasing the level (if applicable - without altering results drastically) of the dominant variable. For example, if the number of factors is unquestionably the dominant variable then the number of factors might be reduced to save computer runs without sacrificing the importance of the results.

The "Pareto Principle" allows the investigator to reduce the number of factors because this principle revolves around the thesis that most systems need roughly 20% of the factors to explain nearly 80% of the simulation results, and that the remaining 80% of the factors account for only 20% of the observed performance of the system. (Ref. 28:153) The difficult part of reducing the number of factors is to actually determine which factors can be placed in the 20% high performance category.

The two main factors in the AMIS simulation are the amount of CPU and I/O time, and the response variables of experimental interest are the gain factor and the CPU utilization. These response variables allow an analysis of the throughput performance to be conducted, but as was mentioned previously, much more needs to be done in determining accurate input data which means that the CPU and I/O factors will produce more reliable results.

V Recommendations

The results from this simulation effort are not extremely accurate and could be improved by the use of cluster analysis. The input data, which currently is made up of about 25% of the workload, could become more representative of all types of AMIS jobs if cluster analysis were to be employed in gathering data.

Cluster Analysis

Cluster analysis is a statistical tool for analyzing data by developing a data classification or identification. The end result of this classification or identification is a collection of data that has a high "natural association" among members of the same group, and a much lesser association between data that has been clustered into different groups. There are several clustering strategies but all methods involve two primary considerations. "First, there is defined a measure of group-density or of inter-group likeness. Examples of the latter type of measure (The so-called 'similarity coefficients').... Secondly, the chosen measure has to be incorporated into a "sorting" strategy" whereby groups of elements are extracted." (Ref. 3:373)

Anderberg (3) briefly discusses four metrics that might be used for clustering data which include the so-called "city block" metric, the Chebychev metric and the familiar Euclidean distance metric.

Euclidean Metric

Anderberg states that any metric must satisfy the following conditions:

- 1) $D(X,Y) = 0$ if and only if $X = y$
- 2) $D(X,Y) \geq 0$ for all X and Y in E ,
- 3) $D(X,Y) = D(Y,X)$ for all X and Y in E ,

4) $D(X,Y) \leq D(X,Z) + D(Y,Z)$ for all X, Y, and X in E.

where E is the symbolic representation for a measurement space and X, Y and Z are any three points in the space E.

The distance between data units can be found by using the Euclidean metric where

$$D_2(x_j, x_k) = \left[\sum_{i=1}^{i=n} (x_{ij} - x_{ik})^2 \right]^{1/2} \quad (3)$$

When dealing with two or more variables, the Euclidean metric can be used to determine clusters and Green (18) has suggested the following steps to compute the cluster analysis.

1. Each variable (characteristic) is transformed so that the data becomes a standard distribution where the following rule converts the variables to a standardized variate with zero mean and standard deviation.

$$z = \frac{x - u}{\sigma} \quad (4)$$

2. Distances between all possible pairs of data are calculated using the Euclidean Metric.

3. The pair of data points with the smallest distance between the two points is chosen as the initial node of the first cluster and the

centroid of this pair is calculated.

4. "Additional points are added to this cluster (based on "closeness" to the last-computed average) until:

- A. Some pre-specified number of points has been clustered, or
- B. The point to be added to the cluster exceeds some pre-specified distance cutoff number." (Ref. 18:391)

Anderberg suggests using the single-linkage method for accomplishing this step in the cluster analysis. As new data points are added to the cluster the new centroid must be calculated so point B is not violated when an additional point is merged or "linked" to the cluster. If it is known that a certain distance between the cluster centroid and the nearest neighbor can not exceed some "coarsening parameter" value, then a new cluster node is established by determining another centroid between the points with the smallest distance between them.

5. The new additional cluster node is used as a basis for building a new cluster and step 4 is repeated.

6. To avoid artificially fine distinctions between clusters, the points may be allowed to be in more than one cluster.

The above has been a general discussion of cluster analysis; the section that follows will suggest a practical use of this technique to refine the data generation employed to drive the simscript model.

Practical Use of Cluster Analysis

A computer program which performs a cluster analysis on the input data (i.e., CPU time, amount of core needed, and need for tape or disc mounts) should be the next step in the analysis of the system. Forgy (31) has provided a simple algorithm consisting of the following steps, which will help in such a cluster analysis:

- 1) Since there are 11 general priority class memberships, the end points, for both CPU time and core needed, within each class should become the initial cluster boundaries.
- 2) Using the two end points in conjunction with the mid-point of the boundaries as seed points, run the data to allocate each data unit to the cluster with the nearest seed point. The seed points must remain fixed during the complete run.
- 3) Compute new seedpoints from the centroids of the resulting clusters.
- 4) Re-run the data using the new centroids as seed points and repeat this iteration until no data units change their cluster membership.

It is not certain just how many runs will have to be made before this iterative process stabilizes, but Forgy suggests that from empirical evidence, this will ordinarily be accomplished within the first 5 runs.

When the cluster analysis has been completed there should be a total of 33 separate clusters, each having a centroid value for the CPU time and core needed. The frequency of each cluster (the number of data points within each cluster) should be calculated and the data is then ready to drive the simulation model.

Now that the data has been reduced to 33 clusters that accurately represent the workload characteristics of the system, the "centroid data"

can be called at random according to the frequency of each cluster. This means that 33 data points, called at random, will suffice for valid inputs. This is much superior to the internal generation of values, the method employed to this point, and allows the analyst to include data from all the different types of jobs run on the system.

VI Conclusions

- 1) Using the data about S2K major production jobs, the baseline simulation run generated a CPU utilization of 86.67%. This value is close to the real system CPU utilization of 85% but it must be remembered that the results in this investigation were not obtained from using data established by the cluster analysis technique in the previous section. In most of the charts in Appendix A the baseline figure for the CPU utilization was higher than the neighboring datum points (sensitivity runs). It appears that although the CPU utilization increased, during a few sensitivity runs, there was a consistent decline in the gain factor (degradation of the throughput or turnaround time). From the results obtained, it appears that a change in the workload will allow an increase in the CPU utilization but poorer turnaround results will also arise for the AMIS users.
- 2) The graph on page 64 supports the idea that the gain factor lowers in conjunction with an increase in CPU utilization. The gain factor for the baseline run was at a maximum, decreasing substantially as the interactive workload increased.
- 3) Never once were there 35 or more interactive terminals being used concurrently. From the results obtained and shown on page 65 a decrease in the number of interactive ports causes a 2% increase in CPU utilization, while at the same time causing the gain factor to fall by .3935.
- 4) If real time updates were to be performed interactively by AMIS users, the simulation results predict that throughput performance would probably decline. This premise is drawn from the data on page 64 where an increase in the interactive workload reflects a steady decline in the gain factor.

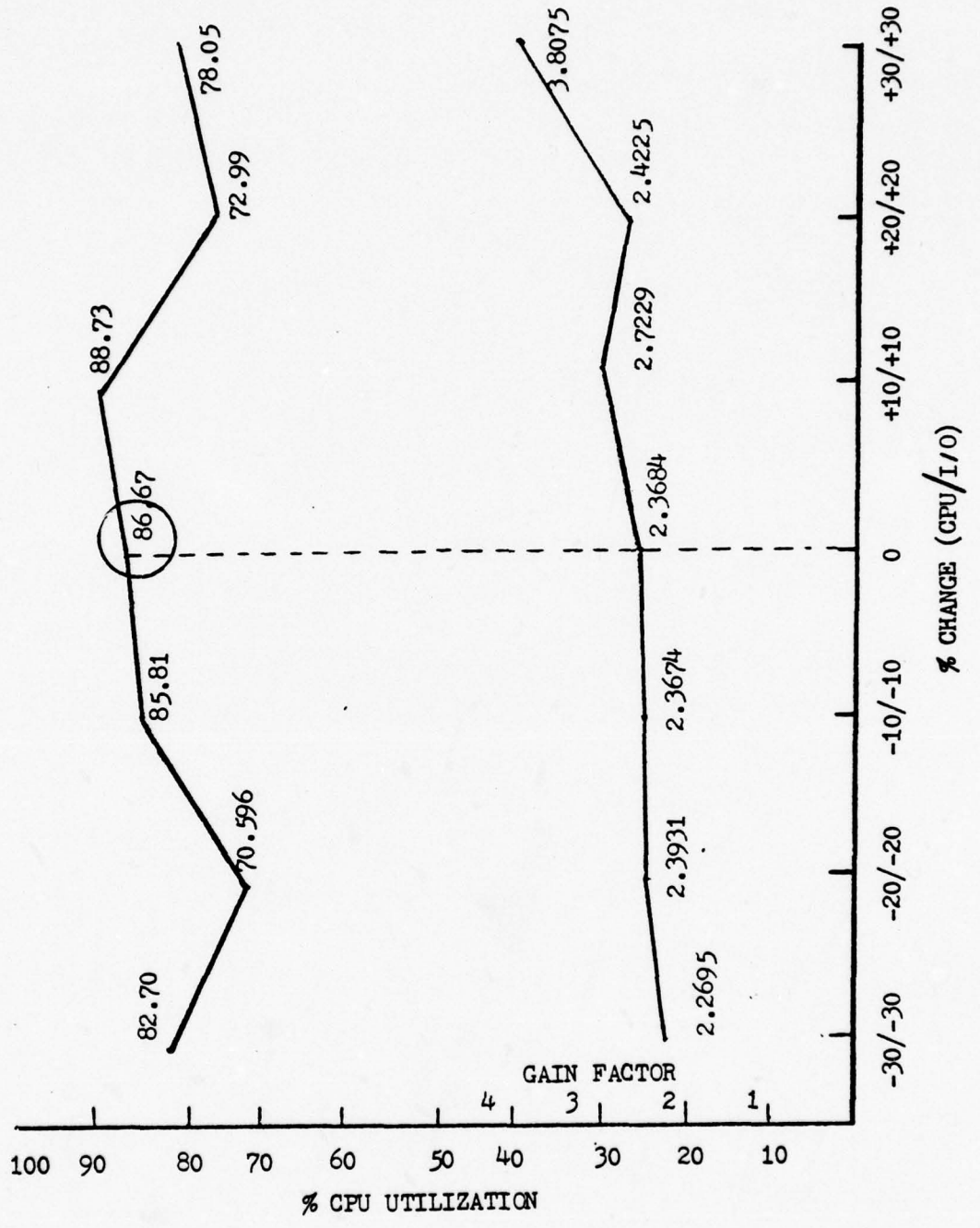
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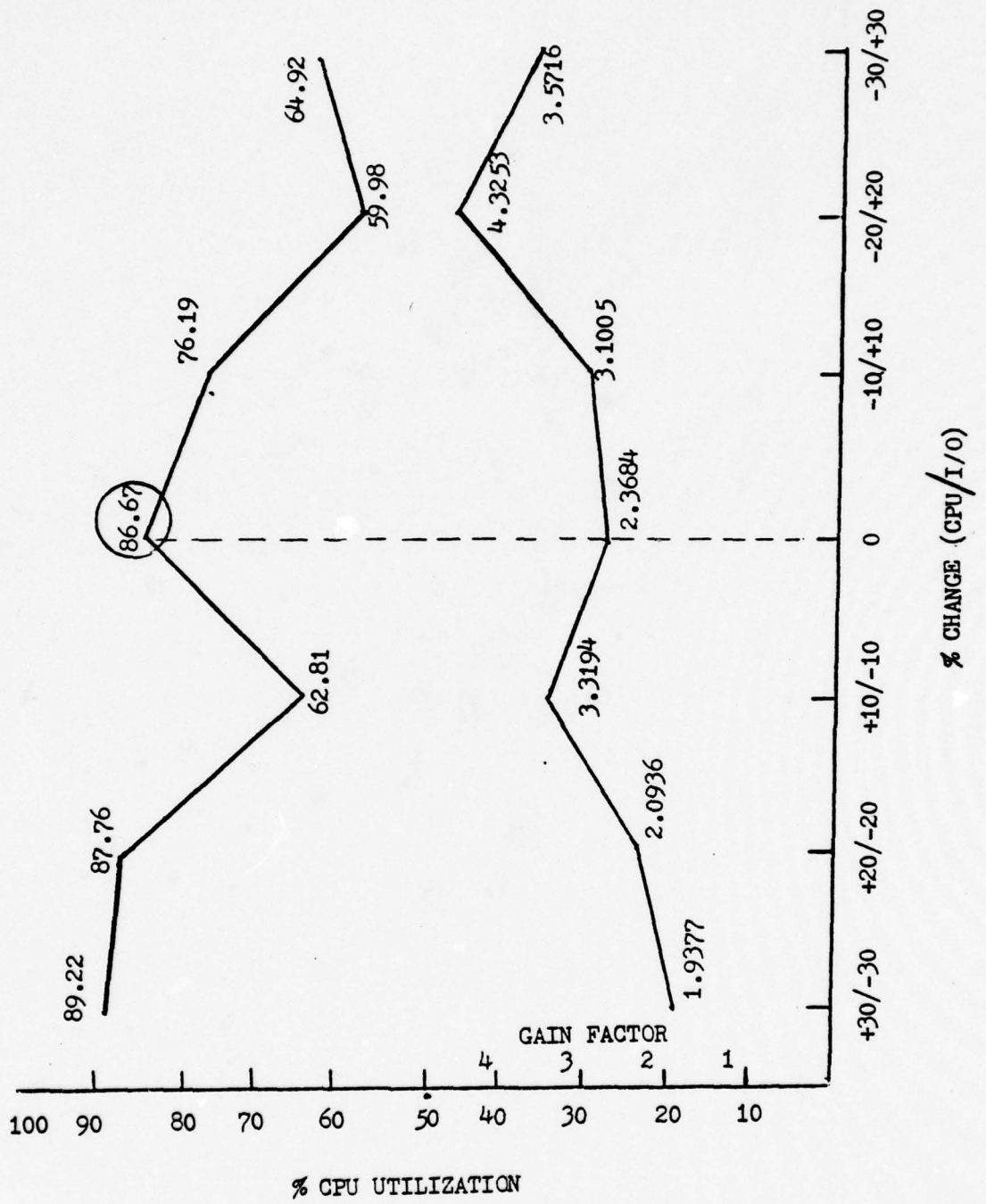
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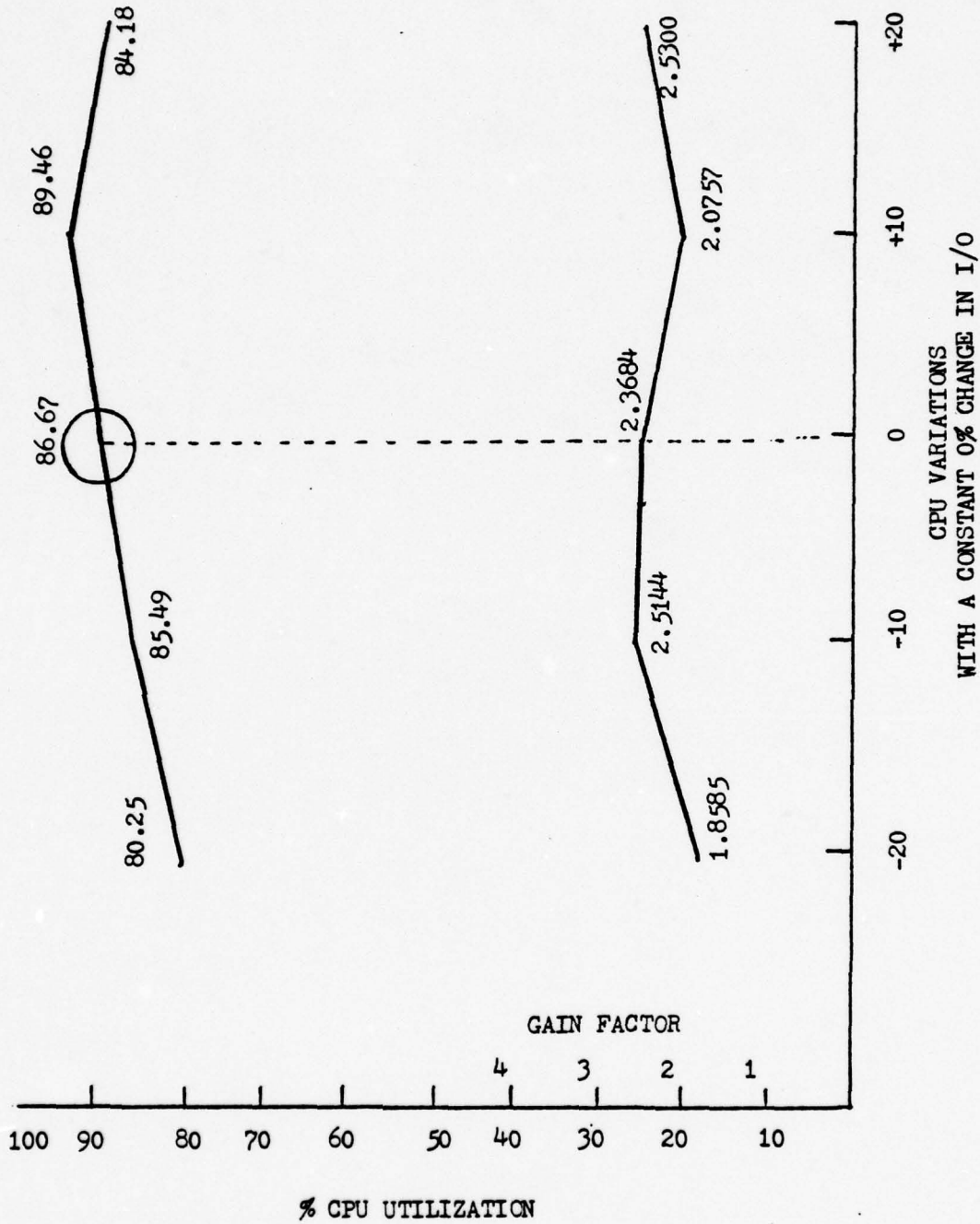
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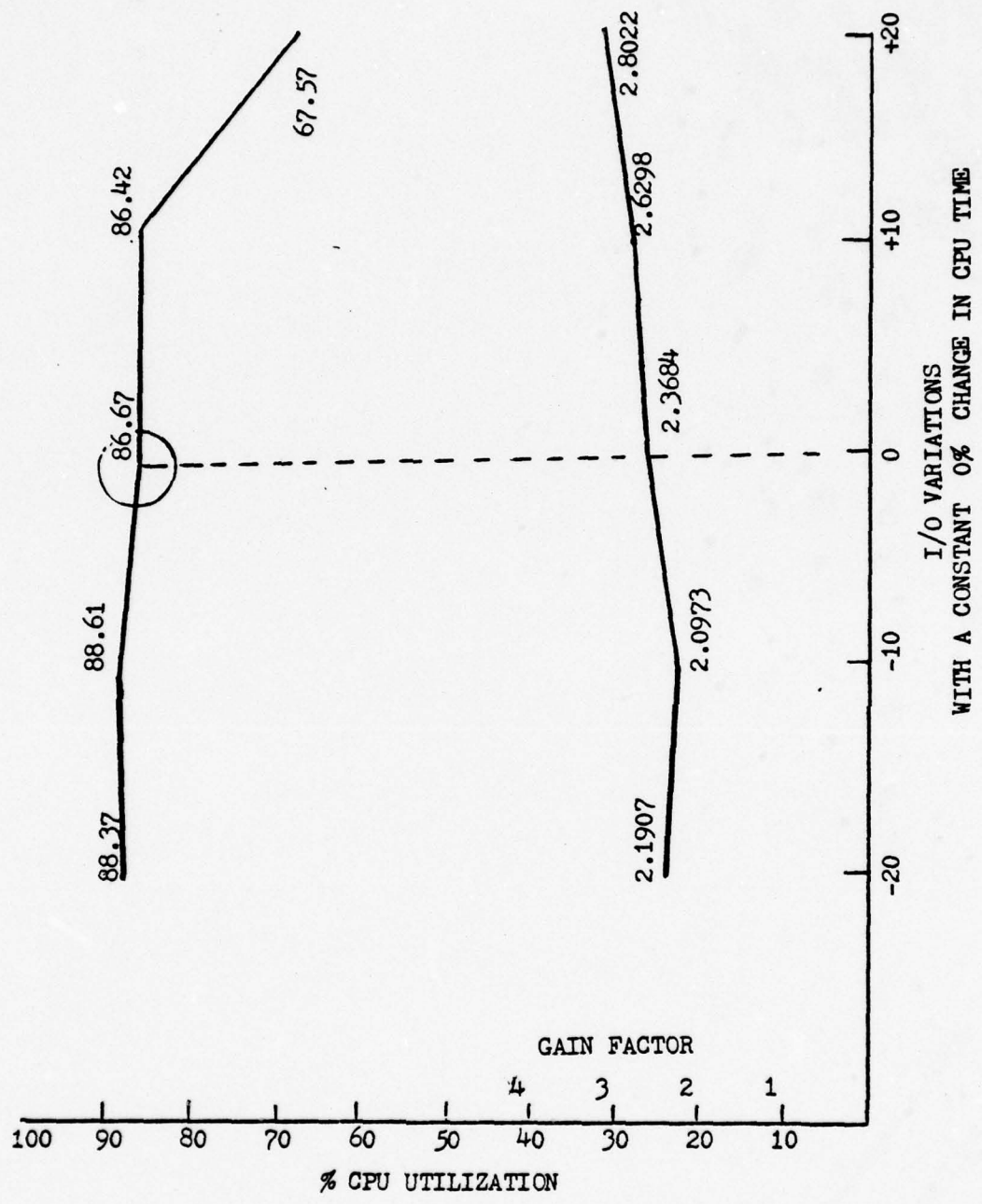
A P P E N D I X A

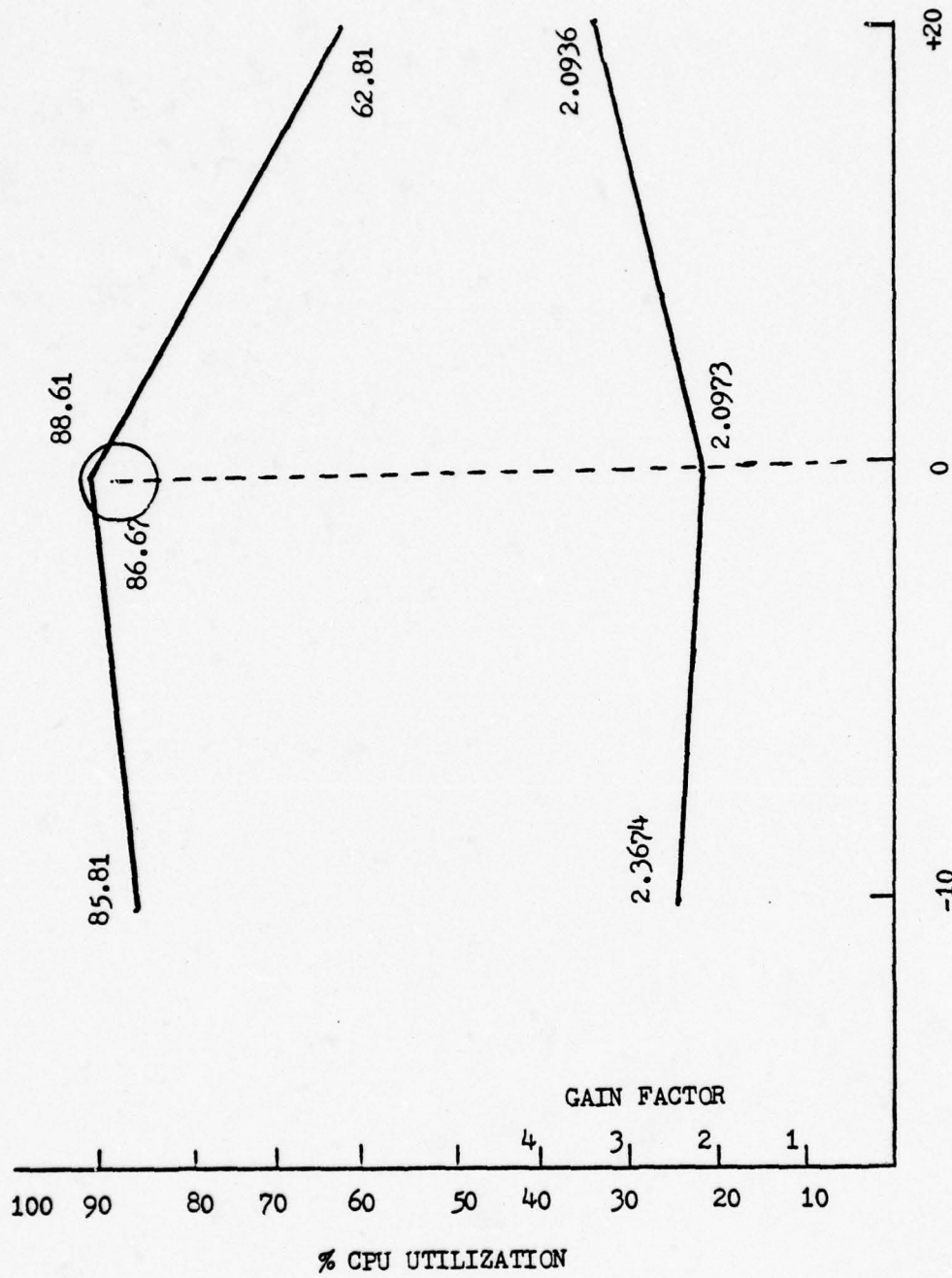
CPU Utilization vs Gain Factor



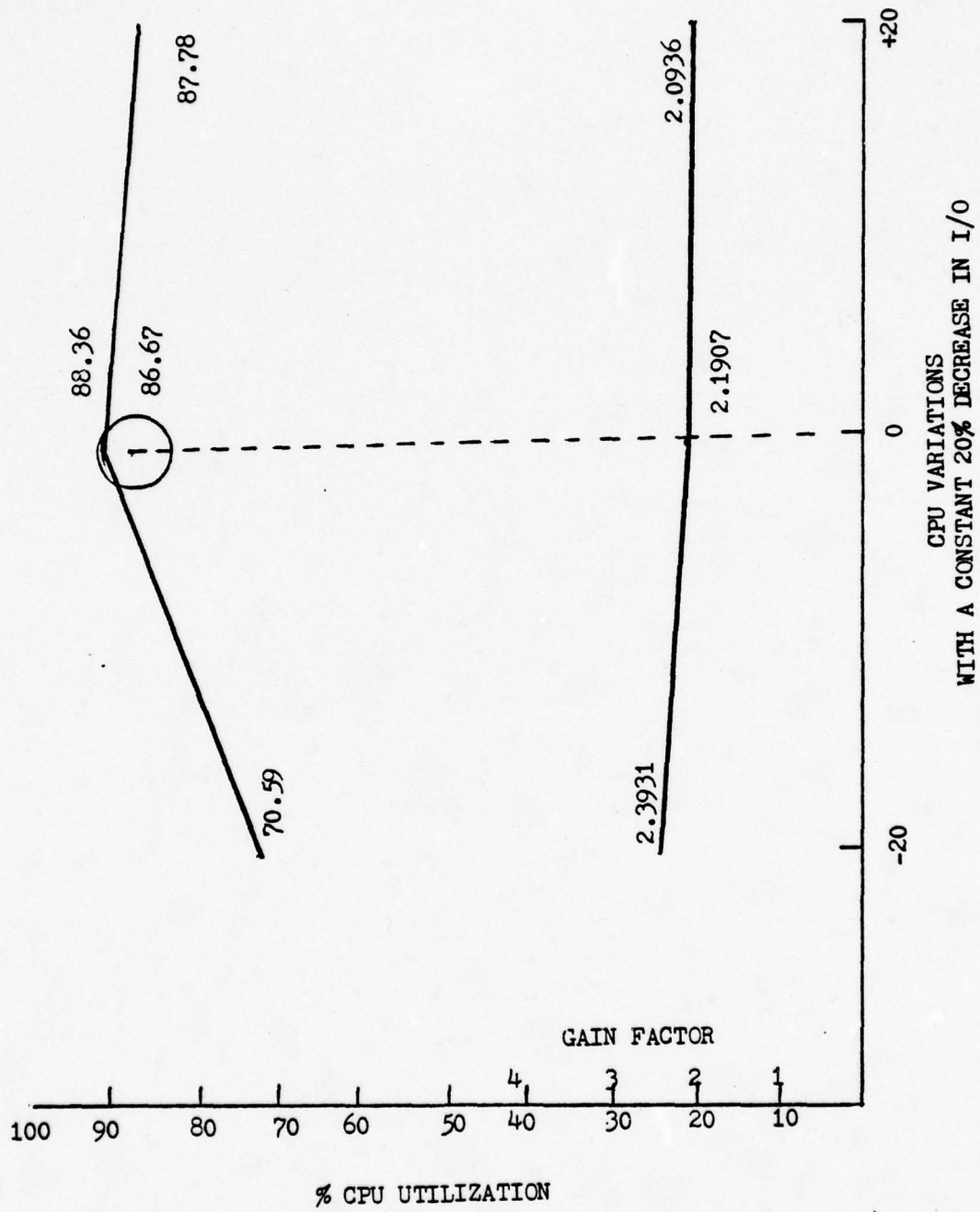


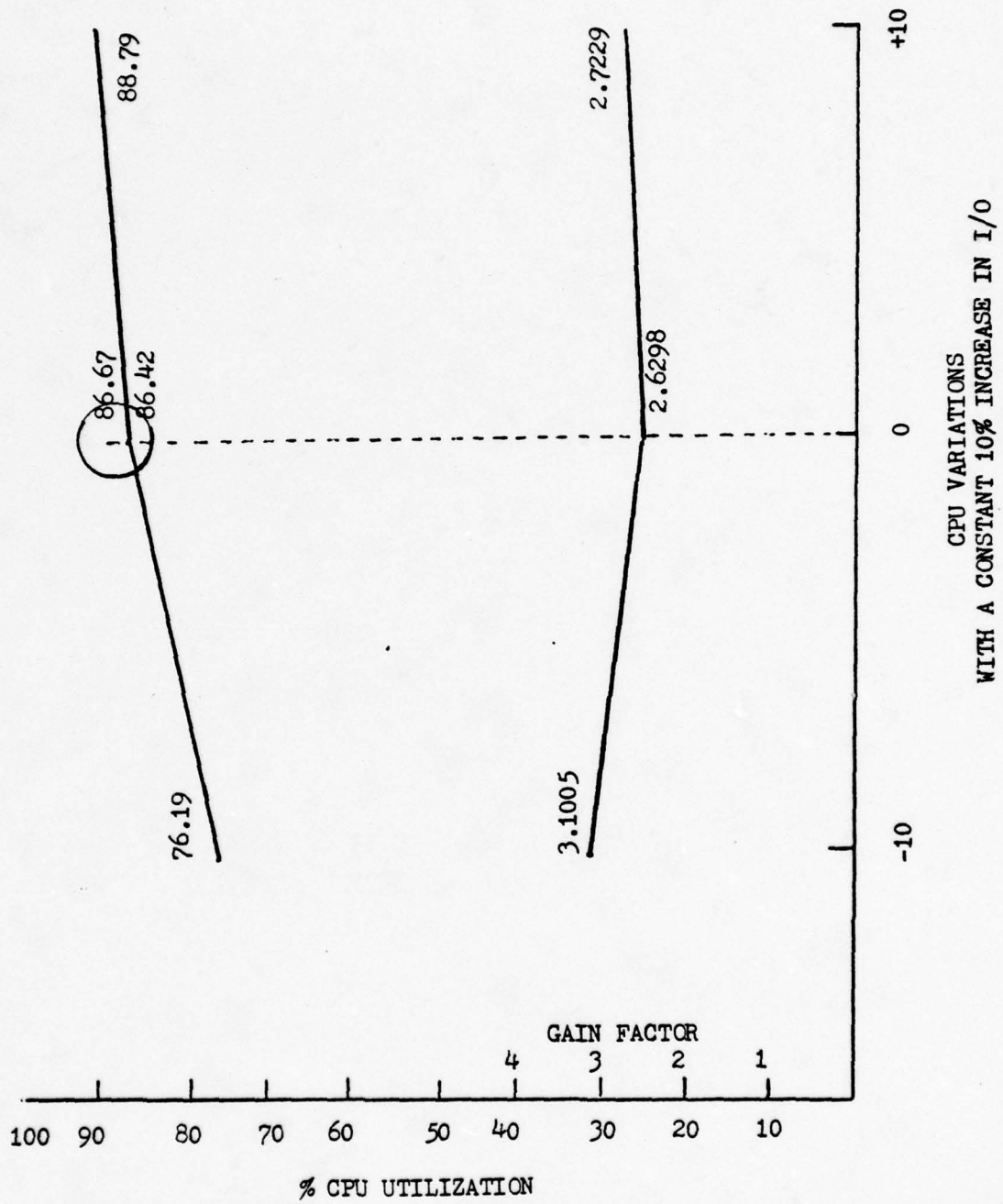


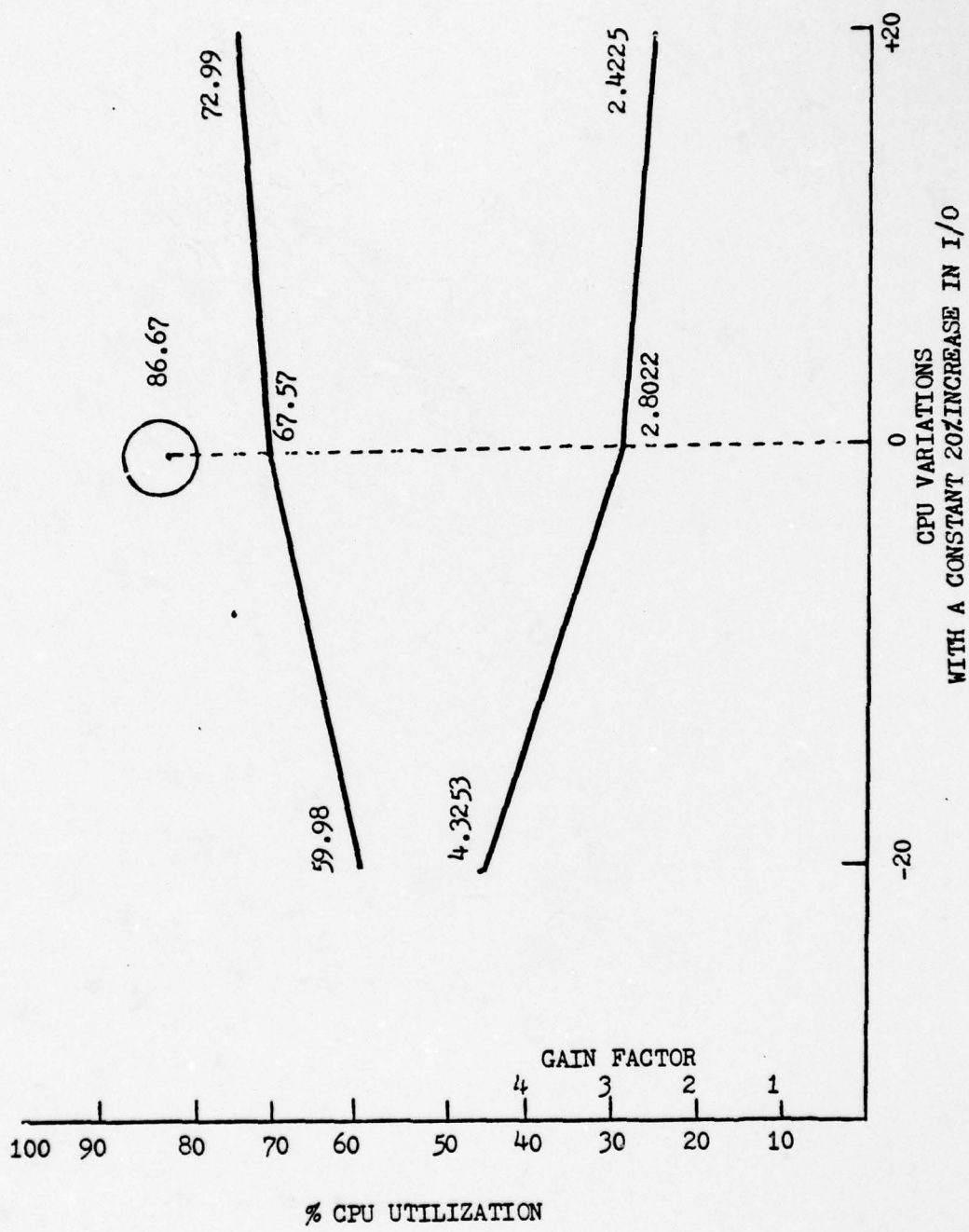


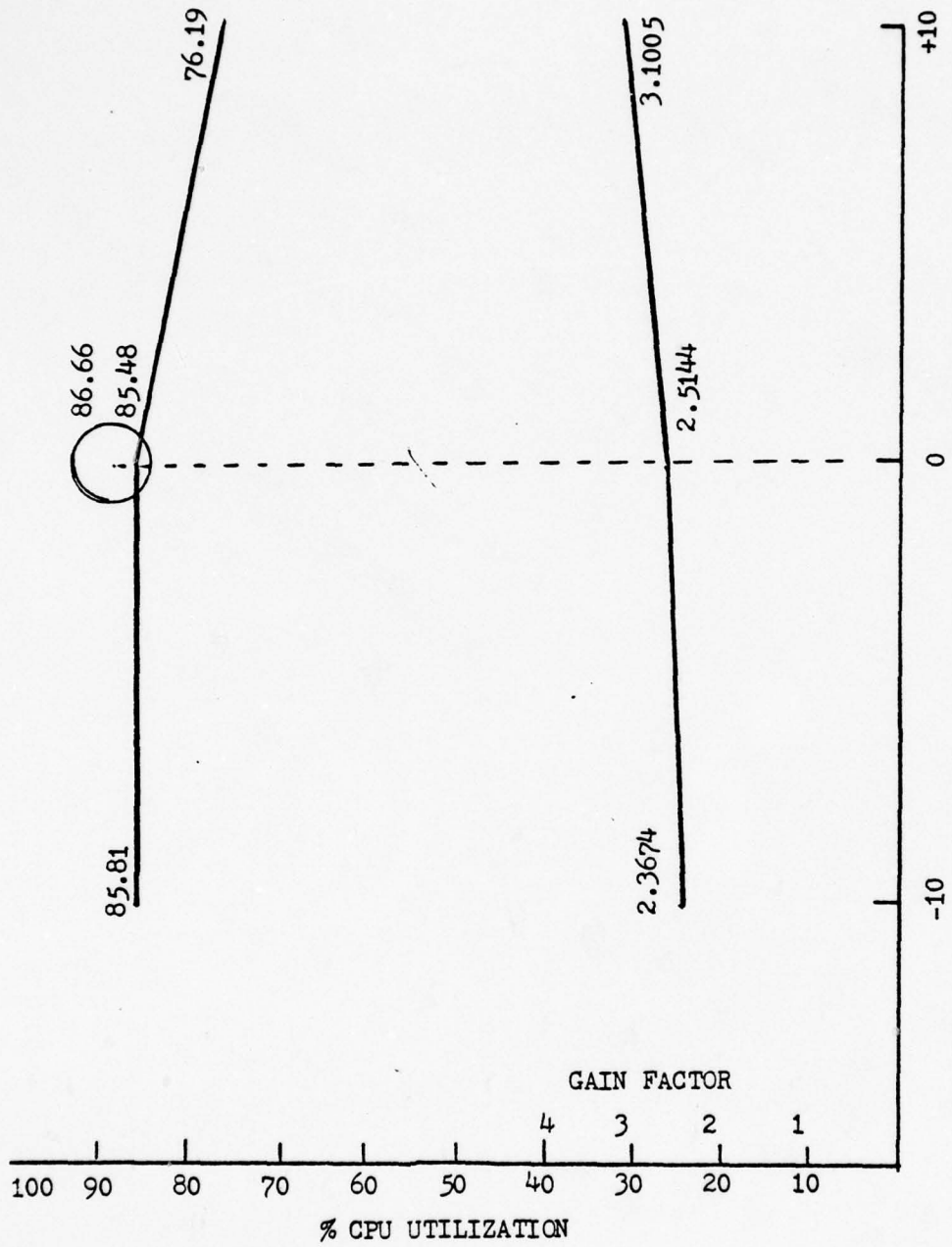


CPU VARIATIONS
WITH A CONSTANT 10% DECREASE IN I/O

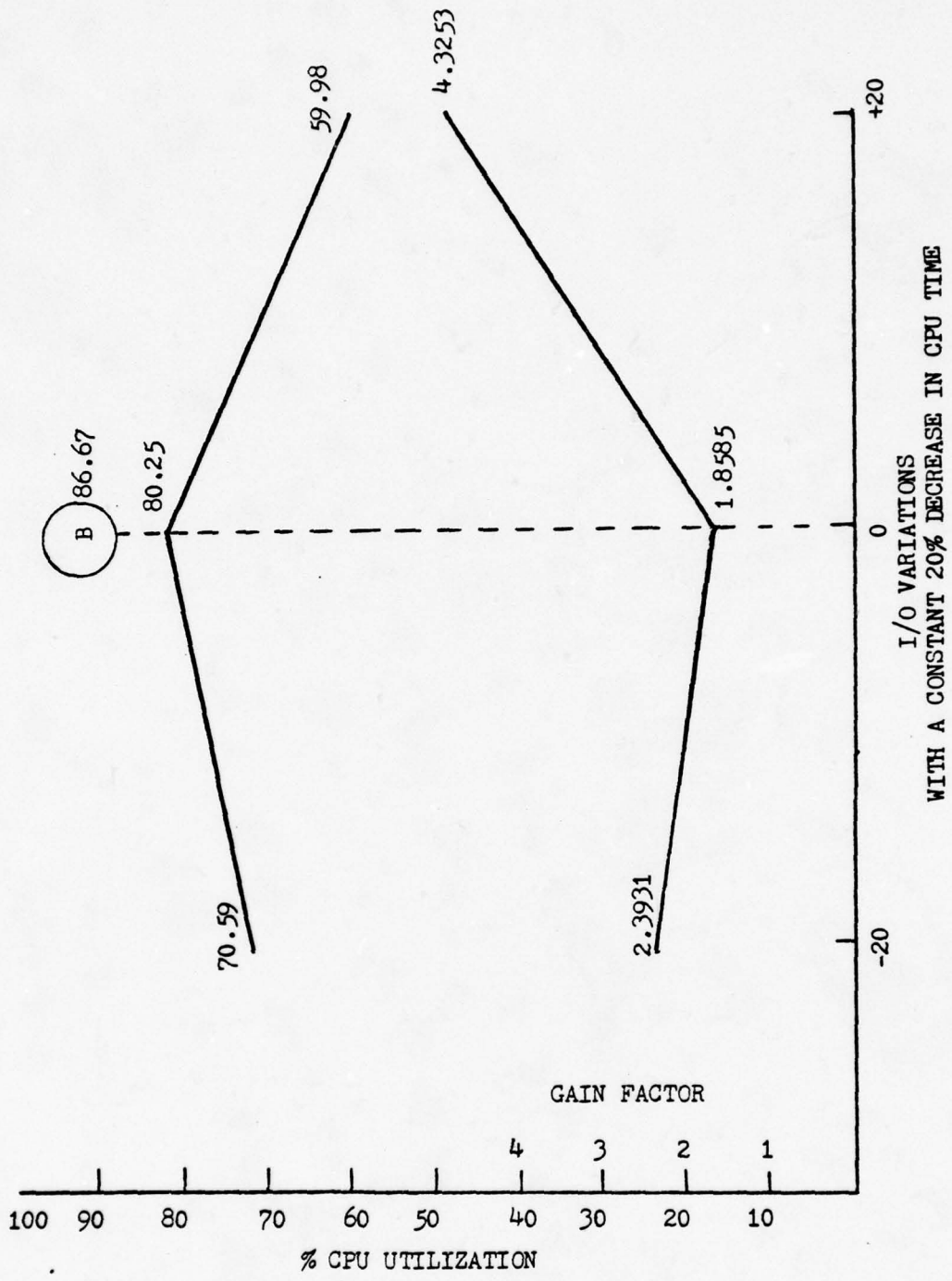


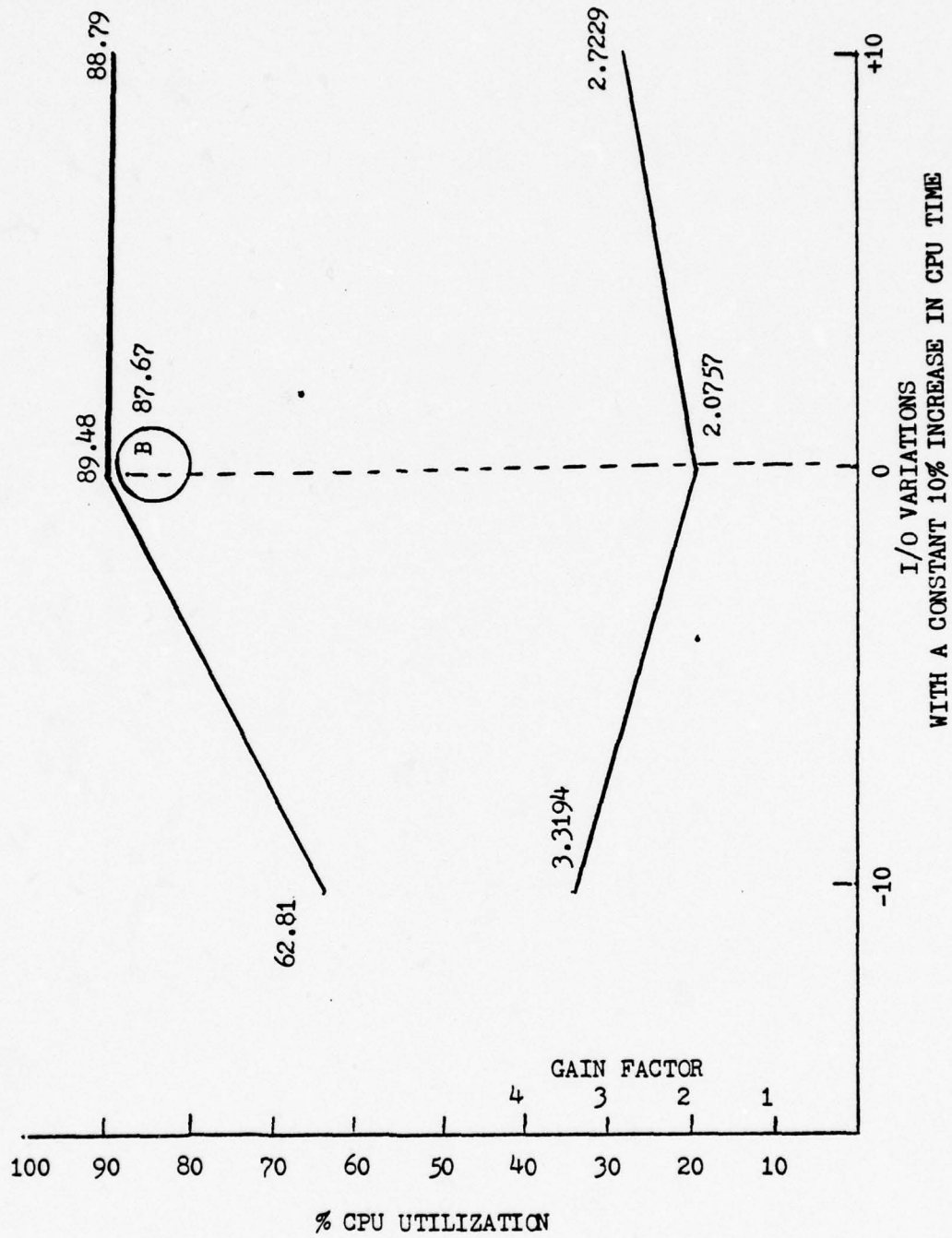


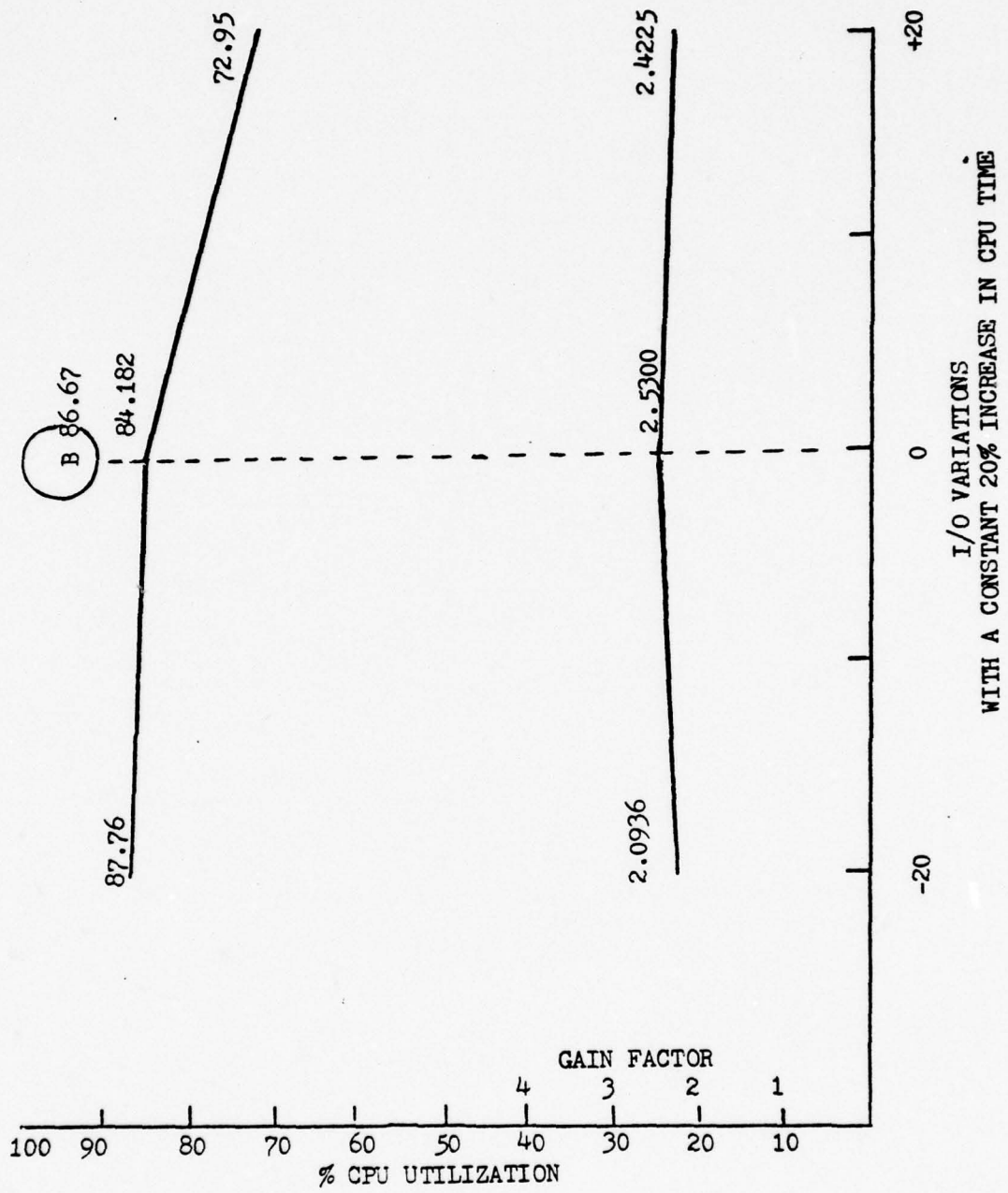


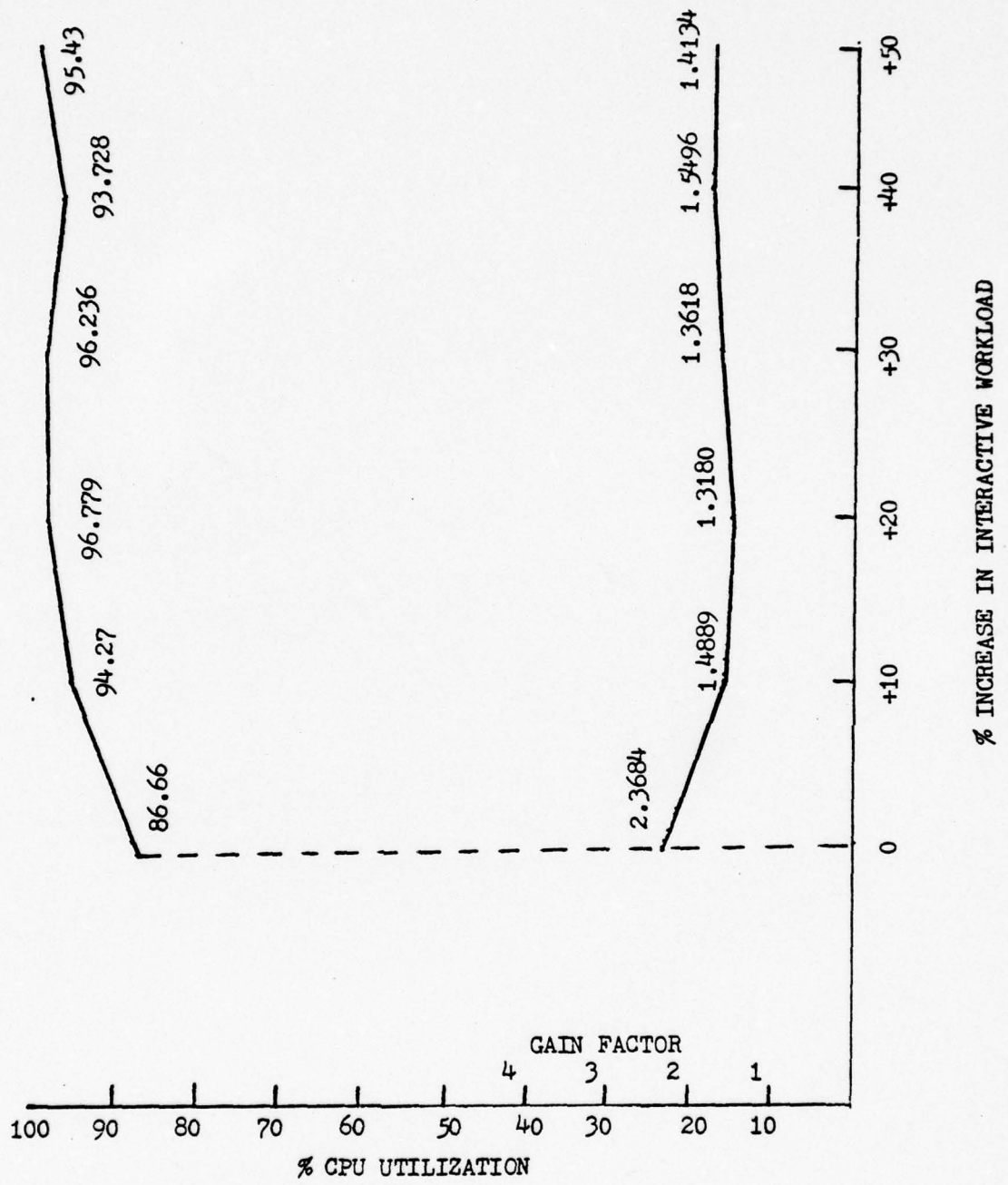


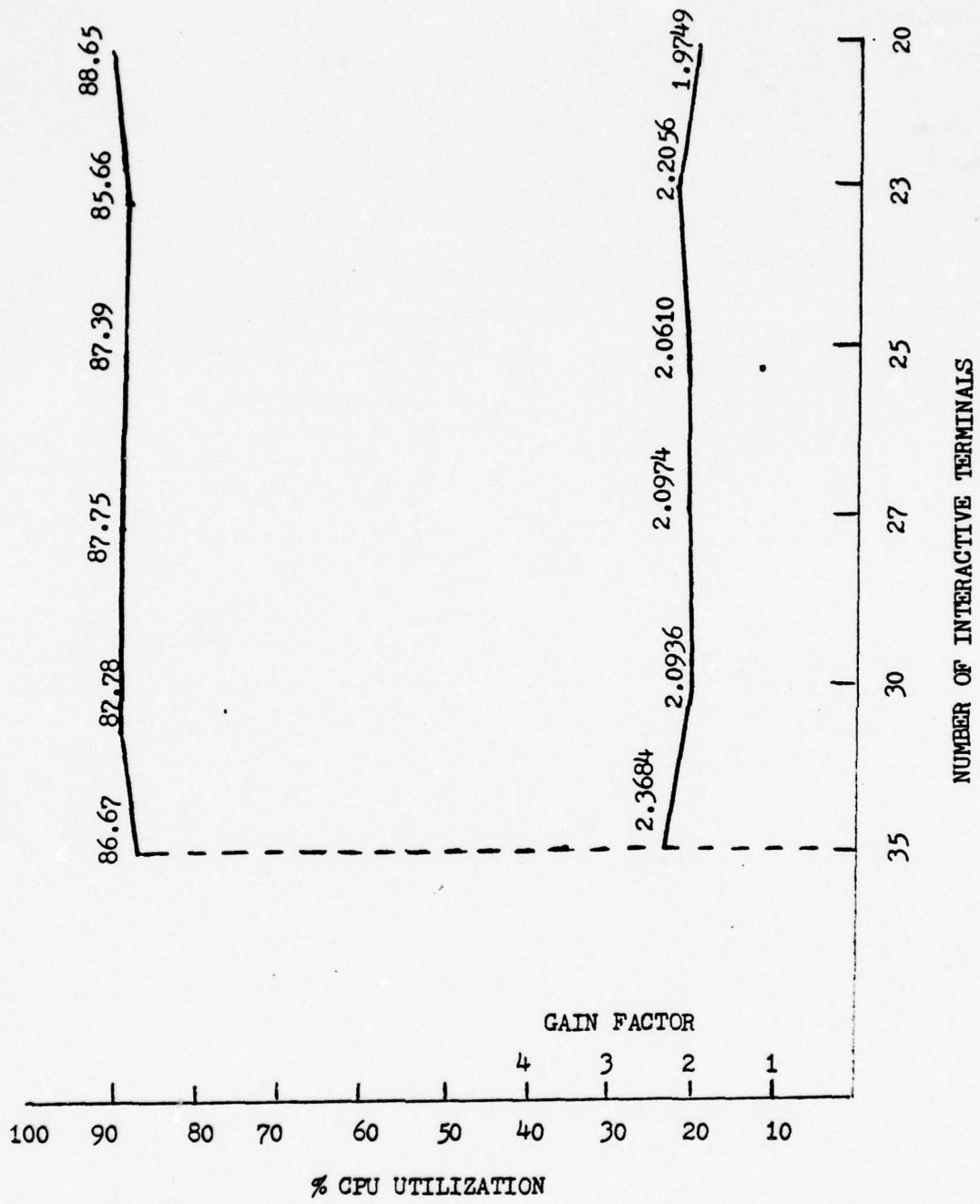
I/O VARIATIONS
WITH A CONSTANT 10% DECREASE IN CPU TIME











A P P E N D I X B
Simsript Simulation Program

This computer program in Appendix B took nearly 60 seconds to compile and each run took about 185 seconds for execution time. The ratio of simulation time to execution time (time to execute each run) is approximately 80 to 1. (It took 185 seconds to execute the simulation program for a simulated period of 4 hours). The model itself is comprised of 23 separate events (similar to subroutines) beginning with the preamble.

The different job attributes, events and parameters are established in the preamble while the actual values are established in the main routine. The variables of interest in each event are given a short discription within the event that they occur. Generally these discriptions are placed between two lines of "stars" which are physically placed near the beginning of each event or routine.

```

PREAMBLE
*****
**JT IS THE JOB TYPE VARIABLE
**
**CPU.SF IS THE CPU STATUS FLAG
** 0 = CPU IDLE 1 = CPU BUSY
**
**NBJ IS THE NUMBER OF BATCH JOBS IN THE SYSTEM
**
**NIJ IS THE NUMBER OF INTERACTIVE JOBS IN THE SYSTEM
**
**NJNPL0 IS THE NUMBER OF INTERACTIVE JOBS NOT PERMITTED TO LOGON
**
**SF.DISK IS THE DISK STATUS FLAG
** 0 = DISK IDLE 1 = BUSY
**
**SECOND IS EQUAL TO 1.0/60.0
**
**MAIN.MEMORY IS THE MAIN MEMORY CORE ARRAY
**
**JIS IS THE NUMBER OF JOB INITIATORS AVAILABLE FOR BATCH JOBS
**JISI IS THE NUMBER OF JOB INITIATORS AVAILABLE FOR INTERACTIVE JOBS
**
**JOB.SIZE INDICATES AMOUNT OF MEMORY NEEDED
**
**CP.T.QUANTUM IS THE TIME QUANTUM FOR BATCH OR INTERACTIVE JOB
**NEEDEN.RESOURCES.F IS A NUMERICAL INDICATION OF RESOURCES NEEDED
**P.OF.RESOURCES IS THE NUMERICAL REPRESENTATION OF RESOURCE WAIT
** TIME USING P&V SEMAPHORES
**GAIN = CP.TIME(AMIS) + IO.TIME(AMIS)
*****
NORMALLY MODE IS INTEGER
DEFINE N.SF AS INTEGER VARIABLE
DEFINE N AS INTEGER VARIABLE
DEFINE K AS INTEGER VARIABLE
DEFINE FLAG AS INTEGER VARIABLE

```

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000200
000210
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000230
000240
000250
000260
000270
000280
000290
000300
000310
000320
000330
000340
000350
000360
000370
000380
000390
000400
000410
000420
000430
000440
000450

```

```

DEFINE STATUS.PRIORITY AS INTEGER VARIABLE
DEFINE DISK.MOUNT AS INTEGER VARIABLE
DEFINE TAPE.MOUNT AS INTEGER VARIABLE
DEFINE JT AS INTEGER VARIABLE
DEFINE CPU.SF AS INTEGER VARIABLE
DEFINE NRJ AS INTEGER VARIABLE
DEFINE JISI AS INTEGER VARIABLE
DEFINE USER.RESPONSE.TIME AS INTEGER VARIABLE
DEFINE NSUI AS INTEGER VARIABLE
DEFINE NSUB AS INTEGER VARIABLE
DEFINE NIJ AS INTEGER VARIABLE
DEFINE NNPLO AS INTEGER VARIABLE
DEFINE SF.DISK AS INTEGER VARIABLE
DEFINE JIS AS INTEGER VARIABLE
DEFINE JJR.SIZE AS INTEGER VARIABLE
DEFINE WARM.UP.FLAG AS INTEGER VARIABLE
DEFINE FF AS INTEGER VARIABLE
DEFINE CHANNEL AS INTEGER VARIABLE
DEFINE LOAD.F AS INTEGER VARIABLE
DEFINE JN AS INTEGER VARIABLE
DEFINE RESOURCES.F AS INTEGER VARIABLE
DEFINE I.MEMORY.REQUEST AS INTEGER VARIABLE
DEFINE R.MEMORY.REQUEST AS INTEGER VARIABLE
DEFINE PFLAG AS INTEGER VARIABLE
DEFINE PJJFLAG AS INTEGER VARIABLE
DEFINE POUTFLAG AS INTEGER VARIABLE
DEFINE PFMFLAS AS INTEGER VARIABLE
DEFINE N.JOBS.PROCESSED AS INTEGER VARIABLE
DEFINE P.COUNTER AS INTEGER VARIABLE
DEFINE I.COUNTER AS INTEGER VARIABLE
DEFINE INUMBER AS INTEGER VARIABLE
DEFINE RNUMBER AS INTEGER VARIABLE

```

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000460
000470
000480
000490
000500
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000520
000530
000540
000550
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000570
000580
000590
000600
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000620
000630
000640
000650
000660
000670
000690
000690
000700
000710
000720
000730
000740
000750
000760
000770

```

```

DEFINE TOT.JOB.IO.TIME AS REAL VARIABLE
DEFINE TOTAL.JOB.CPU.TIME AS REAL VARIABLE
DEFINE BATCH.IO AS REAL VARIABLE
DEFINE TCRUT AS REAL VARIABLE
DEFINE CH AS REAL VARIABLE
DEFINE INTERACTIVE.IO AS REAL VARIABLE
DEFINE GAIN AS REAL VARIABLE
DEFINE GAIN.FACTOR AS REAL VARIABLE
DEFINE IO.OVERHEAD.T AS REAL VARIABLE
DEFINE TS AS REAL VARIABLE
DEFINE ETMF AS REAL VARIABLE
DEFINE I.ETMF AS REAL VARIABLE
DEFINE I.EX.DELAY.FACTOR AS REAL VARIABLE
DEFINE EX.DELAY.FACTOR AS REAL VARIABLE
DEFINE I.CAPABILITY AS REAL VARIABLE
DEFINE T AS REAL VARIABLE
DEFINE R.RT AS REAL VARIABLE
DEFINE I.RT AS REAL VARIABLE
DEFINE I.TURNAROUND.T AS REAL VARIABLE
DEFINE R.TURNAROUND.T AS REAL VARIABLE
DEFINE W.UP.TIME AS REAL VARIABLE
DEFINE I.CPU.TIME AS REAL VARIABLE
DEFINE B.CPU.TIME AS REAL VARIABLE
DEFINE R.CAPABILITY AS REAL VARIABLE
DEFINE CURRENT.UTILIZATION AS REAL VARIABLE
DEFINE T.CPU.TIME AS REAL VARIABLE
DEFINE ST AS REAL VARIABLE
DEFINE P AS REAL VARIABLE
DEFINE SEC.DAYS AS REAL VARIABLE
DEFINE MIL.SEC.DAYS AS REAL VARIABLE
DEFINE SECOND AS REAL VARIABLE
DEFINE MS AS REAL VARIABLE
DEFINE IODTT AS REAL VARIABLE
DEFINE CP.T.QUANTUM AS REAL VARIABLE
DEFINE MAIN.MEMORY AS INTEGER,1-DIMENSIONAL ARRAY

```

```

000840
000950
000950
000870
000890
000890
000900
000910
000920
000930
000940
000950
000950
000970
000980
000990
001000
001010
001020
001030
001040
001050
001060
001070
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001090
001100
001110
001120
001130
001140
001150
001160
001170
001180

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001190
 001200
 001210
 001220
 001230
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 001450
 001460
 001470
 001480
 001490
 001500
 001510
 001520
 001530
 001540
 001550
 001560

```

THE SYSTEM OWNS
  A TEMP.QUEUE,
  A LOAD.QUEUE,
  A WAIT.QUEUE,
  A THINK.STATE,
  A XX.BUFFER.QUEUE,
  A IO.QUEUE,
  A JOB.QUEUE,
  A CM.QUEUE,
  A CPU.REQUEST.QUEUE,
  A AN.OUTPUT.QUEUE,
  A ACM.RELEASE.QUEUE,
  A AH.OUTPUT.QUEUE,
  AND A ACPU

*****
**XX.TOTAL.CPU.TIME IS THE TOTAL TIME REQUESTED BY A SINGLE JOB
**
**ACC.CPU.TIME IS THE TIME THE JOB ACCESSED THE CPU
**
**JOB.IO.REQUESTS IS THE TOTAL NUMBER OF I/O OPERATIONS REQUESTED
**      BY A SINGLE JOB
**
**TLF IS THE TIME LEFT FLAG
**  =0 MEANS THERE IS TIME LEFT OVER FOR CURPENT CPU TASK
**  =1 MEANS THE CPU TASK HAS BEEN FINISHED
**
**XX.CPU.SERVICE TIME IS THE CPU TIME REQUIRED TO PROCESS UNTIL
**      I/O TASK
**
**IO.SERVICE TIME IS THE I/O TIME REQUIRED TO PPROCESS A SINGLE IO TASK
**
**T.LEFT.IN.CPU IS THE AMOUNT OF TIME REMAINING UNTIL NEXT I/O
**
**MEMORY.REQUEST IS THE AMONT OF MEMORY REQUESTED BY A SINGLE JOB
*****

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001740
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001800
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001870
001880
001890
001900
001910
001920

TEMPORARY ENTITIES

EVERY JOB HAS
A TURNAROUND,
A CP.TIME,
A IO.TIME,
A RE.SOURCES,
A MEMORY.REQUEST,
A T.E.CPU,
A START.MEMORY,
A END.MEMORY,
A JOBNUM,
A JOB.TYPE,
A TIME.E.CPJ,
A ARRIVAL.TIME,
A JOB.EXIT.TIME,
A PRIORITY,
A T.LEFT.IN.CPU,
A NEEDED.RESOURCES.F,
A P.OF.RESOURCES,
A XTIME.E.CPU.QUEUE,
A NO.CAR.RETURNS,
A F.LOADER,
A LOADING.T,
A PAGE.LENGTH,
A XX.CPU.SERVICE.TIME,
A XX.CPU.SERVICE.TIME,
A JOB.IO.REQUESTS,
A JOB..IO.REQUESTS,
A JIO,
A BLOCKED.TIME,
A TLF,
A IO.SERVICE.TIME,

001930
 001940
 001950
 001960
 001970
 001980
 001990
 002000
 002010
 002020
 002030
 002040
 002050
 002060
 002070
 002080
 002090
 002100
 002110
 002120
 002130
 002140
 002150
 002160
 002170
 002180
 002190
 002200
 002210
 002220
 002230
 002240
 002250
 002260
 002270
 002280

AND BELONGS TO
 A LOAD.QUEUE,
 A THINK.STATE,
 A WAIT.QUEUE,
 A JOB.QUEUE,
 A CM.QUEUE,
 A IO.QUEUE,
 A CPU.REQUEST.QUEUE,
 A TEMP.QUEUE,
 A ACM.RELEASE.QUEUE,
 A AN.OUTPUT.QUEUE,
 A XX.BUFFER.QUEUE,
 A HH.OUTPUT.QUEUE,
 AND A ACPU

NORMALLY MODE IS REAL
 DEFINE TURNAROUND AS REAL VARIABLE
 DEFINE T.E.CPU AS REAL VARIABLE
 DEFINE JOB.EXIT.TIME AS REAL VARIABLE
 DEFINE IO.TIME AS REAL VARIABLE
 DEFINE CP.TIME AS REAL VARIABLE
 DEFINE ARRIVAL.TIME AS REAL VARIABLE
 DEFINE BLOCKED.TIME AS REAL VARIABLE
 DEFINE T.LEFT.IN.CPU AS REAL VARIABLE
 DEFINE XX.CPU.SERVICE.TIME AS REAL VARIABLE
 DEFINE IO.SERVICE.TIME AS REAL VARIABLE
 DEFINE XX.CPU.SERVICE.TIME AS REAL VARIABLE
 DEFINE XXTIME.E.CPU.QUEUE AS REAL VARIABLE
 DEFINE LOADING.T AS REAL VARIABLE
 DEFINE CPU.R AS REAL VARIABLE

NORMALLY MODE IS INTEGER
 DEFINE HH.OUTPUT.QUEUE AS FIFO SET
 DEFINE LOAD.QUEUE AS FIFO SET
 DEFINE XX.BUFFER.QUEUE AS FIFO SET

002290
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002590
002600
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002640

DEFINE TMRP.QUEUE AS FIFO SET
DEFINE THINK.STATE AS FIFO SET
DEFINE CPU..REQUEST.QUEUE AS A FIFO SET RANKED BY
HIGH PRIORITY
DEFINE CM.QUEUE AS A FIFO SET RANKED BY HIGH PRIORITY
DEFINE JOB.QUEUE AS A FIFO SET RANKED BY HIGH PRIORITY
DEFINE WAIT.QUEUE AS A FIFO SET
DEFINE IO.QUEUE AS FIFO SET
DEFINE ACPU AS A FIFO SET
DEFINE ACM.RELEASE.QUEUE AS A FIFO SET
DEFINE AN.OUTPUT.QUEUE AS A FIFO SET RANKED BY HIGH PRIORITY

EVENT NOTICES INCLUDE

DIAGNOSTIC,
LOAD.DELAY,
R.JOB.REQUEST,
I.JOB.REQUEST,
REQUEST.CM,
XX.TRAFFIC.CONTROLLER,
RELEASE.CPU,
CM.RELEASE,
CPU.PROCESSING,
AN.DISK.RELEASE,
IS.BUFFER.EMPTY,
RESPONSE,
CPU.UTILIZATION,
IO.DELAY,
CPU.REQUEST,
OUTPUT,
XX.OUTPUT,
P.INCREASE,
PJO.INCREASE,
P.OUIP.INCREASE,
PCY.INCREASE,
CPU.ANALYSIS,
AND CEASE.SIM

EVERY RESPONSE HAS A RNAME
 EVERY IO.DELAY HAS A NAME
 TALLY IIO AS THE MEAN OF INTEFACTIVE.IO
 TALLY RIO AS THE MEAN OF RATCH.IO
 TALLY ANSUI AS THE MEAN OF NSUI
 TALLY ENSU3 AS THE MAX OF NSUP
 TALLY FNSUR AS THE MIN OF NSUP
 TALLY BNSUI AS THE MAX OF NSUI
 TALLY CNSUI AS THE MIN OF NSUI
 TALLY ONSUR AS THE MEAN OF NSUI
 TALLY AI.TURNAROUND.T AS MEAN OF I.TURNAROUND.T
 TALLY RI.TURNAROUND.T AS VARIANCE OF I.TURNAROUND.T
 TALLY CI.TURNAROUND.T AS STD.DEV OF I.TURNAROUND.T
 TALLY AR.TURNAROUND.T AS MEAN OF R.TURNAROUND.T
 TALLY RR.TURNAROUND.T AS VARIANCE OF R.TURNAROUND.T
 TALLY CR.TURNAROUND.T AS STD.DEV OF R.TURNAROUND.T
 TALLY AI.CPU.TIME AS MEAN OF I.CPU.TIME
 TALLY RI.CPU.TIME AS VARIANCE OF I.CPU.TIME
 TALLY CI.CPU.TIME AS STD.DEV OF I.CPU.TIME
 TALLY AR.CPU.TIME AS MEAN OF R.CPU.TIME
 TALLY RR.CPU.TIME AS VARIANCE OF R.CPU.TIME
 TALLY CR.CPU.TIME AS STD.DEV OF R.CPU.TIME
 TALLY AA AS MEAN OF I.CAPABILITY
 TALLY RR AS MEAN.SQUARE OF I.CAPABILITY
 TALLY CC AS VARIANCE OF I.CAPABILITY
 TALLY DD AS STD.DEV OF I.CAPABILITY
 TALLY EE AS MEAN OF R.CAPABILITY
 TALLY FFFF AS MEAN.SQUARE OF F.CAPABILITY
 TALLY GG AS VARIANCE OF B.CAPABILITY
 TALLY HH AS STD.DEV OF B.CAPABILITY
 TALLY AAA AS THE MEAN OF I.BT
 TALLY RBT AS THE MEAN.SQUARE OF I.BT
 TALLY CCC AS THE VARIANCE OF I.BT
 TALLY ODD AS THE STD.DEV OF I.BT
 TALLY FEE AS THE MEAN OF R.BT
 TALLY FFF AS THE MEAN.SQUARE OF B.BT

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 002800
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TALLY GGG AS THE VARIANCE OF R.BT
TALLY HHH AS THE STD.DEV OF B.RT
TALLY II AS THE MEAN OF I.MEMORY.REQUEST
TALLY JJ AS THE MEAN.SQUARE OF I.MEMORY.REQUEST
TALLY KK AS THE VARIANCE OF I.MEMORY.REQUEST
TALLY LL AS THE STD.DEV OF I.MEMORY.REQUEST
TALLY MM AS THE MEAN OF B.MEMORY.REQUEST
TALLY NN AS THE MEAN.SQUARE OF B.MEMORY.REQUEST
TALLY OO AS THE VARIANCE OF B.MEMORY.REQUEST
TALLY PP AS THE STD.DEV OF B.MEMORY.REQUEST
TALLY QQ AS THE MEAN OF USER.RESPONSE.TIME
TALLY RR AS THE MEAN.SQUARE OF USER.RESPONSE.TIME
TALLY SS AS THE VARIANCE OF USER.RESPONSE.TIME
TALLY TT AS THE STD.DEV OF USER.RESPONSE.TIME
TALLY UU AS THE MEAN OF ETMF
TALLY VV AS THE VARIANCE OF ETMF
TALLY WW AS THE STD.DEV OF ETMF
TALLY XX AS THE MEAN.SQUARE OF ETMF
TALLY UUU AS THE MEAN OF EX.DELAY.FACTOR
TALLY VVV AS THE VARIANCE OF EX.DELAY.FACTOR
TALLY WWW AS THE STD.DEV OF EX.DELAY.FACTOR
TALLY XXX AS THE MEAN.SQUARE OF EX.DELAY.FACTOR
TALLY I.E AS THE MEAN OF I.ETMF
TALLY I.FA AS THE VARIANCE OF I.ETMF
TALLY I.EB AS THE STD.DEV OF I.ETMF
TALLY I.EC AS THE MEAN.SQUARE OF I.ETMF
TALLY I.EDFA AS THE MEAN OF I.EX.DELAY.FACTOR
TALLY I.EDFB AS THE VARIANCE OF I.EX.DELAY.FACTOR
TALLY I.EDFC AS THE STD.DEV OF I.EX.DELAY.FACTOR
TALLY I.EDFD AS THE MEAN.SQUARE OF I.EX.DELAY.FACTOR
DEFINE TAPE.MOUNT TO MEAN TM
DEFINE DISK.MOUNT TO MEAN DM
END

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003760
003770

```

```

MAIN ** INITIALIZE AND SCHEDULE FIRST JOBS
PRINT 1 LINE THUS
ENTERED MAIN
PRINT 1 LINE WITH TIME.V THUS
SIMULATION TIME IS *****
*****
**NBJ IS THE NUMBER OF PATCH JOBS IN THE SYSTEM
**NIJ IS THE NUMBER OF INTERACTIVE JOBS IN THE SYSTEM
**CPU.SF IS THE CPU STATUS FLAG 0=IDLE,1=RUSY
**ST IS SCHEDULED JOB TIME
**
** *****NJPLO IS THE NUMBER OF JOBS NOT PERMITTED TO LOGIN
** *****SF.DISK IS THE DISK STATUS FLAG 0 = IOLE, 1 = RUSY
** *****JIS IS THE NUMBER OF JB INITIATORS AVAILABLE FOR BATCH JOBS
** *****JISI IS TH NUMBER OF JOB INITIATORS AVAIL FOR INTERACTIVE JOB
*****
DEFINE I AS INTEGER VARIABLE
LET SECOND = 1.0/60.0
LET MS = SECOND/1000.0
LET SEC.DAYS=(1.0/60.0)*(1.0/60.0)*(1.0/24.0)
LET MIL.SEC.DAYS=(1.0/1000.0)*SEC.DAYS
LET FF=0
LET FF=1
LET NJJ = 0
LET NIJ = 0
LET NJPLO = 0
LET JIS = 6
LET JISI = 35
LET JN = 0
LET CPU.SF=0
LET SF.NTSK = 0
LET STATUS.PRIORITY = 0
LET PJOFLAG=0
LET POUTFLAG=0
LET PGMFLAG=0
LET PFLAG = 0

```

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003940
003950
003960
003970
003980
003990
004000
004010
004020
004030
004040
004050
004060
004070
004080
004090
004100
004110
004120
004130

```
LET GAIN = 0.0
LET GAIN.FACTOR = 0.0
LET ETMF = 0.0
LET N.JOBS.PROCESSED = 0
LET LOAD.F=0
LET FLAG=1
LET FLAG = 0
LET WARM.UP.FLAG=0
LET W.UP.TIME=0.0
LET T.CPU.TIME=0.0
LET T.OUTPUT=0.0
LET C1=0.0
LET CURRENT.UTILIZATION=0.0
RESERVE MAIN.MEMORY AS 4000
FOR B=1 TO 5, DO
LET ST=REAL.(RANDI.F(5,60,1))*SEC.DAYS
SCHEDULE A B.JOB.REQUEST IN ST DAYS
LOOP
LET BNUMBER =5
LET B.COUNTER =0
LET INUMBER=0
LET I.COUNTER=0
SCHEDULE A B.JOB.REQUEST NOW
SCHEDULE A I.JOB.REQUEST NOW
SCHEDULE A CPU.UTILIZATION IN 10.0*SECOND MINUTES
SCHEDULE A P.INCREASE IN SECOND MINUTES
SCHEDULE A P.C1.INCREASE IN SECOND MINUTES
SCHEDULE A P.OUTPUT.INCREASE IN SECOND MINUTES
SCHEDULE A P.J1.INCREASE IN SECOND MINUTES
SCHEDULE A CPU.ANALYSIS IN 1.0*SECOND MINUTES
SCHEDULE A DIAGNOSTIC IN .04 DAYS
SCHEDULE A CEASF.SIM IN .162 DAYS
PRINT 1 LINE THUS
EXIT MAIN
START SIMULATION
END
```

004180
004190
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004720
004730

```
EVENT DIAGNOSTIC
LET FLAG = 1
LET FLAG = 0
RETURN
END
EVENT CPU.ANALYSIS SAVING THE EVENT NOTICE
PRINT 1 LINE THUS
      SECOND STATS
PRINT 1 LINE WITH JN AND TIME.V THUS
      * .*****
IF WARM.UP.FLAG GE 1
LET CURRENT.UTILIZATION=(T.CPU.TIME/(TIME.V-W.UP.TIME))*100.0
PRINT 1 LINE WITH CURRENT.UTILIZATION THUS
      ***.**
ELSE
LFT CU=(TCPUT/TIME.V)*100.0
PRINT 1 LINE WITH CU THUS
      ***.**
ALWAYS
SCHEDULE THE CPU.ANALYSIS IN 1.0*SECOND MINUTES
IF TIME.V GT 30.0*SEC.DAYS
CANCEL CPU.ANALYSIS
RETURN
ELSE
RETURN
END
EVENT CPU.UTILIZATION SAVING THE EVENT NOTICE
PRINT 1 LINE THUS
      CURRENT STATS
PRINT 1 LINE WITH JN AND TIME.V THUS
      * .*****
IF FF=0
PRINT 1 LINE WITH T.CPU.TIME THUS
      * .*****
ELSE
ALWAYS
```

```

IF WARM.UP.FLAG GE 1
LET CURRENT.JUTILIZATION=(T.CPU.TIME/(TIME.V-W.UP.TIME))*100.0
PRINT 1 LINE WITH CURRENT.JUTILIZATION THUS
***.**
ELSE
LET CU=(ICPUT/TIME.V)*100.0
PRINT 1 LINE WITH CU THUS
***.**
ALWAYS
SCHEDULE THE CPU.JUTILIZATION IN 10.0*SECOND MINUTES
IF TIME.V GT 500.0*SEC.DAYS
CANCEL CPU.JUTILIZATION
RETURN
ELSE
RETURN
END

EVENT B.JOB.REQUEST SAVING THE EVENT NOTICE
DEFINE PL AS REAL VARIABLE
DEFINE I.O AS REAL VARIABLE
DEFINE CP.T AS REAL VARIABLE
*****
** NEEDED.RESOURCES.F IS A NUMERICAL INDICATION OF RESOURCES
** NEEDED
** NBJ IS THE NUMBER OF BATCH JOBS IN THE SYSTEM
**
** JOB.TYPE 1=BATCH,2=INTERACTIVE
**
** JOB.IO.REQUESTS IS THE TOTAL NUMBER OF I/O OPERATIONS
** REQUESTED BY A SINGLE JOB
**
** IO.SERVICE.TIME IS I/O TIME REQUIRED TO PROCESS A SINGLE I/O
** TASK
**
** XX.CPU.SERVICE.TIME IS THE CPU TIME REQUIRED TO PROCESS
** UNTIL AN I/O
**
** JIS IS THE NUMBER OF JOB INITIATORS AVAILABLE
**

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 005580

```

** MEMORY.REQUEST IS THE AMOUNT OF COPE NEEDED
**
** PAGE.LENGTH IS THE TOTAL PAGES IN THE OUTPUT
**
** T.LEFT.IN.CPU IS THE AMOUNT OF TIME LEFT UNTIL NEXT I/O
** TM = TAPE.MOUNT = 0 MEANS JOB DOES NOT NEED TAPE MOUNTED
**           = 1 MEANS JOB NEEDS TAPE MOUNTED
** OM = DISK.MOUNT = 0 MEANS JOB DOES NOT NEED DIS.MOUNTED
**           = 1 MEANS JOB NEEDS DISK MOUNTED
**
** CLASS          REGION          TIME          TAPE MOUNTS          DISC MOUNT
** A             K<=200          T<5           PERMITTED           NO
** B             K<=260          T<60          NO                   NO
** C             250<K<=500      T<60          NO                   NO
** D             K<=200          T<=5          YES                   NO
** E             K<=250          T<=60         YES                   NO
** F             250<K<=500      T<=60         YES                   NO
** G             K<=200          T<=5          YES                   YES
** H             K<=260          T<=60         YES                   YES
** I             250<K<=500      T<=60         YES                   YES
** J             K<=500          T<=60         YES                   NO
** L             K<=500          T<=60         YES                   YES
** WHERE K IS THE REGION SIZE IN BLOCKS OF 1024 BYTES AND T IS CPU
** TIME IN MINUTES.
*****
IF TIME.V GT .004
LET WARM.UP.FLAG=WARM.UP.FLAG+1
ELSE
ALWAYS
IF WARM.UP.FLAG=1
LET W.UP.TIME=TIME.V
ELSE
ALWAYS
IF NBJ GE JIS
SCHEDULE THE B.JOB.REQUEST IN 1 MINUTES
GO TO 'R'
ELSE

```

```

IF R.COUNTER GE ANUMBRER
SCHEDULE THE B.JOR.REQUEST IN POISSON.F(.025,3) HOURS
ELSE
LET R.COUNTER = R.COUNTER+1
ALWAYS
CREATE A JOB CALLED AMIS
LET JN = JN +1
PRINT 1 LINE WITH JN AND TIME.V THUS
***.
LET JOBNUM(AMIS) = JN
***ESTABLISHES THIS IS A BATCH JOB
LET JOR.TYPE(AMIS) = 1
LET NRJ = NRJ + 1
LET TLF(AMIS) = 0
***ESTABLISHES BATCH CP.TIME
LET CP.TIME(AMIS) =NORMAL.F(26.6,24.0,5)*SEC.DAYS
IF CP.TIME(AMIS) GT 3590.0*SEC.DAYS
LET CP.TIME(AMIS) = 3590.0*SEC.DAYS
ELSE
ALWAYS
IF CP.TIME(AMIS) LT 10.0*SEC.DAYS
LET CP.TIME(AMIS)=10.0*SEC.DAYS
ELSE
ALWAYS
LET T = CP.TIME(AMIS)
LET DISK.MOUNT = RANDI.F(1,100,8)
IF DISK.MOUNT GT 85
LET DISK.MOUNT = 0
ELSE
LET DISK.MOUNT = 1
ALWAYS
LET TAPE.MOUNT = RANDI.F(1,100,7)
IF TAPE.MOUNT GT 70
LET TAPE.MOUNT = 0
ELSE
LET TAPE.MOUNT = 1
ALWAYS

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005880
005890
005900
005910
005920
005930
005940
005950

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'' ****ESTABLISHES IO.TIME
LET IO.TIME(AMIS) = NORMAL.F(40.0, 80.0,2)*SEC.DAYS
IF IO.TIME(AMIS) LE 40.0*SEC.DAYS
LET IO.TIME(AMIS)=40.0*SEC.DAYS
ELSE
ALWAYS
'' **** ESTABLISHING CORE REQUIREMENTS
LET MEMORY.REQUEST(AMIS) = RANDI.F(100,500,7)
LET K = MEMORY.REQUEST(AMIS)
'' ****ESTABLISHES JOB RESOURCES NEEDED
LET RE.SOURCES(AMIS) = 1
LET NEEDED.RESOURCES.F(AMIS) = RANDI.F(10,20,1)
LET P.OF.RESOURCES(AMIS) = NEEDED.RESOURCES.F(AMIS)

'' ****ESTABLISHES THE NUMBER OF I/O REQUESTS
LET JOB.IO.REQUESTS(AMIS)=NORMAL.F(900.0,90.0,3)
IF JOB.IO.REQUESTS(AMIS) LT 5
LET JOB.IO.REQUESTS(AMIS)=RANDI.F(5,30,7)
ELSE
ALWAYS

'' ****ESTABLISHES PAGE LENGTH
LET PL = (JOB.IO.REQUESTS(AMIS)/1)*LOG.NORMAL.F(10.0,5.0,1)
LET PAGE.LENGTH(AMTS) = INT.F(PL) + 1

'' ****ESTABLISHES IO.TIME FOR 1ST I/O TASK
LET I.O = 1/JOB.IO.REQUESTS(AMIS)
LET IO.SERVICE.TIME(AMIS) = I.O * IO.TIME(AMIS)

'' ****ESTABLISHING THE PRIORITY OF THE JOB
'' ****ESTABLISHING CLASS PRIORITY
IF K LE 200 AND T LE 2600*SEC.DAYS AND TM = 0 AND DM = 0
LET PRIORITY(AMIS) = RANDI.F(190,200,1)
GO TO *BYPASS

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005960
005970
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005990
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006320
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006670

```
ELSE  
IF K LE 260 AND T LE 3600*SEC.DAYS AND TM = 0 AND DM = 0  
LET PRIORITY(AMIS) = RANDI.F(180,189,2)  
GO TO 'BYPASS'  
ELSE  
IF K GT 260 AND K LE 500 AND T LE 3600*SEC.DAYS AND TM = 0 AND DM = 0  
LET PRIORITY(AMIS) = RANDI.F(170,179,3)  
GO TO 'BYPASS'  
ELSE  
IF K LE 200 AND T LE 300*SEC.DAYS AND TM = 1 AND DM = 0  
LET PRIORITY(AMIS) = RANDI.F(160,169,4)  
GO TO 'BYPASS'  
ELSE  
IF K LE 260 AND T LE 3600*SEC.DAYS AND TM = 1 AND DM = 0  
LET PRIORITY(AMIS) = RANDI.F(150,159,5)  
GO TO 'BYPASS'  
ELSE  
IF K GT 260 AND K LE 500 AND T LE 3600*SEC.DAYS AND TM = 1 AND DM = 0  
LET PRIORITY(AMIS) = RANDI.F(140,149,6)  
GO TO 'BYPASS'  
ELSE  
IF K LE 200 AND T LE 300*SEC.DAYS AND TM = 1 AND DM = 1  
LET PRIORITY(AMIS) = RANDI.F(130,139,7)  
GO TO 'BYPASS'  
ELSE  
IF K LE 260 AND T LE 3600*SEC.DAYS AND TM = 1 AND DM = 1  
LET PRIORITY(AMIS) = RANDI.F(120,129,8)  
GO TO 'BYPASS'  
ELSE  
IF K GT 260 AND K LE 500 AND T LE 3600*SEC.DAYS AND TM = 1 AND DM = 1  
LET PRIORITY(AMIS) = RANDI.F(110,119,9)  
GO TO 'BYPASS'  
ELSE  
IF K LE 500 AND T LE 3600*SEC.DAYS AND TM = 1 AND DM = 0  
LET PRIORITY(AMIS) = RANDI.F(100,109,1)  
GO TO 'BYPASS'  
ELSE
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 005980
 006990
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 007010
 007020
 007030

```

IF K LE 500 AND T LE 3600*SFC.DAYS AND TM = 1 AND DM = 1
LET PRIORITY(AMIS) = RANDI.F(90,99,2)
ELSE
ALWAYS
LET PRIORITY(AMIS)=80
'BYPASS'
'' *****ESTABLISHES THE JOB STARTING TIME
LET F.LOADER(AMIS)=0
LET ARRIVAL.TIME(AMIS) = TIME.V
'' *****ESTABLISHING THE NEEDED CPU TIME UNTIL 1ST I/O
LET CP.T=1/JOB.IO.REQUESTS(AMIS)
LET XX.CPU.SERVICE.TIME(AMIS) = CP.T * CP.TIME(AMIS)
LET XX.CPU.SERVICE.TIME(AMIS) = XX.CPU.SERVICE.TIME(AMIS)
IF FF=0
PRINT 1 LINE THUS
  BATCH STATS
PRINT 1 LINE WITH JOB.TYPE(AMIS), JORNUM(AMIS), IO.SERVICE.TIME(AMIS),
PRIORITY(AMIS), XX.CPU.SERVICE.TIME(AMIS), XX.CPU.SERVICE.TIME(AMIS),
CP.TIME(AMIS), JOB.IO.REQUESTS(AMIS) AND TIME.V THUS
* . . . . . ***** . . . . . ***** . . . . . ***** . . . . . *****
ELSE
ALWAYS
LET TOTAL.JOB.CPU.TIME=TOT.JOB.CPU.TIME+IO.TIME(AMIS)
LET TOT.JOB.IO.TIME=TOT.JOB.IO.TIME+IO.TIME(AMIS)
LET R.MEMORY.REQUEST = MEMORY.REQUEST(AMIS)
LET NSUB = NBJ
LET R.CPU.TIME = CP.TIME(AMIS)
LET BATCH.IO=IO.TIME(AMIS)
FILE AMIS IN THE JOB.QUEUE
SCHEDULE A REQUEST.CM NOW
'R'
RETURN
END

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007040
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```

EVENT P.INCREASE SAVING THE EVENT NOTICE
*****
**      FLAG = 0 MEANS NO JOB HAS PRIORITY INCREASE
**      1 MEANS A JOB HAS HAD ITS PRIORITY INCREASED
*****
FOR EACH JOB IN THE CPU..REQUEST.QUEUE, WITH
PRIORITY = 5, FIND THE FIRST CASE
IF NONE
  SCHEDULE THE P.INCREASE IN 10 HOURS
  RETURN
ELSE
  IF PFLAG = 1
    SCHEDULE THE P.INCREASE IN 10.0*MIL.SEC.DAYS DAYS
  RETURN
ELSE
  LET PFLAG = 1
  REMOVE JOB FROM CPU..REQUEST.QUEUE
  LET PRIORITY(JOB) = 15
  FILE JOB IN CPU..REQUEST.QUEUE
  REMOVE FIRST AMIS FROM CPU..REQUEST.QUEUE
  FILE AMIS IN CPU..REQUEST.QUEUE
  SCHEDULE THE P.INCREASE IN SECOND MINUTES
  RETURN
END

EVENT P.OUTP.INCREASE SAVING THE EVENT NOTICE
*****
**      POUTFLAG=1 MEANS NO JOB HAS ITS PRIORITY INCREASED
**      1 MEANS A JOB HAS ITS PRIORITY INCREASED
*****
FOR EACH JOB IN THE AM.OUTPUT.QUEUE, WITH
PRIORITY = 5, FIND THE FIRST CASE
IF NONE
  SCHEDULE THE P.OUTP.INCREASE IN 10 HOURS
  RETURN
ELSE

```

```

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007980
007990
008000
008010
008020
008030
008040

IF POUTFLAG=1
SCHEDULE THE P.OUTP.INCREASE IN 10.0*MIL.SEC.DAYS DAYS
RETURN
ELSE
LET FOUTFLAG=1
REMOVE JOB FROM AN.OUTPUT.QUEUE
LET PRIORITY(JOB) = 15
FILE JOB IN AN.OUTPUT.QUEUE
REMOVE FIRST AMIS FROM AN.OUTPUT.QUEUE
FILE AMIS IN THE AN.OUTPUT.QUEUE
SCHEDULE THE P.OUTP.INCREASE IN SECOND MINUTES
RETURN
END

*****
'' PJQFLAG = 0 MEANS NO JOB HAS HAD ITS PRIORITY INCREASED *
''          1 MEANS A JOB HAS ITS PRIORITY INCREASED *
*****
EVENT PJQ.INCREASE SAVING THE EVENT NOTICE
FOR EACH JOB IN THE JOB.QUEUE, WITH
PRIORITY = 5, FIND THE FIRST CASE
IF NCNE
SCHEDULE THE PJQ.INCREASE IN 10 HOURS
RETURN
ELSE
IF PJQFLAG = 1
SCHEDULE THE PJQ.INCREASE IN 10.0*MIL.SEC.DAYS DAYS
RETURN
ELSE
LET PJQFLAG=1
REMOVE JOB FROM JOB.QUEUE
LET PRIORITY(JOB) = 15
FILE JOB IN JOB.QUEUE
REMOVE AMIS FROM JOB.QUEUE
FILE AMIS IN JOB.QUEUE
SCHEDULE THE PJQ.INCREASE IN SECOND MINUTFS
RETURN
END

```

AD-A055 209

AIR FORCE INST OF TECH WRIGHT-PATTERSON AFB OHIO SCH--ETC F/G 5/1
SIMULATION OF THE ACQUISITION MANAGEMENT INFORMATION SYSTEM.(U)
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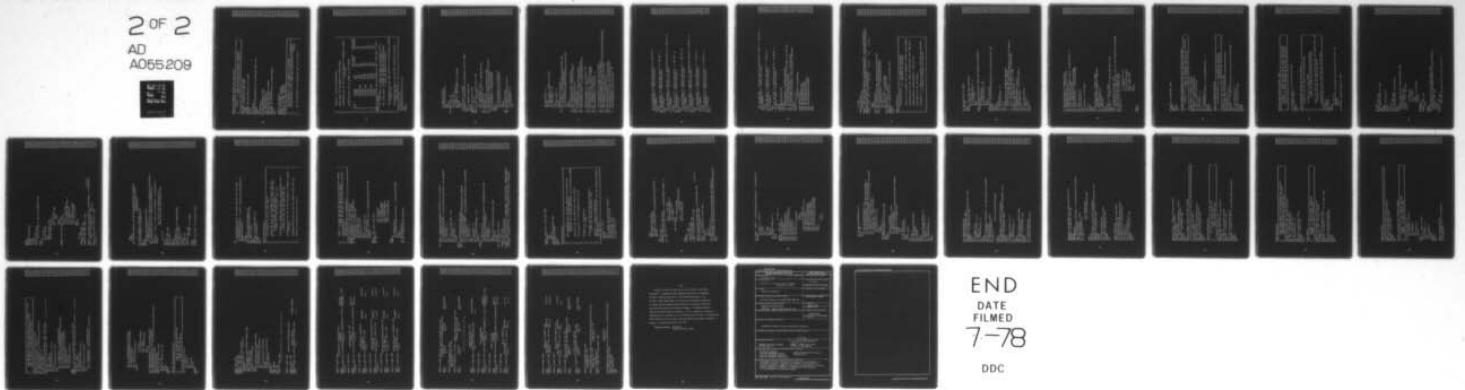
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2 OF 2

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EVENT PCM.INCREASE SAVING THE EVENT NOTICF
*****
**      PCMFLAG = 0 MEANS NO JOB HAS ITS PRIORITY INCREASED
**      1 MEANS A JOB HAS ITS PRIORITY INCREASED
*****
FOR EACH JOB IN THE CM.QUEUE,WITH
PRIORITY = 5,FIND THE FIRST CASE
IF NONE
SCHEDULE THE PCM.INCREASE IN 10 HOURS
RETURN
ELSE
IF PCMFLAG = 1
SCHEDULE THE PCM.INCREASE IN 10.0*MIL.SEC.DAYS DAYS
RETURN
ELSE
LET PCMFLAG=1
REMOVE JOB FROM CM.QUEUE
LET PRIORITY(JOB) = 15
FILE JOB IN CM.QUEUE
REMOVE FIRST AMIS FROM CM.QUEUE
FILE AMIS IN CM.QUEUE
SCHEDULE THE PCM.INCREASE IN SECOND MINUTES
RETURN
END

EVENT I.JOB.REQUEST
DEFINE PL AS REAL VARIABLE
DEFINE I.O AS REAL VARIABLE
DEFINE CP.T AS REAL VARIABLE
*****
**      JOB..IO.REQUESTS = TOTAL NUMBER I/O REQUESTS
**      JOB.IO.REQUESTS = TOTAL NUMBER OF I/O REQUESTS PER CARRIAGE RET
**      NEEDED.RESOURCFS.F IS A NUMERICAL INDICATION OF RESOURCES
**      NEEDED
**      NIJ IS THE NUMBER OF INTERACTIVE JOBS IN THE SYSTEM
**

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IF WARM.UP.FLAG=1
LET W.UP.TIME=TIME.V
ELSE
  ALWAYS
LET NIJ = NIJ + 1
IF NIJ GT 35
LET NIJ = NIJ-1
LFT NJNPLO = NJNPLO + 1
SCHEDULE A I.JOB.REQUEST IN 1 MINUTES
RETURN
ELSE
IF I.COUNTER GE INUMMER
SCHEDULE A I.JOB.REQUEST IN POISSON.F(.1117,8) HOURS
ELSE
LET I.COUNTER=I.COUNTER+1
  ALWAYS
CREATE A JOB CALLED AMIS
LET JN = JN +1
PRINT 1 LINE WITH JN AND TIME.V THUS
  ***.*****
  LET JOBNUM(AMIS) = JN
  *****ESTABLISHES THIS IS A INTERACTIVE JOB
LET JOB.TYPE(AMIS) = 2
LET TLF(AMIS) = 0
  *****ESTABLISHES INTERACTIVE CP.TIME
LFT CP.TIME(AMIS) =NORMAL.F(41.0,48.0,5)*SEC.DAYS
IF CP.TIME(AMIS) GT 3590.0*SEC.DAYS
LET CP.TIME(AMIS) = 3590.0*SEC.DAYS
ELSE
  ALWAYS
IF CP.TIME(AMIS) LT 15.0*SEC.DAYS
LET CP.TIME(AMIS)=15.0*SEC.DAYS
ELSE
  ALWAYS
LET T = CP.TIME(AMIS)
LET DISK.MOUNT = PANDT.F(1,10(,8)
IF DISK.MOUNT GT 85
LFT DISK.MOUNT = 0
ELSE

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009370 LET DISK.MOUNT = 1
009380 ALWAYS
009390 LET TAPE.MOUNT = RANDI.F(1,100,7)
009410 IF TAPE.MOUNT GT 70
009410 LET TAPE.MOUNT = 0
009420 ELSE
009430 LET TAPE.MOUNT = 1
009440 ALWAYS
009450 ***ESTABLISHES IO.TIME
.. LET IO.TIME(AMIS) = NORMAL.F(256.0, 64.0,2)*SEC.DAYS
.. IF IO.TIME(AMIS) LT 50.0*SEC.DAYS
.. LET IO.TIME(AMIS)=UNIFORM.F(50.0,100.0,4)*SEC.DAYS
.. ELSE
.. ALWAYS
.. ***ESTABLISHES CORE REQUIREMENTS
LET MEMORY.REQUEST(AMIS) = RANDI.F(100,500,7)
LET K = MEMORY.REQUEST(AMIS)
LET RE.SOURCES(AMIS) = 1
.. ***ESTABLISHES JOB.RESOURCES NEEDED
LET NEEDED.RESOURCES.F(AMIS) = RANDI.F(10,20,1)
LET P.OF.RESOURCES(AMIS) = NEEDED.RESOURCES.F(AMIS)
..
.. ***ESTABLISHES THE JOB STARTING TIME
LET ARRIVAL.TIME(AMIS) = TIME.V
LET NO.CAR.RETURNS(AMIS) = RANDI.F(5,20,6)
.. ***ESTABLISHES PAGE LENGTH OF OUTPUT
LET PL = LOG.NORMAL.F(20.0,2.0,9)
LET PAGE.LENGTH(AMIS) = INT.F(PL) + 1
..
.. ***ESTABLISHES THE NUMBR OF I/O TASKS
LET JOB..IO.REQUESTS(AMIS)=RANPI.F(2000,4000,9)
LET JOB..IO.REQUESTS(AMIS)=INT.F(JOB..IO.REQUESTS(AMIS)/NO.CAR.RETURNS(AMIS))
.. ***ESTABLISHES IO.TIME FOR 1ST I/O TASK
LET I.O = 1/JOB..IO.REQUESTS(AMIS)
LET IO.SERVICE.TIME(AMIS) = I.O * IO.TIME(AMIS)
LET JOB..IO.REQUESTS(AMIS)=JOB..IO.REQUESTS(AMIS)
.. ***ESTABLISHING PRIORITY OF THIS JOB
.. ***ESTABLISHING CLASS PRIORITY

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010010
010020
010030
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010050
010060

```
IF K LE 200 AND T LE 3600*SEC.DAYS AND TM = 0 AND OM = 0
  LET PRIORITY(AMIS) = RANDI.F(190,200,1)
  GO TO *PASSBY*
ELSE
  IF K LE 260 AND T LE 3600*SEC.DAYS AND TM = 0 AND OM = 0
  LET PRIORITY(AMIS) = RANDI.F(180,189,2)
  GO TO *PASSBY*
ELSE
  IF K GT 260 AND K LE 500 AND T LE 3600*SEC.DAYS AND TM = 0 AND OM = 0
  LET PRIORITY(AMIS) = RANDI.F(170,179,3)
  GO TO *PASSBY*
ELSE
  IF K LE 200 AND T LE 300*SEC.DAYS AND TM = 1 AND OM = 0
  LET PRIORITY(AMIS) = RANDI.F(160,169,4)
  GO TO *PASSBY*
ELSE
  IF K LE 260 AND T LE 3600*SEC.DAYS AND TM = 1 AND OM = 0
  LET PRIORITY(AMIS) = RANDI.F(150,159,5)
  GO TO *PASSBY*
ELSE
  IF K GT 260 AND K LE 500 AND T LE 3600*SEC.DAYS AND TM = 1 AND OM = 0
  LET PRIORITY(AMIS) = RANDI.F(140,149,6)
  GO TO *PASSBY*
ELSE
  IF K LE 200 AND T LE 300*SEC.DAYS AND TM = 1 AND OM = 1
  LET PRIORITY(AMIS) = RANDI.F(130,139,7)
  GO TO *PASSBY*
ELSE
  IF K LE 260 AND T LE 3600*SEC.DAYS AND TM = 1 AND OM = 1
  LET PRIORITY(AMIS) = RANDI.F(120,129,8)
  GO TO *PASSBY*
ELSE
```

```

IF K GT 260 AND K LE 500 AND T LE 3600*SEC.DAYS AND TM = 1 AND DM = 1
  LET PRIORITY(AMIS) = RANDI.F(110,119,9)
  GO TO *PASSBY*
ELSE
  IF K LE 500 AND T LE 3600*SEC.DAYS AND TM = 1 AND DM = 0
  LET PRIORITY(AMIS) = RANDI.F(100,109,1)
  GO TO *PASSBY*
ELSE
  IF K LE 500 AND T LE 3600*SFC.DAYS AND TM = 1 AND DM = 1
  LET PRIORITY(AMIS) = RANDI.F(90,99,2)
  ELSE
  ALWAYS
  LET PRIORITY(AMIS)=80
  *PASSBY*
  LET PRIORITY(AMIS) = PRIORITY(AMIS)+200
  ** **ESTABLISHING THE NEEDED CPU TIME UNTIL 1ST I/O TASK
  LET CP.T=1/JOB..IO.REQUESTS(AMIS)
  LET XX.CPU.SERVICE.TIME(AMIS) = CP.T * CP.TIME(AMIS)
  LET XX..CPU.SERVICE.TIME(AMIS) = XX.CPU.SERVICE.TIME(AMIS)
  F FLAG=1
  PRINT 1 LINE THIS
  INTERACTIVE PARAMETERS
  LIST JOBNUM(AMIS)
  LIST NO.CAR.RETURNS(AMIS)
  LIST JO9.IO.REQUESTS(AMIS)
  LIST JOB..IO.REQUESTS(AMIS)
  LIST XX..CPU.SERVICE.TIME(AMIS)
  LSE
  LWAYS

```

```

010070
010080
010090
010100
010110
010120
010130
010140
010150
010160
010170
010190
010199
010200
010210
010220
010230
010240
010250
010250
010270
010280
010290
010300
010310
010320
010330
010340
010350
010360

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010370
 010380
 010390
 010400
 010410
 010420
 010430
 010440
 010450
 010460
 010470
 010480
 010490
 010500
 010510
 010520
 010530
 010540
 010550
 010560
 010570
 010580
 010590
 010600
 010700
 010710
 010720
 010730
 010740
 010750
 010760
 010770
 010780
 010790
 010800
 010810
 010820

```

IF CF=0
PRINT 1 LINE THUS
INTERACTIVE STAS
PRINT 1 LINE WITH JO3NUM(AMIS), NO.CAR.RETURNS(AMIS), JO3.IO.REQUESTS(AMIS),
JOB..IO.REQUESTS(AMIS), XX..CFU.SERVICE.TIME(AMIS), XX.CPU.SERVICE.TIME(AMIS),
PRIORITY(AMIS), AND IO.TIME(AMIS) THUS
***. ** .***** .***** .*****
ELSE
ALWAYS
LET F.LOADER(AMIS)=0
LET TOTAL.JOB.CPU.TIME=TOTAL.JOB.CPU.TIME+CP.TIME(AMIS)
LET TOT.JOB.IO.TIME=TOT.JOB.IO.TIME+IO.TIME(AMIS)
LET I.MEMORY.REQUEST = MEMORY.REQUEST(AMIS)
LET NSUI = NIJ
LET I.CPU.TIME = CP.TIME(AMIS)
LET INTERACTIVE.IO = IO.TIME(AMIS)
FILE AMIS IN THE JOB.QUEUE
SCHEDULE A REQUEST.CM IN 1.0 *MIL.SEC.DAYS DAYS
RETURN
END

EVENT REQUEST.CM
*****
** RE.SOURCES(AMIS) = 0 - ALL RESOURCES AVAILABLE
** = 1 - RESOURCES NOT AVAILABLE
**
** FLAG.C IS THE COUNTER USED TO DECREMENT CURRENT MEMORY.BEGINS
**
** CONTIGOUS.SPACE IS THE AMOUNT OF CONTIGUOUS SPACE AVAILABLE
**
** MEMORY.BLOCKS = MEMORY.REQUEST = CORE NEEDED
**
** MAIN.MEMORY IS THE CORE ARRAY
**
** RESOURCES.F = 0 MEANS ALL RESOURCES REQUESTED ARE AVAILABLE
** = 1 MEANS ALL RESOURCES NOT AVAILABLE
**
*****

```

```

IF THE JOB.QUEUE IS EMPTY
GO TO 'RE'
ELSE
REMOVE FIRST AMIS FROM JOB.QUEUE
LET PJOFLAG = 0
IF RE.SOURCES(AMIS) = 0
GO TO 'ENTER'
ELSE
CALL TRAFFIC.CONTROLLER GIVEN AMIS YIELDING RESOURCES.F
IF RESOURCES.F = 1
IF PRIORITY(AMIS) = 15
LET PRIORITY(AMIS) = 5
ELSE
ALWAYS
FILE AMIS IN WAIT.QUEUE
SCHEDULE A XX.TRAFFIC.CONTROLLER IN (1.0*MIL.SEC.DAYS) DAYS
RETURN
ELSE
'ENTER'
LET FLAG.C = 0
LET CONTIGUOUS.SPACE = 0
LET X = 1
LET Y = 2
LET MEMORY.BLOCKS = MEMORY.REQUEST(AMIS)
'' *****MRC IS MEMORY NEEDED COUNTER
'' ***** 0 IN MAIN.MEMORY MEANS THIS SECTION OF CORE IS AVAIRLE
'' ***** 1 MEANS THIS SECTION OF CORE IS BEING USED
FOR MRC = 1 TO 4000, UNTIL X EQ Y, DO
IF MAIN.MEMORY(MRC) EQ 0
LET MEMORY.REGINS=MRC - FLAG.C
LET CONTIGUOUS.SPACE = CONTIGUOUS.SPACE+1
LET FLAG.C=FLAG.C+1
ELSE
LET CONTIGUOUS.SPACE=0
LET FLAG.C=0
GO TO 'RET'
ALWAYS

```

```

010830
010840
010850
010860
010870
010880
010890
010900
010910
010920
010930
010940
010950
010960
010970
010980
010990
011000
011010
011020
011030
011040
011050
011060
011070
011080
011090
011100
011110
011120
011130
011140
011150
011160
011170
011180
011190

```

```

IF CCNTIGOUS.SPACE GE MEMORY.FLOCKS
LET START.MAIN.MEMORY = MEMORY.REGINS
LET SMM=START.MAIN.MEMORY
LET FINISH.MAIN.MEMORY= MEMORY.BEGINS + MEMORY.FLOCKS -1
LET FMM = FINISH.MAIN.MEMORY
LET START.MEMORY(AMIS)=START.MAIN.MEMORY
LET END.MEMORY(AMIS) = FINISH.MAIN.MEMORY
LET Y = 1
FOR J=SMM TO FMM, DO
LET MAIN.MEMORY(J)=1
LOOP
LET JOB.SIZE = (FMM-SMM)+1
ELSE
LET Y = 2
ALWAYS
*RET*
LOOP
IF CCNTIGOUS.SPACE LT MEMORY.FLOCKS
** *****THERE ISN'T ENOUGH CM AVAILABLE
IF PRIORITY(AMIS) LE 15
LET PRIORITY(AMIS) = 5
LET PCMFLAG=0
ELSE
ALWAYS
FILE AMIS IN THE CM.QUEUE
SCHEDULE A REQUEST.CM IN (1.0*MIL.SEC.DAYS) DAYS
ELSE
CALL LOADER GIVEN JOB.SIZE, AMIS
** *****XXTIME.E.CPU.QUEUE IS THE TIME THE JOB ENTERED THE CPU.QUEUE
LET XXTIME.E.CPU.QUEUE(AMIS)=TIME.V
FILE AMIS IN THE CPU..RFQUEST.QUEUE
IF FLAG = 1
IF JOBNUM(AMIS) EQ 2
LIST TIME.V
LIST STATUS.PRIORITY
PRINT 1 LINE THUS
EVTN REQEST.CM
ELSE
ALWAYS
ELSE
ALWAYS

```

```

011200
011210
011220
011230
011240
011250
011260
011270
011280
011290
011300
011310
011320
011330
011340
011350
011360
011370
011380
011390
011400
011410
011420
011430
011440
011450
011460
011470
011480
011490
011500
011510
011520
011530
011540
011550
011560
011570
011580

```

011600
 011610
 011620
 011630
 011640
 011650
 011660
 011670
 011680
 011690
 011700
 011710
 011720
 011730
 011740
 011750
 011760
 011770
 011780
 011790
 011800
 011810
 011820
 011830
 011840
 011850
 011860
 011870
 011880
 011890
 011900
 011910
 011920
 011930
 012030
 012040
 012050
 012060
 012070
 012080
 012090
 012100
 012110
 012120
 012130
 012140
 012150
 012160
 012170
 012180
 012190
 012190

```

SCHEDULE A CPU.REQUEST NOW
ALWAYS
  RE
RETURN
END

ROUTINE TRAFFIC.CONTROLLER GIVEN AMIS YIELDING RESOURCES.F
*****
'' RESOURCES.F = 0 MEANS ALL RESOURCES REQUESTED ARE AVAILABLE *
'' RESOURCES.F = 1 MEANS ALL RESOURCES NOT AVAILABLE *
*****
DEFINE UPPER AS INTEGER VARIABLE
DEFINE RE.F AS INTEGER VARIABLE
LET UPPER = P.OF.RESOURCES(AMIS)
LET RE.F = RANDI.F(1, UPPER, 3)
IF RE.F LE 8
LET RESOURCES.F = 0
ELSE
LET RESOURCES.F = 1
LET P.OF.RESOURCES(AMIS) = P.OF.RESOURCES(AMIS) - 2
ALWAYS
RETURN
END

EVENT XX.TRAFFIC.CONTROLLER
*****
'' RE.SOURCES(AMIS) = 0 - ALL RESOURCES AVAILABLE *
'' RE.SOURCES(AMIS) = 1 - RESOURCES NOT AVAILABLE *
*****
REMOVE FIRST AMIS FROM WAIT.QUEUE
CALL TRAFFIC.CONTROLLER GIVEN AMIS YIELDING RESOURCES.F
IF RESOURCES.F = 1
FILE AMIS IN WAIT.QUEUE
SCHEDULE A XX.TRAFFIC.CONTROLLER IN (1.0*MIL*SEC.DAYS) DAYS
ELSE
LET RE.SOURCES(AMIS) = 0
FILE AMIS IN JOB.QUEUE
SCHEDULE A REQUEST.CM NOW
ALWAYS
RETURN
END
  
```

012200
 012300
 012310
 012320
 012330
 012340
 012350
 012360
 012370
 012380
 012390
 012400
 012410
 012510
 012520
 012530
 012540
 012550
 012560
 012570
 012580
 012590
 012600
 012610
 012620
 012630
 012640
 012650
 012660
 012670
 012680
 012690
 012700
 012710
 012720
 012730

```

ROUTINE LOADER GIVEN JOB.SIZE, AMIS
*****
** F.LOADER = 1 MEANS JOB NEEDS TO BE LOADED INTO MEMORY **
** F.LOADER = 0 MEANS JOB NEED NOT BE LOADED INTO MEMORY **
** LOAD.T IS THE TIME NECESSARY TO LOAD JOB INTO MEMORY **
*****
LET LOADING.T(AMIS) = (1.0/10(0.0))*SEC.DAYS
LET F.LOADER (AMIS) = 1
RETURN
END

EVENT CPU.REQUEST SAVING THE EVENT NOTICE
*****
** CPU.SF IS THE CPU STATUS FLAG 1=BUSY,0=IDLE **
**
** F.LOADER IS THE FLAG LOADER **
** 1 MEANS JOB NEEDS TO BE LOADED INTO MEMORY **
** 0 MEANS JOB NEEDS CPU FOR ACTUAL PROCESSING **
**
** LOAD.F = 0 MEANS THE JOB DOES NOT NEED CPU FOR LOADING INTO MEM*
** =1 MEANS JOB MUST BE LOADED INTO MEMORY **
**
** BLOCKED.TIME IS THE TOTAL TIME IN THE CPU.QUEUE **
*****
IF FLAG = 1
LIST CPU.SF
LIST LOAD.F
ELSE
ALWAYS
IF LOAD.F=1
SCHEDULE THE CPU.REQUEST IN (6.0/100(0.0))*SEC.DAYS DAYS
RETURN
ELSE
IF THE CPU.REQUEST.QUEUE IS EMPTY

```

```

IF FLAG = 1
PRINT 1 LINE THUS
CPU..REQUEST QUEUE IS EMPTY
ELSE
ALWAYS
DESTROY CPU.REQUEST
RETURN
ELSE
IF FLAG = 1
PRINT 1 LINE THUS
CPU..REQUEST.QUEUE IS NOT EMPTY
ELSE
ALWAYS
FOR EACH AMIS IN CPU..REQUEST.QUEUE, WITH
F.LOADER(AMIS) EQ 1
FIND THE FIRST CASE
IF NONE GO TO 'PROCESS'
ELSE
IF CPU.SF=0
LET LOAD.F=1
GO TO 'CKS'
ELSE
DESTROY CPU.REQUEST
RETURN
'PROCESS'
IF FLAG = 1
PRINT 1 LINE THUS
3RR3
ELSE
ALWAYS
FOR EACH AMIS IN CPU..REQUEST.QUEUE, WITH
F.LOADER(AMIS)=0 AND PRIORITY(AMIS) GT STATUS.PRIORITY,
FIND THE FIRST CASE
IF NONE
IF FLAG = 1
PRINT 1 LINE THUS
THERE ARE NO JOBS THAT MEET THIS REQUIREMENT

```

```

012740
012750
012760
012770
012780
012790
012800
012810
012820
012830
012840
012850
012860
012870
012880
012890
012900
012910
012920
012930
012940
012950
012960
012970
012980
012990
013000
013010
013020
013030
013040
013050
013060
013070
013080
013090
013100

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013510
 013520
 013530
 013540
 013550
 013560
 013570
 013580
 013590
 013600
 013610
 013620
 013630
 013640
 013650
 013660
 013670
 013680
 013690
 013700
 013710
 013720
 013730
 013740
 013750
 013760
 013770
 013780
 013790
 013800
 013810
 013820
 013830
 013840
 013850
 013860
 013870

```

ELSE
  ALWAYS
  LET XX.CPU.SERVICE.TIME(ANDY)=XX.CPU.SERVICE.TIME(ANDY)
    -(TIME.V-T.E.CPU(ANDY))
    -IO.OVERHEAD.T
IF XX.CPU.SERVICE.TIME(ANDY) LE 0.0
  LET XX.CPU.SERVICE.TIME(ANDY)=0.0
ELSE
  ALWAYS
  LET XXTIME.E.CPU.QUEUE(ANDY)=TIME.V
    IF WARM.UP.FLAG GE 1
  LET T.CPU.TIME=T.CPU.TIME+((TIME.V-T.E.CPU(ANDY))-IO.OVERHEAD.T)
  LET GAIN=GAIN+((TIME.V-T.E.CPU(ANDY))-IO.OVERHEAD.T)
ELSE
  LET TPUT = TPUT + (TIME.V-T.E.CPU(ANDY) - IO.OVERHEAD.T)
  ALWAYS
  FILE ANDY IN CPU..REQUEST.QUEUE
  REMOVE AMIS FROM CPU..REQUEST.QUEUE
  GO TO 'C'
'CKS'
IF FLAG = 1
  PRINT 1 LINE THUS
  JOB AT CKS
ELSE
  ALWAYS
  LET CPU.SF = 1
  REMOVE AMIS FROM CPU..REQUEST.QUEUE
  FILE AMIS IN ACPU
  GO TO 'SCH'
'CHECK'
LET CPU.SF=1
IF FLAG = 1
  PRINT 1 LINE THUS
    JOB MADE IT TO CHECK
ELSE
  ALWAYS
  REMOVE AMIS FROM CPU..REQUEST.QUEUE

```

```

013880
013890
013900
013910
013920
013930
013940
013950
013960
013970
013980
013990
014000
014010
014020
014030
014130
014140
014150
014160
014170
014180
014190
014200
014210
014220
014230
014240
014250
014260
014270
014280
014290
014300
014310
014320

'C' *****BLOCKED.TIME IS THE TOTAL TIME IN THE CPU QUEUE
'' IF WARM.UP.FLAG GE 1
LET BLOCKED.TIME (AMIS) = TIME.V - XXTIME.E.CPU.QUEUE (AMIS) +
BLOCKED.TIME (AMIS)
ELSE
ALWAYS
LET XXTIME.E.CPU.QUEUE (AMIS) = (.0
LET STATUS.PRIORITY = PRIORITY (AMIS)
FILE AMIS IN ACPU
'SCH'
SCHEDULE A CPU.PROCESSING NOW
DESTROY CPU.REQUEST
RETURN
END

EVENT CPU.PROCESSING
*****
'' CPU.T.QUANTUM IS THE CPU TIME QUANTUM
''
'' T.LEFT.IN.CPU IS THE REMAINING TIME UNTIL NEXT I/O
'' TLF IS THE TIME LEFT FOR I/O TASK (TIME LEFT FLAG)
'' 0 MEANS NO MORE CPU TIME NEEDED
'' 1 MEANS MORE TIME NEEDED UNTIL NEXT I/O
''
'' TIME.FLAG = 1 MEANS CPU.R EQUAL TO XX.CPU.SERVICE.TIME
'' 1 MEANS THE CPU TASK HAS BEEN COMPLETED
'' 0 MEANS THAT THE CPU TASK NEEDS MORE TIME
''
'' IO.SERVICE.TIME IS THE IO.TIME REQUIRED TO PROCESS A SINGLE
'' I/O TASK
''
'' XX.CPU.SERVICE.TIME IS THE TIME NEEDED UNTIL NEXT I/O
''
'' CPU.R IS THE TIME TO RELEASE THE CPU
''

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014330
 014340
 014350
 014360
 014370
 014380
 014390
 014400
 014410
 014420
 014430
 014440
 014450
 014460
 014470
 014480
 014490
 014500
 014510
 014520
 014530
 014540
 014550
 014560
 014570
 014580
 014590
 014600
 014610
 014620
 014630
 014640
 014650
 014660
 014670
 014680

```

** LOAD.F = 0 MEANS THE JOB DOES NOT NEED LOADING INTO MEMORY *
** 1 MEANS JOB MUST BE LOADED INTO MEMORY *
** TIME.E.CPU.QUEUE IS THE TIME THE JOB MOST RECENTLY ENTERED *
** THE CPU QUEUE *
** JOB.TYPE 1=RATCH, 2 = INTERACTIVE *
*****
DEFINE TIME.FLAG AS INTEGFR VARIABLE *****
DEFINE IO.OVERHEAD.T AS REAL VAFIABLE *****
LET IO.OVERHEAD.T = (1.0/100.0)*SEC.DAYS *****
FOR EACH AMIS IN ACPU, WITH *****
F.LOADER(AMIS) EQ 1 *****
FINO THE FIRST CASE *****
IF NONE GO TO 'A' *****
ELSE *****
REMOVE AMIS FROM ACPU *****
FILE AMIS IN ACPU *****
SCHEDULE A RELEASE.CPU IN (5.0/1000.0)*SEC.DAYS *****
RETURN *****
'A' *****
REMOVE FIRST AMIS FROM ACPU *****
IF FLAG = 1 *****
LIST TLF(AMIS) *****
LIST TIME.V *****
LIST PRIORITY(AMIS) *****
LIST JORNUM(AMIS) *****
LIST STATUS.PRIORITY *****
PRINT 1 LINE THUS *****
EVENT CPU.FPROCESSING *****

ELSE *****
ALWAYS *****
IF TLF(AMIS) EQ 1 *****
LET XX.CPU.SERVICE.TIME(AMIS) = T.LEFT.IN.CPU(AMIS) *****
ELSE *****
ALWAYS *****
IF JOB.TYPE(AMIS) = 1 *****
LET CP.T.QUANTUM=4000.0*SEC.DAYS *****

```



```

ELSE
  REMOVE AMIS FROM ACPU
  FILE AMIS IN LOAD.QUEUE
  SCHEDULE A LOAD.DELAY IN (5.0/1000.0)*SEC.DAYS DAYS
  RETURN
*8*
LET CPU.SF=0
REMOVE FIRST AMIS FROM ACPU
  LET JOB.IO.REQUESTS(AMIS) = JOB.IO.REQUESTS(AMIS) -1
IF FLAG = 1
  PRINT 1 LINE THUS
  RELEASE CPU SECTION A
  LIST JOBNUM(AMIS)
  LIST TIME.V
  LIST YLF(AMIS)
  LIST PRIORITY(AMIS)
  LIST STATUS.PRIORITY
  PRINT 1 LINE THUS
  EVENT RELEASE.CPU
ELSE
  ALWAYS
  IF FF=0
  PRINT 1 LINE THUS
  STATS IN RELEASE CPU
  PRINT 1 LINE WITH TIME.V THUS
  PRINT 1 LINE WITH JOBNUM(AMIS), JOB.TYPE(AMIS) AND
  XX.CPU.SERVICE.TIME(AMIS) THUS
  *** * .*****
  PRINT 1 LINE WITH T.CPU.TIME THUS
  * .*****
  ELSE
  ALWAYS
  IF WARM.UP.FLAG GE 1
  LET GAIN=GAIN + (XX.CPU.SERVICE.TIME(AMIS) -IO.OVERHEAD.T)
  LET T.CPU.TIME=T.CPU.TIME+(XX.CPU.SERVICE.TIME(AMIS) -IO.OVERHEAD.T)
  ELSE

```

LET TCPU = TCPU + (XX.CPU.SERVICE.TIME(AMIS) - IO.OVERHEAD.T)
ALWAYS

IF TIME.V LT .00232

LET FF=0

LET FF=1

ELSE

LET FF=0

LET FF=1

ALWAYS

IF JOB.TYPE(AMIS) EQ 1

GO TO 'E'

ELSE

IF FLAG = 1

PRINT 1 LINE THUS

RELEASE CPU SECTION B

LIST JOBNUM(AMIS)

LIST NO.CAR.RETURNS(AMIS)

LIST JOB.IO.REQUESTS(AMIS)

LIST XX.CPU.SERVICE.TIME(AMIS)

LIST STATUS.PRIORITY

LIST PRIORITY(AMIS)

LIST TLF(AMIS)

LIST JOB..IO.REQUESTS(AMIS)

ELSE

ALWAYS

IF FLAG = 1

IF NO.CAR.RETURNS(AMIS) LE 0

LIST JOBNUM(AMIS)

LIST NO.CAR.RETURNS(AMIS)

LIST JOB.IO.REQUESTS(AMIS)

LIST JOB..IO.REQUESTS(AMIS)

LIST XX.CPU.SERVICE.TIME(AMIS)

LIST XX.CPU.SERVICE.TIME(AMIS)

ELSE

ALWAYS

FLSF

ALWAYS

015930
015940
015950
015960
015970
015980
015990
016000
016010
016020
016030
016040
016050
016060
016070
016080
016090
016100
016110
016120
016130
016140
016150
016160
016170
016180
016190
016200
016210
016220
016230
016240
016250
016260
016270
016280
016290

```

IF NO.CAR.RETURNS(AMIS) LE 0
GO TO 'EE'
ELSE
  IF JOB.IO.REQUESTS(AMIS) LE 0
  LET NO.CAR.RETURNS(AMIS) = NO.CAR.RETURNS(AMIS) -1
  LET JOB.IO.REQUESTS(AMIS)=JOB.IO.REQUESTS(AMIS)
  LET XX.CPU.SERVICE.TIME(AMIS) = XX.CPU.SERVICE.TIME(AMIS)
  LET THINK.TIME = RANDI.F(2,15,3)
  LET SPECIAL.CASE = RANDI.F(1,20,3)
  IF SPECIAL.CASE GT 12
  LET THINK.TIME = 240
  ELSE
  ALWAYS
  IF HARM.UP.FLAG GF 1
  LET USER.RESPONSE.TIME=THINK.TIME
  ELSE
  ALWAYS
  CREATE A RESPONSE
  LET RNAME(RESPONSE) = AMIS
  LET STATUS.PRIORITY=0
  SCHEDULE THE RESPONSE IN (THINK.TIME*SECOND) MINUTES
  SCHEDULE A CPU.REQUEST NOW
  RETURN
  ELSE
  IF TLF(AMIS) EQ 0
  GO TO 'REESTABLISH'
  ELSE
  GO TO 'FILE'
  'E'
  IF FLAG = 1
  LIST JOB.IO.REQUESTS(AMIS)
  LIST TLF(AMIS)
  LIST CPU.SF
  ELSE
  ALWAYS
  IF JOB.IO.REQUESTS(AMIS) EQ 0
  IF PRIORITY (AMIS) LE 15
  LET PFLAG =0
  LET PRIORITY(AMIS) = 5
  ELSE
  ALWAYS

```

```

016300
016310
016320
016330
016340
016350
016360
016370
016380
016390
016400
016410
016420
016430
016440
016450
016460
016470
016480
016490
016500
016510
016520
016530
016540
016550
016560
016570
016580
016590
016500
016510
016520
016630
016640
016650
016660
016670
016680
016690
016700

```

```

LET STATUS.PRIORITY=0
FILE AMIS IN XX.RUFFER.QUEUE
SCHEDULE A CPU.REQUEST IN 5.0*MIL.SEC.DAYS DAYS
SCHEDULE A IS.RUFFER.EMPTY NOW
RETURN
ELSE
IF TLF(AMIS) EQ 0
IF PRIORITY (AMIS) LE 15
LET PFLAG =0
LET PRIORITY(AMIS) = 5
ELSE
ALWAYS
LET XX.CPU.SERVICE.TIME(AMIS) = XX.CPU.SERVICE.TIME(AMIS)
CREATE A IO.DELAY
LET NAME(IO.DELAY) = AMIS
LET STATUS.PRIORITY=0
IF FLAG = 1
LIST IO.SERVICE.TIME(AMIS)
ELSE
ALWAYS
LET GAIN=GAIN + IO.SERVICE.TIME(AMIS)
SCHEDULE THE IO.DELAY IN IO.SERVICE.TIME(AMIS) DAYS
SCHEDULE A CPU.REQUEST NOW
RETURN
ELSE
IF PRIORITY (AMIS) LE 15
LET PFLAG =0
LET PRIORITY(AMIS) = 5
ELSE
ALWAYS
FILE AMIS IN CPU.REQUEST.QUEUE
LET STATUS.PRIORITY=0
SCHEDULE A CPU.REQUEST NOW
RETURN
'EF'
LET STATUS.PRIORITY=0
FILE AMIS IN XX.RUFFER.QUEUE

```

```

016710
016720
016730
016740
016750
016760
016770
016780
016790
016800
016810
016820
016830
016840
016850
016860
016870
016880
016890
016900
016910
016920
016930
016940
016950
016960
016970
016980
016990
017000
017010
017020
017030
017040
017050
017060
017070

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017080
017090
017100
017110
017120
017130
017140
017150
017160
017170
017180
017190
017200
017210
017220
017230
017240
017250
017250
017270
017280
017290
017300
017310
017400
017410
017420
017430
017440
017450
017460
017470
017480
017490
017500
017510

SCHEDULE A CPU.REQUEST IN 5.0*MIL.SEC.DAYS DAYS
SCHEDULE A IS.BUFFER.EMPTY NOW
RETURN
'REESTABLISH'
LET XX.CPU.SERVICE.TIME(AMIS) = XX..CPU.SERVICE.TIME(AMIS)
CREATE A IO.DELAY
LET NAME(IO.DELAY) = AMIS
LET STATUS.PRIORITY=0
IF FLAG = 1
    LIST JOBNUM(AMIS)
    LIST IO.SERVICE.TIME(AMIS)
ELSE
ALWAYS
LET GAIN=GAIN + IO.SERVICE.TIME(AMIS)
SCHEDULE THE IO.DELAY IN IO.SERVICE.TIME(AMIS) DAYS
SCHEDULE A CPU.REQUEST NOW
RETURN
'FILE'
LET STATUS.PRIORITY=0
FILE AMIS IN CPU..REQUEST.QUEUE
SCHEDULE A CPU.REQUEST NOW
RETURN
END

EVENT LOAD.DELAY
REMOVE FIRST AMIS FROM LOAD.QUEUE
LET STATUS.PRIORITY=0
LET XXTIME.E.CPU.QUEUE(AMIS)=TIME.V
FILE AMIS IN CPU..REQUEST.QUEUE
LET CPU.SF = 0
LET LOAD.F = 0
LET F.LOADER(AMIS) = 0
SCHEDULE A CPU.REQUEST NOW
RETURN
END

```

```

EVENT IO.DELAY SAVING THE EVENT NOTICE
LET AMIS= NAME(IO.DELAY)
DESTROY IO.DELAY
LET XXTIME.E.CPU.QUEUE(AMIS) = TIME.V
FILE AMIS IN CPU..REQUEST.QUEUE
SCHEDULE A CPU..REQUEST NOW
RETURN
END
EVENT RESPONSE SAVING THE EVENT NOTICE
*****
*****
*****
*****
LET AMIS = RNAME(RESPONSE)
DESTROY RESPONSE
LET XXTIME.E.CPU.QUEUE(AMIS)=TIME.V
FILE AMIS IN CPU..REQUEST.QUEUE
SCHEDULE A CPU..REQUEST NOW
RETURN
END
EVENT IS.BUFFER.EMPTY
*****
**      B.SF IS THE BUFFER STATUS FLAG
**      0 MEANS BUFFER EMPTY
**      1 MEANS BUFFER NOT EMPTY
*****
CALL BUFFER YIELDING B.SF
REMOVE FIRST AMIS FROM XX.BUFFER.QUEUE
IF B.SF = 0
FILE AMIS IN ACM.RELEASE.QUEUE
SCHEDULE A CM.RELEASE NOW
ELSE
FILE AMIS IN XX.BUFFER.QUEUE
SCHEDULE A IS.BUFFER.EMPTY IN (1.0/1000.0)*SEC.DAYS DAYS
ALWAYS
RETURN
END

```

```

017590
017600
017610
017620
017630
017640
017650
017650
017770
017790
017790
017900
017810
017820
017830
017840
017850
017850
017870
017950
017950
017970
017980
017990
018000
018010
018020
018030
018040
018050
018060
018070
018080
018090
018100
018110
018120

```

018200
 018210
 018220
 018230
 018240
 018250
 018260
 018270
 018280
 018290
 018300
 018310
 018320
 018330
 018340
 018350
 018450
 018460
 018470
 018480
 018490
 018500
 018510
 018520
 018530
 018540
 018550
 018560
 018570
 018580
 018590
 018600
 018610
 018620

```

EVENT CM.RELEASE
*****
** START MAIN.MEMORY IS THE BEGINNING CORE LOCATION *
** FINISH.MAIN.MEMORY IS THE LAST CORE LOCATION *
** MAIN.MEMORY IS THE MAIN MEMORY ARRAY *
*****
REMOVE FIRST AMIS FROM ACM.PELEASE.QUEUE
LET SMM = START.MEMORY(AMIS)
LET FMM = END.MEMORY(AMIS)
FOR I = SMM TO FMM, 10
LET MAIN.MEMORY(I) = 0
LOOP
FILE AMIS IN AN.OUTPUT.QUEUE
SCHEDULE A OUTPUT NOW
RETURN
END

EVENT OUTPUT
*****
** CHANNEL = 0 MEANS CHANNEL AVAILABLE *
** CHANNEL = 1 MEANS CHANNEL BUSY *
*****
REMOVE FIRST AMIS FROM AN.OUTPUT.QUEUE
CALL IO GIVEN AMIS YIELDING CHANNEL
IF CHANNEL = 1
FILE AMIS IN HH.OUTPUT.QUEUE
SCHEDULE A XX.OUTPUT IN 1.0 * MIL.SEC.DAYS DAYS
ELSE
FILE AMIS IN AN.OUTPUT.QUEUE
SCHEDULE A AN.DISK.RELEASE NOW
ALWAYS
RETURN
END
  
```


019250
 019260
 019270
 019280
 019290
 019300
 019310
 019320
 019330
 019340
 019350
 019360
 019370
 019380
 019390
 019400
 019410
 019420
 019430
 019440
 019450
 019460
 019470
 019480
 019490
 019500
 019510
 019520
 019530
 019540
 019550
 019560
 019570
 019580
 019590
 019600

```

*****
**   DISK CONTROL ROUTINE WILL RELEASE THE APPROPRIATE DISKS   **
**   NSUI IS NUMBER OF SIMULTANEOUS INTERACTIVE USERS         **
**   NSUR IS NUMBER OF SIMULTANEOUS BATCH USEPS               **
*****
REMOVE FIRST AMIS FROM AN.OUTPUT.QUEUE
LET JOB.EXIT.TIME(AMIS) = TIME.V
LET TURNAROUND(AMIS) = TIME.V-APRIVAL.TIME(AMIS)
LET STATUS.PRIORITY = 0
IF JOB.TYPE(AMIS) = 1
PRINT 1 LINE WITH JOBNUM(AMIS) AND TIME.V THUS
RATCH JOB# ***. MADE IT OUT AT TIME * *****
  LET NRJ = NBJ-1
  IF WARM.UP.FLAG=0
  LET CU=(TCPUT/TIME.V)*100.0
  LIST CII
  GO TO 'WARMING.UP.PERIOD'
ELSE
  LET B.TURNAROUND.T = TURNAROUND(AMIS)
  LET B.CAPABILITY=TURNAROUND(AMIS)/CP.TIME(AMIS)
  LET V = TIME.V-ARRIVAL.TIME(AMIS)
  LET W=(CP.TIME(AMIS)+IO.TIME(AMIS))
  LET Y=CP.TIME(AMIS)
  LET ETMF=V/W
  LET EX.DELAY.FACTOR=V/Y
ELSE
PRINT 1 LINE WITH JOBNUM(AMIS) AND TIME.V THUS
INTERACTIVE JOB NUM ***. MADE IT OUT AT TIME * *****
  LET NIJ = NIJ-1
  IF WARM.UP.FLAG=0
  LET CU=(TCPUT/TIME.V)*100.0
  LIST CII
  GO TO 'WARMING.UP.PERIOD'
ELSE
  LET I.TURNAROUND.T = TURNAROUND(AMIS)
  LET I.CAPABILITY= TURNAROUND(AMIS)/CP.TIME(AMIS)

```

019610
 019620
 019630
 019640
 019650
 019660
 019670
 019680
 019690
 019700
 019710
 019720
 019730
 019740
 019750
 019760
 019770
 019780
 019790
 019800
 019810
 019820
 019830
 019840
 019850
 019950
 019960
 019970
 019980
 019990
 020000
 020010
 020020
 020030
 020040
 020050
 020060
 020070
 020080

```

LET V = TIME.V-ARRIVAL.TIME(AMIS)
LET W=(CP.TIME(AMIS)+IO.TIME(AMIS))
LET Y=CP.TIME(AMIS)
LET I.ETMF = V/W
LET I.EX.DELAY.FACTOR=V/Y
ALWAYS
  
```

```

IF FLAG=1
LIST STATUS.PRIORITY
LIST TIME.V
LIST PRIORITY(AMIS)
ELSE
  ALWAYS
  
```

```

LET N.JOBS.PROCESSED = N.JOBS.PROCESSED+1
LET GAIN.FACTOR=GAIN/(TIME.V-W.UP.TIME)
LET CURRENT.UTILIZATION =(T.CFU.TIME/(TIME.V-W.UP.TIME))*100.0
LIST CURRENT.UTILIZATION
*WARMING.UP.PERIOD*
DESTROY THE JOB CALLED AMIS
IF JN GT 200
  SCHEDULE A CEASE.SIM NOW
  ELSE
  ALWAYS
  RETURN
  END
  
```

```

ROUTINE BUFFER YIELDING B.SF
*****
'' B.SF IS THE BUFFER STATUS FLAG
'' 0 MEANS BUFFER EMPTY
'' 1 MEANS BUFFER NOT EMPTY
*****
LET B.SF=RANDI.F(1,10,4)
IF B.SF LT 4
  LET B.SF=1
  ELSE
  LET B.SF=0
  ALWAYS
  RETURN
  END
  
```

020190
 020200
 020210
 020220
 020230
 020240
 020250
 020260
 020270
 020280
 020290
 020300
 020310
 020320
 020330
 020340
 020350
 020360
 020370
 020380
 020390
 020400
 020410
 020420
 020430
 020440
 020450
 020460
 020470
 020480
 020490
 020500
 020510
 020520
 020530

```

EVENT CEASE.SIM
PRINT 1 LINE THUS
ENTERED CEASE.SIM
PRINT 1 LINE WITH TIME.V THUS
SIMULATION TIME IS *****
'RECHECK'
FOR EACH AMIS IN CPU..REQUEST.QUEUE, WITH
JOB.TYPE(AMIS) GT 0,
FIND THE FIRST CASE
IF NONE GO TO 'CONTINUE'
ELSE
REMOVE AMIS FROM CPU..REQUEST.QUEUE
IF JOB.TYPE(AMIS) EQ 1
LET B.BT=BLOCKED.TIME(AMIS)
ELSE
LET I.BT=BLOCKED.TIME(AMIS)
ALWAYS
DESTROY THE JOB CALLED AMIS
GO TO 'RECHECK'
'CONTINUE'
LIST JN
LIST NJNPLO
LIST NIJ
LIST NRJ
PRINT 3 LINES THUS
INTERACTIVE STATISTICS

PRINT 1 LINE THUS
NUMBER OF INTERACTIVE USERS
PRINT 1 LINE THUS
MEAN VALUE
PRINT 2 LINES WITH AMSUI, ANSUI, AND CNSUI THUS
***.
MAXIMUM NUMBER
MINIMUM NUMBER
***.
  
```

```

020540
020550
020560
020570
020580
020590
020600
020610
020620
020630
020640
020650
020660
020670
020680
020690
020700
020710
020720
020730
020740
020750
020760
020770
020780
020790
020800
020810
020820
020830
020840
020850
020860
020870
020880
020890

PRINT 2 LINES THUS
INTERACTIVE TURNAROUND TIME
MEAN
PRINT 2 LINES WITH AI.TURNAROUND.T, BI.TURNAROUND.T, AND CI.TURNAROUND.T THUS
*.*****
          STD.DEV
          *.*.*.*.*

PRINT 2 LINES THUS
INTERACTIVE CPU TIME
MEAN
PRINT 2 LINES WITH AI.CPU.TIME, BI.CPU.TIME, AND CI.CPU.TIME THUS
*.*****
          STD.DEV
          *.*.*.*.*

PRINT 2 LINES THUS
INTERACTIVE CAPABILITY
MEAN
PRINT 2 LINES WITH AA, BB, CC, AND DD THUS
*.*****
          VARIANCE
          *.*.*.*.*
          STD.DEV
          *.*.*.*.*

PRINT 2 LINES THUS
INTERACTIVE CPU BLOCKED TIME
MEAN
PRINT 2 LINES WITH AAA, BBB, CCC, AND DDD THUS
*.*****
          VARIANCE
          *.*.*.*.*
          STD.DEV
          *.*.*.*.*

PRINT 2 LINES THUS
INTERACTIVE MEMORY REQUESTS IN K
MEAN
PRINT 2 LINES WITH II, JJ, KK, AND LL THUS
*.*****
          VARIANCE
          *.*.*.*.*
          STD.DEV
          *.*.*.*.*

PRINT 2 LINES THUS
INTERACTIVE USER RESPONSE TIME
MEAN
PRINT 2 LINES WITH OO, RR, SS, AND TT THUS
*.*****
          VARIANCE
          *.*.*.*.*
          STD.DEV
          *.*.*.*.*

```


VITA

Alfred H. Linder, III was born on 14 June 1948 in San Diego, California. He graduated from Flagstaff High School in Flagstaff, Arizona in 1966 and enlisted in the United States Navy in June of 1966. After three years in the Navy, as an electronic technician, he joined the Navy Reserve while working for the Sante Fe Railroad and then the Mountain Bell Telephone Company. He attended Northern Arizona University where he received, in 1976, a Bachelor of Science in Mathematics and a commission in the United States Air Force. He entered the AFIT residence school in August 1976 and received the degree of Master of Science in Computer Science in Dec 1978.

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