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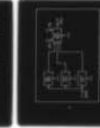
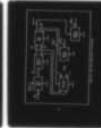
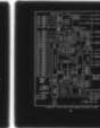
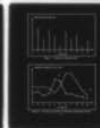
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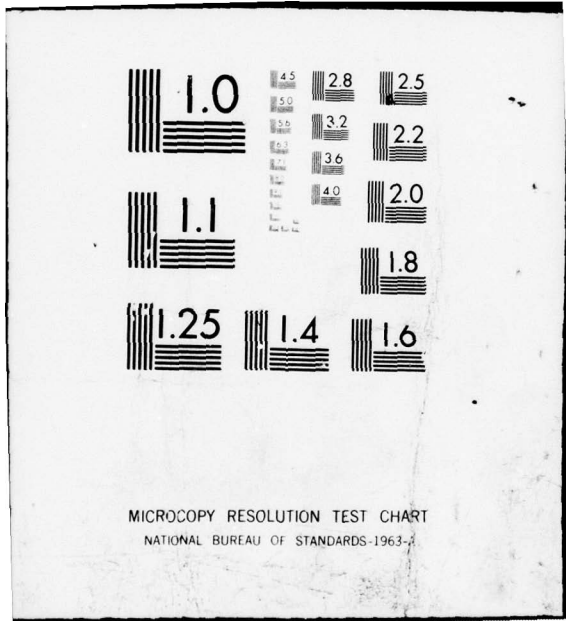
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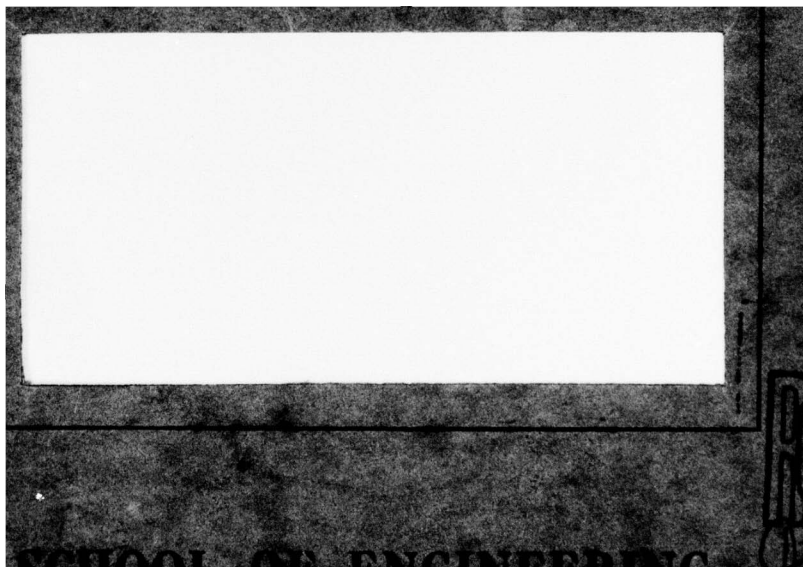
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MICROPROCESSOR BASED DATA ACQUISITION  
AND PROCESSING SYSTEM

THESIS

AFIT/GE/EE/78-29

Saleem Iftikhar  
Flt/Lt PAF

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## Preface

This report is the result of my attempt to develop a microprocessor based data acquisition and processing system for the Air Force Flight Dynamics Laboratory, Wright-Patterson Air Force Base, Ohio. The report describes the three phases of system development: conceptual phase, design phase, and implementation and test phase. Structured analysis and design technique was used in the conceptual and design phases. In my opinion, this graphical technique makes it possible to follow a top-down design approach. A bench-type model was the result of the implementation and test phase.

I wish to express my appreciation to Dr. Gary B. Lamont, my thesis advisor, for his many hours of help and advise. I greatly appreciate his detailed reading of the report to bring to light errors which I had made and his suggestions for improvement. I would like to thank the committee members, Capt Gregg L. Vaughn and Dr. Thomas C. Hartrum for their valuable guidance during the course of the project.

Also, I would like to thank Dr. Lynn C. Rogers and Mr. Robert W. Gordon of the Air Force Flight Dynamics Laboratory who provided the thesis topic. Their interest in my work and assistance during the project were invaluable.

I would like to mention my appreciation to Lieutenant Peter D. Summers, III, and Mr. Frank E. Beitel of the Air Force Materials Laboratory, Computer Activities Division, Wright-Patterson Air Force Base, Ohio, for extending their help during the software development.

The manuscript in its present form would not be possible without the efforts of my typist, Miss Joyce L. Wilson.

Finally, I wish to acknowledge my indebtedness to my wife and baby for their understanding, patience, and encouragement during the course of this project.

Saleem Iftkhar

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Abstract

↘ The Air Force Flight Dynamics Laboratory at Wright-Patterson Air Force Base, Ohio, plans to design a constrained layer damping treatment for aircraft components. This design requires the knowledge of vibration and temperature variations that a component encounters in service.

The development of a microprocessor based data acquisition and processing system (MIBDAPS) to monitor the vibration and temperature variations of a component under test is discussed in this paper. The designed system is capable of acquiring and processing data in real time and generating cumulative damage factor versus frequency and temperature contours. These contours are used in the design of the damping treatment.

The system uses an LSI-11M microcomputer as its central processor. The support hardware for this microcomputer includes an analog-to-digital converter, core memory, real time clock, and serial line interface module. The software developed for this system handles the acquisition and processing of data. The real time processing includes the computation of the FFT and power spectrum of the vibration signal and subsequent calculation of the cumulative damage factor versus frequency and temperature contours. The bench model of the designed system was implemented and tested for functional performance.

↙

## I. Introduction

The Air Force Flight Dynamics Laboratory (AFFD), Wright-Patterson Air Force Base, Ohio, plans to develop a constrained layer damping treatment for aircraft components such as the airframe, avionics hardware, and other accessories. The specific damping layer treatment design requires knowledge of temperature and vibration variations encountered by the component. To generate this information, a data acquisition and processing system is required to serve as an in-flight monitor for the component under test. This report describes the design, development, and bench model implementation of such a system.

### Background

Research has established that high-cycle fatigue damage due to resonant vibration is one of the major reasons for structural failure of components (Ref 13). A constrained layer damping treatment would be cost effective in extending the life of the component. In order to determine the correct type of damping layer, the temperatures, vibration nodes, and frequencies at which damage occurs in service must be known.

It has been determined that there are approximately ten vibration frequency bands of interest (Ref 13). They lie in the range from 100 HZ to 1000 HZ and are spaced 100 HZ apart (Figure 1). A damage factor (DF) at each frequency of interest can be calculated using the following relationship (Ref 13):

$$DF = \frac{K(\bar{A})^\beta}{(f)^{2\beta-1}} \quad (1)$$

Here,  $\beta$  is approximately equal to 3.2, and  $\bar{A}$  is the root-mean-square (RMS) value of acceleration in the frequency band of interest.

A summation of the DF over a predetermined period of time would yield a cumulative damage factor (CDF) and curves as shown in Figure 2 could be drawn.

Existing in-flight data acquisition techniques employ an analog signal conditioner and a multichannel analog tape recorder (Ref 13). The analog signals from the vibration and temperature sensors are conditioned and recorded on a magnetic tape. This system is bulky and takes excessive setup time. Skilled personnel are required to install the equipment on an aircraft, and calibration is difficult and time-consuming.

To simplify this process, the design of a self-contained Microprocessor Based Data and Processing System (MIBDAPS) appeared as an attractive alternative. Such a system should be easy to install and remove, and only minimum operator intervention during its operation should be necessary.

#### System Concept

The underlying concept of MIBDAPS was a self-contained system that would monitor the analog signals from the temperature and vibration sensors, digitize these signals, process the data in real time, and store the results on an appropriate media. The real time processing would include computation of the power spectrum of the vibration signal and calculation of DF at each frequency of interest. The results of processing would be stored in a CDF versus frequency and temperature array, respectively. Therefore, at the termination of system operation, the CDF versus frequency and temperature contours (Figure 2) should be available.

MIBDAPS should be capable of either 15 to 20 hours of continuous operation (for gathering data on transport aircraft hardware) or time

segmented operation (for gathering data on fighter aircraft hardware which has limited flying endurance). During a time segmented operation, the system should be able to operate for a total aggregate of 15 to 20 hours. The data processed during this time interval would enable determination of vibration nodes at which the damage occurs.

### Approach

The development of MIBDAPS was carried out in three phases: the conceptual phase, the design phase, and the implementation and testing phase. During the conceptual phase, the detailed requirements placed on the system were defined using a structured analysis and design technique (SADT) (Ref 1). The design phase included the software/hardware partition of the system, selection of hardware components, and design of software and hardware modules to implement different functions of MIBDAPS. During the implementation and testing phase, the software and hardware modules were given a functional check. The system integration and performance check of the integrated system was also completed during this phase.

### The Development Environment

The LSI-11M\* microcomputer was used as the system processor because of its powerful instruction set, speed, and availability of support hardware and software. The use of the LSI-11 microcomputer system for MIBDAPS determined the general development requirements for both hardware and software modules of MIBDAPS. The off-the-shelf hardware modules used were analog-to-digital converter (ADC), real time clock (RTC), and the 4K by 16-bits core memory. All of these modules are manufactured by the Digital Equipment Corporation (DEC).

\*LSI-11M is the militarized version that is manufactured by Norden Systems, Inc.

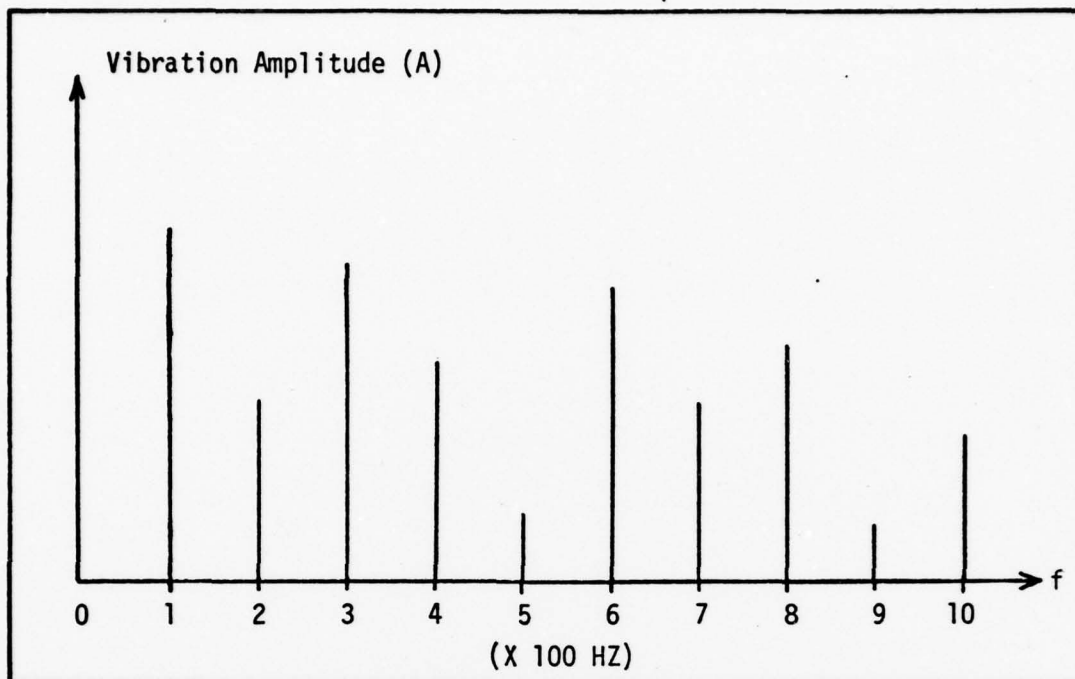


Figure 1. Vibration Frequency Bands

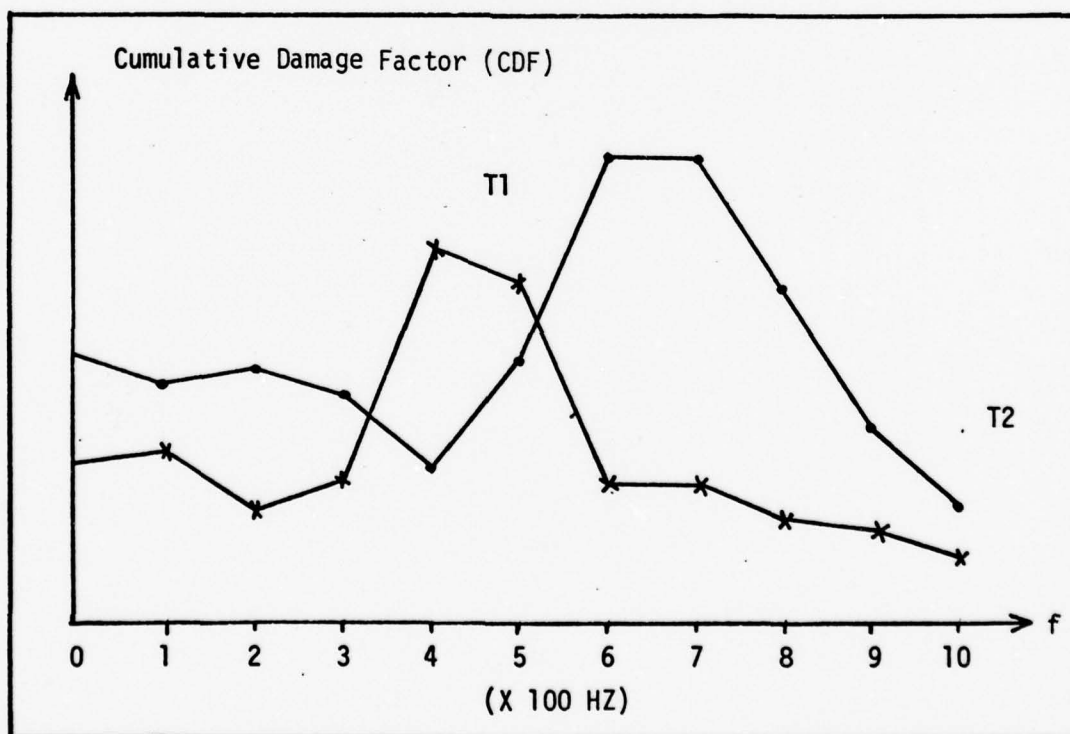


Figure 2. CDF Versus Frequency at Different Temperature Values

The software development system used was a PDP-11/03 minicomputer which has the LSI-11 as the processor. The RXV011 floppy disk based operating system was available on this minicomputer. The development system was located at the Air Force Materials Laboratory (AFML), Wright-Patterson Air Force Base, Ohio.

#### Scope of Thesis

The structured analysis and design technique was used to develop the activity model of MIBDAPS (Ref Chap II), and the LSI-11M microcomputer was used as the processor for MIBDAPS (Ref Chap III). Hardware and software modules were designed and tested (Ref Chaps IV and V). The hardware/software modules were then integrated into a bench model of MIBDAPS which was then functionally tested (Ref Chap VI). The bench model of MIBDAPS performed the required functions of data acquisition and processing. Testing of the integrated system for computational accuracy has not been completed at this time. Recommendations for further work have been generated in order to complete the implementation of the prototype module (Ref Chap VII).

## II. System Requirements for MIBDAPS

### Introduction

The conceptual idea is not enough to form the basis of design and implementation of any complex system. The design can be facilitated if the complete system is divided into functional subunits. These subunits are easier to analyze and design. The division of a system into functional subunits according to the requirements placed on the system can be called the "requirement definition" phase of the design. A carefully prepared requirement definition ensures that ambiguities in the design and implementation phase are minimized.

### Analysis and Design Methodology

Structured Analysis and Design Techniques (SADT) is a design methodology for performing the functional analysis and design of complex systems (Ref 1). A functional model of the system is created first which helps in understanding what activities the system will perform, followed by a design model to show how the system will be implemented to perform the requisite functions (Ref 1). The functional model also helps to clarify the requirement placed on the system.

Appendix A describes SADT methodology and its application rules. Some changes in convention are made to suit the design of MIBDAPS; these are also included in Appendix A.

The SADT was chosen as a functional analysis and design tool because it is a graphical method of analyzing the functional requirements of a system. The activity model, which shows different functions which the system is required to perform, is independent of the hardware/

software implementation of the system. Thus, any changes in the activity model can be made during this phase of the design without affecting the final shape of the product.

Once the activity model is complete, it is easy to review and correct any oversights that might have occurred during the functional analysis of the system.

### Activity Model

Before the activity model is drawn, the system is divided into different nodes. The node index given in Table I shows the hierarchy of the system from activity viewpoint. This node index is based on the initial problem definition by the Air Force Avionics Laboratory (AFAL), Wright-Patterson Air Force Base, Ohio, and subsequent discussions with the sponsor of this thesis. The node index also indicates how the author visualized the system was required to function.

TABLE I  
Node Index

A-1	Data Collection and Analysis
A-0	Acquire and Process Data
A0	Acquire and Process Data
A1	Control All Activity
A2	Acquire Data
A3	Process Data
A4	Store Data

The explanation of each node follows.

Node A-1, Data Collection and Analysis (Figure 3). This node shows the overall research activity and the place where MIBDAPS fits in the overall function of the organization. Block 3, Plot Data, activity will be carried out in the laboratories.

Node A-0, Acquire and Process Data (Figure 4). Node A-0 is the model of the MIBDAPS. The node shows the system requirements. Data input comes in the form of analog signals— analog vibration signal (I1) and analog temperature signal (I2). The output O3 of the system is a three-dimensional array. This array contains the cumulative damage factor (CDF) as a function of frequency and temperature. Output O2 is the scale factor for the array elements, and output O1 is the total operational time for which the system was operational.

The MIBDAPS will be able to continuously acquire and process data for a total flying time of 15 to 20 hours. This is the requirement placed by AFAL because the CDF versus temperature and frequency array for this duration will be a good estimate of the fatigue damage profile of the test piece (Ref 13). Input signal specifications are:

Analog Vibration Signal (I1)	0-5V
Analog Temperature Signal (I2)	0-5V

Node A0, Acquire and Process Data (Figure 5). The MIBDAPS is expected to acquire data, process data, and store the results. Therefore, node A0 shows the decomposition of the system into four primary activities:

- Acquire Data
- Process Data
- Store Data
- Control All Processes

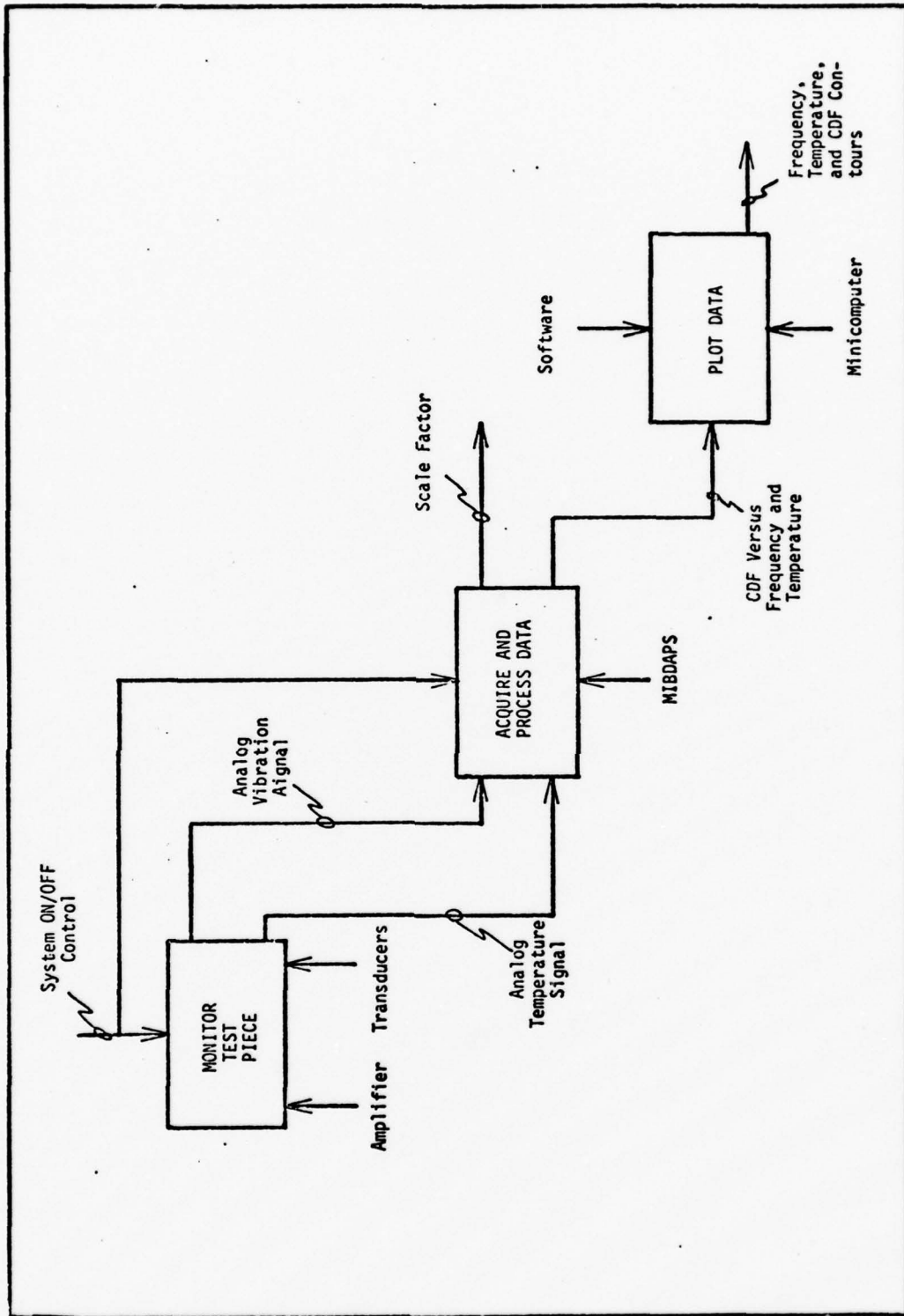


Figure 3. Node A-1, Data Collection and Analysis

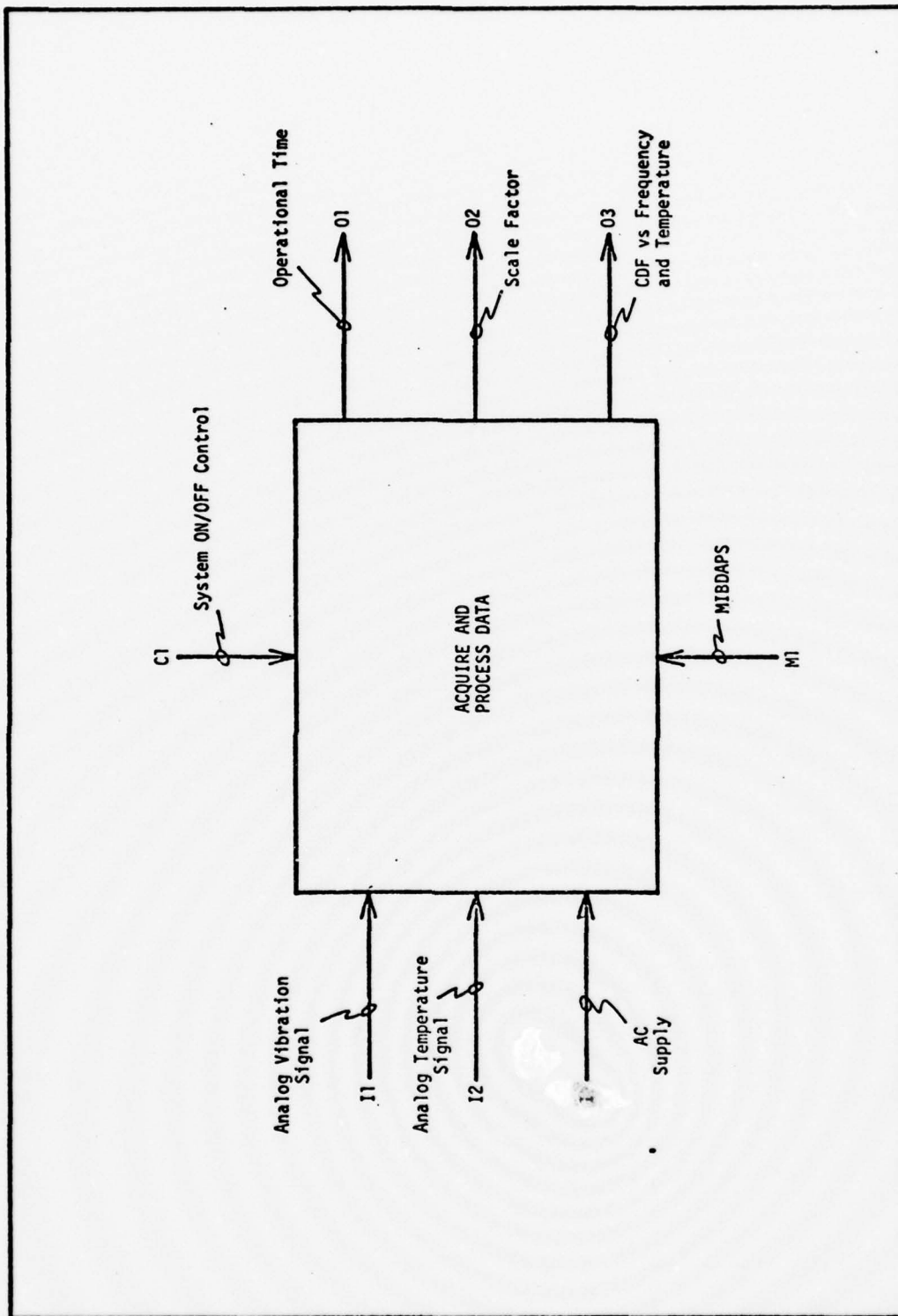


Figure 4. Node A-0, Acquire and Process Data

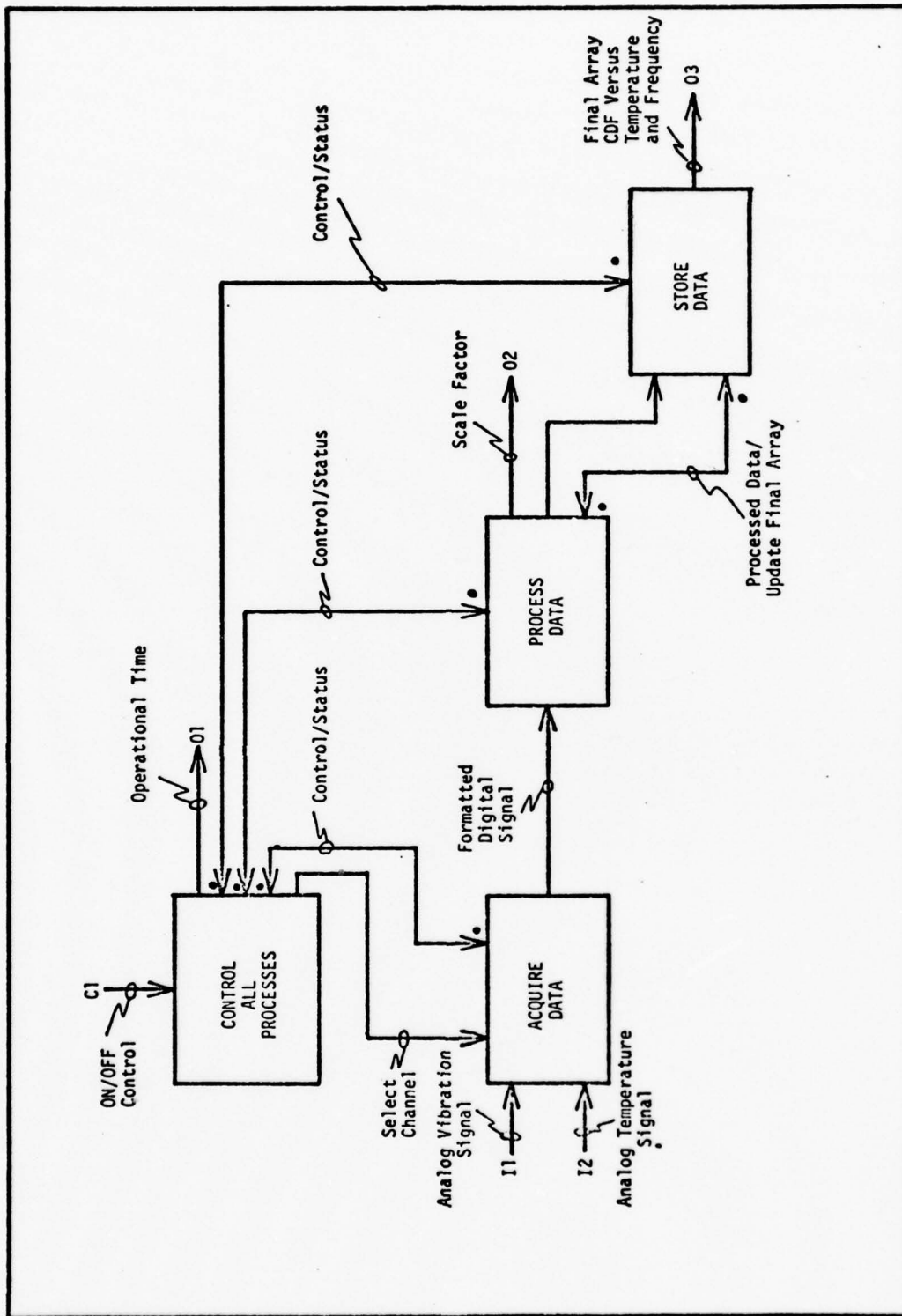


Figure 5. Node A0, Acquire and Process Data

An explanation of how a signal travels through each subnode is now given. The activity Acquire Data continuously digitizes the input signal under the control from node A1. The digitized signal is queued and sent to the processing node (A2).

The power spectrum of the signal is computed by the processing node. For each frequency of interest, a damage factor is calculated (Ref Chap I) and stored in an array. The activity Store Data adds the most recent set of DF values to the previously stored values. This results in a storage of cumulative damage factors.

Node A01, Control All Processes, keeps all events in the system synchronized. It initiates appropriate commands to ensure that data is processed and stored in proper order. The input and output signal specifications are given below:

- Input Signal

- a. Analog Vibration Signal (I1): 0-5V.
- b. Analog Temperature Signal (I2): 0-5V.
- c. AC Supply (I3): 110V, 60 HZ.

- Output Signal

- a. Operational Time (O1): Displayed on the front panel.
- b. Scale Factor (O2).
- c. CDF Versus Temperature and Frequency Array (O3).

Node A1, Control All Activity (Figure 6). Node A1 subdivides the overall control of the system into the following functions:

Control Acquisition of Data  
Control Processing of Data  
Control Storage of Data  
Control Power to the System  
Record Operational Time

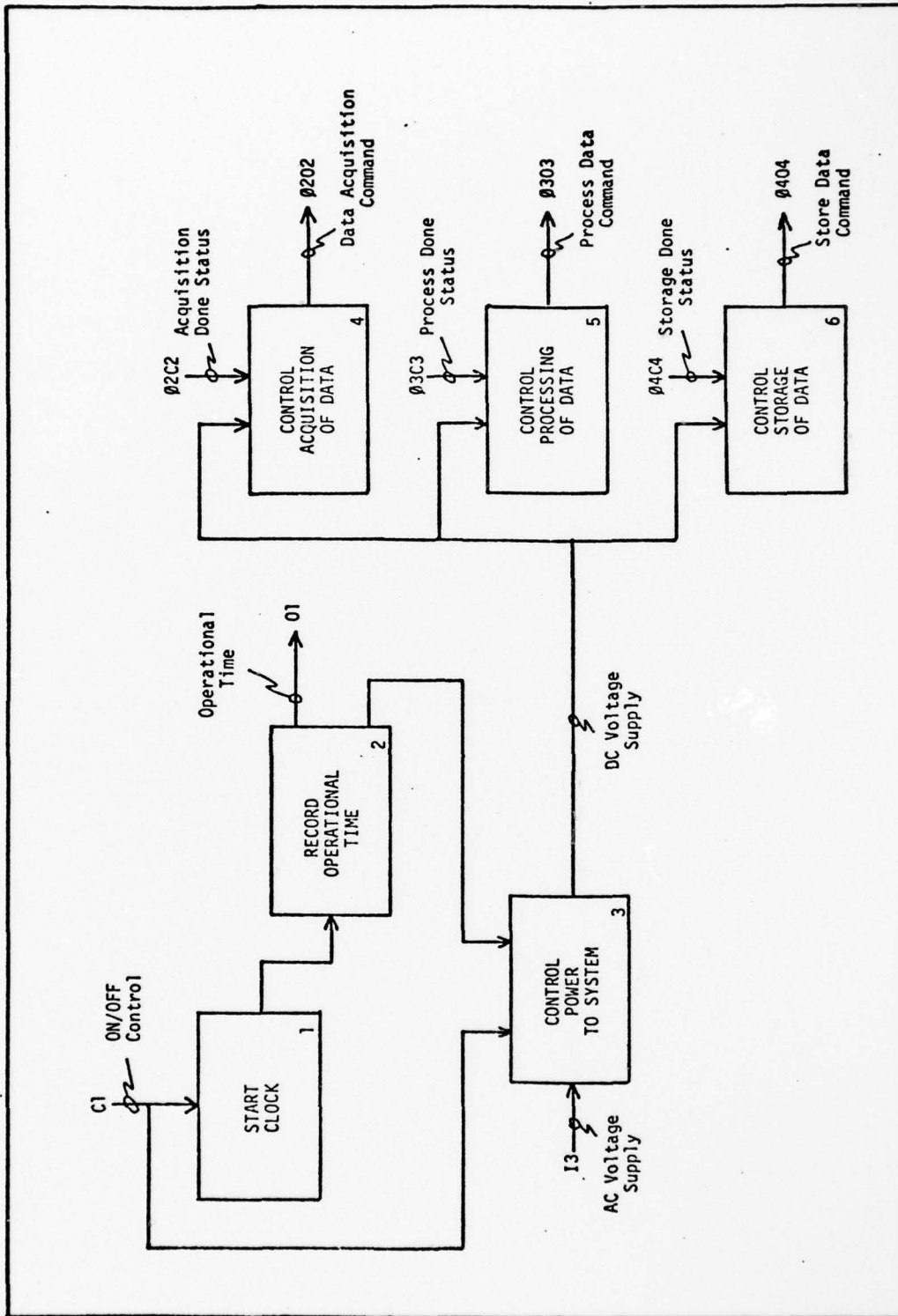


Figure 6. Node A1; Control A11 Activity

Control signal C1 (switch ON/OFF) activates the system. The power to the system is also controlled by total operational time; therefore, at the end of 15 to 20 hours of operation, the system halts automatically (Ref 13). The output and control signals are listed below along with their explanation.

- Output Signal

a. Operational Time (01): Four binary coded decimal (BCD) numbers, each 4-bit long, display the operational time elapsed in hours and minutes.

b. Data Acquisition Command (0202): This control signal initiates the acquisition of N-data points. This signal acts as a control signal to node A2.

c. Process Data Command (0302): This signal starts the processing of either the raw data or update of CDF versus temperature and frequency array.

d. Store Data Command (0403): This command signal initiates either the storage of the N-processed data element or storage of updated CDF versus temperature and frequency array.

- Control Signal

a. Switch ON/OFF Control (C1): This is a logic "1" when the input switch is closed and a logic "0" when the input switch is open.

b. Acquisition Done Status (02C2): This signal indicates node A14 whenever data acquisition of N-data points is complete.

c. Process Done Status (02C3): This control signal indicates the completion of processing by node A3. The processing is initiated by node A15.

d. Storage Done Status (Ø4C4): Whenever the storage of data is required, node A16 initiates this activity. Once the storage of data is complete, this signal indicates the status.

Node A2, Acquire Data (Figure 7). This node converts the analog vibration and analog temperature signals to an equivalent digital signal. The acquisition of data is initialized by node A1, and either the analog vibration signal or the analog temperature signal is selected. The selected signal is digitized, and the final output is an equivalent digital value of the input signal along with a signal identifier. The input/output and control signals specifications are:

- Input Signal

- a. Analog Vibration Signal (I1): 0-5V.
- b. Analog Temperature Signal (I2): 0-5V.

- Output Signal

a. Acquisition Done Status (Ø101): This signal indicates the status of the data acquisition process. Whenever N-data elements have been acquired, this status signal informs node A1 about the completion of data acquisition.

b. Formatted Data (Ø302): This is the digitized data along with an identifier which indicates whether the data point corresponds to the analog vibration signal or the analog temperature signal.

- Control Signal

Data Acquisition Control (Ø1C1): This signal initiates the acquisition of N-data elements. This signal controls node A2 in the following manner:

(1) It selects either the analog vibration signal or the analog temperature signal.

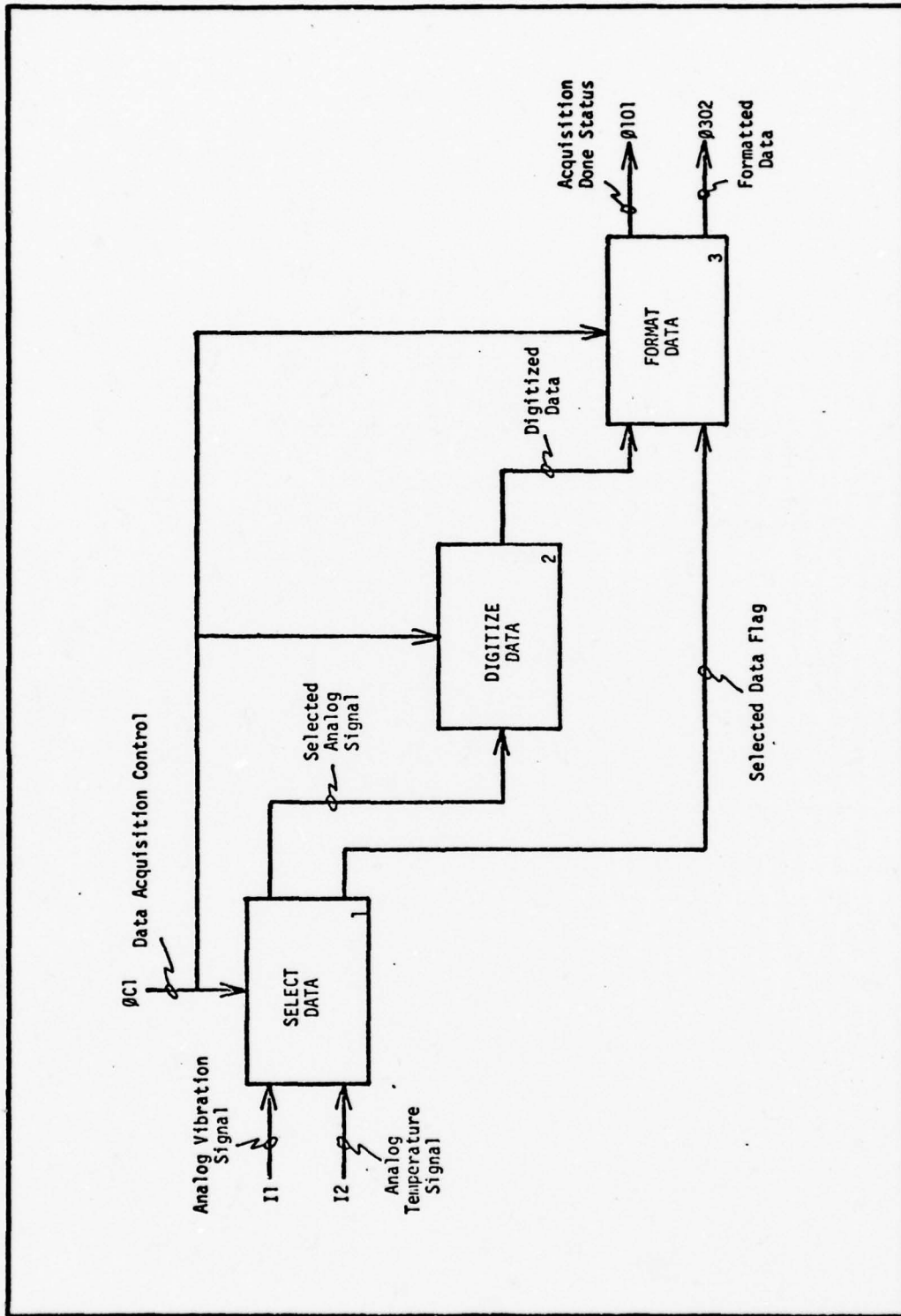


Figure 7. Node A2, Acquire Data

(2) It initiates digitizing of N-samples of an analog vibration signal.

(3) It initiates digitizing of one sample of an analog temperature signal for every N-samples of an analog vibration signal.

(4) It initiates formatting of each digitized data point.

Node A3, Process Data (Figure 8). The activities for this node are listed below:

a. N-elements of raw data array are taken from the raw data storage area and processed. The processing involves Fourier transform of N-raw data points. The power spectrum of first M-frequency components is computed next (M is less than N). The DF is calculated for each of the M-frequency components.

b. M-elements of the final data array (CDF versus temperature and frequency) are updated whenever a new set of DF array is available after processing of N-raw data points.

The input/output and control signals specifications are:

- Input Signal

a. Raw Data From Storage (Ø4I1): The N-element array from the Raw Data Storage area is used by the Process Raw Data module to give an M-element processed array.

b. New Data Array (Ø4I2): This is the data array which has resulted from processing N-element raw data array.

c. Old Data Array (Ø4I3): CDF versus temperature and frequency array which has to be updated by adding in new data array. The Update Data Module creates updated data array.

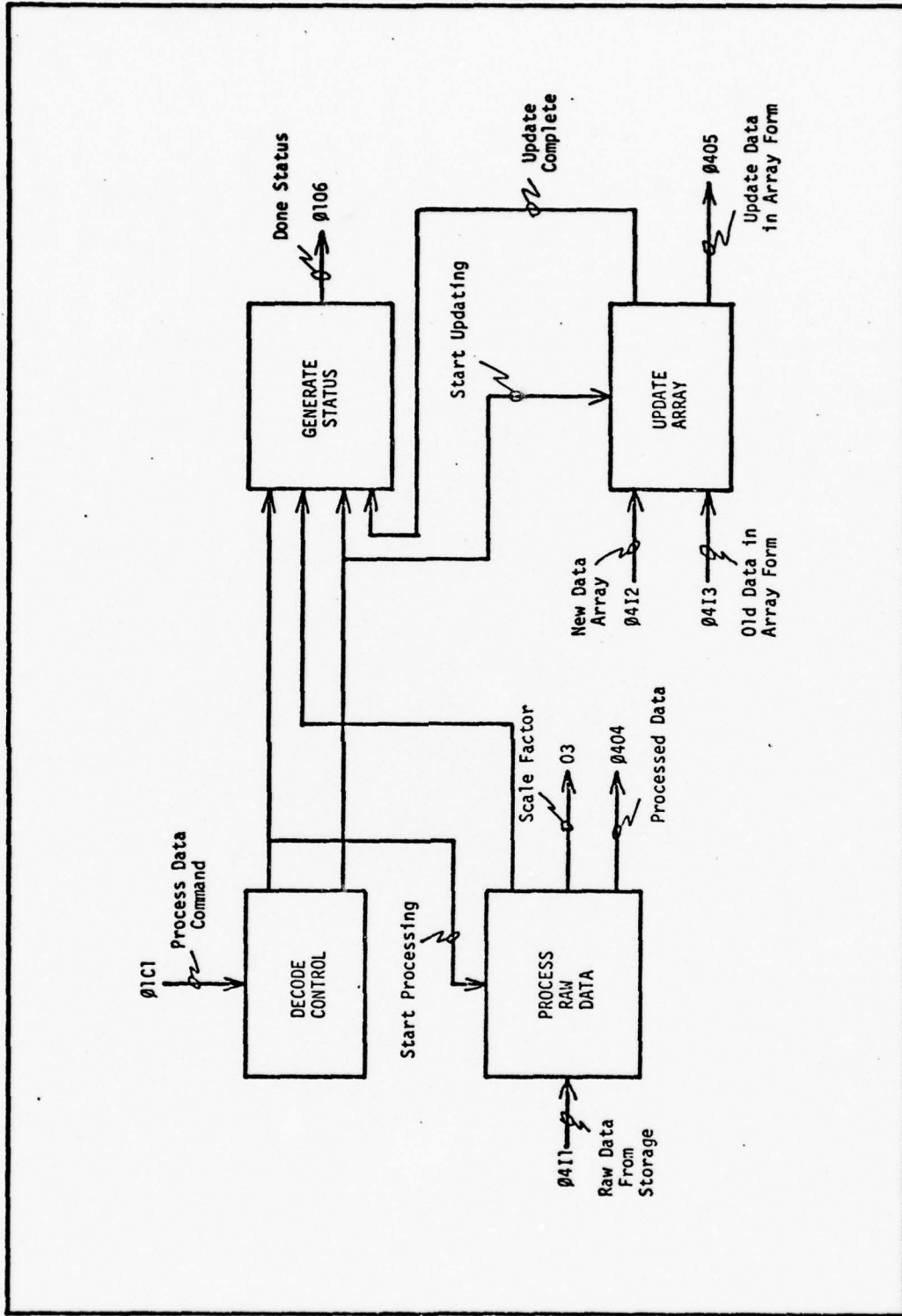


Figure 8. Node A3, Process Data

- Output Signal

- a. Scale Factor (03): This is the value which tells by what amount the elements in the final arrays have been scaled.
- b. Processed Data (0404): This is an M-element data array where the Process Raw Data module stores the results after operating on raw data array.
- c. Updated Data Array (0405): This is an M-element array. The value of CDF as a function of frequency and temperature is stored in this array. Every time a new set of M-processed data elements is available, the final array is updated by adding these new values to the old array.

- Control Signal

Process Data Command (01C1): This command initiates the processing of N-raw data elements. After the processing is completed, the CDF versus temperature and frequency array is updated. The update takes place by adding the newly processed data array to the existing final array (CDF versus temperature and frequency).

Node A4, Store Data (Figure 9). This node also performs multiple activities on the data. The different activities are enumerated below.

- a. The Formatted Raw Data points are stored sequentially in an N-point array. This array is called Raw Data Array, and it serves as a temporary storage area.
- b. Whenever raw data has to be processed, N-elements of the raw data array are transferred to node A3 for processing.
- c. The processed data is stored in M-elements processed data array. This storage of processed data points is temporary.
- d. Whenever the final output arrays have to be updated, M-data elements from the Process Data Array are transferred to node A3.

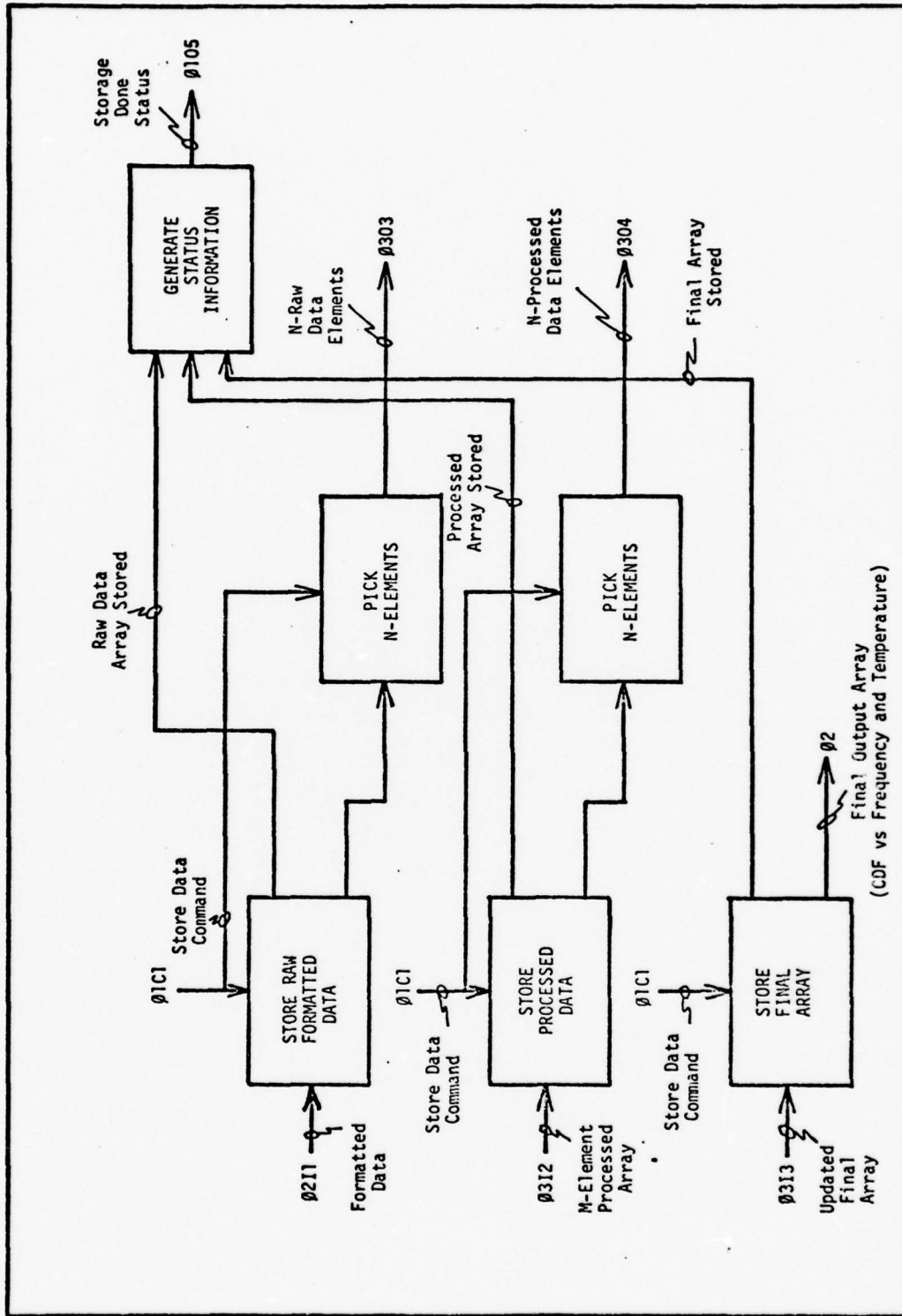


Figure 9. Node A4, Store Data

e. The Update Final Array (Ø3I3) is sent from node A3, and it is stored as the final CDF versus temperature and frequency array (Ø2).

The input/output and control signal specifications are:

- Input Signal

a. Formatted Data Input (Ø2I1): This is the formatted digitized version of input analog signals. The data points at this stage are also known as Raw Data Elements.

b. M-Element Processed Array (Ø3I2). This is an M-element processed array, which is the output from node A3. This array is generated every time N-row data elements are processed.

c. Updated Final Array (Ø3I3): The final output array (CDF versus frequency and temperature) has to be updated whenever a new M-element processed array is ready. This input is the result of updating the existing final output array.

- Output Signal

a. N-Raw Data Elements (Ø3Ø3): Node A3 requires N-row data elements whenever processing has to be done. This output is actually an N-row data element array, which is sent to node A3 for processing.

b. M-Processed Data Elements (Ø3Ø4): This output is used to update the final output array, so the CDF versus frequency and temperature is continuously updated.

c. Storage Done Status (Ø1Ø5): This signal is used to synchronize different storage operations. The Control A11 Process, node A1, is informed about the storage status of node A4.

d. Final Output Array (Ø2): This array stores CDF as a function of frequency and temperature and is the final output for MIBDAPS.

### Summary

In this chapter, the overall requirements placed on MIBDAPS are defined. All possible activities which MIBDAPS is expected to perform, along with their interrelationships, have been consistently and completely documented. At this stage, engineering decisions about the mechanization are not taken; however, the next phase of the design would be the engineering decision regarding the mechanization of the system. Hardware and software subdivision would also take place in the next phase.

### III. Software/Hardware Division of MIBDAPS

#### Introduction

The functional analysis model for MIBDAPS, which was developed in Chapter II, does not discuss mechanization of the system. The choice of the hardware for mechanization of MIBDAPS will, to a large extent, determine the subdivision of the system functions. In this chapter, the engineering decisions that define the realization of MIBDAPS are presented. The implementation of different modules of MIBDAPS and the subdivision of different modules into software and hardware activities are also presented. SADT is again employed in this development of MIBDAPS' activities.

#### Choice of Microcomputer System

One of the objectives in the development of MIBDAPS was a self-contained microprocessor based data acquisition and processing system. Thus, there were two possibilities for the mechanization of MIBDAPS:

i) Selection of an off-the-shelf microprocessor with the requisite word length and instruction set, followed by the development of a useable microcomputer system with the required amount of read/write memory (RAM), input/output (I/O) interface circuits, and development of application software.

ii) Selection of an industry standard microcomputer board with already developed I/O interface circuits, on-board memory and memory expansion capability, and software development tools.

At this stage, the approach taken for implementation of MIBDAPS was, in part, guided by the time available for the project. The first

possibility where the system could be developed around a selected microprocessor would encompass design of I/O interfaces, memory boards, and finally the development of software. This would have taken longer than the allotted time. Therefore, it was decided that an industry standard microcomputer board along with its existing I/O interface circuits, memory options, and software development tools would be used in the implementation of MIBDAPS.

#### Selection of Microcomputer Board

The main features which were considered in the selection of the microcomputer board are given below:

Word Size

I/O Capability

Instruction Set

Memory Capacity

Software Development Tools

Since the MIBDAPS is to process the input signals in real time, the choice of the microcomputer word size and the instruction set would determine the processing speed of the system. The word size determines the number of memory accesses required for execution of each instruction. For the same instruction set, a machine with a smaller word size requires more memory access as compared with a machine with a larger word size (Ref 4). Thus, a 16-bit microcomputer was an obvious choice. The signal processing function (Ref Chap I) of MIBDAPS required that multiply and divide instructions should be a part of the instruction set. Although subroutines can be written to carry out multiplication/division of binary numbers, this technique is orders of magnitude slower than hardwired multiply/divide

instructions. This is because the hardwired instructions are micro-programmed, which is a lower primitive than the software implemented subroutine. Table II lists the candidate microcomputer boards. Texas Instruments' TMS 9900 and the Digital Equipment Corporation's LSI-11 have multiply/divide instructions as part of their instruction set. Also, both of these microcomputers have a 16-bit word length. Therefore, the choice was narrowed down to these two boards.

The data acquisition function of MIBDAPS (Ref Chap I) dictated the need for an analog-to-digital converter (ADC) board which would be compatible with the selected microcomputer. Also, the need for synchronization and timing of real time signal processing made it desirable to have a hardwired system to perform this function. The Digital Equipment Corporation (DEC) manufactures both an ADC board and a real time clock (RTC) board that are compatible with the LSI-11. This feature of the LSI-11 made it an obvious choice. Recently, Norden Systems Incorporated, which is a subsidiary of United Technology Incorporated, announced the manufacturer of a bit-sliced based microcomputer board (designated LSI-11M) which is six times faster than its commercial counterpart (see Table II). This board also conforms to military specifications (mil specs). Thus, the mil spec LSI-11M was selected as the microcomputer board to be used for MIBDAPS.

#### LSI-11 Microcomputer System

The LSI-11 microcomputer system is a microprocessor based version of DEC's PDP-11 series minicomputer. The PDP-11 development software (assemblers, debuggers, loaders, compilers) is useable on the LSI-11 based system.

TABLE II

## Candidate Microcomputer Boards

Manufacturer	Candidate Microcomputer Board	Word Length	Resident RAM	Multiply Instruction Time ( $\mu$ sec)	Memory Expansion	I/O Lines	Vectored Interrupt (levels)
Norden Systems Inc.	LSI-11M*	16-bits	4K	12.5F**	32K	17	Yes (1)
Digital Equipment Corp.	LSI-11 (with KEV-11)	16-bits	4K	75/120F**	32K	17	Yes (1)
Texas Instruments Inc.	TI 990/4	16-bits	4K	18.0 $\mu$ sec	20K	16	Yes (8)
Intel Corp.	SBC 80/XX	8-bits	256-4K	None	Yes	22-48	Yes (4-8)
Data General Corp	$\mu$ nova 61	16-bits	2K/4K	None	32K	-	Yes

\*Reference Norden Systems Product Brochure on LSI-11M.

\*\*F = fixed-point arithmetic only.

(Adapted From Ref 4)

In the following paragraphs, a brief description of the LSI-11 processor, interface/memory options, and software development tools that were used in the implementation of MIBDAPS is given. These options were selected because they support either the hardware or the software realization of different modules of MIBDAPS. The detailed description of each of the above modules and the installation instructions are given in Ref 11.

KD11-F Microcomputer (Ref 11:87-99). The KD11-F microcomputer is contained on a single 8.5-inch by 10-inch printed circuit board. The module also contains resident dynamic 4K by 16-bit semiconductor RAM. Figure 10 is a functional block diagram of the microcomputer board. The salient features of this microcomputer are given below:

- This is a 16-bit microcomputer board using four custom LSI integrated circuit chips.
- Direct addressing of 32K 16-bit words or 64K 8-bit bytes ( $K = 1024$ ).
- Efficient processing of 8-bit bytes without the need to rotate, swap, or mask.
- Hardware memory stack for handling structured data, subroutines, and interrupts.
- Direct memory access (DMA) for high data rate devices inherent in the bus architecture.
- Eight general purpose registers that are available for data storage, pointers, and accumulators; two are dedicated—stack pointer (SP) and program counter (PC).
- A bus structure that provides position-dependent priority to the peripheral devices that are connected to the I/O bus.
- Vectored interrupt facility.
- Typical operating speed is 400 nsec based on a 10 MHz clock.

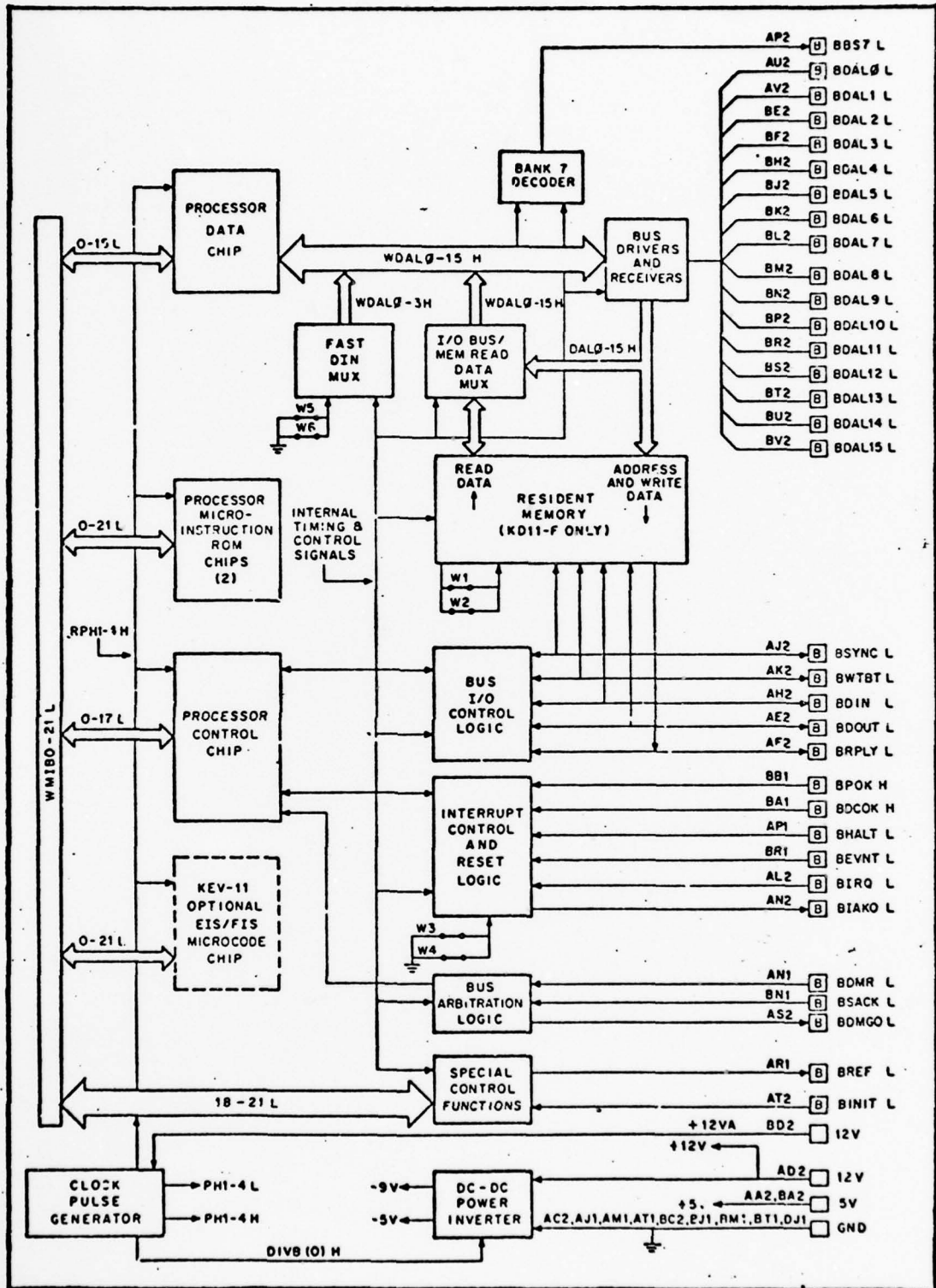


Figure 10. KD11-F Microcomputer Logic Block Diagram (From Ref 11:38)

- The power supply requirements are:

+5V +5% at 2.4A maximum current

+12V +5% at 1.10A maximum current

MMV11 4K by 16-Bit Core Memory (Ref 11:101-102). The MMV11-A 4K by 16-bit core memory option provides a nonvolatile read/write storage of the user's program and data (Figure 11). This optional module was used so that system software which would be developed for MIBDAPS could be stored permanently. This would eliminate the need to reload the program whenever the system is switched ON. The processed data would also be stored in the core memory for either updating or retrieval at a later time. The salient features of the MMV11-A core system are:

- 4096 by 16-bit capacity.

- Typical access time is equal to 425 nsec (475 nsec maximum); full read/restore cycle time is equal to 1.15 nsec.

- User-selected bank address; this allows the user to assign any 4K bank address.

- Power requirements are:

+5V +5% at 2.0A maximum

+12V +5% at 0.56A maximum

ADV11-A Analog-to-Digital Converter (Ref 11:186-189). The ADV11-A is a 12-bit successive approximation analog-to-digital converter with a built-in multiplexer and sample-and-hold for use on the LSI-11 bus (Figure 12). The multiplexer selection accommodates 16 single-ended or 8 quasi-differential inputs.

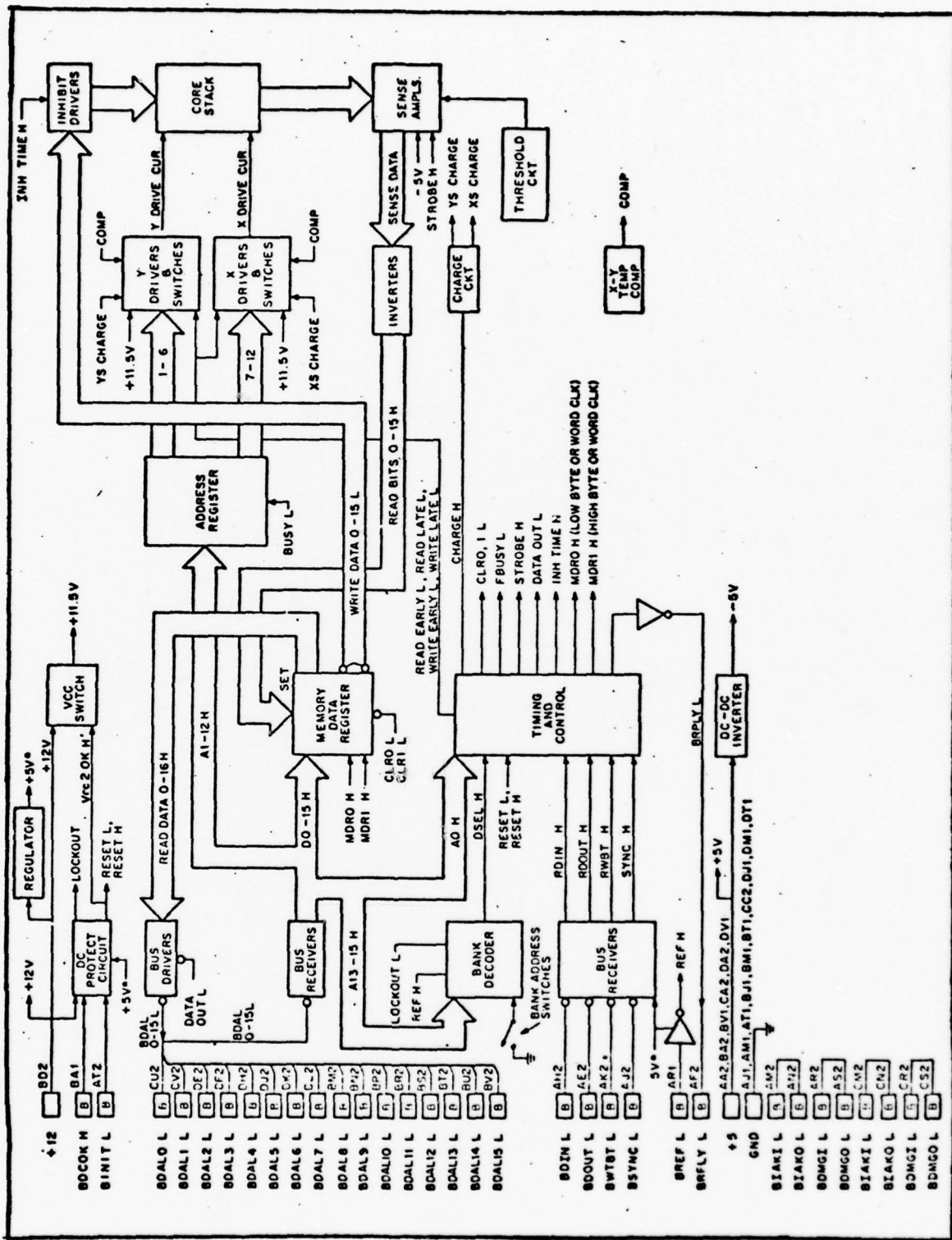


Figure 11. NW11-A Logic Block Diagram (From Ref 11:105)

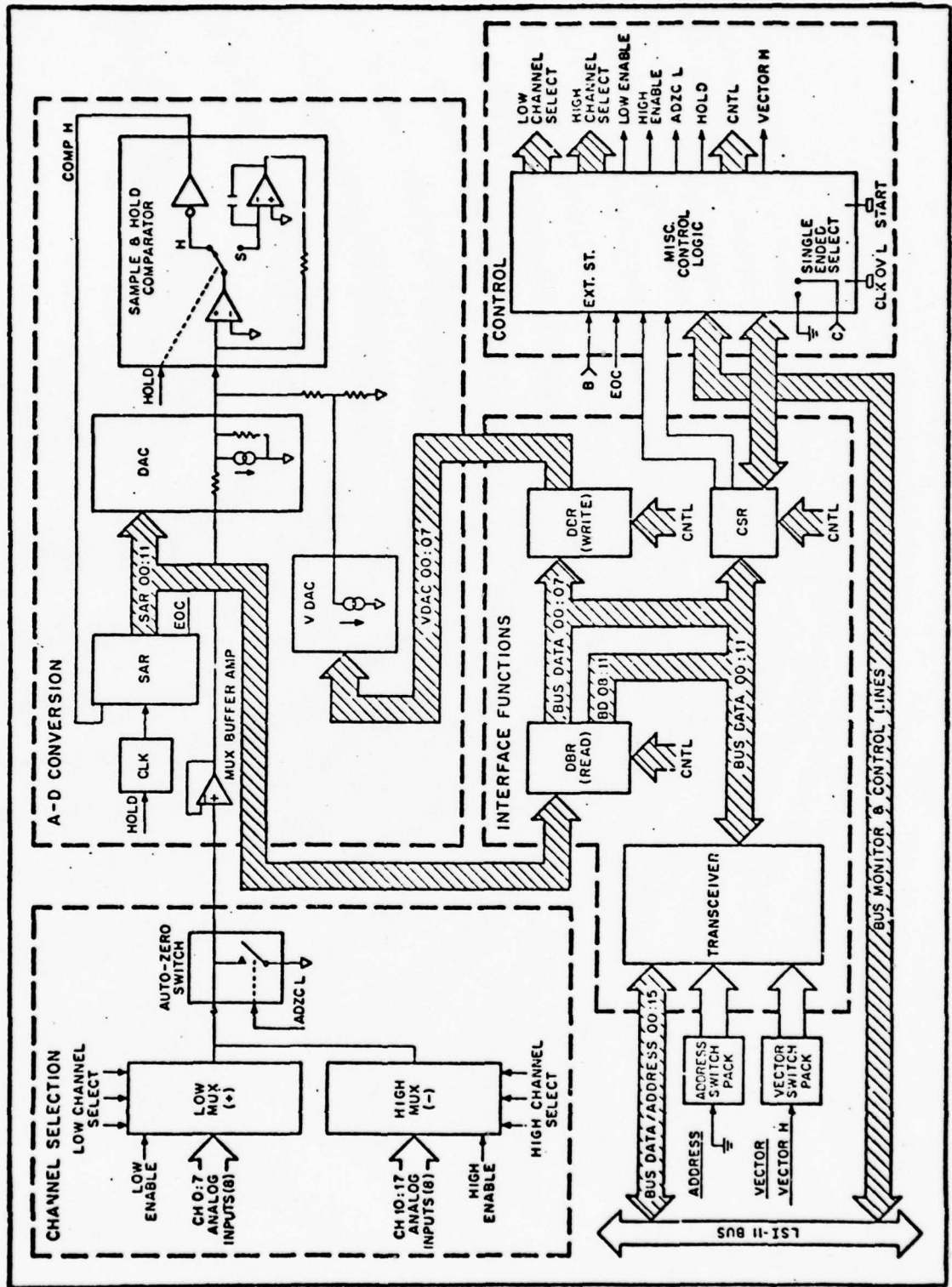


Figure 12. ADV11-A Function Block Diagram (From Ref 11:191)

This A/D converter was adequate for MIBDAPS requirements because only two analog signals needed to be digitized. Also, it was required that the conversion error should not exceed +1% of the full scale (Ref 13). For analog signals of 0-5V, this requirement translates to an error of +5mv. The one least significant bit (LSB) value for the ADV11-A is 2.5mv which means that the maximum error that can be introduced is +2.5mv, which is within the requirement.

The main features of ADV11-A are:

- Analog input full scale range is 10.24V bipolar (-5.12V to +5.12V).
- Logic input voltages are:

low (logic "0") = 0.0 to +0.7V

high (logic "1") = +2V to +5V

- Resolution of ADV11-A 12 12-bit binary weighted. The format of the digitized signal is parallel offset binary right justified.

- The timing requirements are:

- a. External start should be a low-level pulse, 50 nsec minimum duration and 10  $\mu$ sec maximum duration. Conversion starts on the leading edge.

- b. Conversion time is 16 times the clock period. The maximum clock frequency is 100 KHZ.

- Power requirements are:

+5V +5% at 2.0A maximum current

+12V +5% at 450mA maximum current

KW11-A Programmable Real Time Clock (Ref 11:211-212). The functional diagram for the real time clock is shown in Figure 13. This programmable clock/counter provides a variety of means for determining time

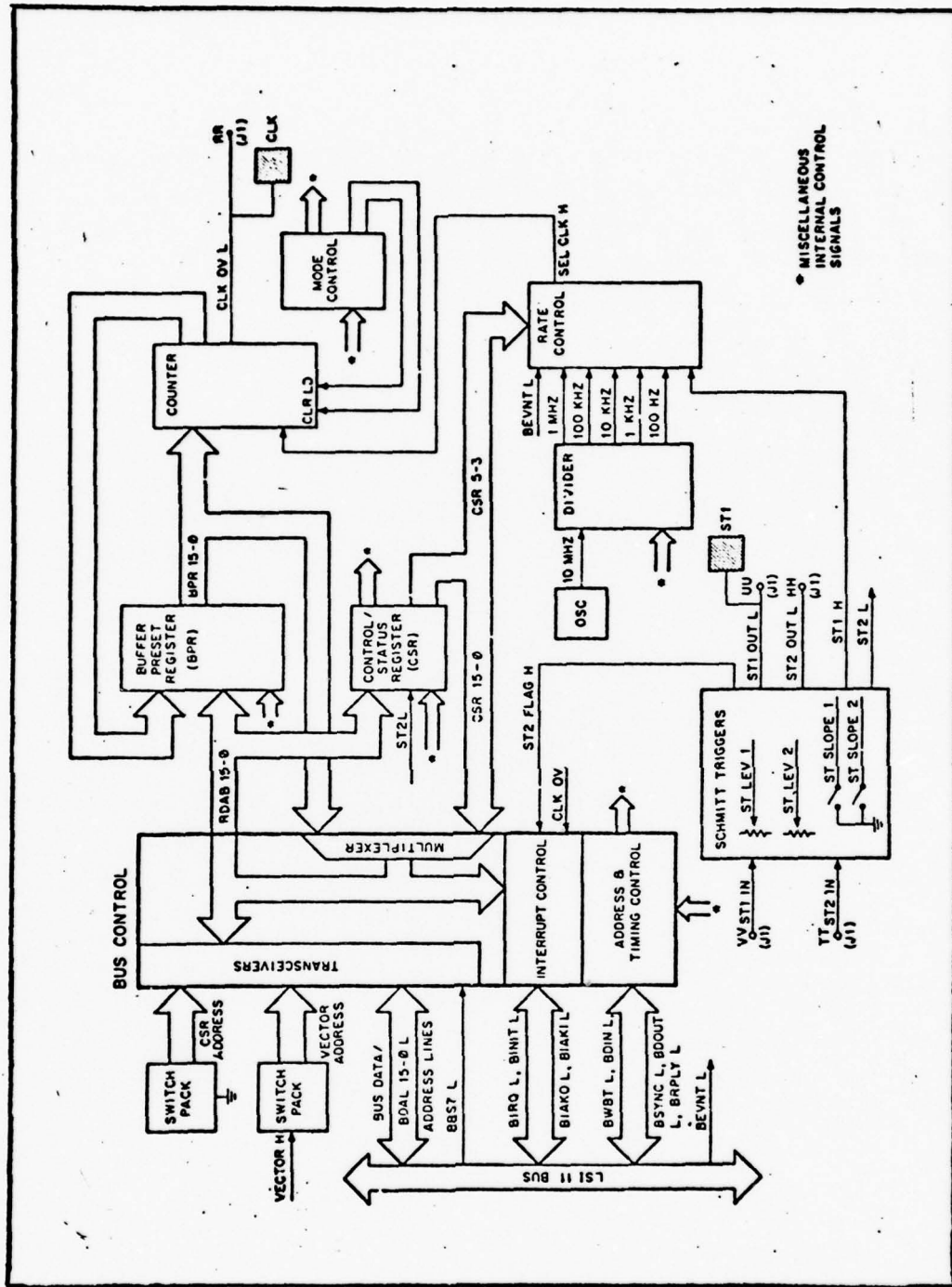


Figure 13. KxV11-A Real Time Clock (From Ref 11:213)

intervals or counting levels. It can be used to generate interrupts to the LSI-11 processor at predetermined intervals to synchronize the processor to external events or to measure time intervals or establish programmed ratios between input and output events. It can also be used to start the ADV11A analog-to-digital converter either by clock counter overflow or by the firing of a Schmitt trigger circuit.

The clock counter has a resolution of 16 bits and can be driven from any five internal crystal-controlled frequencies (100 HZ to 1 MHZ). It can also be driven from a line frequency input or from a Schmitt trigger fired by an external clock input. The performance specifications for the KW11-A are:

- Internal Clock

The internal clock has a frequency accuracy of  $\pm 0.01\%$ . The base frequency is 10 MHZ, and it is divided into five selectable rates (1 MHZ, 100 KHZ, 10 KHZ, 1 KHZ, and 100 HZ).

- External Clock

The external clock can either be the line frequency input or a clock generator having the desired frequency.

- Output Signals

a. A clock overflow signal occurs whenever a certain timing event is completed. The characteristics of this signal are: 500 nsec low asserted level pulse which is TTL compatible.

b. Two trigger outputs, Schmitt trigger 1 and Schmitt trigger 2, are also available for either being used as a trigger for the ADV11-A (A/D converter board) or for any other external device. The characteristics are the same as that for the clock overflow signal.

DLV11 Serial Line Unit (Ref 11:147-148). The DLV11 is the basic interface module used for connecting asynchronous serial line devices to the LSI-11 bus. The functional diagram for the DLV11 is shown in Figure 14. The salient features of the DVL11 unit are:

- Either an EIA RS232 or a 20mA current loop device can be interfaced to an LSI-11 bus.
- There are 13 baud rates that can be selected: 50, 75, 110, 134.5, 150, 200, 300, 600, 1200, 1800, 2400, 4800, and 9600.
- The LSI-11 bus interface and control logic for interrupt processing and vector generation are also available on the board.
- The control/status register (CSR) and data buffer register (DBR) for programmed or interrupt driven transfer of data are available.

Software Development Tools. One of the reasons why the LSI-11 micro-computer system was chosen for mechanization of MIBDAPS was because it is software compatible with the PDP-11 computer family. The following systems that were available on Wright-Patterson Air Force Base could, therefore, be used for software development:

- The Department of Electrical Engineering at the Air Force Institute of Technology (AFIT), Wright-Patterson Air Force Base, Ohio, has a PDP-11/20 along with a paper tape software system. (A single floppy disk drive has been added; therefore, the RXV11 operating system could also be used.)
- The AFAL has a cross-assembler for PDP-11 assembly language on its DEC-10 computer system. Programs could be created on the DEC-10 system using text editors. The PDP-11 cross-assembler (MACY11) can then be used to create absolute loader formatted paper tapes. There is, however, no facility to simulate execution of these programs.

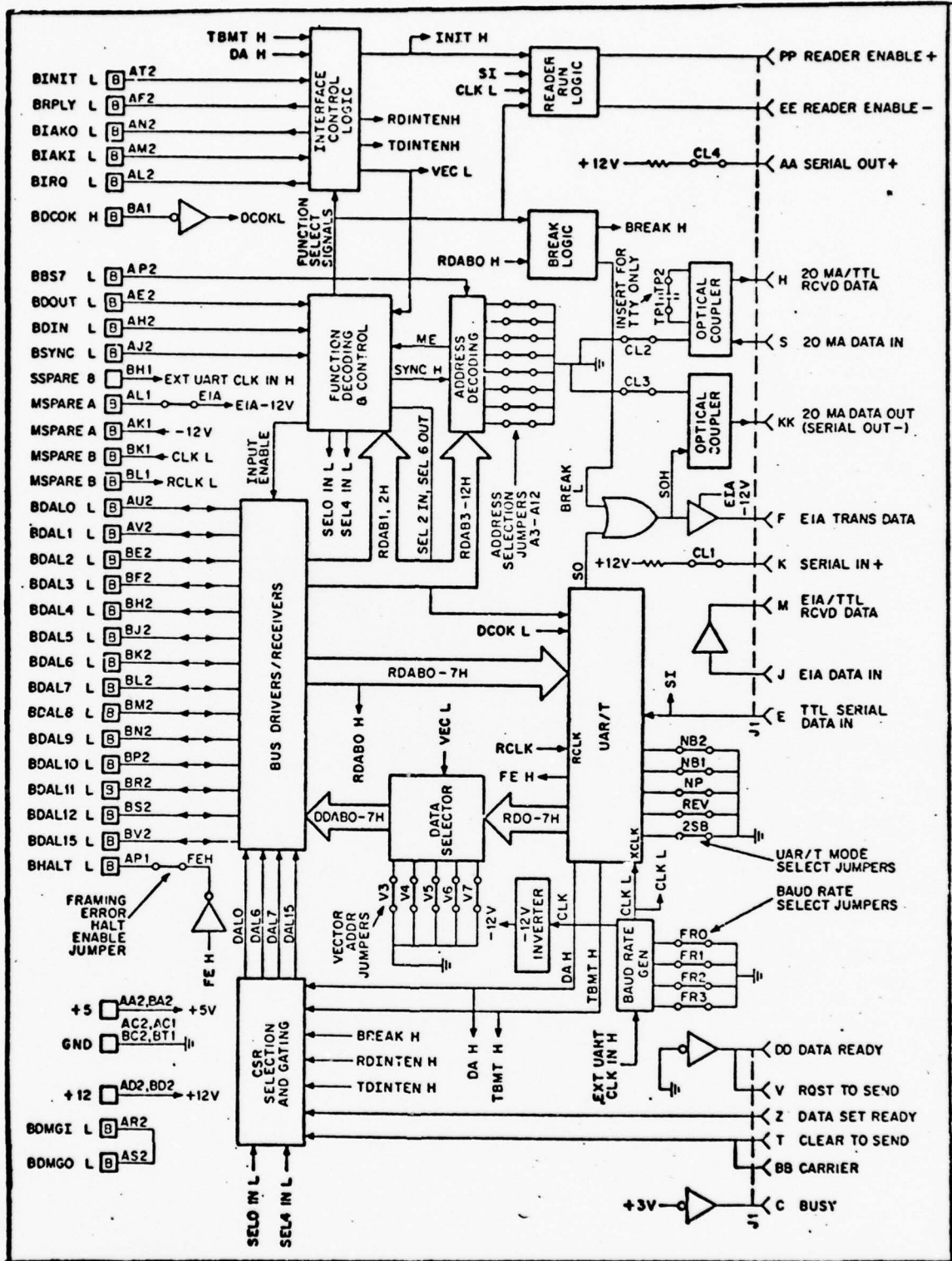


Figure 14. DLV11 Logic Block Diagram (From Ref 11:150)

- The AFML has a PDP-11/03 with dual floppy disk drive and an RXV11 operating system. This system uses an LSI-11 microcomputer and all interface boards. The RXV11 operating system has a text editor, macro assembler, linker, FORTRAN compiler, and peripheral interchange program (PIP) available on a floppy disk. This system was suitable not only for developing and debugging programs, but also for real time program execution.

The PDP-11/03 system at AFML was eventually used for the development of software for MIBDAPS.

#### Analysis and Design of MIBDAPS

The decisions about the implementation of different functional modules of MIBDAPS, either in software or hardware, were guided by the following factors:

- Capabilities and limitations of the KDF-11 processor module along with 4K words of core memory used as the microcomputer system.
- Capabilities and limitations of support hardware modules, in particular ADV11-A (ADC) and KWV11 (RTC).
- Speed of software execution.

The KDF-11 processor along with 4K words of core memory (subsequently referred to as "microcomputer system") would be used to implement the controlling and processing functions of MIBDAPS in software. The ADC board (ADV11-A) and RTC (KWV11) along with a system clock circuit (which would be designed and fabricated) were visualized as the only hardware needed to implement the data acquisition, signal processing, and synchronization functions of MIBDAPS. Keeping these two factors as guidelines, another in-depth study of the system requirement model developed in Chapter II was carried out. SADT was employed to carry out a functional analysis

of the system and a new activity model, keeping implementation in view, was developed. The node index for this activity model is given in Table III. It should be noted that this time the mechanization (M) for each node is also defined as hardware or software, or both. Software implementation implies programs which would be executed on the micro-computer system. Hardware implementation implies use of ADC, RTC, or an external system clock.

TABLE III

Node Index for New Activity Model

A-0	Acquire and Process Data (same as in Chap II)
A0	Acquire and Process Data
A1	Control All Activity
A13	Switch Power to System
A14	Stop Execution
A15	Start Execution
A16	Start Acquisition and Processing of Data
A162	Initiate Data Acquisition
A163	Initiate Signal Processing
A2	Acquire Data
A3	Process Data
A32	Compute Fourier Transform
A33	Compute Power Spectrum
A34	Compute Damage Factor
A35	Update CDF Array

The explanation of each node along with data structures used and input/output signal specifications follows.

Node A0, Acquire and Process Data (Figure 15). The difference between node A0 as compared with node A0\* (Figure 4) is the absence of module A4 (Store Data). This difference has come about because the Process Data Module will perform the task of storing and updating result that are stored. The various activities taking place in this node are listed below.

- The Control All Activity module performs the function of managing the control and execution of all other activities. It also provides the DC supply to all hardware modules.

- The Acquire Data module continuously samples and digitizes the input analog signals, analog vibration signal and analog temperature signal. Whenever a set of data values have been digitized, it sets the Acquisition Done flag.

- The Process Data module takes in the input raw data and processes it to give an array of damage factor versus frequency. This array is added on to one of the cumulative data factor versus frequency arrays. There are 11 CDF versus frequency arrays, one for each expected temperature intervals (Ref Chap II).

The input/output and control signal specifications are:

- Input Signal

- a. Analog vibration signal (I1): 0-5 volts.
- b. Analog temperature signal (I2): 0-5 volts.
- c. AC power (I3): 110 volts, 60 cycles.

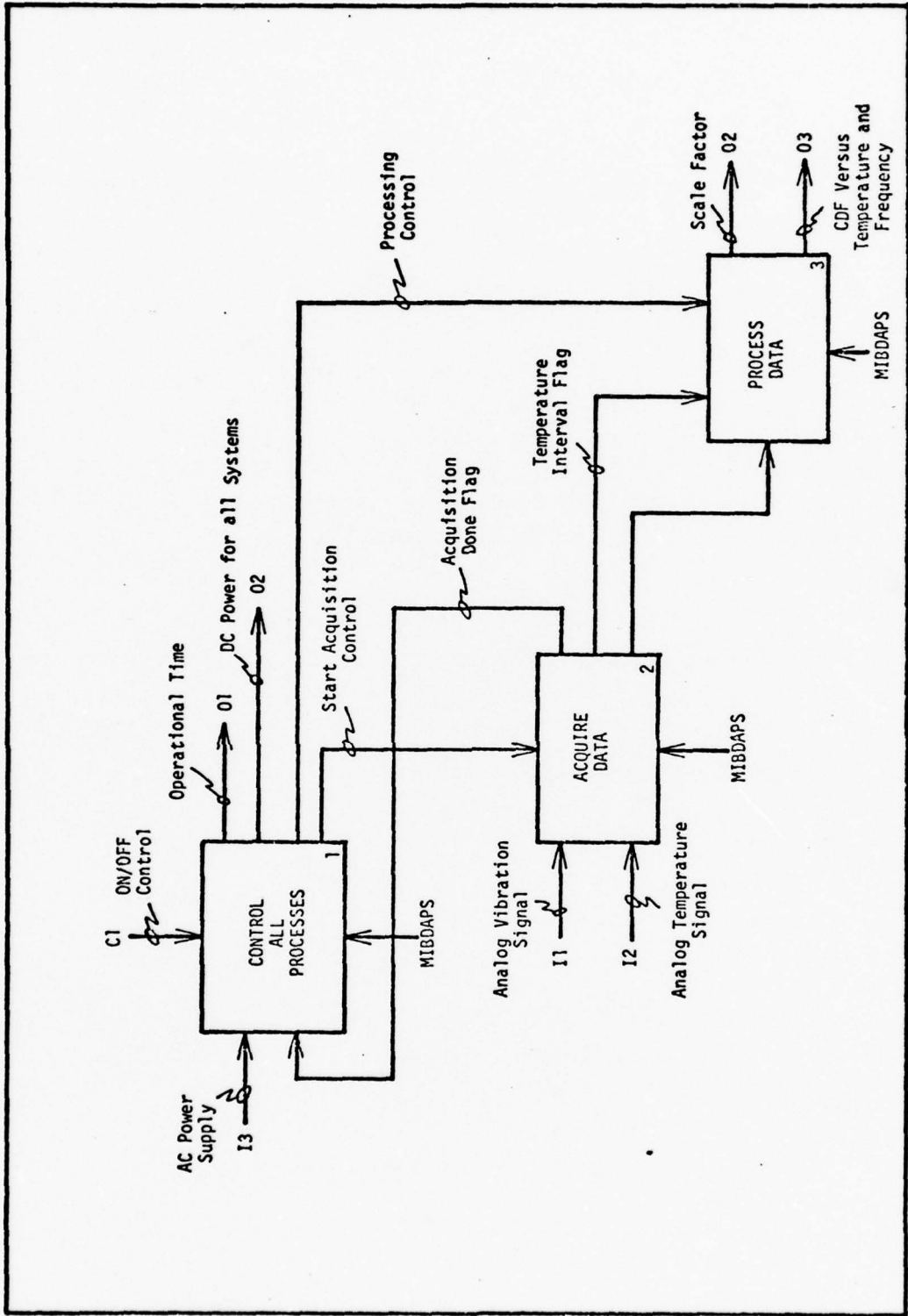


Figure 15. Node A0, Acquire and Process Data

- Output Signal

- a. Total Operational Time (01): A display of operational time in hours and minutes.
- b. DC Power to System (02):
- c. Cumulative Scale Factor (03): This is a 16-bit octal number which gives the scale factor for the output arrays.
- d. CDF Versus Frequency and Temperature Arrays (03): These are 13-word long arrays having the data structure as shown in Figure 16.

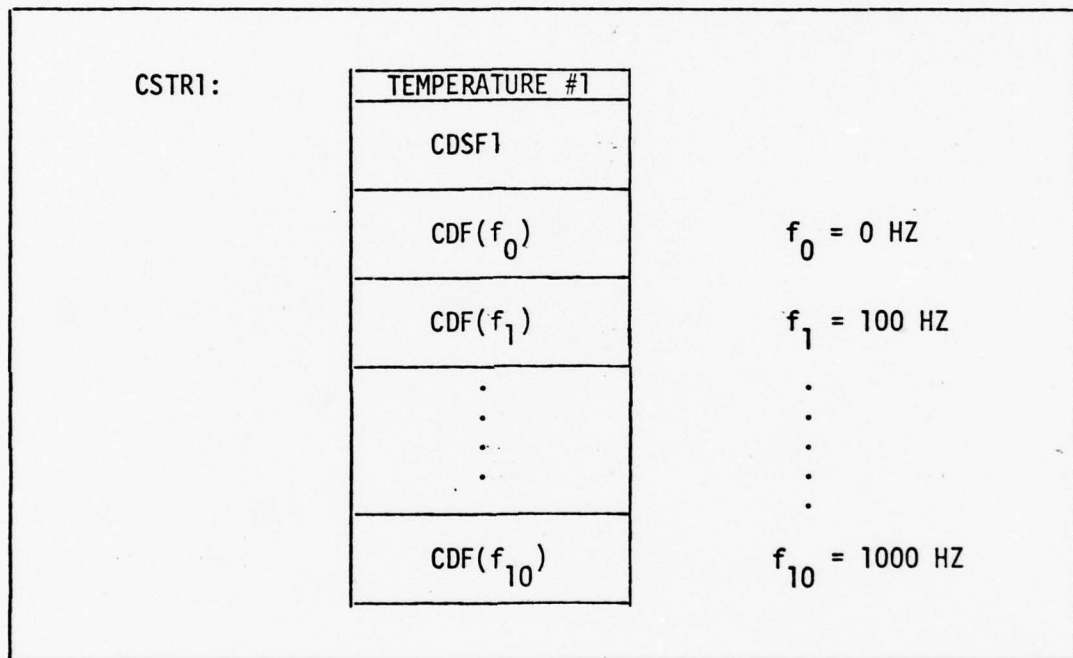


Figure 16. Data Structure for CDF Vs Frequency and Temperature Array

- Control Signal

System ON/OFF Control (C1): Once the system is switched ON and the AC power is available, the system starts functioning.

Node A1, Control All Activity (Figure 17). This node performs the activities necessary to control the synchronization of all other activities

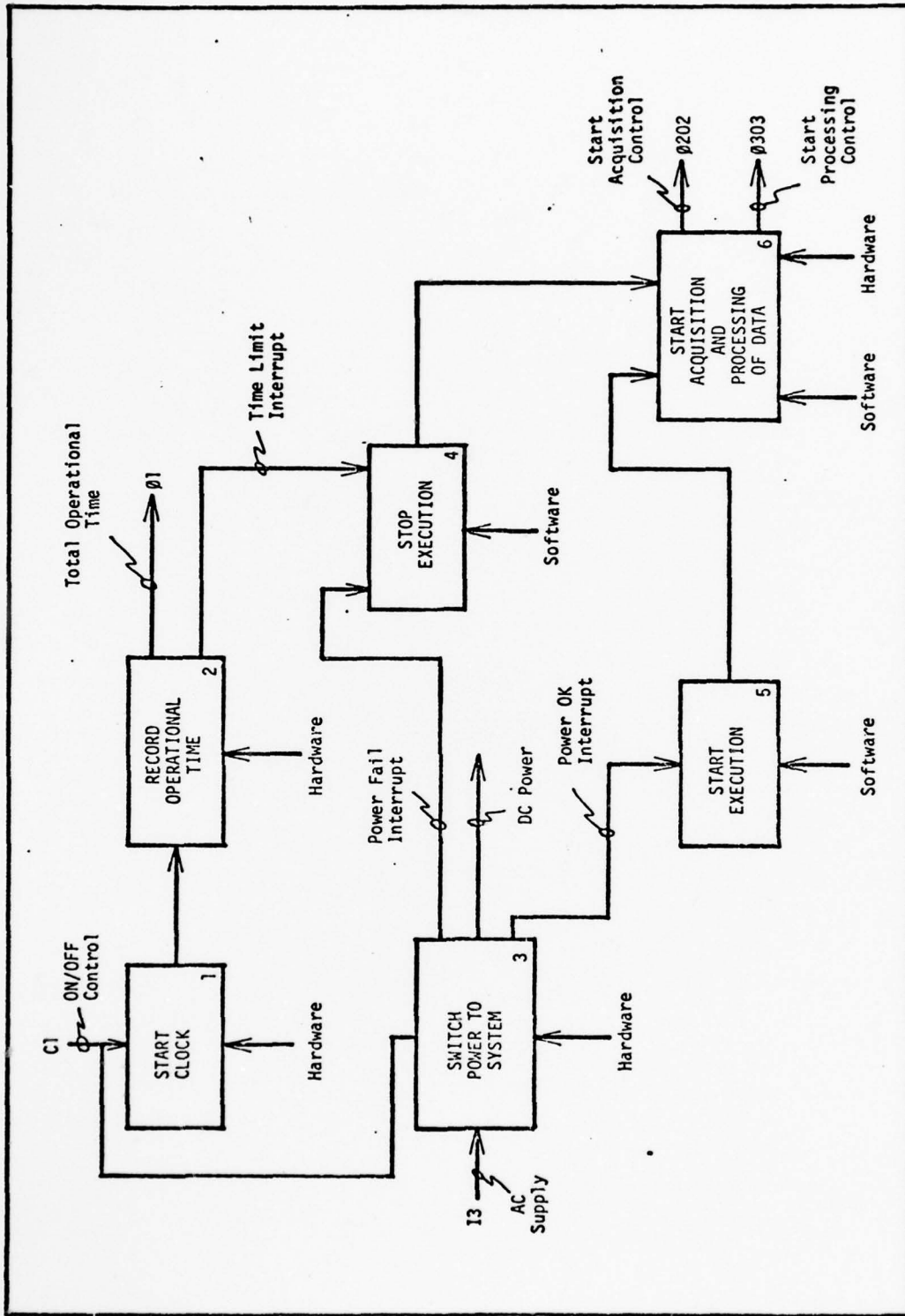


Figure 17. Node A1, Control A11 Activity

of MIBDAPS. The mechanization of different activities in this node are either done by hardware or software or a combination of both. It is assumed that the LSI-11 computer is used to execute the software; however, hardware mechanization includes various other modules which are being interfaced to the microcomputer. The various activities performed in this node are listed below.

- The DC power required by different modules of the LSI-11 micro-computer system is provided by this node. The 110-volt AC supply is converted to +5V and +12V DC supplies. Whenever the system's ON/OFF switch is selected to the ON position, the DC supply is switched ON. This is done by a Switch Power to System module. This module also generates an interrupt in case the DC power is within an acceptable limit. In case of a power failure, an interrupt is also generated.

- Total operational time for the system is also computed by the Record Operational Time module. In case the operational time reaches the selected time limit (15 to 20 hours), an interrupt is generated.

- The Start Execution module, after receiving the Power OK interrupt, initializes the system and sends a signal to Start Acquisition and Processing of Data.

- The Stop Execution module stops execution of the system function in case the following interrupts occur:

  - Power Fail Interrupt

  - Time Limit Interrupt

- The Start Acquisition and Processing of Data module initializes different variables which are used by the software. It then signals nodes A2 and A3 to start acquisition and processing of data.

Specifications for the input/output and control signals are:

- Input Signal

AC Supply (I3): 110 volts, 60 HZ.

- Output Signal

a. Total Operational Time (O1): Four 4-bit BCD numbers.

b. Start Acquisition Control (Ø202): A software-generated command. The control/status register of the ADV11-A A/D converter (ADC) is loaded with the appropriate value to enable the ADC. This initiates the data acquisition cycle.

c. Start Processing Control (Ø303): This is also a software-generated command, which activates the processing modules for input data processing.

- Control Signal

System ON/OFF Control (C1): A signal which initiates the operation of the system. A logic "1" value is available when the system is turned ON and logic "0" value when the system is turned OFF.

Node A13, Switch Power to the System (Figure 18). This node carries out the following activities:

- Converts the AC power supply to +5 volts and +12 volts DC.
- Generates an interrupt to the microcomputer when DC power is available.
- Detects a power failure and generates a Power Fail interrupt to the microcomputer.

The input and output signal specifications are:

- Input Signal

AC Supply (I3): This is the only input to this node. The system operates on 110V AC, 60 HZ.

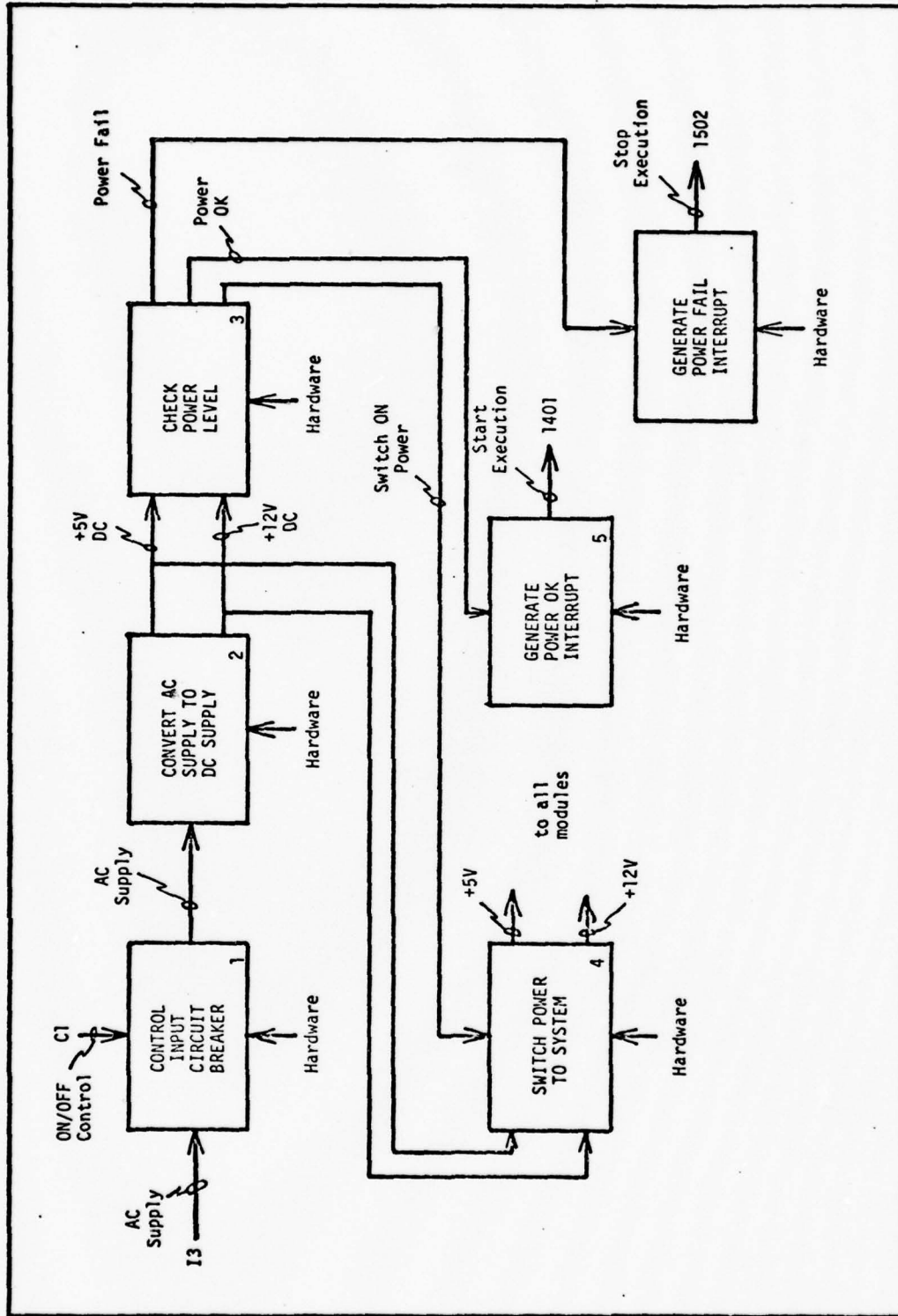


Figure 18. Node A13, Switch Power to the System

- Output Signal

a. DC Power Supply: The +5V and +12V DC power supply is routed to all modules of MIBDAPS.

b. Start Execution (1401): A software-generated signal which is used to start execution of the program. Bit-0 of variable STEXEC is set to initiate data acquisition and processing.

c. Power Fail Interrupt (1501): A software-generated signal which stops execution of the program. Bit-0 of variable STPXC is set to stop execution of the program.

Node A14, Stop Execution (Figure 19). This node receives the Power Fail interrupt; it then accesses the vector interrupt address for the service routine. The service routine is executed and bit-0 of the variable STPXC is set to a logic "1." This variable is sent as an output signal to stop execution of the system execution.

Node A15, Start Execution (Figure 20). This node receives the Power OK interrupt; it then accesses the vector interrupt address for the service routine. The service routine is executed and bit-0 STEXEC switch is set to logic "1." This variable is sent as to the GO Command to start system execution.

Node A16, Start Acquisition and Processing of Data (Figure 21). Once the system has been turned ON and DC power is made available to the system, this node starts execution of the data acquisition and data processing functions of MIBDAPS.

The Stop Execution Command (14C2) and Start Execution Command (14C1) are used to start and halt the execution of data activity. The input/output signal specifications are:

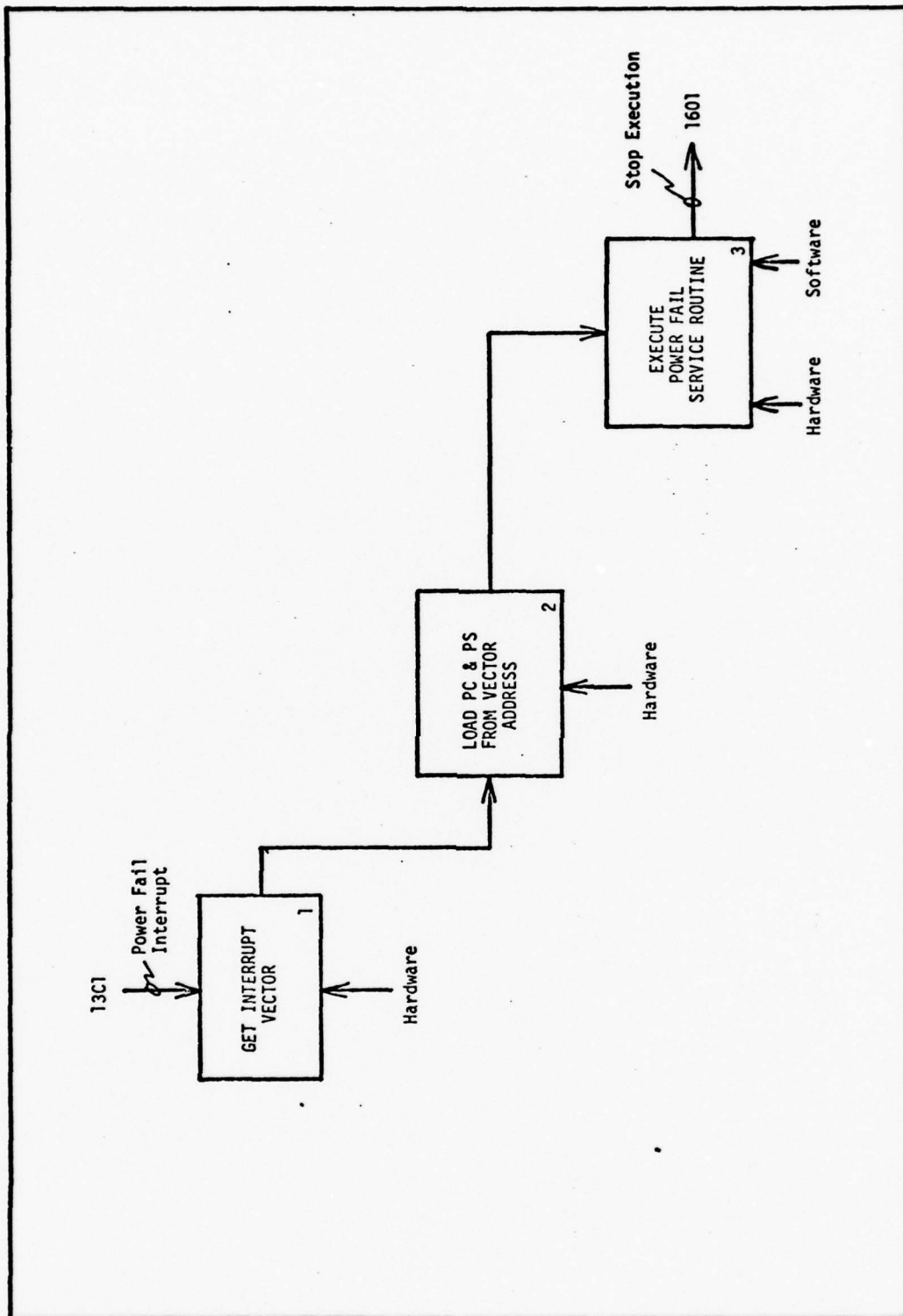


Figure 19. Node A14, Stop Execution

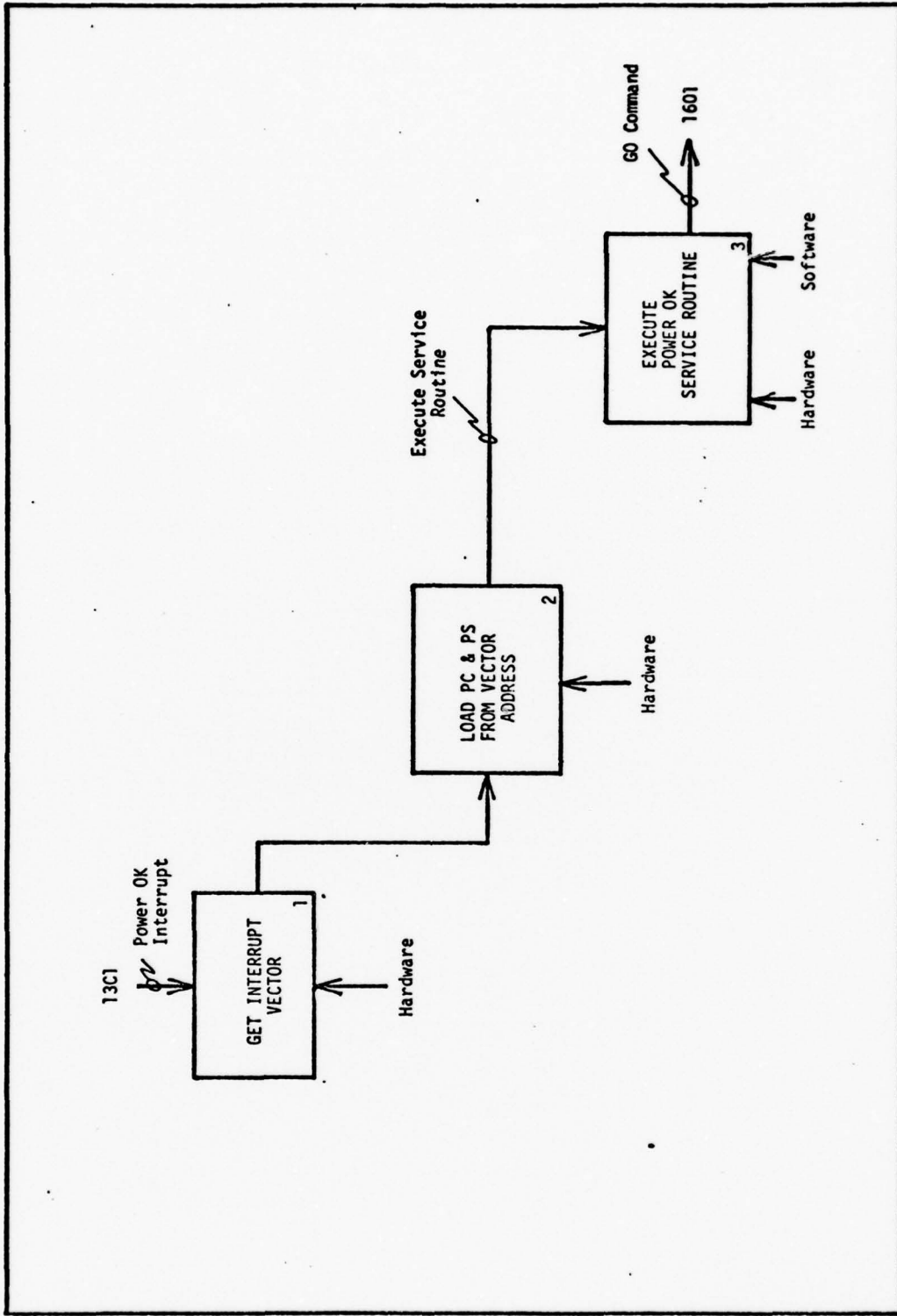


Figure 20. Node A15, Start Execution

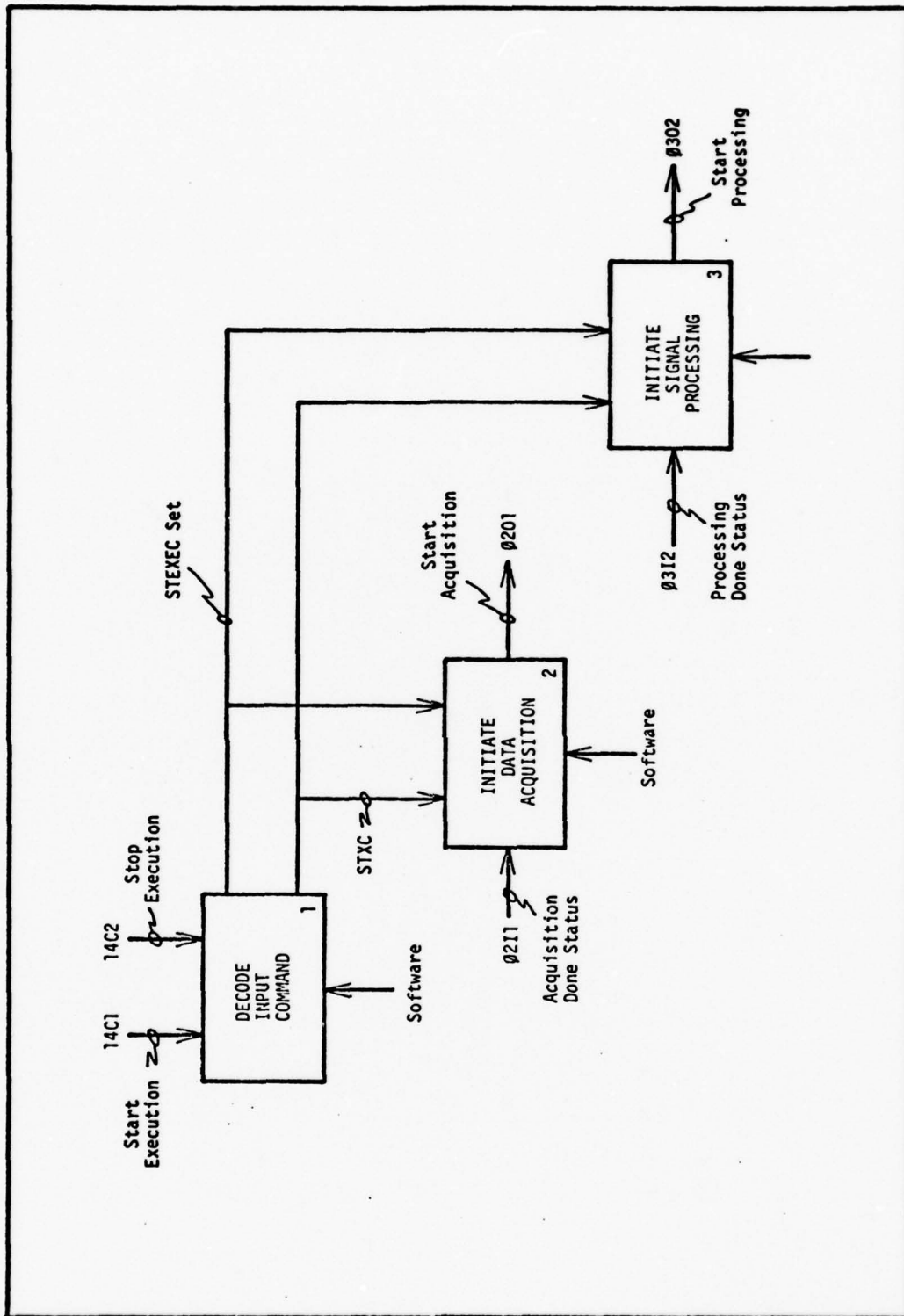


Figure 21. Node A16, Start Acquisition and Processing of Data

- Input Signal

a. Acquisition Done Status (Ø2I1): Whenever the acquisition of n-data points is complete, the variable ADFLG is set to logic "1." This variable along with ADCNT (number of data points to be converted) and CDATA (starting address of data buffer where data is to store) are passed to the Initiate Data Acquisition module. All of these variables need to be initialized whenever data acquisition has to be reinitiated.

b. Processing Done Status (Ø3I2): The following variables are passed to the Initiate Processing module for initialization: DATA - Contains starting address of data to be processed. SFACTR, CUMSF1, and CDSFi (i=1, 2...) - Scale factors used in node A3, which are explained later; however, each has to be cleared before processing.

- Output Signal

a. Start Acquisition (Ø2O1): This signal is generated when the control/status register (ADCSR) for the A/D converter is loaded with the appropriate bit pattern. Once the ADCSR is loaded, the acquisition of data starts.

b. Start Processing (Ø3O2): This is a software-generated command, and it initiates the processing of data.

- Control Signal

a. Start Execution (14C1): This is a software switch, STExec, which is set to logic "1" whenever DC power is available and the system has been initialized.

b. Stop Execution (14C2): This is also a software switch, STPXC, which is set to logic "1" whenever a power failure occurs.

Node 162, Initiate Data Acquisition (Figure 22). This node shows explicitly the functional decomposition of the module Initiate Data

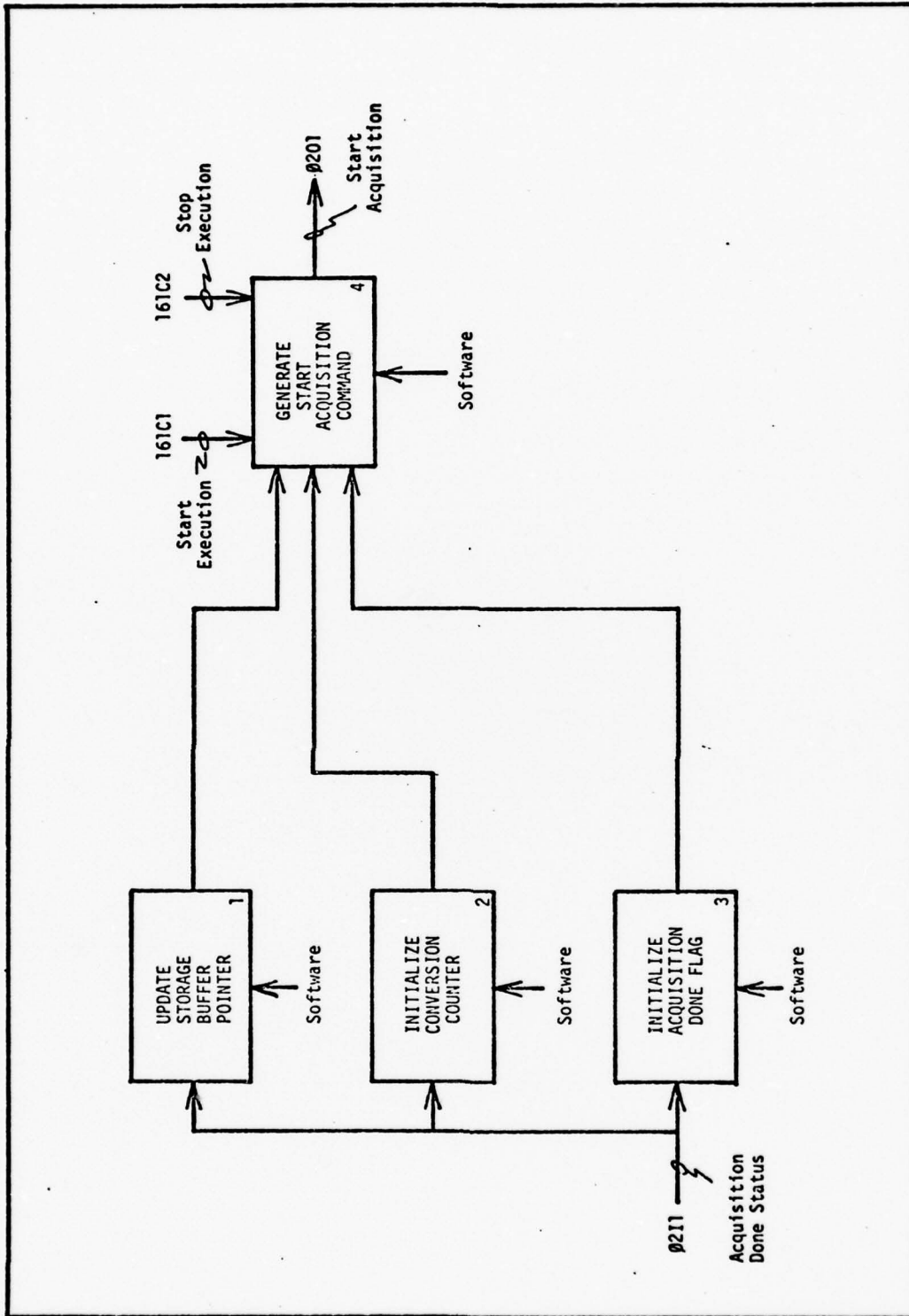


Figure 22. Node A162, Initiate Data Acquisition

Acquisition of node A16. The STEEXEC switch (161C1) and STPXC switch (161C2) control the activity in this node. The input signal Ø2I1—CDATA, ADCNT, and ADFLG—is initialized by modules 1, 2, and 3, respectively. The Generate Start Acquisition Command loads the ADCSR with appropriate bits to initiate acquisition of data by ADC, only after all of the above mentioned switches are initialized and STEEXEC is set to logic "1."

Node 163, Initiate Signal Processing (Figure 23). This node shows the functional decomposition of the Initiate Signal Processing module of node A16. The activities in this node are similar to node A162. The input signal Processing Done Status (Ø321) and output signal Start Processing (Ø301) are different. The input includes SFACTR, CUMSF1, CDSFi SPCTRA, and DATA switches. These are initialized in this node. The Start Processing Command generated in this node starts the execution of the processing of acquired data.

Node A2, Acquire Data (Figure 24). This node performs the activity of acquiring data and storing it in temporary storage. Once the acquisition of N-data points is complete, the acquisition is stopped and the ADFLG switch is set to logic "1." When the analog temperature signal is selected for acquisition, the temperature interval flag, TFLG, is loaded with a number from Ø to 11. The temperature interval number that is loaded in TFLG depends on the value of the temperature signal at the time of conversion. The input and output signal specifications are:

- Input Signal

- a. Analog vibration signal: 0-5V.
- b. Analog temperature signal: 0-5V.

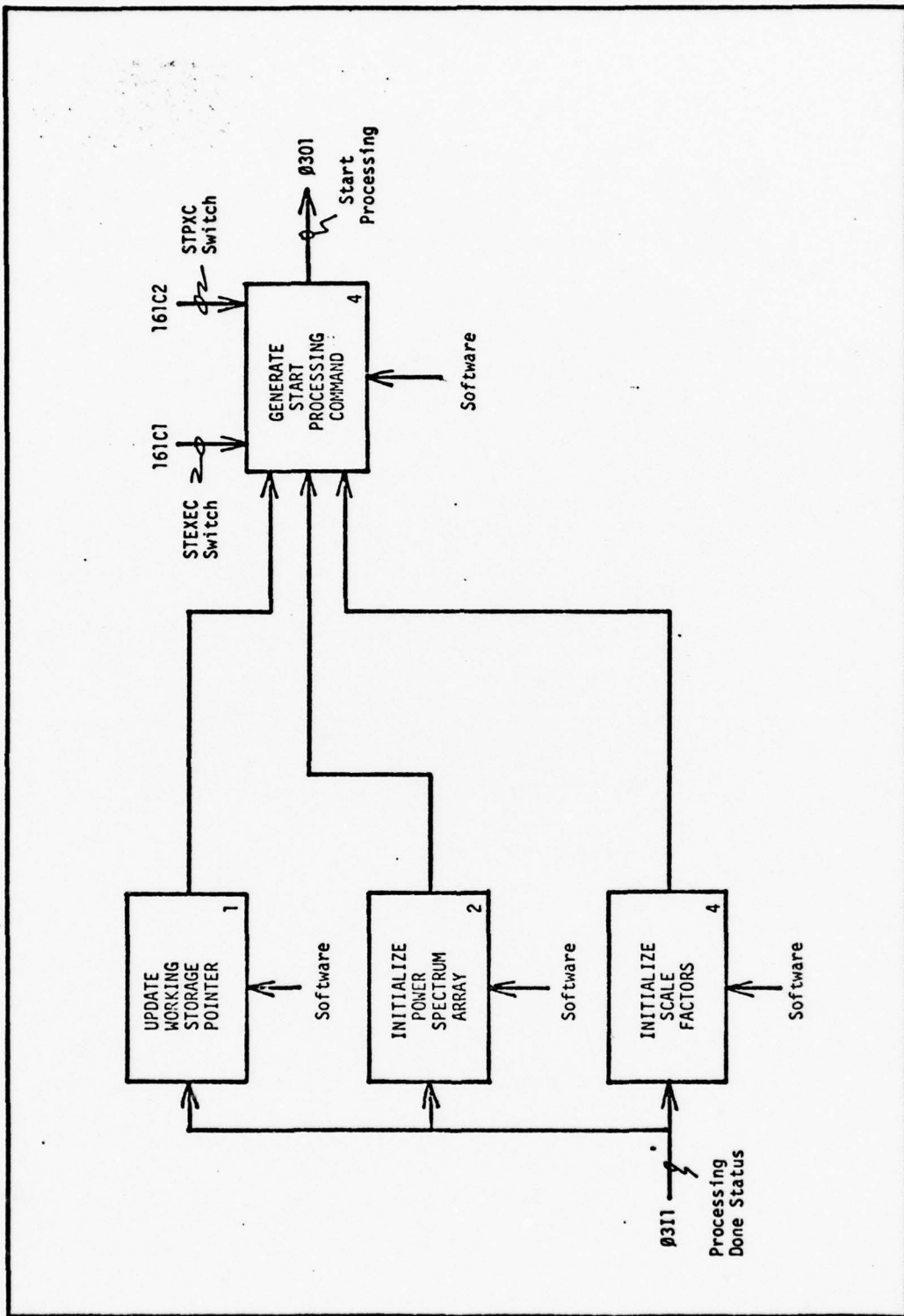


Figure 23. Node 163, Initiate Signal Processing

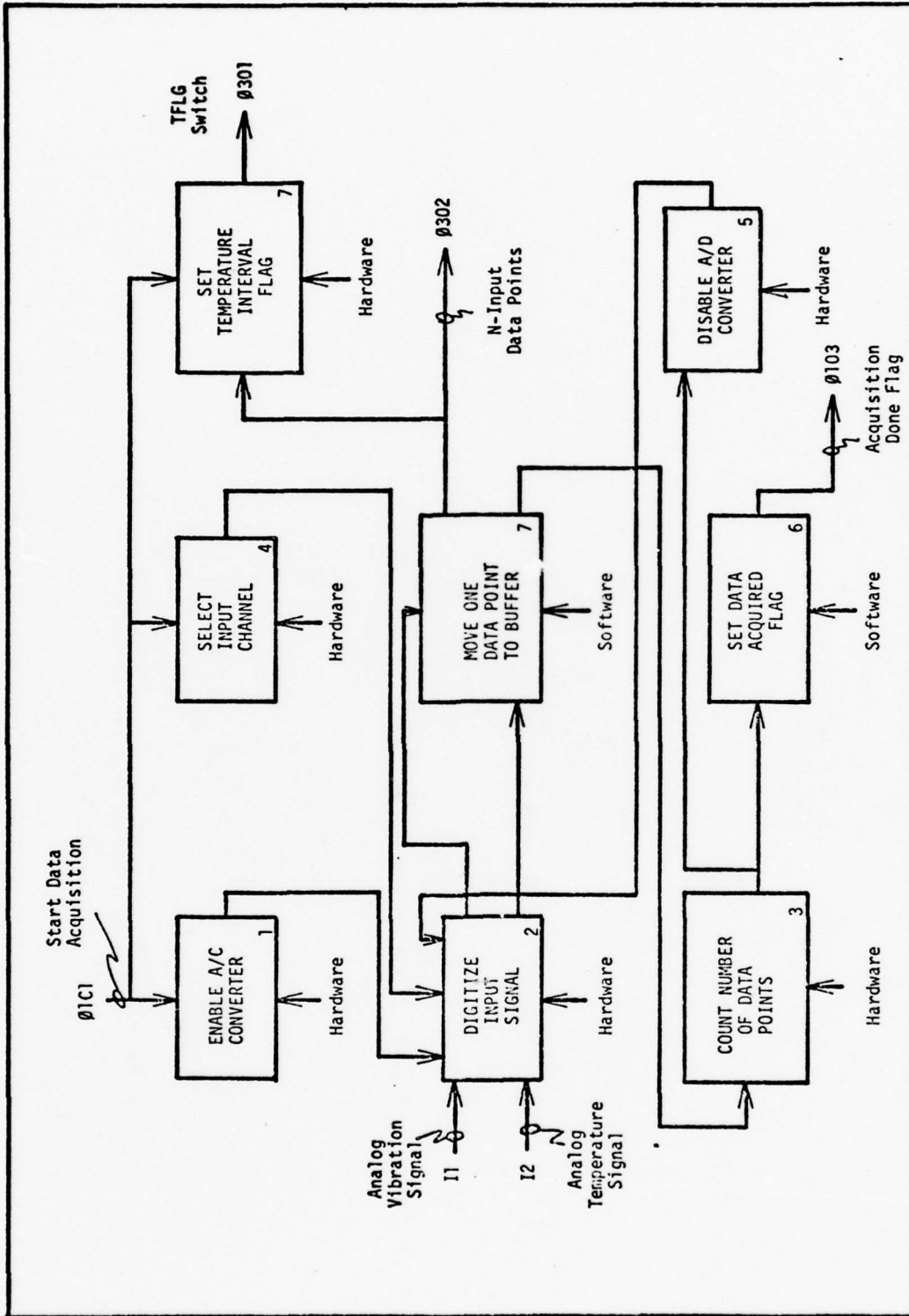


Figure 24. Node A2, Acquire Data

- Output Signal

a. Temperature Interval Flag (Ø301): A decimal number between 0 and 11 is loaded in TFLG. This number depends on the value of the temperature signal at the time of acquisition.

b. N-Input Data Points (Ø302): The data structure for the data buffer storage is given in Figure 25. CDATA is the switch which contains the starting address of the data buffer to be used. ISTR1 and ISTR2 are the starting addresses of the two buffers. Each buffer is N-words long.

c. Acquisition Done Flag (Ø103): Bit-0 of switch ADFLG is set to logic "1" whenever data acquisition of N-points is complete.

Node A3, Process Data (Figure 26). The activities performed by this will be easier to understand if the flow of data through each module is explained. The Input Data (Ø211) is N-points of acquired data which are stored in a buffer area. The starting address of this buffer is available in switch DATA. The fourier transform of these N-points is computed, and the results are stored back in the same buffer area. Switch SFACTR contains the scale factor for the resulting array. The fourier transformed array serves as an input to the Computer Power Spectrum module. This module computes the power spectrum for first M-frequency components (including DC term). The output for this module is M-element power spectrum array and switch CUMSF1, which has the scale factor for this array. The M-element DF array is added to one of the existing CDF versus frequency arrays. The value of switch TFLG is used to select the final array, where the DF array is added. The data structure associated with each module of this node is explained in the subordinate nodes.

Node A32, Computer Fourier Transform (Figure 27). This node takes the input data from the buffer storage, computes the fourier transform

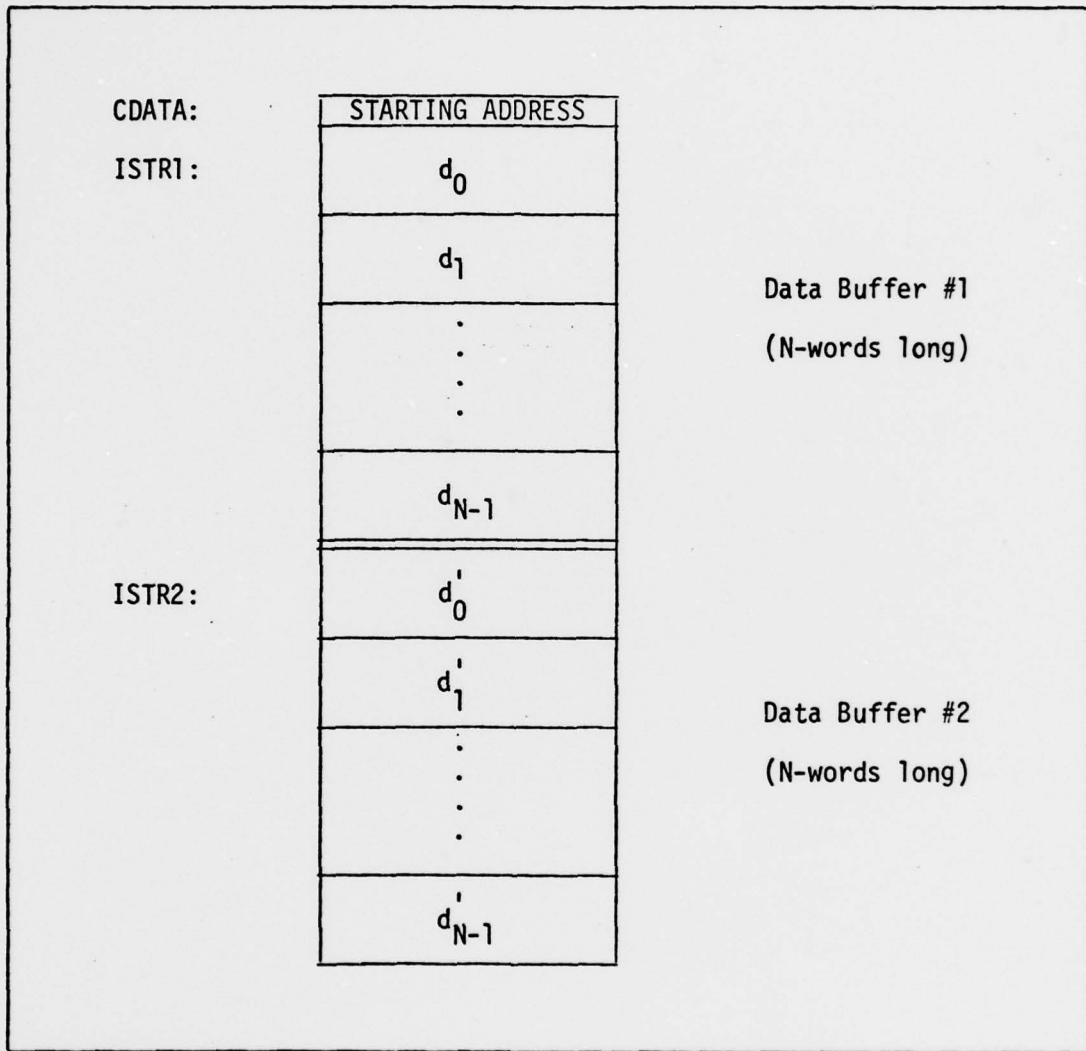


Figure 25. Data Structure for Buffer Storage

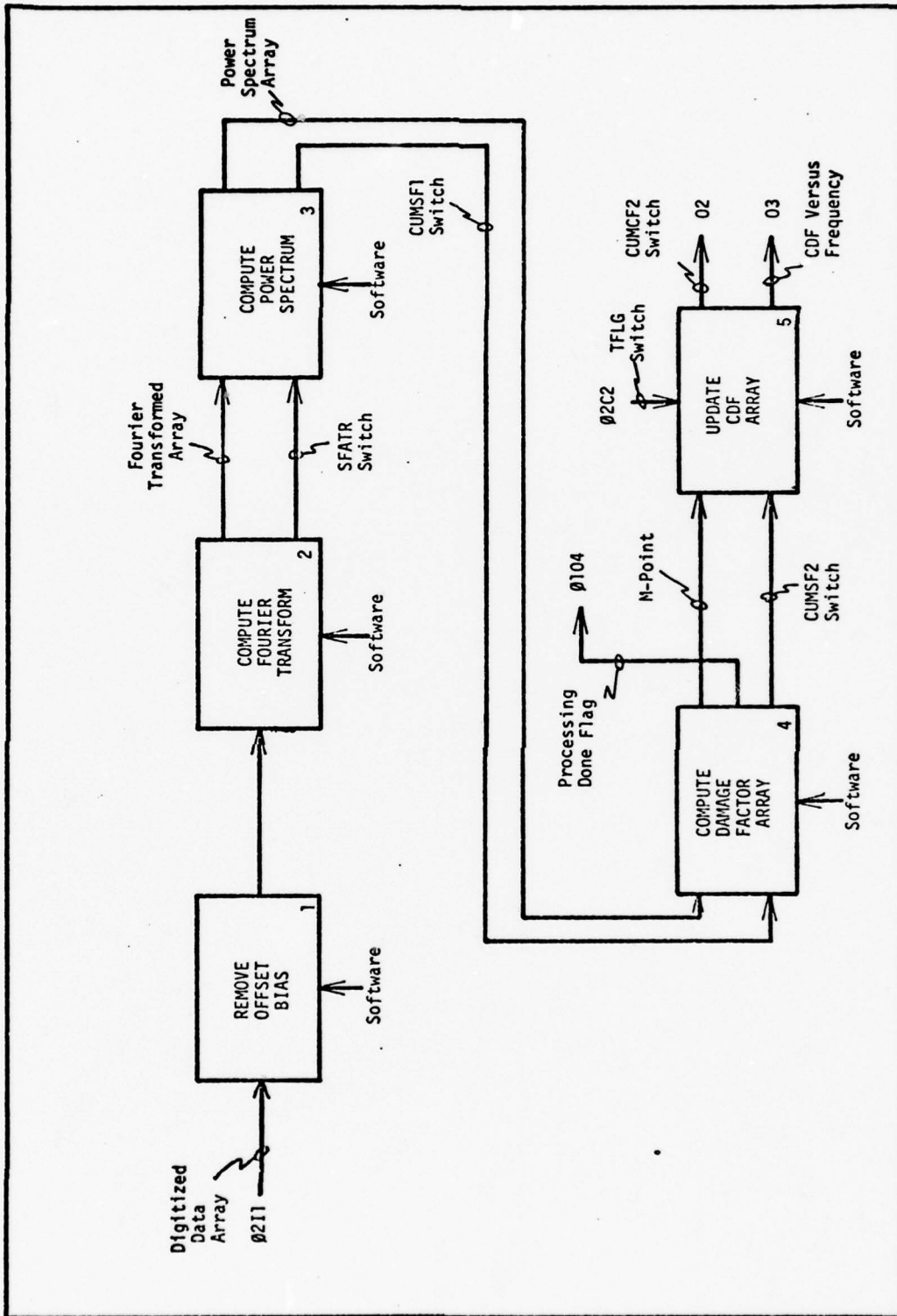


Figure 26. Node A3, Process Data

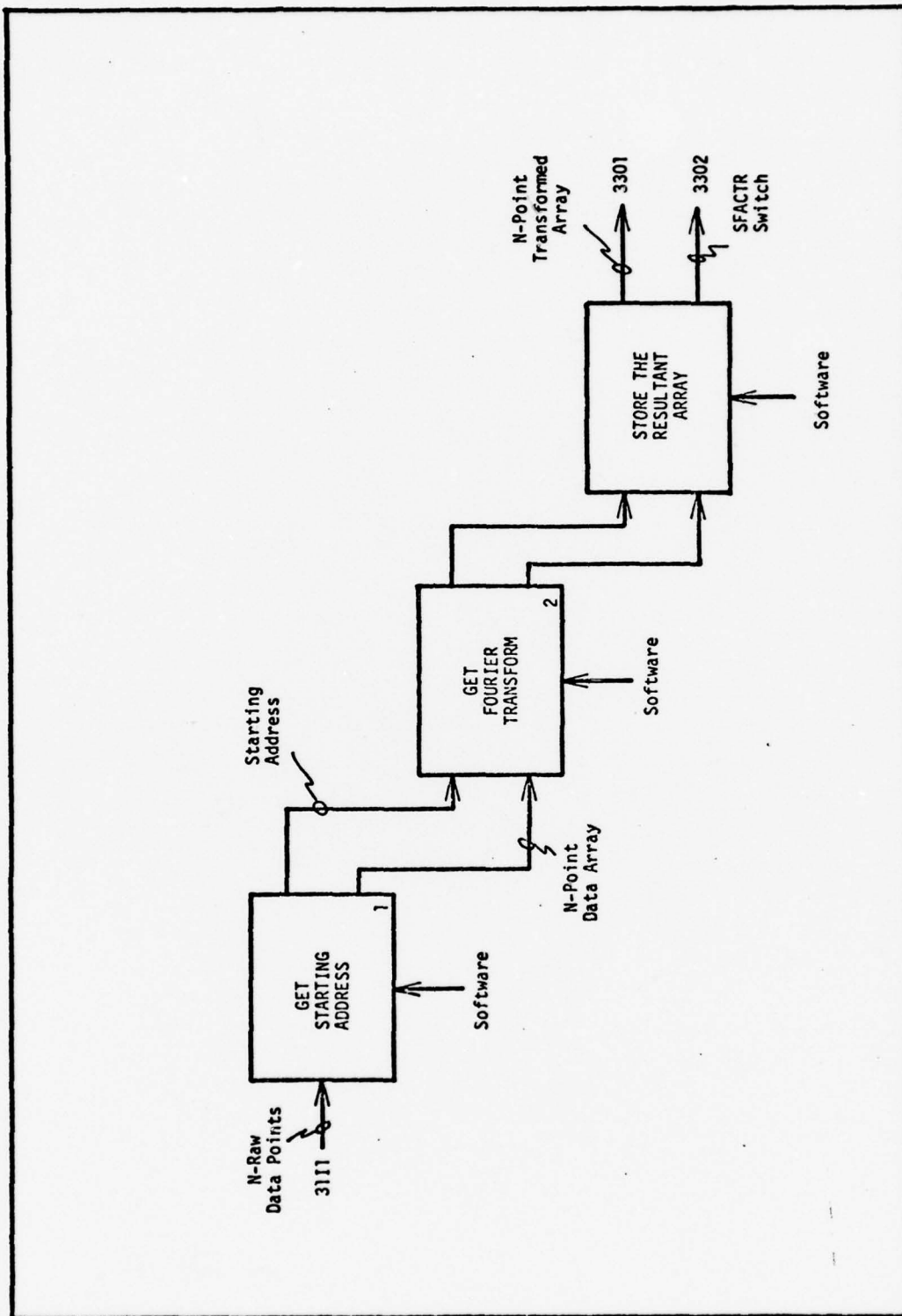


Figure 27. Node A32, Computer Fourier Transform

of the N-data points, and stores the results in the original data buffer array. The input/output data structure is shown in Figure 28. Switch DATA contains the starting address for the input/output data buffer area. Each data buffer area is N-words long.

Node A33, Computer Power Spectrum (Figure 29). This node takes the N-point transform array and computes the power spectrum for first M-frequency components. The resultant M-point array is stored in another temporary data storage area whose starting address is stored in switch SPCTRA. The input/output data structure is shown in Figure 30.

Node A34, Computer Damage Factor (Figure 31). The M-point power spectrum array is input to this node (3311). It takes the starting address from the SPCTRA switch and computes the damage factor for each M-data points in the input array. The resultant DF array is passed on to Node A35.

Node A35, Update CDF Array (Figure 32). This node takes the damage factors computed by node A34 and updates one of the 11 CDF versus frequency arrays. Switch TFLG (02C1) is used to select the starting address of the output array which corresponds to the correct temperature interval. The output data structure for CDF versus frequency arrays is given in Figure 33.

### Summary

In this chapter, the selection of a microcomputer system to implement MIBDAPS was <sup>accomplished</sup> made. Based on the selected microcomputer system, the partitioning of MIBDAPS into hardware or software realizable modules was performed. A detailed design of MIBDAPS, including mechanization of different modules, input/output data structures, and breakdown of different modules into primitives, was carried out using SADT.

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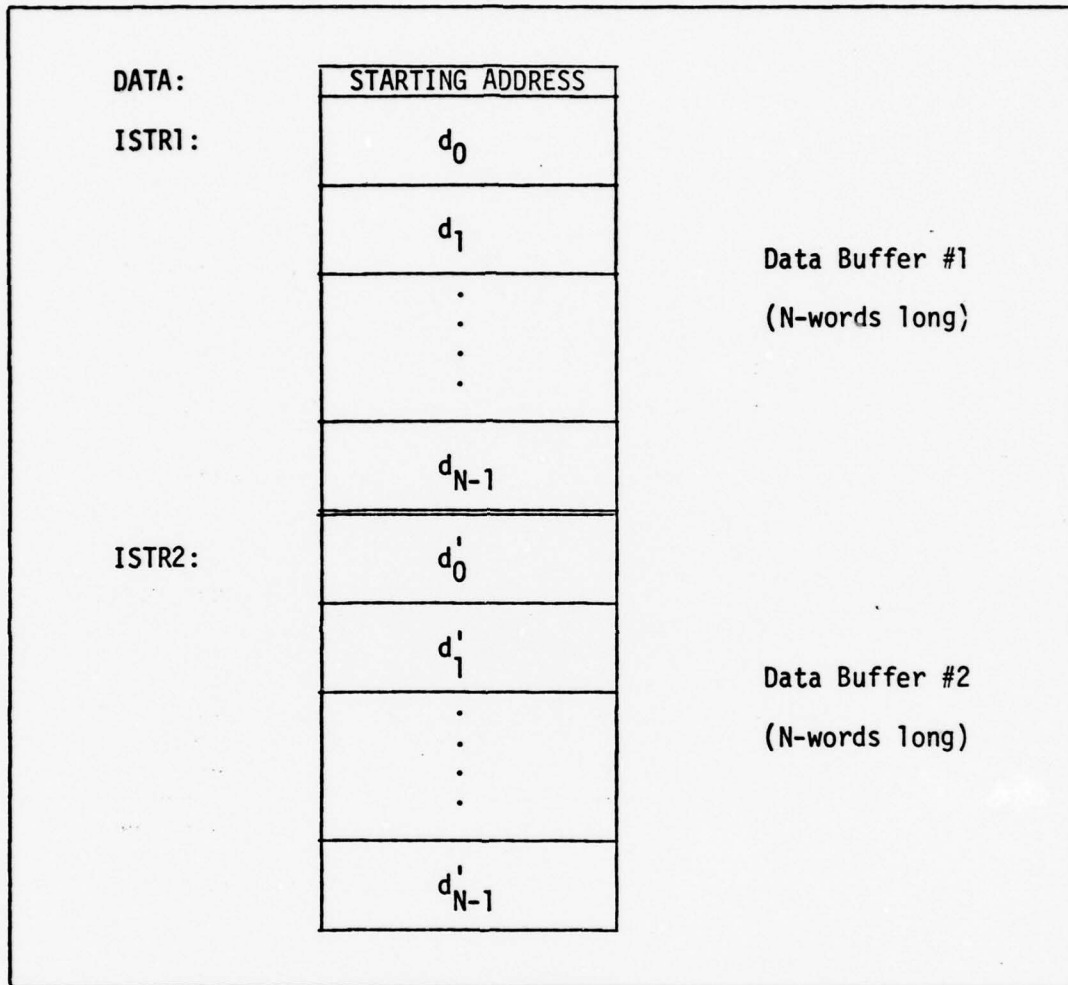


Figure 28. Input/Output Data Storage for Node A32

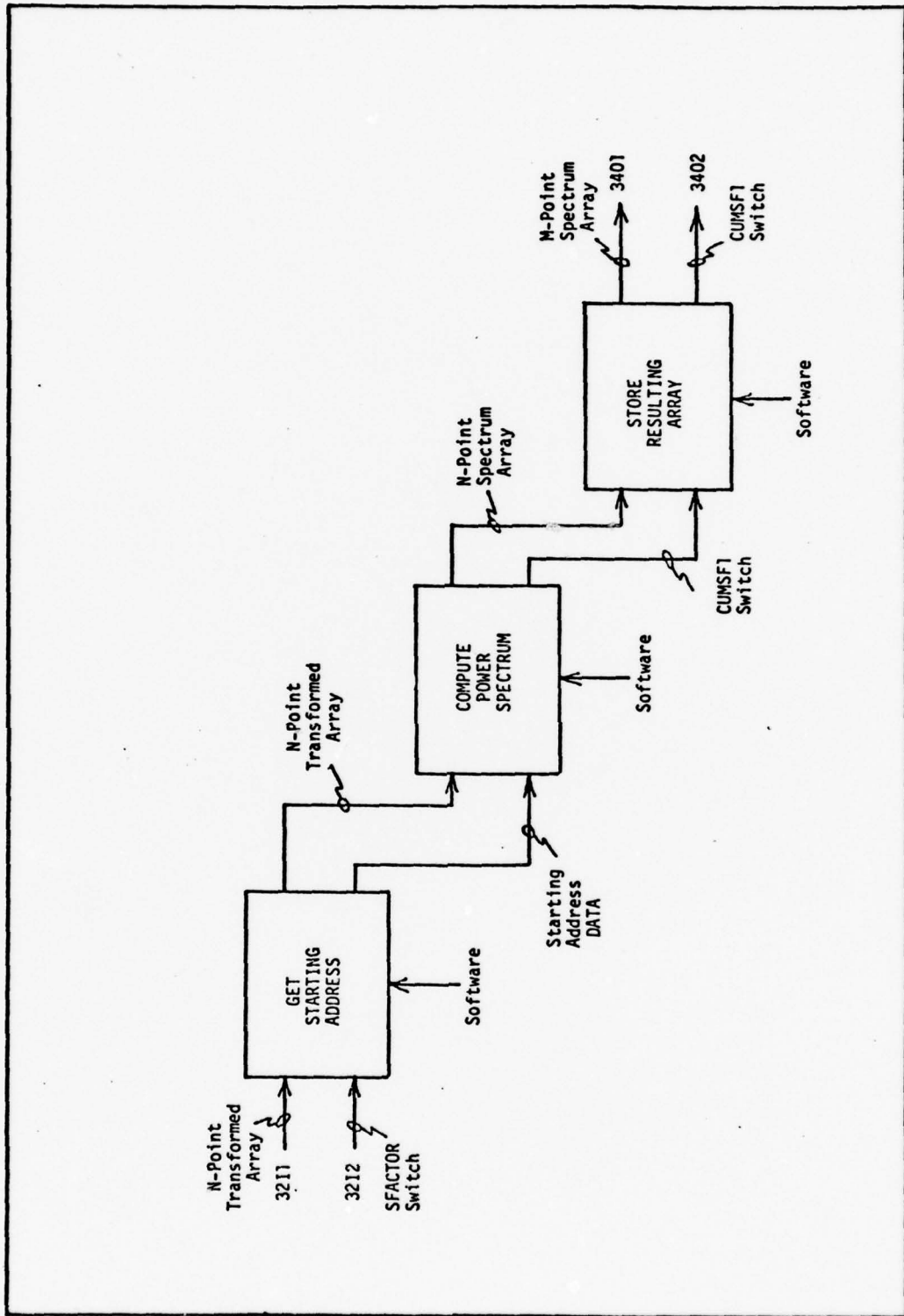


Figure 29. Node A33, Computer Power Spectrum

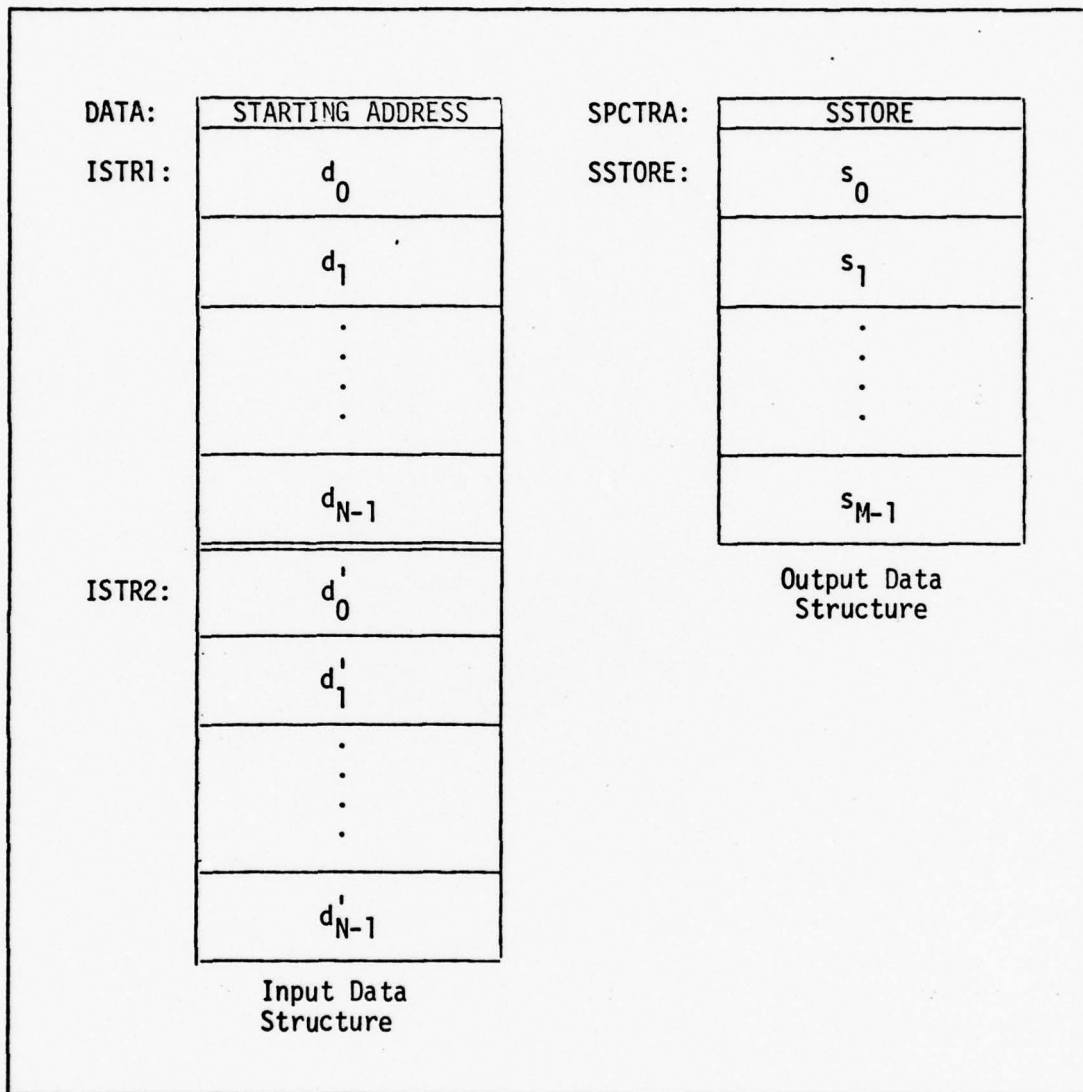


Figure 30. Input/Output Data Structure for Node A33.

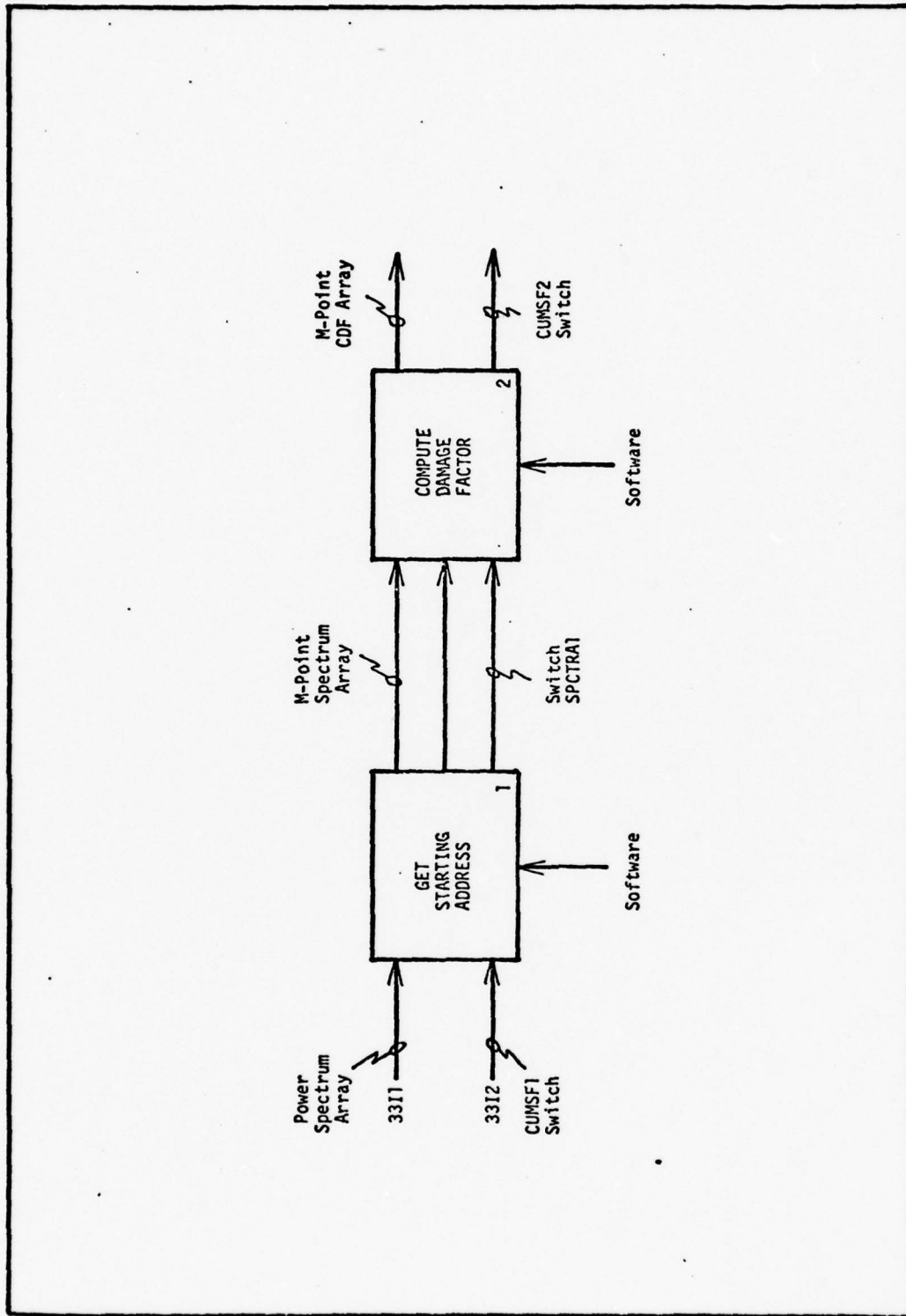


Figure 31. Node A34, Computer Damage Factor

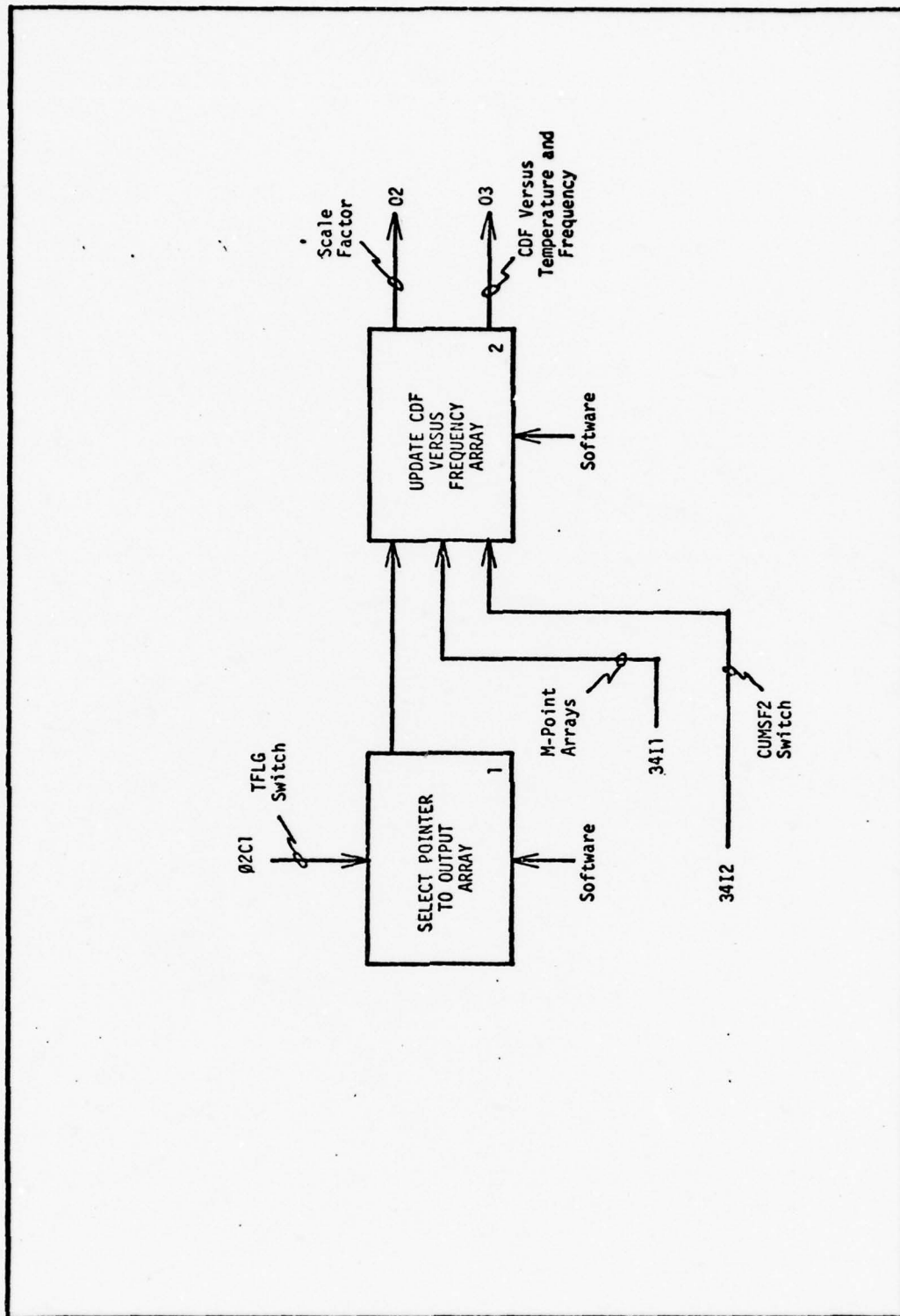


Figure 32. Node A35, Update CDF Array

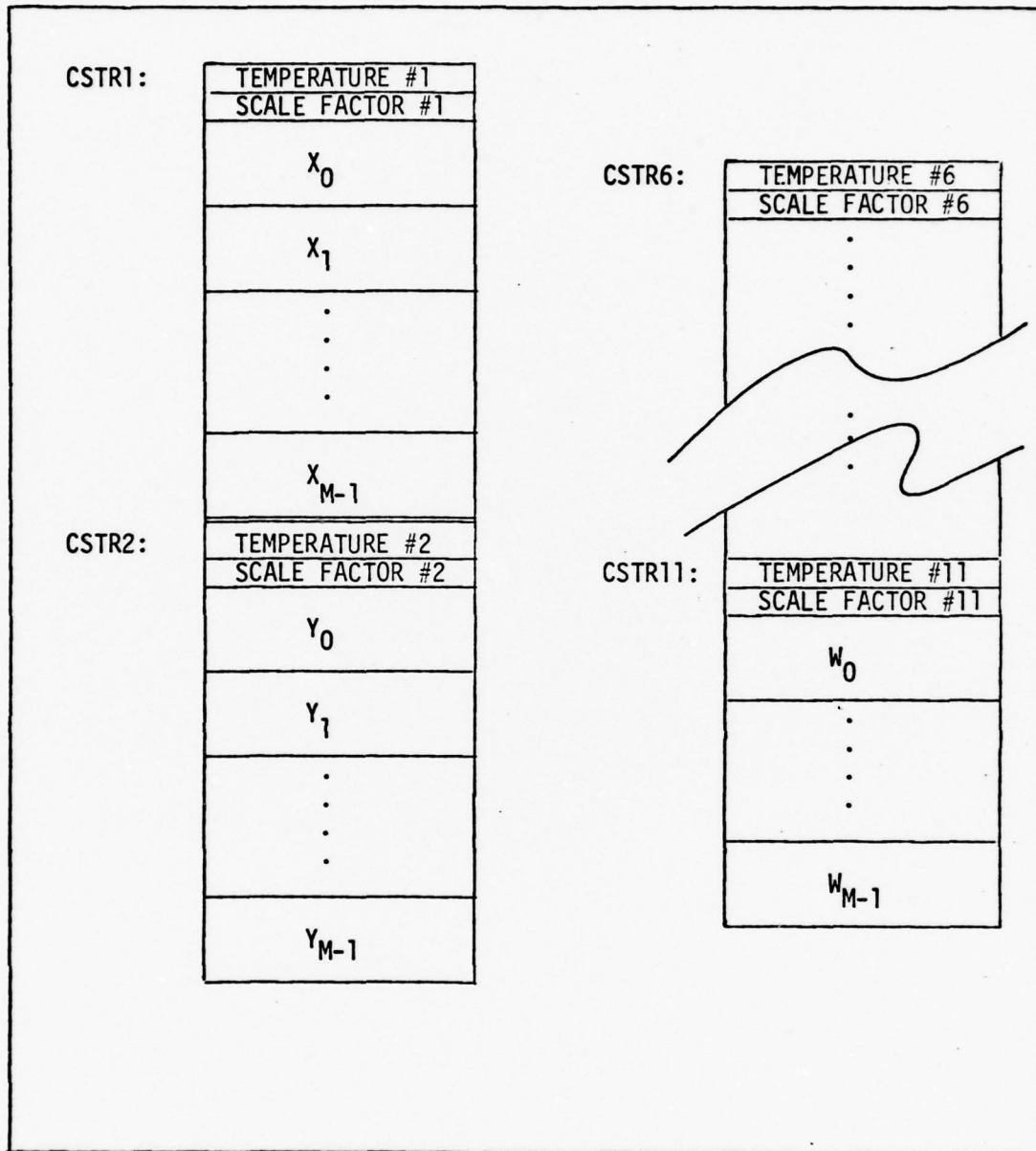


Figure 33. Output Data Structure for Node A35

#### IV. Design and Implementation of Software

The general design of MIBDAPS discussed in Chapter III presented the foundation for system implementation. The division of the overall system into software and hardware modules was also accomplished. In this chapter, the algorithm for each software module is developed. Subsequently, the LSI-11 assembly language code is written for the designed algorithms. The source listing for each software module is included in Appendices E, F, and G. Before the software algorithms are designed, a mathematical analysis of the signal processing function of MIBDAPS is presented. This analysis is performed to determine system parameters and their variation during the signal processing of data.

##### Determination of System Parameters

The development of any software module requires a priori knowledge of the computation taking place within the module and the relationship of input and output parameters. Therefore, a mathematical analysis of the data acquisition and signal processing functions of MIBDAPS is required. This analysis should result in the type and amount of computation required within each software module of MIBDAPS and the different input/output parameters. Because the processing of the acquired data is to be done in real time (Ref Chap I), timing constraints placed on the system would also be determined in this analysis.

Data Acquisition Analysis. The MIBDAPS is required to digitize the analog vibration signal and the analog temperature before any signal processing is carried out. The acquisition analysis for the analog vibration signal is carried out first. This would be followed by the acquisition analysis for the analog temperature signal.

The analog-to-digital conversion of a continuous signal is a two-step process, as shown in Figure 34. The analog signal is first sampled, and then followed by quantization of the amplitudes of the sampled signal (Ref 2:155).

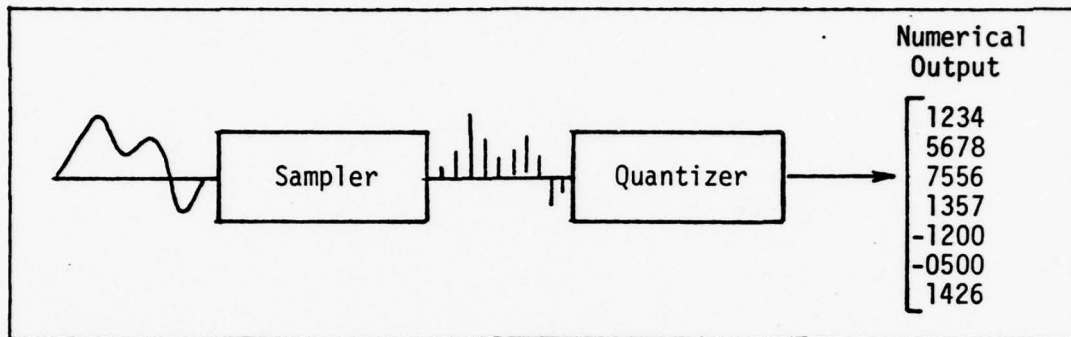


Figure 34. Sampling and Quantization of an Analog (From Ref 2:155)

The limitations and applicability of these two steps will first be considered and, particularly, the consequences arising from the selection of a given sampling rate and number of quantizing levels.

The first problem that arises from sampling a continuous signal is that of aliasing. The nature of this problem is illustrated in Figure 35 which shows that the same set of sampled data points can describe a number of time series histories, which are indistinguishable to the digital computer.

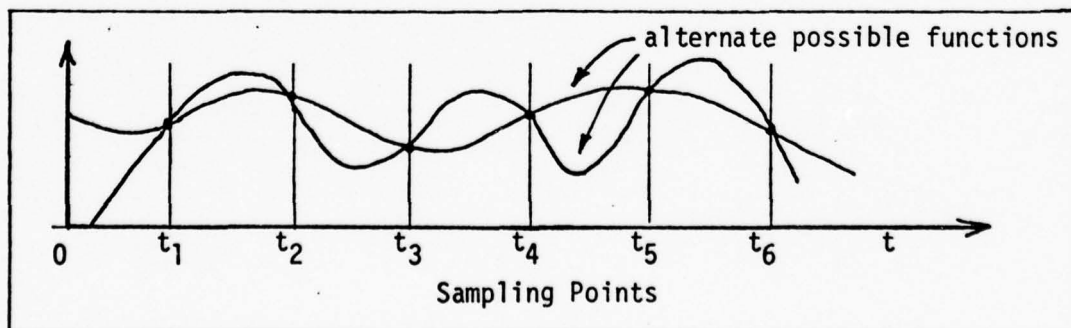


Figure 35. Aliasing (From Ref 2:157)

To overcome this problem, the continuous signal which has a finite band width up to and including B rad/sec should be sampled at  $\pi/B$  sec intervals (Ref 2:160). Therefore, the sampling interval  $\Delta T$  is given as:

$$\Delta T = \pi/B \quad (2)$$

$$B = 2\pi f_m \quad (3)$$

where

$f_m$  = maximum frequency content of the signal

From Eq (2) and Eq (3) we get

$$\Delta T = 1/2f_m \quad (4)$$

The sampling frequency F, which is the reciprocal of  $\Delta T$ , is given by

$$F = 1/\Delta T = 2f_m \quad (5)$$

Equation (5) defines the sampling frequency required to reduce the effect of aliasing. This rate is also known as the Nyquist sampling rate (Ref 6: 225). The analog vibration signal input to MIBDAPS is expected to have a finite band width extending from OHZ to 1000 HZ (Ref Chap I); therefore, Eq (5) specifies that the sampling frequency for this signal should be equal to or greater than 2000 HZ.

The second step in data acquisition is quantization. "Quantization" is defined as the representation of a variable amplitude series of discrete sample values as an equivalent series of discrete numbers representing their amplitude values. This process can only be an approximation because the number of bits in a digital representation is limited, even though the analog signal can assume an infinite number of values. The

numerical values of the quantized variable may be represented by some form of binary code to permit entry into the digital computer (Ref 2:162). Figure 36 shows the transfer characteristics of the quantizer which gives a 3-bit code for an input analog signal. The analog values are quantized by partitioning the continuum into eight discrete ranges. All analog values within a given range are represented by the same digital code, which corresponds to the nominal mid-range value. There is, therefore, an inherent quantization uncertainty of  $\pm \frac{1}{2}$  LSB. The only sure way to reduce this quantization uncertainty is to increase the number of bits for the representation of output and have more quantization levels (Ref 8:69).

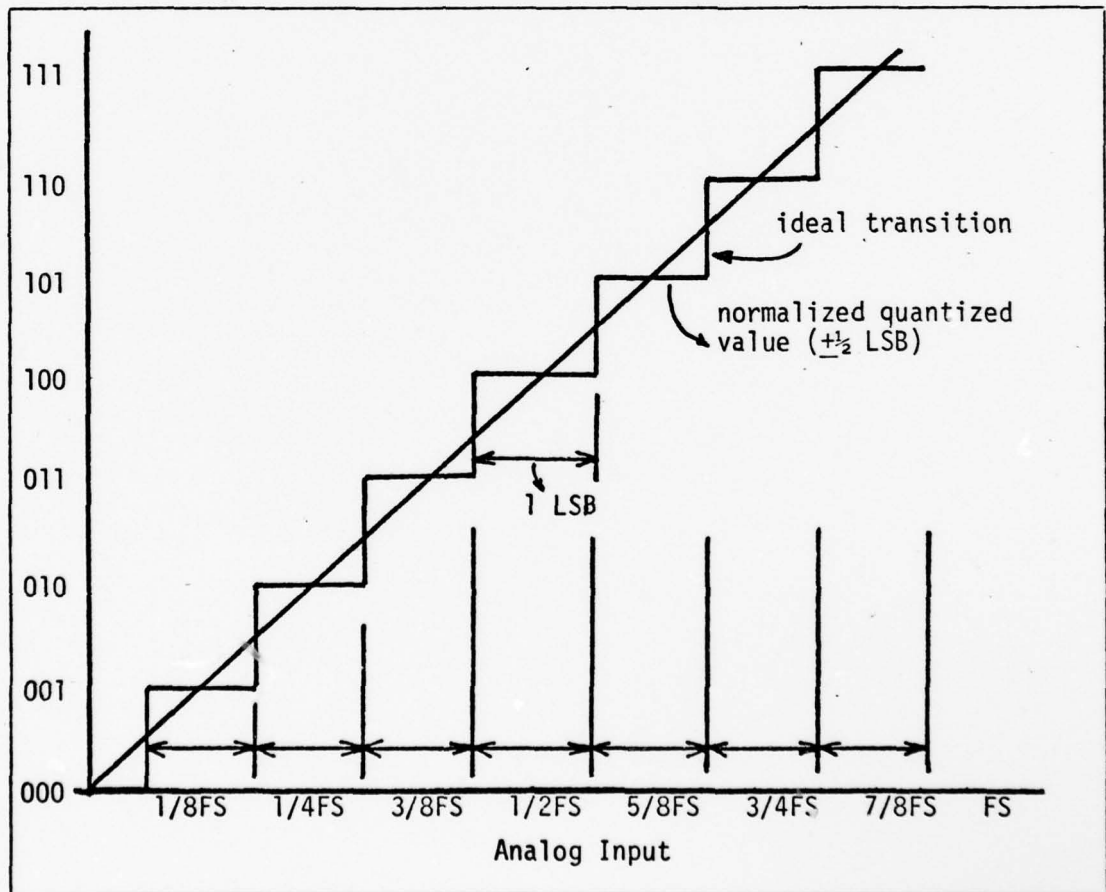


Figure 36. Quantizer Transfer Characteristics (From Ref 8:69)

The ADC used for data acquisition in MIBDAPS is a 12-bit successive approximation converter which has one LSB value equal to 25mv (Ref Chap III). Thus the quantization error, which is  $\pm\frac{1}{2}$  LSB, is equal to  $\pm 12.5$ mv. The analog vibration and temperature signal will have a full-scale value of 5 volts; therefore, the quantization error will be  $\pm 0.25$  percent of the full-scale value.

Signal Processing Analysis. The signal processing function of MIBDAPS has been broken down into three sub-functions (Ref Chap III, Node A0). The mathematical analysis of the signal processing in each subunit will now be carried out.

1. Processing an Analog Vibration Signal. The power spectral density of the digitized vibration signal needs to be computed first before the damage factor for each frequency of interest is computed (Ref Chap III). This can be accomplished by first determining the frequency domain spectrum of the input time domain signal. The Fourier transform is a method that converts a signal from time to frequency domain. Since the input signal is a sampled version of the continuous time domain signal, the discrete fourier transform (DFT) method was used for the transformation. Equation (6) gives the expression for the conversion of N sampled data points from time domain to frequency domain:

$$F(n) = \sum_{k=0}^{N-1} x(k) \exp\left(\frac{-2\pi n k}{N}\right), \quad n = 0, 1, 2 \dots N-1 \quad (6)$$

where  $F(n)$  is the nth point in the frequency domain and  $n(k)$  the kth point in the time domain. The  $x(k)$  may be complex and the  $F(n)$ 's are always complex. We can rewrite Eq (6) as

$$F(n) = \sum_{k=0}^{N-1} x(k)W^{nk}, \quad n = 0, 1, 2 \dots N-1 \quad (7)$$

where

$$W = \exp\left(\frac{2\pi i}{N}\right) \quad (8)$$

Figure 37 shows the required frequency domain spectrum (Ref Chap I) from 0 HZ to 1000 HZ. The frequency interval  $\Delta f$ , which is 100 HZ, will be used to determine the number of time domain samples required for the Fourier transformation. Equation (9) gives the relationship between  $\Delta f$  and the sampling interval  $T$  (Ref 3:87).

$$\Delta f = \frac{1}{T} \quad (9)$$

or 
$$T = \frac{1}{\Delta f} \quad (10)$$

also 
$$T = M \times \Delta T \quad (11)$$

therefore 
$$M = \frac{T}{\Delta T} \quad (12)$$

Equation (12) gives the number of sampled points  $M$  required to get a frequency spectrum shown in Figure 37. The relationship between the signal in the frequency domain and time domain is shown in Figures 37 and 38, respectively. From Eq (4)  $\Delta T$  is computed equal to 0.5 msec, and Eq (10) gives  $T$  equal to 10 msec. Therefore, the number of sampled points  $M$  required to get the desired frequency spectrum is 20, which is calculated using Eq (12). Once the frequency spectrum of the time domain signal is computed, the next step in signal processing is to compute the power spectral density for each frequency component. After the Fourier transformation, each frequency component  $F(n)$  in Eq (6) is a complex number as given in Eq (13):

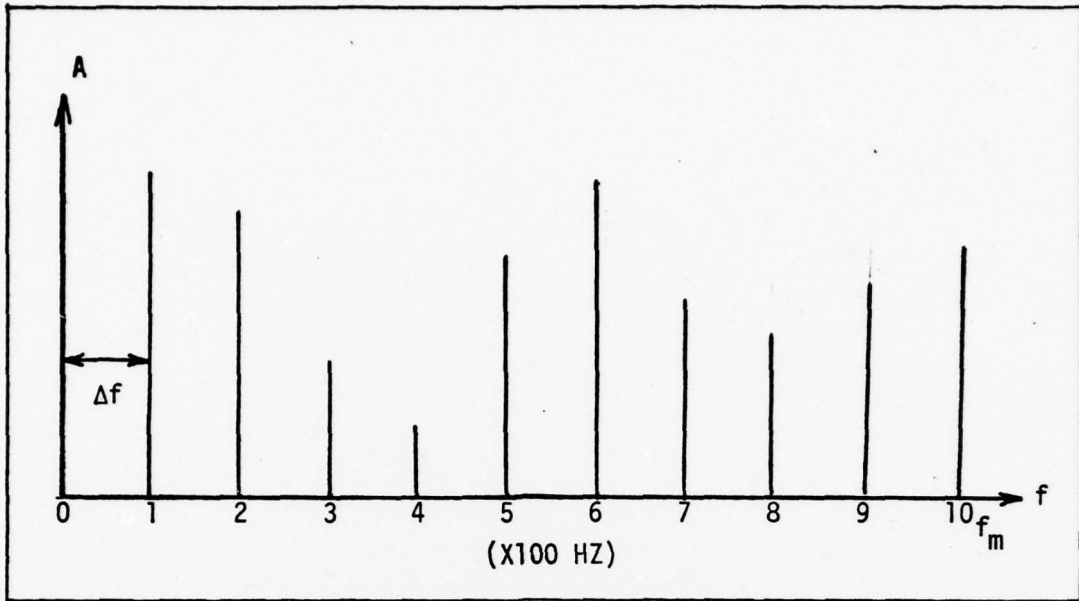


Figure 37. Frequency Domain

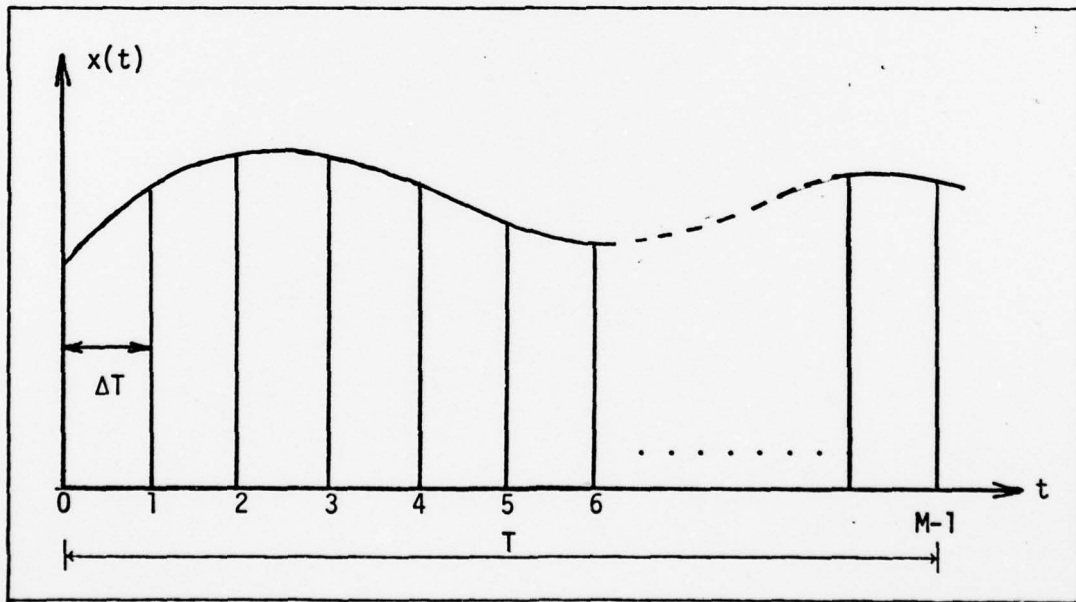


Figure 38. Time Domain

$$F(n) = x_n + jy_n \quad (13)$$

the power spectral density can be calculated using Eq (14):

$$A(n) = (x_n)^2 + (y_n)^2 \quad (14)$$

However, the nature of the input signal is random (Ref 13) and to get a good representation of the power spectral density a mean value for the P set of spectral density points is taken (Ref 2:195). This is expressed mathematically in Eq (15):

$$\tilde{A}(n) = \frac{1}{P} \sum_{k=0}^{P-1} \frac{A(n)}{K}, \quad n = 0, 1, \dots, M-1 \quad (15)$$

where subscript k denotes a unique set of M spectral density values, and  $\tilde{A}$  is the mean power spectral density.

Once the mean power spectral density is computed, the damage factor at each frequency component can be evaluated using Eq (16) (Ref 13):

$$DF(n) = \frac{k[\tilde{A}(n)]^{3.2}}{(f_n)^{5.4}}, \quad n = 0, 1 \dots M-1 \quad (16)$$

where

$$\bar{A}(n) = \sqrt{\tilde{A}(n)} \quad (17)$$

Here  $f_n$  is the frequency at which the damage factor is being calculated.

From Eqs (16) and (17) the following relationship is derived:

$$DF(n) = \frac{k[\tilde{A}(n)]^{3.2}}{(f_n)^{5.4}}, \quad n = 0, 1 \dots M-1 \quad (18)$$

or

$$DF(n) = \frac{k[\tilde{A}(n)]^{1.6}}{(f_n)^{5.4}}, \quad n = 0, 1 \dots M-1 \quad (19)$$

Equation (19) eliminates the need to compute the root mean square (RMS) of the average power spectral density. The CDF, which is a function of frequency and temperature, is obtained by continuous summation of DF values over a predetermined interval. Equation (20) gives the required relationship:

$$\text{CDF}(t_i, m) = \sum_{j=1}^R \text{DF}_j(m), \quad m = 0, 1 \dots M-1 \quad (20)$$

Here,  $t_i$  represents the  $i$ th temperature interval for which the CDF is computed, and subscript  $j$  indicates the number of DF value sets that have been summed together.

2. Analog Temperature Signal. The temperature signal is sampled every time a set of 20 data points of the vibration signal is processed and the power spectral density has been computed. Since  $P$  set of power spectral density values are summed [Eq (15)], there are  $P$  temperature values that can be summed up to give an average temperature value  $\bar{T}$ . The average temperature value would be available when the DF is computed and the final output array (CDF versus frequency and temperature) is to be updated. The following equation gives the mathematical relationship for  $\bar{T}$  (average temperature):

$$\bar{T} = \frac{1}{P} \sum_{k=0}^{P-1} T(k) \quad (21)$$

This concludes the mathematical analysis of the data acquisition and signal processing functions of MIBDAPS. The software modules to implement different functional activities will now be developed.

## Implementation of Software Modules

The software modules which will be used to perform different functional activities are developed in subsequent paragraphs.

Fourier Transform Module. The Fourier transform is a method to convert the signal spectrum from the time domain to a spectrum in the frequency domain. Since MIBDAPS uses digitized data for processing, the DFT algorithm can be used to transform the digitized vibration signal from the time domain to the frequency domain. Equation (7) gives the mathematical relationship to convert a digitized signal to a frequency domain. A computer program to perform the summation given in Eq (7) can be written, but it turns out that this can be a very slow and time-consuming process if the equation is implemented as given (Ref 6:267). The reason for this inefficiency is that the transformation of N-data points requires  $N^2$  complex multiplications and  $N(N-1)$  complex additions (Ref 3:151). To speed up the DFT implementation on the computer, Cooley and Tukey (Ref 3:151) developed a fast Fourier transform (FFT) algorithm, which for  $N = 2^x$  points reduces the complex multiplication to  $Nx/2$  and the complex addition to  $Nx$  (Ref 3:151).

Appendix C gives the theoretical development of the base 2 FFT algorithms. It was determined earlier that 20 points of DFT are required to get the required frequency spectrum of the vibration signal.  $N = 20$  does not satisfy the relationship  $N = 2^x$  where  $x$  is an integer; therefore, a base 2 FFT algorithm would not be used. FFT algorithms for arbitrary factors can also be written (Ref 3:184). An FFT algorithm for 20-point DFT was written, but it was found that 125 complex multiplications and 108 complex additions are required. On the other hand, a 32-point FFT algorithm requires 80 complex multiplications and 160 complex additions,

so it was decided to use a 32-point FFT algorithm. This would give a frequency spectrum from 0 HZ to 1500 HZ, and by picking up the first 11 frequency components, the desired spectrum for MIBDAPS could be obtained. Figure 39 shows the relationship between the time domain signal and the resultant frequency domain signal. The sampling frequency  $F$  should now be 3200 HZ [Ref Eq (5)], and the sampling period  $\Delta T$  becomes equal to 0.3125 msec.

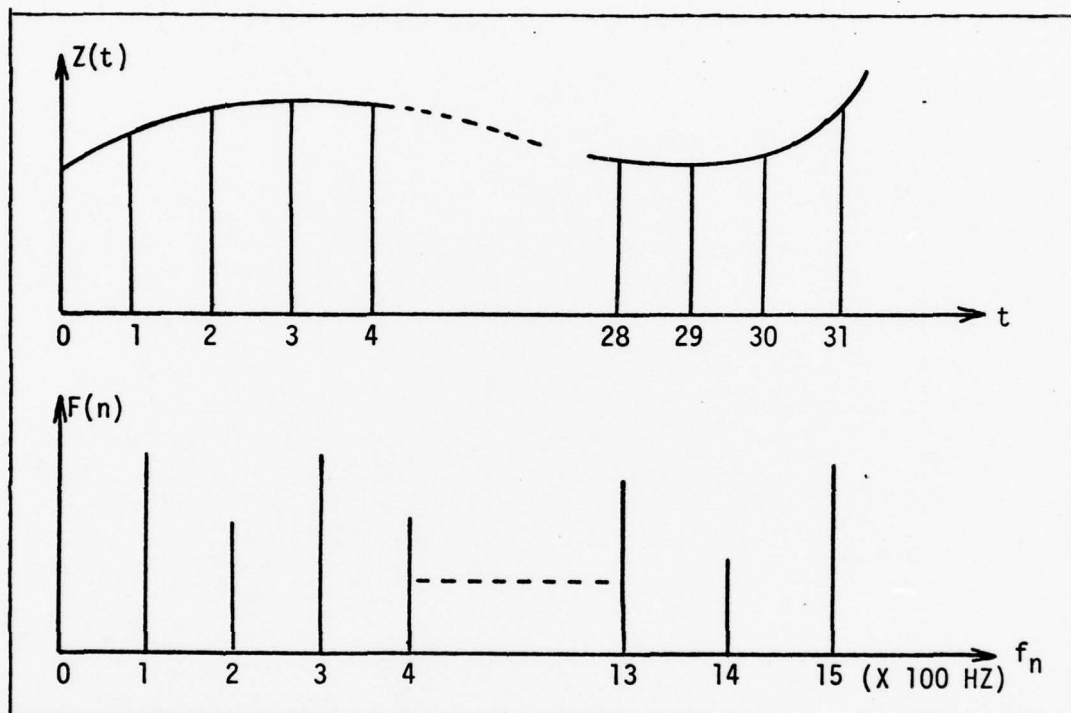


Figure 39. Time and Frequency Domain Signals for 32 Samples

The FFT algorithm developed in Appendix C can handle both real and complex input data. However, the input to MIBDAPS is real data, and computation time would be halved if  $2N$  points of real data are treated as  $N$  points of complex data. This would enable the use of an  $N$ -point FFT algorithm. The theoretical background for transforming  $2N$  points of real data using an  $N$ -point FFT is given in Appendix D; this algorithm was eventually

used for implementation of the Fourier transform module. The description of the algorithm is given in the following paragraphs.

1. Algorithm to Implement Fourier Transform Module. Figure 40 shows the flow chart of the algorithm that was implemented. The input data array is assumed to have 32 real data points in contiguous memory locations. Since the FFT algorithm described in Appendix D would be implemented, it is assumed that the input array has 16 complex data points (all odd points are treated as real points and all even points are treated as complex points). The bit-inversion subroutine is called first so that the data elements are bit-reversed before being processed by the FFT subroutine. This is done to ensure that the output of the FFT subroutine would be in the correct order. The FFT computation is performed on the bit-inversed array; the resultant array would be 16 complex Fourier transformed points. To get the correct transformed array for real data points, the post-processing computation (Ref Appendix D) is performed; this would give the 16 complex values corresponding to frequencies from 0 HZ to 1600 HZ. Although a 32-point real data array should result in 32 complex point Fourier transformed arrays, the algorithm being implemented only gives 16 complex Fourier transformed points. This discrepancy is because the other 16 complex transformed points are the complex conjugates of these points and are obtainable easily. However, these points are not required for further processing in MIBDAPS and, therefore, are not computed.

2. Description of Different Subroutines.

a. Bit-Inversion Subroutine. The flow diagram for the bit-inversion (BIV) subroutine is given in Figure 41. The starting address of the input data array is stored in pointer DATA. The five LSB of the

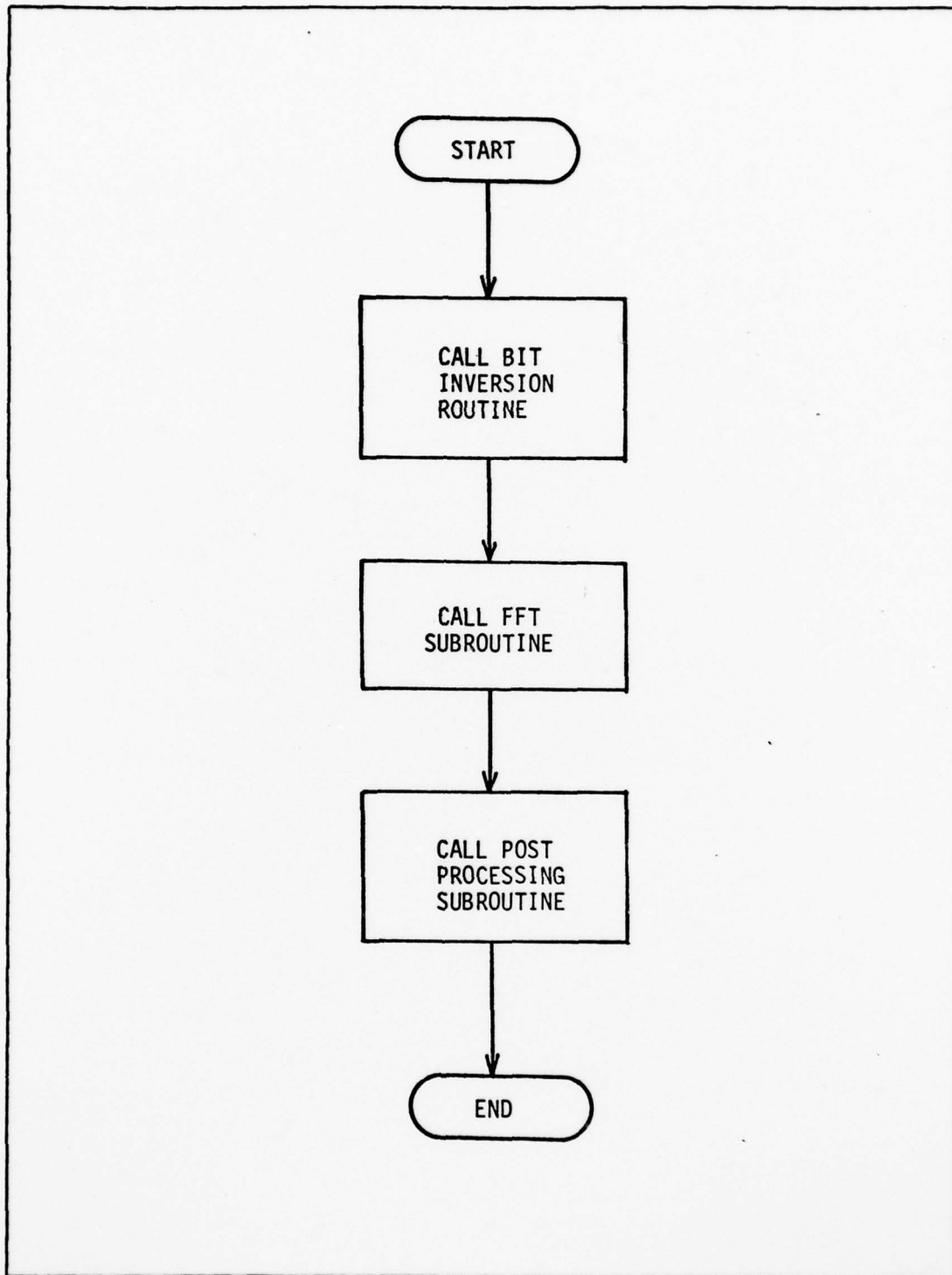


Figure 40. Flow Diagram from Fourier Transform Module

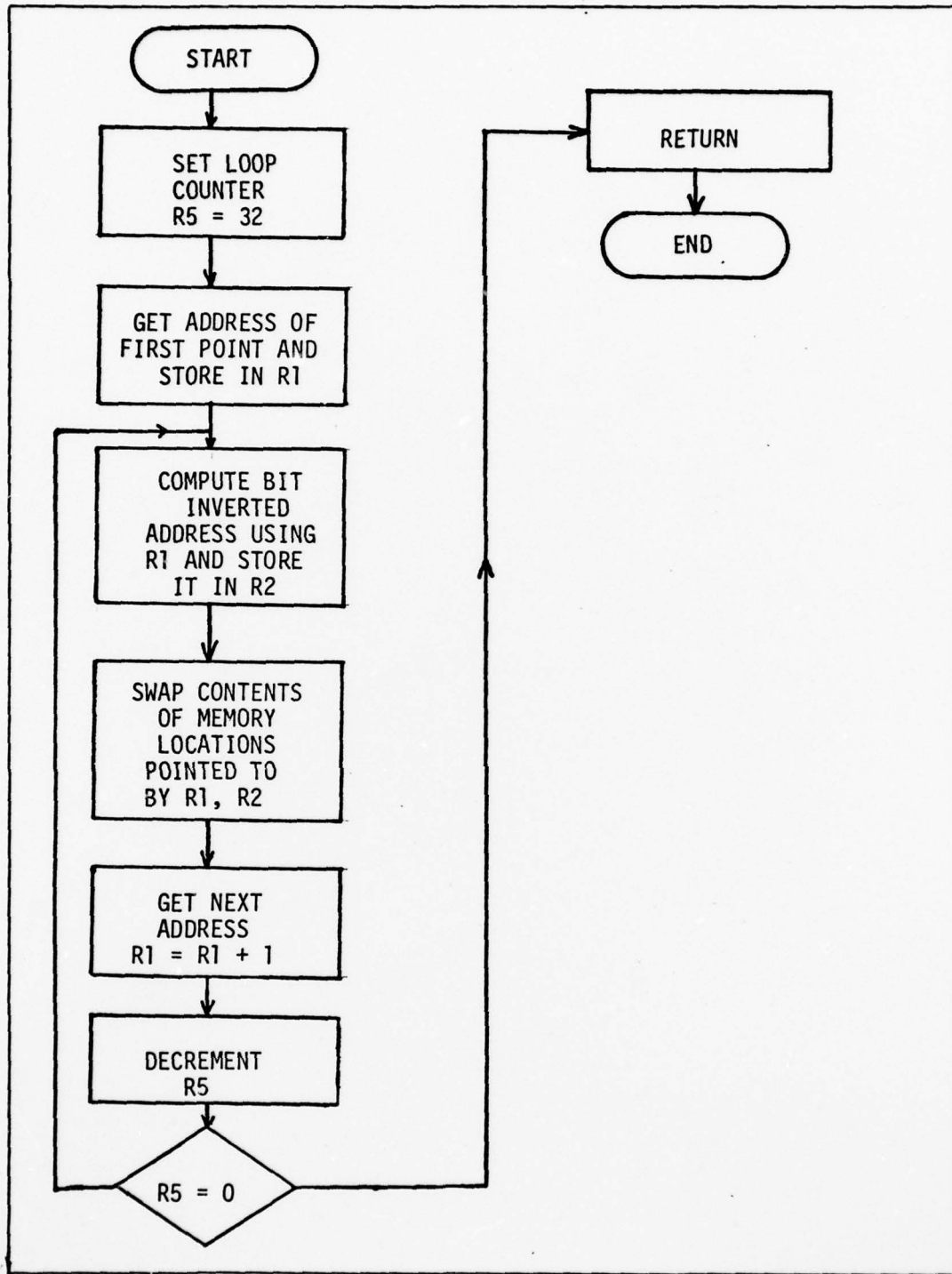


Figure 41. Flow Diagram for Bit-Inversion Subroutine

address for any data elements are used to compute the bit-inverted address. These two memory location contents are then swapped. This process is repeated for all data elements, and the resulting array is in bit-reversed order.

b. FFT Subroutine. The transformation of input data is carried out in four steps (or passes). The manipulation of data within this subroutine follows the scheme given in the signal flow graph shown in Figure C-2, Appendix C. During the first pass, there are no complex multiplications (Ref Table C-I, Appendix C); therefore, the following equations need to be implemented:

$$R(m)' = R(m) + R(n) \quad (22)$$

$$I(m)' = I(m) + I(n) \quad (23)$$

$$R(n)' = R(m) - R(n) \quad (24)$$

$$I(n)' = I(m) - I(n) \quad (25)$$

The letter within the parenthesis signifies a pair of points which are used in the calculation.

The second, third, and fourth pass require complex multiplications because of the  $W_N^{nk}$  factor. The equations implemented in these passes are:

$$X(m) = X(m) + X(n)W^Y \quad (26)$$

$$X(n) = X(m) + X(n)W^Z \quad (27)$$

Here

$$X(m) = R(m) + jI(m)$$

$$X(n) = R(n) + jI(n)$$

$$W^Y = \exp\left(\frac{-2\pi in}{N}\right) = W^Z$$

Also, the exponents of  $W$ ,  $y$  and  $z$ , are  $N/2$  apart, such that the following relationship holds:

$$W^y = -W^{N/2+y} = -W^z \quad (28)$$

Equations (26) and (27) can be rewritten using Eq (28).

$$X(m) = X(m) + X(n)W^y \quad (29)$$

$$X(n) = X(m) - X(n)W^y \quad (30)$$

or

$$R(m) + jI(m) = R(m) + jI(m) + [R(n) + jI(n)](\text{Cos}y + j\text{Sin}y) \quad (31)$$

$$R(n) + jI(n) = R(m) + jI(m) - [R(n) + jI(n)](\text{Cos}y + j\text{Sin}y) \quad (32)$$

Collecting terms and rearranging Eqs (31) and (32), the following relationship evolves:

$$R(m)' = R(n)\text{Cos}y + I(n)\text{Sin}y + R(m) \quad (33)$$

$$R(n)' = -R(n)\text{Cos}y - I(n)\text{Sin}y + R(m) \quad (34)$$

$$I(m)' = -R(n)\text{Sin}y + I(n)\text{Cos}y + I(m) \quad (35)$$

$$I(n)' = R(n)\text{Sin}y - I(n)\text{Cos}y + I(m) \quad (36)$$

Equations (33), (34), (35), and (36) are finally used in computation during the second, third, and fourth pass. The flow diagram for the FFT subroutine is given in Figure 42. The flow diagram shows a call to subroutine SCALE; this subroutine is used to ensure that the data array is scaled to avoid any arithmetic overflow during computations in any pass. The flow diagram for the subroutine SCALE is shown in Figure 43. Before the subroutine is called, it is assumed that the starting address is loaded in Register R0, and the number of points in the array is stored in Register R1.

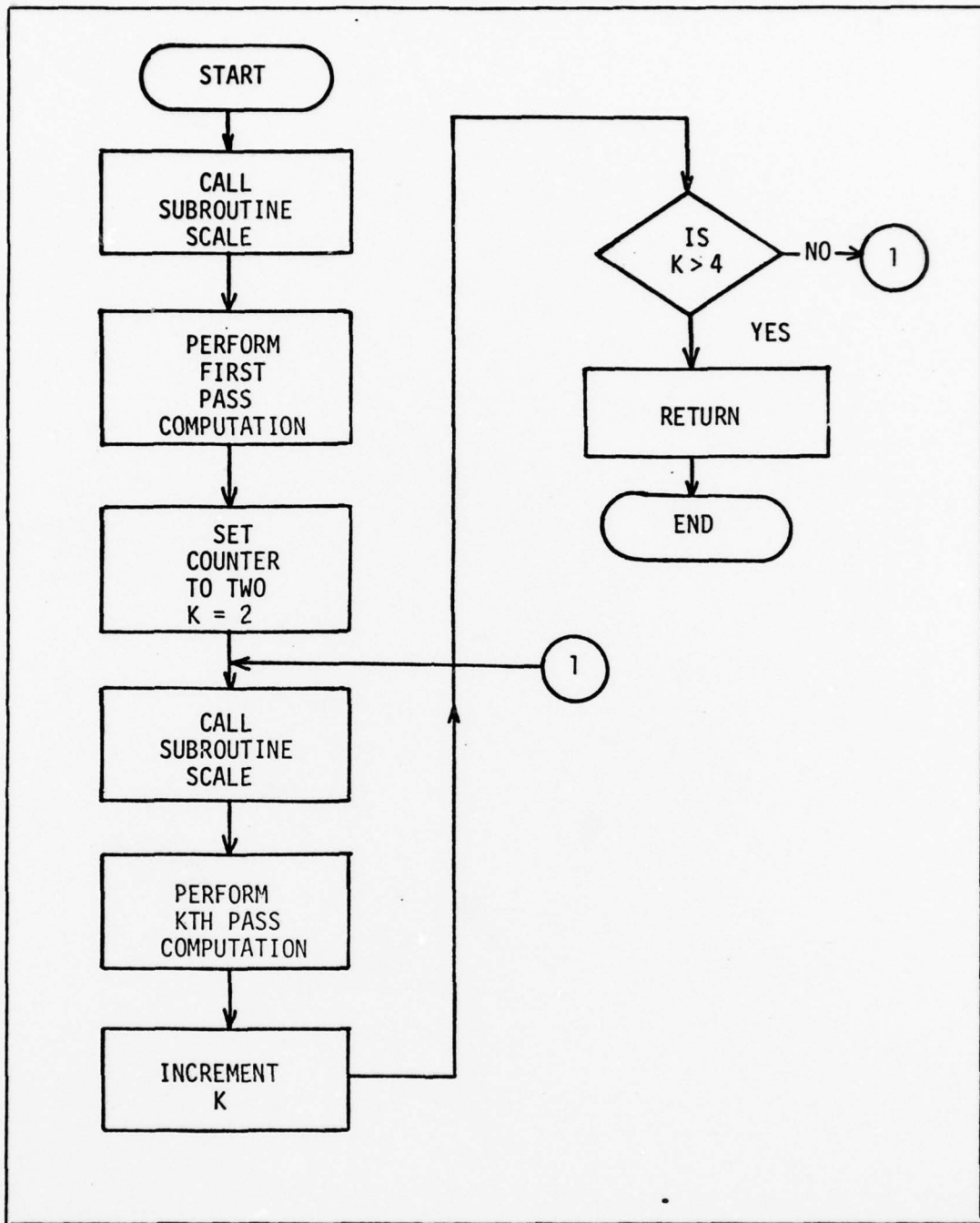


Figure 42. Flow Diagram for FFT Subroutine

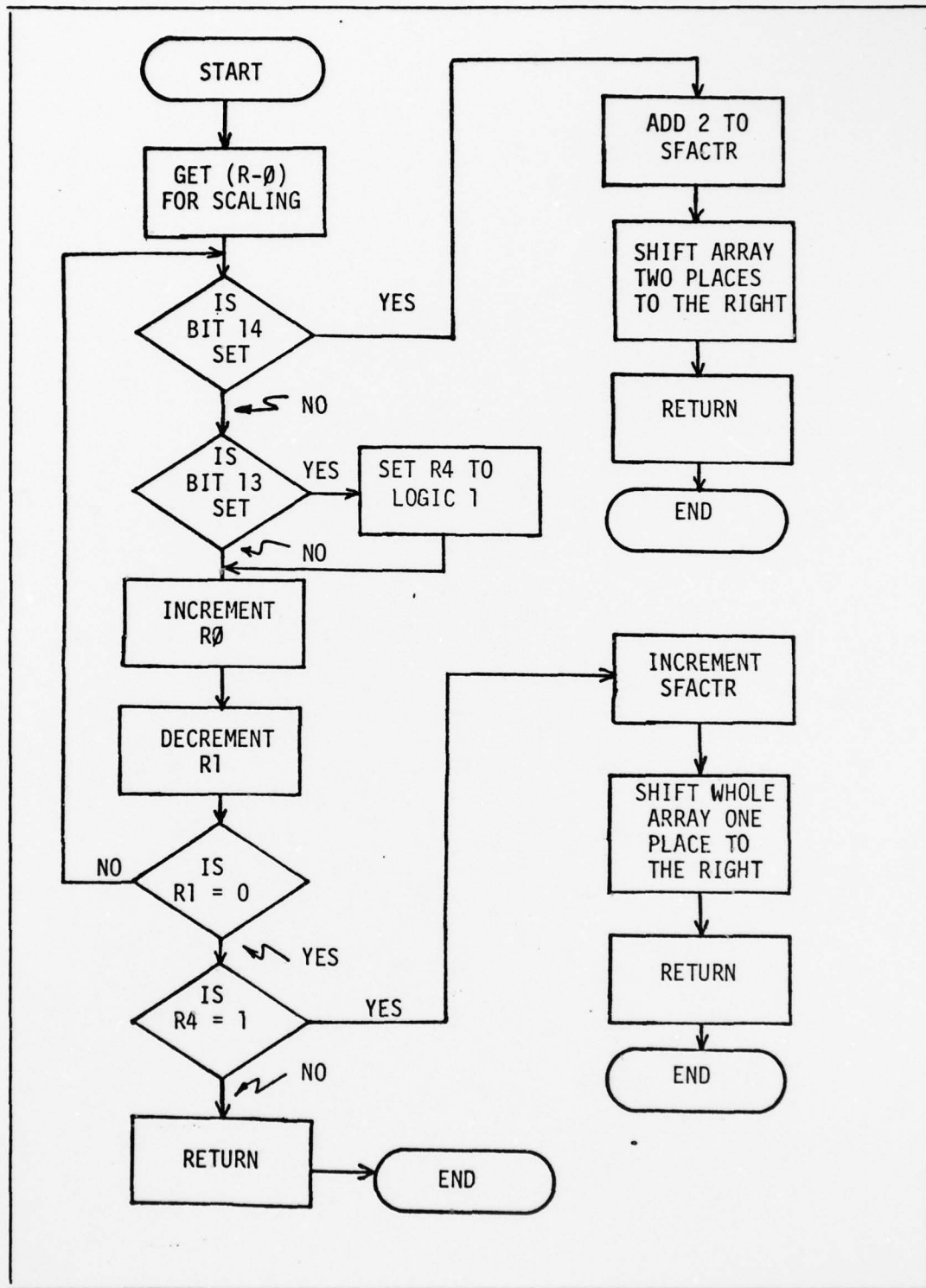


Figure 43. Flow Diagram for Subroutine SCALE

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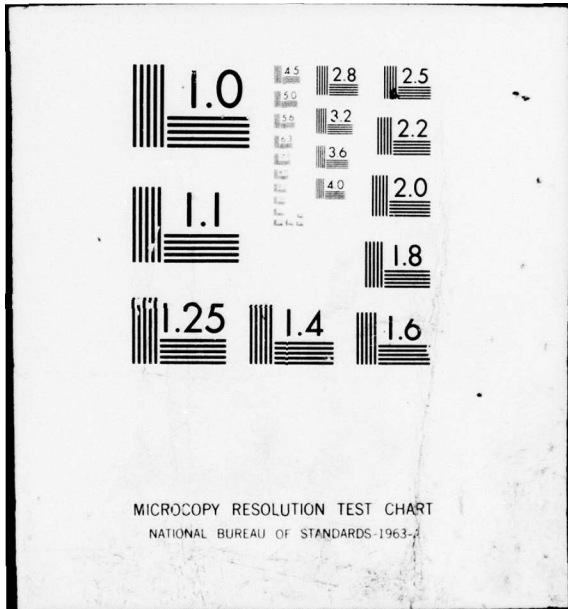
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c. Post-Processing Subroutine. After the FFT of input data is completed, the post-processing of the resulting array is carried out to get the correct output array (Ref Appendix D). Equation (D-8), Appendix D, is rewritten below:

$$X(n) = \frac{1}{2}[V(n) + V^*(N-n)] + \frac{1}{2}j[V(n) - V^*(N-n)]W_{2N}^n \quad (37)$$

where

$$V(n) = Y(n) + jZ(n) \quad (38)$$

Substituting Eq (38) in Eq (37) and separating the real and imaginary parts, the following relationships are derived.

$$\begin{aligned} \text{Re}[X(n)] &= \frac{1}{2}[Y(N-n)] + \text{Cos}\left(\frac{\pi n}{N}\right)[Z(n) + Z(N-n)] \\ &\quad - \text{Sin}\left(\frac{\pi n}{N}\right)[Y(n) - Y(N-n)] \end{aligned} \quad (39)$$

$$\begin{aligned} \text{Im}[X(n)] &= \frac{1}{2}[Z(n) - Z(N-n)] - \text{Cos}\left(\frac{\pi n}{N}\right)[Y(n) - Y(N-n)] \\ &\quad - \text{Sin}\left(\frac{\pi n}{N}\right)[Z(n) + Z(N-n)] \end{aligned} \quad (40)$$

If Rp, Ip, Rn, and Im are defined as follows

$$R_p = Y(n) + Y(N-n) \quad (41)$$

$$I_p = Z(n) + Z(N-n) \quad (42)$$

$$R_m = Y(n) - Y(N-n) \quad (43)$$

$$I_m = Z(n) - Z(N-n) \quad (44)$$

and using the fact that  $\text{Sin}\left[\frac{(N-n)\pi}{N}\right] = \text{Sin}\left(\frac{n\pi}{N}\right)$  and  $\text{Cos}\left[\frac{(N-n)\pi}{N}\right] = -\text{Cos}\left(\frac{n\pi}{N}\right)$ , Eqs (39) and (40) can be rewritten:

$$\text{Re}[X(n)] = R_p + I_p \text{Cos}\left(\frac{n\pi}{N}\right) - R_m \text{Sin}\left(\frac{n\pi}{N}\right) \quad (45)$$

$$\text{Im}[X(n)] = I_m - I_p \text{Sin}\left(\frac{n\pi}{N}\right) - R_m \text{Cos}\left(\frac{n\pi}{N}\right) \quad (46)$$

$$\text{Re}[X(N-n)] = R_p - I_p \cos\left(\frac{n\pi}{N}\right) - R_m \sin\left(\frac{n\pi}{N}\right) \quad (47)$$

$$\text{Im}[X(N-n)] = I_m - I_p \sin\left(\frac{n\pi}{N}\right) + R_m \cos\left(\frac{n\pi}{N}\right) \quad (48)$$

Equations (45), (46), (47), and (48) were eventually used in the implementation of the post-processing subroutine. The flow diagram for the post-processing subroutine is given in Figure 44.

3. Assembly Language Code for FFT. The FFT algorithm was coded in LSI-11 assembly language. A deviation from the flow diagram of Figure 42 was made and all other subroutines (SCALE, BIT INVERSION, POST-PROCESSING) were included in the FFT subroutine. Appendix E gives a source listing of FFT subroutines. The complete analysis of the subroutine execution time is also given in Appendix E.

Power Spectrum Module. The Fourier transformed array, which is the output from the Fourier transform module, contains 16 complex frequency components. The power spectrum for the first 11 frequency components (from 0 HZ to 1000 HZ) needs to be computed. This is accomplished using a subroutine called PRSPEC (Power Spectrum). The flow diagram from the subroutine PRSPEC is shown in Figure 45, which shows that the first 11 frequency components of the power spectrum are calculated by squaring the real and imaginary parts and summing them together. The output is added to the previous elements of an array whose starting address is stored at pointer SPCTRA. This feature helps in computing the average power spectrum which would then be used for computing the damage factor at each frequency of interest [Ref Eq (15)]. The assembly language code for PRSPEC is given in Appendix F; this appendix also contains the analysis of execution time for this subroutine.

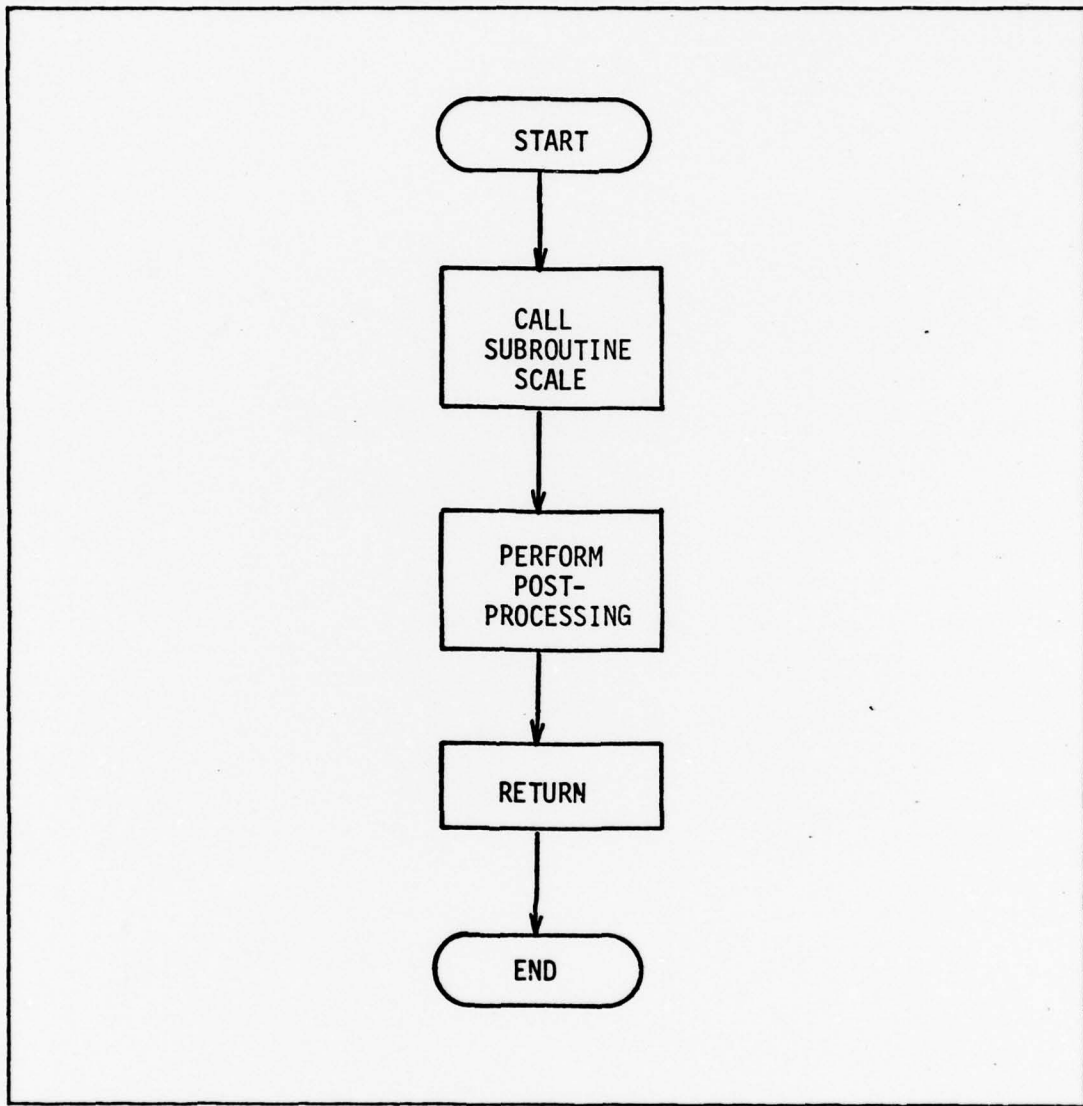


Figure 44. Flow Diagram for Post-Processing Subroutine

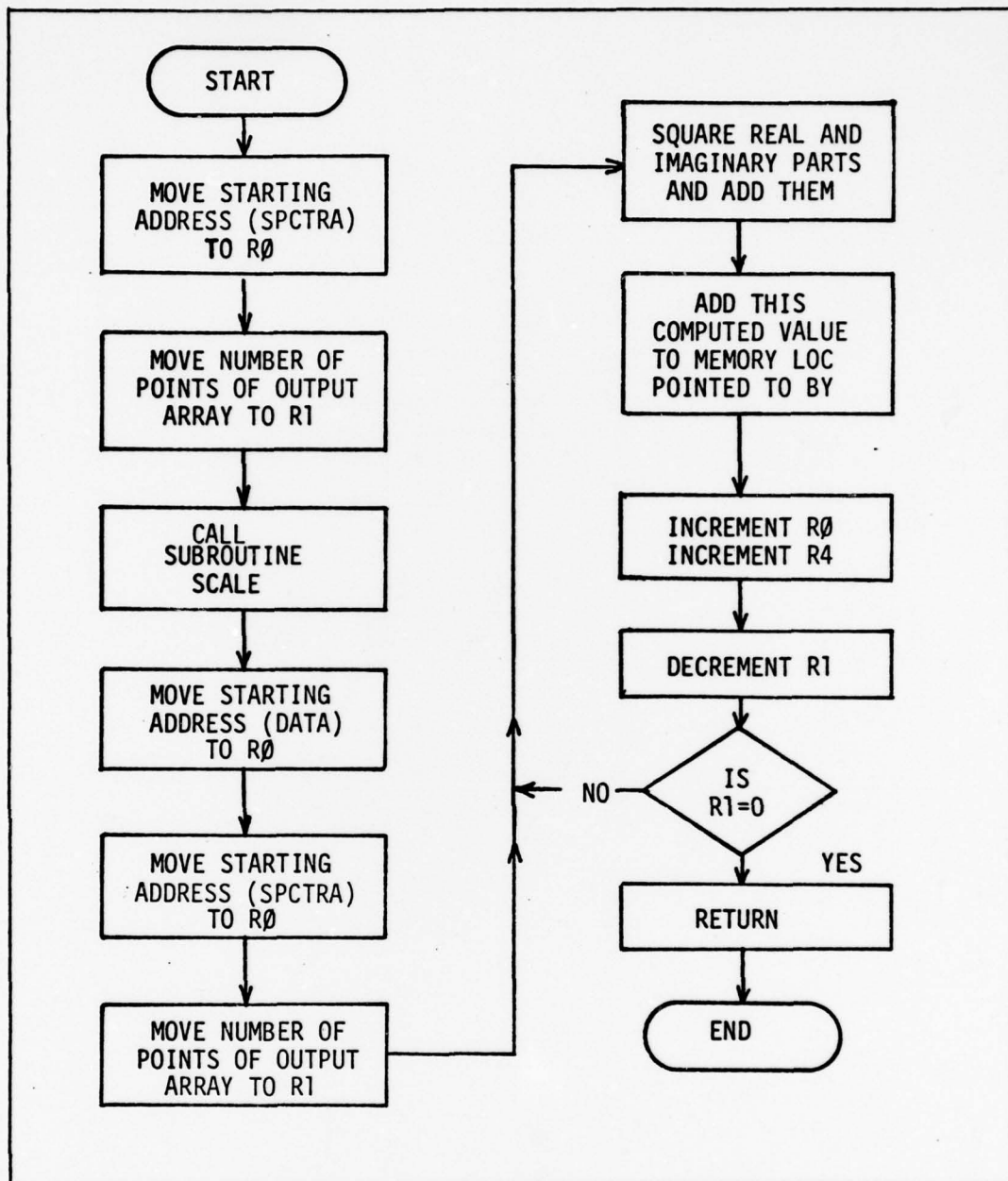


Figure 45. Flow Chart for Subroutine PRSPEC

Damage Factor Module. The subroutine to calculate the DF at each frequency of interest implements Equation (19), which is rewritten:

$$DF(n) = \frac{K[\tilde{A}(n)]^{1.6}}{f(n)^{5.4}}, \quad n = 0, 1, \dots, 10 \quad (49)$$

or

$$DF(n) = K'[\tilde{A}(n)]^{1.6} \quad (50)$$

where

$$K' = \frac{K}{(f_n)^{5.4}}$$

Since the factor  $K'$  is constant for each frequency, the function implemented in this subroutine is  $[\tilde{A}(n)]^{1.6}$ . To implement this exponential factor, the expected values of  $\tilde{A}(n)$ , which are going to be less than equal to 1.0, were divided into 10 intervals as shown in Figure 46.

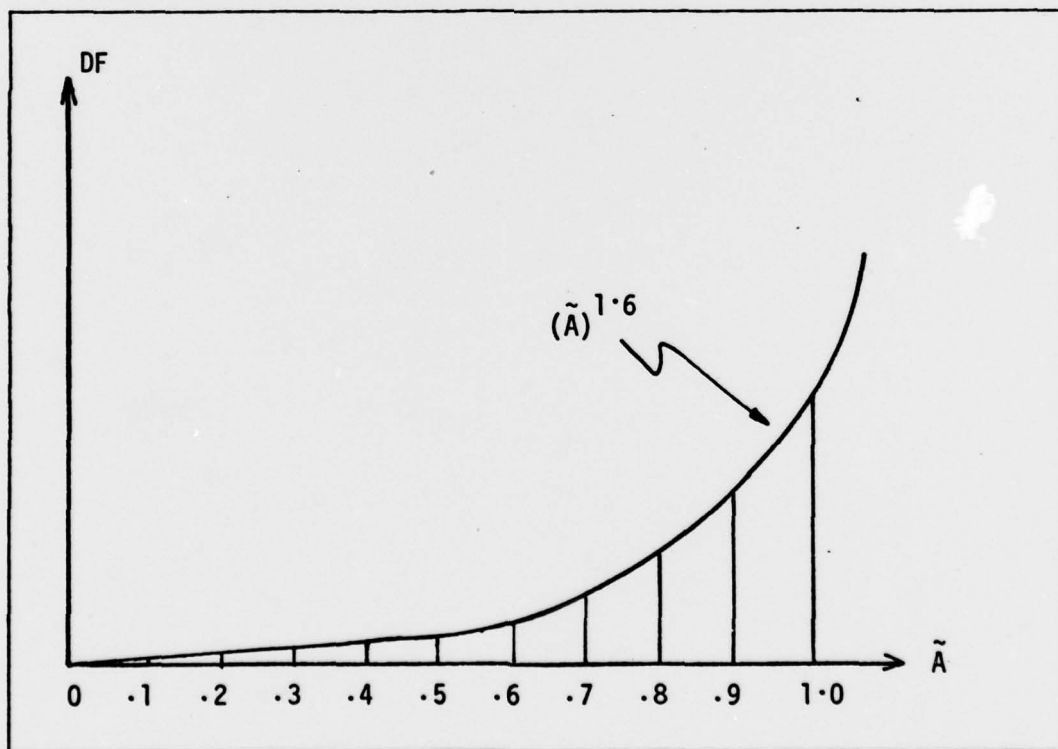


Figure 46. Damage Factor Versus  $\bar{A}$  Curve Divided Into 10 Intervals

A polynomial curve fit program PLSCF, available as an AFIT Library subroutine on the CDC 6600 computer (Ref 15:A-5), was used to fit polynomial curves to each of the 10 intervals shown in Figure 46. It was experimentally determined that if the ordinate (DF) was divided into 8 intervals, shown in Figure 47, a second-order polynomial gave the least RMS error for these intervals as compared with RMS errors for other combinations of intervals. Appendix K gives a description of the curve fitting program. The equation for a second-order polynomial and the coefficients for the 8 intervals (Figure 47) are also given in Appendix K.

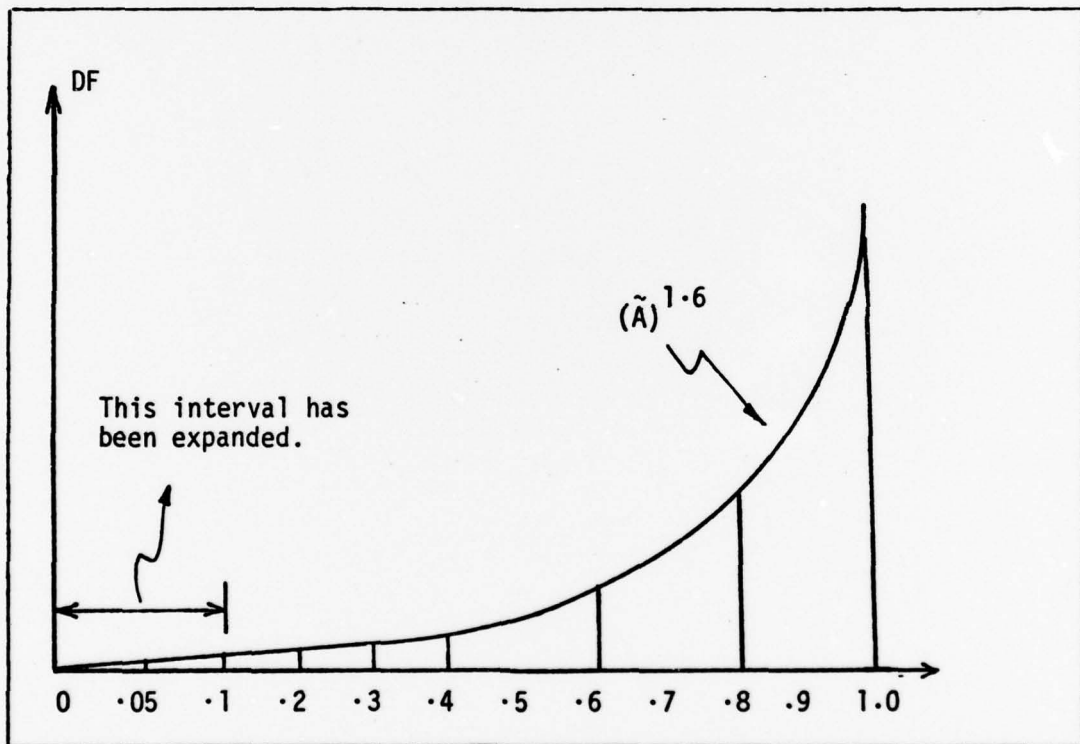


Figure 47. Damage Factor Versus  $\bar{A}$  Curve with the Selected Intervals for Curve Fitting Program (8 Intervals)

The flow diagram for the subroutine to compute the damage factor is given in Figure 48; this subroutine is called POLYN. Note that the

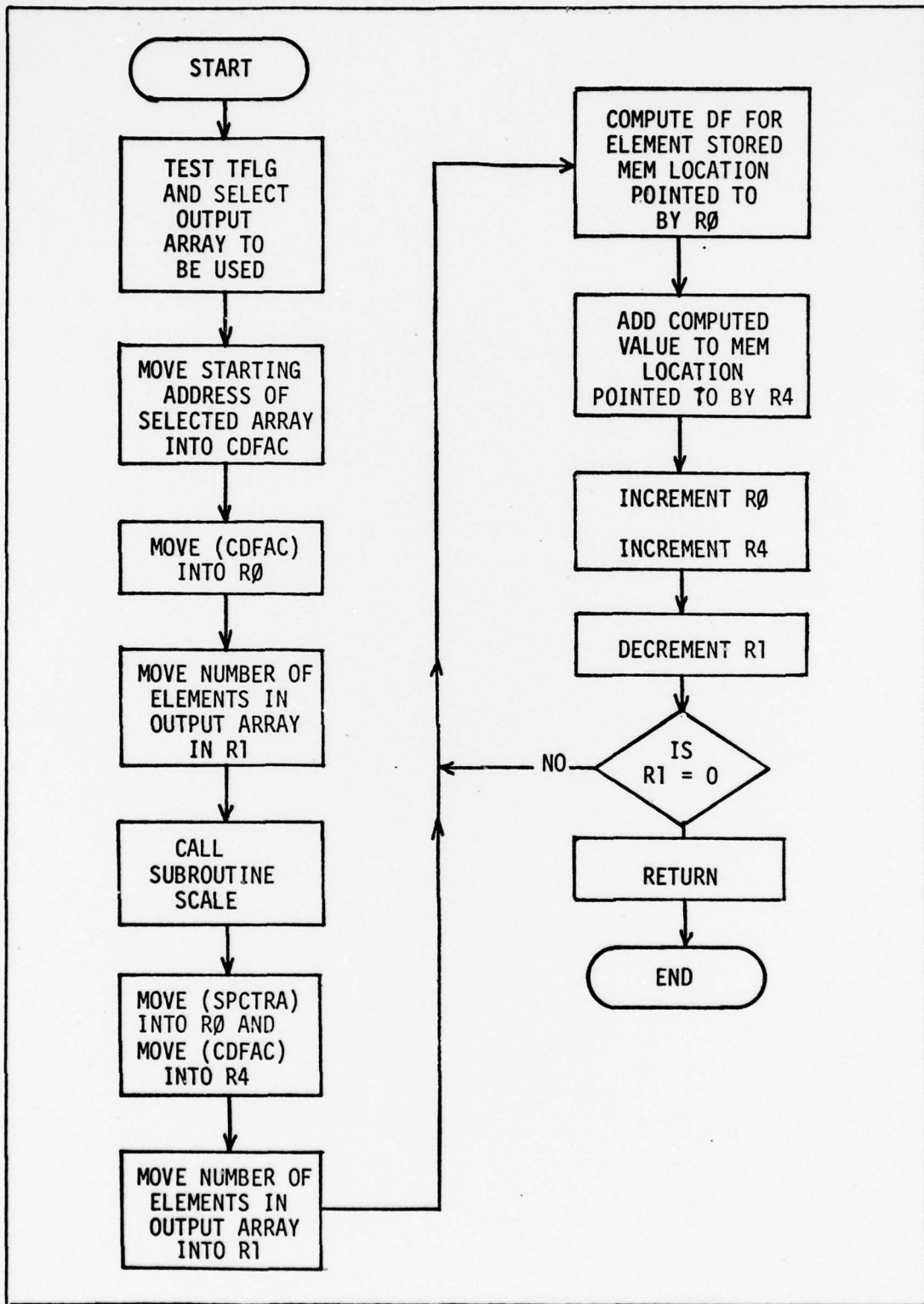


Figure 48. Flow Diagram for Subroutine POLYN

resulting DF values are added to the existing elements of the selected output array (Ref Chap III). This is done to get the CDF as a function of temperature and frequency. The assembly language code for subroutine POLYN is given in Appendix G. This appendix also gives an analysis of execution time for this subroutine.

Control Executive (EXEC). The Control Executive is the systems program that handles the execution of all the modules in MIBDAPS. The control and synchronization of different activities is also handled by this program.

The state transition diagram that shows the initial switching on sequence is given in Figure 49. Once the power is switched ON to the system, the microcomputer executes its own interval microcode and jumps to location  $(17300)_8$  where it finds a WAIT instruction. Once the POWER OK or POWER FAIL interrupt occurs, the service routine for either of these interrupts is executed. In case the POWER ON sequence is successful, the execution of the system starts.

Figure 50 shows the flow diagram for the continuous execution of the system. The acquisition of data is done using the interrupt feature of ADV11-A. Therefore, the data acquisition occurs while the processing of previously acquired data is going on. The real time clock would be used for synchronization for the entire data acquisition and processing cycle, but this feature has not yet been included in the flow diagram. The "dashed" block in Figure 50 (page 94) would accomplish the synchronization of different processing cycles.

The assembly language code for EXEC is given in Appendix H. The execution time analysis is also included in Appendix H.

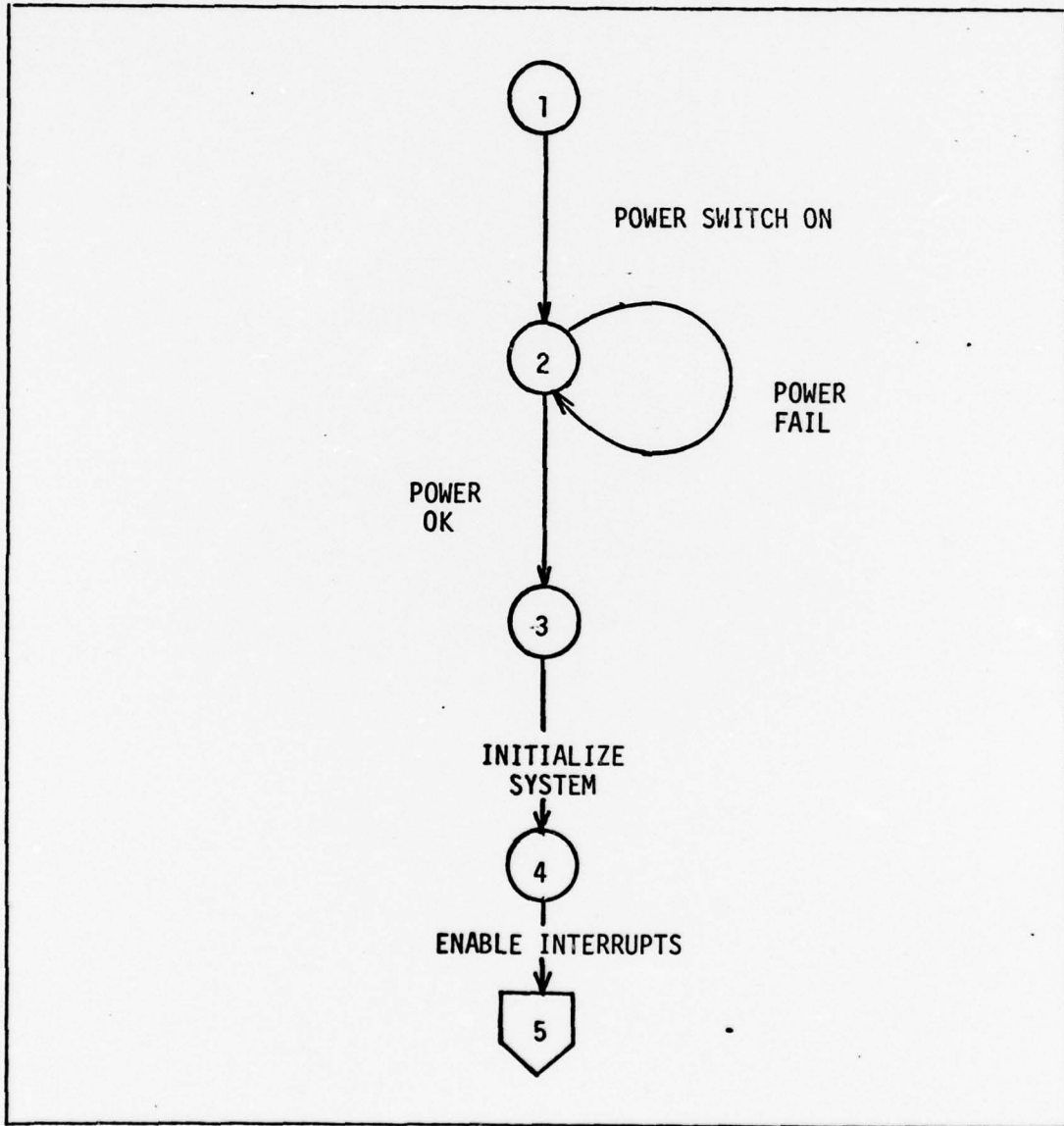


Figure 49. State Transition Diagram for Initial Power ON Sequence

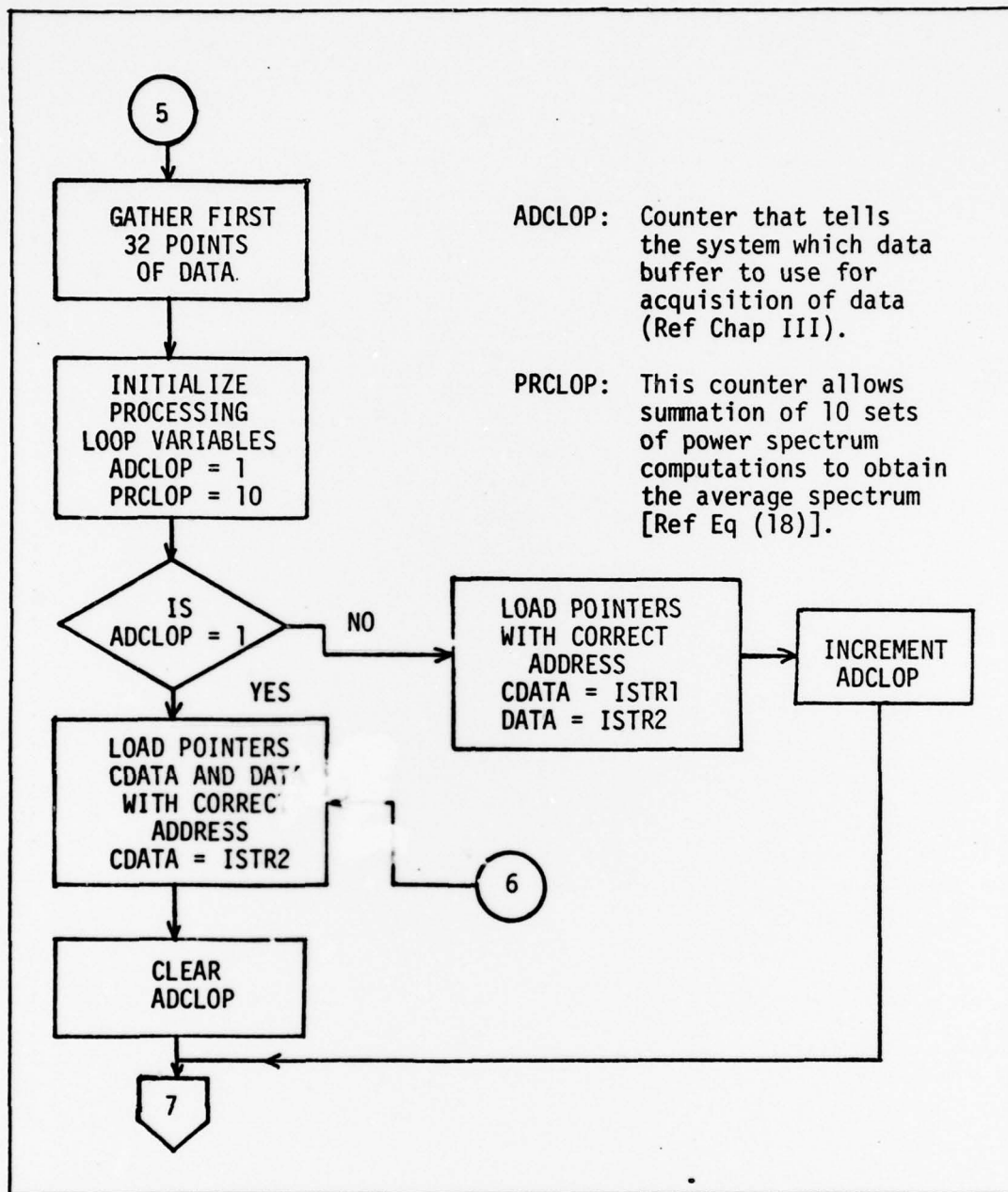


Figure 50. Flow Diagram for Continuous Execution

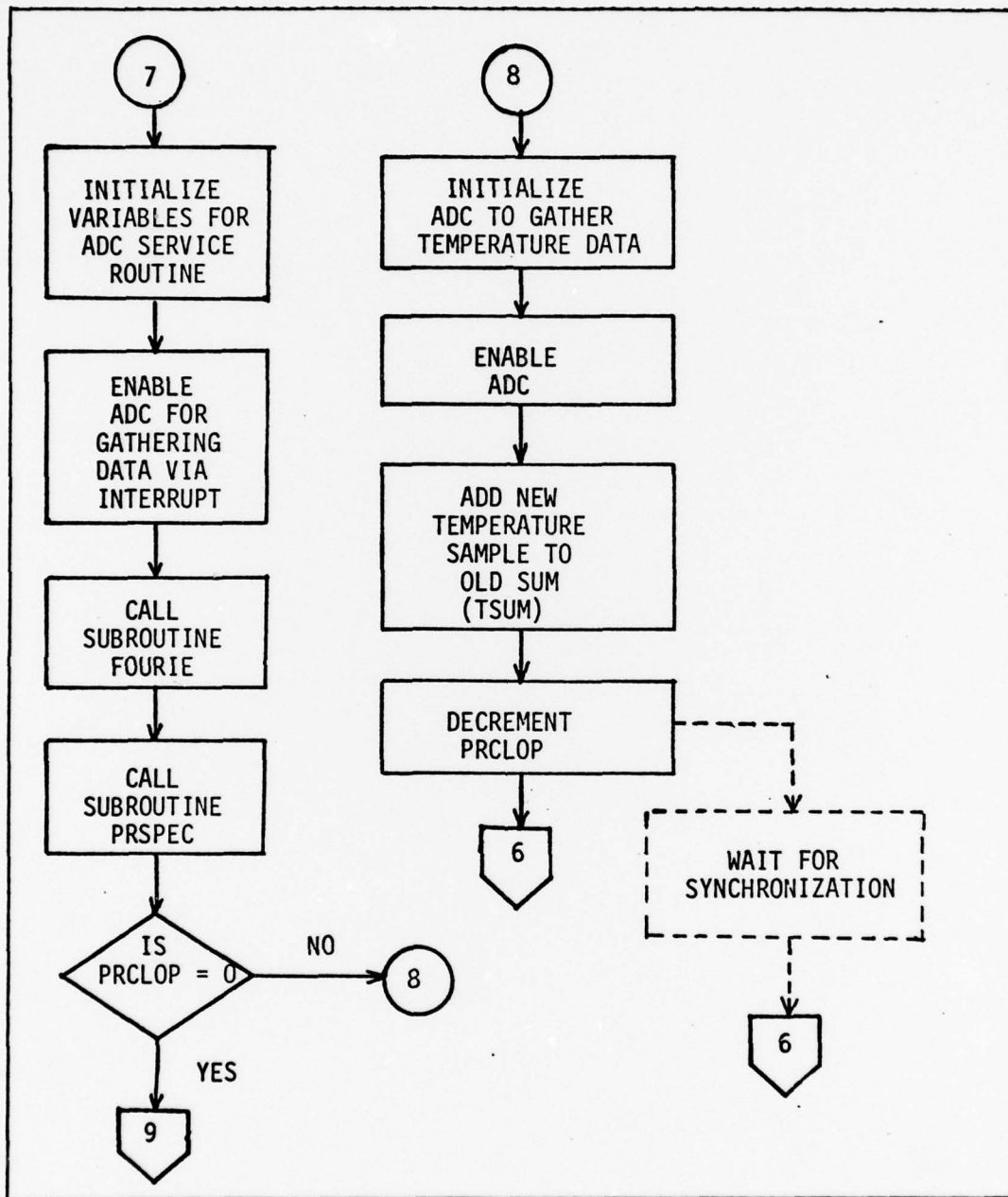


Figure 50. Flow Diagram for Continuous Execution (cont)

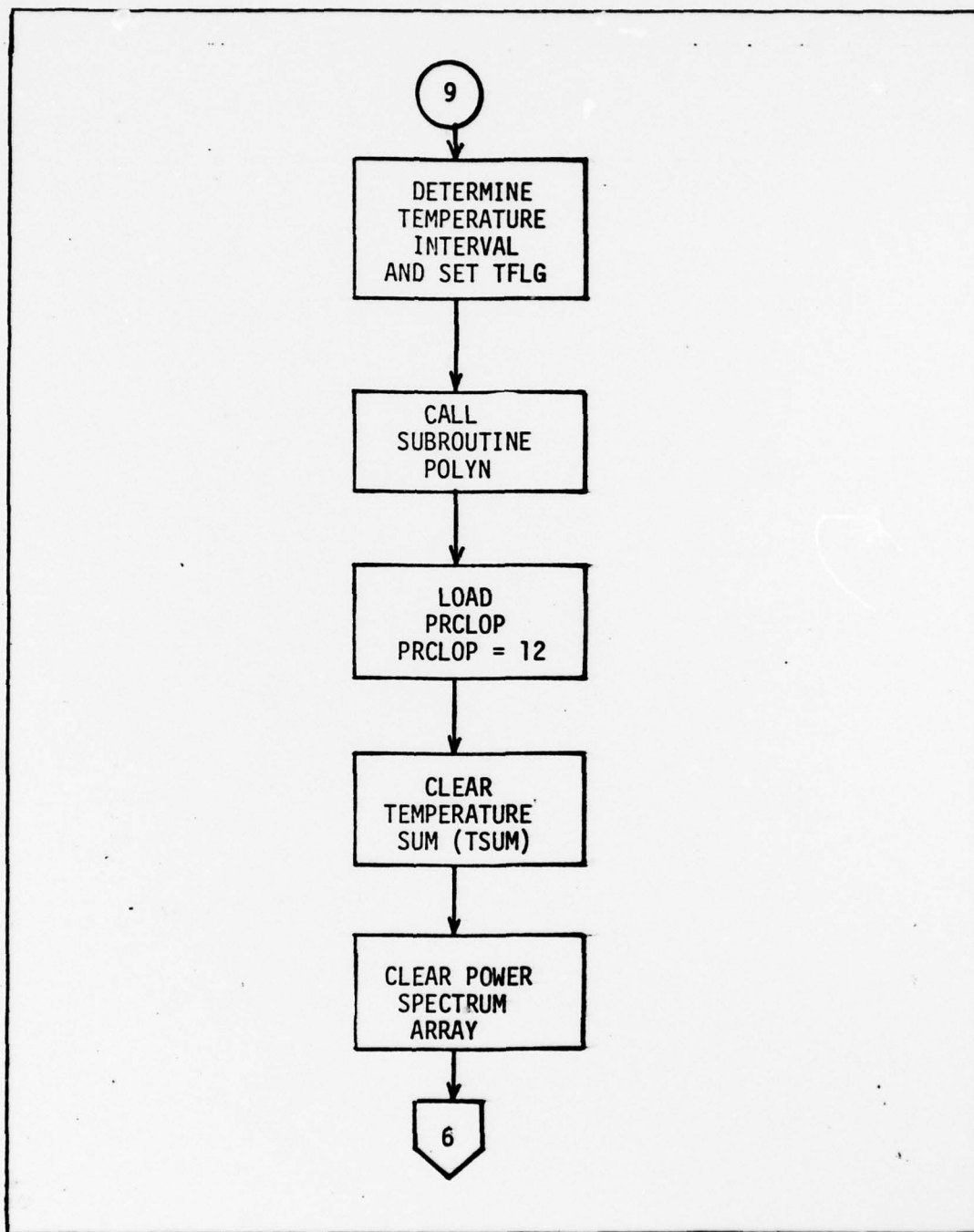


Figure 50. Flow Diagram for Continuous Execution (cont)

### Summary

In this chapter, the mathematical analysis of data acquisition and processing functions of MIBDAPS was presented. The theoretical background necessary for the development of each software module was discussed. This was followed by the generation of a flow diagram and assembly language code for each software module. The design of hardware realizable modules of MIBDAPS is undertaken in the next chapter.

## V. Design and Implementation of Hardware

### Introduction

The design and implementation of software realizable modules of MIBDAPS was presented in Chapter IV. The next step in the development of MIBDAPS was the design and implementation of hardware realizable modules (Ref Chap III). In this chapter, the general design approach is discussed and the implementation of the designed circuits using standard off-the-shelf integrated circuits is explained. At the end of this chapter, the DC power supply requirements are discussed and the use of standard DC power supply modules is suggested. The author's conception of the prototype module of MIBDAPS along with its physical dimensions is also presented.

### Design of Hardware Circuits.

The design of different hardware realizable modules is described in the following paragraphs. The design is based on the functional requirements spelled out in Chapter III.

System Clock Generation Circuit. The data acquisition function of MIBDAPS required that the vibration analog signal be sampled at a frequency of 3200 HZ (Ref Chap IV). Therefore, to gather data at this sampling rate, it was necessary to trigger the ADC (ADV11-A) by a 3200 HZ external clock (Ref 11:335-338). The same clock was also required to synchronize different processing loops during system execution (Ref Chap IV). Figure 51 is the schematic that shows the hardware configuration where the system clock was used. It should be noted that the output of the clock is fed to the RTC before being applied to the ADC. This has

been done to ensure that the clock is conditioned by the Schmitt trigger circuit on the RTC board (Ref 11:211-212) before being applied to the ADC.

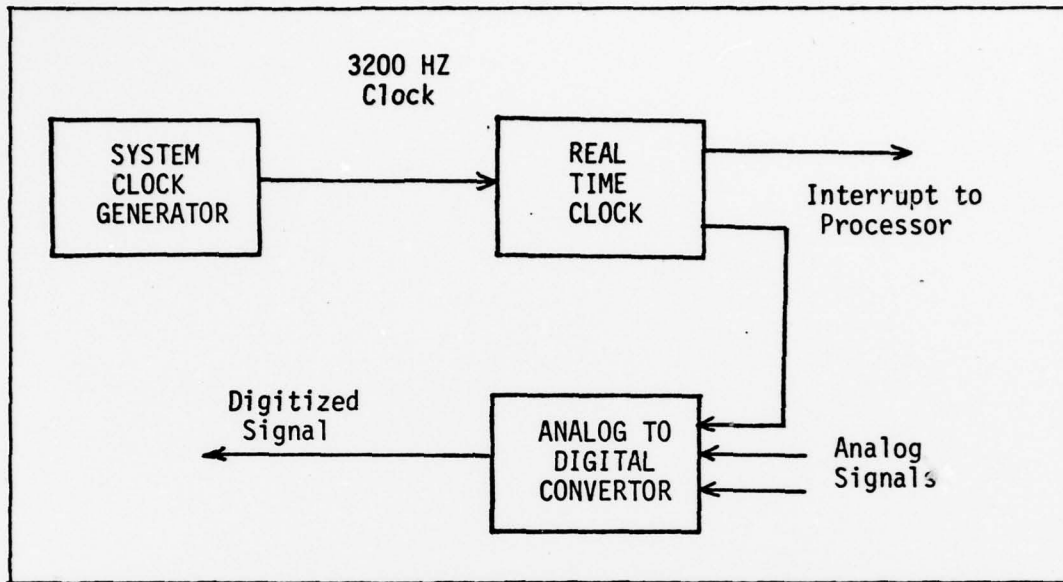


Figure 51. Hardware Configuration Showing System Clock Generator

To design a stable 3200 HZ clock, SN74LS424 clock generator IC and a 64K HZ crystal were used as the reference oscillator (the choice of crystal frequency was arbitrary). This reference frequency was divided down to get the required 3200 HZ clock. Two SN7490 ICs were configured to accomplish the appropriate frequency division. Figure 52 shows the circuit diagram of the designed clock generator.

POWER ON Interrupt Circuit. The LSI-11 bus foundation module (DRV11-P), which is manufactured by DEC, can be used for user defined interfaces (Ref 11:169). This module is supplied with the logic necessary for interfacing to the LSI-11 bus. This logic includes bus transceivers, a device address comparator, protocol logic, interrupt logic, vector address comparator, and bus receivers and invertors (Ref 11:169). The

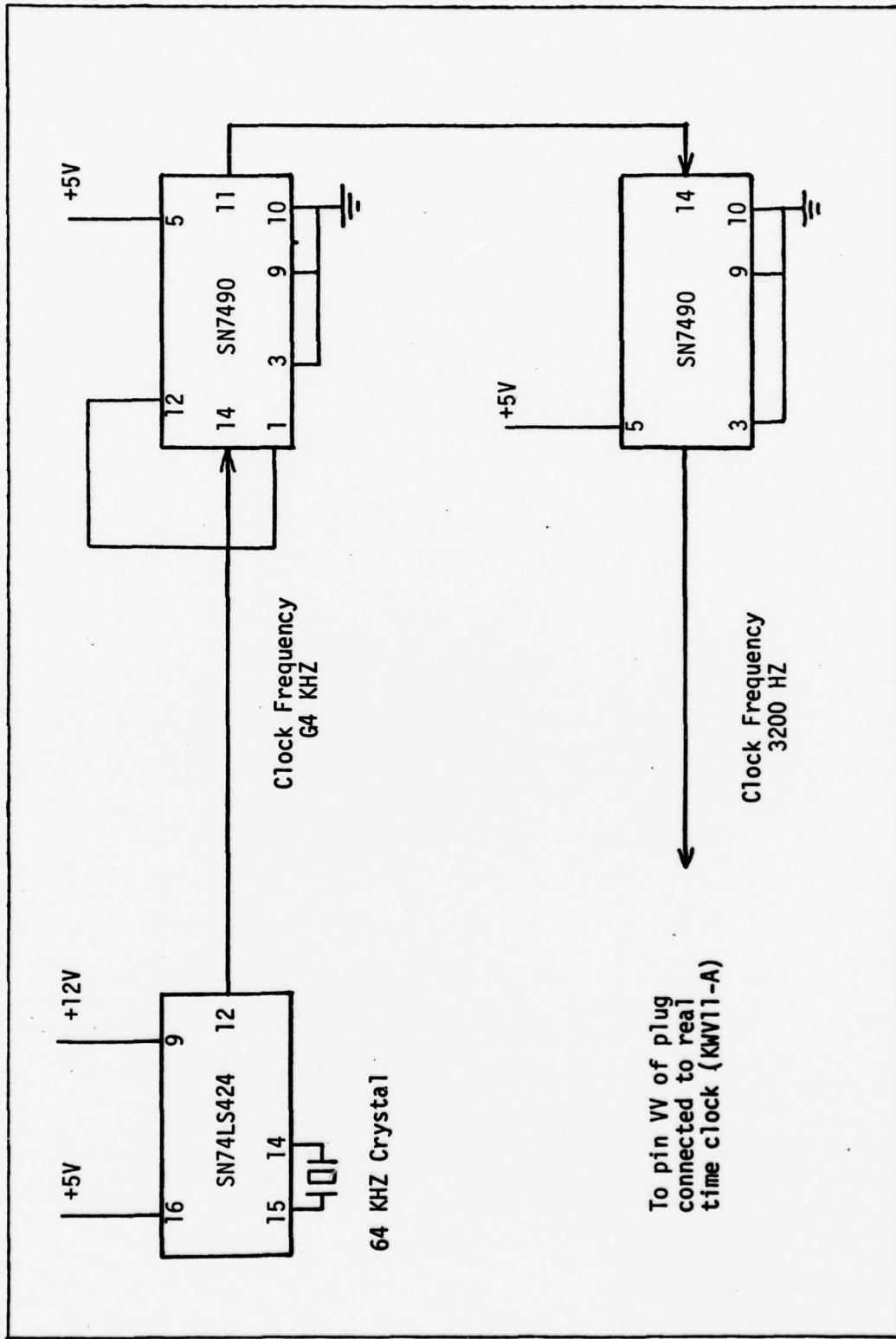


Figure 52. Circuit for System Clock Generator

device interrupt generator and vectors address generation feature of this module were used in the design of the POWER ON interrupt circuit.

The POWER OK interrupt occurs whenever the DC power to the system has been switched ON and +5V, +12V DC supplies are within acceptable limits. The POWER ON transition and a signal from the power sense circuit would be used to generate the requisite sequence of signals necessary to activate the interrupt and vector address generator feature on the DRV11-P. (The reader is referred to Ref 11:177 for an understanding of the logic requirement placed on the requesting device.)

The block diagram of the designed circuit is given in Figure 53.

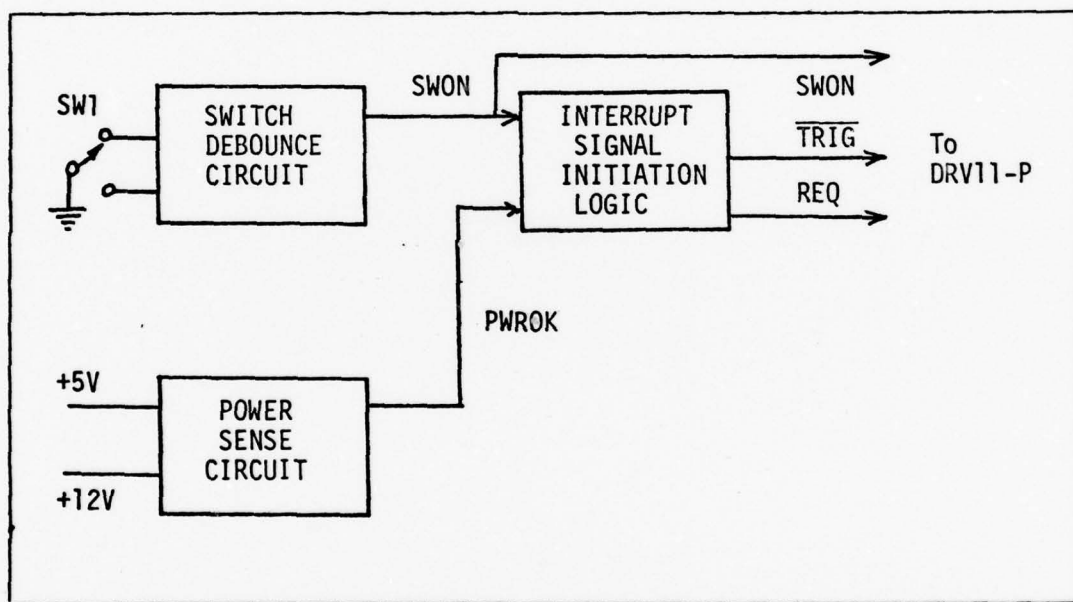
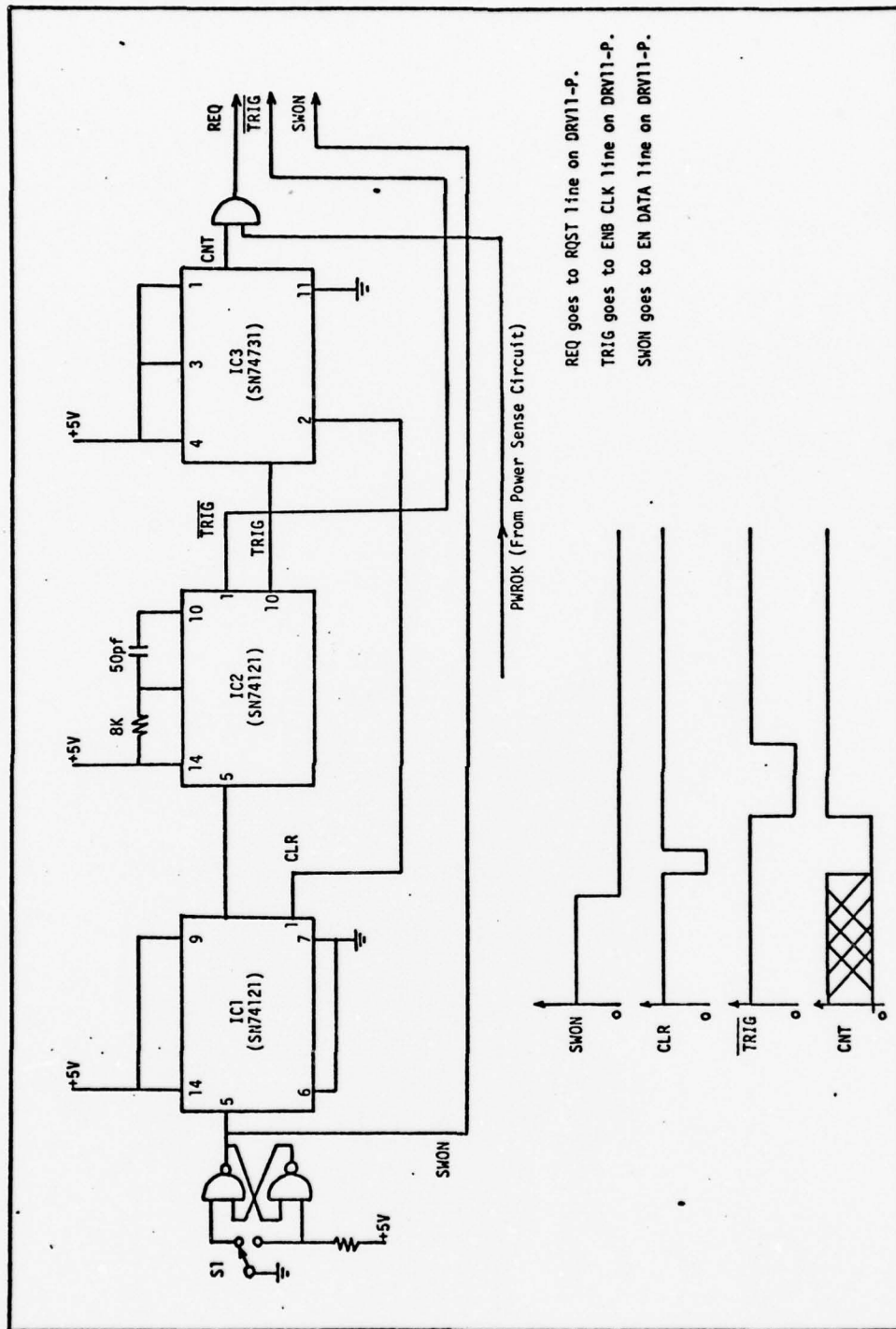


Figure 53. Block Diagram of PWROK Interrupt Generator

This circuit assumes that the +5V supply is available to itself and the DRV11-P module before the POWER ON switch (SW1) makes the OFF-to-ON transition. This assumption would create a problem if there is only one power ON/OFF switch in the system. The solution to overcome this problem



REQ goes to RQST line on DRV11-P.  
 TRIG goes to ENB CLK line on DRV11-P.  
 SWON goes to EN DATA line on DRV11-P.

Figure 54. Circuit Diagram for POWER OK Interrupt Generator

is to have a RUN/STANDBY switch, which is put to the STANDBY position before the POWER ON/ORR switch is switched ON. The RUN/STANDBY switch is then switched to the RUN position, which would generate the appropriate signal sequence for the interrupt and interrupt vector generation logic on the DRV11-P. Figure 54 shows the designed circuit diagram along with the timing diagram of the logic signals generated.

POWER FAIL Interrupt Generator. This circuit would be used to generate an interrupt whenever a power failure occurs or when the RUN/STANDBY switch makes a RUN-to-STANDBY transition. Figure 55 shows the block diagram of the hardware configuration required. This circuit was not designed due to the shortage of time. The DRV11-P could be used to generate the POWER FAIL interrupt; however, a logic circuit to initiate the interrupt request and interrupt vector cycles would have to be designed.

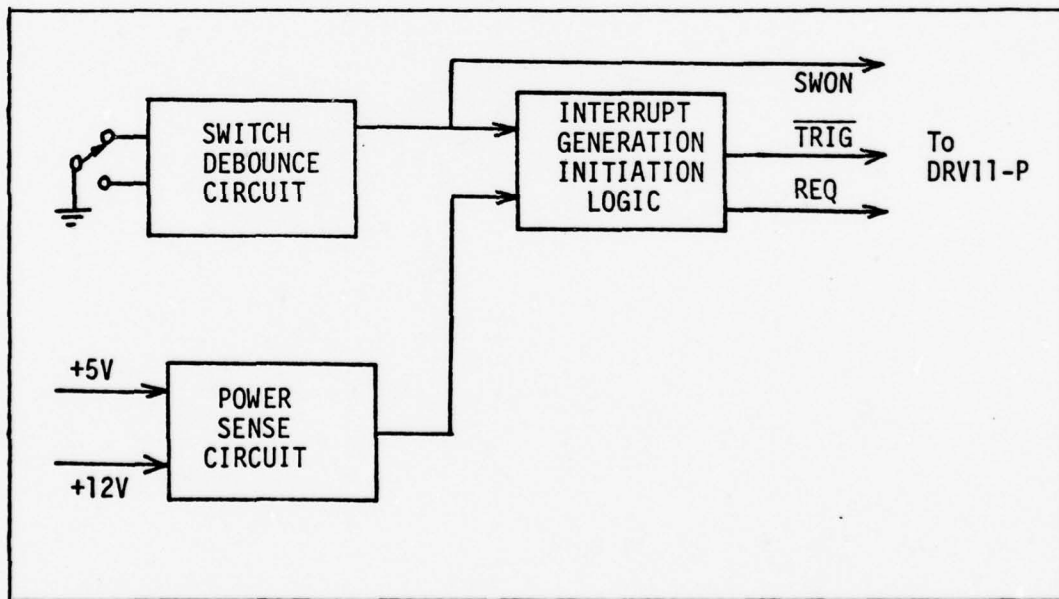


Figure 55. POWER FAIL Interrupt Generation Logic

Operational Time Recorder. The block diagram for the operational time recorder is given in Figure 56. The system clock is fed to this

circuit, and it keeps track of the time for which the system remains operational. Whenever the operational time reaches the 15-hour limit, an interrupt is generated and the system executives a "halt" cycle.

The operational time is also displayed on the front panel. In case either a power failure occurs or the system is switched to STANDBY mode, the current operational time would be stored in the core memory. This is done to obtain the total operational time for which the system monitors a test piece in case the system was operated in different time segments.

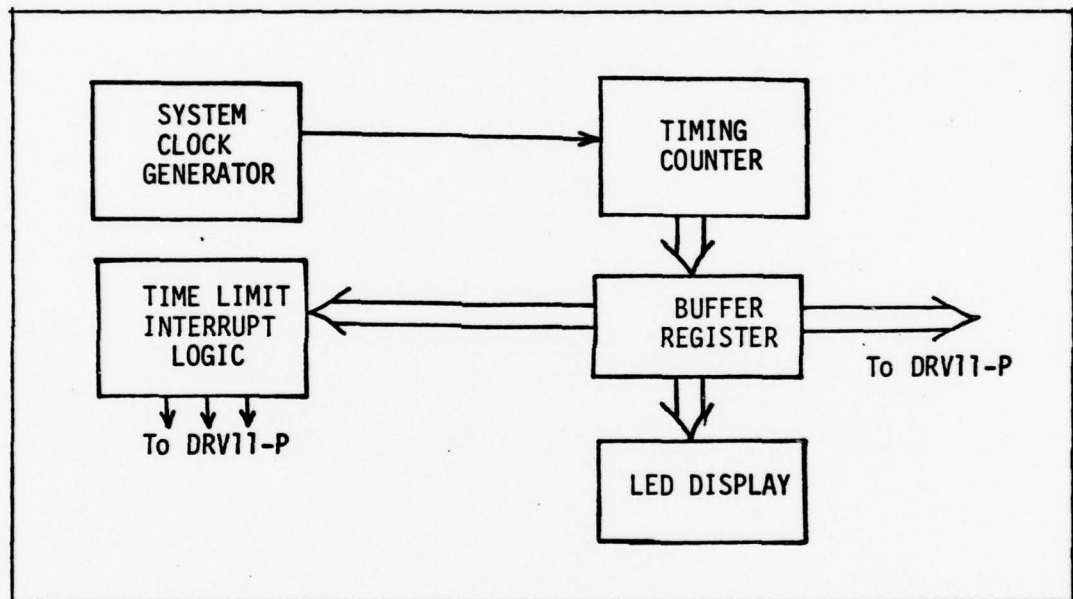


Figure 56. Operational Time Recorder

#### Power Supply Requirement

The power requirement for each hardware module used in MIBDAPS is tabulated in Table IV. The +5V DC supply is required to source a maximum current of 16.63A, and the +12V is required to source 2.21A maximum. A comparative analysis of the +5V and +12V DC supplies available from different manufacturers was done. The following characteristics were compared during this analysis: maximum output current, load regulation, ripple and

noise, case size, and weight. Based on this analysis, model numbers HE 252 and HE 212, the +5V DC and +12V DC supplies, manufactured by Computer Products Incorporated, are the recommended DC power supplies for MIBDAPS. Table V lists the salient features of these recommended supplies; it should be noted that the maximum rated current for the supplies is more than the required maximum current for MIBDAPS (see Table IV). This would enable future additions of the required hardware modules.

#### Mechanical Description of MIBDAPS

The author's conception of the final shape of MIBDAPS (called MIBDAPS-I) is given in Figure 57. The detailed diagram showing the layout of different subunits and their dimensions is given in Figure 58. It should be noted that the rear view shows an AC inlet socket and an RS-232 outlet. The RS-232 would be provided to connect the system to a TTY, which would be used as an I/O device to get a printout of CDF versus frequency and temperature arrays.

#### Summary

In this chapter, the design and implementation of two hardware modules was presented along with the block diagram descriptions of the other two hardware modules which need to be designed. The power supply requirements was determined and a recommendation for off-the-shelf DC power supplies was made.

The physical dimensions for the mainframe and the location of subunits within the mainframe were made. Diagrams depicting the physical shape and dimensions of MIBDAPS have also been drawn. Integration of the system was accomplished next. The results of testing the integrated system are discussed in the following chapter.

TABLE IV

## Power Requirement for MIBDAPS

Description of Modules	Current Required From +5V Supply (nominal/maximum)	Current Required From +12V Supply (nominal/maximum)
LSI-11M Processor	1.8A/2.4A	0.8A/1.1A
DLV11 Serial Line Unit	1.6A/1.6A	0.18A/0.25A
ADV11-A A/D Convertor	2.0A/2.0A	0.45A/0.45A
MMV11-A 4K Core Memory	3.0A/7.0A	0.6A/0.6A
REV11-C Bootstrap ROM and Terminator	1.0A/1.88A	--
KWV11-A Real Time Clock	1.75A/1.75A	0.01A/0.01A
Total current requirement for +5V supply = 11.15A/16.63A. Total current requirement for +12V supply = 2.04A/2.21A.		

TABLE V

## Recommended Power Supply Characteristics

Description	Output Current (maximum)	Load Reg. Maximum (NL-FL)	Ripple and Noise	Case Size (in.) L x W x H	Weight (lbs.)
+5V HE 252 (110V AC output)	20.0A	$\pm 0.1\%$	50MV P-P (13 MV RMS)	6.5x4.5x	3.25
+12V HE 212 (110V AC output)	3.0A	$\pm 0.1\%$	20MV P-P (2MV RMS)	6.5x4.5x 1.8	1.7

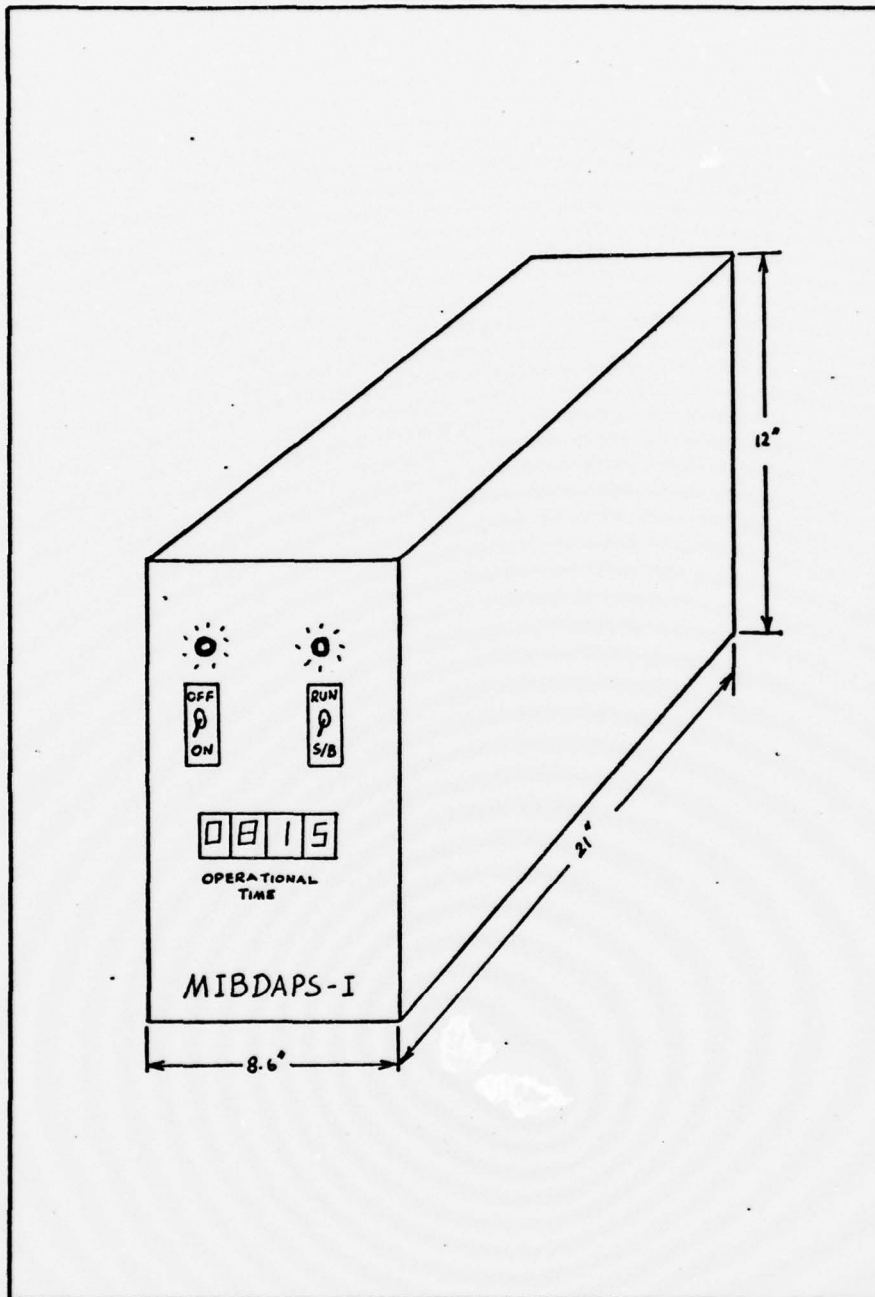


Figure 57. Author's Conception of the Final Shape of MIBDAPS

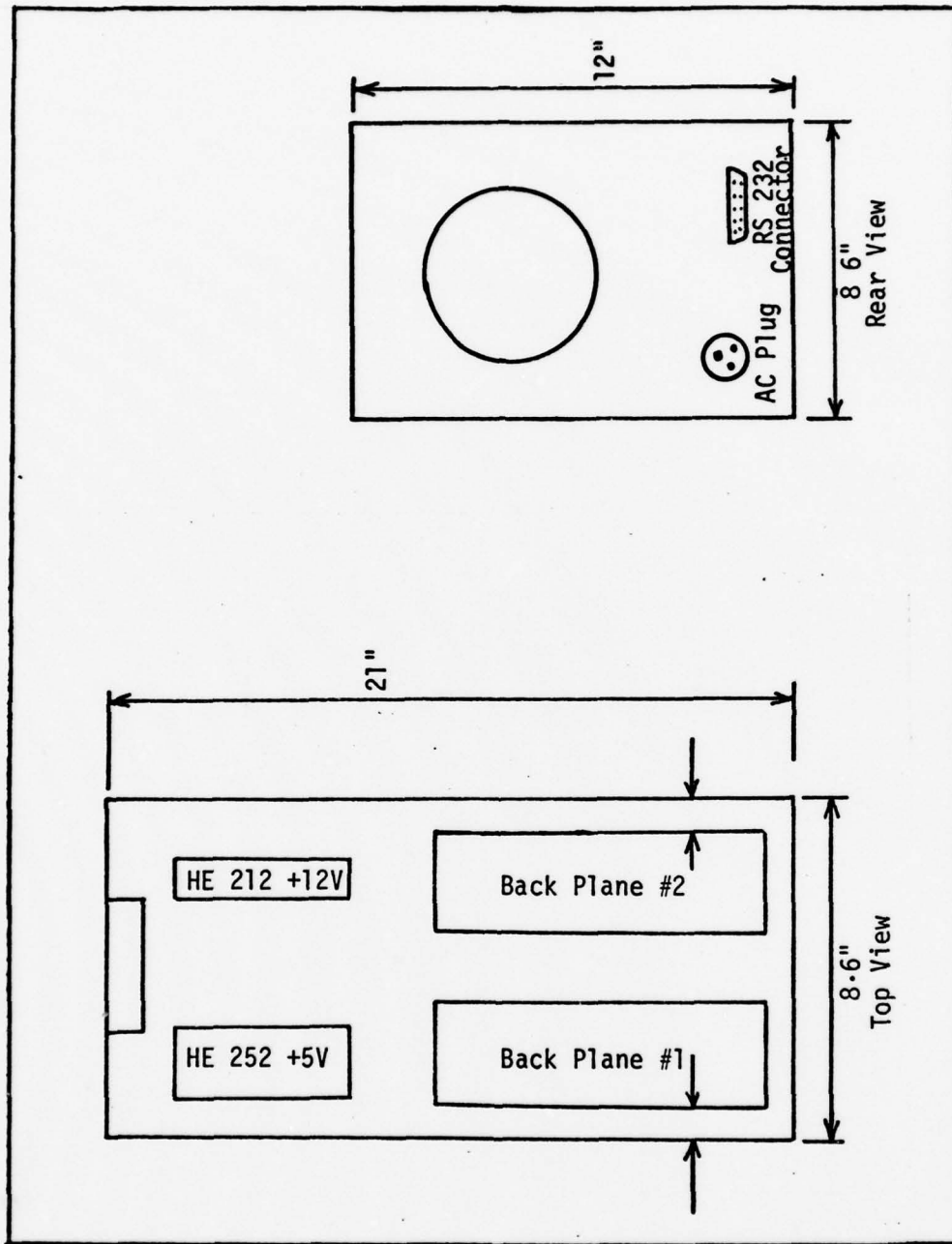


Figure 58. Location of Subassemblies in the Mainframe

## VI. Testing and Experimental Results

### Introduction

The testing of the software during the development stage, the integration of hardware and software, and the testing of the integrated system are discussed in this chapter. The experimental setup used in both the software development stage and the system integration stage is explained, and the results of the experimentation and software developed for testing are presented at the end of the chapter.

### Software Debugging and Testing

The software algorithms that were designed and coded in Chapter IV were debugged and tested on a PDP-11/03 minicomputer system. The configuration of the development system used is given in Figure 59. A three-step approach was taken to develop software. First, the assembly language code was assembled and syntax errors were removed. Second, the logical errors detected during program execution were debugged. And, lastly, the testing of software for computational accuracy was done. The flow diagram given in Figure 60 shows the complete software development cycle.

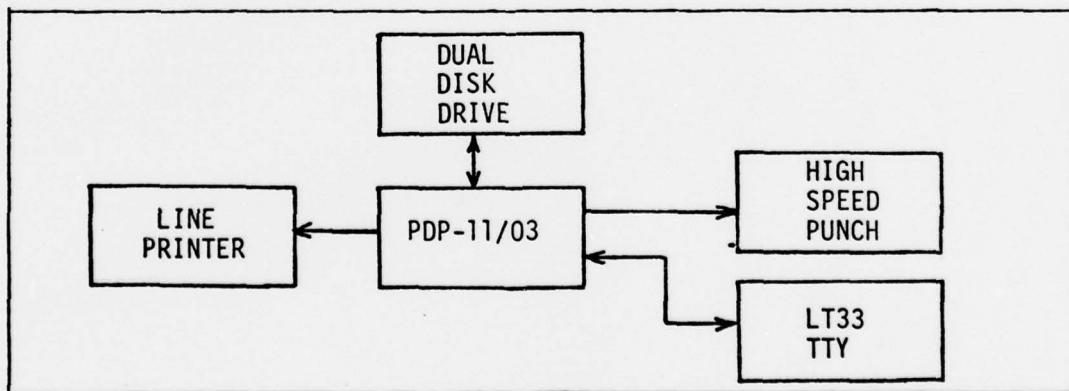


Figure 59. Configuration of the Software Development System

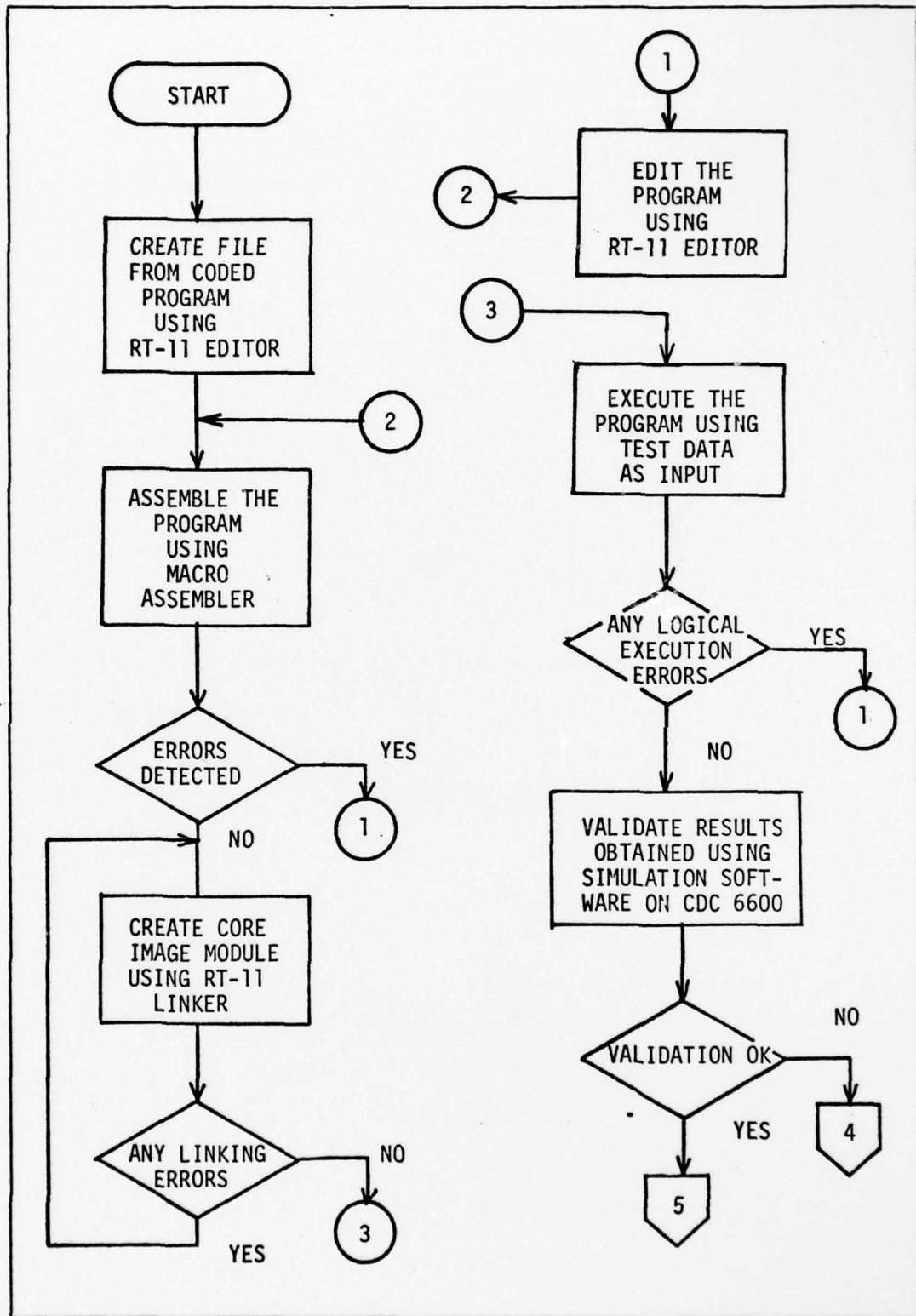


Figure 60. Software Development Cycle

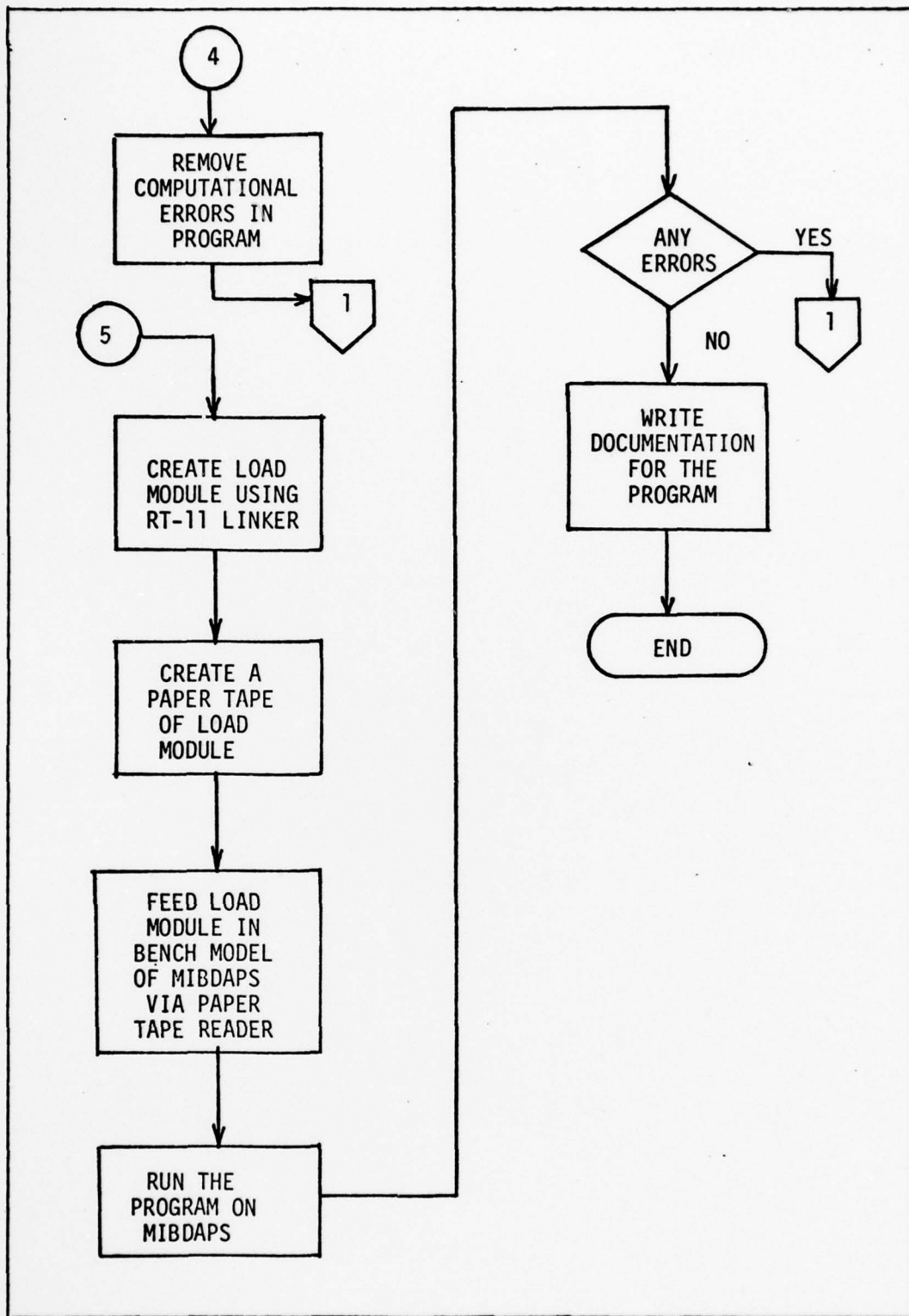


Figure 60. Software Development Cycle (cont)

Testing of Software Modules. After the syntax/semantics and logical error had been removed, the testing of each software module was undertaken. The approach followed was to write FORTRAN simulated programs which duplicated the computations performed by the software module being tested. These simulation programs were executed on the CDC 6600. The input data to the simulation programs was similar in format and value as that used during the execution of the actual software module. The results obtained from the simulation were compared with earlier results from the system software and, in case of discrepancies, the system software code was checked for computational errors. Appendix J gives the simulation program used to check the subroutines FOURIE, PRSPEC, and POLYN (Ref Chap IV). The test data used and the results obtained from the simulation programs and system programs are also tabulated for comparison (Ref Appendix J). These results are discussed at the end of this chapter.

#### Integration of Software and Hardware

The system integration was done once the software modules had been developed and tested. The system integration phase required setting up a bench model for MIBDAPS and testing the control execution. The timing constraints for both data acquisition and processing were also determined at this stage.

Bench Model for MIBDAPS. The configuration of the bench model for MIBDAPS is given in Figure 61. The function generator was used to simulate the vibration signal, and the square wave generator was used as the system clock. The bench model was tested by running the control executive program which exercised the data acquisition and processing function

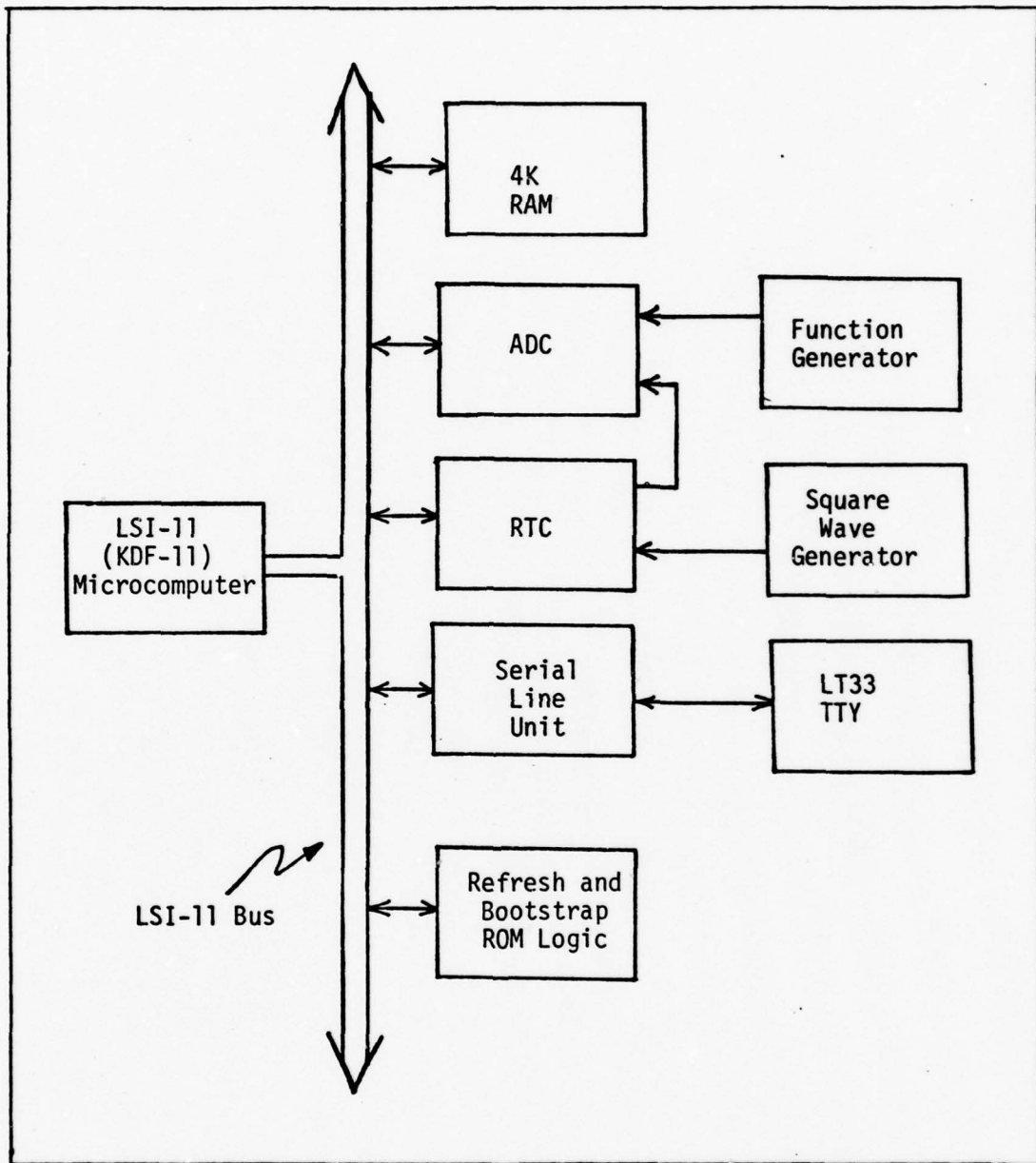


Figure 61. Configuration of the Bench Model for MIBDAPS

of MIBDAPS. To check the results at different stages of acquisition and processing, a decimal print-out subroutine called DECPR was included; this subroutine was called by a macro called DECMAC. The source listing for subroutines DECPR and macro DECMAC is included in Appendix I. The final source listing for the control executive, which is given in Appendix H does not contain this print-out feature because during the actual operation of MIBDAPS the intermediate results are not required to be printed. The results of the control executive operation are discussed later in the chapter.

### Results of the Experimentation

The result obtained from testing individual software modules and the operation of the control executive are presented in the following sections.

Software Test Results. The results of software testing are presented as the root means square (RMS) error between the output from the simulation program and the output from the actual software module. These results are given in the tables in Appendix J. The RMS error criterion was used because an estimate of the average error for each software module was desired. The mean of average RMS error is .0387 (see Appendix J). This error is acceptable because it is much less than the specified tolerance of  $\pm 10\%$  (Ref 13).

Execution Results for Control Executive. The two aspects that are important from the functional standpoint of MIBDAPS are the execution time for the acquisition and processing loop and the accuracy of time computations.

The execution times for each software module were computed separately (Ref Chap IV), and the total execution time for the control

executive was computed using the flow diagram given in Figure 62. The loop that computes the power spectrum for a new set of data points takes 29.391 msec, and the time taken to compute the damage factor once the average power spectrum has been computed (Ref Chap IV) is equal to 4.596 msec. Figure 63 shows the timing relationship between the actual acquisition and processing time and the desired time as defined in the initial requirement (Ref Chap I). The designed system, therefore, would update the CDF versus frequency and temperature array every 298.48 msec. For a new set of 32 data points, the power spectrum computation takes 29.36 msec; the average over 10 such intervals would yield a good estimate of the power spectral density of the input vibration signal. This estimate would not be in error as compared with the average power spectral density obtained if the power spectrum computed had taken less than 10 msec (as originally stipulated). This statement is true because the input vibration signal can be assumed to be wide-sense stationary (Ref 13), and the estimate of the average power spectral density of such a signal depends on the number of power spectrum samples that are averaged and not on the rate at which these samples are available (Ref 6).

Interpreting the Output Array. The CDF versus frequency and temperature array would contain fractional binary numbers that have to be correlated with actual physical data. An analysis of the scaling introduced during data acquisition and processing of data is now discussed.

Assume that an accelerometer is used as the vibration signal sensor which gives out a voltage of  $K_1$  for every 32 ft/sec<sup>2</sup> of acceleration. (This acceleration is referred to as one "g.") Therefore, the input voltage from the vibration signal is given by the equation:

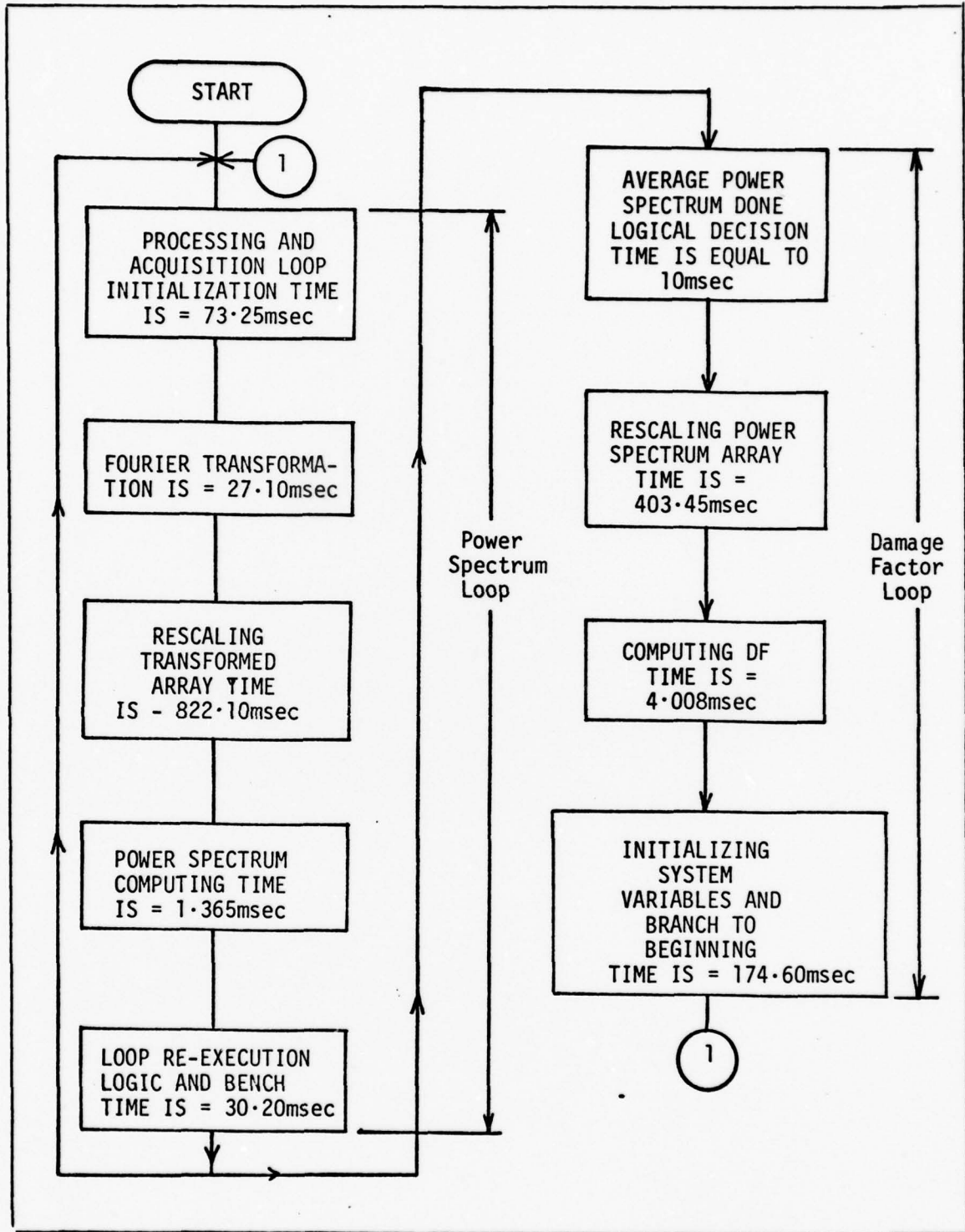


Figure 62. Execution Time for Control Executive

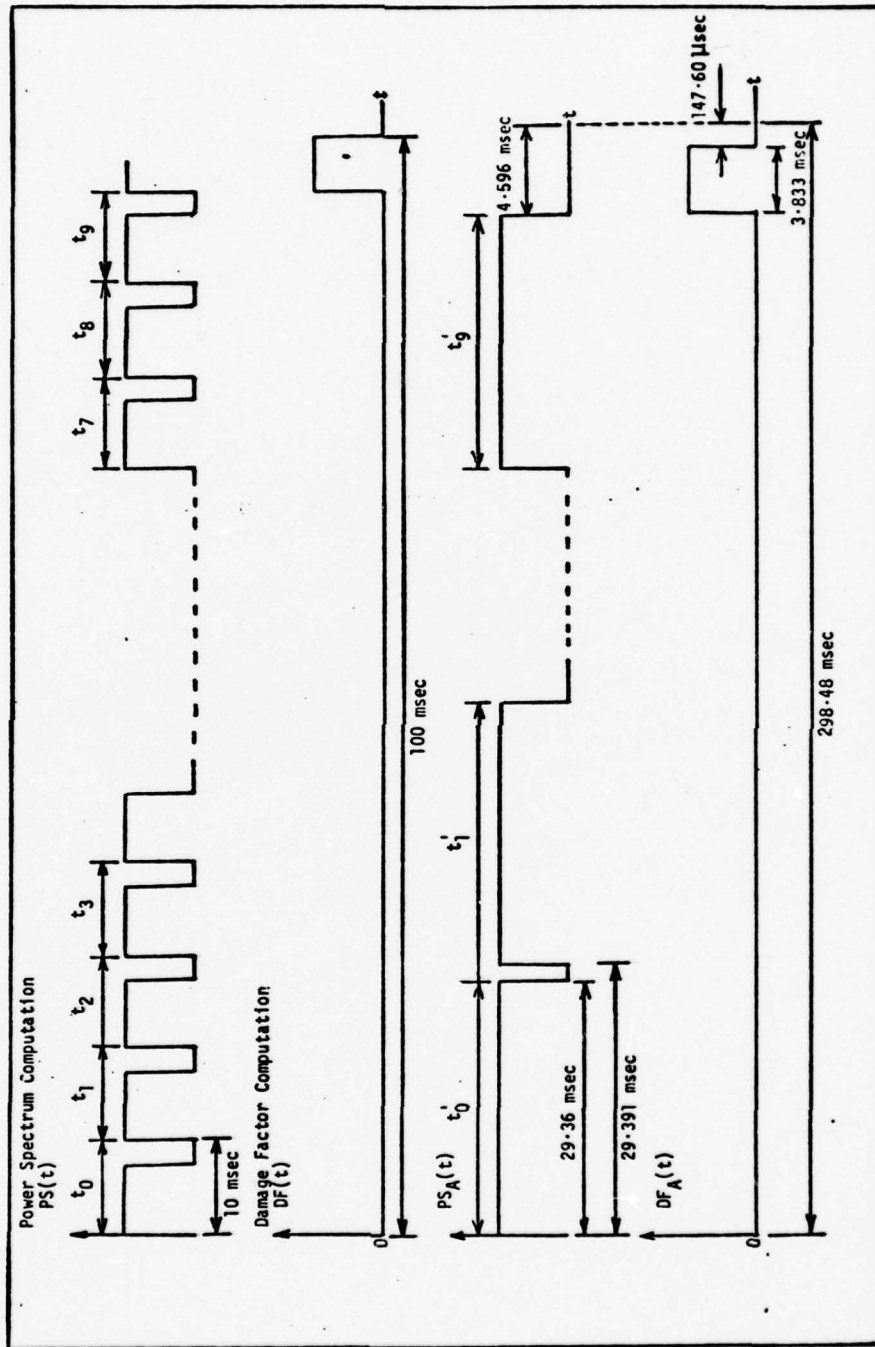


Figure 63. Timing Relationship Between Actual and Desired Execution Times

$$e(k) = K_1 G(k) \quad (51)$$

Here,  $G(k)$  is the acceleration encountered by the test piece.

The analog-to-digital convertor, ADV11-A, gives a 12-bit binary representation of the input voltage. Since the computation carried out in MIBDAPS assumes fractional numbers (Ref Chap IV), an input voltage of one volt would be represented by the fraction  $K_2$  which is defined below.

$$K_2 = \frac{(007777)_8 \times 2^{-15}}{5.12} \quad (52)$$

Therefore, each data point  $x(n)$  is given by the following equation:

$$x(k) = K_2 e(k) \quad (53)$$

The Fourier transformation of these data points yields the equivalent value in the frequency domain. Using Eq (6), we get

$$F(n) = \sum_{k=0}^{N-1} x(k) W^{nk}, \quad n = 0, 1, \dots, N-1 \quad (54)$$

since the Fourier transformation module (subroutine FOURIE) scales the data array to avoid memory overflow (Ref Chap IV). Therefore, the resultant Fourier transformed array is given by the following:

$$Z(n) = (SFACTR)F(n) \quad (55)$$

Here, SFACTR is the value by which the Fourier transformed array has been scaled. (This value in subroutine FOURIE has been fixed equal to 3.)

The power spectrum subroutine PRSPEC introduces another scale factor called CUMSF1 (Ref Chap IV). Also, since an average power spectrum density is computed, Eq (15) can be modified to give:

$$\tilde{A}'(n) = K3\bar{A}(n) \quad (56)$$

where  $K3 = \text{CUMSF}_i / p \quad (57)$

and  $A(n) = |Z(n)|^2 \quad (58)$

The cumulative damage factor computation introduced a scale factor called CDSFi (here  $i = 1, 2, \dots, 11$ ), and Eq (19) can be modified and rewritten as:

$$\text{CDF}(n) = K4[\tilde{A}']^{1.6} \quad (59)$$

where  $K4 = \frac{\text{CDSFi}}{[f(n)]^{5.4}} \quad (60)$

The equations derived here can be used to interpret the output array CDF versus frequency and temperature.

### Summary

In this chapter, the testing procedures used to check the software modules as well as the integrated system were discussed. The experimental results obtained for individual software modules and also for the overall system were presented. The scale factors introduced by each software module and the equations necessary to interpret the output array were also derived.

## VII. Recommendations and Conclusions

### Introduction

At present, a bench model of MIBDAPS exists; the development of a prototype model will require additional work. This effort would include developing software to test the system's accuracy, configuring the hardware in a mainframe, and incorporating the system error detection and recovery features. This chapter presents recommendations for future development and conclusions of the investigation.

### Recommendations

The recommendations for future work are listed in the following paragraphs:

- Software to test the performance and accuracy of results should be developed. This software should include provisions to generate input test signals which duplicate real world signals. The software should gather memory overflow statistics for varying operating intervals. These statistics would enable the determination of the fail-safe operating period.

- The software to output the CDF versus frequency and temperature arrays, at the conclusion of the system operation, needs to be developed. This software should either print the output arrays in a tabular form or plot CDF versus frequency and temperature contours. (The availability of a plotter is essential for the latter.)

- At present, the control executive includes a latent code which needs to be activated in the future. This latent code has been written to handle the acquisition of a temperature signal. It would also

compute the average temperature during the interval when the power spectral density is being computed.

- The ADC used in the bench model can handle 16 single-ended inputs. At present, only two inputs are being used for acquiring the analog temperature and analog vibration signals. Therefore, similar data from seven additional test pieces can be acquired and processed. The possibility of increasing the number of test pieces monitored by MIBDAPS should be explored. Changes in existing control executive would be required for handling the additional data.

- The need to minimize the existing MIBDAPS configuration is vital to those applications where size and weight of MIBDAPS is critical (Ref 13). In case this venue is explored for developing the prototype, then the following changes are suggested to the existing system:

a. Use of static RAM in place of core memory would decrease the weight and power requirement.

b. Use of custom-designed dual channel ADC would eliminate the need for the DEC manufactured ADC board.

c. Design of a dedicated FFT hardwired module would decrease the memory requirements for the software. The FFT integrated circuit chip R5601\*, manufactured by Reticon, would be a candidate device. If designed, the FFT module should be bus-compatible with the LSI-11M microcomputer and act as a DMA device.

- The DC power sense circuit which would sense DC supply variations needs to be designed. At switch ON, if the DC supply is within limits, this circuit would send a signal to the POWER OK interrupt generator.

\*Ref Preliminary Data Sheet: Quad Chipped Transversal Fitter, 27 March 1978.

In the event of power failure, this module would signal the POWER FAIL interrupt generator circuit which, in turn, would interrupt the processor (Ref Chap IV).

- An operational time recorder is required to keep track of the time for which the system remains in operation. There should be provisions to move the current value in the buffer register (Figure 56) to memory. This would enable the permanent record of total operational time. There should also be provisions to move the previously stored operational time back to the buffer register and recording the additional operational time. This would enable the system to record the aggregate operational time even if the system is required to monitor the same test piece at different intervals.

- The DC power backup for use in case of a main power failure should also be designed for the prototype. The backup power should maintain the system for the duration it takes to execute the power fail service routine.

### Conclusions

The bench model of MIBDAPS as implemented is the first step towards developing a prototype system. The initial design and implementation of different functions of MIBDAPS were successfully carried out. The choice of the LSI-11M microcomputer as the processor for MIBDAPS and the development of software to acquire and process data fulfilled the initial requirements placed on MIBDAPS (Ref Chap II).

The time constraints placed on the processing function of MIBDAPS were exceeded because of the software implementation of the FFT module. The individual software modules have been tested and found accurate

within the bounds of a 16-bit machine. Therefore, the testing of the integrated system would be a fairly easy task. The incorporation of the recommendation discussed earlier in this chapter would result in a prototype model. This prototype system would then accomplish the acquisition and processing of requisite data for damping the layer breakout design as envisioned by AFFDL.

The computation accuracy of the data acquisition and processing functions of the integrated system (MIBDAPS bench model) could not be determined because of the shortage of time. But because the performance of the individual software modules was acceptable, it is a fair assumption that the integrated system would perform accurately, too.

The designed system in its present form shows the practicability of using the state-of-the-art technology to accomplish the task of data acquisition and processing.

The application of such a system is not limited to monitoring aircraft hardware for vibration and temperature data only. It can be configured to act as an environmental recorder where additional analog signals need to be monitored and processed. This could be accomplished by making appropriate changes in the software. In its present form, the system is capable of computing the FFT and power spectrum of a signal; therefore, this feature can be used for spectrum analysis functions. This investigation has successfully demonstrated the feasibility of developing a microcomputer based system from the conceptual phase to a working bench-type model.

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APPENDIX A

STRUCTURED ANALYSIS AND DESIGN TECHNIQUE

## APPENDIX A

### Structured Analysis and Design Technique

This appendix (Ref 9) gives a short description of how structured analysis (SA) models are constructed and explains the SA diagram conventions used in this paper. It must be noted that the format used to present the models in this paper is not standard according to the rules developed by SofTech. The changes were made to present the models in a manner which is more familiar to readers who have no experience with SA models. Although the format is not that used by SofTech, the diagrams of the models are organized and related according to SofTech procedures, and the conventions used to construct individual diagrams are standard.

The structured analysis design technique (SADT) is a general purpose top-down, modular technique for modeling functions. The functions may be as varied as farming or manufacturing, but SA was developed primarily as a software requirements definition and design tool. Although a complete SA model actually consists of two models, one for activities and one for data, this paper employs only activity models so the conventions described here are those which apply to activity models.

An SA activity model consists of a series of diagrams which present in progressively more detail the activities necessary to perform some function. Each diagram represents a self-contained activity which is part of the overall function. A diagram shows how its activity is decomposed into subactivities and how the subactivities are related to each other. The subactivities in each diagram may then be decomposed on separate diagrams which leads to a tree structure of several levels. At the top is one diagram which represents the whole function, and at the bottom

Figure A-1 shows how an SA model would appear if all the diagrams were on one page. Of course, in real SA diagrams only one level of decomposition is shown, but the figure demonstrates the top-down nature of SA and the way activities are grouped into modules. In the figure, as in real models, one large box represents the whole function, and that is decomposed into successive levels of related activities. The decomposition process continues until the desired amount of detail has been developed, which may require more levels than shown in Figure A-1. Another thing to note is that while the figure shows only three subactivities in each decomposition, any number from three to six is acceptable.

From Figure A-1, it should be apparent that SA diagrams are constructed with boxes and arrows. In an activity model, each box represents an activity and is called a "node." Arrows represent data where the word "data" is used in a very general sense to include anything that is not an activity. Figure A-2 shows the different meanings given to arrows depending on which side of a box they enter or leave. An input is data that is modified by the activity to produce an output. A control is data which may or may not be converted into output, but which in some way restricts the activity (starts or stops it, for example). Every box must have at least one control arrow. A mechanism is a person or thing which acts as a processor. Mechanism arrows are often omitted when the processor is the same for all nodes. No limit is placed on the number of arrows which may interface with a side of a box, but it is common practice to group related types of data.

Between boxes, arrows may split and join. In general, all branches of an arrow contain the same data unless a branch is given a separate label. This convention is summarized in Figure A-3, which also gives two

forms of OR branches. The OR branches are used to show that data follows one path or the other, but not both.

When two nodes are related so that the output of each is a control for the other, a special two-way arrow may be used. Figure A-4 shows a mutual control situation with a two-way arrow and the equivalent form with normal allows. An arrow showing mutual control has two labels separated by a slash; the first label identifies data going forward, and the second is the feedback data.

A special numbering system is used to distinguish between nodes at different levels and between nodes at the same level. In an activity model, node numbers are prefixed with the letter "A." For preliminary nodes, A is followed by a dash and a number. Node A-1 may be used to show the model in relationship to other functions (Figure 3). Node A-0 serves as a cover sheet for the model; the node is simply a box showing inputs, outputs, controls, and mechanisms for the function which the model is to describe (Figure 4). Decomposition begins in node A0. Note in Figure 5 that each box of the decomposition is numbered; the boxes on all decomposition diagrams are numbered, and this number is used to form the node number. For the activities subordinated to node A0, the node number is simply the box number on A0, Process Data in Figure 15, for example, becomes node A3. From this level on, the node number is a combination of the node number of the parent diagram and the box number of the subordinate. As an example, the decomposition of Process Data is given in Figure 26. Box 2, Compute Fourier Transform, is assigned the node number A32. Subordinates of A32, if diagrammed, would have the numbers A321, A322, and so on through the last box number.

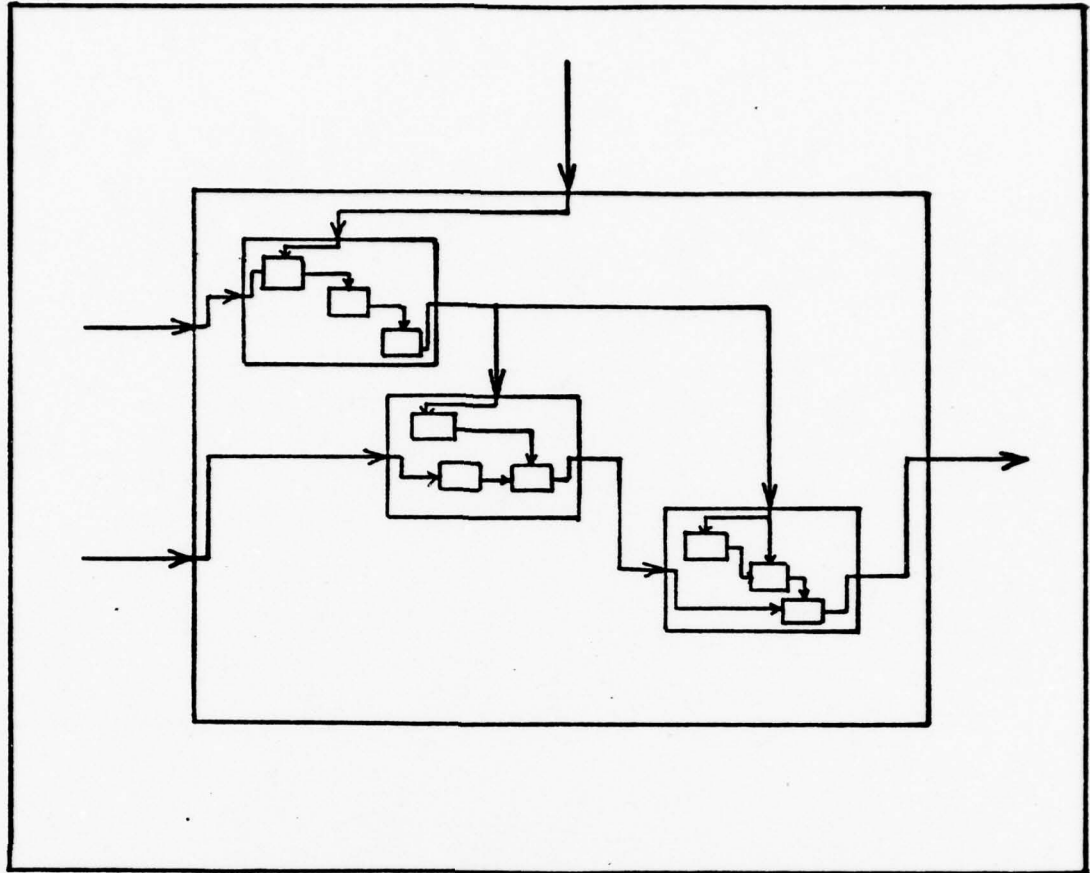


Figure A-1. Top-down View of an SA Model (Ref 9)

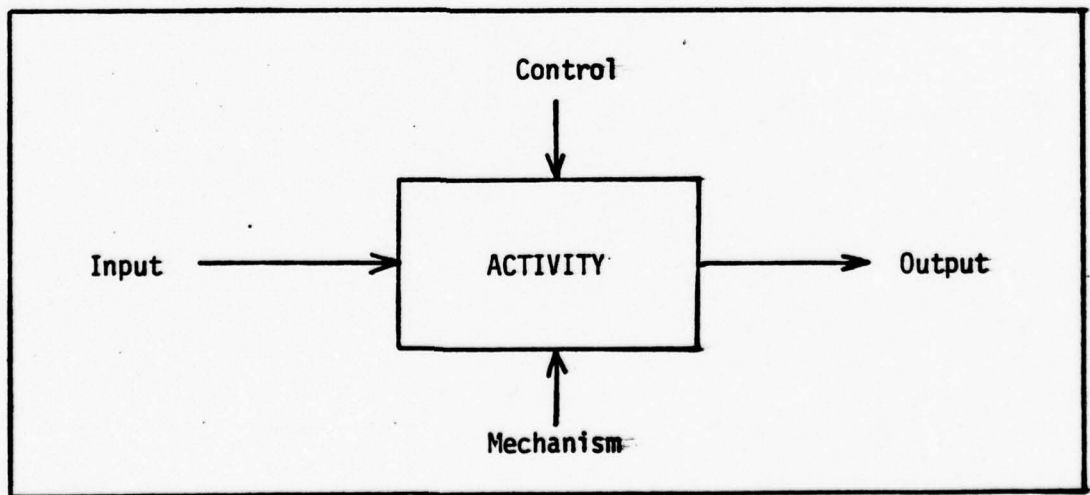


Figure A-2. Arrow Definitions (Ref 9)

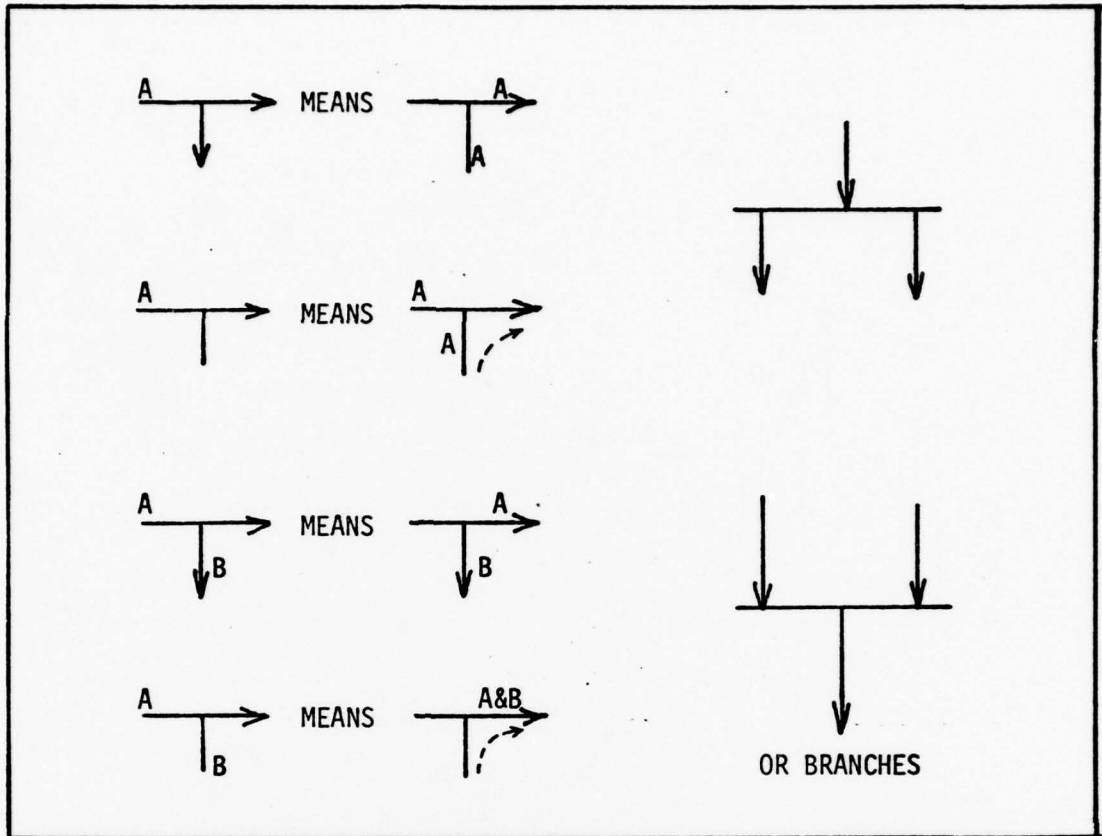


Figure A-3. Arrow Branches

(Ref 9)

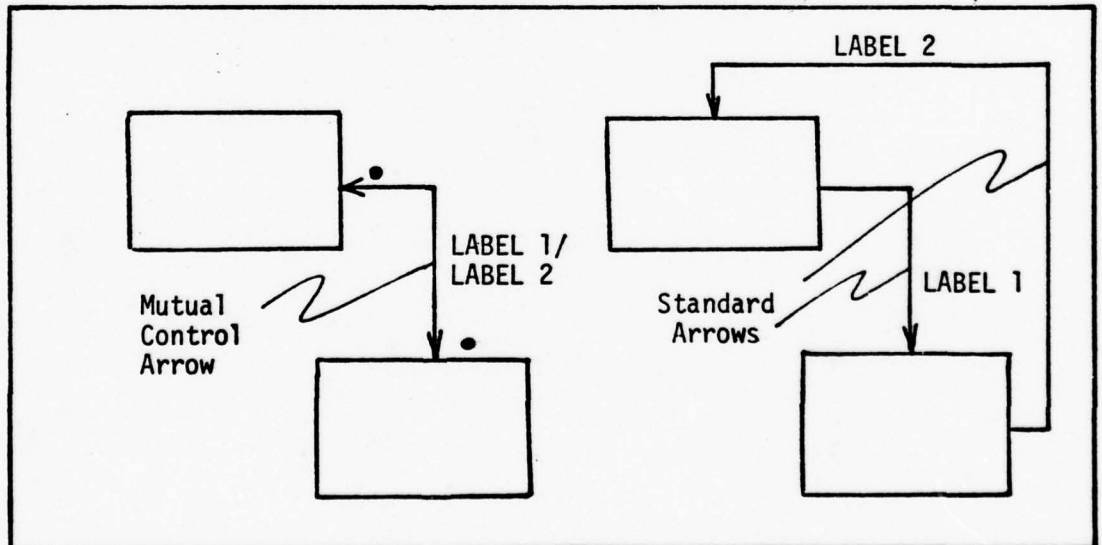


Figure A-4. Arrows Showing Mutual Control

(Ref 9)

A special code called an ICOM code (input, control, output, mechanism) is used to identify arrows. The code contains a number, a letter, and another number (Figure A-5).

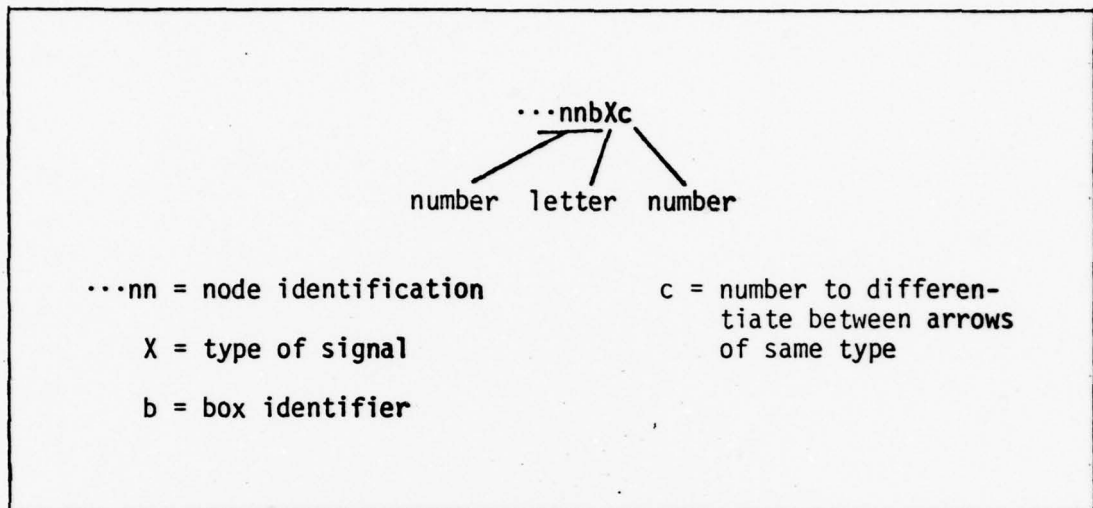


Figure A-5. ICOM Code

The first number has two parts: the digit immediately preceding the letter is that of source or destination box, and the digits preceding this number identify the node on which the box is located. The letter refers to the type of arrow: I for input, C for control, O for output, and M for mechanism. The last number distinguishes between arrows of the same type on a box. One exception to the rule is node A $\emptyset$  where the ICOM code has no number preceding the letter; this has been done to differentiate between signals within the system and signals which are input, output, or control from the outside world. If any signal that appears on node A $\emptyset$  is also present on a subordinate node, the identifying ICOM code remains unchanged.

A few important points about the text describing each SA diagram must be included here. The text is intended to point out the highlights

of a diagram and not repeat all the details. As an aid to following the discussion, the ICOM code is included in parenthesis following any reference to specific diagram features. Finally, the text describing the diagrams in this paper includes I/O specifications, which are not a standard SA practice.

APPENDIX B

INSTRUCTION TIMING

## APPENDIX B

### Instruction Timing

#### LSI-11M Instruction Execution Time (Ref 10)

The execution time for an instruction depends on the instruction itself, the modes of addressing used, and the type of memory referenced. In most cases, the instruction execution time is the sum of a basic time, a source address (SRC) time, and a destination address (DST) time.

$$\text{Instr Time} = \text{Basic Time} + \text{SRC Time} + \text{DST Time}$$

$$(\text{Basic Time} = \text{Fetch Time} + \text{Decode Time} + \text{Execute Time})$$

Some of the instructions require only some of these times. All timing information is in microseconds, unless otherwise noted. Times are typical; processing timing can vary +20%.

The following tables give the basic time for the LSI-11M instruction set.

TABLE B-I  
Source and Destination Times

Mode	SRC Time (Word)	SRC Time (Byte)	DST Time (Word)	DST Time (Byte)
0	0	0	0	0
1	1.05 $\mu$ s	1.05 $\mu$ s	1.30 $\mu$ s	1.30 $\mu$ s
2	1.55	1.80	1.55	1.80
3	2.55	2.55	2.55	2.55
4	1.80	1.80*	1.80	1.80*
5	3.35	3.35	3.35	3.35
6	2.85	2.85	2.85	2.85
7	4.35	4.35	4.35	4.35

\*If R6 or R7 is used with Mode 4 and Byte operation, add 0.25 $\mu$ s to SRC time and 0.25 $\mu$ s to DST time.

TABLE B-II  
Basic Time (DOP)

Double Operand Instruction	SM0 DM0	SM1-7 DM0	SM0 DM1-7	SM1-7 DM1-7
MOV, ADD	1.75 $\mu$ s	2.55 $\mu$ s	4.55 $\mu$ s	4.55 $\mu$ s
MOVB	2.25 <sup>(1)</sup>	4.05 <sup>(1)</sup>	4.55 <sup>(2)</sup>	5.55 <sup>(2)</sup>
SUB	2.80	2.80	4.80	4.80
CMP, BIT	3.05	3.05	3.05	3.05
CMPB, BITB	2.55	3.55	3.80	2.80
BIC	2.80	2.80	2.80	2.80
BIS	3.05	3.05	4.30	4.30
BICB, BISB	2.80	3.80	5.55 <sup>(2)</sup>	6.55 <sup>(2)</sup>
XOR	2.30	-	4.05	-

(1) Add 0.25 $\mu$ s if effective SRC byte operand is negative.  
(2) Add 0.25 $\mu$ s if DST byte address is odd.

TABLE B-III  
Basic Time (SOP)

Single Operand Instruction	DM0	DM1-7
CLR, INC, DEC, SWAB	2.05 $\mu$ s	4.05 $\mu$ s
ADC	2.30	3.05
SBC	2.30 <sup>(5)</sup>	3.05 <sup>(7)</sup>
COM, NEG, ROL, ASL, ROR, ASR, CLRB, SXT	2.30	4.30
TST	2.30	2.30
INCB	2.55 <sup>(1)</sup>	5.05 <sup>(3)</sup>
ADCB	2.30 <sup>(5)(9)</sup>	3.80 <sup>(7)(10)</sup>
DECB	2.55 <sup>(2)</sup>	5.05 <sup>(3)</sup>
SBCB	2.30 <sup>(5)(11)</sup>	3.80 <sup>(8)(10)</sup>
COMB	2.30	5.30 <sup>(3)</sup>
NEGB	2.80 <sup>(2)</sup>	5.30 <sup>(3)</sup>
ROLB, ASLB, RORB, ASRB	3.55	5.55 <sup>(4)</sup>
TSTB	3.05	3.05
MFPS (1067DD)	3.05	4.30 <sup>(3)</sup>
MTPS (1064SS)*	3.55	4.05
JMP	10.44	2.50
JRS	10.54	5.00

\*For MTPS, use byte DST times, not SRC times.

(1) Add 0.25 $\mu$ s if register byte operand = 377 octal.  
(2) Add 0.25 $\mu$ s if register byte operand = 0.  
(3) Add 0.25 $\mu$ s if DST byte address is odd.  
(4) Add 0.25 $\mu$ s if DST byte address is even.  
(5) Add 0.25 $\mu$ s if C-bit is set.  
(6) Add 1.50 $\mu$ s if C-bit is set.  
(7) Add 1.75 $\mu$ s if C-bit is set.  
(8) Add 2.00 $\mu$ s if C-bit is set.  
(9) Add 0.25 $\mu$ s if C-bit is set and register byte operand = 377<sub>8</sub>.  
(10) Add 0.25 $\mu$ s if C-bit is set and DST byte address is odd.  
(11) Add 0.25 $\mu$ s if C-bit is set and register byte operand = 0.

TABLE B-IV  
Miscellaneous Instruction

Instruction	Basic Time
SOB (BRANCH)	3.05 $\mu$ s
SOB (NO BRANCH)	2.55
SET CC	2.05
CLEAR CC	2.30
NOP	2.30
RTS	3.30
MARK	4.10
RTI	5.60 <sup>(1)</sup>
RTT	5.60 <sup>(2)</sup>
TRAP, EMT, IOT, BPT	10.15
WAIT	2.75 <sup>(1)</sup>
RESET	3.25+10 $\mu$ sINIT+90 $\mu$ s

(1) Add 0.50 $\mu$ s if NEW PS has bit 4 (T-bit) set.  
(2) Add 0.25 $\mu$ s if NEW PS has bit 4 (T-bit) set.

TABLE B-V  
Basic Time (Branch)

Branch Instruction	Branch <sup>(1)</sup>	No Branch
BR	2.80	-
BNE, BEQ, BPL, BMI, BVC, BVS, BCC/BLO, BCS/BHOS	3.05	2.00
BGE	3.30 <sup>(2)</sup>	2.25 <sup>(2)</sup>
BLT	3.30 <sup>(3)</sup>	2.25 <sup>(3)</sup>
BGT	3.80 <sup>(2)</sup>	2.00 <sup>(4)</sup>
BLE	3.05 <sup>(5)</sup>	2.75 <sup>(3)</sup>
BHI	3.55	2.00 <sup>(6)</sup>
BLOS	3.05 <sup>(6)</sup>	2.50
<p>(1) Add 0.25<math>\mu</math>s if offset is negative.            (2) Add 0.25<math>\mu</math>s if N-bit is set.            (3) Add 0.25<math>\mu</math>s if V-bit is clear.            (4) If Z-bit clear: add 0.75<math>\mu</math>s if N-bit is clear; add 1.00<math>\mu</math>s if N-bit is set.            (5) If Z-bit clear: add 0.75<math>\mu</math>s if V-bit is set; add 1.00<math>\mu</math>s if V-bit is clear.            (6) Add 0.05<math>\mu</math>s if Z-bit is clear.</p>		

TABLE B-VI  
Basic Time (EIS)

EIS Instruction*	Basic Time
MUL	16.70 $\mu$ s worst case
DIV	23.40 $\mu$ s worst case
ASH (RIGHT)	5.80 $\mu$ s + 0.25 $\mu$ s per shift
ASH (LEFT)	5.55 $\mu$ s + 0.75 $\mu$ s per shift
ASHC (RIGHT)	6.60 $\mu$ s + 0.25 $\mu$ s per shift
ASHC (LEFT)	6.35 $\mu$ s + 0.75 $\mu$ s per shift

\*Use word DST times, not SRC times.

TABLE B-VII  
Basic Time (FIS)\*

FIS Instruction	Basic Time
FADD	20.15 typical 22.80 worst case
FSUB	20.40 typical 23.00 worst case
FMUL	50.25 typical 69.40 worst case
FDIV	69.40 typical 70.10 worst case

\*Inst Time (FIS) = Basic Time + Shift Time for Binary Points + Shift Time for normalization.  
Binary Point Alignment: 0.25 $\mu$ s per shift.  
Normalization: 0.75 $\mu$ s per shift.

APPENDIX C

THEORETICAL DEVELOPMENT OF A BASE-2 FFT ALGORITHM

## APPENDIX C

### Theoretical Development of a Base-2 FFT Algorithm

#### Development of the Cooley-Tukey FFT Algorithm (Ref 2:257-265)

The DFT of N-data can be written as:

$$X(n) = \sum_{k=0}^{N-1} x(k)W^{nk}, \quad n = 0, 1, \dots, N-1 \quad (C-1)$$

where

$$W^{nk} = \exp \frac{-2\pi ink}{N}, \quad n = 0, 1, \dots, N-1, \quad k = 0, 1, \dots, N-1$$

The direct evaluation of Eq (C-1) requires the solution of the following set of equations.

$$X(0)^1 = W^0 x(0) + W^0 x(1) + W^0 x(2) \dots + W^0 x(N-1)$$

$$X(1) = W^0 x(0) + W^1 x(1) + W^2 x(2) \dots + W^{N-1} x(N-1)$$

$$X(2) = W^0 x(0) + W^2 x(1) + W^4 x(2) \dots + W^{2N-2} x(N-1)$$

$$X(3) = W^0 x(0) + W^3 x(1) + W^6 x(2) \dots + W^{3N-3} x(N-1)$$

⋮

$$X(N-1) = W^0 x(0) + W^{N-1} x(1) + W^{2N-2} x(2) \dots + W^{(N-1)N-(N-1)} x(N-1)$$

The above set of equations can also be written in the matrix form given below:

$$\begin{bmatrix} X(0) \\ X(1) \\ X(2) \\ X(3) \\ \vdots \\ X(N-1) \end{bmatrix} = \begin{bmatrix} W^0 & W^0 & W^0 & \dots & W^0 \\ W^0 & W^1 & W^2 & \dots & W^{N-1} \\ W^0 & W^2 & W^4 & \dots & W^{2N-2} \\ W^0 & W^3 & W^6 & \dots & W^{3N-3} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ W & W^{N-1} & W^{2N-2} & \dots & W^{(N-1)N-(N-1)} \end{bmatrix} \begin{bmatrix} X(0) \\ X(1) \\ X(2) \\ X(3) \\ \vdots \\ X(N-1) \end{bmatrix}$$

Therefore, in order to compute a DFT of N-data points, we have to solve the matrix given by Eq (C-2), all terms of this matrix being complex:

$$[X(n)] = [W^{nk}][x(k)] \quad (C-2)$$

This direct evaluation requires  $N^2$  complex multiplications and  $N(N-1)$  complex additions (Ref-3:151). The FFT algorithm developed by Cooley and Tukey (Ref 3:151) reduces the number of complex multiplications and additions, and computer execution of this algorithm results in faster computation time as compared with the evaluation of Eq (C-1) (Ref 3:151).

In case N is limited to a power of 2 (i.e.,  $N = 2^P$ ), it is possible to express k and n in terms of the index p, as a binary weighted series:

$$k = k_{p-1}2^{p-1} + k_{p-2}2^{p-2} + \dots + k_12^1 + k_0 \quad (C-3)$$

$$n = n_{p-1}2^{p-1} + n_{p-2}2^{p-2} + \dots + n_12^1 + n_0 \quad (C-4)$$

Here  $k_{p-1}, k_{p-2}, \dots, k_0$  and  $n_{p-1}, n_{p-2}, \dots, n_0$  can take a value of either 0 or 1. In this way, we can express all the N possible values of the indices in terms of a binary number and, hence, facilitate the consideration of storage in the digital computer. Using this convention, Eq (C-1) can be rewritten in the following form:

$$X(n_{p-1}, n_{p-2}, \dots, n_0) = \sum_{k_0=0}^1 \sum_{k_1=0}^1 \dots \sum_{k_{p-1}=0}^1 x(k_{p-1}, k_{p-2}, \dots, k_0) \times W^{n(k_{p-1}2^{p-1} + k_{p-2}2^{p-2} + \dots + k_0)} \quad (C-5)$$

If the first term of the exponential for Eq (C-3) is expanded, Eq (C-6) is the resulting equation:

$$W^{nk_{p-1}2^{p-1}} = W^{n_{p-1}2^{p-1}k_{p-1}2^{p-1}} + W^{n_{p-2}2^{p-2}k_{p-1}2^{p-1}} + \dots + W^{n_12^1k_{p-1}2^{p-1}} + W^{n_0k_{p-1}2^{p-1}} \quad (C-6)$$

This can be simplified to give Eq (C-7) because all other terms go to unity. This enables

$$W^{nk_{p-1}2^{p-1}} = W^{n_0 k_{p-1} 2^{p-1}} \quad (C-7)$$

the innermost sum of Eq (C-5) over  $k_{p-1}$  to be written as a shorter Fourier transform:

$$X_1(n_0, k_{p-2}, \dots, k_0) = \sum_{k_{p-1}=0}^{N-1} x_1(k_{p-1}, k_{p-2}, \dots, k_0) W^{n_0 k_{p-1} 2^{p-1}} \quad (C-8)$$

Unlike the complete Fourier transform, this sum consists of a set of  $N$  numbers only, each calculated from two of the original data points. Subsequent sums proceeding outwards in Eq (C-5) can be calculated using a generalized recursive expression for the exponential term:

$$W^{nk_{p-q}2^{p-q}} = W^{(n_{q-1}2^{q-1} + n_{q-2}2^{q-2} + \dots + n_0)k_{p-q}2^{p-q}} \quad (C-9)$$

here

$$q = 1, 2, 3 \dots p$$

The successive sums are evaluated according to the equation

$$X_q(n_0, n_1, \dots, n_{q-1}, k_{p-q-1}, k_{p-q-2}, \dots, k_0) = \sum_{k_{p-q}=0}^{N-1} x_{q-1}(n_0, n_1, \dots, n_{q-2}, k_{p-q}, \dots, k_0) \times W^{(n_{q-1}2^{q-1} + n_{q-2}2^{q-2} + \dots + n_0)k_{p-q}2^{p-q}} \quad (C-10)$$

here

$$q = 1, 2, 3 \dots p$$

To apply this recursive formulae, the initial set of data,  $x(k)$ , is first made equal to  $x_0(k)$  so that  $q = 1$ , thus

$$x(k) = x(k_{p-1}, k_{p-2} \dots k_0) = x_0(k_{p-1}, k_{p-2} \dots k_0) \quad (C-11)$$

This leads to the derivation of succeeding arrays in  $x_q$  so that the final array will be in  $X(n)$ . Equation (C-12) given below shows that the elements of  $X(n)$  will be calculated in incorrect order, so some sort of shuffling is required to get the result in correct order. (This point is clarified when the algorithm for a 16-point DFT using this FFT method is worked.)

$$X(n_{p-1}, n_{p-2} \dots n_0) = x_p(n_0, \dots, n_{p-2}, n_{p-1}) \quad (C-12)$$

#### A 16-Point FFT Algorithm

In case the number of data points to be transferred is 16, the value of  $P = 4$ . Therefore,  $k, n$  can be expressed as 4-bit binary numbers as given in the following equations:

$$k = 8k_3 + 4k_2 + 2k_1 + k_0 \quad (C-13)$$

$$n = 8n_3 + 4n_2 + 2n_1 + n_0 \quad (C-14)$$

Using Eqs (C-13) and (C-14), Eq (C-5) can be written as

$$X(n_3, n_2, n_1, n_0) = \sum_{k_0=0}^1 \sum_{k_1=0}^1 \sum_{k_2=0}^1 \sum_{k_3=0}^1 x(k_3, k_2, k_1, k_0) \times W^{n(8k_3 + 4k_2 + 2k_1 + k_0)} \quad (C-15)$$

Using Eqs (C-7) and (C-8), the innermost summation can be written as

$$x_1(n_0, k_2, k_1, k_0) = \sum_{k_3=0}^1 x_0(k_3, k_2, k_1, k_0) W^{8k_3 n_0} \quad (C-16)$$

Subsequent summations can be written using Eqs (C-9) and (C-10):

$$x_2(n_0, n_1, k_1, k_0) = \sum_{k_2=0} x_1(n_0, k_2, k_1, k_0) W^{4k_2(2n_1 + n_0)} \quad (C-17)$$

$$x_3(n_0, n_1, n_2, k_0) = \sum_{k_1=0} x_2(n_0, n_1, k_1, k_0) W^{2k_1(4n_2 + 2n_1 + n_0)} \quad (C-18)$$

$$x_4(n_0, n_1, n_2, n_3) = \sum_{k_0=0} x_3(n_0, n_1, n_2, k_0) W^{k_0(8n_3 + 4n_2 + 2n_1 + n_0)} \quad (C-19)$$

$$X(n_3, n_2, n_1, n_0) = x_4(n_0, n_1, n_2, n_3) \quad (C-20)$$

The set of equations that result when summations of Eqs (C-16), (C-17), (C-18), (C-19), and (C-20) are expanded is given in tables C-I, C-II, C-III, C-IV, and C-V, respectively.

Table C-V shows that the final array for  $X(n)$  is in bit-reversed order; therefore, to get the result in order, a post-shuffling needs to be done. Another way to overcome this problem is to bit invert the input array  $[x_0(k)]$ , and the output will automatically be in order (Ref 3:178). The signal flow graph for the 16-point FFT without bit inverting the data is shown in Figure C-1, and Figure C-2 shows the signal flow graph for data which has been bit inverted before being transformed.

TABLE C-I

Expansion Results of Eq (C-16)

$$\begin{aligned}
 x_1(0, 0, 0, 0) &= x_0(0, 0, 0, 0)w^0 + x_0(1, 0, 0, 0)w^0 \\
 x_1(0, 0, 0, 1) &= x_0(0, 0, 0, 1)w^0 + x_0(1, 0, 0, 1)w^0 \\
 x_1(0, 0, 1, 0) &= x_0(0, 0, 1, 0)w^0 + x_0(1, 0, 1, 0)w^0 \\
 x_1(0, 0, 1, 1) &= x_0(0, 0, 1, 1)w^0 + x_0(1, 0, 1, 1)w^0 \\
 x_1(0, 1, 0, 0) &= x_0(0, 1, 0, 0)w^0 + x_0(1, 1, 0, 0)w^0 \\
 x_1(0, 1, 0, 1) &= x_0(0, 1, 0, 1)w^0 + x_0(1, 1, 0, 1)w^0 \\
 x_1(0, 1, 1, 0) &= x_0(0, 1, 1, 0)w^0 + x_0(1, 1, 1, 0)w^0 \\
 x_1(0, 1, 1, 1) &= x_0(0, 1, 1, 1)w^0 + x_0(1, 1, 1, 1)w^0 \\
 x_1(1, 0, 0, 0) &= x_0(0, 0, 0, 0)w^0 + x_0(1, 0, 0, 0)w^8 \\
 x_1(1, 0, 0, 1) &= x_0(0, 0, 0, 1)w^0 + x_0(1, 0, 0, 1)w^8 \\
 x_1(1, 0, 1, 0) &= x_0(0, 0, 1, 0)w^0 + x_0(1, 0, 1, 0)w^8 \\
 x_1(1, 0, 1, 1) &= x_0(0, 0, 1, 1)w^0 + x_0(1, 0, 1, 1)w^8 \\
 x_1(1, 1, 0, 0) &= x_0(0, 1, 0, 0)w^0 + x_0(1, 1, 0, 0)w^8 \\
 x_1(1, 1, 0, 1) &= x_0(0, 1, 0, 1)w^0 + x_0(1, 1, 0, 1)w^8 \\
 x_1(1, 1, 1, 0) &= x_0(0, 1, 1, 0)w^0 + x_0(1, 1, 1, 0)w^8 \\
 x_1(1, 1, 1, 1) &= x_0(0, 1, 1, 1)w^0 + x_0(1, 1, 1, 1)w^8
 \end{aligned}$$

TABLE C-II  
Expansion Results of Eq (C-17)

$x_2(0, 0, 0, 0) = x_1(0, 0, 0, 0)w^0 + x_1(0, 1, 0, 0)w^0$
$x_2(0, 0, 0, 1) = x_1(0, 0, 0, 1)w^0 + x_1(0, 1, 0, 1)w^0$
$x_2(0, 0, 1, 0) = x_1(0, 0, 1, 0)w^0 + x_1(0, 1, 1, 0)w^0$
$x_2(0, 0, 1, 1) = x_1(0, 0, 1, 1)w^0 + x_1(0, 1, 1, 1)w^0$
$x_2(0, 1, 0, 0) = x_1(0, 0, 0, 0)w^0 + x_1(0, 1, 0, 0)w^8$
$x_2(0, 1, 0, 1) = x_1(0, 0, 0, 1)w^0 + x_1(0, 1, 0, 1)w^8$
$x_2(0, 1, 1, 0) = x_1(0, 0, 1, 0)w^0 + x_1(0, 1, 1, 0)w^8$
$x_2(0, 1, 1, 1) = x_1(0, 0, 1, 1)w^0 + x_1(0, 1, 1, 1)w^8$
$x_2(1, 0, 0, 0) = x_1(1, 0, 0, 0)w^0 + x_1(1, 1, 0, 0)w^4$
$x_2(1, 0, 0, 1) = x_1(1, 0, 0, 1)w^0 + x_1(1, 1, 0, 1)w^4$
$x_2(1, 0, 1, 0) = x_1(1, 0, 1, 0)w^0 + x_1(1, 1, 1, 0)w^4$
$x_2(1, 0, 1, 1) = x_1(1, 0, 1, 1)w^0 + x_1(1, 1, 1, 1)w^4$
$x_2(1, 1, 0, 0) = x_1(1, 0, 0, 0)w^0 + x_1(1, 1, 0, 0)w^{12}$
$x_2(1, 1, 0, 1) = x_1(1, 0, 0, 1)w^0 + x_1(1, 1, 0, 1)w^{12}$
$x_2(1, 1, 1, 0) = x_1(1, 0, 1, 0)w^0 + x_1(1, 1, 1, 0)w^{12}$
$x_2(1, 1, 1, 1) = x_1(1, 0, 1, 1)w^0 + x_1(1, 1, 1, 1)w^{12}$

TABLE C-III

Expansion Result of Eq (C-18)

$$\begin{aligned}
x_3(0, 0, 0, 0) &= x_2(0, 0, 0, 0)w^0 + x_2(0, 0, 1, 0)w^0 \\
x_3(0, 0, 0, 1) &= x_2(0, 0, 0, 1)w^0 + x_2(0, 0, 1, 1)w^0 \\
x_3(0, 0, 1, 0) &= x_2(0, 0, 0, 0)w^0 + x_2(0, 0, 1, 0)w^8 \\
x_3(0, 0, 1, 1) &= x_2(0, 0, 0, 1)w^0 + x_2(0, 0, 1, 1)w^8 \\
x_3(0, 1, 0, 0) &= x_2(0, 1, 0, 0)w^0 + x_2(0, 1, 1, 0)w^4 \\
x_3(0, 1, 0, 1) &= x_2(0, 1, 0, 1)w^0 + x_2(0, 1, 1, 1)w^4 \\
x_3(0, 1, 1, 0) &= x_2(0, 1, 0, 0)w^0 + x_2(0, 1, 1, 0)w^{12} \\
x_3(0, 1, 1, 1) &= x_2(0, 1, 0, 1)w^0 + x_2(0, 1, 1, 1)w^{12} \\
x_3(1, 0, 0, 0) &= x_2(1, 0, 0, 0)w^0 + x_2(1, 0, 1, 0)w^2 \\
x_3(1, 0, 0, 1) &= x_2(1, 0, 0, 1)w^0 + x_2(1, 0, 1, 1)w^2 \\
x_3(1, 0, 1, 0) &= x_2(1, 0, 0, 0)w^0 + x_2(1, 0, 1, 0)w^{10} \\
x_3(1, 0, 1, 1) &= x_2(1, 0, 0, 1)w^0 + x_2(1, 0, 1, 1)w^{10} \\
x_3(1, 1, 0, 0) &= x_2(1, 1, 0, 0)w^0 + x_2(1, 1, 1, 0)w^6 \\
x_3(1, 1, 0, 1) &= x_2(1, 1, 0, 1)w^0 + x_2(1, 1, 1, 1)w^6 \\
x_3(1, 1, 1, 0) &= x_2(1, 1, 0, 0)w^0 + x_2(1, 1, 1, 0)w^{14} \\
x_3(1, 1, 1, 1) &= x_2(1, 1, 0, 1)w^0 + x_2(1, 1, 1, 1)w^{14}
\end{aligned}$$

TABLE C-IV

Expansion Result of Eq (C-19)

$$\begin{aligned}
x_4(0, 0, 0, 0) &= x_3(0, 0, 0, 0)w^0 + x_3(0, 0, 0, 1)w^0 \\
x_4(0, 0, 0, 1) &= x_3(0, 0, 0, 0)w^0 + x_3(0, 0, 0, 1)w^8 \\
x_4(0, 0, 1, 0) &= x_3(0, 0, 1, 0)w^0 + x_3(0, 0, 1, 1)w^4 \\
x_4(0, 0, 1, 1) &= x_3(0, 0, 1, 0)w^0 + x_3(0, 0, 1, 1)w^{12} \\
x_4(0, 1, 0, 0) &= x_3(0, 1, 0, 0)w^0 + x_3(0, 1, 0, 1)w^2 \\
x_4(0, 1, 0, 1) &= x_3(0, 1, 0, 0)w^0 + x_3(0, 1, 0, 1)w^{10} \\
x_4(0, 1, 1, 0) &= x_3(0, 1, 1, 0)w^0 + x_3(0, 1, 1, 1)w^6 \\
x_4(0, 1, 1, 1) &= x_3(0, 1, 1, 0)w^0 + x_3(0, 1, 1, 1)w^{14} \\
x_4(1, 0, 0, 0) &= x_3(1, 0, 0, 0)w^0 + x_3(1, 0, 0, 1)w^1 \\
x_4(1, 0, 0, 1) &= x_3(1, 0, 0, 0)w^0 + x_3(1, 0, 0, 1)w^9 \\
x_4(1, 0, 1, 0) &= x_3(1, 0, 1, 0)w^0 + x_3(1, 0, 1, 1)w^5 \\
x_4(1, 0, 1, 1) &= x_3(1, 0, 1, 0)w^0 + x_3(1, 0, 1, 1)w^{13} \\
x_4(1, 1, 0, 0) &= x_3(1, 1, 0, 0)w^0 + x_3(1, 1, 0, 1)w^3 \\
x_4(1, 1, 0, 1) &= x_3(1, 1, 0, 0)w^0 + x_3(1, 1, 1, 1)w^{11} \\
x_4(1, 1, 1, 0) &= x_3(1, 1, 1, 0)w^0 + x_3(1, 1, 1, 1)w^7 \\
x_4(1, 1, 1, 1) &= x_3(1, 1, 1, 0)w^0 + x_3(1, 1, 1, 1)w^{15}
\end{aligned}$$

TABLE C-V

Expansion Result of Eq (C-20)

$$x(0, 0, 0, 0) = x_4(0, 0, 0, 0)$$

$$x(0, 0, 0, 1) = x_4(1, 0, 0, 0)$$

$$x(0, 0, 1, 0) = x_4(0, 1, 0, 0)$$

$$x(0, 0, 1, 1) = x_4(1, 1, 0, 0)$$

$$x(0, 1, 0, 0) = x_4(0, 0, 1, 0)$$

$$x(0, 1, 0, 1) = x_4(1, 0, 1, 0)$$

$$x(0, 1, 1, 0) = x_4(0, 1, 1, 0)$$

$$x(0, 1, 1, 1) = x_4(1, 1, 1, 0)$$

$$x(1, 0, 0, 0) = x_4(0, 0, 0, 1)$$

$$x(1, 0, 0, 1) = x_4(1, 0, 0, 1)$$

$$x(1, 0, 1, 0) = x_4(0, 1, 0, 1)$$

$$x(1, 0, 1, 1) = x_4(1, 1, 0, 1)$$

$$x(1, 1, 0, 0) = x_4(0, 0, 1, 1)$$

$$x(1, 1, 0, 1) = x_4(1, 0, 1, 1)$$

$$x(1, 1, 1, 0) = x_4(0, 1, 1, 1)$$

$$x(1, 1, 1, 1) = x_4(1, 1, 1, 1)$$

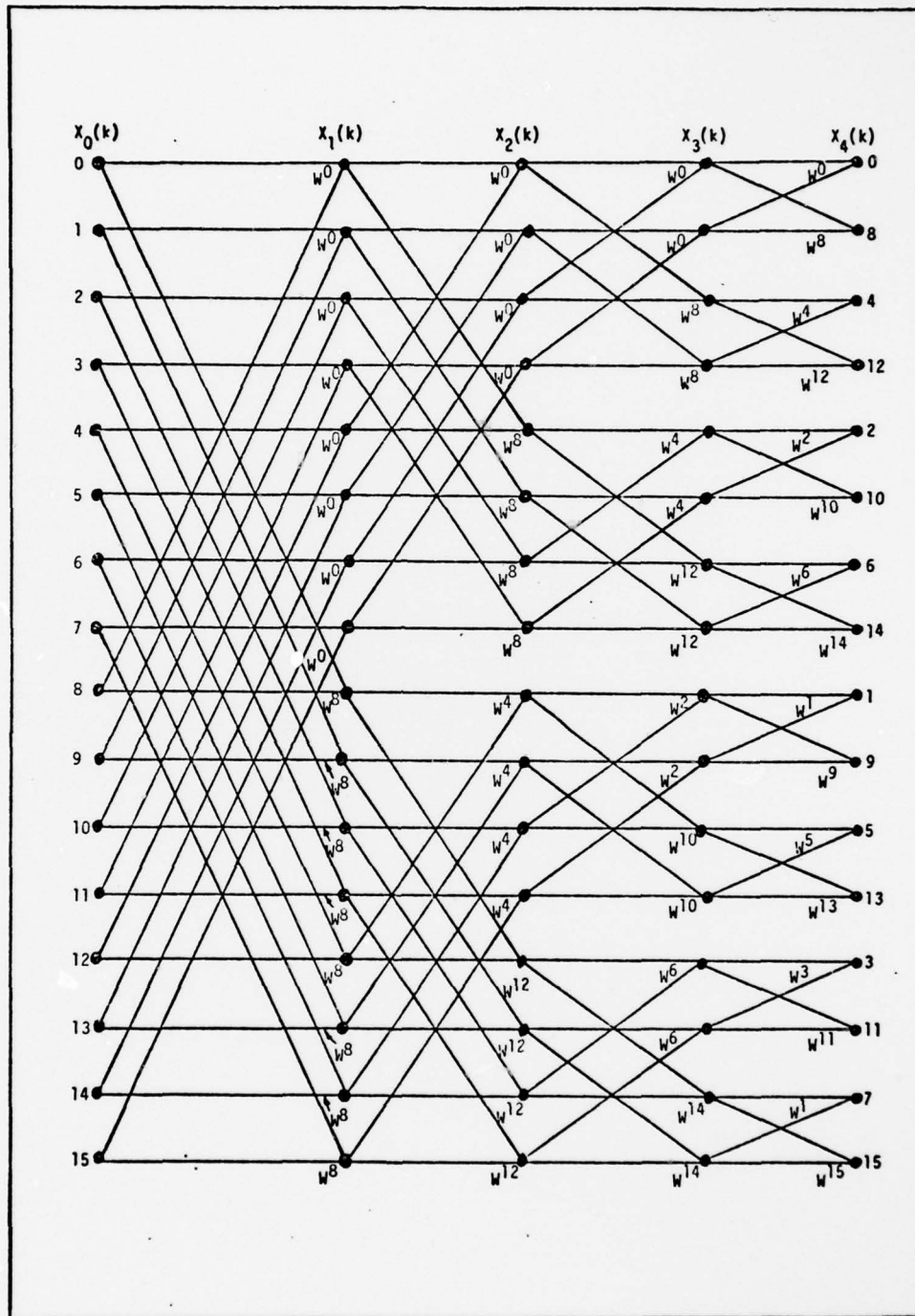


Figure C-1. Signal Flow Graph Without Bit-Inversion of Data

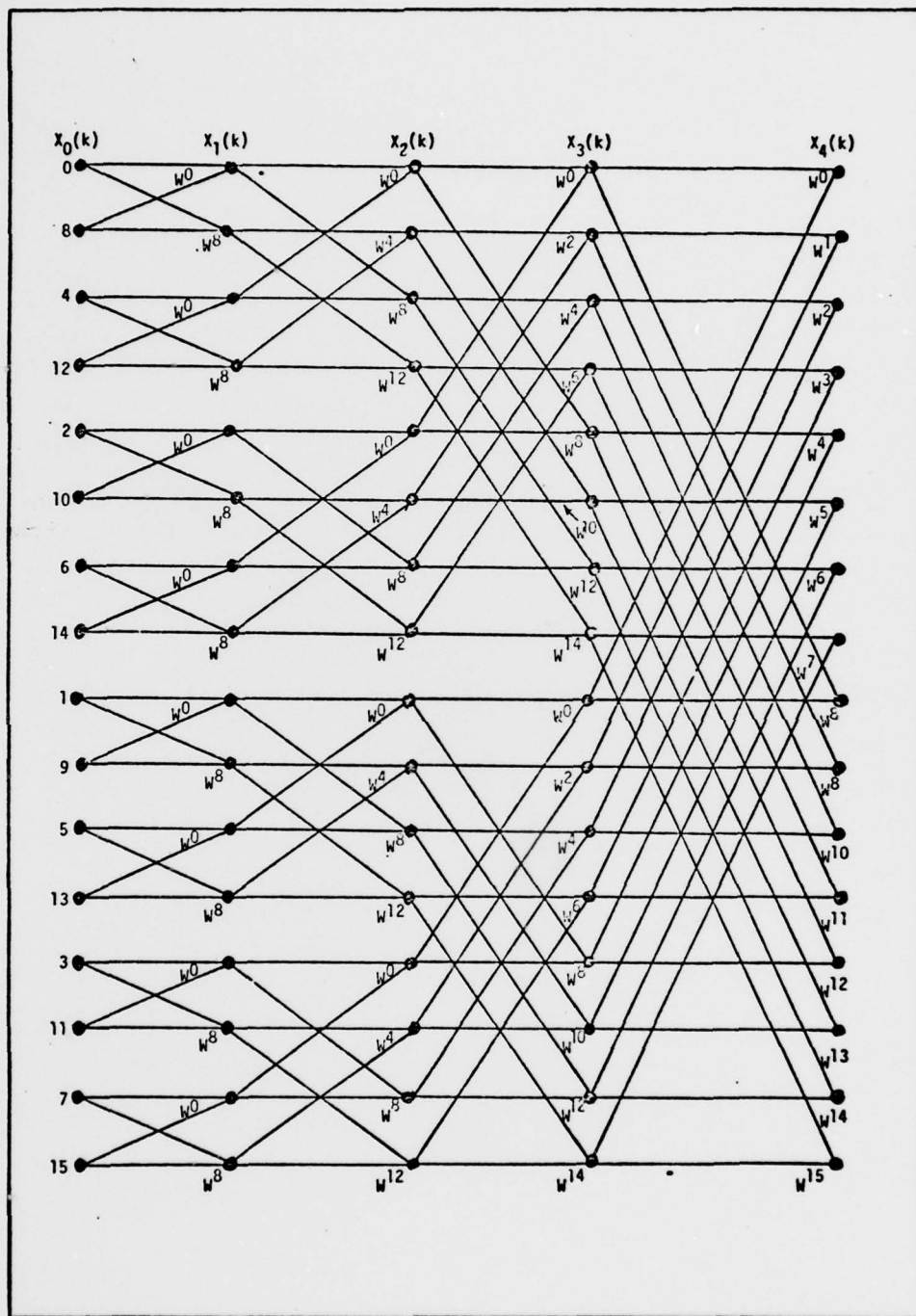


Figure C-2. Signal Flow Graph With Bit-Inversion of Data

APPENDIX D

CALCULATION OF DFT OF REAL DATA

## APPENDIX D

### Calculation of DFT of Real Data

Consider a real  $2N$  point sequence  $x(k)$ ; the DFT of this sequence can be carried out using Eq (D-1):

$$X(n) = \sum_{k=0}^{2N-1} x(k) W_{2N}^{nk} \quad n = 0, 1, \dots, 2N-1 \quad (D-1)$$

$$W_{2N}^{nk} = \exp \frac{-2\pi n i k}{2N}, \dots, n = 0, 1 \dots 2N-1 \quad (D-2)$$

Equation (D-1) can be rewritten in the following form:

$$X(n) = \sum_{r=0}^{N-1} x(2r) W_{2N}^{2rn} + \sum_{r=0}^{N-1} x(2r+1) W_{2N}^{(2r+1)n} \quad (D-3)$$

$$X(n) = \sum_{r=0}^{N-1} x(2r) W_N^{Rn} + W_N^n \left[ \sum_{r=0}^{N-1} x(2r+1) W_N^{rn} \right] \quad (D-4)$$

Equation (D-4) has resulted by splitting  $x(k)$  in (D-1) into even and odd sequences as shown in Equations (D-5) and D-6) respectively.

$$y(r) = x(2r) \quad \text{even sequence} \quad (D-5)$$

$$z(r) = x(2r+1) \quad \text{odd sequence} \quad (D-6)$$

A complex sequence  $v(r)$  is now defined such that the even points of  $x(k)$  [Equation (D-5)] becomes the real part of  $v(r)$  and the odd point of  $x(k)$  [Equation (D-6)] becomes the imaginary part of  $v(r)$ .

$$v(r) \equiv y(r) + jz(r) \quad (D-7)$$

An N-point DFT of  $V(r)$  yields  $V(n)$ , which is given in Eq (D-8):

$$V(n) = \sum_{r=0}^{N-1} v(r)W_N^{rn} \quad (D-8)$$

or

$$V(n) = \sum_{r=0}^{N-1} y(r)W_N^{rn} + j \sum_{r=0}^{N-1} z(r)W_N^{rn} \quad (D-9)$$

or

$$V(n) = Y(n) + jZ(n) \quad (D-10)$$

The DFT of real and imaginary parts can be written in terms of  $V(n)$  as (Ref 5:319).

$$Y(n) = 1/2[V(n) + V^*(N-n)] \quad (D-11)$$

and

$$Z(n) = 1/2j[V(n) - V^*(N-n)] \quad (D-12)$$

here

$$V^*(N-n) = Y(N-n) - jZ(N-n) \quad (D-13)$$

The result obtained in Eqs (D-11) and (D-12) can be used to rewrite Eq (D-4) as:

$$X(n) = Y(n) + W_{2N}^n Z(n) \quad (D-14)$$

or

$$X(n) = 1/2[V(n) + V^*(N-n)] + 1/2j[V(n) - V^*(N-n)] \quad (D-15)$$

This equation is true for  $0 \leq n \leq N$ . [Note that  $V(N) = V(0)$  ].

(Ref 5:320-321)

APPENDIX E

LSI-11 ASSEMBLY LANGUAGE CODE FOR  
FFT SUBROUTINE (FOURIE) AND  
EXECUTION TIME COMPUTATION

APPENDIX E

LSI-11 Assembly Language Code for FFT Subroutine  
 (Fourie) and Execution Time Computation  
 (Ref 6:288-293)

FOURIE --FFT SUBROUTINE RT-11 MACRO VMD2-12 30-OCT-78 23:42:04 PAGE 1

```

1 000000'
2
3 .TITLE FOURIE --FFT SUBROUTINE
4 .SBTTL INTRODUCTION --INTKO: TO FOURIE AND ITS USE.
5 .CSECT
6
7
8
9
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23 000000 000040
24
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26 000002 000020
27
28 000004 000004
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32 000006 000000
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41
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43 000010 000000
44 000112' 000000
45 000112 000000
46 000214' 000214'
47
48
49

```

1. DESCRIPTION:  
 -----  
 THIS SUBROUTINE CAN BE USED TO TRANSFORM 32 REAL DATA POINTS. IT ASSUMES THE DATA EQUIVALENT TO 16 COMPLEX POINTS (THE ODD REAL POINTS ARE CONSIDERED REAL PARTS AND THE EVEN PARTS ARE CONSIDERED COMPLEX). AFTER THE TRANSFORM A POST PROCESSING OF DATA GIVES 16 FOURIER TRANSFORMED FREQUENCY COMPONENTS (THE OTH -ER 16 COMPONENTS ARE THE COMPLEX CONJUGATES OF THESE VALUES).

2. DEFINITION OF PARAMETERS USED  
 -----  
 THE FOLLOWING ARE THE DIFFERENT PARAMETERS USED IN THE FFT SUBROUTINE "FOURIE":  
 N: 40 IN IS THE # OF REAL DATA POINTS IN OCTAL FORMAT  
 NZ: 20 INZ IS THE ASSUMED # OF COMPLEX POINTS IN OCTAL.  
 B: 4 B IS THE EXPONENT OF 2 WHICH SATISFIES THE RELATIONSHIP N2=2\*\*B.  
 DATA: 0 DATA IS POINTER THAT WILL CONTAIN THE STARTING ADDRESS OF INPUT REAL POINTS THAT ARE TO BE TRANSFORMED.

3. DATA STORAGE AREA ALLOCATION:  
 -----  
 TWO DATA BUFFER AREAS ARE USED SO THAT DATA CAN BE ACQUIRED AND STORED IN THE 2ND BUFFER WHILE PROCESSING IS BEING DONE ON THE 1ST BUFFER.  
 ISTR1: 0  
 .+100  
 ISTR2: 0  
 .+100  
 NOTE: EACH BUFFER IS 32 WORDS LONG.

1  
2  
3  
4  
5  
6  
7  
8  
9  
10  
11  
12  
13  
14  
15  
16  
17  
18  
19  
20  
21  
22  
23

4. GLOBAL SYMBOLS:  
-----  
THE FOLLOWING SYMBOLS OF "FOURIE" ARE USED BY OTHER  
SUBROUTINES THAT ARE CALLED BY THE EXECUTIVE.  
.GLOBL FOURIE,DATA,ISTR1,ISTR2,N,NZ

5. REGISTER DEFINITION:  
-----  
.MCALL .RECDEF  
.RECDEF

15 000214  
16  
17  
18  
19  
20  
21  
22  
23

CONTINUED ON NEXT PAGE....



```

1
2
3
4 000306 005200
5 000310 010001
6
7 000312 077532
8
9
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11
12
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16
17
18
19
20
21
22
23
24
25
26 000314 016700 177466
27
28 000320 016701 177454
29 000324 004767 001126
30
31 000330 016700 177452
32 000334 010001
33 000336 022121
34 000340 012705 000010
35
36 000344 011002
37 000346 061120
38 000350 161102
39 000352 010221
40 000354 011002
41 000356 061120
42 000360 161102
43 000362 010221
44 000364 022021
45 000366 022021
46 000370 077513
47 000372 000412
48
49
50
51
52
  
```

```

      INC      R0
      MOV      R0,R1
      SOB      R5,BIV1D
      NOTE:
      -----
  
```

AT THIS POINT THE BIT INVERSION IS COMPLETE  
 AND DATA IS READY TO BE TRANSFORMED.

2. FAST FOURIER TRANSFORMATION (FIRST PASS):  
 -----  
 DURING THIS PASS THERE ARE NO MULTIPLICATIONS AND  
 THE FOLLOWING EQUATIONS ARE IMPLEMENTED:

$$\begin{aligned}
 R(M) &= R(M)+R(N) \\
 I(M) &= I(M)+I(N) \\
 R(N) &= R(M)-R(N) \\
 I(N) &= I(M)-I(N)
 \end{aligned}$$

```

FPASS:  MOV      DATA,R0
        MOV      N,R1
        JSR      PC,SCALE
        MOV      DATA,R0
        MOV      R0,R1
        CMP      (R1)+,(R1)+
        MOV      #10,R5
        MOV      (R0),R2
        ADD      (R1),(R0)+
        SUB      (R1),R2
        MOV      R2,(R1)+
        MOV      (R0),R2
        ADD      (R1),(R0)+
        SUB      (R1),R2
        MOV      R2,(R1)+
        CMP      (R0)+,(R1)+
        CMP      (R0)+,(R1)+
        SOB      R5,PASS1
        BR       SPASS
  
```

```

;LOAD STARTING ADDRESS
;OF DATA FOR SCALING.
;LOAD # OF POINTS.
;ALWAYS SCALE DATA TO
;PREVENT OVERFLOW.
;FIRST POINT R(M)
;2ND POINT R(N)
;COUNTER=NZ/2
;SAVE R(M)
;R(M)=-R(M)+R(N)
;R(N)=-R(M)-R(N)
;I(M)=-I(M)+I(N)
;I(N)=-I(M)-I(N)
;GO TO NEXT PAIR
;GO BACK NZ/2 TIMES
  
```

```

      NOTE:
      -----
      FIRST PASS IS OVER NOW GO TO SECOND PASS.
  
```

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 10 000374 000000  
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 18  
 19 000376 000000  
 20  
 21 000400 000000  
 22  
 23 000402 000000  
 24  
 25 000404 000000  
 26  
 27  
 28  
 29 000406 000000  
 30  
 31 000410 000000  
 32  
 33  
 34 000412 000000  
 35  
 36  
 37 000414 000000  
 38  
 39  
 40  
 41 000416 000000  
 42  
 43  
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 46  
 47

3. FAST FOURIER TRANSFORM (2ND PASS AND ONWARD):  
 -----  
 THE 2ND,3RD AND 4TH PASS USE THE SAME LOOP,  
 HOWEVER IN EACH PASS THE PARAMETERS DEFINED  
 BELOW TAKE DIFFERENT VALUES:

DELTA: 0    I SIZE OF INCREMENT FOR EXPO  
 -MENT OF "M" WHICH SATISF  
 -IES THE FOLLOWING RELATIO  
 -NSHIP:  
 $M = EXP(IZ * PI * K / N)$   
 $K = 0, 1, 2, \dots, N-1$   
 $N = \# \text{ OF POINTS}$

CELNUM: 0    I# OF CELLS IN EACH PASS.  
 PAIRN: 0    I# OF PAIRS IN EACH PASS.  
 CELDIS: 0    I DISTANCE BETWEEN PAIR OF POINTS  
 I IN A CELL.

POINT: 0    I TEMPORARY POINTER TO STARTING  
 ADDRESS OF DATA TO TRANSFORMED.

CELCNT: 0    I CELL COUNTER TO KEEP TRACK  
 OF # OF CELLS IN EACH PASS.

PARCNT: 0    I COUNTER TO KEEP TRACK OF  
 # OF PAIRS IN EACH CELL.

SINE: 0    I TEMPORARY STORAGE TO HOLD  
 VALUE FOR SIN(Y).

COS: 0    I TEMPORARY STORAGE TO HOLD  
 VALUE FOR COS(Y).

YN: 0    I THIS INDICATOR HOLDS THE  
 VALUE OF "Y" BEFORE SINE  
 COSINE LOOK UP TABLES ARE  
 SEARCHED.

CONTINUED ON NEXT PAGE.....

```

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14
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16
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19
20
21
22
23
24 000420 012767 000004 177746
25 000426 012767 000004 177742
26 000434 012767 000002 177736
27 000442 012767 000010 177732
28
29
30
31
32
33 000450 016700 177332
34
35
36 000454 016701 177320
37 000460 004767 000772
38
39 000464 016767 177706 177714
40 000472 016767 177310 177704
41 000500 000404
42
43
44
45
46
    
```

4. CODE FOR 2ND PASS AND ONWARD:  
 -----  
 THE EQUATIONS THAT ARE IMPLEMENTED IN 2ND,3RD  
 AND 4TH PASS ARE GIVEN BELOW:

$$\begin{aligned}
 R1' &= R1 + R2 * \cos(Y) + I1 * \sin(Y) \\
 I1' &= I1 - R2 * \sin(Y) + I2 * \cos(Y) \\
 R2' &= R1 - R2 * \cos(Y) - I2 * \sin(Y) \\
 I2' &= I1 + R2 * \sin(Y) - I2 * \cos(Y)
 \end{aligned}$$

WHERE Y=2\*PI\*K/N, N=# OF COMPLEX POINTS, AND  
 K=0,1,2,...,N/2-1

NOTE:  
 -----

3RD AND 4TH PASS START FROM LABEL "NENPASS"  
 WHEREAS 2ND PASS STARTS AT LABEL "SPASS".

SPASS: MOV #4, DELTAY  
 MOV #4, CELNUM  
 MOV #2, PAIRNM  
 MOV #10, CELDIS

NENPASS: MOV DATA, RD  
 N,R1  
 PC, SCALE

MOV CELNUM, CELCNT  
 MOV DATA, POINT  
 BR NENCEL

ICET STARTING ADDRESS  
 AND # OF POINTS BEFORE  
 CALLING SUBI\*SCALE\*

ISET UP CELL COUNTER.

CONTINUED ON NEXT PAGE.....

```

1
2
3
4
5
6
7
8 000502 000000
9 000504 000000
10 000506 000000
11 000510 000000
12
13
14
15 000512 016767 177662
16 000520 005067 177672
17 000524 016700 177654
18 000530 010001
19 000532 066701 177644
20 000536 016702 177654
21 000542 004767 001060
22
23 000546 010267 177640
24 000552 062702 040000
25 000556 004767 001044
26 000562 010267 177626
27 000566 011102
28 000570 070267 177620
29 000574 073227 000001
30
31 000600 010267 177704
32 000604 011102
33 000606 070267 177600
34 000612 073227 000001
35 000616 010267 177662
36 000622 010104
37 000624 005724
38 000676 011402
39 000630 070267 177556
40 000634 073227 000001
41 000640 010267 177636
42 000644 011402
43 000646 070267 177542
44 000652 073227 000001
45 000656 010267 177624
46
47
48
49
50

```

NOTE: EACH NEW CELL WITHIN A PASS STARTS HERE.  
 ---- SOME TEMPORARY STORAGE VARIABLES ARE DEFINED BELOW.

```

I2*SINY: 0
R2*SINY: 0
I2*ICOS(Y): 0
R2*ICOS(Y): 0
I2*ICOS(Y): 0
R2*ICOS(Y): 0

```

```

NENCL: MOV PAIFNM,PARCNT
CLR YN
NENC10: MOV POINT,RO
MOV RO,R1
ADD CELDIS,R1
MOV YN,R2
JSR PC,SNL0K2

FT10: MOV R2,SINE
ADD #40000,R2
JSR PC,SNL0K2
MOV R2,COS
MOV (R1),R2
MUL COS,R2
ASHC #1,R2

MOV R2,R2COSY
MOV (R1),R2
MUL SINE,R2
ASHC #1,R2
MOV R2,R2SINY
MOV (R1),R2
MUL SINE,R2
ASHC #1,R2
MOV R2,I2SINY
MOV (R1),R2
MUL COS,R2
ASHC #1,R2
MOV R2,I2COSY

```

CONTINUED ON NEXT PAGE.....



FOURIE --FFT SUBROUTINE RT-11 MACRO VMD2-12 30-OCT-76 23142104 PAGE 9  
 FOURIE --DIFFERENT PARTS AND THEIR CODE.

```

1  |
2  |
3  |
4  |
5  |
6  |
7  |
8  |
9  |
10 |
11 |
12 | 000770 006267 177402
13 | 000774 001420
14 |
15 | 000776 006367 177376
16 |
17 | 001002 006367 177374
18 | 001006 006267 177362
19 | 001012 000411
20 |
21 |
22 |
23 |
24 |
25 |
26 |
  
```

NOTE1  
 -----

ONE PASS IS DONE NOW GO ON TO NEXT PASS.  
 SET NEW VALUES OF PARAMETERS DEFINED  
 IN PARA 3.

```

ASR  CELNUM
BEQ  POST
ASL  PAIRNM
ASL  CELDIS
ASR  DELTAY
BR   POST

!LESS CELLS THIS TIME.
!ALL DONE WHEN VARIABLE
  CELNUM=0.
!THIS TIME THERE ARE
  MAKE PAIRS IN EACH CELL.
!PAIRS ARE FURTHER APART.
!LESS INC OF YN PER CELL.
  
```

END OF FFT ALGORITHM.

.SBTTL POST PROCESSING ALGORITHM

1. INTRODUCTION TO POST PROCESSING ALGORITHM:

-----  
 POST PROCESSING OF THE TRANSFORMED ARRAY IS DONE  
 BECAUSE THE INPUT REAL DATA WAS ASSUMED TO BE COMPLEX  
 AND THE OUTPUT OF THE FFT IS NOT IN CORRECT FORM.  
 THE EQUATIONS THAT ARE IMPLEMENTED IN THIS PASS  
 ARE GIVEN BELOW:

$$\begin{aligned} AR(M) &= RP + IP * COS(T) - RM * SIN(T) \\ AI(M) &= IM - IP * SIN(T) - RM * COS(T) \\ AR(N-M) &= RP - IP * COS(T) + RM * SIN(T) \\ AI(N-M) &= IM - IP * SIN(T) - RM * COS(T) \end{aligned}$$

WHERE T=PI\*M/N, N=16 AND M=0,1,2,...,N/2-1, AND  
 RP=I(M)+I(N-M), RM=R(M)+R(N-M), IM=I(M)-I(N-M) AND  
 RM=R(M)-R(N-M).

SOME TEMPORARY VARIABLES THAT ARE USED IN  
 THIS PASS ARE DEFINED BELOW:

RP: 0  
 IP: 0  
 RM: 0  
 IM: 0  
 ARM: 0  
 ARNM=AR(N-M)  
 ARNI=0  
 AI=AI(M)  
 AIIN=AI(N-M)  
 INX: 0  
 IINDEX TO GET SIN,COS VALUES.

```

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16
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18
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20
21
22
23
24
25 001014 000000
26 001016 000000
27 001020 000000
28 001022 000000
29 001024 000000
30 001026 000000
31 001030 000000
32 001032 000000
33 001034 000000
34
35
36
37
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39
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41
42
43
44
    
```

CONTINUED ON NEXT PAGE.....

```

1
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3
4
5
6
7
8
9
10
11
12
13
14 001036 016700 176744
15
16
17 001042 016701 176732
18 001046 004767 000404
19 001052 016700 176730
20 001056 012701 000010
21 001062 012767 000001
22 001070 012704 000074
23
24 001074 066704 176706
25
26
27
28
29
30
31
32
33
34 001100 016002 000002
35
36 001104 060210
37 001106 006320
38
39 001110 005020
40
41
42
43
44
45
46
47
48
49
50 001112 016702 177716
51 001116 004767 000600
52
53
54
55
    
```

2. CODE FOR POST PROCESSING :

NOTE!

SET UP ADDRESS POINTERS AND OTHER RELATED COUNTERS.

POST1 MOV DATA,R0  
 ;SET ADDRESS POINTER  
 AND # OF POINTS BEFORE  
 SCALING DATA ARRAY.

MOV N,R1  
 JSR PC,SCALE  
 MOV DATA,R0  
 MOV #10,R1  
 MOV #1,INX  
 MOV #74,R4

;SET COUNTER TO DECIMAL 8,  
 INCREMENT FOR SINE POINTER,  
 OFFSET FOR ADDRESS OF  
 POINT.

ADD DATA,R4

NOTE!

THE POST PROCESSING OF FIRST POINT IS DONE SEPERATELY BECAUSE WE DO NOT PHYSICALLY GENERATE THE (N+1)ST POINT.

POST1 MOV 2(R0),R2  
 ADD R2,(R0)  
 ASL (R0)+  
 CLR (R0)+

IGET IMAG PART OF 1ST POINT.  
 IADD IT TO REAL PART.  
 IDOUBLE THIS RESULT TO MAINTAIN SCALE.  
 ICLEAR THE LOCATION WHERE THE IMAG WAS STORED.

NOTE!

NOW THE REST OF THE POINTS ARE PROCESSED,SO GET RM,IM,RP AND IP FOR ALL EXCEPT THE FIRST POINT.

POST5 MOV INX,R2  
 JSR PC,SNL0K1

IGET THE VALUE OF SINE.

CONTINUED ON NEXT PAGE.....



POST PROCESSING ALGORITHM

```

1 |
2 |
3 |
4 |
5 | 001402 016710 177416
6 | 001406 016714 177414
7 | 001412 016760 177412 000002
8 | 001420 016764 177406 000002
9 |
10 |
11 |
12 |
13 |
14 |
15 |
16 |
17 | 001426 005267 177402
18 | 001432 022020
19 | 001434 024444
20 | 001436 005301
21 | 001440 001402
22 | 001442 000167 177444
23 | 001446 005367 000002
24 | 001452 000207
25 |
26 |
27 |
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29 |
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33 |
34 |
35 |

```

```

POSTS: MOV ARM, (RD)
MOV ARNM, (R4)
MOV AIM, 2(RD)
MOV AINM, 2(R4)

```

NOTE:

-----  
ONE PAIR OF POINTS HAS BEEN DONE NOW  
GO ON TO NEXT PAIR OF POINTS.

```

INC INX
CMP (RD)+, (RD)+
CMP -(R4), -(R4)
DEC R1
BEQ POEXT
JMP POSTS
RTS SFACR PC

INX INX
INCR RD BY 2
INCR R4 BY 2
R1<-R1-1
IF R1=0 EXIT

```

END OF POST PROCESSING ALGORITHM.

.SBTTL SUBROUTINE SCALE

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1. INTRODUCTION TO SUBROUTINE "SCALE":

-----  
THIS SUBROUTINE CHECKS BITS 13-14 AND SHIFTS THE WHOLE ARRAY RIGHT BY TWO BITS IF BIT-14 IS SET AND SHIFTS THE ARRAY ONE BIT RIGHT IF BIT-13 IS SET. BEFORE CALLING THIS SUBROUTINE IT SHOULD BE ENSURED THAT THE STARTING ADDRESS OF THE ARRAY IS LOADED IN R0 AND THE LENGTH OF THE ARRAY IS LOADED IN R1. THE SUBROUTINE IS CALLED BY GIVING THE COMMAND "JSR PC,SCALE".  
THE VARIABLE "SFACTR" WHICH IS DEFINED BELOW KEEPS TRACK OF THE # OF TIMES THE ARRAY HAS BEEN SHIFTED RIGHT BY ONE BIT.

SFACTR: 0 SFACTR IS DEFINED HERE.

2. CODE FOR SUBROUTINE SCALE:

```
SCALE:  MOV  R0,R5          ;GET STARTING ADDRESS
        MOV  R1,R3          ;AND STORE IN REG: R5.
        CLR  R4             ;GET ARRAY LENGTH AND
                          ;STORE IN REG: R3.
        MOV  (R0)+,R2      ;CLEAR R4 WHICH IS A FLAG
                          ;TO KEEP TRACK WHETHER BIT
                          ;-14 OF BIT-13 IS SET.
        BPL  SCAL2         ;GET EACH POINT.
                          ;SKIP NEXT INSTRUCTION IF
                          ;IF VALUE IS POSITIVE.
        NEG  R2            ;GET ABS VALUE.
        BNE  RDUCE2       ;IS BIT-14 SET?
        BEQ  SCAL3        ;YES,SCALE TWICE.
        MOV  #20000,R2    ;IS BIT-13 SET?
        BEQ  SCAL3        ;NO,GET ANOTHER POINT.
        MOV  #1,R4        ;YES,SET FLAG AND CHECK
                          ;NEXT POINT.
```

CONTINUED ON NEXT PAGE.....

FOURTE --FFT SUBROUTINE RT-11 MACRO VMD2-12 30-OCT-78 23142104 PAGE 15  
 SUBROUTINE SCALE

```

1 |
2 |
3 |
4 |
5 | 001512 077114
6 | 001514 005704
7 | 001516 001004
8 | 001520 000207
9 |
10 |
11 |
12 |
13 |
14 |
15 |
16 |
17 | 001522 005004
18 | 001524 005267
19 | 001530 005267
20 | 001534 005704
21 | 001536 001004
22 | 001540 006215
23 |
24 | 001542 006225
25 | 001544 077303
26 | 001546 000207
27 | 001550 006225
28 | 001552 077302
29 | 001554 000207
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```

SCAL3: SOB R1, SCAL1  
 TST R4  
 BNE RDUCE1  
 RTS PC

NOTE:  
 -----  
 SHIFT ENTIRE ARRAY EITHER TWO BITS OR ONE  
 BIT TO THE RIGHT.

RDUCE2: CLR R4  
 INC SFACTR  
 RDUCE1: INC SFACTR  
 TST R4  
 BNE RDUCE4  
 RDUCE3: ASR (R5)  
 ASR (R5)+  
 SOB R3, RDUCE3  
 RTS PC

RDUCE4: ASR (R5)+  
 SOB R3, RDUCE4  
 RTS PC

SHIFT EACH POINT TWO  
 BITS TO THE RIGHT.

ALL POINTS CHECKED?  
 IF TEST IF SCALING IS REQUIRED,  
 YES IT IS REQUIRED.

ALL POINTS SCALED?  
 IF YES

END OF SCALING SUBROUTINE.

FOURIE ---FFT SUBROUTINE RT-11 MACRO VM02-12 30-OCT-78 23:42:104 PAGE 16  
SUBROUTINE "SWAP"

.SBTTL SUBROUTINE "SWAP"

1. SUBROUTINE "SWAP" (INTRODUCTION AND CODE):

-----  
THIS SUBROUTINE IS USED TO SWAP THE CONTENTS OF  
MEMORY LOCATIONS POINTED TO BY REGISTERS R1 AND R2.  
BEFORE CALLING THIS SUBROUTINE IT SHOULD BE ENSURED  
THAT THE ADDRESS OF MEMORY LOCATIONS THAT ARE TO SWAP  
--PED ARE LOADED IN R1 AND R2.

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20  
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23  
24  
25  
26  
27  
28  
29  
30  
31
```

001556 011103  
001560 011211  
001562 010312  
001564 000207

SWAP:

MOV (R1),R3  
MOV (R2),(R1)  
MOV R3,(R2)  
RTS PC

END OF SUBROUTINE SWAP.

.SBTTL SUBROUTINE "SNLOKZ"

1. INTRODUCTION TO "SNLOKZ":

-----  
 THIS SUBROUTINE IS CALLED FROM FOURIE WHEN THE ALGORITHM IS EXECUTING 2ND PASS AND ONWARD. THIS SUB-ROUTINE IS USED TO RETURN VALUES FOR SINE AND COSINE BY SEARCHING TABLE "SNTAB2". THE ONLY DIFFERENCE BETWEEN THIS SUB: AND "SNLOK1" IS THE SINE-COSINE LOOK UP TABLE. TABLE "SNTAB2" HAS SIXTEEN VALUES STORED, THE FIRST EIGHT VALUES ARE FOR SIN(PI\*M/8) AND THE NEXT EIGHT ARE FOR COS(PI\*M/8). IN THIS CASE M=0,1,2,...,N/2-1. THE SINE-COSINE TABLE IS GIVEN BELOW:

SNTAB2: 000000,030373,055202,073101  
 077777,073101,055202,030373  
 077777,073101,055202,030373  
 000000,147404,122575,104676

2. CODE FOR "SNLOKZ":

-----  
 SNLOKZ: BIT 440000,R2  
 BNE CSLOKZ  
 BR LOKZ  
 CSLOKZ: BIC #17770,R2  
 LOKZ: ADD #10,R2  
 ASL R2  
 MOV SNTAB2(R2),R2  
 R16 PC

```

1 2
2 3
3 4
4 5
5 6
6 7
7 8
8 9
9 10
10 11
11 12
12 13
13 14
14 15
15 16
16 17
17 18
18 19
19 20
20 21
21 22
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23 24
24 25
25 26
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31 32
32 33
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34 35
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36 37
37 38
38 39
39 40
40 41
41 42
42 43
43 44
44 45
45 46
46 47
47 48
48 49
49 50
  
```

END OF SUBROUTINE "SNLOKZ"

.SBTTL SUBROUTINE "SNLOK1"

1. INTRODUCTION TO SUBROUTINE "SNLOK1"

-----  
 SUBROUTINE "SNLOK1" IS USED TO LOOK UP THE VALUE  
 OF SINE AND COSINE TO THE CALLING PROGRAM BY SEARCH-  
 ING A SINE COSINE TABLE "SNTAB1". WHEN CALLING PROGRAM  
 LOADS R2 WITH A 4-BIT BINARY # CORRESPONDING TO THE DES-  
 IRED VALUE OF ANGLE. SINCE THIS SUBROUTINE IS USED BY  
 POST PROCESSING ALGORITHM, TABLE "SNTAB1" HAS SINE AND  
 COSINE VALUES STORED FOR  $T=PI*M/16$  WHERE  $M=0,1,...,N/2$ .  
 THE SINE COSINE TABLE "SNTAB1" IS GIVEN BELOW, THE FIRST  
 NINE VALUES ARE FOR  $SIN(PI*M/16)$  AND THE NEXT NINE  
 VALUES ARE FOR  $COS(PI*M/16)$ !

SNTAB1: 000000,014370,030373  
 043434,055202,065155  
 073101,076612,077777  
 077777,076612,073101  
 065155,055202,043434  
 030373,014370,000000

2. CODE FOR SUBROUTINE "SNLOK1":  
 -----

```

SNLOK1: BIT      #40000,R2          ;IS BIT-14 SET?
        BNE      CSLOK1           ;YES, GET VALUE OF COSINE
        BR       LOK1            ;CLEAR BITS 3-15
        BIC      #17760,R2       ;ADD OFFSET FOR COSINE VALUES.
        ADD      #11,R2          ;SHIFT LEFT FOR BYTE ADDRESSING.
        ASL      R2              ;GET VALUE FROM SNTAB1.
        MOV      SNTAB1(R2),R2
        RTS      PC
    
```

.END

```

1
2
3
4
5
6
7
8
9
10
11
12
13
14
15
16
17
18
19
20
21 001656 000000 014370 030373
22 001664 043434 055202 065155
23 001672 073101 076612 077777
24
25 001700 077777 076612 073101
26 001706 065155 055202 043434
27 001714 030373 014370 000000
28
29
30
31
32
33
34
35 001722 032702 040000
36 001726 001001
37 001730 000404
38 001732 042702 17760
39 001736 062702 000011
40 001742 006302
41 001744 016202 001656
42 001750 000207
43
44
45
46 000001
    
```

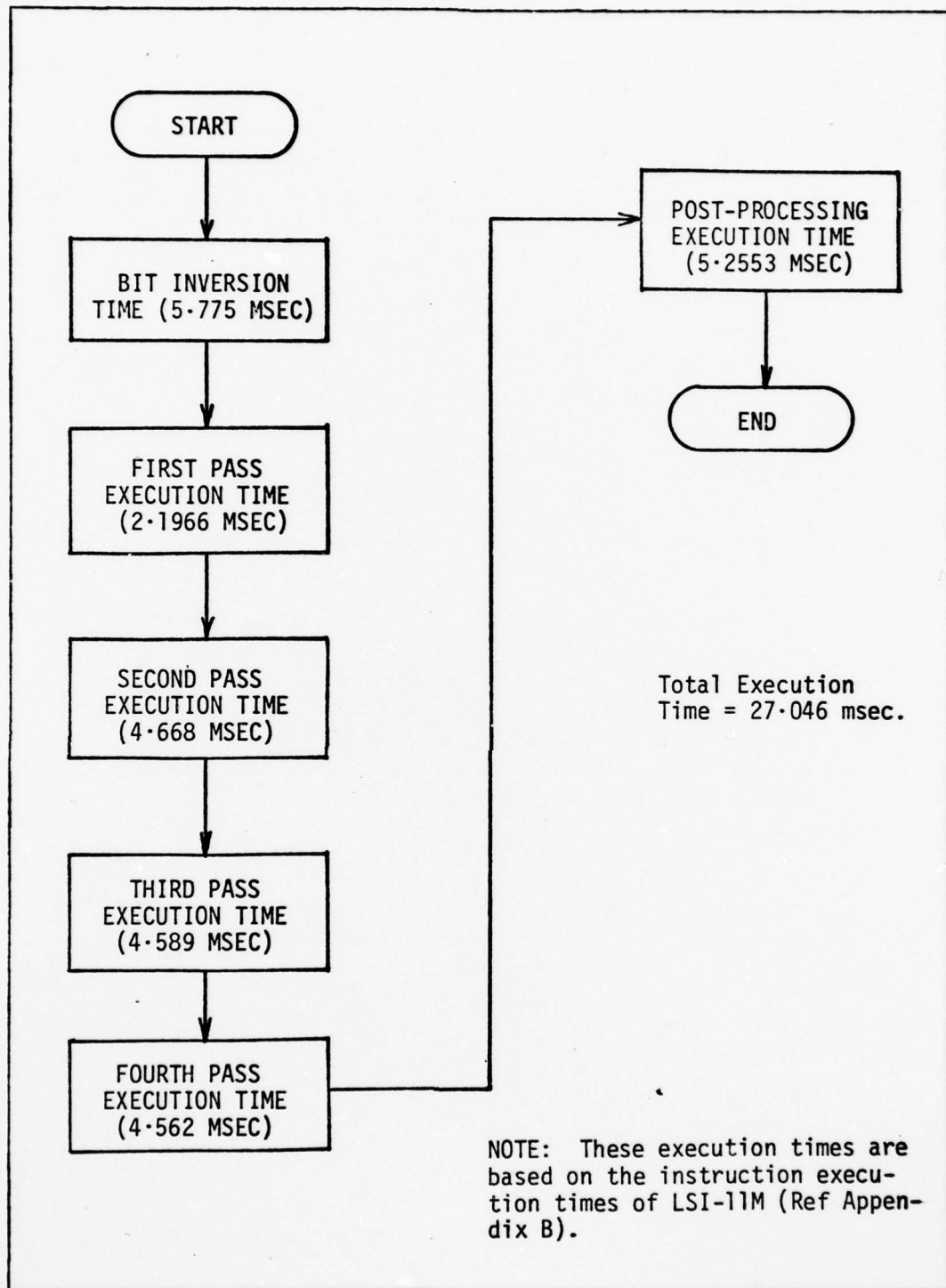


Figure E-1. Execution Time for Subroutine FOURIE

APPENDIX F

LSI-11 ASSEMBLY LANGUAGE CODE FOR  
POWER SPECTRUM SUBROUTINE (PRSPEC)  
AND EXECUTION TIME COMPUTATION

APPENDIX F

LSI-11 Assembly Language Code for Power Spectrum Subroutine (PRSPEC) and Execution Time Computation

PRSPEC---POWER SPECTRUM SUBROUT RT-11 MACRO VM02-12 01:28:18 PAGE 1

1 .TITLE PRSPEC---POWER SPECTRUM SUBROUTINE.  
2 .SECT INTRODUCTION TO "PRSPEC" SUBROUTINE.

3  
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1. INTRODUCTION TO "PRSPEC" SUBROUTINE:  
-----  
THIS SUBROUTINE CALCULATES THE POWER SPECTRUM OF AN INPUT ARRAY BY SQUARING AND ADDING THE REAL AND IMAGINARY PARTS OF EACH DATA POINT. THE INPUT ARRAY IS ASSUMED TO HAVE REAL AND IMAGINARY PARTS OF EACH DATA POINTS IN CONTIGUOUS MEMORY LOCATIONS. THE INPUT ARRAY IS ASSUMED TO HAVE 32 DATA POINTS, AND IT IS ASSUMED BY THIS SUBROUTINE THAT THE STARTIN -C ADDRESS OF THE INPUT ARRAY IS STORED IN LOCATION "DATA".  
THE POWER SPECTRUM OF ONLY THE FIRST 11 DATA POINTS IS CALCULATED, BECAUSE THE FREQUENCY OF INTEREST LIES BETWEEN 0-HZ AND 1000-HZ.  
THE GLOBAL VARIABLES AND MACRO CALLS ARE DEFINED BELOW:

.GLOBL SCALE,DATA,SPECTRA,SSTORE  
.GLOBL PRSPEC,SFACTR,N,CUNSF1

.MCALL .RECDEF  
.RECDEF

.CSECT

000000

000000'

.SBTTL DATA STRUCTURE AND VARIABLE DEFINITION.

```

1
2
3
4
5
6
7
8
9
10
11 000000 000000
12
13
14
15 000002 000000
16
17
18
19
20
21
22
23
24 000004 000006
25
26 000006 000024
27
28
29 000104
30
31
  
```

2-DATA STRUCTURE AND VARIABLES USED:

-----  
 THE VARIABLES USED IN THIS SUBROUTINE ARE DEFINED  
 BELOW:

```

CUMSF1: 0      ! CUMULATIVE SCALE FACTOR
              ! WHICH KEEPS TRACK OF THE
              ! NUMBER OF TIMES THE OUTPUT
              ! ARRAY HAS BEEN SHIFTED ONE
              ! PLACE TO THE RIGHT.
              ! TEMPORARY VARIABLE THAT
              ! STORES THE SQUARE OF REAL
              ! PART OF A DATA POINT.

RLSOR:  0
  
```

THE DATA STRUCTURE FOR THE OUTPUT ARRAY IS GIVEN  
 BELOW. THE OUTPUT ARRAY IS CALLED "SFCTRA"  
 AND IS 11 WORDS LONG:

```

SFCTRA: .WORD SSTORE      ! STARTING ADDRESS OF
                          ! THE ARRAY "SFCTRA".
SSTORE: .REPT 24
        .WORD 0
        .ENDM
        .=.+26
  
```

```

1 11 000104 016700 177674
2
3
4
5
6
7
8
9
10
11 000110 016767 177664 000000G
12
13 000116 012701 000013
14
15
16
17 000122 004767 000000G
18 000126 016767 000000G 177644
19 000134 016700 000000G
20
21
22 000140 016704 177640
23
24 000144 012701 000013
25 000150 012002
26 000152 070202
27
28 000154 073227 000001
29 000160 010267 177616
30 000164 012002
31 000166 070202
32 000170 073227 000001
33 000174 066702 177602
34
35 000200 040224
36
37
38 000202 077116
39
40 000204 000207
41
42
43 000001*
  
```

.SETTL CODE FOR "PRSPEC".

3. CODE FOR "PRSPEC":

-----  
 THE CODE FOR POWER SPECTRUM SUBROUTINE STARTS  
 BELOW:

```

PRSPEC: MOV     SFCTRA,R0      ;LOAD THE STARTING
                                ADDRESS OF THE
                                OUTPUT ARRAY.
      MOV     CUMSF1,SFACTR   ;MOVE THE # OF POINTS
                                IN THE OUTPUT ARRAY
                                TO R1.
      MOV     #13,R1          ;SCALE THIS ARRAY.
      JSR     PC,SCALE        ;GET STARTING ADDRESS
                                OF INPUT ARRAY.
      MOV     SFACTR,CUMSF1   ;GET STARTING ADDRESS
                                OF THE OUTPUT ARRAY.
      MOV     DATA,PD        ;REAL PART.
      MOV     SFCTRA,R4       ;SQUARE THE REAL
                                PART.
      MOV     #13,R1          ;STORE THIS VALUE.
      MOV     (R0)+,R2        ;IMAG PART.
      MUL     R2,R2           ;SQUARE THE IMAG PART.
      ASHC   #1,R2           ;ADD SQUARE OF REAL
                                AND IMAG PART.
      MOV     R2,RLSOR        ;ADD THIS VALUE TO
                                THE PREVIOUS VALUE IN
                                THE ARRAY.
      MUL     (R0)+,R2        ;SQUARE THE IMAG PART.
      ASHC   #1,R2           ;ADD SQUARE OF REAL
                                AND IMAG PART.
      ADD     R2,(R4)+        ;ADD THIS VALUE TO
                                THE PREVIOUS VALUE IN
                                THE ARRAY.
      SOB    R1,PR$10        ;ALL POINTS DONE?
      RTS    PC              ;NO!GO AND DO MORE.
                                ;YES..RETURN.
      .END
  
```

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AIR FORCE INST OF TECH WRIGHT-PATTERSON AFB OHIO SCH--ETC F/G 9/2  
MICROPROCESSOR BASED DATA ACQUISITION AND PROCESSING SYSTEM.(U)  
DEC 78 S IFTEKHAR

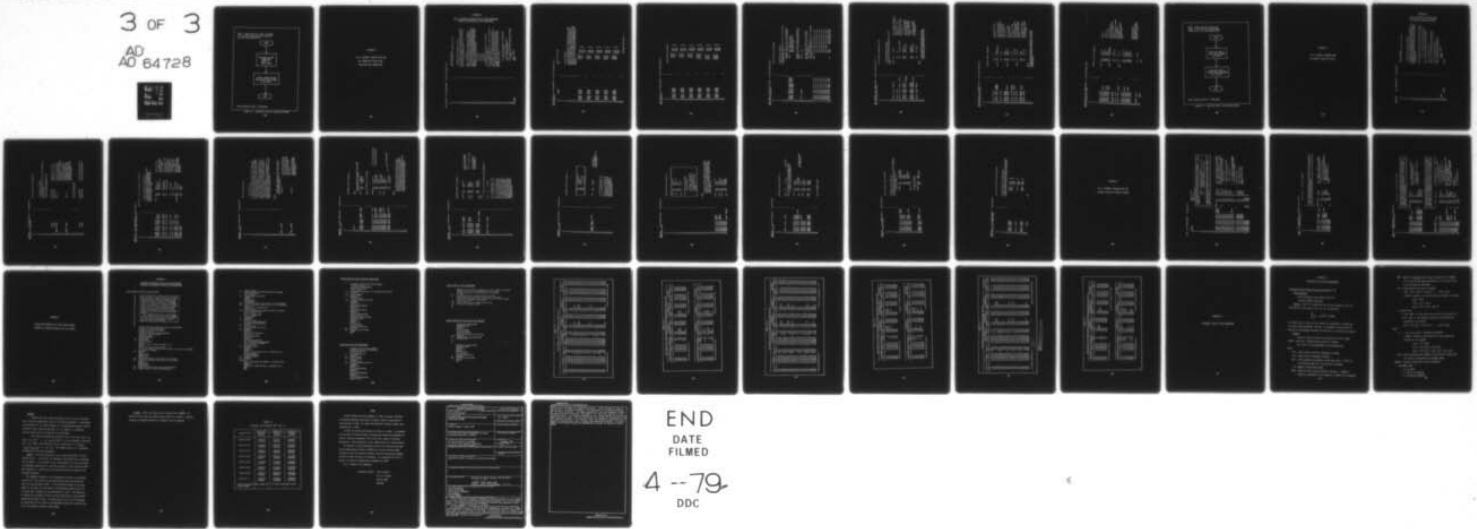
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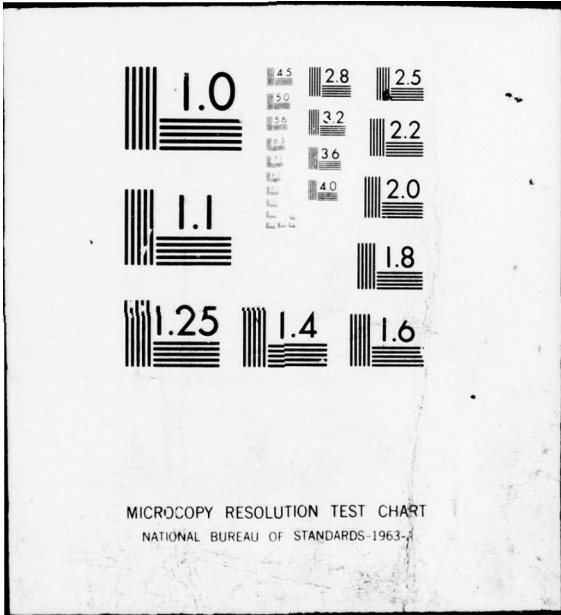
3 OF 3

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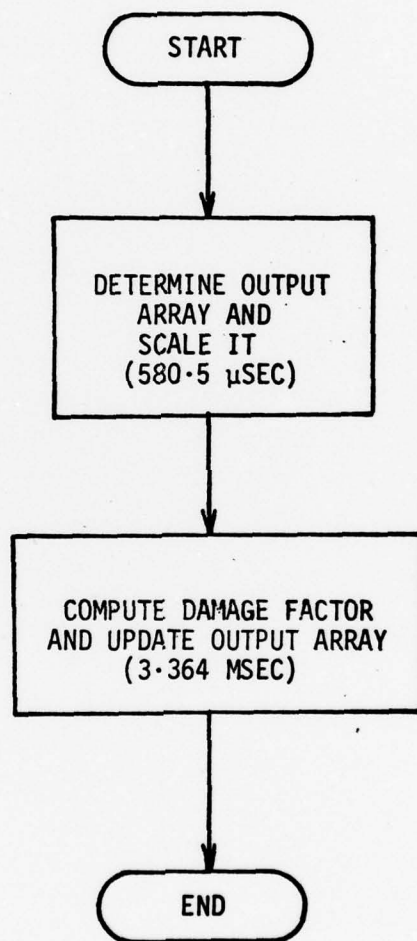


END  
DATE  
FILMED

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DDC



NOTE: These execution times are based on the instruction execution times of LSI-11M (Ref Appendix B).



Total Execution Time = 3.964 msec.

Figure F-1. Execution Time for Subroutine PRSPEC

APPENDIX G

LSI-11 ASSEMBLY LANGUAGE CODE FOR  
CDF SUBROUTINE (POLYN) AND  
EXECUTION TIME COMPUTATION

APPENDIX G

LSI-11 Assembly Language Code for CDF Subroutine  
(POLYN) and Execution Time Computation

POLY CURVE FIT RT-11 MACRO VM02-12 30-OCT-78 23:58:05 PAGE 1

```

1 .TITLE POLY CURVE FIT
2 .SEttl INTRODUCTION TO SUBROUTINE "POLYN".
3
4
5
6
7
8
9
10 1. INTRODUCTION TO SUBROUTINE "POLYN":
11 -----
12 THIS SUBROUTINE IS USED TO APPROXIMATE THE FUNCTION
13 THAT IS USED TO CALCULATE THE DAMAGE FACTOR(DF) AT
14 THE FREQUENCY OF INTEREST. THE DF CAN BE CALCULATED
15 USING THE FOLLOWING RELATIONSHIP:
16
17
18
19
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21
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49
50
51
52
53
54

```

DF=(X)\*\*1.6  
WHERE X IS POWER DENSITY VALUE AT THE FREQUENCY OF  
INTEREST. THE VALUE OF X LIES BETWEEN 0.0 AND 1.0  
BECAUSE THE SYSTEM WORKS WITH FRACTIONAL NUMBERS.  
A POLYNOMIAL CURVE FIT PROGRAM WAS USED TO APPROX-  
-IMATE THE DF CURVE, AND IT WAS DETERMINED THAT THE  
CURVE SHOULD BE DIVIDED INTO FOLLOWING INTERVALS.EACH  
INTERVAL CAN THEN BE APPROXIMATED BY A SECOND DEGREE  
POLYNOMIAL WITH MINIMUM RMS ERROR:  
0.0 TO 0.1  
0.1 TO 0.2  
0.2 TO 0.3  
0.3 TO 0.4  
0.4 TO 0.6  
0.6 TO 0.8  
0.8 TO 1.0  
THE POLYNOMIAL EQUATION USED TO APPROXIMATE  
DF FUNCTION IS GIVEN BELOW:  
Y=A\*\*2+B\*\*X+C  
FOR EACH INTERVAL DEFINED ABOVE THERE IS A UNIQUE SET  
OF COEFFICIENTS A,B,C.  
THESE COEFF: ARE STORED AT LOCATION STARTING AT  
LABEL "COEFF:". THE FIRST VALUE FOR EACH INTERVAL IS  
USED TO TEST FOR THE INTERVAL IN WHICH X LIES.  
ALSO IF THE VALUE FOR X IS LESS THAN 0.04 THEN  
ANOTHER SET OF COEFF: A,B,C STORED AT LABEL "CF1"  
ARE USED. THIS HAS BEEN DONE TO MINIMIZE THE ERROR  
FOR VERY SMALL VALUES.

NOTE:THE GLOBAL VARIABLES AND SUBROUTINES  
ARE DEFINED BELOW.  
.GLOBL SSTORE,TFLG,SCALE,SPCTRA  
.GLOBL POLYN,SFACR,CDFAC  
NOTE:THE REGISTER DEFINITION MACRO IS DEFINED  
BELOW!

```

53 .MCALL .REGDEF
54 .REGDEF
54 000000

```

DATA STRUCTURE.

.SBTTL DATA STRUCTURE.  
.CSECT

```

1 000000
2
3
4
5
6
7
8
9
10
11
12
13
14
15
16
17
18
19 000000 000000
20 000002 000000
21 000004 000024
22
23
24
25
26 000054 000001
27 000056 000000
28 000060 000024
29
30
31
32
33 000130 000000
34 000132 000002
35 000134 000024
36
37
38
39
40 000204 000003
41 000206 000000
42 000210 000024
43
44
45
46
47 000260 000004
48 000262 000000
49 000264 000024
50
51
52
53
54
55
    
```

2. DATA STRUCTURE:

-----  
THE DATA STRUCTURE FOR THE FINAL ARRAY, CDF VERSUS FREQUENCY AND TEMPERATURE, IS GIVEN BELOW. THERE ARE 11 ARRAYS CORRESPONDING TO 11 TEMPERATURE INTERVALS OF INTEREST. THE FIRST ELEMENT IN EACH ARRAY GIVES THE VALUE FOR TEMPERATURE INTERVAL, THE 2ND VALUE GIVES THE SCALE FACTOR FOR THAT ARRAY, AND THE REMAINING 11 ELEMENTS ARE THE CDF VALUES FOR THE FREQUENCY RANGE OF INTEREST:

```

ARRAY0  .WORD  0
        .TEMP0  0
        .DFSF01 0
        .CSTR0: 24
        .WORD  0
        .ENDM

ARRAY1  .WORD  1
        .TEMP1: 0
        .DFSF1: 0
        .CSTR1: 24
        .WORD  0
        .ENDM

ARRAY2  .WORD  2
        .TEMP2: 0
        .DFSF2: 24
        .CSTR2: 0
        .WORD  0
        .ENDM

ARRAY3  .WORD  3
        .TEMP3: 0
        .DFSF3: 0
        .CSTR3: 24
        .WORD  0
        .ENDM

ARRAY4  .WORD  4
        .TEMP4: 0
        .DFSF4: 24
        .CSTR4: 0
        .WORD  0
        .ENDM
    
```

CONTINUED ON NEXT PAGE.....

```

1 ;
2 ;
3 ;
4 ;
5 ;
6 ;
7 ;
8 000334 000005
9 000336 000000
10 000340 000024
11
12
13
14
15 000410 000006
16 000412 000000
17 000414 000024
18
19
20
21 000444 000007
22 000466 000000
23 000466 000000
24 000470 000024
25
26
27
28
29 000540 000010
30 000542 000000
31 000544 000024
32
33
34
35
36 000614 000011
37 000616 000000
38 000620 000024
39
40
41
42
43 000670 000012
44 000672 000000
45 000674 000024
46
47
48
49
50
  
```

DATA STRUCTURE CONTINUED.....

```

ARRAY5
TEMP5: .WORD 5
DFS5: .WORD 0
CSTR5: .REPT 24
        .WORD 0
        .ENDM

ARRAY6
TEMP6: .WORD 6
DFS6: .WORD 0
CSTR6: .REPT 24
        .WORD 0
        .ENDM

ARRAY7
TEMP7: .WORD 7
DFS7: .WORD 0
CSTR7: .REPT 24
        .WORD 0
        .ENDM

ARRAY10
TEMP10: .WORD 10
DFS10: .WORD 0
CSTR10: .REPT 24
        .WORD 0
        .ENDM

ARRAY11
TEMP11: .WORD 11
DFS11: .WORD 0
CSTR11: .REPT 24
        .WORD 0
        .ENDM

ARRAY12
TEMP12: .WORD 12
DFS12: .WORD 0
CSTR12: .REPT 24
        .WORD 0
        .ENDM
  
```

POLY CURVE FIT RT-11 MACRO VM02-12 30-OCT-78 23:58:05 PAGE 4  
 DEFINITION OF VARIABLES

1				
2				
3				
4				
5				
6				
7				
8				
9				
10				
11				
12				
13				
14	000744	000004	000060	000134
15	000752	000210	000264	000340
16	000760	000414	000470	000544
17	000766	000620	000674	
18				
19				
20				
21				
22				
23				
24	000772	000000		
25				
26				
27				
28				
29				
30				
31				
32				
33				
34	000774	006314	062746	003136
35	001002	177765		
36	001004	014631	040405	006470
37	001012	177641		
38	001014	023146	032444	010766
39	001022	177477		
40	001024	031463	027205	012763
41	001032	177214		
42	001034	046314	024157	015416
43	001042	176571		
44	001044	063146	021533	020441
45	001052	175646		
46	001054	017774	023205	174575
47	001062	044601	000726	177776

.SBTTL DEFINITION OF VARIABLES

3. DEFINITION OF VARIABLES:

-----  
 THE VARIABLES AND COEFFICIENTS USED IN THIS SUROU  
 -TINE ARE DEFINED BELOW:

VARIABLE "CDFAC" IS USED AS AN POINTER THAT  
 CONTAINS THE ADDRESS OF THE OUTPUT ARRAY.

CDFAC: .WORD CSTR0,CSTR1,CSTR2  
 .WORD CSTR3,CSTR4,CSTR5  
 .WORD CSTR6,CSTR7,CSTR10  
 .WORD CSTR11,CSTR12

VARIABLE "CSTSAV" IS USED AS A TEMPORARY  
 STORAGE FOR THE CALCULATED ADDRESS OF OUTPUT  
 ARRAY.

CSTSAV: .WORD 0

THE COEFFICIENTS USED IN THE POLYNOMIAL  
 ARE GIVEN BELOW. THEY ARE STORED IN THE  
 FOLLOWING MANNER:

.WORD FLAG,C,B,A  
 "FLAG" IS THE VALUE WHICH IS USED TO DETERMINE  
 THE INTERVAL IN WHICH X LIES.

COEFF: .WORD 006314,062746,003136,177765  
 .WORD 014631,040405,006470,177641  
 .WORD 023146,032444,010766,177477  
 .WORD 031463,027205,012763,177214  
 .WORD 046314,024157,015416,176571  
 .WORD 063146,021533,020441,175646  
 .WORD 017774,023205,174575  
 .WORD 044601,000726,177776

CF:





```

1
2
3
4
5
6
7
8
9
10 001214 012402
11 001216 070200
12 001220 073227 000001
13 001224 062402
14 001226 070200
15 001230 073227 000001
16 001234 062402
17 001236 060265 000004*
18
19
20 001242 005725
21 001244 077136
22
23 001246 000207
24
25 001250 012704 001062*
26
27 001254 012402
28 001256 070200
29 001260 073227 000001
30 001264 062402
31 001266 070200
32 001270 073227 000001
33 001274 061402
34 001276 006302
35
36
37
38
39
40 001300 000756
41
42
43 000001*
  
```

CODE FOR "POLYN" CONTINUED.....

```

PLY30: MOV (R4)+,R2
      MUL R0,R2
      ASHC #1,R2
      ADD (R4)+,R2
      MUL R0,R2
      ASHC #1,R2
      ADD (R4)+,R2
      R2,CSTRD(R5)

PLY35: ADD (R5)+
      TST R1,PLY10
      SOB RTS
      MOV #CF,R4

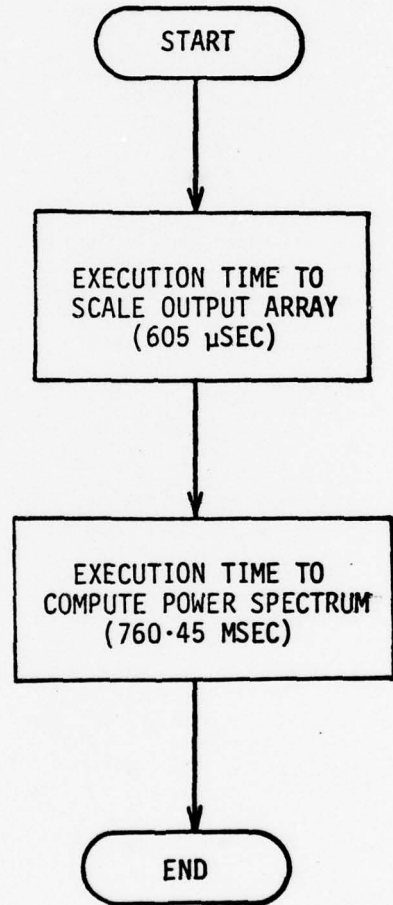
      MOV (R4)+,R2
      MUL R0,R2
      ASHC #1,R2
      ADD (R4)+,R2
      MUL R0,R2
      ASHC #1,R2
      ADD (R4)+,R2
      ASL R2

      BR PLY35

      .END
  
```

!THE CODE HERE IS  
 SAME AS IN PLY30  
 EXCEPT THAT THE FINAL  
 FINAL RESULT IS HALVED.  
 THEREFORE WE DOUBLE IT..

NOTE: These execution times are based on the instruction execution times of LSI-11M (Ref Appendix B).



Total Execution Time = 1.3655 msec.

Figure G-1. Execution Time for Subroutine POLYN

APPENDIX H

LSI-11 ASSEMBLY LANGUAGE CODE  
FOR CONTROL EXECUTIVE (EXEC)

APPENDIX H

LSI-11 Assembly Language Code  
for Control Executive (EXEC)

EXECUTIVE RT-11 MACRO VM02-12 00:06:08 PAGE 1

```

1 .TITLE EXECUTIVE
2 .SBTTL INTRODUCTION TO CONTROL EXECUTIVE
3
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37
38
39
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```

1. INTRODUCTION TO CONTROL EXECUTIVE:

-----

THIS PART OF THE CONTROL EXECUTIVE GETS CONTROL OF THE SYSTEM ONCE POWER HAS BEEN SWITCHED ON AND THE RUN/STANDBY SWITCH HAS BEEN PLACED IN RUN POSITION. THE EXECUTIVE WILL RELINQUISH CONTROL IF EITHER THE POWER HAS BEEN SWITCHED OFF OR THE SYSTEM HAS BEEN SWITCHED TO STANDBY MODE.

NOTE: THE GLOBL DATA STORAGE AREAS ARE DEFINED BELOW:

-----

```

.GLOBL DATA,ISTR1,ISTR2,CDATA,SPCTRA,TFLG

```

NOTE: THE GLOBAL SUBROUTINES ARE DEFINED BELOW:

-----

```

.GLOBL FOURIE,PRSPEC,POLYN

```

NOTE: THE GLOBAL VARIABLES ARE DEFINED BELOW:

-----

```

.GLOBL SFACR,CUMSF1,TSUM,TFLG

```

NOTE: THE REGISTER ARE DEFINED BELOW:

-----

```

.MCALL .REGDEF
.REGDEF

```

.CSECT

```

000000
000000*

```

```

1
2
3
4
5
6
7
8
9
10
11
12
13
14
15
16
17
18
19
20
21
22
23
24
25
26
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28
29
30
31
32
33
34
35
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37
38
39
40
41
42
43

        .SBTTL DEFINITION OF DEVICES USED

2. DEFINITION OF DEVICES:
-----
THE DEVICES THAT ARE CALLED IN BY THE EXECUTIVE
ARE DEFINED HERE.
-----
ANALOG TO DIGITAL CONVERTOR (ADC)
-----
        ADCSR=170400
        ADBFR=170402
        ADFLG: 0

        ADCNT: 0

        REAL TIME CLOCK (RTC)
        -----
        RTCSR=170420
        RTBFR=170422

        ADC CONTROL AND STATUS REGISTER.
        ADC DATA BUFFER REGISTER
        ADC FLAG WILL BE USED
        TO INDICATE THAT 32 DATA
        POINTS HAVE BEEN READ IN
        HAND STORED AT THE ALLOCATED
        STORAGE AREA.
        THIS FLAG IS CLEARED BY
        THE EXECUTIVE BEFORE ADC
        INTERRUPT IS ENABLED
        THIS COUNTER KEEPS TRACK
        OF NUMBER OF POINTS TO BE
        CONVERTED AND STORED. IT
        IS LOADED BY THE EXEC BEFO
        RE ADC IS ENABLED.

        RTC CONTROL AND STATUS REGI
        RTC DATA BUFFER REGISTER
    
```

```

1
2
3
4
5
6
7
8
9
10
11
12
13
14
15 000004 005067 177770
16 000010 012767 000520' 000476
17 000016 012767 000040 177756
18 000024 012737 000500' 000440
19
20
21 000032 012737 000100 000442
22 000040 012767 177740 170422'
23
24 000046 012737 000426' 000400
25
26
27 000054 012737 000100 000402
28
29 000062 106427 000100
30
31 000066 012767 007520 170400'
32 000074 000001
33 000076 032767 000001 177674
34
35 000104 001773
36 000106 000402
  
```

.SBTTL FPASS --FIRST PASS THROUGH THE EXECUTIVE

3. FIRST PASS THROUGH THE EXECUTIVE:

-----  
 FIRST TIME THROUGH THERE ARE NO DATA POINTS  
 AVAILABLE THEREFORE THE PROCESSOR REMAINS  
 IDLE AND ADC GATHERS 32 DATA POINTS. THE  
 DATA POINTS ARE STILL BEING CONVERTED AND  
 STORED VIA INTERRUPT SCHEME.

```

EXEC: CLR      ADFLC      ADCLR ADC FLAG
      MOV     #ISTR1.CDATA  !SET COUNTER TO 32.
      MOV     #40.ADCNT    !SET ADDRESS OF RTC
      MOV     #RTCSR.V.0#440 !SERVICE ROUTINE AT
                                   !VECTOR ADDRESS
      MOV     #100.0#442    !SET RTC BUFFER WITH
                                   !REQUIRED COUNT.
      MOV     #177740.RTBFR !SET ADDRESS OF ADC
      MOV     #ADCSR.V.0#400 !SERVICE ROUTINE AT
                                   !VECTOR ADDRESS
      MOV     #100.0#402    !SET PROCESSOR Prio
                                   !--RITY TO 0
      MTPS    #100         !SET PROCESSOR Prio
                                   !--RITY TO 0
      MOV     #7520.ADCSR   !ENABLE ADC
      MOV     #START2.WAIT !WAIT FOR INTERRUPTS.
      BIT     #1.ADFLGL    !HAVE 32 POINTS BEEN
                                   !CONVERTED?
      BEG    START2
      BR     CONT1
  
```

.SBTTL PROCESSING CONTROL

4. PROCESSING CONTROL:

```

-----
EACH SET OF 32 DATA POINTS ARE FIRST FOURIE TRANS
FORMED USING SUBROUTINE "FOURIE". THIS IS FOLLOWED
BY COMPUTATION OF POWER SPECTRAL DENSITY. THIS PROCESS
IS REPEATED FOR TEN SETS OF DATA POINTS AND AN AVERAGE
POWER SPECTRUM OBTAINED. NEXT THE DAMAGE FACTOR IS CAL
CULATED AT EACH FREQUENCY.
EVERY TIME THE POWER SPECTRUM IS CALCULATED THE
TEMPERATURE SIGNAL IS SAMPLED AND ACQUIRED. THIS ACQUI
RED VALUE IS SUMMED IN A STORAGE VARIABLE "TSUM".
THEREFORE AFTER THE AVERAGE POWER SPECTRUM IS OBTAINED
(AFTER 10 LOOP EXECUTIONS) THE AVERAGE TEMPERATURE
VALUE IS STORED IN "TSUM".
EVERY TIME THE PROCESSING LOOP IS STARTED THE ADC IS
ENABLED AND NEXT 32 DATA POINTS ARE GATHERED VIA
INTERRUPT SCHEME. THE REAL TIME CLOCK IS USED TO SYNCHRO
NIZE DIFFERENT PROCESSING LOOPS.
    
```

NOTE:

```

-----
BEFORE THESE POINTS ARE PROCESSED THE PROCESS
LOOP COUNTERS (PRCLOP AND ADCLOP) ARE DEFINED
AND INITIALIZED BELOW:
    
```

```

33 000110 000000
34
35
36
37
38
39
40
41 000112 000000
42
43
44
    
```

```

ADCLP: 0
    
```

```

;THIS COUNTER IS USED TO
;TELL THE ADC WHICH BUFFER
;STORAGE SHOULD BE USED TO
;STORE CONVERTED DATA. IT
;IS ALSO USED TO DETERMINE
;WHICH 32 POINTS NEED
;PROCESSING.
    
```

```

PRCLOP: 0
    
```

```

;THIS COUNTER IS USED TO
;DETERMINE HOW MANY POWER
;SPECTRUM SETS HAVE TO BE
;ADDED TOGETHER.
    
```

```

1
2
3
4
5
6
7
8
9
10 000114
11
12
13
14
15 000114 012767 000001 177766
16 000122 012767 000010 177762
17 000130 000400
18
19
20
21
22
23
24
25 000132 032767 000001 177750
26 000140 001411
27 000142 012767 000714' 000344
28 000150 012767 000520' 000340
29 000156 005067 177726'
30 000162 000410
31 000164 012767 000520' 000322
32 000172 012767 000714' 000316
33 000200 005267 177704
34 000204 012767 000040 177570
35 000212 005067 177562
36 000216 106427 000100
37
38 000222 012767 007520 170400'
39 000230 004767 000000G
40
41
42
43
44
45
46
47
48
49
50
51
52
53
    
```

PROCESSING CONTROL CONTINUED.....

NOTE:

-----  
 THE LOOP COUNTERS ARE INITIALIZED BEFORE  
 PROCESSING IS INITIALIZED.

CONT1: MOV #1,ADCL0P  
 BR #10,PRCL0P  
 PRCL0P

NOTE: PROCESSING LOOP STARTS NOW.

PROCESS: BIT #1,ADCL0P ;IS IT FIRST DATA BUFFER?  
 BEQ PKC2 ;YES,NOW USE ZDN BUFFER  
 MOV #1STR2,CDATA  
 MOV #1STR1,DATA  
 CLR ADCL0P  
 BR PKC3

PRC2: MOV #1STR1,CDATA  
 MOV #1STR2,DATA  
 INC ADCL0P

PRC3: MOV #40,ADCNT  
 CLR ADFLG  
 MTPS #100 ;SET PROCESSOR PRIO  
 -RITY TO 0

NOTE:

-----  
 AT THIS POINT THE VALUE OF SFACTR WILL SHOW  
 THE NUMBER BITS THE FOURIER TRANSFORMED ARRAY  
 HAS BEEN SHIFTED 1-BIT TO THE RIGHT. TO HAVE A  
 FIXED VALUE OF SCALE FACTOR THE TRANSFORMED AR  
 -RAY IS RESCALED SO THAT SFACTR IS EQUAL TO  
 3. THIS VALUE IS CHOSEN AFTER EXPERIMENTATION AND  
 IT CAN BE CHANGED IF FUTURE DATA WARRANTS A  
 DIFFERENT VALUE FOR SCALE FACTOR (SFACTR).

```

1
2
3
4
5
6
7
8
9
10 000234 016701 000000G
11 000240 162701 000003
12
13
14 000244 001410
15 000246 012703 000040
16 000252 016700 000240
17 000256 011004
18 000260 072401
19
20
21 000262 010420
22 000264 077304
23 000266 012767 000003 000000G
24
25
26
27
28
29
30 000274 004767 000000G
31
32
33
34
35
36
37
38
39
40
41
42
43
44
45
46
47
48
49
50
    
```

PROCESSING CONTROL CONTINUED.....

```

RESCL: MOV SFACTR,R1
SUB #3,R1
;STORE THE DIFFERENCE
;BETWEEN SFACTR AND 3
;IN R1.

REQ PR20
MOV #40,R3
MOV DATA,R0
MOV (R0),R4
ASH R1,R4
;GET # OF POINTS IN R3.

MOV R4,(R0)+
SOB R3,PR15
MOV #3,SFACTR
PR20: MOV
NOTE:
-----
PR21: JSR PC,PRSPEC
    
```

AT THIS STAGE THE SFACTR IS ALWAYS GOING TO 3.

```

*****
* IMPORTANT NOTE *
*****
AFTER THE POWER SPECTRUM OF
THE FIRST 11 FREQUENCY COMPO
-NENTS HAS BEEN CALCULATED
THE ACQUISITION OF TEMPERATU
-RE SIGNAL IS CARRIED OUT.
IN "STORE" AND LATER ADDED
TO VARIABLE "TSUM", THUS AN
AVERAGE TEMPT: VALUE WOULD
BE AVAILABLE BEFORE DF IS CA
-LCULATED. THE CODE GIVEN ON
NEXT PAGE HAS NOT YET BEEN
IMPLEMENTED BUT IT CAN BE AC
-TIVATED IN FUTURE.
    
```

1  
 2  
 3  
 4  
 5  
 6  
 7  
 8  
 9  
 10  
 11  
 12  
 13  
 14  
 15  
 16  
 17  
 18  
 19  
 20  
 21  
 22  
 23  
 24  
 25 000300 005767 177606  
 26 000304 001042  
 27  
 28  
 29  
 30  
 31  
 32  
 33  
 34  
 35  
 36  
 37  
 38  
 39  
 40  
 41  
 42

PROCESSING CONTROL CONTINUED.....

```

*****
* MOV #TSTORE,CDATA
* MOV #1,ADCNT
* CLR ADFLC
* MOV #7520,ADCSR
* WAIT
* ADD TSTORE,TSUM
*****
    
```

```

PR22:  TST  PRCL0P
        BNE  PR34
*****
        ;IS PRCL0P=0
        ;IF NOT GO BACK
        COMPUTE ANOTHER SET OF
        POWER SPECTRUM.
    
```

```

*****
* IMPORTANT NOTE
*****
    
```

BEFORE DF IS CALCULATED THE  
 THE TEMPERATURE INTERVAL IS  
 DETERMINED AND "IFLG" IS LO  
 -ADED WITH THE APPROPRIATE  
 VALUE. THE CODE GIVEN ON  
 NEXT PAGE IS INCLUDED FOR  
 FUTURE ACTIVATION.

```

1
2
3
4
5
6
7
8
9
10
11
12
13
14
15
16
17
18
19
20
21
22
23
24
25
26
27
28
29
30
31
32
33
34
35
36
37
38
39
40
41
42
43 000306 016701 000000G
44 000312 162701 000004
45 000316 001410
46 000320 012703 000013
47 000324 016700 000000G
48 000330 011004
49 000332 072401
50 000334 010420
51 000336 077304
52 000340 012767 000004 000000G
    
```

PROCESSING CONTROL CONTINUED.....

```

*****
* INTL: MOV #INTVL,R1
* #13,R2
* CLR TFLG
* INJ: CMP TSUM,(R1)+
* BLE. INS
* INC TFLG
* SOB R2,IN3
* INS: BR PR25
*
*
* INTVL: 0000,1111,2222,3333
* 4444,5555,6666,7777
* 8888,9999,10000
*
* (THE ABOVE 11 VALUES ARE
* FICTITIOUS. THEY HAVE TO
* BE FILLED UP WITH THE OC
* -TAL EQUIVALENT OF THE
* BOUNDARY VALUE FOR EACH
* INTERVAL).
*****
    
```

NOTE:  
 -----

THE VALUE STORED IN "CUMSF1" TELLS THE AMOUNT BY WHICH THE ARRAY "SCTRA" HAS BEEN SCALED. EXPLICITLY IT HAS BEEN DETERMINED THAT IF "CUMSF1" IS FIXED EQUAL TO 4, NO SIGNIFICANT LOSS IN ACCURACY OCCURS. THE CODE THAT FOLLOWS ACCOMPLISHES THIS.

```

RSCL2: MOV CUMSF1,R1
SUB #4,R1
BEQ PR23
PR24: MOV #13,R3
MOV SPCTRA,R0
RSCLP: MOV (R0),R4
ASH R1,R4
MOV R4,(R0)+
SOB R3,RSCLP
PR23: MOV #4,CUMSF1
    
```

```

1
2
3
4
5
6
7
8
9 000346 004767 000000G
10
11
12
13
14
15 000352 005067 000534
16 000356 005067 000526
17 000362 005067 000000G
18 000366 012767 000012 177516
19 000374 016700 000000G
20
21
22 000400 012701 000013
23 000404 005020
24 000406 077102
25 000410 000402
26
27 000412 005367 177474
28 000416 005067 000000G
29 000422 000167 177504
30
31
32
33
34
35
    
```

PROCESSING CONTROL CONTINUED.....

PR25: JSR PC:FOLYN ; COMPUTE DF

NOTE: ONCE THE DF IS COMPUTED THE WHOLE LOOP  
 ----- IS REPEATED AGAIN. THE FOLLOWING CLEARS  
 THE REQUISITE VARIABLES AND ARRAYS.

```

PCLR: CLR TSUM
      CLR TSTORE
      CLR CUMSFI
      MOV #12,FRCLOP
      MOV SPECTRA,RO
      ; THE POWER SPECTRUM
      ; IS CLEARED NOW.

PR30: MOV #13,R1
      CLR (R0)+
      SOB R1,PR30
      BR PR35

PR34: DEC PRCLOP
PR35: CLR SFACTR
      JMP PROCESS
    
```

END OF PROCESSING CONTROL

```

1
2
3
4
5
6
7
8
9
10
11
12
13 000426 016777 170402' 000060
14 000434 162777 004000 000052
15 000442 005267 000046
16 000446 005267 000042
17 000452 005367 177324
18 000456 001007
19 000460 005267 177314
20 000464 042767 000300 170400'
21 000472 106427 000200
22 000476 000002
23
24
25
26
27
28 000500 106427 000200
29 000504 012767 000100 177766
30 000512 000002
31
  
```

.SBTTL INTERRUPT SERVICE ROUTINES.

5. INTERRUPT SERVICE ROUTINES:

THE SERVICE ROUTINES TO HANDLE INTERRUPT GENERATED  
 BY THE ADC AND RTC ARE GIVEN BELOW:

ADC SERVICE ROUTINE:

```

ADCSR: MOV  ADRF,ACDATA
SUB  #4000,ACDATA
INC  CDATA
INC  CDATA
DEC  ADCNT
BNE  ADRTI
INC  ADFLG
BIC  #300,ADCSR
MPS  #200
ADRTI: RTI
  
```

RTC SERVICE ROUTINE:

```

RTCSR: MTPS #200
MOV  #100,RTCSR
RTI
  
```

.SBTTL DATA STRUCTURE AND VARIABLES USED

```

1
2
3
4
5
6
7
8
9
10
11
12 000514 000000
13 000516 000000
14 000520 000076
15
16
17
18 000714 000076
19
20
21
22 001110 000000
23 001112 000000
24
25
26
27 000001
  
```

6. DATA STRUCTURE AND VARIABLES USED:  
 -----  
 THE DATA STRUCTURE WHICH IS USED TO STORE ACQUIRED  
 DATA AND DIFFERENT VARIABLES USED ARE GIVEN BELOW:

```

CDATA: .WORD 0
DATA: .WORD 0
1STR1: .REPT 76
      .WORD 0
      .ENDM
1STR2: .REPT 76
      .WORD 0
      .ENDM
TSTORE: .WORD 0
TSUM: .WORD 0
  
```

.END

APPENDIX I

LSI-11 ASSEMBLY LANGUAGE CODE FOR  
DECIMAL AND OCTAL PRINTOUT (DECPR)





```

1  .SBTIL INTPT -- INTEGER PRINT ROUTINE
2
3
4 *****
5 ** PRINT DECIMAL INTEGER ROUTINE CALLED BY JSR PC,INTPT. USES
6 ** STANDARD CALLING SEQUENCE WITH ONE PARAMETER WHICH CONTAINS THE
7 ** BINARY INTEGER TO BE CONVERTED AND PRINTED. THE OUTPUT FORMAT IS
8 ** EQUIVALENT TO FORTRAN (IX,I7).
9 *****
10
11
12 INTPT: MOV 02(R5),R3
13 MOV #5,R1 ;R1 ← NUMBER OF DIGITS TO CONVERT
14 MOVB #40,6# ;SIGN ← BLANK FOR POSITIVE
15 TST R3 ;IF NUMBER IS NEGATIVE, THEN
16 BGE 1#
17 NEG R3 ; MAKE IT POSITIVE
18 MOVB #'-.6# ; REMEMBER TO PRINT A MINUS SIGN
19 MOV #20040,5# ;BLANK THE BUFFER
20 MOV #20040,5#+2
21 MOV #20040,5#+4
22 MOV #20060,5#+6
23
24
25 *****
26 ** THE FOLLOWING LOOP CONVERTS THE NUMBER TO A CHARACTER STRING ONE
27 ** DIGIT AT A TIME STARTING WITH THE UNITS DIGIT, THEN THE 10'S DIGIT,
28 ** 100'S DIGIT, ETC. THESE DIGITS ARE CONVERTED TO ASCII AND STORED
29 ** FROM RIGHT TO LEFT IN A BUFFER. THE BUFFER IS LATER PRINTED.
30 *****
31
32
33 CLR R2 ; FOR EACH DIGIT FROM 10**0 TO 10**5, DO
34 DIV #10,,R2 ; R2 ← NUMBER/10.
35
36
37 BIS #60,R3 ; R3 ← LEAST SIGNIFICANT DECIMAL
38 MOVB R3,5#+2(R1) ; CONVERT THE DIGIT TO ASCII
39 MOV R2,R3 ; PUT IT IN THE BUFFER
40 BEQ 3# ; IF ALL HIGHER ORDER DIGITS = 0,
41 ; THEN EXIT LOOP
42 INC R1,2# ; GO BACK FOR NEXT DIGIT.
43 MOVB 6#+5#+1(R1) ;R1 ← 1 (SOB DECREMENTS ONCE TOO OFTEN)
44 MOVB #5#,R2 ;BUFFER(R1) ← SIGN CHARACTER
45 MOV #10,,R1 ;R2 ← ADDRESS OF BUFFER
46 MOVB (R2)+,RO ;R1 ← NUMBER OF CHARACTERS TO PRINT
47 JSR PC,TYPE ;FOR I = 1 TO NUMBER OF CHARACTERS,
48 SOB R1,4# ; PRINT BUFFER(I)
49 RTS PC ;RETURN
50 ;**** BUFFER ***
51 ;*** SIGN CHARACTER ***
52 .END

```

APPENDIX J

SIMULATION PROGRAM FOR TESTING COMPUTATIONAL  
ACCURACY OF SOFTWARE MODULES AND TEST RESULTS

APPENDIX J

Simulation Program for Testing Computational Accuracy of Software Modules and Test Results

Source Code for Test Program (TFFT)

```

C *****
C * THIS PROGRAM IS USED TO TEST THE COMPUTATI- *
C * ONAL ACCURACY OF DIFFERENT MODULES OF THE *
C * MICROPROCESSOR BASED DATA ACQUISITION AND *
C * PROCESSING SYSTEM (MIRDAPS). THE PROGRAM *
C * ACCEPTS 32 DECIMAL DATA VALUES AND PERFORM *
C * AS THE SAME PROCESSING OF DATA AS MIRDAPS. *
C * SINCE THIS PROGRAM WAS EXECUTED ON A CDC- *
C * 6600, WHICH IS A 60-BIT MACHINE THEREFORE *
C * THE COMPUTATIONAL RESULTS WOULD BE MORE *
C * ACCURATE AS COMPARED WITH THE RESULTS OBT- *
C * AINED ON PDP-11/03, WHICH IS A 15-BIT MA- *
C * CHINE. *
C * THE RESULTS PRINTED OUT BY THIS PROGRA- *
C * M ARE IN DECIMAL AS WELL AS IN OCTAL FORM- *
C * AT. THIS HAS BEEN DONE FOR EASY COMPARISON *
C * WITH THE RESULT FROM ASSEMBLY LANGUAGE *
C * SOFTWARE WRITTEN FOR MIRDAPS. *
C *****
C
150 PROGRAM TFFT(INPUT,OUTPUT,TAPE5=JUTPJT,TAPE6=INPUT)
      DIMENSION IPDATA(32),PRS(12)
      DIMENSION F(32),PF(32),CCF(32)
      DIMENSION CONF(32)
      DIMENSION ERR1(32),ERR2(10),ERR3(10)
      DIMENSION LDATA(32),CF(12)
      DIMENSION X(32)
      COMPLEX RDATA(32)
      WRITE(5,150)
      FORMAT(1H1)
      DO 444 I=1,29,4
      PRINT*," "
      PRINT*," INPUT FOUR DATA VALUES....."
      PRINT*," "
      READ*, X(I),X(I+1),X(I+2),X(I+3)
      X(I)=X(I)*.99999 $ X(I+1)=X(I+1)*.999999 $ X(I+2)=X(I+2)*.999999
      X(I+3)=X(I+3)*.999999
444 CONTINUE
      M=32
      WRITE(5,222)
      PRINT*," "
      WRITE(5,232)
222 FORMAT(10X,30HTABLE BELOW SHOWS 32 REAL DATA)
232 FORMAT(10X,32HPPOINTS WHICH ARE FED IN AS INPUT)
      PRINT*," "
      PRINT*," "
      WRITE(5,121)
      WRITE(5,131)
121 FORMAT(3X,11HDATA POINTS,19X,11HDATA POINTS)
131 FORMAT(3X,10HIN DECIMAL,20X,8HIN OCTAL)
      -----

```

```

WRITE (5,242)
242  FORMAT(3X,3HODD,9X,4HEVEN,9X,3HODD,5X,4HEVEN)
      CALL OPRINT(X,M)
      DO 10 I=1,32
10    RDATA(I)=CMPLX(X(I),0.0)
      PRINT*," "
      PRINT*," "
      WRITE(5,252)
      WRITE(5,262)
252  FORMAT(3X,34HTABLE BELOW SHOWS THE COMPLEX DATA)
262  FORMAT(8X,36HARRAY WHICH IS FED TO FFT SUBROUTINE)
      PRINT*," "
      WRITE(5,121) & WRITE(5,131)
      WRITE(5,272)
272  FORMAT(7X,4HREAL,9X,4HIMAG,9X,4HREAL,5X,4HIMAG)
      CALL OPRINT(RDATA,64)
      N=5 & XI=1.
      CALL FFT(RDATA,N,XI)
      DO 111 J=1, 32
111  READ*,F(J)
      CONTINUE
      DO 276 I=1,16
      CONF(2*I-1)=REAL(RDATA(I))
      CONF(2*I)=AIMAG(RDATA(I))
275  CONTINUE
      DO 277 J=1, 32
      ERR1(J)=(CONF(J)-F(J))**2
277  CONTINUE
      PRINT*," "
      PRINT*," "
      CALL OPRINT1(RDATA,32)
      PRINT*," "
      CALL OPRINT1(F,32)
      PRINT*," "
      CALL OPRINT1(ERR1,32)
      CALL PRSPEC(RDATA,PRS,32,ERR2)
      CALL POLY(PRS,CF,11,ERR3)
      DO 777 I=1,32
777  SUM1=SUM1+ERR1(I)
      CONTINUE
      PRINT*," "
      PRINT*," "
      PRINT*," RMS ERROR FOR FFT=" ,(SUM1/32.)**.5
      DO 888 I=1,10
      SUM2=SUM2+ERR2(I)
      SUM3=SUM3+ERR3(I)
888  CONTINUE
      PRINT*," "
PRINT*," "
PRINT*," RMS ERROR FOR PWSPEC=" ,(SUM2/10.)**.5
PRINT*," "
PRINT*," "
PRINT*," PWS ERROR FOR CDF=" ,(SUM3/10.)**.5
STOP
END

```

### Source Code for Power Spectrum Subroutine

```
C
SUBROUTINE PRSPEC(RDATA,PRS,LEN,ERR2)
DIMENSION PRS(11)
DIMENSION PF(10),ERR2(10)
COMPLEX RDATA(LEN)
DO 10 K=1,11
PRS(K)=(REAL(RDATA(K)))**2.+(AIMAG(RDATA(K)))**2.
10 CONTINUE
DO 333 J=1,10
333 READ*,PF(J)
CONTINUE
DO 444 J=1,10
444 ERR2(J)=(PRS(J)-PF(J))**2
CONTINUE
PRINT*," "
PRINT*," "
CALL OPRINT1(FRS,10)
PRINT*," "
PRINT*," "
CALL OPRINT1(PF,10)
PRINT*," "
PRINT*," "
CALL OPRINT1(ERR2,10)
RETURN & END
SUBROUTINE OPRINT1(X,NP)
DIMENSION X(NP)
PRINT*," "
PRINT*," "
PRINT*," "
DO 10 I=1,NP
I1=X(I)*2.**15
WRITE(5,100) X(I),I1
10 CONTINUE
100 FORMAT(4X,G13.4,4X,06)
RETURN
END
```

### Source Code for CDF Subroutine

```
SUBROUTINE POLY(PRS,CF,NP,ERR3)
DIMENSION PRS(NP),CF(NP),CCF(10)
DIMENSION ERR3(10)
DO 10 K=1,NP
CF(K)=PRS(K)**1.6
10 CONTINUE
DO 555 J=1,10
555 READ*,CCF(J)
CONTINUE
DO 666 J=1,10
666 ERR3(J)=(CF(J)-CCF(J))**2
CONTINUE
PRINT*," "
PRINT*," "
CALL OPRINT1(CF,10)
PRINT*," "
PRINT*," "
CALL OPRINT1(CCF,10)
PRINT*," "
PRINT*," "
CALL OPRINT1(ERR3,10)
RETURN & END
```

### Source Code for FFT Subroutine

```
SUBROUTINE FFT(X,M,XI) $ COMPLEX X(1),J,M,T $ N=2**4 $ NV2=N/2
NM1=N-1 $ J=1 $ PI=3.1415926535698 $ DO 7 I=1,NM1
IF(I.GE.J) GO TO 5 $ T=X(J) $ X(J)=X(I) $ X(I)=T
5 K = NV2
IF(K.GE.J) GO TO 7 $ J=J-K $ K=K/2 $ GO TO 6
7 J=J+K $ DO 20 L=1,M $ LF=2**L $ LE1=-E/2 $ U=(1.,0.)
W=CEXP(CMPLX(0.,-XI**PI/LE1)) $ DO 20 J=1,LE1 $ DO 10 I=J,N,LE
IP=I+LE1 $ T=X(IP)*U $ X(IP)=X(I)-T
10 X(I) = X(I) + T
20 U=U*W $ IF(XI.GT.0.) RETURN $ DO 30 I=1,N
30 X(I)=X(I)/N $ RETURN $ END
C
C
```

### Source Code for Printing Out Test Results.

```
SUBROUTINE OPRINT(X,NP)
DIMENSION X(NP)
PRINT*," "
PRINT*," "
DO 10 I=1,NP,2
I1=X(I)**2.**15
I2=X(I+1)**2.**15
WRITE (5,100) X(I),X(I+1),I1,I2
10 CONTINUE
100 FORMAT (3X,2G13.4,4X,2(3X,06))
RETURN $ END
C
C
```

```
SUBROUTINE OPRINT1(X,NP)
DIMENSION X(NP)
PRINT*," "
PRINT*," "
PRINT*," "
DO 10 I=1,NP
I1=X(I)**2.**15
WRITE(5,100) X(I),I1
10 CONTINUE
100 FORMAT (4X,G13.4,4X,06)
RETURN
END
```

TABLE J-1

Test Results for FFT Computation (Data Set #1)

Mem Loc	INPUT DATA		OUTPUT FROM SUBROUTINE FOURIE (F(n))*		DESCALED ARRAY X(n) = SFACTR x F(n)		FFT OUTPUT FROM TEST PROGRAM Y(n)		ERROR SQUARE  Y(n) - X(n)  <sup>2</sup>	
	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	value in Decimal	value in Octal
4172	0.1153E+00	7301	0.3959E+00	31254	.7918	062531	.7922	062546	.1635E-06	000300
4174	0.5356E+01	3333	0.0000E+00	0	0.	000000	0.	000000	0.	000300
4176	0.1094E+00	7001	0.1454E+00	11233	.2908	022470	.2908	022471	.5092E-10	000300
4200	0.9134E-01	5661	0.6821E-01	4273	.1354	019565	.1359	019545	.2507E-05	000300
4202	0.7809E-01	4777	0.2376E+00	17151	.4752	035323	.4756	036337	.1446E-06	000300
4204	-0.2536E-01	176301	-0.1590E+00	165647	-.3180	753513	-.3181	753510	.1291E-07	000300
4206	0.8575E-01	5372	-0.3909E-01	175377	-.7810E-01	773000	-.7833E-01	772770	.6466E-07	000300
4210	-0.7751E-02	177402	-0.5646E-01	174306	-.1129	770614	-.1123	770520	.1639E-07	000300
4212	-0.1240E+00	170042	-0.1348E+00	167300	-.2635	756575	-.2698	756566	.4977E-07	000300
4214	0.1001E+00	6320	0.3842E-01	2353	.7680E-01	004724	.7690E-01	004727	.9710E-08	000300
4216	-0.8002E-01	172702	-0.3128E-01	175777	-.6250E-01	773777	-.6241E-01	774002	.7225E-08	000300
4220	-0.9378E-01	171777	-0.1467E+00	166470	-.5934	755161	-.5938	755145	.1364E-06	000300
4222	0.9982E-01	6307	-0.1513E+00	166241	-.3025	755704	-.3024	754513	.4076E-07	000300
4224	-0.1094E+00	170777	0.1396E+00	10740	.2792	021074	.2792	021674	.1055E-09	000300
4226	0.8841E-01	5521	-0.2086E+00	162514	-.6172	746331	-.6171	745234	.7810E-08	000300
4230	0.1071E+00	6666	-0.3230E+00	153251	-.5465	725517	-.5460	725516	.8946E-09	000300
4232	-0.3107E-01	176006	0.1386E+00	10676	.2772	021573	.2770	021565	.3032E-07	000300
4234	0.5466E-01	3377	-0.2229E+00	161571	-.4458	743350	-.4459	743156	.3667E-08	000300
4236	-0.7178E-01	173320	0.1016E+00	6402	.2032	015002	.2034	015011	.9880E-07	000300
4240	0.7224E-01	4477	0.4092E-01	2475	.8180E-01	005170	.8155E-01	005160	.5719E-07	000300
4242	0.1190E+00	7474	0.9055E-01	5627	.1811	013456	.1815	013472	.1352E-06	000300
4244	0.7571E-01	4661	0.1775E+00	13270	.3550	025560	.3550	026560	.5432E-11	000300
4246	-0.2927E-01	176101	0.2647E+00	20741	.5294	041703	.5294	041704	.2247E-08	000300
4250	-0.9546E-01	171710	-0.1681E+00	165174	-.1362	752367	-.1362	752367	.5230E-10	000300
4252	0.6244E-01	3776	-0.2488E+00	180046	-.4975	740116	-.4973	740106	.5212E-07	000300
4254	0.8575E-01	5372	-0.2912E+00	155273	-.5824	732563	-.5825	732561	.6523E-08	000300
4256	-0.4657E-01	175012	0.1426E+00	11101	.2852	022201	.2854	022210	.4339E-07	000300
4260	-0.7928E-01	172732	0.2236E+00	16236	.4472	034475	.4469	034465	.5930E-07	000300
4262	0.1094E+00	7000	0.1145E+00	7247	.2290	016517	.2293	016531	.9259E-07	000300
4264	0.9375E-01	6000	-0.1453E+00	166547	-.2906	755315	-.2908	755306	.5073E-07	000300
4266	0.9583E-01	6104	0.2098E+00	15333	.4196	032665	.4197	032672	.2073E-07	000300
4270	-0.1115E+00	170671	-0.1205E-01	177165	-.2410E-01	776352	-.2427E-01	775344	.2834E-07	000300

\*Scale factor SFACTOR equals 2.0.

The RMS error for FOURIE equals .00022

TABLE J-II  
Test Result for Power Spectrum and CDF Computation (Data Set #1)

Mem Loc	OUTPUT FROM SUBROUTINE PRSPEC [P(n)]*		DESCALED ARRAY R(n) = CURSFI x P(n)		POWER SPECTRUM RESULT FROM TEST PROGRAM [S(n)]		ERROR SQUARE [S(n) - R(n)] <sup>2</sup>	
	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal
23030	0.3915E-01	2403	.5264	050055	.6276	050124	.1411E-05	000000
23032	0.6439E-02	323	.1030	006457	.1030	006460	.1399E-08	000000
23034	0.2042E-01	1235	.3267	024721	.3274	024747	.4528E-06	000000
23036	0.1160E-02	46	.1850E-01	001136	.1866E-01	001151	.1274E-06	000000
23040	0.4883E-02	240	.7810E-01	004777	.7872E-01	005023	.7819E-0F	000000
23042	0.5615E-02	270	.8980E-01	005576	.9020E-01	005613	.1558E-06	000000
23044	0.1056E-01	532	.1689	012636	.1694	012656	.2416E-0E	000000
23046	0.3693E-01	2272	.5908	045637	.5913	045660	.2881E-0E	000000
23050	0.1721E-01	1064	.2753	021675	.2753	021504	.5521E-07	000000
23052	0.2960E-02	141	.4730E-01	007015	.4803E-01	003045	.5362E-05	000000

\*Scale factor CUMSFI equals 16.

The RMS error for PDSPEC equals .0006

Mem Loc	OUTPUT FROM SUBROUTINE POLYN [U(n)]*		DESCALED ARRAY V(n) = U(n) x CUMSF2		DF OUTPUT FROM TEST PROGRAM [U(n)]		ERROR SQUARE [R(n) - V(n)] <sup>2</sup>	
	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal
23310	0.2747E-02	132	.4639	035541	.4745	036275	.1174E-03	000003
23312	0.6104E-04	2	.1030E-01	000521	.2635E-01	001537	.2576E-03	000010
23314	0.9155E-03	36	.1546	011711	.1675	012661	.1669E-03	000005
23316	-0.1221E-03	177774	-.2060E-01	776534	.1741E-02	000071	.4991E-03	000020
23320	0.0000E+00	0	0.	000000	.1713E-01	001061	.2934E-03	000011
23322	0.6104E-04	2	.1030E-01	000521	.2130E-01	001271	.1209E-03	000003
23324	0.3052E-03	12	.5150E-01	003227	.5837E-01	007570	.4727E-04	000001
23326	0.2441E-02	120	.4122	032302	.4315	033471	.3708E-03	000014
23330	0.6714E-03	26	.1133	007200	.1271	001106	.1915E-03	000006
23332	-0.6104E-04	177776	-.1030E-01	777256	.7771E-02	000376	.3265E-03	000012

\*Scale factor CUMSF2 equals 168.89.

The RMS error for POLYN equals .0154

TABLE J-III  
Test Results for FFT Computation (Data Set #2)

Item Loc	INPUT DATA		OUTPUT FROM SUBROUTINE		ENRIE [F(n)]*		DESCALED ARRAY		FFT OUTPUT FROM TEST PROGRAM		ERROR SQUARE	
	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal
4172	0.1093E+00	6777	0.2899E-01	1666	.1159	00732E	.1348	010500	.3570E-03	000013	000013	
4174	-0.1098E+00	170761	0.0000E+00	0	0.	000000	0.	000000	0.	000000	000000	
4176	-0.7639E-01	173071	0.2252E-01	1342	0.9090E-01	005505	.8650E-01	905422	.1228E-04	000000	000000	
4200	0.1111E+00	7072	-0.1263E-01	177142	-5.050E-01	774641	-3.321E-01	775677	.2389E-03	000011	000011	
4202	0.1084E+00	6737	0.3854E-01	2357	.1541	011671	.1375	010632	.2743E-03	000011	000011	
4204	-0.9833E-01	171552	-0.6106E-01	174061	-.2240	761523	-.2516	757714	.7807E-03	000030	000030	
4206	-0.4468E-01	175110	-0.1013E+00	171411	-.4052	745042	-.3952	746550	.9927E-04	000007	000007	
4210	0.1048E+00	6552	0.1048E+00	6552	4.192	032650	4.038	031657	.2359E-03	000007	000007	
4212	0.9349E-01	5776	0.1535E-01	767	.6140E-01	003733	.7433E-01	304503	.1671E-03	000007	000007	
4214	0.1089E+00	7021	0.1743E-01	1073	.5370E-01	004453	.8197E-01	005175	.1505E-03	000004	000004	
4216	-0.1094E+00	170777	0.5524E-01	3422	.2209	014106	.2062	015145	.7155E-03	000007	000007	
4220	-0.1231E+00	170077	-0.2560E+00	157472	-1.024	676495	-1.015	677040	.9900E-04	000002	000002	
4222	-0.1233E+00	170070	-0.9064E-02	177327	-.3620E-01	775535	-.4314E-01	775172	.6817E-04	000001	000001	
4224	0.1559E-01	777	0.2274E+00	16433	.9095	072155	.8928	071110	.2807E-03	000011	000011	
4226	0.1529E-01	765	0.5081E-01	3201	.2032	015002	.2209	016105	.3121E-03	000012	000012	
4230	0.9375E-01	6000	0.8942E-01	5562	.3576	025705	.3539	025513	.1399E-04	000000	000000	
4232	0.9821E-01	6272	0.2652E+00	20762	1.060	107656	1.061	107713	.3000E-06	000000	000000	
4234	-0.7648E-01	173066	0.1554E+00	11752	.6224	047652	.6404	050767	.3225E-03	000012	000012	
4236	-0.1182E+00	170337	0.3433E-01	2145	.1373	010823	.1196	007517	.7177E-07	000012	000012	
4240	0.1086E+00	6745	0.1398E+00	10744	.5592	043623	.5552	043417	.1525E-04	000009	000009	
4242	0.1225E+00	7655	0.7077E-01	4417	.2330	022071	.2901	022442	.5050E-04	000001	000001	
4244	0.7806E-01	4776	0.5545E-01	3431	.2218	016143	.2049	015072	.2957E-03	000011	000011	
4246	-0.1096E+00	170770	-0.4193E-01	175242	-.1677	765210	-.1528	766160	.2218E-03	000007	000007	
4250	-0.1206E+00	170220	-0.3931E-01	175370	-.1572	765740	-.1475	766476	.9401E-04	000002	000002	
4252	-0.1562E-01	177000	0.1993E-01	1215	.7970E-01	005063	.6883E-01	004215	.1555E-03	000002	000002	
4254	-0.3125E-01	176000	-0.2713E-01	176207	-.1035	771034	-.9611E-01	771567	.1535E-03	000002	000002	
4256	-0.2049E-01	174532	-0.5417E-01	174421	-.2166	762106	-.2264	761403	.9633E-04	000012	000012	
4260	0.1133E+00	7200	0.3912E-01	2402	.1564	012004	.1413	011025	.2825E-03	000007	000007	
4262	0.1074E+00	6700	0.2908E-01	1671	.1163	007342	.1332	010416	.2872E-03	000011	000011	
4264	-0.1106E+00	170731	0.1370E-01	701	.5480E-01	003403	.8767E-01	003032	.5080E-04	000001	000001	
4266	-0.9645E-01	171641	0.5658E-01	3476	.2263	016367	.2300	016560	.1369E-04	000000	000000	
4270	0.1116E+00	7112	-0.2606E-01	176252	-.1042	771251	-.0707E-01	772332	.2933E-03	000011	000011	

\*Scale factor SFACTOR equals 4.0.

The RMS error for FOURIE equals .0027.

TABLE J-IV  
Test Results for Power Spectrum and CDF Computation (Data Set #2)

Mem Loc	OUTPUT FROM SUBROUTINE PRSPEC [P(n)]*		DESCALED ARRAY R(n) = CUMSF1 x F(n)		POWER SPECTRUM RESULT FROM TEST PROGRAM [S(n)]		ERROR SQUARE [S(n) - R(n)] <sup>2</sup>	
	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal
23030	0.1831E-03	6	.1170E-01	000577	.1817E-01	001123	.4195E-04	000001
23032	0.1527E-03	5	.9000E-03	00003F	.5585E-03	000431	.5905E-04	000001
23034	0.1282E-02	52	.8200E-01	005176	.8221E-01	005205	.4425E-07	000000
23036	0.5780E-02	255	.3379	025500	.3193	024335	.3470E-03	000013
23040	0.9155E-04	3	.5000E-03	000020	.1224E-01	000621	.1379E-03	000004
23042	0.1712E-01	1061	1.096	106074	1.072	104466	.5581E-03	000022
23044	0.1291E-01	647	.8262	064700	.7990	067106	.7381E-03	000030
23046	0.2625E-02	126	.1680	012601	.1740	013105	.3610E-04	000001
23050	0.2362E-01	1406	1.512	140574	1.535	142215	.5739E-03	000022
23052	0.5157E-02	251	.3300	025075	.3225	024510	.5591E-04	000001

\*Scale factor CUMSF1 equals 64.

The RMS error for PRSPEC equals .0032.

Mem Loc	OUTPUT FROM SUBROUTINE POLYN [U(n)]*		DESCALED ARRAY V(n) = U(n) x CUMSF2		DF OUTPUT FROM TEST PROGRAM [X(n)]		ERROR SQUARE [X(n) - V(n)] <sup>2</sup>	
	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal
23310	-0.1221E-03	177774	-.1895	767676	.1540E-02	000065	.3653E-01	002255
23312	-0.1221E-03	177774	-.1895	767676	.4942E-03	000020	.3510E-01	002236
23314	-0.1221E-03	177774	-.1895	763676	.1836E-01	001131	.4321E-01	002507
23316	0.6104E-04	2	.9470E-01	006037	.1809	012231	.4387E-02	000217
23320	-0.1221E-03	177774	.1895	014101	.8723E-03	000034	.3558E-01	002215
23322	0.6714E-03	26	1.042	102540	1.118	107416	.5719E-02	000273
23324	0.4272E-03	16	.6530	052335	.6984	054545	.1253E-02	000051
23326	-0.6104E-04	177776	-.9470E-01	771740	.6094E-01	007714	.2622E-01	001431
23330	0.1160E-02	46	1.800	163163	1.986	177074	.7453E-01	002153
23332	0.0000E+00	0	0.	000000	.1636	312357	.2675E-01	001554

\*Scale factor CUMSF2 equals 1552.09.

The RMS error for POLYN equals .1515.

TABLE J-V  
Test Results for FFT Computation (Data Set #3)

Mem Loc	INPUT DATA		OUTPUT FROM SUBROUTINE		FOURIE [F(n)]*		X(n) = SFACTOR x F(n)		TEST PROGRAM [Y(n)]		ERROR SQUARE	
	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal
4172	-0.1105E+00	170717	-0.2319E+00	161120	-0.4639	747242	-0.4632	747265	0.	000000	.3532E-06	900000
4174	-0.1194E+00	170271	0.0000E+00	0	0.	000000	0.	000000	0.	000000	0.	000000
4176	-0.3125E-01	176000	-0.1302E+00	167524	-0.2604	757253	-0.2607	757240	.4083E-06	000000	.4083E-06	000000
4200	0.7755E-01	4755	-0.1003E+00	171451	-0.2006	763122	-0.2012	767075	.8608E-06	000000	.8608E-06	000000
4202	0.1048E+00	6655	0.5188E-03	21	-0.1000E-03	000003	-0.1033E-02	000041	.8608E-06	000000	.8608E-06	000000
4204	0.1158E+00	7321	-0.2823E-01	176143	-0.5640E-01	776307	-0.5656E-01	774277	.5656E-07	000000	.5656E-07	000000
4206	-0.8743E-01	172317	-0.1031E+00	171315	-0.2062	762633	-0.2063	762626	.2173E-07	000000	.2173E-07	000000
4210	-0.7394E-01	173211	-0.1165E+00	170425	-0.2330	761056	-0.2334	761036	.1935E-06	000000	.1935E-06	000000
4212	-0.1208E+00	170211	-0.5333E+00	135676	-4.067	674571	-4.067	674560	.7093E-07	000000	.7093E-07	000000
4214	-0.7785E-01	173011	-0.8028E-01	172672	-0.1605	765564	-0.1608	765551	.1131E-05	000000	.1131E-05	000000
4216	0.6665E-01	4210	-0.3154E+00	153642	-0.5308	727501	-0.5309	727530	.1102E-06	000000	.1102E-06	000000
4220	0.1050E+00	6560	0.8075E-01	5126	-0.1615	012254	-0.1611	012235	.1885E-06	000000	.1885E-06	000000
4222	0.1018E+00	6407	0.3546E+00	26544	-0.7092	058307	-0.7093	058313	.1577E-07	000000	.1577E-07	000000
4224	-0.7620E-01	173077	0.1705E+00	12723	.1410	025645	.1410	025647	.1860E-08	000000	.1860E-08	000000
4226	-0.1172E+00	170377	-0.1853E+00	164107	-0.3705	750220	-0.3708	750210	.5388E-07	000000	.5388E-07	000000
4230	0.1013E+00	6366	-0.1202E-01	177166	-0.2434E-01	776354	-0.2417E-01	776347	.1762E-07	000000	.1762E-07	000000
4232	0.7751E-01	4754	0.5621E-01	3462	-0.1124	007143	-0.1125	007145	.8585E-06	000000	.8585E-06	000000
4234	-0.7885E-01	172747	0.1501E+00	11466	.3002	023154	.3004	023155	.5245E-07	000000	.5245E-07	000000
4236	-0.9576E-01	171676	0.1373E-02	55	-0.2000E-03	000006	-0.2896E-02	000136	.7270E-05	000000	.7270E-05	000000
4240	0.7886E-01	5030	-0.3140E+00	153720	-0.5280	727635	-0.5281	727632	.7433E-06	000000	.7433E-06	000000
4242	-0.5355E-01	174444	0.1441E+00	11162	.2882	022743	.2881	022742	.2565E-06	000000	.2565E-06	000000
4244	0.6250E-01	4000	0.1393E+00	10725	.2786	021651	.2789	021660	.4932E-07	000000	.4932E-07	000000
4246	0.6442E-01	4077	-0.6567E-01	173430	-0.1313	767461	-0.1311	767467	.3305E-07	000000	.3305E-07	000000
4250	-0.1161E+00	170444	-0.7465E-01	173162	-0.1493	766343	-0.1492	766346	.8725E-05	000000	.8725E-05	000000
4252	-0.1139E+00	170555	0.2025E+00	14782	.4050	031727	.4052	031734	.2991E-07	000000	.2991E-07	000000
4254	-0.6253E-01	173777	0.1321E+00	10352	.2642	020721	.2644	020730	.5167E-07	000000	.5167E-07	000000
4256	0.7782E-02	377	0.1547E+00	11715	.3094	023632	.3099	023652	.2428E-06	000000	.2428E-06	000000
4260	0.7059E-01	4411	0.2433E+00	17445	.4866	037110	.4864	037101	.4605E-07	000000	.4605E-07	000000
4262	-0.9225E-01	172061	-0.9668E-01	171641	-0.1933	763501	-0.1933	763502	.4195E-09	000000	.4195E-09	000000
4264	0.1163E+00	7344	0.2798E-01	1625	-0.5590E-01	003447	-0.5576E-01	003443	.1914E-07	000000	.1914E-07	000000
4266	-0.1250E+00	170001	-0.1101E+00	170750	-0.2202	761720	-0.2202	761721	.4724E-09	000000	.4724E-09	000000
4270	-0.6308E-01	173755	0.8505E-01	5343	.1701	012705	.1697	012672	.1232E-06	000000	.1232E-06	000000

\*Scale factor SFACTOR equals 20. THIS PAGE IS BEST QUALITY PRINTOUT FROM COPY FURNISHED TO DDG. The RMS error for FOURIE equals .00057

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FROM COPY FURNISHED TO DDG

TABLE J-VI

Test Results for Power Spectrum and CDF Computation (Data Set #3)

Item Loc	OUTPUT FROM SUBROUTINE FRSPEC [P(n)]*		DESCALED ARRAY R(n) = CUMSFT x P(n)		POWER SPECTRUM RESULT FROM TEST PROGRAM [S(n)]		ERROR SQUARE [S(n) - R(n)] <sup>2</sup>	
	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal
23030	0.3337E-02	156	.2148	015576	.2146	015566	.5784E-07	000000
23032	0.1648E-02	66	.1054	065575	.1055	00F742	.9461E-05	000000
23034	0.3052E-04	1	.1000E-03	000003	.1005	000151	.9679E-05	000000
23036	0.1465E-02	60	.9370E-01	005776	.3211E-02	000151	.1135E-04	000000
23040	0.1816E-01	1123	1.162	112302	.9707E-01	006154	.3505E-05	000000
23042	0.6592E-02	330	.4218	032775	1.164	112430	.4999E-05	000000
23044	0.9644E-02	474	.5172	047400	.4239	033103	.5102E-05	000000
23046	0.2136E-02	106	.1167	010577	.6195	047512	.1346E-05	000000
23050	0.1587E-02	64	.1015	006375	.1301	010655	.2009E-05	000000
23052	0.6134E-02	311	.3924	031072	.3945	031177	.4413E-05	000000

\*Scale factor CUMSFT equals 64.

The RMS error for PPSPEC equals .00228

Item Loc	OUTPUT FROM SUBROUTINE PCLYN [U(n)]*		DESCALED ARRAY V(n) = U(n) x CUMSFT <sup>2</sup>		DF OUTPUT FROM TEST PROGRAM [M(n)]		ERROR SQUARE [M(n) - V(n)] <sup>2</sup>	
	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal	Value In Decimal	Value In Octal
23310	-0.6104E-04	177776	-.9470E-01	771740	.8521E-01	005350	.1237E-01	007044
23312	-0.1221E-03	177774	-.1695	763676	.2851E-01	001651	.4757E-01	007026
23314	-0.1221E-03	177774	-.1895	763676	.1025E-03	000003	.3595E-01	002231
23316	-0.1221E-03	177774	-.1895	763676	.2395E-01	001420	.4555E-01	002724
23320	0.7324E-03	30	1.137	110577	1.275	121470	.1917E-01	001154
23322	0.6104E-04	2	.9470E-01	006037	.2533	020154	.2515E-01	001470
23324	0.2441E-03	10	.3788	030174	.4648	035574	.7357E-02	000362
23326	-0.6104E-04	177776	-.9470E-01	771740	.4210E-01	002543	.1871E-01	001145
23330	-0.1221E-03	177774	-.1894	763700	.2530E-01	001539	.4653E-01	002764
23332	0.6104E-04	2	.9470E-01	006037	.2258	016346	.1718E-01	001052

\*Scale factor CUMSFT equals 1552.09.

The RMS error for POLYN equals .1719

APPENDIX K

POLYNOMIAL CURVE FITTING SUBROUTINE

APPENDIX K

Polynomial Curve Fitting Subroutine

Polynomial Curve Fitting Subroutine (PLSCF)(Ref 15)

Identification:

PLSCF Polynomial Least Squares Curve Fit

CDC 6600 FORTRAN Subroutine

Purpose. PLSCF will compute the best fitting polynomial  $p(x)$  to a discrete set of data points  $(x_i, y_i)$ ,  $i=1, NP$  such that

$$\sum_{i=1}^{NP} [y_i - p(x_i)]^2 = \text{minimum}$$

The coefficients of  $p(x)$  will be returned for polynomials of degree  $\leq 6$ . For higher degree polynomial fits and, if requested, for polynomials of degree  $\leq 6$ , the Chebyshev polynomial multipliers are returned (see remarks).

Control:

Dimension  $X(NP), Y(NP), W(NP), C(NC), \text{WORK}(N+1)(N+6)/2$ ) where  $N=NDEG$ , and  $NC=N+1$  if  $NDEG > 0$ ,  $NC=(N+1)(N+2)/2$  if  $NDEG < 0$

CALL PLSCF (X,Y,W,NP,NDEG,NMAX,C,IN,XD,XO,WORK,IER)

where

$X(I)$  = Input array of distinct independent variables.

$Y(I)$  = Input array of dependent variables.

$W(I)$  = Input weights corresponding to each data point. If  $W(1) = 0$ , all  $W(I)$  are taken as +1 and need not be defined.

$NP$  = Number of given data points.

$NDEG$  = Degree of least square polynomial requested. If  $NDEG$  is negative, polynomials of all degrees 0 to  $NMAX$  will be computed.

NMAX = Degree of calculated least square polynomial (0 to |NDEG|).

This is the maximum degree polynomial fit for which no loss of significance was indicated.

C(I) = Output coefficient array as follows:

$$p(x) = C(1) + C(2) x + C(3) x^2 + \dots + C(N + 1) x^2$$

If NDEG is negative, the coefficients are returned as follows:

$$P_0(x) = C(1)$$

$$P_1(x) = C(2) + C(3) x$$

$$P_2(s) = C(4) + C(5) + C(6) x^2$$

IN Input Flag

= 0 and NDEG  $\leq 6$ , the coefficients of  $p(x)$  are returned in C.

$\neq 0$  or NDEG  $\geq 6$ , the Chebyshev polynomial multipliers are returned in C as follows:

$$p(x) = C(1) T_0(t) + C(2) T_1(t) + \dots + C(N+1) T_N(t)$$

where

$t = x \cdot XD + X$  and  $T_i(t)$  = Chebyshev polynomials

If NDEG is negative, the multipliers of each degree are returned in C as follows:

$$p_0(x) = C(1) T_0(t)$$

$$p_1(x) = C(2) T_0(t) + C(3) T_1(t)$$

$$p_2(x) = C(4) T_0(t) + C(5) T_1(t) + C(6) T_2(t)$$

XD, X0 = Output multiplicative constant XD and additive constant X0 for lineary transformation of argument range.

WORK(I) = Working storage area of  $((N+1) (N+6)/2)$

IER Output Flag

= 0, no errors

= 1, errors in dimension

= 2, coinciding arguments

### Methods.

- Determining least squares polynomial fits to a set of discrete data is generally meaningful only for low degree polynomials. Determining the polynomial fit for higher degrees in a straightforward manner involves solving a linear system of equations  $AC = B$ , where A is a positive definite matrix which is frequently ill-conditioned.

- This routine computes the polynomial  $p(x)$  in the form  $p(x) = c_1 T_0(t) + c_2 T_1(t) + \dots + c_m T_{m-1}(t)$  where  $T_i$  are the Chebyshev polynomials and  $t$  is a linear transformation of the original data,  $(t = (2x - x_{\max} - x_{\min})) / (x_{\max} - x_{\min}) = x \cdot XD + XO$ ). This method results in a remarkable improvement of the normal equations.

Remarks. With this subroutine, one can obtain polynomial fits for degrees up to 6. In addition, the Chebyshev coefficients may be obtained for any degree. If one wishes to use a high degree fit, one can evaluate the Chebyshev expansion for a specified argument  $x$  using subroutine CHEB with argument  $t = x \cdot XD + XO$  and the calculated coefficient vector of the Chebyshev expansion.

The Chebyshev expansion of the polynomial  $p(x)$  gives a much better indication of the accuracy of the approximation than the coefficient vector of the polynomial itself. If the specified degree of the polynomial is too high, the last terms of the Chebyshev expansion will be uniformly small compared to the coefficients in front. The degree may be reduced by the number of small trailing coefficients without unduly enlarging the overall error. An upper bound for the error introduced by neglecting the last terms of the Chebyshev expansion is given by the sum of the absolute values of these terms.

Storage. PLSCF uses  $672_8$  words of storage and no COMMON. In addition to the input and output vector arrays X,Y, W and C, a working storage area  $WORK((N+1)(N-6)/2)$ ,  $N=|NDEG|$ , must be supplied.

TABLE K-I  
 Constants for Polynomial  $Ax^2 + Bx + C$

Interval for x	Value of A	Value of B	Value of C
	Decimal (Octal)*	Decimal (Octal)*	Decimal (Octal)*
0.00 to 0.04	-.3765E-04 (177776)	.1435E-01 (000726)	.5743E+00 (044601)
0.04 to 0.10	-.3263E-03 (177765)	.4974E-01 (003136)	.7961E+00 (062746)
0.10 to 0.20	-.2890E-02 (177641)	.1033E+00 (006470)	.5080E+00 (040405)
0.20 to 0.30	-.6618E-02 (177477)	.1403E+00 (010766)	.4152E+00 (032444)
0.30 to 0.40	-.1134E-01 (177214)	.1715E+00 (012763)	.3635E+00 (027205)
0.40 to 0.60	-.1973E-01 (176571)	.2114E+00 (015416)	.3159E+00 (024157)
0.60 to 0.80	-.3397E-01 (175646)	.2588E+00 (020441)	.2762E+00 (021533)
0.80 to 1.0	.5085E-01 (174575)	.3010E+00 (023205)	.2499E+00 (017774)

\*These octal equivalent values for A, B, and C were used in sub-routine POLYN.

### Vita

Saleem Iftkhar was born November 11, 1949 at Karachi, Pakistan. He attended Cathedral High School at Lahore, where he completed his matriculation in 1965. He joined the Government College, Lahore, and completed FSc in 1967.

In 1967, he joined the Pakistan Air Force as a cadet. He attended the Pakistan Air Force College of Aeronautical Engineering (PAFCAE) at Karachi, where he graduated in 1971 with a B.E. degree in Avionics Engineering. Upon graduation, he was commissioned as a flying officer.

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SECURITY CLASSIFICATION OF THIS PAGE (When Data Entered)

REPORT DOCUMENTATION PAGE		READ INSTRUCTIONS BEFORE COMPLETING FORM
1. REPORT NUMBER AFIT-GE/EE/78D-29	2. GOVT ACCESSION NO.	3. RECIPIENT'S CATALOG NUMBER
4. TITLE (and Subtitle) MICROPROCESSOR BASED DATA ACQUISITION AND PROCESSING SYSTEM		5. TYPE OF REPORT & PERIOD COVERED M.S. Thesis
7. AUTHOR(s) Saleem Iftekhhar, Flt/Lt, PAF		6. PERFORMING ORG. REPORT NUMBER
9. PERFORMING ORGANIZATION NAME AND ADDRESS Air Force Institute of Technology (AFIT-EN) Wright-Patterson AFB, OH 45433		8. CONTRACT OR GRANT NUMBER(s)
11. CONTROLLING OFFICE NAME AND ADDRESS Structural Mecahnics Division Air Force Flight Dynamics Laboratory Wright-Patterson AFB OH 45433		10. PROGRAM ELEMENT, PROJECT, TASK AREA & WORK UNIT NUMBERS
14. MONITORING AGENCY NAME & ADDRESS (if different from Controlling Office)		12. REPORT DATE December 1978
		13. NUMBER OF PAGES 237
		15. SECURITY CLASS. (of this report)
		15a. DECLASSIFICATION/DOWNGRADING SCHEDULE
16. DISTRIBUTION STATEMENT (of this Report) Approved for public release; distribution unlimited.		
17. DISTRIBUTION STATEMENT (of the abstract entered in Block 20, if different from Report)		
18. SUPPLEMENTARY NOTES Approved for public release; IAW AFR 190-17  JOSEPH P. HIPPS, Major, USAF Director, Office of Information 1-23-79		
19. KEY WORDS (Continue on reverse side if necessary and identify by block number) Data Acquisition Signal Processing Fast Fourier Transform Microcomputer Fatigue Damage		
20. ABSTRACT (Continue on reverse side if necessary and identify by block number) The Air Force Flight Dynamics Laboratory at Wright-Patterson Air Force Base, Ohio, plans to design a constrained layer damping treatment for aircraft components. This design requires the knowledge of vibration and temperature variations that a component encounters in service. The development of a microprocessor based data acquisition and processing system (MIBDAPS) to monitor the vibration and temperature variations of a component under test is discussed in this paper. The designed system is capable of acquiring and processing data in real time and generating cumulative damage (cont)		

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factor versus frequency and temperature contours. These contours are used in the design of the damping treatment. The system uses an LSI-11M microcomputer as its central processor. The support hardware for this microcomputer includes an analog-to-digital converter, core memory, real time clock, and serial line interface module. The software developed for this system handles the acquisition and processing of data. The real time processing includes the computation of the FFT and power spectrum of the vibration signal and subsequent calculation of the cumulative damage factor versus frequency and temperature contours. The bench model of the designed system was implemented and tested for functional performance.

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