

AD-A069 543

OHIO STATE UNIV RESEARCH FOUNDATION COLUMBUS
A SIMULATION LANGUAGE FOR EVALUATING INFORMATION PROCESSING SYS--ETC(U)
SEP 78 T 6 DELUTIS

F/G 9/4

DAAG29-77-6-0203

UNCLASSIFIED

ARO-15356.1-A-M

NL

3
OF
AD
A069543



ARO 15356.1-A-M

RF Project 760792/710632
Final Report

LEVEL *125*

DDC
RECEIVED
JUN 7 1979
C

BEST AVAILABLE COPY

AD A 069543

**the
ohio
state
university**

research foundation

1314 kinnear road
columbus, ohio
43212

A SIMULATION LANGUAGE FOR EVALUATING
INFORMATION PROCESSING SYSTEMS FROM A DEMS
AND USER PERSPECTIVE: EXTENSIONS TO THE
INFORMATION PROCESSING SYSTEM SIMULATOR

Thomas G. DeLutis
Department of Computer and Information Sciences

For the Period
August 1, 1977 - January 31, 1978

U.S. ARMY RESEARCH OFFICE
P.O. Box 12211
Research Triangle Park, North Carolina 27709

Grant No. DAAG29-77-G-0203

DDC FILE COPY

This document has been approved
for public release and sale; its
distribution is unlimited.

March, 1979

29 v6 05 031

Unclassified

SECURITY CLASSIFICATION OF THIS PAGE(When Data Entered)

Block 20 (continued)

An integral part of the research was the modeling of a complex information system architecture to test the appropriateness of the information processing system simulator design facility (IPSS) as a design tool.

This document describes the investigative and modeling efforts and reports on the following conclusions:

1. IPSS is a useful tool for modeling in a topdown, modular manner complex information system architectures, and
2. The proposed new capabilities would significantly reduce modeling effort.

Unclassified

SECURITY CLASSIFICATION OF THIS PAGE(When Data Entered)

6 A SIMULATION LANGUAGE FOR EVALUATING INFORMATION PROCESSING SYSTEMS FROM A DBMS AND USER PERSPECTIVE: EXTENSIONS TO THE INFORMATION PROCESSING SYSTEM SIMULATOR.

9 FINAL rept. 1 Aug 77-31 Jan 78,

10 THOMAS G. DELUTIS, PH.D.

11 1 Sep 78

SEPTEMBER 1, 1978

12 270 p.

U. S. ARMY RESEARCH OFFICE

18 ARD

19 153 56.2-A-M

15 DAAG 29-77-G-0203

Accession For	
NTIS GMAI	<input checked="" type="checkbox"/>
DDC TAB	<input type="checkbox"/>
Unannounced	<input type="checkbox"/>
Justification	<input type="checkbox"/>
By _____	
Distribution/	
Availability Codes	
Dist	Avail and/or special
A	

THE OHIO STATE UNIVERSITY

APPROVED FOR PUBLIC RELEASE; DISTRIBUTION UNLIMITED.

THE FINDINGS IN THIS REPORT ARE NOT TO BE CONSTRUED AS AN OFFICIAL DEPARTMENT OF THE ARMY POSITION, UNLESS SO DESIGNATED BY OTHER AUTHORIZED DOCUMENTS.

79 06 05 031

267 360

alt

ACKNOWLEDGEMENTS

This research activity represents the cumulative efforts and support of a number of individuals. John Chandler and Joe Brownsmith spent countless hours analyzing SIDPERS and IDMS, synthesizing IPSS models and identifying and correcting program bugs in the IPSS. Without their dedicated support this research would not have been possible. Finally, a heartfelt thank you to Dr. Lawrence Rose for the hours he has devoted to the editing of this report.

The support and assistance provided by the U.S. Army Computer Systems Command personnel in support of the modeling experiments was especially appreciated. Mr. Alan Curry provided overall project guidance and Walt Kwiatek provided the liaison to the Command. Mr. Brownsmith and Chandler have asked that the following individuals also be recognized: Mr. Frank Hunter for assistance on IDMS and SIDPERS details; Mr. Milt Schul for PDP-11/70 assistance; Captain Jerry Jones for additional computer time; the personnel of the Systems Integration Directorate for technical operational support; and the personnel of the Technical Training and library Division for providing office space.

TABLE OF CONTENTS

	<u>Page</u>
Abstract	
Acknowledgements	ii
Table of Contents	iii
Table of Figures	iv
Tables of Contents: Appendices	v
Table of Figures: Appendices	vii
Table of Tables: Appendices	viii
1. Synopsis of the Research	1
1.1 Introduction to the Problem	1
1.2 Research Objectives	4
1.3 Results of the Research	5
2. The Information Processing System Simulator (IPSS)	10
3. Background to the Research	21
3.1 Methodologies for the Evaluation of User Goals	21
3.2 Methodologies for the Evaluation of Data Base Management Systems	25
4. Experiments	29
4.1 The Modeling Effort	30
4.2 SIDPERS/IDMS Models	31
4.3 Insights Gained from the Modeling	37
5. Future Extensions to the IPSS	40
6. Bibliography	44
Appendices	
Appendix A: Background to the Modeling Experiments	49
Appendix B: Summary of Model Statistics	99
Appendix C: Results of the Research Activities	188

TABLE OF FIGURES

<u>Figure Number</u>	<u>Figure Title</u>	<u>Page</u>
2-1	Hierarchical View of Information Processing Systems	11
2-2	IPSS Modeling Philosophy	12
2-3	Relation Between IPSS Methodology, Language and Execution Facility	14
2-4	Components of an IPSS Model (with DBS Extension)	15
4.2-1	Structure of Final SIDPERS/IPSS Model	33
4.2-2	IPSS SIDPERS/IPMS Models	34

TABLE OF CONTENTS: APPENDICES

		<u>Page</u>
Appendix A	Background to the Modeling Experiments	49
A1	The Problem Environment	49
A1.1	SIDPERS	50
A1.2	Prior Analysis Activities of SIDPERS	51
A1.3	IPSS Model of SIDPERS	54
A2	The Hardware/Software Environment	60
A2.1	The Host Machine	60
A2.2	The Backend Machine	63
A2.3	Integrated Database Management System (IDMS)	67
A3	Modeling Using the IPSS Design Facility	76
A3.1	A Functional View of the Modeling Design	76
A3.2	IPSS Model of the Host Computer	79
A3.3	IPSS Model of the Backend Computer	87
A3.4	IDMS Models	92
Appendix B	Summary of Model Statistics	99
B1	SIDPERS/IDMS Workload Model Statistics	99
B2	Host/SIDPERS Submodel Statistics	110
B3	IDMS/Backend Submodel Statistics	115
B4	IPSS Functional Model of the SIDPERS/IDMS	122

TABLE OF CONTENTS: APPENDICES (continued)

		<u>Page</u>
Appendix C	Results of the Research Activities	188
C1	IPSS Language Statements for Characterizing Data Base Systems	189
C1.1	The Data Base Structure Component	190
C1.2	The Data Base Access Component	194
C1.3	IPSS/DBS Oriented Built-in Facilities	195
C2	IPSS Performance Measures	205
C2.1	IPSS Operational Statistics	206
C2.2	IPSS Request Stream Statistics	210
C2.3	IPSS I/O Activity Statistics	213
C2.4	IPSS Queueing Statistics	217
C2.5	IPSS Utilization Statistics	225
C2.6	IPSS Wait Statistics	233
C2.7	IPSS Service Related Statistics	240
C2.8	IPSS Task & Activity Statistics	248
C3	IPSS Extensions to Characterize Computer Networks	250
C3.1	Multiple Facility Pools	251
C3.2	Changes to IPSS Simulation Nucleus	253

TABLE OF FIGURES: APPENDICES

<u>Figure Number</u>	<u>Figure Title</u>	<u>Page</u>
A1.3-1	Services Represented in the IPSS Submodel of SIDPERS Transaction Processing	55
A2-1	Experimental Version of VIABLE SIDPERS Processing Environment	59
A2.1-1	Modeled Hardware Configuration	61
A2.1-2	Modeled Host System Resident Software	62
A2.2-1	USACSC PDP 11/70 Configuration	65
A2.2-2	Modeled Data Base System Resident Software	66
A2.3-1	Modeled IDMS Schema for SIDPERS Data Base	68
A2.3-2	SIDPERS Data Diagram	70
A3.1-1	Relation of Submodel to Total Model	77
A3.1-2	Modeled Operating System Functions	80
A3.2-1	Assumed Host Hardware Configuration	82
A3.2-2	Relation of Host Machine Processes to SIDPERS Processes	84
A3.2-3	IPSS Realization of Host/SIDPERS Model	85
A3.3-1	Relation of Data Base Machine Processes to IDMS Processes	89
A3.3-2	IPSS Realization of Backend/IDMS Model	90
A3.4-1	Schematic of the Modeled IDMS	94
A3.4-2	Schematic of IPSS/DBS Model of IDMS/ SIDPERS Data Base	97
B1-1	SIDPERS/IDMS Workload Model Service Invocation Sequence	101
B3-1	USACSC PDP 11/70 Configuration	116
C1-1	IPSS System Architecture with DBMS Extensions	191
C3.1-1a	Basic Structure of IPSS Simulation Nucleus	252
C3.1-1b	Structure of Model Nucleus to Support Networking	252

TABLE OF TABLES: APPENDICES

<u>Table Number</u>	<u>Table Title</u>	<u>Page</u>
A1.2-1	The Results of the SIDPERS Interaction Test	52
A1.2-2	Average Entering Time	53
A2.3-1	Record Occurrence Characteristics	71
A2.3-2	Set Occurrence Characteristics	73
A2.3-3	Page Utilization Within the SIDPERS Areas	74
A2.3-4	SIDPERS IDMS DML Types	75
B1-1	Index to Modeled SIDPERS/IDMS Processes	101
C1-1	IPSS/DBS Facilities	192
C1-2	IPSS/DBS Built-in Procedures	196
C2.1-1	IPSS Operational Statistics	207
C2.1-2	IPSS Operational Statistics - Required Model Statements	208
C2.1-3	IPSS Operational Statistics - Defini- tions	209
C2.2-1	IPSS Request Stream Statistics	210
C2.2-2	IPSS Request Stream Statistics - Required Model Statements	211
C2.2-3	IPSS Request Stream Statistics - Definitions	212
C2.3	IPSS I/O Activity Statistics	214
C2.4-1	IPSS Queueing Statistics	218
C2.4-2	IPSS Queueing Statistics - Required Model Statements	220
C2.4-3	IPSS Queueing Statistics - Definitions	222
C2.5-1	IPSS Utilization Statistics	226
C2.5-2	IPSS Utilization Statistics - Required Model Statements	228
C2.5-3	IPSS Utilization Statistics - Defini- tions	230

TABLE OF TABLES: APPENDICES (continued)

<u>Table Number</u>	<u>Table Title</u>	<u>Page</u>
C2.6-1	IPSS Wait Statistics	234
C2.6-2	IPSS Wait Statistics - Required Model Statements	236
C2.6-3	IPSS Wait Statistics - Definitions	238
C2.7-1	IPSS Service Related Activities	242
C2.7-2	IPSS Service Related Activities - Required Model Statements	244
C2.7-3	IPSS Service Related Activities - Definitions	246
C2.8-1	IPSS Task and Activity Statistics	249

1.0 SYNOPSIS OF THE RESEARCH

1.1 INTRODUCTION TO THE PROBLEM

In the past quarter century, computer technology has had a profound impact upon the manner in which data is processed. With each advance in computer hardware technology has come associated advances in software capability both in respect to the processing of data into more meaningful packets and in the management of the computer hardware. A seemingly never ending cyclic phenomenon has developed with regard to computer based information processing systems. Each expansion in the processing capabilities only increases the appetite of the user community and consequently new demands are placed on the system designers and analysts to increase system throughput, reduce response time and provide more services while maintaining or reducing costs. This has mandated the better utilization of expensive computer hardware and software resources and associated personnel and thus, the development of complex system architectures incorporating sophisticated supervisors to handle the job scheduling, task management and resource allocation functions, multi-processors to increase throughput and integrated distributed data bases to meet the user's information needs.

Contemporary computer based information processing system architectures are reflective of the information systems they support. As the information needs become more sophisticated, so does the computer system. In part, this phenomenon is the result of new technologies which permit increased volumes of data to be incorporated into information systems at reduced operating cost and faster retrieval time per quantum of data stored. Unfortunately, the complexities of modern society have not only increased the enterprise's need for information but also the transiency of data's value. Furthermore, needed information is more integrative in nature in order to reflect the interdependencies of modern organizations. Ill-designed systems which do not meet their operational requirements are expensive in terms of the time, manpower and dollars lost in their design and implementation.

The design and analysis problem is a complex one, and its complexity is augmented by the costs associated with such systems. Furthermore, the analyst cannot, in most cases, rely on past experiences since each system must be tailored to the particular environment in which it will operate. The need exists, therefore, for modeling technologies which provide the analyst/designer with insights into the behavior of complex on-line, real-time information processing system. Additionally, such technologies must be attuned to the descriptive and evaluative needs of the information processing system so as to minimize the cost and effort for their use both with regard to synthesizing models and evaluating statistics resulting from their execution.

Ideally, models provide the designer and analyst with the ability to identify and characterize an information system's user, data base, software and hardware components and their logical, control and data interfaces. During the lifecycle of an information system, modeling activities provide a feedback capability to the designer and analyst which enhances their ability to attain and maintain performance objectives. By being representations of systems or subsystems, models serve three purposes -- as an aid to design by providing insight into the behavior of selected system components; as a predictor of performance changes that would result for alterations to existing systems; and as a guide to monitoring activities by identifying components to be monitored and statistics to be collected.

Discrete event simulation models can provide a number of benefits to system designers and analysts. They can also be expensive endeavors which remain unused. Among the important benefits to be achieved from this form of performance assessment are:

1. Insight gained into the behavior of the constituents comprising a system,
2. Identification of performance sensitive components,
3. Identification of system and component performance in response to user activity (e.g., volume and mix), and

4. Estimates of performance changes which would result from new hardware, software⁽¹⁾, and data base architectures and system procedures.

Simulation models serve as predictors of system activity; they are cost effective only when the costs associated with model synthesis, validation and evaluation are significantly lower than the cost of building benchmarks or prototype systems.

The modeling of contemporary systems requires the inclusion of submodels for data base, data base access software and support facility components. Without a systematic and modular approach to the total modeling effort, costs can become prohibitive. Therefore, the first step in achieving a modeling capability is the identification of a set of functional components common to the information processing system architectures to be investigated. A second requirement for a cost effective modeling capability is the availability of models which permit the analyses of alternative architectures and/or system activity with a minimum of modification. Third is the ability to incorporate new evaluative technologies and procedures into existing models. Fourth, simulation models must produce meaningful statistics from model evaluation; in particular, it is important that the statistics they generate are useful to the analyst and designer and that they can be validated. Finally, the performance assessment system should allow the reuse of component characterizations in order to effect technology transfer and thus reduce overall cost.

To be effective, simulation facilities must be compatible with the designer and analyst's evaluation goals. This implies the ability to perform a spectrum of modeling tasks from macroscopic analyses of preliminary designs to detailed assessments of instruction execution characteristics. Model synthesis must be consistent with the information known about the system and output data commensurate with the evaluation needs. Required is an integrated view of information system

⁽¹⁾ Software includes operating system modules, application programs and general data base access software.

activities and interfaces which permits a top down, modular approach to model synthesis and is capable of accommodating present and anticipated hardware and software architectures.

1.2 RESEARCH OBJECTIVES

The Information Processing System Simulator (IPSS) has been developed⁽²⁾ with four general objectives in mind:

1. Ease of model synthesis,
2. Modularity in model construction,
3. System self-containment, and
4. Portability.

Additionally, it is recognized that the system must maintain executional efficiency and must be usable by a wide spectrum of users, from the novice to the sophisticate. A continuing goal is to further the development of IPSS within the framework of the above objectives.

The purpose of this research was to investigate methods for extending the IPSS language and statistical facilities, focusing upon those facilities which could increase the definitional and evaluative capabilities of IPSS. The research has led to the incorporation of new language and statistics gathering features into the IPSS. These include (a) DBMS -- associated definitional statements and performance measures, and (b) user and task-oriented evaluative capabilities which provide:

1. Definitional statements to ascertain the performance of a model relative to user goals;
2. Definitional statements specific to network and hierarchical type DBMS architectures;
3. Enhanced standard statistical outputs to provide additional insights into model behavior; and
4. Interfaces to internal model facilities thereby permitting additional model control and statistics generation.

⁽²⁾ IPSS's development was sponsored in part through research grants GN-36622 and SIS-75-21648 from the National Science Foundation. (Office of Science Information Services) to The Ohio State University.

These new features permit modeling activities to be more closely identified with actual analysis and design functions. Additionally, they reduce model artificiality and provide more meaningful and identifiable statistics. The benefit to the modeler is a reduction in time and effort required to synthesize and validate models of complex information processing systems.

1.3 RESULTS OF THE RESEARCH

The research proposal identified modeling constructs and statistics as candidates for inclusion into the IPSS design facility. The vehicle for assessing the usefulness of these extensions was the modeling of a complex system architecture to be described in Section 3. Details of the evaluative efforts and the results established are found in Appendices A, B and C.

For evaluative purposes, the proposed IPSS extensions were placed in the following categories:

- A. MODEL STRUCTURE RELATED
 - 1. Data Base Subsystem Submodels
 - a. IPSS Data Base Access Component
 - b. IPSS Data Base Structure Component
 - 2. Network of Submodels as a Model
- B. LANGUAGE EXTENSION RELATED
 - 1. Data Base Subsystem Definition
 - 2. System Effectiveness Evaluation Definition
 - 3. Network Definition
- C. STATISTICS RELATED
 - 1. System Effectiveness Evaluation
 - 2. Data Base Subsystem Behavior
 - 3. General

In Appendix C the results of the evaluation are discussed. In all 37 statements and 123 statistics identified in the proposal were evaluated. The results of the evaluation are summarized in the following table.

Table 1.3-1. Summary of the Evaluation of Proposed IPSS Language Extensions

ASSESSMENT CONCLUSION	DISPOSITION CODE	TOTAL	
		Statements	Statistics
Not Examined	NE	6	20
Examined and Rejected	ER	0	0
Examined and Found to be Useful -			
Has been implemented as proposed	EU-I	16	0
Has been implemented but merged into another facility	EU-IM	6	7
Has been defined and not implemented	EU-D	0	93
Has been defined, used in paper models, but not implemented	EU-DP	9	0

The modeling activities also demonstrated the feasibility and suitability of applying IPSS and its underlying methodology to the synthesis of models of complex information systems. The system modeled to demonstrate these capabilities was a host/back-end processor configuration which supported interactive application processing and a data base management system. The modeled software processes included a detailed characterization of application loading, DBMS processing, and the operating system function of task management, job scheduling and resource allocation.

The IPSS methodology proved capable of providing the modeler with a modular, top-down approach to system analysis and model synthesis thus facilitating a solution to this complex problem. Several models were constructed, each of which reflected system details of a particular aspect of the system. The ease of model expandability in both breadth and depth was demonstrated through the construction of independent models and their

integration into an overall complete model. The SIDPERS/IDMS workload model represents the software processes of the applications and IDMS in detail while characterizing the operating systems and computer hardware as black boxes. From this basic model, two subsequent models were developed by splitting the software processes at the host/back-end interface and adding characterizations of the appropriate hardware and operating system software functions. The synthesis of the final overall model was easily accomplished since the system interfaces were clearly identified.

These models were independently tested, each producing a set of statistics reflective of the modeled processes and computing environment. In all cases queueing and utilization statistics were obtained for the modeled software and hardware resources. These statistics provide an easily understood synopsis of activities since one or more IPSS services were used to represent major units of processing (e.g., application programs, TP line scheduling, IDMS processing). Statistics were also automatically collected on each of the hardware facilities indicating both potential bottlenecks and under-utilized components.

The models were all internally verified. That is, through experimentation, the internal processing consistency was verified to be accurately reflective of the real world processes. Due to the lack of operational data, however, it was not possible to validate them through comparison to actual system performance. Thus, although some confidence can be placed in the relative values of the statistics reported in Appendix B no validation has been made with respect to the magnitude of the data values.

The experiment demonstrated the suitability of IPSS to model complex systems while requiring minimal time for model synthesis. The modeling effort required approximately four man months. The resulting models have proven to be versatile as well as adaptable to changing requirements indicating that significant portions can be used for other application areas. For example, the DMBS and operating system

processes as well as hardware facilities need not be modified when another major application area is modeled, thus resulting in a substantial reduction in future related modeling efforts.

In addition, this experiment has served to focus future development work as well as to validate current goals. In particular, the experiment demonstrated the need for simulation facilities specifically attuned to the characterization of integrated data structures as represented in IPSS/DBS. These initial models did not use these features since they were not fully operational, and thus the logical structure of the data base was characterized at a more abstract level than desired. Furthermore, experimentation with IDMS record mapping policies was not attempted because of the complexities involved in relation to the time available.

A number of experiments on the basic models became readily apparent. These were:

1. Identification of the maximum number of terminals which could be supported as the system approaches saturation;
2. Response time behavior for system processes and resources as loading increases and/or mix of transaction types is varied;
3. Effect on throughput as the transaction error rate is varied;
4. Response time behavior as a function of transaction arrival rate;
5. Effect on system throughput and response time with a m-way, multi-threaded DBMS;
6. Impact of presorting transactions on system throughput.

The initial model did not contain operating system or computer hardware characteristics. However, later models did have these characterizations. Thus the effect of changes in their characterizations on throughput and response time could also be examined in future analyses.

The modeling was done using the original IPSS, now designated IPSS/BASIC. The IPSS/DBS features were not employed. As a result detailed models of DBMS behavior could not be easily modeled. However, paper models using the proposed syntax were written comparatively quickly

(one man-month which included review of the IDMS literature), This exercise demonstrated the value of the IPSS/DBS capabilities to reduce the modeling effort needed to characterize DBMS software and data base scheme.

Finally, the incorporation of the IPSS/DBS into the current IPSS has extended the concept of information systems modeling into the area of networks of simulation models. Since the IPSS/DBS was designed to execute either as a stand alone or asynchronously with IPSS/BASIC models, it became necessary to devise a mechanism to accommodate this concept. This mechanism has been extended so that an arbitrary collection of nodes⁽³⁾ can be incorporated into an IPSS model. The implementation requirements for this concept are discussed in Appendix C.

⁽³⁾An IPSS model node is defined to be a uniquely identified submodel comprised of IPSS/BASIC, IPSS/DBS or IPSS/BASIC-IPSS/DBS components.

2.0 THE INFORMATION PROCESSING SYSTEM SIMULATOR (IPSS)

The Information Processing System Simulator (IPSS) is a special purpose digital simulation facility designed to facilitate the performance assessment of complex computer based information processing systems. It represents the realization of a generalized methodology for the performance assessment of information systems and has been designed to be applicable to every phase of the information processing system life cycle. Additionally, models developed for one phase can be easily and inexpensively modified in both scope and depth of detail. This enables the model to evolve in parallel with the system's progression through its lifecycle.

The methodology employed by IPSS divides the elements of an information processing system into three categories according to the function they serve: data bases, services which access the data bases, and support facilities. As illustrated in Figure 2-1, a system is viewed as a hierarchy of services and data bases with each level supporting the service and data needs of the next higher level and requesting support from the adjacent lower level. Performance measures identify the behavior of services relative to their data base access activity, acquisition and use of support facilities, and use by other services. The methodology emphasizes data base access activity and has identified those system elements at each data base level which contribute to this activity. Within IPSS, the characterization of system elements has been formulated in a manner such that their interaction during model evaluation automatically produces the desired performance measures. An illustration of the modeling process is shown in Figure 2-2.

IPSS is more than a language in the sense of a GPSS or a SIMSCRIPT; it is a complete system. In addition to statements defining resources and tasks for the system under investigation, IPSS contains special statements directing the IPSS system in a number of auxiliary functions including the use and management of user defined library facilities, model compilation, and model execution sequencing.

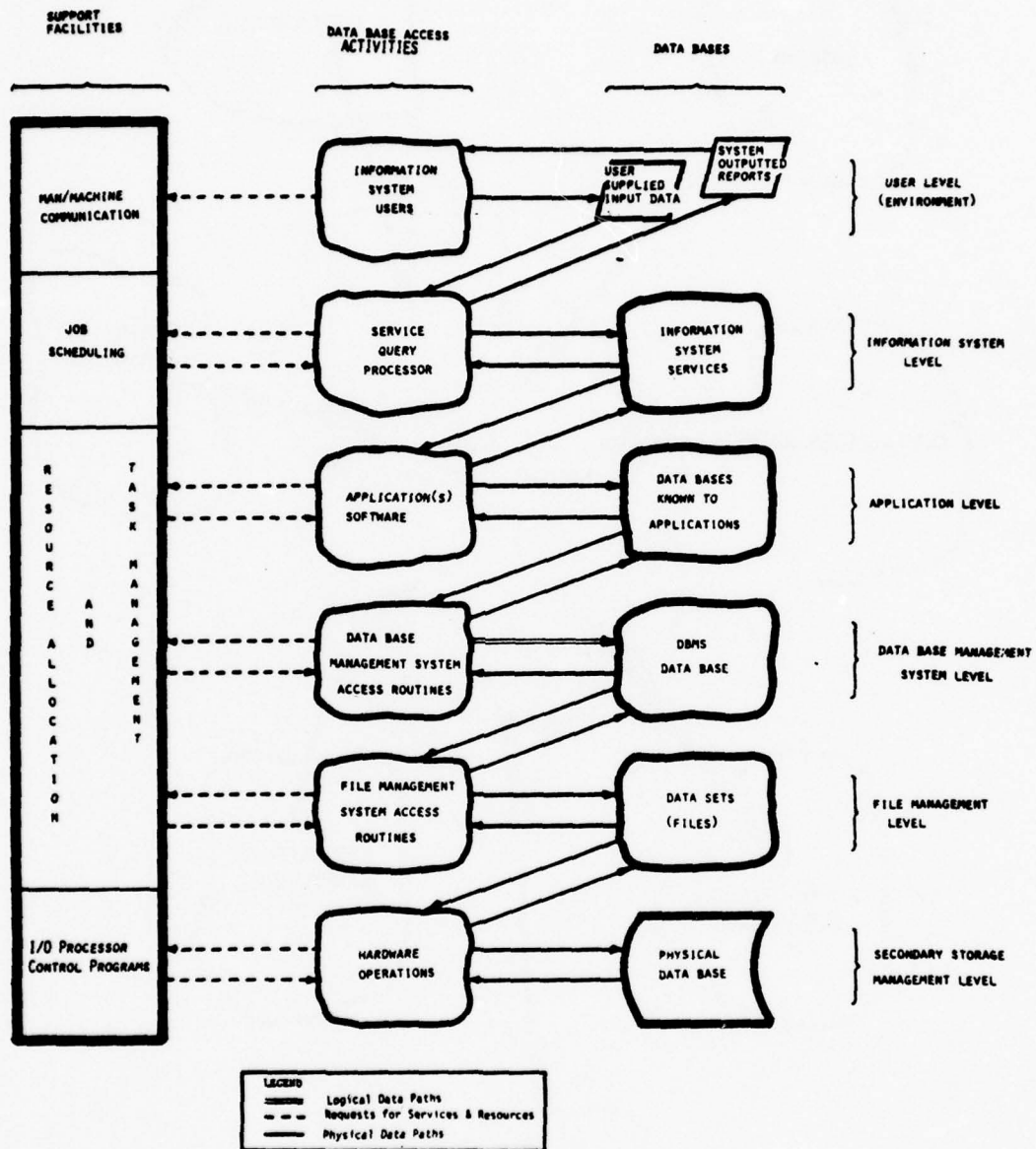


Figure 2-1 Hierarchical View of Information Processing Systems

IPSS MODELING PHILOSOPHY

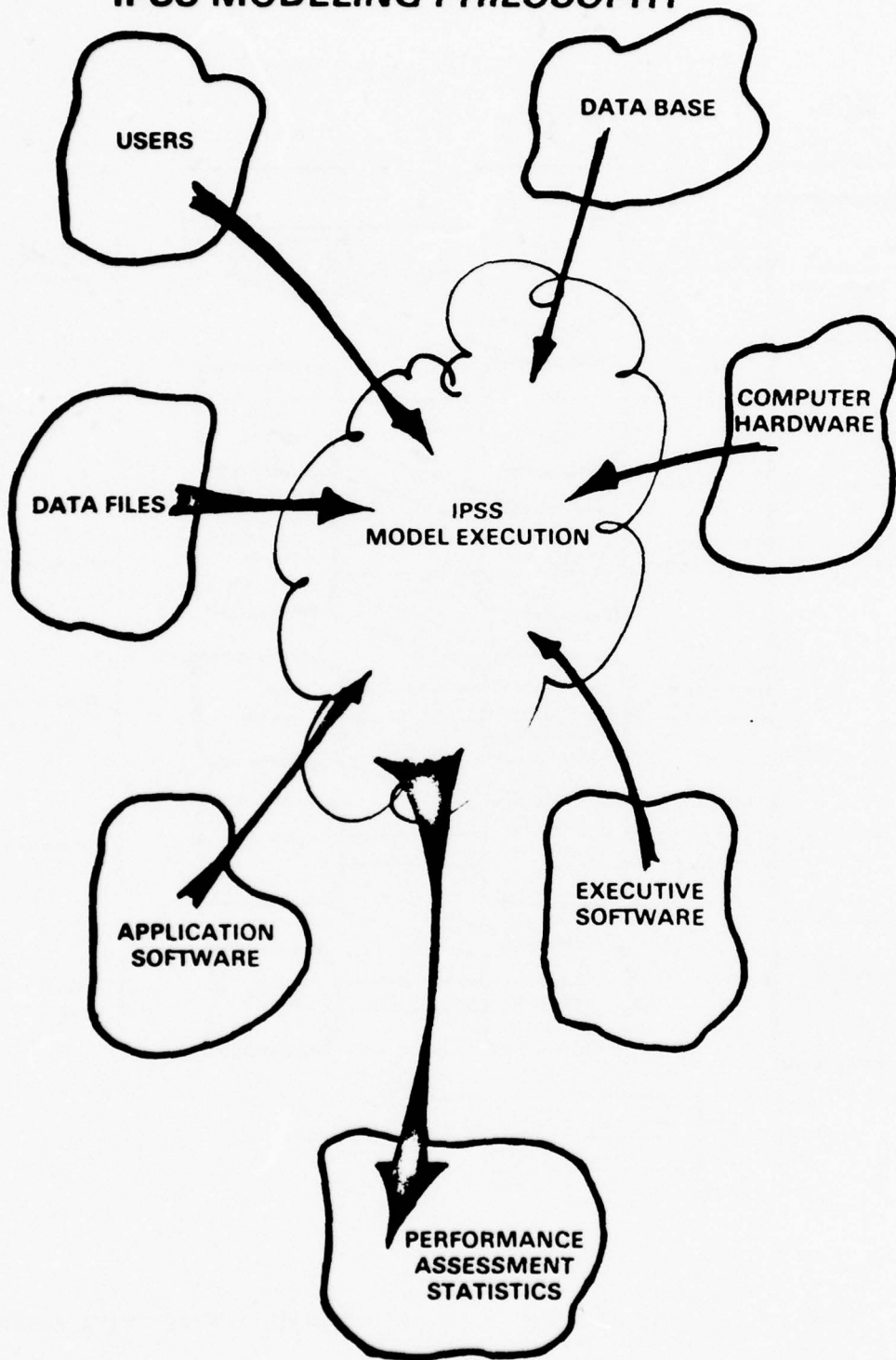


Figure 2-2 IPSS Modeling Philosophy

The IPSS language facilities have been carefully designed to focus the modeler's synthesis activities on the IPSS methodological underpinnings. In this manner the development of models of complex systems can be done in a systematic, top down manner. This affords the modeler the same advantages that these techniques provide in information system analysis, design and implementation. Figure 2-3 illustrates the three phases in the use of IPSS for IPSS model synthesis. At the beginning all that is available to the modeler⁽⁴⁾ is general knowledge of the information system (at some stage in its life cycle). Phase I is an implicit one during which the modeler employs the IPSS methodology to separate overall system knowledge into specific knowledge of the following four functional components:

- A. System users and their request characteristics;
- B. Services to be provided to users and to other services, and their inter service connections;
- C. Data base resources and configurations;
- D. Hardware resources and configurations.

The IPSS language has been specifically formulated to characterize these functional components.

During Phase II the above functional components are characterized via the IPSS language into independently defined model components. Presently five distinct model components are possible to describe an information system as shown in Figure 2-4. (The sixth component, the model director, is used to control simulation activities during Phase III, the model execution phase.) Not all are required to model a particular computer facility. Additionally, more than one of each component (except model director) can appear in an IPSS model. A synopsis of the role of each component follows:

1. SYSTEM RESOURCES -- contains definitions for all hardware, software and data resources comprising an information system model. Included in the SYSTEM RESOURCES component is the IPSS supplied clockwork mechanism to schedule and control simulated events and to determine

⁽⁴⁾A modeler is assumed to be a systems/analyst or a systems designer.

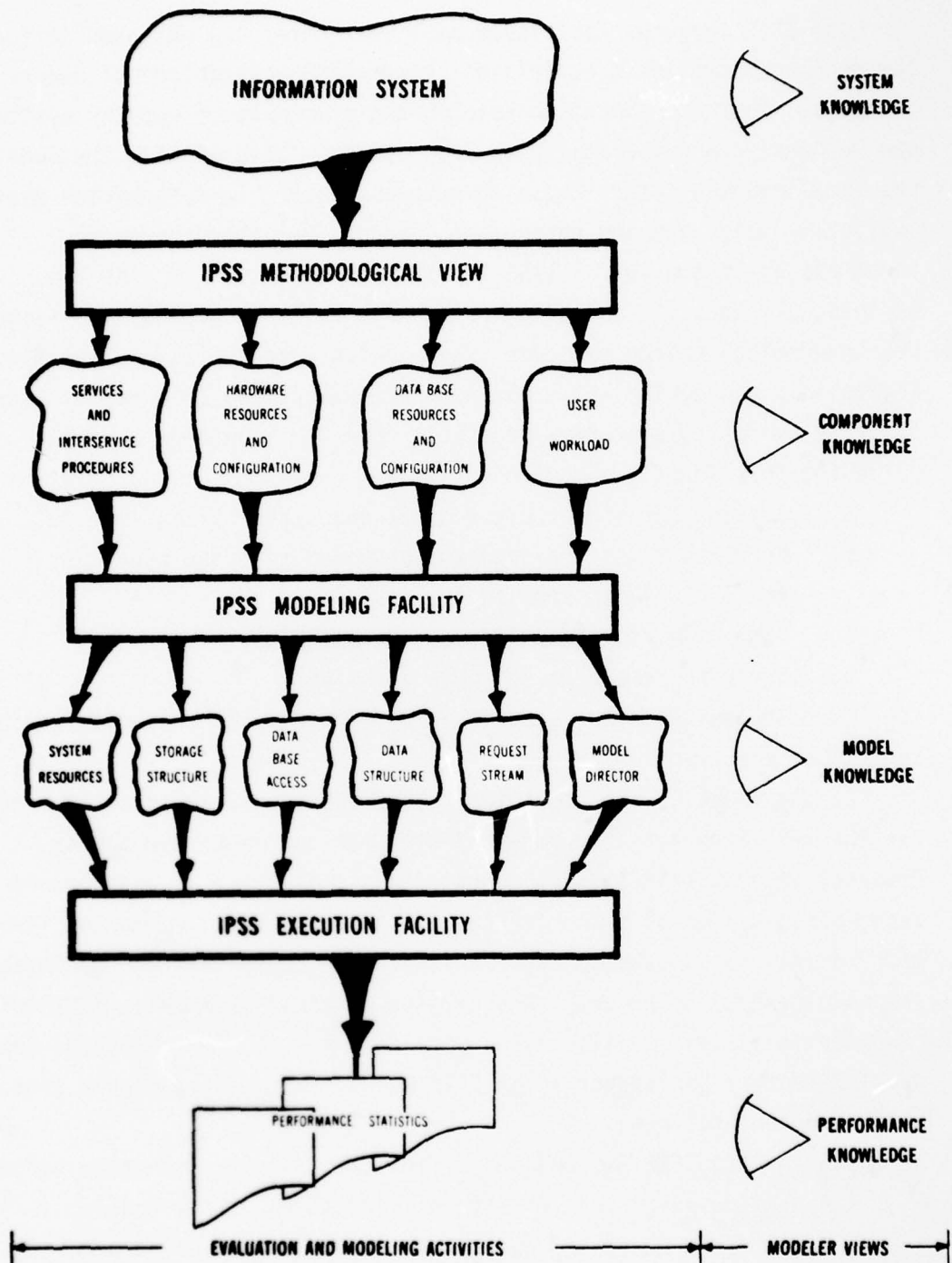


Figure 2-3 Relation Between IPSS Methodology, Language and Execution Facility

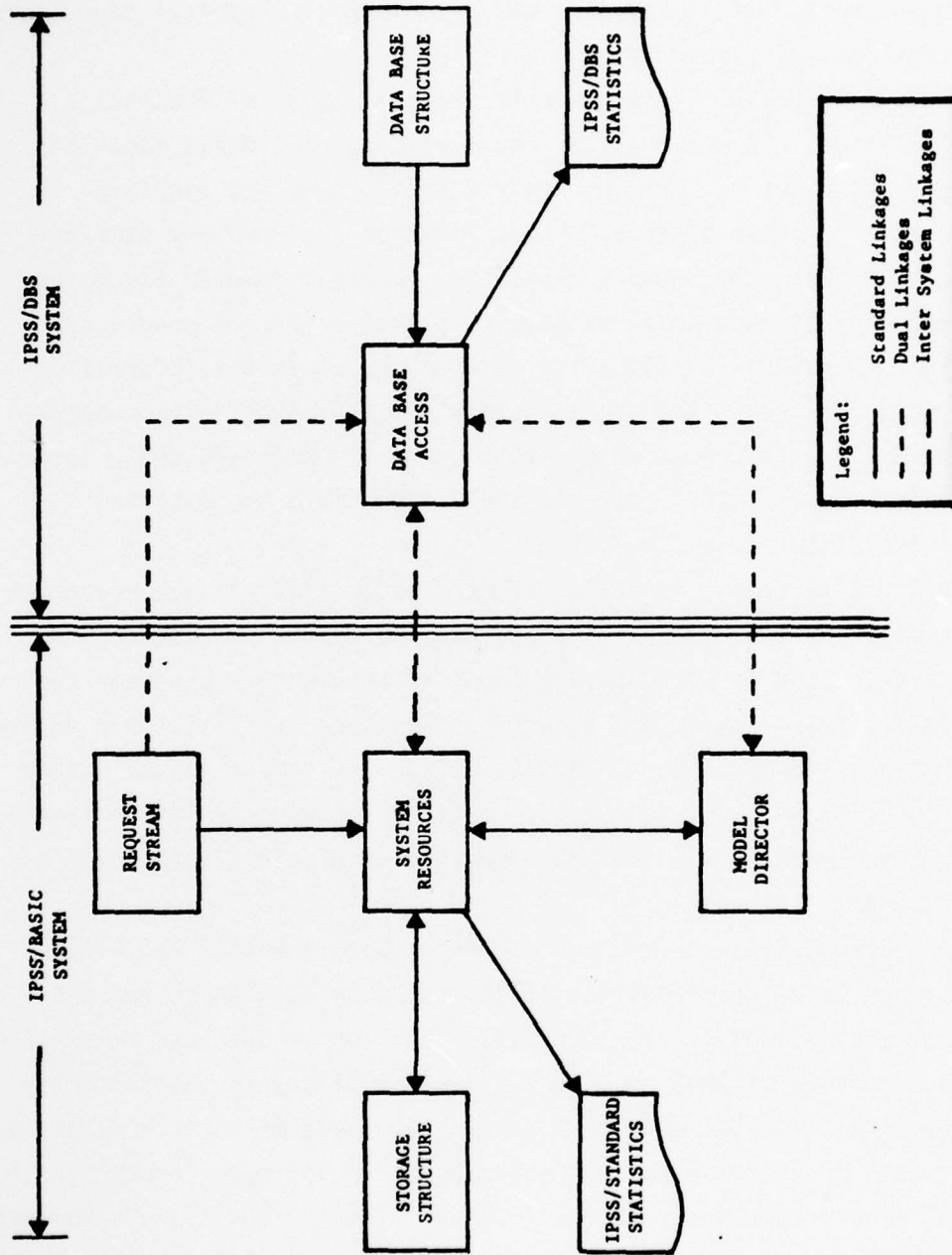


Figure 2-4 Components of an IPSS Model (with DBS Extension)

when the simulation is to terminate. The clockwork logic is based on the next most immediate event philosophy for controlling discrete event digital simulations;

2. STORAGE STRUCTURE -- supplements the definition of data set facilities. It provides the model with general definitions for data set files in terms of their formats, size and location relative to each other and their location in secondary storage. Additionally, the STORAGE STRUCTURE component is used to define secondary storage space management policies for the model;
3. REQUEST STREAM -- defines the external requests for information system services; that is, the users of the system being modeled. The modeler is required to define types of users and their inter-arrival times. IPSS converts these times into a composite arrival time stream;
4. DATA BASE ACCESS -- contains the definitions of all the resources required by the DBMS. These include the hardware resources of buffers and user work areas as well as application programs and DBMS software. All DBMS-related entity type facilities are defined within the component. DATA BASE ACCESS is similar to the SYSTEM RESOURCES component in that it contains its own simulation clockwork mechanism similar in purpose to that of the REQUEST STREAM component.
5. DATA STRUCTURE -- provides the modeler with a set of facilities which allows the definition of logical data structures and the characterization of relationships among them. This can be applied to a variety of DBMS architectures and application environments. The DATA STRUCTURE component permits the modeler to investigate the effects on system behavior caused by alternate set, record type, and access path definitions. The definitional facilities provided enable the modeler to investigate a wide spectrum of logical data structure organizations and allocation policies.

Also illustrated in Figure 2-4 is the division of the IPSS components into two subsystems -- IPSS/BASIC and IPSS/DBS⁽⁵⁾. This has been done to permit a modeler to concentrate his efforts on a particular aspect of the problem. Thus, the capability to model only the data base or the system aspects of a total system is provided the modeler.

Associated with each component are IPSS language facilities which permit the modeler to define and characterize information system entities and attributes, gather statistics and control model behavior. Over 100 such statements are available to the modeler.

The function of Phase III is to assemble the model, initiate its execution and display the generated performance assessment statistics. IPSS provides a modeler with a number of statistics concerning the behavior of modeler defined entities and IPSS supplied built-in information system services.⁽⁶⁾ These statistics fall into eight general categories:

1. Operational Statistics,
2. Request Stream Statistics,
3. I/O Activity,
4. Queueing Statistics,
5. Utilization Statistics,
6. Wait Statistics,
7. Service Statistics, and
8. Task/Activity Statistics.

Additionally, the modeler can employ the complete facilities of the Fortran language to develop his own statistics. Statistics are printed automatically at the conclusion of each model simulation unless explicitly inhibited.

⁽⁵⁾ This configuration of IPSS components is currently being implemented. The CSC model reported here employed only IPSS/BASIC components. However, the IDMS was also written using IPSS/DBS facilities but has yet to be validated.

⁽⁶⁾ These services are characterizations of the basic I/O operations for tape, DASD, unit record, and terminal devices.

Finally, IPSS contains a library facility which permits the modeler to retain previously defined models, model components and code segments for later referral. This feature can greatly reduce modeling time.

An implicit goal of the IPSS is to achieve the maximum applicability of the methodology and the IPSS facility. Therefore, the following criteria have served as guidelines throughout its development.

1. Breadth of Applicability -- the ability to model the behavior of contemporary and foreseeable system architectures and operating environment;
2. Functional View of Systems -- the ability to identify and characterize system components and activities based on their function independently of a particular architecture or environment;
3. Top Down, Modular Model Synthesis -- the ability to model to a level of detail commensurate with the desired analysis of research objectives;
4. Extensible Structure -- the capability to incorporate new, higher level descriptive facilities and performance measures into the methodology and simulation system; and
5. Flexibility of Use -- the ability to be used by a wide spectrum of modelers from the experienced system analyst/designer and researcher to the practitioner and student.

Examples illustrating the broad flexibility and applicability of the extended IPSS system follow.

Modeling Activities

Modeling activities are those activities concerned with the development of evaluative capabilities for examining the behavior of existing and proposed software, hardware and data base architectures.

1. Models of specific data base management systems for the purpose of examining their behavior under a variety of different criteria such as loading, data structure, and data mapping.

2. Models for evaluating the influence of data structures on system performance over a broad spectrum of DBMS loadings and support architectures.
3. Models of systems where the central site computing facilities are comprised of computer networks in order to examine the effects of alternative scheduling and tasking philosophies.
4. Models for evaluating the interaction of front-end and back-end computing capabilities.
5. Models for investigating the effects of queue scheduling policies for basic system activities such as job scheduling, task management, resource allocation, and paging strategies for virtual memory systems.
6. Models to evaluate data base security procedures; in particular, the identification of authorized users to the system and the impact of such activities on overall performance.
7. Development of a simulation capability and support library which permits the modeling of a wide range of contemporary architectures for the purpose of evaluating new technologies as they become available.

Performance Measurement Research

Performance measures are those measures which assist designers and analysts in their performance evaluation activities.

1. Measures of system performance based on user oriented goals.
2. Measures of the impact of new applications on current systems.
3. Development of higher level measures which permit the identification of critical system activities based on applications performance objectives.
4. Development of measures and procedures which aid in identifying (from a set of alternatives) the "best" alternative to achieve system performance objectives.

5. Development of measures to assist the data base administrator. (It is assumed that data base administration is an on-going activity and its major function is to maintain and assure the integrity of the data base and to improve its performance via changes in its data structure.)

These activities were possible due to the unique design of IPSS which permits a top down modular view of information processing systems and the special language constructs and facilities developed specifically for the above purposes.

The IPSS system is designed to analyze information processing systems in their totality. Initially, IPSS mechanism for the definition and evaluation of hardware, computer system software and storage level data structures were developed and have been extensively tested. The next phase in the continued development of IPSS was to incorporate features for the description, manipulation and evaluation of logical data structures of the kind normally encountered in modern data base management systems. While the current IPSS capabilities substantially reduce the effort required to model system and data base level software, hardware and dataware, modelers can model all aspects of an information processing system. A high level query processing capability is currently being implemented and development is planned for a telecommunication specific extension.

3.0 BACKGROUND OF THE RESEARCH NEEDS

The main goals of the research are the extension of the current IPSS system to include language constructs and statistics attuned to data base management system modeling activities, and to the measurement of system performance from a user perspective. The following discussion provides a background to evaluative methodologies associated with the research needs for these areas.

3.1 METHODOLOGIES FOR THE EVALUATION OF USER GOALS

As the complexity of current and anticipated systems increases, the analyst's need for more sophisticated evaluative techniques becomes more apparent. The types and number of services being provided by an information system are growing rapidly. Where once was a homogenous set of similiar user classes which could be evaluated by relatively simple measures, now operates a multitude of user classes and services demanding various conflicting performance requirements and lacking a single common measure of performance. Thus, from an external performance point of view, the analyst must evaluate information systems against a multiple set of user criteria.

From an internal performance view, current and anticipated demands force the analyst to also evaluate multiple criteria. As the number of services provided increases, the size of the delivery system must also increase. This implies that the number of resources required to build the system will increase. Information system resources -- hardware, software and dataware are expensive to buy, install and operate. Thus, the analyst must also be concerned with the cost effective use of these resources.

Current evaluative technologies focus on only one criteria in the system design equation, either the use or the system's resource performance. The ability to simultaneously ascertain the impact of resource performance on user goals or vice versa is not readily achievable through these methodologies. One of the purposes of this evaluative

expansion is to establish a formal liaison between the evaluation of user goals as a function of system behavior and the analysis of resource performance as a function of user activity.

User oriented analyses with objective functions based on response time, throughput, and cost have been (are continuing to be) reported in the literature. Most frequently, analytic approaches use queueing models as their basis (GAV76, BUZ73, KLE76 are representative of this type of analysis). Due to the necessity to maintain tractable models, many simplifications are required for a model's analytical solution. Simulation models have also been applied to user oriented analysis (CON65, ROE70). Unfortunately, these models yield only average and/or aggregate measures of system response. As a result of these simplifications, the analyses produced by both of the approaches fails in many cases to identify the relationship between users and resources. Therefore, they are suspect when used to predict the impact on system performances of modifying the current environment.

Alternatively, performance analyses can be made from the system's standpoint, treating the user and his goals in the aggregate. The most common approach is a subsystem study, where a particular part of the information system complex is isolated, with the subsystem user(s) represented by a stochastic generator, both analytic and simulative. The most emphasized areas of research has been the I/O subsystem (ABA68, NAH73, BOU72) and CPU utilization (KLE72, HEL70, MLS76). The problem with this level of evaluation is that, although providing valuable local intuitive insight, these models rarely relate to the ultimate information system user, and, therefore, do not provide realistic insight into global performance.

Examinations of complete systems have also been made. Exhaustive hardware/software measurements have been analyzed by Lindsay and Cole (LIN76, COL72) while simulation models, including an aggregate user component, have been built by Reeves and Pooch, Norland, and Lum (REE75, NOR71, LUM70). Although results of the evaluations include resource utilization statistics and user oriented measures such as response time, there is little attempt in these models to relate particular resource

usage to the effort on user goal attainment. But from practical experience it is evident that there is indeed a relationship between user goals and resource usage. In fact, Buzen (BUZ76) has recently proposed some fundamental laws for computer performance which relate resource activity to global system/user measures such as response time and throughput.

Thus, current evaluative efforts have not been able to provide an acceptable approach to handling the multiple criteria environment. Because of the inherent conflicting nature of the criteria an optimal design is not possible, and therefore, good system design is assumed to satisfy both performance related criteria. The approach embodied in this evaluative expansion was to first establish a causal relationship (liaison) between user goal attainment and system resource activity. Then after measuring the resultant behavior, a multi-objective technique was employed to produce a satisfied evaluation. The liaison will be specifically maintained in IPSS via the Task and Activity facilities. A mathematical programming technique, multiple goal programming (MGP) has been adapted to evaluate the multiple performance criteria (CHN77).

The importance of providing a task modeling capability is evident by examining current multi-programming system design philosophies. There has been a continuing trend toward hierararchic design of operating systems (DIJ68, ZUR68, MAD76, SDB77), both top-down and bottom-up. Such designs imply that particular functions are assigned to particular levels within the hierarchy, e.g., job scheduling at one level, hardware I/O control at another, etc. Each of the functions (levels) also in a sense, manage a class of resources via allocation, usage and de-allocation. In order to model and evaluate the effect of user demand on each of these functional levels, and thus, on their subordinate resources, one must be able to model the tasking function.

The overall modeling philosophy of IPSS, coupled with the addition of the TASK and ACTIVITY facilities allow the modeler to accomplish this. The purpose of the TASK facility is to maintain control over the resultant resource activity with respect to its causal user demand. This resource activity is further broken down by activities which characterize the individual system functions.

The control that the TASK facility exerts over its subordinate Activities and resources is essentially its maintenance of statistics. In IPSS the modeling resource usage automatically produces an extensive set of queueing and utilization statistics. The characterization of software processing also produces, automatically, service by service statistics on service time and hierarchy interactions. The TASK facility is associated with a particular user class to complete the user-resource liaison. It is the function of the TASK facility to maintain a set of statistics measuring the behavior of its subordinate resources. These statistics, collected dynamically, essentially partition of the aggregate queueing, utilization and service time statistics to each causal TASK (user class). Furthermore, the collection and maintenance of these statistics is performed automatically by IPSS. Thus, the first part of this approach to information systems establishes the link from user demand to resource activity and automatically produces statistics.

The second part of the approach to be implemented in this expansion is an evaluation of the multiple performance criteria. To be most effective, this evaluation should relate resource activity to user oriented performance goals. This entails the production of higher level performance measures, that may not be solely time related such as cost. This completes the two-way liaison between user demand and resource activity by measuring the contributions of resources to the satisfaction of user goals. This is, in part, accomplished via the TASK/ACTIVITY liaison, but the evaluation, however, is performed by a multiple goal programming based procedure.

The evaluative view of an information system formulated by the MGP procedure is identical to the modeling view characterized by the TASK and ACTIVITY facilities whereas using the TASK and ACTIVITY aids in the creation of a top-down modular performance evaluation (PE) model. In fact, the variables of the model are the TASK and ACTIVITIES. This analogy was designed to allow the automatic collection of statistics by the TASK facility during normal modeling be directly applicable to the PE model for evaluation.

The result of evaluating this PE model is two-fold. First, the desired evaluation of the user criteria is produced, on a global goal basis and on an individual goal basis (Task basis). Second, each activity is also evaluated relative to its service time performance, across all tasks (user classes). This evaluation can determine if a given activity is being under-utilized, saturated or used to its proper capacity, where proper is with respect to all tasks. These activity measures can be used as guidelines when deciding which designs/alternatives to follow. The most probable performance inhibiting activities can be identified in this way. An overview of the underlying methodology of this multiple criteria evaluation can be found in the paper by Chandler and DeLutis (CHN77).

3.2 METHODOLOGIES FOR THE EVALUATION OF DATA BASE MANAGEMENT SYSTEMS

Sable (SAB71) outlines a specific evaluation procedure for data base systems which is based on establishing the cost effectiveness of each candidate system with respect to a specific problem. He proposes a multistage method for "scoring" or grading DBMS features. The scores are normalized according to an assigned nominal requirement for each feature and multiplied by a weighting factor which represents relative importance. The overall system score is the sum of the weighted results. This methodology relies heavily upon the intuition of the analyst to select the important factors and to assign correct weightings. In addition, it presupposes that the evaluation data are available.

While methodologies based upon this approach are useful for establishing guidelines and controlling the overall evaluation process, they are static in nature and therefore are not useful in measuring dynamic DBMS behavior. However, the current DBMS literature contains several methodologies specifically designed for the measurement of this type of behavior. These methodologies all provide a set of descriptive facilities for the definition of both DBMS software and computer hardware. In addition, they all provide a special purpose, discrete event simulation tool for the analyst/designer.

Phase II (OWE70) is one such tool for modeling hierarchical base management systems. Its objective is to allow the modeler to study the interaction of such DBMS variables as storage structure characteristics, data distribution, device interaction, and access methods. The modeler specifies the characteristics of the data, the data accessing strategies, and the underlying hardware configuration through a special purpose user-oriented language. Data field contents are characterized by distribution functions and data fields are related to each other to form what is called segment (the logical structure of the data). The logical data structure is then mapped to a physical storage structure by partitioning the records into files which exist on hardware devices. The accessing strategy allows the modeler to specify the lists of records which are to be accessed, the accessing method (i.e., sequential, indexed, or direct), and the order in which the accesses are to occur. Phase II is limited to the I/O analysis of single level hierarchical data base management systems. Its design objectives are intentionally modest and thus it is not suitable for modelling more complex software interactions.

A methodology for evaluating data base scheme, physical data organization, and data base loading is described by Hulten and Soderlund (HUL76). Their simulator is called ARTS/DB (Analysis of Real Time Systems/Data Base oriented systems). It is implemented in Simula and executes on the DEC10 system. The descriptive facilities which they provide allow the characterization of data base systems in terms of CODASYL-DBTG features. An ARTS/DB model consists of four levels, i.e., definitions of the application, the DBMS, the system software, and the hardware. This provides the analyst with a useful mechanism for performing successive experiments with only slight modifications to the model.

A data base simulator based on the Data-Independent Accessing Model I (DIAM I) (SEN72) has been developed with two objectives in mind: the comparison and selection of existing data base systems and the design and evaluation of new data base management techniques

(SCH76). This simulator consists of four subsystems similar to the ARTS/DB approach. In this case the subsystems correspond to the levels of the DIAM and are: a query generator, a query translator which uses defined access paths, and I/O generator, and the host computer system.

A three stage DBMS model building approach as well as a simulator system utilizing this approach has been described by Reiter and Finkel (REI75, REI76). First, high level user-supplied logical data view descriptions are translated into an intermediate standard form. Second, an instruction stream is generated which refers to data by block address. In this stage no assumptions are made concerning the location of data in physical storage. Third, the dynamic elements of the system are employed to produce performance measures. The particular DBMS under study is parameterized with respect to data representation, resulting in a specific model of that DBMS. The model is constructed of static and dynamic parts which results in a similarity to operational DBMS and also in executional efficiencies.

DeLutis (DEL72) presents a methodology for the analysis of auxiliary storage systems (AUXSIM). In this system the modeler describes hardware devices and software processes using language statements which are highly descriptive of these devices and processes. The AUXSIM translator produces Fortran source code which executes as a discrete event simulation of the system. This research was the predecessor to the Information Processing System Simulator (IPSS). The research proposed here was, in part, a further extension of the IPSS methodology to data base management systems.

In addition to these generalized evaluation techniques, several specific DBMS models have been developed using a variety of simulation languages and techniques. The primary evaluation criteria used by these simulations is response time and throughput.

Ghosh and Tuel (GH076) designed an experiment to model a large IMS data base to determine the effect of varying such DBMS features as the retrieval sequence and logical data structures. They collected operational data by executing IMS in a stand-alone environment. Their measurement criteria was the elapsed time required to execute a

sequence of calls. Griffith (GRI75) constructed a GPSS simulation model of Univac's DMS-1100 also using response time as the evaluation criteria. The simulation results reported by Nakamura (NAK75) include response time, resource utilization (cpu, channel, disk), as well as statistics related to transactions and queues. Hall (HAL74) describes a DBMS simulation model of a hierarchical data base whose objective was to minimize response time and costs to the user.

4.0 EXPERIMENTS

A series of modeling experiments were performed between mid June and mid September 1977 for the purpose of; (a) assessing the ease of employing the IPSS methodology and language in modeling complex systems, and (b) evaluating the usefulness of the proposed IPSS language and statistics extensions. The system modeled was an on-line real-time system comprised of; (a) data entry terminals; (b) a host main frame for the application software and user interface; and (c) a CODASYL-like data base residing on a backend mini computer. Section 4.1 summarizes the modeling activities. (The reader may refer to Appendix A for a complete discussion of each of these models.) The following subsection provides background as to why such a complex system was chosen.

The purpose of the experiment, initiated by the U.S. Army Computer Systems Command, was to evaluate a complex information system architecture via the Information Processing System Simulator (IPSS). The primary objectives of the experiments were three-fold:

1. To demonstrate the capability of IPSS and its underlying methodology to accurately model a complex system,
2. To ascertain its ease of use, and
3. To develop a generalized basic model of the following systems:
 - a. an interactive version of the SIDPERS Personnel system resident in an idealized Host computer, and
 - b. The IDMS data management system resident in a PDP-11/70 serving as the back-end computer;

and through the model judge the appropriateness of the proposed IPSS language and statistics extensions.

Results indicate that IPSS is well suited for the task. Through the use of a top down modeling philosophy a complete model of the application and IDMS processing was written and tested in less than two man months and model statistics tracked well to the available performance measures. The following subsection provides an overview of the salient features of the modeling effort.

4.1 THE MODELING EFFORT

The purpose of the modeling activities was to demonstrate the feasibility and suitability of the IPSS to model large, complex information processing systems. Reflection on the modeling activities associated with the SIDPERS/IDMS experimentation provides positive evidence that the IPSS is equal to the task. One important reason that IPSS achieves its goal is that its underlying methodology provides a top down, modular approach to system analysis and model synthesis thus facilitating answers to complex analysis and design questions.

Several models were constructed, each of which reflected system details suited to the multi-stage development process which was adopted. The ease of model expandability in both breadth and depth was demonstrated through the construction of independent models and their integration into an overall complete model. Appendix A summarized their major features. The SIDPERS/IDMS workload model represents the software processes of the applications and the IDMS in detail while characterizing the operating systems and computer hardware as a block box. After this basic model was tested, subsequent models were developed by splitting the software processes at the host/back-end interface and adding characterizations of the appropriate hardware and operating system software functions. The synthesis of the final overall model was easily accomplished since the system interfaces were clearly identified.

These models were independently tested, each producing a set of statistics reflective of the modeled processes and computing environment. These results are summarized in Section B.2. In all cases queueing and utilization statistics were obtained for the modeled software and hardware resources. These statistics provide an easily understood synopsis of activities since one or more IPSS services were used to represent major units of processing (e.g., application programs, TP line scheduling, IDMS processing). Statistics were also automatically collected on each of the hardware facilities indicated.

The modeling activities were performed with approximately six man months of effort. All the models to be described in the next section were completed. However, three conditions occurred which prohibited the creation of the sought-after final models. These conditions were:

1. The IPSS/DBS parsers were not completed in time for use. However, models of the data base reflecting the IDMS scheme have been written in the IPSS/DBS syntax and semantics specifications. Furthermore, these submodels have now been executed under the IPSS/DBS.
2. Combination of all the submodels (excluding the IDMS schema) could not be successfully compiled. Investigations into the cause of the problem has lead to the detection of an error in the IPSS nucleus. The nucleus did not detect an IPSS library manager message signaling the lack of library space. Thus model code was lost and was undetected until final model assembly. This condition has been recently fixed and the models have been successfully executed.
3. The interfacing mechanism discussed in Appendix C3 has not been completely implemented. This is due to the unexpected quantity of code which had to be modified. Therefore, the IPSS and IPSS/DBS cannot currently execute together. This problem is currently being alleviated.

Section 4.2 discusses the complete modeling program envisioned for the modeling experiment. Appendix B discusses the statistics for the submodels actually completed during the experiment.

4.2 SIDPERS/IDMS MODELS

The purpose of this section is to provide an overview of the system analysis and model design activities related to the IPSS modeling of the interactive SIDPERS/IDMS system.⁽⁷⁾ For analysis purposes the system was divided in the following logical model components:

⁽⁷⁾The IPSS methodology provides a modular, top-down approach to system analysis and model synthesis thus permitting the solution discussed in this section.

Hardware

Host machine configuration
Back-end machine configuration

Software

Operating system for the host machine
Operating system for the back-end machine
Terminal interaction modules
SIDPERS transactions
Interface modules
IDMS

These logical model components were used to define five inter-connected system activities each representing specific SIDPERS/IDMS functions. This organization is shown in Figure 4.2-1. A number of models were written which incorporated a one or more of the modeling components and represented one or more system activities. Figure 4.2-2 identifies the submodels and incorporated model components.

Because of the magnitude and complexity of the system, a phased approach was taken in model development. The model synthesis and validation is implicit in IPSS, thus, the approach posed no difficulties. In Phase I a model of the SIDPERS transactions as processed by IDMS (system activities 1 and 5) was constructed. However, it is important to note that the entire structure of the final model was represented in this preliminary model. This was accomplished through IPSS services for system activities 2, 3, and 4 which were essentially null routines. That is, they were invoked by other services but their internal processing function remained unspecified. In particular, terminal interaction and secondary storage accesses were characterized merely by advancing the simulator clock. This simplified approach allowed the SIDPERS transactions and the IDMS processing functions to be modeled and validated quickly. Once this workload flow model was executing successfully and generating statistics, the operating system and computer hardware model components were added in succeeding modeling phases.

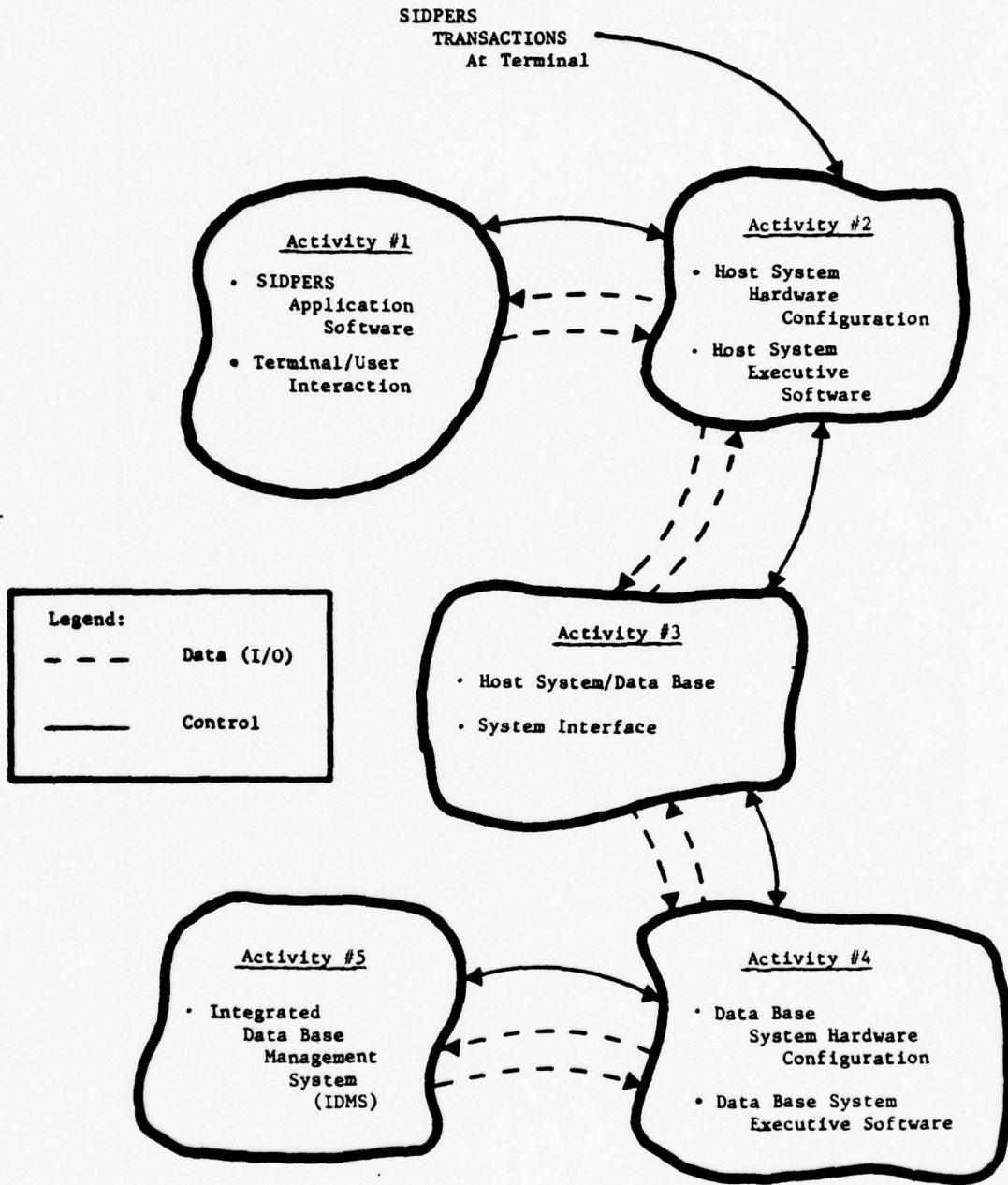


Figure 4.2-1 Structure of Final SIDPERS/IPSS Model

System Facilities Represented	Models		SIDPERS/IDMS Workload Model	Host Submodel		Back-end Submodels		Host/Back-end Models	IDMS	Total
	Host Submodel			Back-end Submodels						
	V1	V2		V1	V2					
<u>Application Environment</u>										
<u>Activity #1</u>										
● IDPERS Application Software	●	●	●					●		●
● Terminal User Interface	●	●	●					●		●
<u>Activity #2</u>										
● Cost System Hardware	●	●	●					●		●
● Cost System Executive Software	●	●	●					●		●
<u>Activity #3</u>										
● Cost to PDP 11/70 Interface		●						●		●
<u>Data base Environment (PDP 11/70)</u>										
<u>Activity #3</u>										
● PDP 11/70 to Host Interface								●		●
<u>Activity #4</u>										
● PDP 11/70 Hardware				●				●		●
● PDP 11/70 Executive Software				●				●		●
<u>Activity #5</u>										
● DMS General Software			●					●	●	●
● DMS DML Statements								●	●	●
● Data Base - Logical Structure								●	●	●
● Data base - Physical Structure								●		●

Figure 4.2-2 IPSS SIDPERS/IPMS Models

Phase II focused on the host environment. The base model was system Activity #1 as represented in the Phase I model with the backend treated as a black box. To this base model was added successively more detailed characterizations of the host system. The final model in this phase included the characterization for the backend interface hardware and software. A synopsis of Phase II activities is given below. Appendix A2 presents detail consideration of the modeling activities in this phase.

Phase III was concerned with the characterization of the backend system. Again the modeling proceeded in an evolutionary manner. A sequence of models was planned which would add the PDP 11/70 hardware and operating system software to the IDMS workload model developed in Phase I. For these activities the front end was treated as a white box, i.e., an external source of back end requests. Specifically omitted was any attempt to model IDMS DML's or traversal activities for the SIDPERS logical data base (as defined via the IDMS DDL's). These latter functions were modeled using the IPSS/DBS facilities. A synopsis of this phase is given below. Details are found in Appendix A3.

The goal of Phase IV was to combine the Phase II and III models into a single model. As mentioned at the beginning of this appendix, a minor IPSS programming error, since corrected prevented their phase to be completed within the scheduled experimentation time line. Results of this modeling phase will be reported at a latter date.

Phase V was planned whose purpose was to model the IDMS DML's and the DDL described SIDPERS data base. A paper model was completed. However, execution was delayed until completion of the IPSS/DBS parsers. As with Phase IV the activity was completed after the termination of the modeling activities. This model is being reported in Mr. Brownsmith's dissertation, titled "A Methodology for the Performance Assessment of Data Base Systems: An Extension to the Information Processing System Simulator Methodology", which is to be available Spring, 1979 from Ohio State University.

The goal of Phase VI was to integrate the Phase V model with the final Phase IV model. Again this activity has to await modifications to the current IPSS and IPSS/DBS.

A Synopsis of the Phase II Activities (SIDPERS/HOST Models)

The SIDPERS application model was built assuming a computer free environment. The purpose of the Host submodel was to provide both the hardware configuration and the operating system of a host computer. Since no particular computer system had been identified as the host computer for the SIDPERS/IDMS project, a generalized host model was constructed. The elements of this model were characterizations of those operating system functions that support the SIDPERS application and the Host/Back-end interface. These include the job scheduling, task management and resource allocation, assumed to be available in the Phase I model, were now characterized. The hardware configuration that was modeled consisted of only these hardware components required to support the modeled host software.

The approach to construction of the host model followed the premise that the host operated in a reactive mode. The existing SIDPERS model was analyzed and when an operating system function was required, a module that characterized its behavior was added to the Phase II models. This procedure was followed for all levels of applications represented in the SIDPERS 1, and for the Host/Back-end interface needs. As a result a model was created of a host computer environment that satisfied the SIDPERS operating needs. The operating system functions included in the host model were characterized in a general, modular parameterized manner, such that at a later time, the mechanisms of the functions of a particular host computer could be easily incorporated into the existing model. Thus, although system activities 1 and 2 were communicating continuously, modifications to the latter could be made without affecting the existing SIDPERS application model. The SIDPERS/HOST submodels were tested by "black boxing" the back-end machine.

Synopsis of Phase III Activities (BACK-END/IDMS Models)

The IDMS model was built like the SIDPERS application model, assuming no particular computer environment. The purpose of the Back-end system

model was to provide the hardware configuration and the operating system of the back-end environment. Previous Army studies had designated the PDP-11/70 as the specific computer to be employed as the back-end computer. The functions of this computer system that needed to be modeled were the operating system functions to support IDMS and Host/Back-end interface routines. As with the host submodel, the back-end submodel was assumed to operate in a reactive mode, responding to requests for its services. The hardware configuration that was modeled was the complete USACSC PDP-11/70 configuration, recently moved with AIRMICS to the Georgia Institute of Technology in Atlanta.

The back-end models were built modularly, adding characteristics of operating system or interfacing functions as they were needed. Although these characterizations were made specific to the AIRMICS PDP-11/70 computing environment, they were modeled in a parameterized, modular manner so that if different PDP-11/70 strategies were to be substituted, the IDMS and interface submodels would not be affected. The IDMS/Back-end models were tested using the SIDPERS application model as the work load generator.

4.3 INSIGHTS GAINED FROM THE MODELING

The purpose of the modeling effort reported in this document was to demonstrate the feasibility and suitability of the Information Processing System Simulator (IPSS) with respect to modeling large complex information systems. The system modeled to demonstrate these capabilities was a host/back-end processor configuration which supported interactive application processing and a data base management system. The modeled software processes included a detailed characterization of application loading, DBMS processing, and the operating system functions of task management, job scheduling and resource allocation.

The IPSS methodology provided a modular, top-down approach to system analysis and model synthesis thus facilitating a solution to this complex problem. Several models were constructed, each of which

reflected system details suited to the multi-stage development process which was adopted. The ease of model expandability in both breadth and depth was demonstrated through the construction of independent models and their integration into an overall complete model. The SIDPERS/IDMS workload model represents the software processes of the applications and IDMS in detail while characterizing the operating systems and computer hardware as a black box. From this basis model, two subsequent models were developed by splitting the software processes at the host/back-end interface and adding characterizations of the appropriate hardware and operating system software functions. The synthesis of the final overall model was easily accomplished since the system interfaces were clearly identified.

These models were independently tested, each producing a set of statistics reflective of the modeled processes and computing environment. These results are summarized in Appendix B. In all cases queueing and utilization statistics were obtained for the modeled software and hardware resources. These statistics provide an easily understood synopsis of activities since one or more IPSS services were used to represent major units of processing (e.g., application programs, TP line scheduling, IDMS processing). Statistics were also automatically collected on each of the hardware facilities indicating both potential bottlenecks and under-utilized components.

These models were all internally verified. That is, through experimentation, the internal processing consistency was verified to be accurately reflective of the real world processes. Due to the lack of operational data, however, it was not possible to validate them through comparison to actual system performance. Thus, although some confidence can be placed in the relative values of the statistics reported in the Appendix B, no validation has been made with respect to the magnitude of the data values.

This project has demonstrated the suitability of IPSS to model complex systems while requiring a minimal time for model synthesis. The modeling effort as reported here required approximately four man months. The resulting models have proven to be versatile as well as

adaptable to changing requirements indicating that significant portions can be used for other application areas. For example, the DBMS and operating system processes as well as hardware facilities need not be modified when another major application area is modeled, thus resulting in a substantial reduction in future related modeling efforts.

In addition, this project has served to focus future development work as well as to validate current goals. In particular, this project has demonstrated the need for simulation facilities specifically attuned to the characterization of integrated data structures as represented in IPSS/DBS. These models did not use these features since they were not fully operational, and thus the logical structure of the data base was characterized at a more abstract level than desired. Furthermore, experimentation with record mapping policies was not attempted because of the complexities involved. Finally, several other research endeavors were identified. In particular, large complex models of computer networks could easily be constructed from models of individual nodes if a multi-simulator approach were adopted. Work is currently underway to allow multiple node simulations through extension of IPSS facilities and concepts. Extensions for modeling teleprocessing activities are also planned.

5.0 FUTURE EXTENSIONS TO THE IPSS

This research activity is just another step in a continuing IPSS research, development and engineering program whose goal is to create a comprehensive computer aided design facility appropriate to all stages in an information system's lifecycle. Three steps in the RD&E plan have now been completed. These include

1. 1967 - 1972: Initial Methodological and Simulator Development - this activity resulted in the Auxiliary Storage System Simulator (AUXSYM). Details are described in "A Methodology for the Analysis of Auxiliary Storage Systems" available from University Microfilms.
2. 1972 - 1975: Extension to the methodology and complete redesign of the simulator. These activities resulted in the IPSS; details are reported in "A Methodology for the Performance Evaluation of Information Processing Systems" (DEL77) and "The Information Processing System Simulator (IPSS): Language Syntax and Semantics".
3. 1975 - 1977: Continued methodological development to include the following
 - a. DBS specific features;
 - b. Assessment of system effectiveness with respect to competing user goals;
 - c. Addition of mechanisms to determine costing policies.
 - d. Assessment of system loading created by multi key query processing
 - e. Development of comparative evaluative techniques for simulation languages vis-a-vis performance assessment needs.
4. 1976 - 1979: Verification and Validation - The purpose of these activities has been to demonstrate the effectiveness of the IPSS as a modeling facility. Validation has included paper models of different vendor hardware

software and file structures, models of idealized systems, models of existing complex systems (this report describes one such system.) A proposal by Chandler is designed to extend this activity into more specific validation activities where extrapolation can be validated.

As a result of the current and past research experiences, an extensive number of continuing research and experimental activities have been identified. They are presented here to illustrate the potential for the IPSS and its underlying methodology and to identify a means of activities which must be accomplished in order to achieve the ultimate research objective - that being to create a computer assisted design facility.

A. Extensions to the IPSS Methodology

The goal of these research activities is to continue the extension of the methodology's scope of applicability to (a) higher views of information system processes (b) other stages in the information process system lifecycle and (c) analysis and design processes. Five research activities have been identified for this area. They are:

1. Telecommunications - extension of the IPSS the permit characterization of data communication subsystems in a manner familiar to their structure and behavior.
2. Requirement Analysis - extension of the IPSS to interface with requirements analysis methodologies such as Tiechrow's PSL/PSA and TRW Corp's SCREM.⁽⁸⁾
3. Sensitivity Analysis - development of a methodology which permits the assessment of models to identify performance sensitive components.
4. Estimating System Lifecycle Costs - development of a methodology to examine IPSS models for the purpose of (a) estimating the number and complexity of components comprising the modeled system, and (b) using these data as input to lifecycle estimating techniques.

⁽⁸⁾ PSL/PSA is acronym for Problem Statement Language/Problem Statement Analyzer SREM stands for System Requirement Engineering Methodology.

5. Automated Component Selection - formulates a methodology which is based on a predefined view of system components, and on one or more design criteria will automatically alter IPSS models thus identifying feasible designs for additional analyses by the modeler.

B. Extensions to the IPSS Language

The goal of this development activity is to expand the model synthesis and statistical display facilities of the current IPSS in order to reduce modeler effort. Five development activities have been identified. They are:

1. Statistics Generation - development of language extensions to cause the automatic gathering of standard statistics for sequential ("serially reusable") processes.
2. Verification and Validation - extends existing and develop new statements to support model verification and validation activities.
3. Operating System Characterizations - extend existing and develop new facilities attuned to; (a) the modeling of information system job scheduling, resource allocation and task management functions and (b) characterizing main memory management, CPU interrupt activity and CPU instruction times.

C. Extensions to IPSS Simulation Facility

The purpose of these activities is to increase the ease of use of the IPSS simulator facility by the modeler. The following four items are proposed as extensions in this area:

1. On-Line Modeling - extend the processing environment for the IPSS to include; (a) on-line entry of IPSS language statements comprising an IPSS input stream, and (b) an on-line, real-time prompter facility to guide the modeler in the use of IPSS language syntax. The modeler will be able to switch between model stream definitions and stream processing as desired.

2. MACRO Library - the IPSS macro and library facilities currently permits the modeler to retain models and/or model components for future reference. This activity would develop a set of generalized structures presenting the salient features of the following software:
 - a. Operating System: task management, job scheduling, resource allocation and memory management;
 - b. File Utilities: sort/merge, file copy, and opening and closing of files;
 - c. Data Base Management System: buffer management, high level DML's, schema structures.
3. Graphical Display - the IPSS would be enhanced to permit graphical interaction between the system and the modeler. Initially this will involve only output of items such as:
 - a. Model statistics including frequency listographics.
 - b. Model structure showing the network of services described for the model and the data and file structures.
 - c. Dynamic model structure showing the sequencing of service invocation and their hierachy during a task's execution.

Ultimately, the goal of graphics use is to prove a modeler/system communication interface which will permit interactive support of model synthesis and the display of model structure, model verification data, and model behavior statistics.

REFERENCES

- ABA68 Abate, J., Dubner, H., and Weinberg, S.B., "Queueing Analysis of the IBM 2314 Disk Storage Facility", JACM, Vol. 15, No. 4, (1968), pp. 577-589.
- BOU72 Boutross, King and Rutledge, "Discrete Simulation Model of DASD", IBM Research (RC3698), Jan, 1972.
- BUZ73 Buzen, J.P., "Computer Algorithms for Closed Queueing Networks with Exponential Servers", CACM, Vol. 16, No. 9, (1973), pp. 527-531.
- BUZ76 Buzen, J.P., "Fundamental Laws of Computer Performance", in Proceedings of International Symposium on Computer Performance Modeling, Measurement and Evaluation (CPMME), (1976) pp. 200-210.
- CHN77 Chandler, J.S. and DeLutis, T.G., "A Multi-Criteria Approach to Information Systems Design", to be published in Proceedings of National Computer Conference, (1977).
- COL72 Cole, Gerald, D., "Performance Measurements on the ARPA Computer Network", IEEE Transaction on Communication, (1972) pp. 630-636.
- CON65 Conger, C.R., "The Simulation and Evaluation of Information Retrieval Systems", Report 352-R-17, (April 1965).
- DEL72 DeLutis, T.G., "A Methodology for the Analysis of Auxiliary Storage Systems", Ph.D. Dissertation, Purdue University, 1972.
- DIJ68 Dijkstra, E.W., "The Structure of 'T.H.E' Multiprogramming System", CACM, Vol. 11, No. 5, (1968), pp. 341-346.
- GAV76 Gaver, D.P., and Hunfeld, G., "Multitype Multiprogramming: Probability Models and Numerical Procedures", Proceedings of GPMME, (1976), pp. 38-43.
- GHO76 Ghosh, S.P., and Tuel, W.G., "A Design of an Experiment to Model Data Base System Performance", IEEE Transactions on Software Engineering, Vol. SE-2, No. 2, (1976).
- GRI75 Griffith, W.G., Ingerman, D., and Price, C.E., "A Simulation Model of Univac's D,S-1100 -- More than Just a Performance Evaluation Tool", Symposium on the Simulating Computer Systems, III, (1975), pp. 91-99.
- HAL74 Hall, W.A., "A Simulation Model to Aid in the Design and Timing of Hierarchical Databases", Winter Simulation Conference, (1974), pp. 277-284.

- HEL70 Hellerman, H.R., and Smith, H.J., "Throughput Analysis of Some Idealized Input, Output and Computer Overlap Configurations", Computing Surveys, Vol. 2, No. 2, (1970), pp. 111-118.
- HUL75 Hultén, C., and Söderlund, L., "A Tool for Performance Evaluation of Data Base Oriented Information Systems", Project MAD, Research Group CADIS, Department of Information Processing and Computer Science, The Royal Institute of Technology and the University of Stockholm, Sweden, (1975).
- KLE72 Kleinrock, L., and Muntz, R.R., "Processor-Sharing Queueing Models of Mixed Scheduling Disciplines for Time-Shared Systems", JACM, Vol. 19, No. 3, (1972), pp. 464-482.
- KLE76 Kleinrock, L., Queueing Systems, Volume 1: Theory and Volume 2: Computer Applications, New York: John Wiley & Sons, Inc, (1976).
- LIN76 Lindsay, D.S., "A Hardware Monitor Study of a CDC KRONOS System", Proceedings of CPMME, (1976), pp. 179-186.
- LUM70 Lum, V.Y., Ling, H., and Senko, M.E., "Analysis of a Complex Data Management Access Method by Simulation Modeling", FJCC, (1970), pp. 211-222.
- MAD74 Madnick, S., and Donovan, J., Operating Systems, New York: McGraw-Hill, 1974.
- MLS76 Mills, P.M., "The Overlapping of CPU and I/O Processing in Multiprogramming", in Proceedings of CMG-VII, (1976), pp. 200-206.
- NAH73 Nahourii, E., "Direct Access Device Simulation", IBM Systems Journal, Vol. 13, No. 1, (1973), pp. 19-31.
- NAK75 Nakamura, F., Yoshida, I., and Kondo, H., "A Simulation Model for Data Base System Performance Evaluation", National Computer Conference, (1975), pp. 459-465.
- NOR71 Norland, K.E., and Bulgren, W.C., "A Simulation Model of GECOS III", Proceedings of ACM, (1971), pp. 596-612.
- OWE70 Owens, P.N., "Phase II: A Data Base Management Modeling System", Proceedings of IFIP, (1970).
- REE75 Reeves, T.E., and Pooch, U.W., "A Multiple Subsystem Simulation of Processor Scheduling", Symposium on the Simulation of Computer Systems, III, (1975), pp. 129-135.

- REI75 Reiter, A., "On Performance Modeling of Data Base Management Systems -- An Inductive Approach", MRC Technical Summary Report, Mathematics Research Center, University of Wisconsin, Madison, (1975).
- REI76 Reiter, A., and Finkel, B., "Simulating a Virtual Data Machine", MRC Technical Summary Report, Mathematics Research Center, University of Wisconsin, Madison (1976).
- ROE70 Roehrkaese, R.C., and Smith, D., "Simulation of Operating Systems", Technical Report GITIS-70-11, (1970), Georgia Institute of Technology.
- SAB71 Sable, J.D., "Generalized Data Management Systems: Query and Language Processing", Proceedings of the International Seminar on Information Storage and Retrieval, (1971), pp. 93-111.
- SCH76 Schneider, L.S., and Connolly, T.W., "A Generalized Database System Simulator Based on the Data Independent Accessing Model I", Winter Simulation Conference, (1976), pp. 513-524.
- SDB77 SIGBDP, "Guidelines for Cost Accounting Practices for Data Processing", DATA BASE, Vol. 8, No. 3, 1977.
- SEN70 Senko, M.E., Altman, E.B., Astrahan, M.M., Fehder, P.L., and Wang, C.P., "A Data-Independent Architecture Model 1: Four Levels of Description from Logical Structures to Physical Search Structures", IBM Research Report RJ982, San Jose, California, February 1972.
- ZUR69 Zurcher, F.W., and Randall, B., "Iterative Multi-Level Modeling: A Methodology for Computer System Design", IFIP 68, pp. 867-871.

ADDITIONAL REFERENCES

- DEL76 Delutis, T.G., and Smith, J.D., "IPSS/DBMS: A Simulator for Evaluating DBMS Performance", Winter Simulation Conference, 1976, pp. 527-536.
- DEL77 Delutis, T.G., "A Methodology for the Performance Evaluation of Information Processing Systems", Final Report to the National Science Foundation, Grant Number GN 36622, 1977, Available via NTIS.

REFERENCES USED IN CONJUNCTION
WITH MODELING EXPERIMENTS

Cullinane Corporation, IDMS Data Base Design and Definition Guide, Release 4.0, June 1976.

Cullinane Corporation, IDMS DDL Data Definition Languages, Utilities and GCI Reference Guide, Release 3.1, Revision 1, April 1975.

Cullinane Corporation, IDMS DML Data Manipulation Language Programmer's Reference Guide, Release 3.1, April 1975.

Cullinane, J., Goldman, R., Mueurer, T., Nawara, R., "Commercial Data Management Processor Study", Cullinane Corporation, Wellesley Mass., December 1975.

Digital Equipment Corporation, IAS User's Guide, DEC-11-OIUGA-B-D, 1976.

Digital Equipment Corporation, Microprocessor Handbook, 1976.

Digital Equipment Corporation, PDP11 IAS/R SX-11 I/O Operations Reference Manual, DEC11-OIORA-A-D, 1975.

Digital Equipment Corporation, PDP11 Peripherals Handbook, 1976.

Digital Equipment Corporation, PDP11/70 Processor Handbook, 1976.

Digital Equipment Corporation, PDP-11 Computer Family Products and Services, 1976.

Hunter, F., Personal Communication, June 1977.

Hunter, F., Mercier, R., Peckham, G., Sernovitz, L., "Report on the Conversion of a File Oriented Interactive System to an Interactive Data Base Management System", U.S. Army Computer Systems Command, Fort Belvoir, VA, 1977.

Mercier, R., Mills, R., Price, J., Sernovitz, L., "Report of an Experimental On-line System", U.S. Army Computer Systems Command, Fort Belvoir, VA 1976.

Schaff, H., "Description of the VIABLE Transactions for the DIMUI Model", U.S. Army Computer Systems Command, Fort Belvoir, VA, 1976.

PRECEDING PAGE BLANK-NOT FILMED

APPENDICES

APPENDIX A

BACKGROUND TO THE MODELING EXPERIMENTS

The objective of this experiment was to provide a realistic problem environment against which the IPSS design facility (simulator and methodology) could be applied. The system modeled was chosen by Mr. Allan Curry of AIRMICS* specifically to test IPSS's design capabilities, ease of implementation, and adequacy and accuracy of the resultant statistics. The chosen system architecturally is of great interest to the U.S.A. Computer Systems Command as a replacement for exists BASOP systems. Section A1 of this appendix discusses the problem environment including a full breakdown of the hardware and software characteristics of the system. Section A2 defines the IPSS models generated to characterize the system.

A1. THE PROBLEM ENVIRONMENT

The purpose of the following discussion is to provide a background to the system modeled. The system modeled evolved from a series of R&D efforts performed by the U.S. Army Computer Systems Command (USACSC) at the request of the Department of the Army, Director of Management Information Systems (DA DMIS) in direct support of their "Vertical Installation Automation Baseline" (VIABLE) Project. The VIABLE project was established to upgrade both hardware and software at Army data processing installations. Included in its scope is the competitive procurement of new ADP equipment, and the extension of functional software capabilities to include management information system requirements as well as interactive processing.

The initial USACSC research plan was developed in September 1975. Its purpose was to explore and demonstrate those technologies which would conceivably be utilized in Army data processing systems developed

*AIRMICS stands for Army Institute for Research in Management Information and Computer Sciences.

over the next five to eight years. These areas included on-line processing, formal data base management systems, mini and micro computers and distributed data processing systems.

One aspect of the plan was to select a single functional area and build a baseline system which could be modified repetitively to add new features and migrate from one environment to another. The Standard Installation/Division Personnel System (SIDPERS), which processes Military Personnel actions was selected.

The computer system modeled was composed of a Front End Module (host machine) which processes SIDPERS transactions in a real-time interactive mode, attached to a Back End Module utilizing Cullinane Corporation's Integrated Data Base Management System (IDMS) to interface with the personnel data base. The backend machine modeled was the DEC PDP-11/70 System.

A1.1 SIDPERS

SIDPERS (Standard Installation/Division Personnel System) is the batch oriented multi-command personnel data processing system. It is a large system, consisting of approximately 800 types of transactions which perform update and retrieval operations on a personnel data base. However, since six of these transactions account for almost eighty percent of the total transaction activity, only these transactions were selected for incorporation into the baseline system. These transactions perform the following functions: adding a soldier to the data base, establishing duty status when a soldier arrives at an installation, changing the duty status and grade of a soldier, deleting a soldier from the data base on departure from an installation and acquiring data on a soldier and/or unit. These commands are denoted ADD, ARRIVAL, DUTY CHANGE, GRADE CHANGE, DEPARTURE, and INQUIRY respectively.

The SIDPERS test system required the creation of a small data base of sample records. Approximately 4,500 records representing 12 different

record types were selected and organized into an IDMS processable data base. The details of this data base and its organization are given in the next section.

The SIDPERS test system facilities for interactive transaction processing have been described by Schaaf (SCH77)⁽¹⁾. Basically, special terminal interface processing modules were designed which allowed a terminal user to enter the above transactions in either a batch or tutorial mode of operation. In batch mode, all the data for a complete transaction is entered at once. In the tutorial mode, the terminal user is prompted for every data item of the transaction. This requires an access to a VALIDITY record which contains the text for prompting. Items that fail to pass the validity check are returned to the terminal operator for re-submittal. For either mode of operation, once all the data items of a transaction pass this test, the corresponding application program is invoked. Essentially, there is one application program for each of the transactions listed above. As part of application processing, compatibility checks are performed. If a data item doesn't pass this check, several accesses are required to correct the error; namely an access to a COMPATIBLE record, and several accesses to VALIDITY records. Such errors cannot always be resolved and in these cases the transaction is terminated. Otherwise the application program issues a series of DML commands to the IDMS Modules in order to store, modify, and retrieve records in the data base. During the IDMS Processing of each DML, application processing is suspended. When the application is completed, control is returned for further interaction.

A1.2 Prior Analysis Activities of SIDPERS

The data required to characterize these activities was provided from a previous study of VIABLE transactions (SHA77). (See Tables A1.2-1 and A1.2-2 for further specifics). This document provided the following data for each of the six types of transactions:

⁽¹⁾ NOTE. All references are found at the end of the main body of this report.

Table A1.2-1 The Results of the SIDPERS Interaction Test

	No. of entered data items	Data no. of transactions	No. of Cancelled transactions	No. of Transactions in Direct mode	No. of transactions in tutorial mode	No. of validity errors	% of data items with val. errors	No. of compatibility errors	% of data items with comp. errors
ARRIVAL	9	158	25 / 16%	93 / 59%	65 / 41%	29		31	
DEPARTURE	12	193	34 / 18%	111 / 57%	82 / 43%	60		101	
DUTY CHANGE	8	331	23 / 7%	208 / 63%	123 / 37%	91		170	
GRADE CHANGE	10	52	11 / 21%	31 / 60%	21 / 40%	15		34	
ADD	70	6	0 / 0%	0 / 0%	6 / 100%	17		6	
INQUIRY	3	38	0 / 0%	31 / 82%	7 / 18%	2		12	
TOTALS		778	93 / 12%	474 / 61%	304 / 39%	214	2.9%	354	4.8%

Table A1.2-2 Average Entering Time

Transaction Type	Entering Time Per Transaction		Transaction Characteristics		Entering Time Per Data Item	
	Tutorial Min:Sec	Direct Min:Sec	No. of Entered Data Items	Errors (about 8%)	Tutorial (sec)	Direct (sec)
ARRIVAL	5:23	2:49	9	1	32.3	16.9
DEPARTURE	6:03	2:55	12	1	27.9	13.7
DUTY CHANGE	3:58	1:40	8	1	26.4	11.1
GRADE CHANGE	5:55	1:48	10	1	32.3	3.8
ADD	25:57		70	6	20.5	
INQUIRY	1:23	0:54	3	-	27.7	18
AVERAGE OVER ALL TRANSACTIONS					27.9	13.9

1. the relative frequency of occurrence,
2. the percent direct mode,
3. the number of data items per transaction entry,
4. the percent data item validity error,
5. the percent data item compatibility error,
6. the time for direct mode data items,
7. the time for tutorial mode data items, and
8. the percent of transactions canceled.

This data was incorporated into the IPSS model of SIDPERS/IDMS through Fortran DATA statements and was used to "drive" the model through representative terminal sessions. In addition, this document contains for each transaction, a list of the DMLs that are invoked and their sequence within the application program. This data was used to invoke lower level DBMS processes as described in the next section.

A1.3 IPSS Model of SIDPERS

The IPSS SIDPERS/IDMS transaction model was designed to reflect the arrival validation and DML activities of each of the six previously identified SIDPERS transactions. Figure A1,3-1 identifies the major functional areas of this model. Each functional area is represented by an IPSS service or services within the model. The function of each of these services are defined as follows:

System Loading

Within the IPSS model, the arrival of a transaction at a work station is represented by an exogenous event (defined within the Exogenous Event Stream Component). This event was scheduled to occur every 5.89 minutes plus or minus 5 minutes. This is based upon the following assumptions and parameters. First, the data base represents 203 soldiers; second, that there are an average of 8 transactions per month per soldier; third, that data base updating activities through

terminal interaction is possible 8 hours per day, 20 days per month. These figures were used to derive the transaction arrival event frequency of 5.80 minutes. The deviation of 5 minutes was selected arbitrarily to represent variation in the arrivals. In order to model a larger data base or different arrival patterns, a single statement need be changed in the procedure associated with the event. This fact is documented within the procedure and the statement identified for ease of modeler experimentation.

Transaction Arrival and Sequencing

This service represents the arrival of a transaction at a work station. The model assumes a single terminal and hence the transaction may be queued if the previous one has not completed. The model also assumes a FIFO (first-in/first-out) queueing discipline. The time that a transaction waits in this queue is a statistic automatically provided by IPSS. When a transaction is selected for processing, it is removed from the queue, the transaction Selection Service is invoked, and the service waits its completion. The total elapsed time from transaction selection to completion is also maintained by IPSS. Upon completion, the next transaction is selected from the queue and the service becomes idle. Idle time and busy period statistics are also kept. All the statistics mentioned above are maintained for each of the services and will not be further documented in this section.

This service can easily be modified to reflect an arbitrary number of terminals simply by establishing the maximum number of terminals permitted to be in operation at once, and keeping track of the current number. Only those transactions which arrive when the current number is equal to the maximum number need be queued. An infinite number of terminals can be modeled simply by not comparing these two variables and hence never queueing a transaction.

Transaction Selection

This service selects the type of transaction that has arrived and associates with it a set of parameters,⁽²⁾ The selection is made by a probability distribution function based on relative frequency data provided by the Schaaff (SCH77) study. Once this selection is made, this service invokes the terminal interaction and validation service and waits for its completion. When the processing associated with this lower level service has completed, the Transaction Selection Service, determines if the transaction has been canceled (via a return code) and, if not, invokes the proper Application Service and waits for its completion.

Terminal Interaction and Validation

This service represents all the activities associated with the entry and validation of a transaction by a terminal operator; and can be conceptually divided into two parts, namely those activities that occur in obtaining valid transaction through terminal interaction in direct mode and those activities required to process and complete the transaction in tutorial mode. These processes are both characterized through tabular data and probability distribution functions. Input from the terminal, for example, is represented by advancing the simulator clock. In separate models, facilities representing the hardware and software processes are characterized and thus these activities are more accurately measured. The tutorial input mode requires several more parameters than the direct mode to characterize the events. The final activity of this service is to determine whether the transaction is to be canceled. This is communicated to the invoking service by a parameter value.

⁽²⁾ These parameters represent the transaction characteristics and are derived from the data in Tables A1.2-1 and A1.2-2.

Application Services

The model contains a set of six functionally similar services, each of which represents the processing performed by an application program for one of the transactions. These services are all invoked by the Transaction Selection Service and they all at some time invoke the IDMS Interface Service. The processing performed by these transactions is represented by a series of DML invocations to this lower level service. The internal data manipulation activities are not represented in the model since their impact on performance is assumed to be minimal. The DMLs are identical in type and sequencing to those reported by Schaaf (SCH77); their format is further discussed in the next section.

While not all performance evaluations characterize application processing to the detail represented in this model, this approach was deemed best for this model for the following reasons: (1) the data was readily available, (2) the mechanics of model construction was facilitated by IPSS language modeling features and (3) the model architecture emulates the processing performed by the real system. Thus, when the model is used in the future it will be evident what processes are being modeled.

Within IPSS, a service is dynamic facility and as such statistics are automatically collected for its use. For each of the services identified in Figure A1.3-1, utilization statistics are produced. These include: busy period, idle period, concurrency, seizure, and transaction transmit time statistics. In addition, queueing statistics are also produced. However, because of the nature of the model, only the Transaction Sequencing Service has queueing associated with it. These queueing statistics include: busy period, idle period, queue length, queue entry, and transaction transmit time statistics.

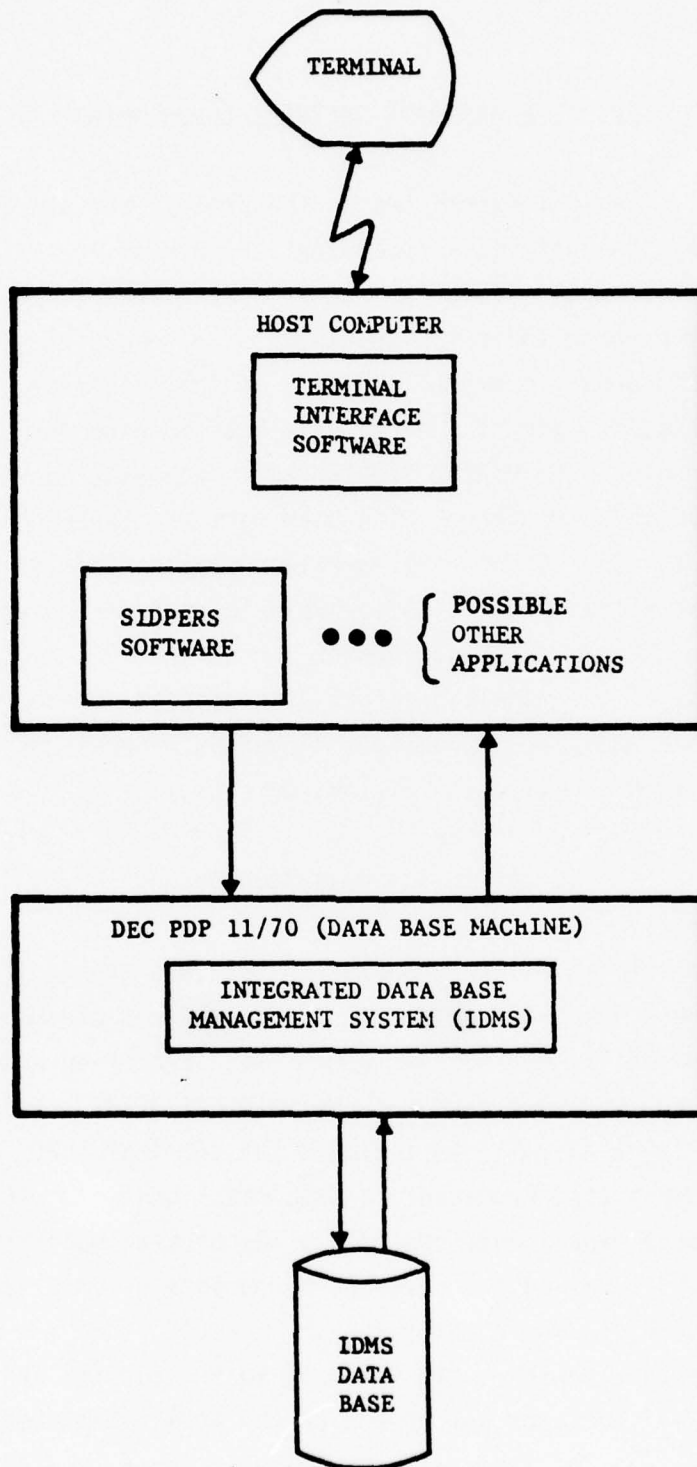


Figure A2-1 Experimental Version of VIABLE SIDPERS Processing Environment

A2. THE HARDWARE/SOFTWARE ENVIRONMENT

The configuration established as the prototype for the VIABLE SIDPERS project employs the host/backend concept (Figure A2-1). The IDMS SIDPERS data base would reside on and be controlled by the backend machine. The purpose for this structure is to reduce the load on the host machine in order to extend the life of the mainframe. Previous studies have not designated a particular host machine for use by USACSC. Currently, the VIABLE SIDPERS test data base is maintained on the USACSC HCC IBM 370 system, with both host and backend machines residing there. One of the most immediate applications of the models developed during this project is to evaluate different host environments, in particular, the IBM 370 environment versus the PDP-11/70 environment. Thus, the goal of the modeling effort with respect to the host machine is to develop a model of a machine independent generalized host machine, allowing for each substitution of environments.

A2.1 The Host Machine

To accurately model the IBM 370/165 host was deemed to be too time consuming a task with regard to the objectives of the project. Therefore, an idealized host environment was identified which paralleled the hardware and software needs of the VIABLE/SIDPERS. Host hardware is shown in Figure A2.1-1. It included the terminal, CPU, secondary storage equipment (I/O Processor, disk control unit and eight IBM 3330 type disk units), and a data channel to the backend machine. Also shown in this figure are the hardware components assumed for the backend computer.

Host software functions in addition to the SIDPERS are shown in Figure A2.1-2. The major components of the application/host environment are the teleprocessing (TP) interface with the user, the TP interface with the backend machine, the host's disk I/O software, and the SIDPERS

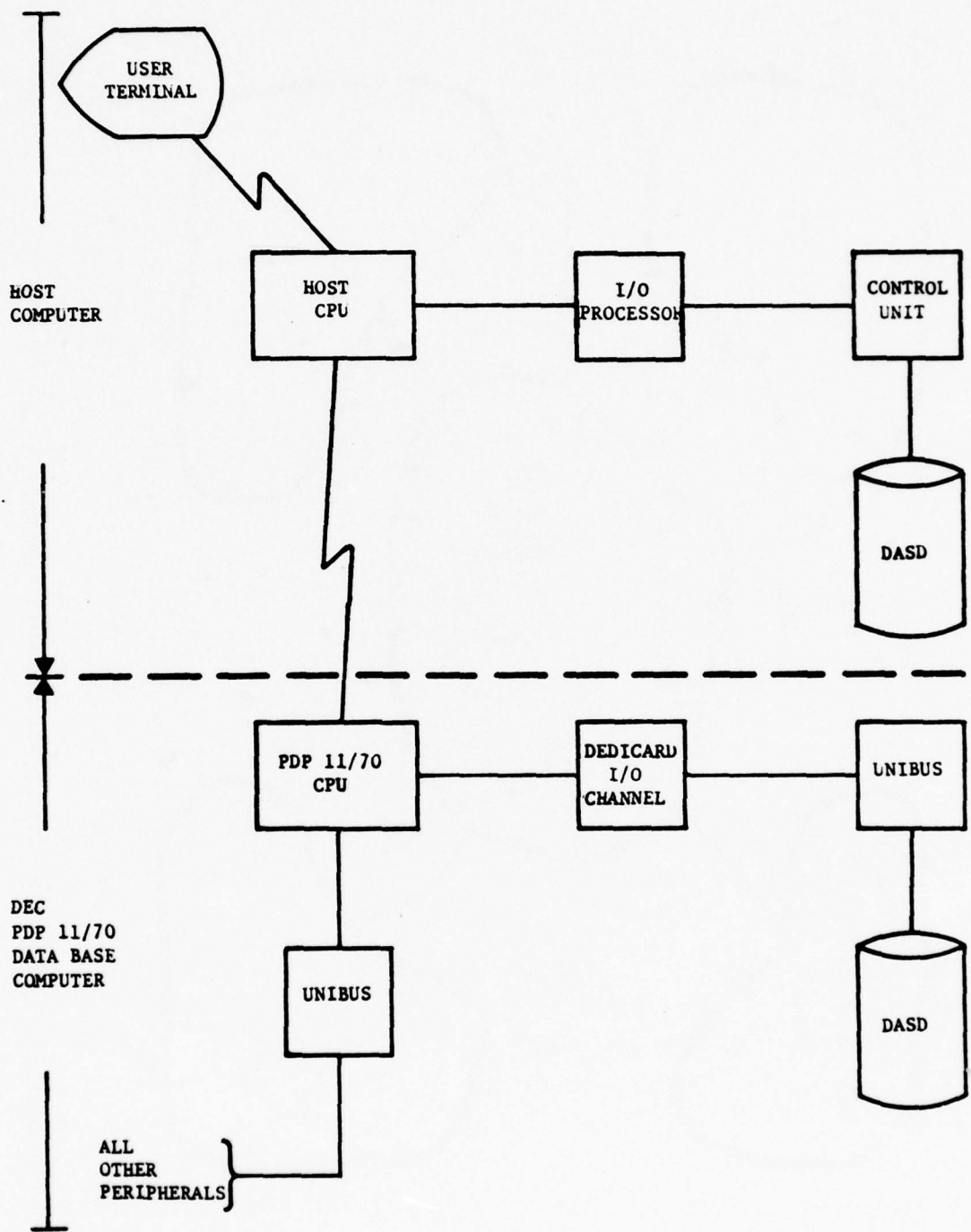


Figure A2.1-1 Modeled Hardware Configuration

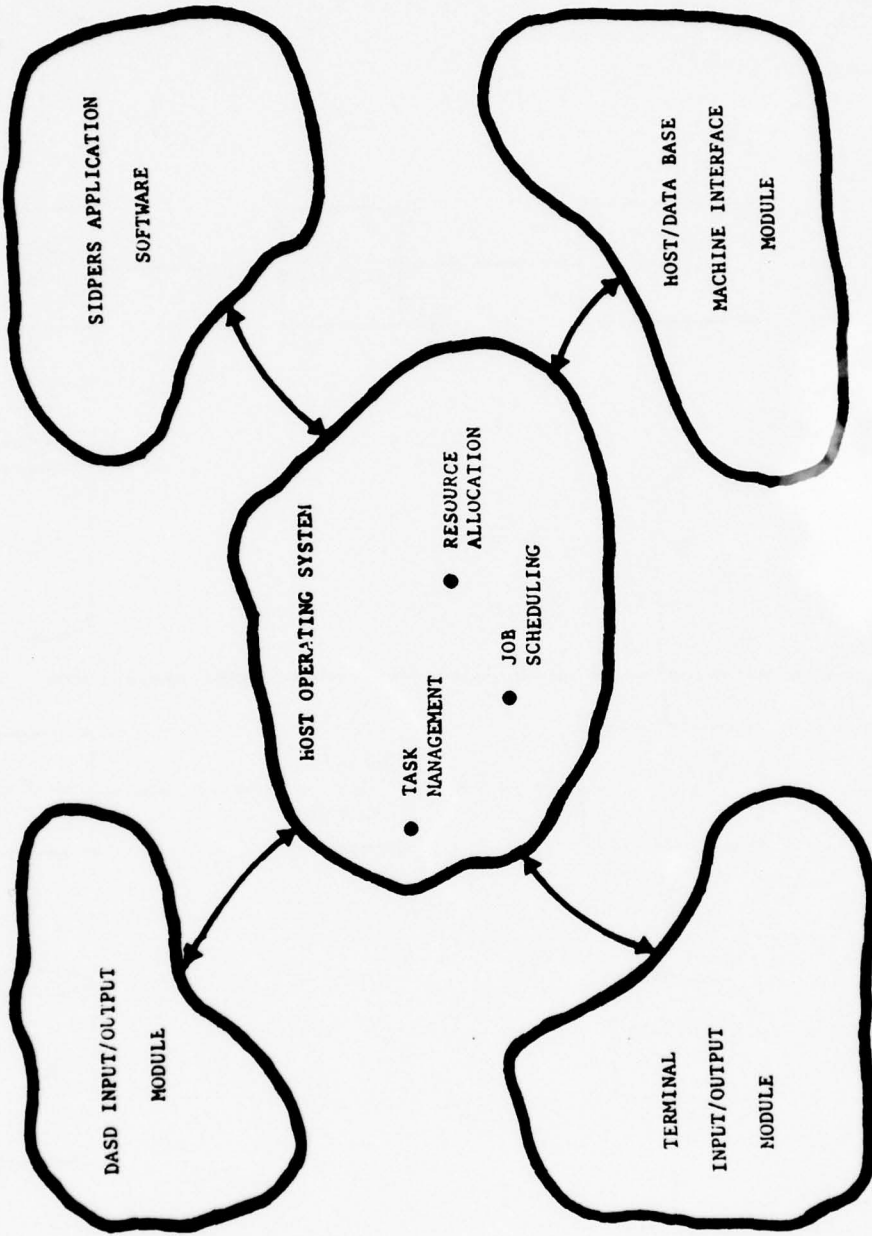


Figure A2.1-2 Modeled Host System Resident Software

application itself. The generalized host operating system (OS) is assumed to be reactive in operation, i.e., executing only to respond to the needs of the other software components.

Since VIABLE SIDPERS is interactive, communication with the user is maintained throughout each step in the process. In this capacity, the host performs the function of scheduling the communication, managing the TP line and sending/receiving the messages. After the initial log-on procedure the host operating system schedules the initiation of SIDPERS processing. The input and validation of SIDPERS data involves, in addition to the host's TP functions, disk I/O for the purpose of retrieving error messages for transmission to the terminal.

The crucial element of this model is the liaison between the application in the host machine and the DBMS in the backend machine. The DML request must be transmitted to the backend for processing and the retrieved record must be received. To accomplish this, the host supports an interface subsystem which formats the DML request, reformats the resultant record, handles the asynchronous transmission and reception of messages to and from the backend, and controls the necessary devices for this communication link. It is assumed that each of the host modules rely on the basic operating system functions of task management and resource allocation. As in the actual system, the performance of these host functions is transparent to the processing of the application.

In order to make the model of the host environment machine independent, several conventions were employed. When scheduling was required (job, disk, or I/O), a FIFO queue was assumed for that host module. Resource allocation algorithms for TP and IOCS needs were also assumed to be FIFO. Any algorithm, however, can be substituted by changing only the host scheduling services without effecting the model's application services.

A2.2 The Backend Machine

The particular computer system selected by USACSC for the backend machine is the PDP-11/70. As a result of this prior decision a more

detailed and specific model could be built. The hardware configuration for the existing USACSC AIRMICS PDP-11/70 was modeled. Secondary storage characterization conformed to the IDMS a fixed page environment, i.e., the physical data records of the data base with a page size of 1,024 bytes and a direct file organization. Figure A2.2-1 shows the hardware configuration modeled.

The crucial aspect of modeling the backend machine for this project was to identify and characterize the interactions between the PDP-11/70 operating system and the IDMS applications. The main application oriented functions to be coordinated by the resident PDP-11/70 OS are illustrated in Figure A2.2-2. The internal functions of the backend machine are essentially the same as the generalized host except for UNIBUS management. The concept of the UNIBUS (a bidirectional data path through which all resources, including CP and memory, communicate) is central to PDP-11/70 processing. But its effect on processing is transparent to the application. The backend PDP-11/70 has an interface in the host. Because the PDP-11/70 is not dedicated to data base applications, and because USACSC has plans to support more than one data base per backend machine, each invocation of IDMS is treated as a separate subtask. The problem of concurrent updating by multiple users, however, forces the need for a single entry point or interface, called CAMP (Central Access Monitor Program). Even though entry to IDMS is controlled via the application, the actual operation of IDMS necessitates many requests for task management, resource allocation and I/O control. The model of the PDP-11/70 operates, as was assumed for the model of the host machine, in a reactive mode.

The object of the IDMS processing, retrieval of the specified physical record, necessitated the scheduling and use of I/O hardware which comes under the preview of the PDP-11/70 analog of the host's disk I/O module. Passing retrieved records to the host requires the aid of the OS to handle the sequencing of completed IDMS tasks and de-allocation of resources,

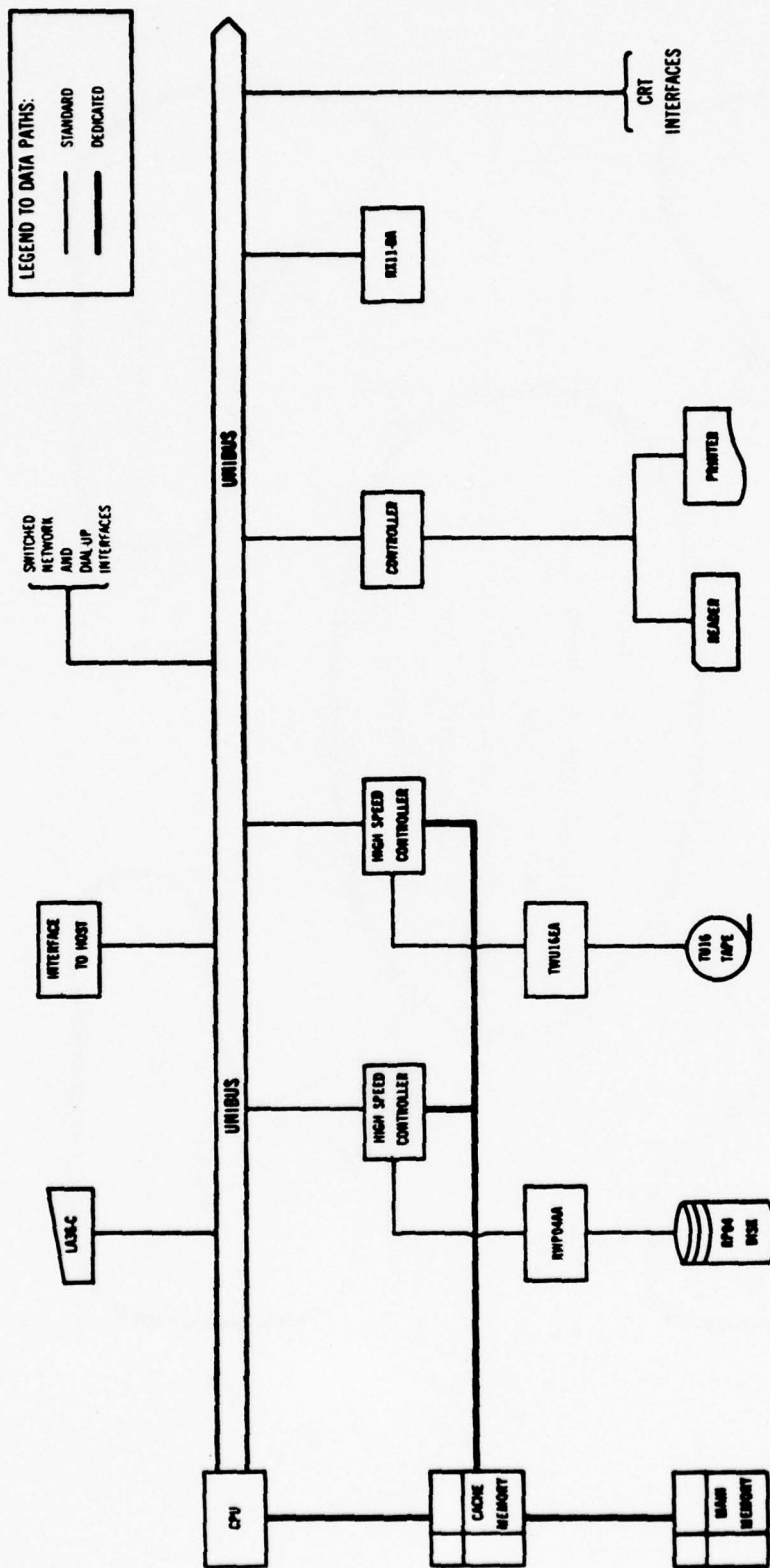


Figure A2.2-1 USACSC PJP 11/70 Configuration

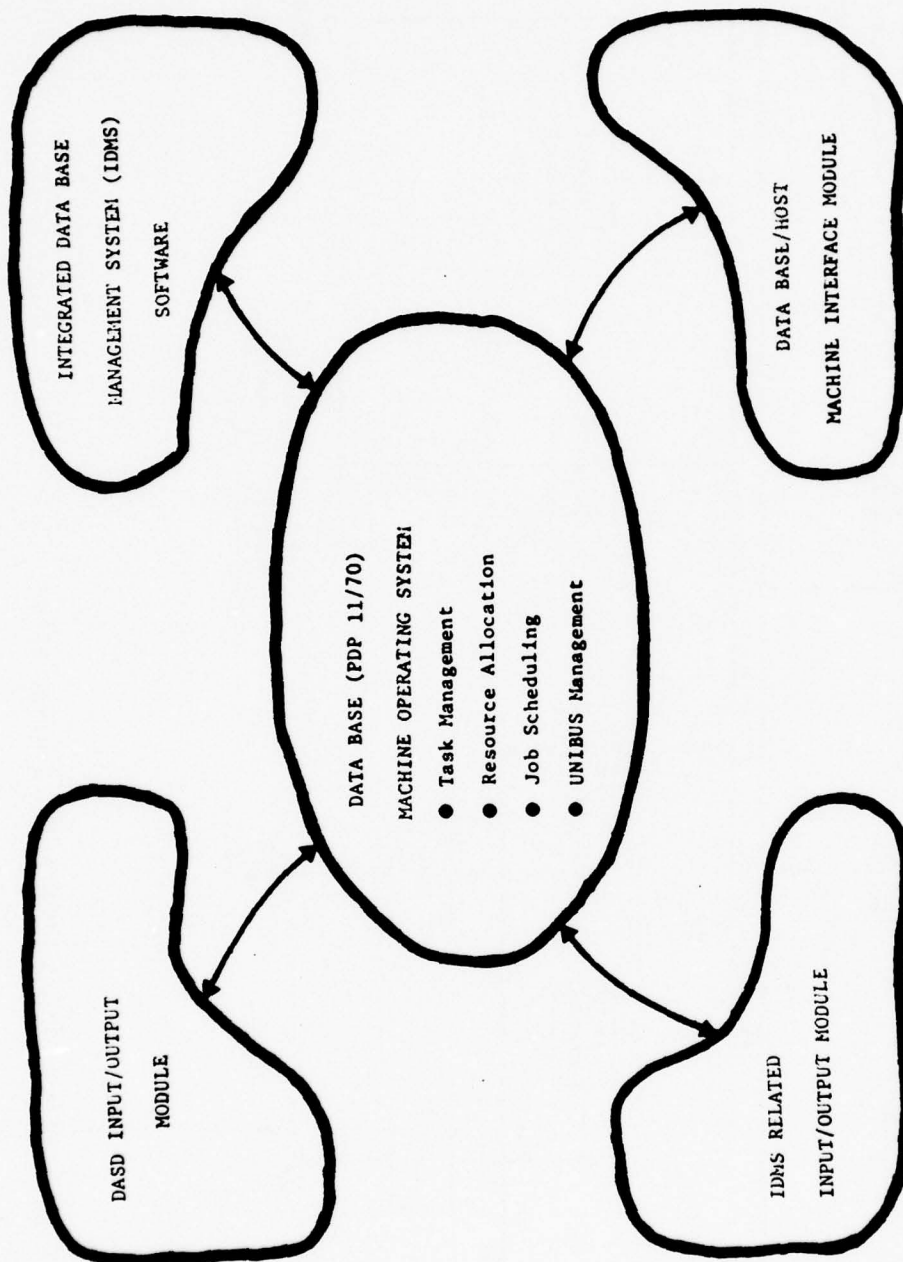


Figure A2.2-2 Modeled Data Base System Resident Software

The job and task scheduling disciplines employed by PDP-11/70 OS are priority queues based on the immediacy of use of the UNIBUS and were so modeled. All tasks performed by the PDP-11/70 priority scheme resulted in, essentially, FIFO queues.

A2.3 Integrated Database Management System - IDMS

The Integrated Database Management System is a network oriented implementation of a subset of the 1971 CODASYL Database Task Group specifications. The following activities were involved in making the conversion from sequential, batch processing to the IDMS on-line integrated database. First, the architecture of the integrated structure of the database was determined. Each of the files identified in Figure A2.3-1 became (in IDMS terminology) "record types" whose interconnections are represented by "sets". Second, the COBOL application programs for the six transaction types was rewritten to reflect IDMS processing. The major impact was in the incorporation of a centralized description of the database and in changing the I/O statements to standard IDMS Data Manipulation Language (DML) commands.

IDMS performs the following functions in response to an application DML. First it validates the request using the sub-schema definitions. Next it translates the request into a virtual page reference, determines if the page is present in its internal buffer, and requests the page from the underlying file management system when it is not already present. Finally, it places the subschema record into the application's work area and returns control to the application.

IDMS has its origins at the B.F. Goodrich Company where it was initially developed. The Cullinane Corporation (Wellesley, MA) now maintains product responsibility. IDMS operates on IBM 360/370 hardware. In addition, it is available through the Digital Equipment Corporation as DBMS-11 which operates on PDP-11/70 hardware.

The following paragraphs briefly introduce the major IDMS system components which were incorporated into the IPSS model (CUL76):

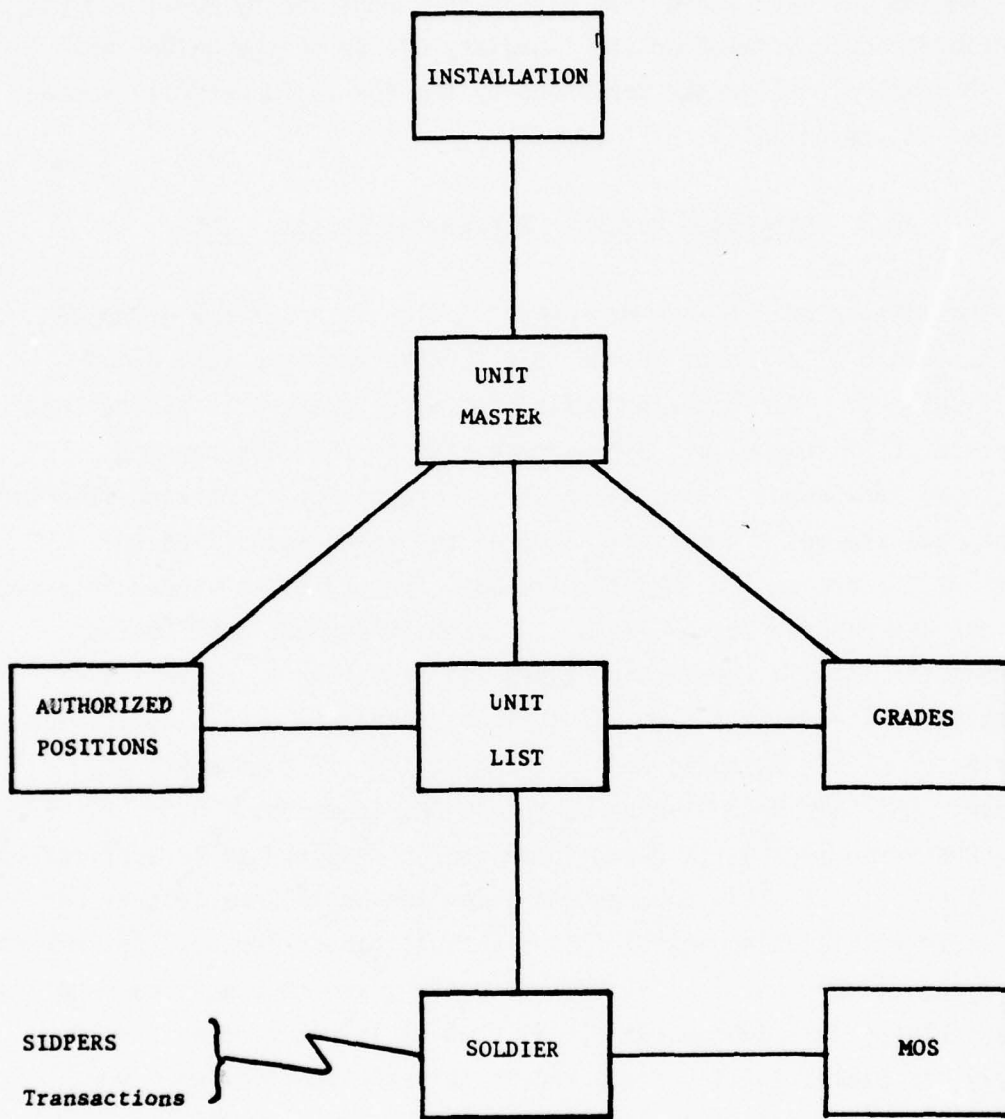


Figure A2.3-1 Modeled IDMS Schema for SIDPERS Data Base

Data Description Language (DDL) -- A facility for describing data independent of the programs that process it. Within IDMS the DDL consists of three parts:

1. Schema DDL -- provides a complete description of the database including names and attributes of all data items, records, sets, and areas.
2. Device Media Control Language (DMCL) -- provides control over all the disk files which are to be on-line.
3. Subschema DDL -- provides a description of those data items, records, sets and areas to be used by one or more application programs.

Data Manipulation Language (DML) -- A facility for accessing base. It provides a wide range of commands for requesting database services which store data in the database or returns data to the user's program.

Schema Description for SIDPERS

The SIDPERS test database was designed to support the six transaction types identified earlier. The overall system architecture reflects twelve record types organized into four IDMS areas. These areas reflect functional differences in the data that is stored and its use. PERSONNEL area contains the individual soldier and unit records of an installation. The MSCAREA contains diagnostic and error messages. The ADAREAL and ADAREA2 Areas contain additional personnel data. The overall structure is shown in Figure A2.3-2.

The record types belonging to each of the areas are identified in Table A2.3-1. The record occurrence and record size attributes of each of the record types are also identified in this table. The personnel area is of primary interest since most of the SIDPERS transaction activity references it record types and sets. This area consists of 3,909 records distributed among 7 record types. Although the GRADE records account for 68% of the total, the IDENT records (5.2% of the

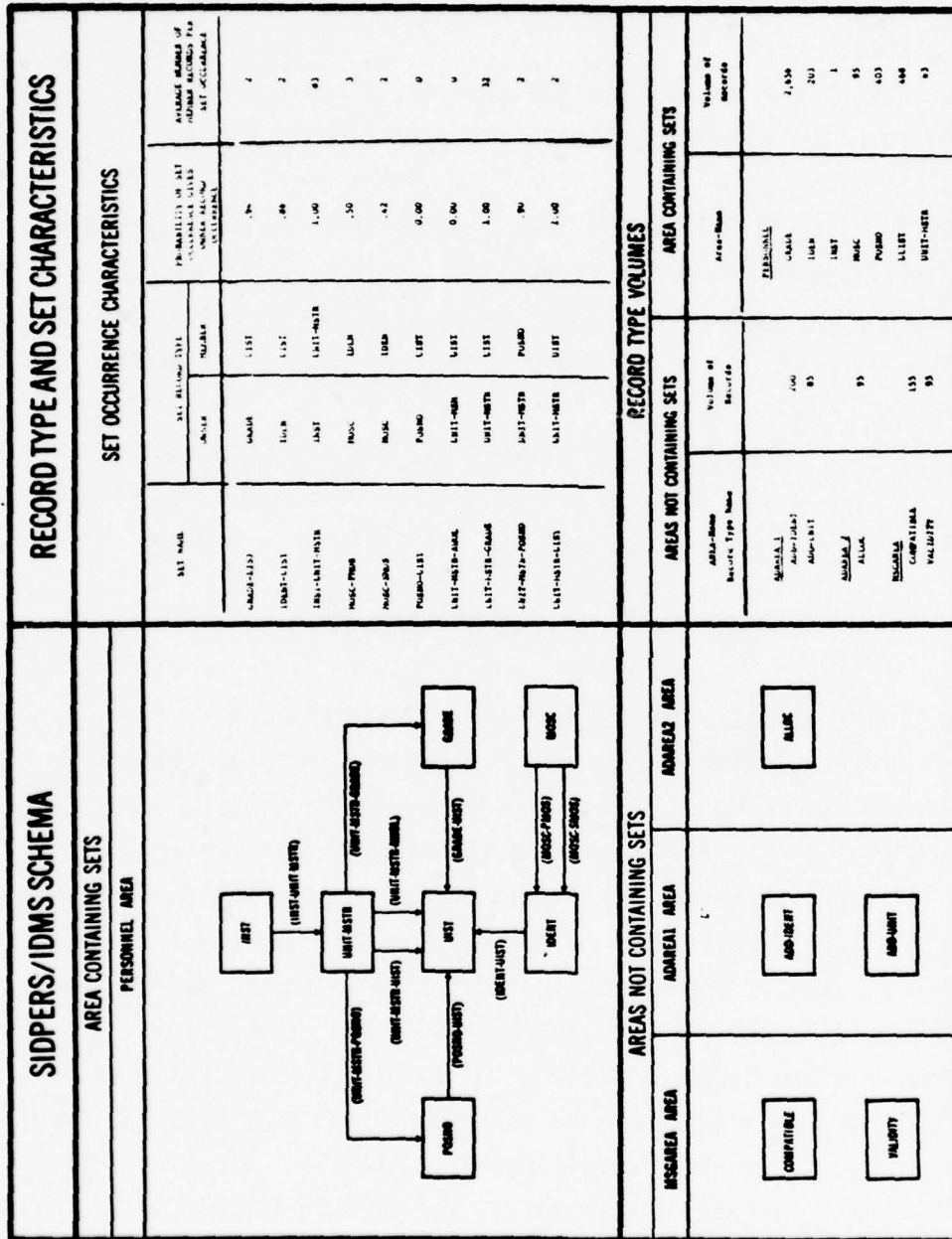


Figure A2.3-2 SIDPERS Data Diagram

Table A2.3-1 Record Occurrence Characteristics

Record Type Within Area	Record Occurrence		Record Size	
	Number	% of Total (within area)	Length* (Bytes)	Percent Used of Space Available Within Area
<u>PERSONNEL AREA</u>				
INST	1	0.0	48	.01
UNIT-MSTR	83	2.1	52	.79
IDENT	203	5.2	120	4.12
UIST	468	12.0	84	6.82
GRADE	2,656	68.0	20	11.78
POSNO	403	10.3	140	9.45
MOSE	95	2.4	236	3.67
<u>ADAREAL AREA</u>				
ADD-IDENT	200	70.2	284	23.13
ADD-UNIT	85	29.8	552	18.85
<u>MSGAREA AREA</u>				
COMPATIBLE	155	62.0	296	14.93
VALIDITY	95	38.0	512	15.65
<u>ADAREA2 AREA</u>				
AALOC	95	100.0	96	.15
*Includes IDMS Overhead of 8 bytes per record.				

total) are most frequently accessed. Most of the transactions are oriented toward updating an individual soldier's data and thus the most common entry point is the IDENT record. This is accessed via the CALC set using the soldier's last name.

Figure A2.3-2 identified the organization of the personnel area record types into sets. Ten distinct set types (excluding the CALC sets) are employed to provide the inter-record traversal capabilities required by the SIDPERS transactions. Table A2.3-2 identifies these sets and displays their characteristics. The sets are identified by both an "external" name which is the set name used in the SIDPERS subset, and also by an "internal" name which is the set identifier within the IPSS models. For modeling purposes, two data items per set type were derived from analysis of the sample data base. These were:

1. the probability of an owner record type having a non-null set of member records associated with it (i.e., the probability of a set occurrence), and
2. the number of member records in a set occurrence for each set type.

Table A2.3-2 also identifies these data items for each set of the PERSONNEL AREA. Although the table shows the average number of member records, for modeling purposes probability distribution functions were used to more accurately characterize the database.

The records of the SIDPERS data base reside within specially formatted IDMS pages. These IDMS pages are virtual pages since they represent the data as known by IDMS and as such need not be reflective of secondary storage record organization. The page size itself is a function of the IDMS processing environment. Interactive SIDPERS with IDMS operating on IBM 370/165 equipment uses a page size of 3,156 bytes, whereas the page size for IDMS operating on a PDP-11/70 is projected to be 1,024 bytes. The test version of the SIDPERS data base has been loaded onto IDMS pages as shown in Table A2.3-2. The Personnel Area occupies 200 pages and is 36.6% full. The IPSS simulation model of this data base assumes a PDP-11/70 environment as the number of pages and page range was adjusted to keep the percent utilization the same while changing the page size to 1,024 bytes. This data is also reflected in Table A2.3-3.

Table A2.3-2 Set Occurrence Characteristics

Set Name		Record Type Participation		Probability of Set	Average Number of
External ID	Internal	Owner	Member	Occurrence Given Owner Record	Member Records
INST-UNIT-MSTR	INUM	INST	UNMA	1.0	83.0
UNIT-MSTR-POSNO	UMPOS	UNMA	POSN	.9	5.7
UNIT-MSTR-UIST	UMUI	UNMA	UIST	1.0	2.6
UNIT-MSTR-AWOL	UMAW	UNMA	UIST	0.0	0.0
UNIT-MSTR-GRADE	UMGR	UNMA	GRADE	1.0	32.0
POSNO-UIST	POUI	POSN	UIST	0.0	0.0
GRADE-UIST	GRUI	GRADE	UIST	.94	1.6
IDENT-UIST	IDUI	IDEN	UIST	1.0	3.0
MOSC-PMOS	PMID	MOSC	IDEN	.86	1.0
MOSC-SMOS	SMID	MOSC	IDEN	.57	2.0

Table A2.3-3 Page Utilization Within the SIDPERS Areas

Area	Page Size = 3,156 Bytes			Page Size = 1,024 Bytes	
	Number of Pages	Page Range	% Utilized	Number of Pages	Page Range
Personnel	200	1-200	36.64	617	1-617
ADAREA1	80	201-280	41.98	246	618-863
MSG2EA	100	301-400	30.58	308	925-1233
ADAREA2	20	281-300	.15	60	864-924

DML Description for SIDPERS

A previous study of a VIABLE transaction has specified the number and type of IDMS DMLs that are issued as well as their sequence within the transaction (SHA77). These DML have been classified by type as shown in Table A2.3-4.

The function of each of these DML Commands is as follows:

FIND: Locates the selected record in the data base and returns its identity to the application program.

OBTAIN: Locates the selected record in the data base and returns the record as well as its identity to the application program.

MODIFY: Causes the specified record to be updated.

STORE: Causes the specified record to be stored in the data base.

ERASE: Causes the specified record to be physically removed from the data base.

CONNECT: Causes the specified record to be inserted into an identified set.

DISCONNECT: Causes the specified record to be removed from the identified set.

Table A2.3-4 SIDPERS IDMS DML Types

Category	DML
FIND/OBTAIN	FIND record-name RECORD FIND NEXT record-name RECORD of set-name SET FIND OWNER RECORD OF set-name SET OBTAIN record-name RECORD OBTAIN NEXT record-name RECORD of set-name SET OBTAIN OWNER RECORD OF set-name SET
MODIFY/STORE/ERASE	MODIFY record-name RECORD STORE record-name RECORD ERASE record-name
CONNECT/DISCONNECT	CONNECT record-name TO set-name DISCONNECT record-name FROM set-name

Within the IPSS model of VIABLE SIDPERS, each IDMS DML is represented by statements which invoke a lower level service and wait for its completion. Four parameters that are required to pass the DML information. These are:

1. DML Type (FIND, OBTAIN, etc.);
2. Qualifier (NEXT, OWNER, etc.);
3. Record Type Identified (ADDI, IDEN, etc.); and
4. Set Type Identifier (GALC, GRUI, etc.),

A3. MODELING USING THE IPSS DESIGN FACILITY

A3.1 A Functional View of the Modeling Design

The purpose of the host and the backend submodels is to provide specific computer system environments for the SIDPERS application and IDMS submodels. The relationship between these submodels can be seen in Figure A3.1-1. The SIDPERS submodel was constructed by characterizing the functional SIDPERS processes required by a SIDPERS transaction from the time it enters the host system until the completion of processing. These functional processes included such activities as input validation, backend interaction, and diagnostics transmissions. The SIDPERS submodel characterized these processes, however, in a computer system free environment, concerning itself only with the impact of these various processes on SIDPERS processing. Similarly, the IDMS submodel was constructed with no assumed computer system environment. Functions such as buffer and page management, DML invocation and host interfacing were characterized as part of the IDMS submodel, at the level of complexity of the IDMS request. Both the host and backend computing environments were treated as black boxes in these models. Thus, the initial preliminary model had the SIDPERS submodel interacting with the IDMS submodels (arrows A & B).

The purpose of the host and backend submodels were to provide a specific computer environment, including hardware, software and data elements, within which the SIDPERS and IDMS submodels could operate. The operating system functions that had been assumed to be operational in the application oriented submodels, but had been aggregated, were now characterized. For example, transmitting SIDPERS diagnostic messages to the user at the terminal involves accessing the message from the host's disks (including I/O scheduling, disk allocation and I/O operations) and the actual transmission (which requires TP scheduling, TP line allocation and I/O operations). With respect to the IDMS submodel, the function of page replacement, for example, involves task scheduling, resource allocation, and memory management. It is the characterization of these support functions that the host and backend submodels provided.

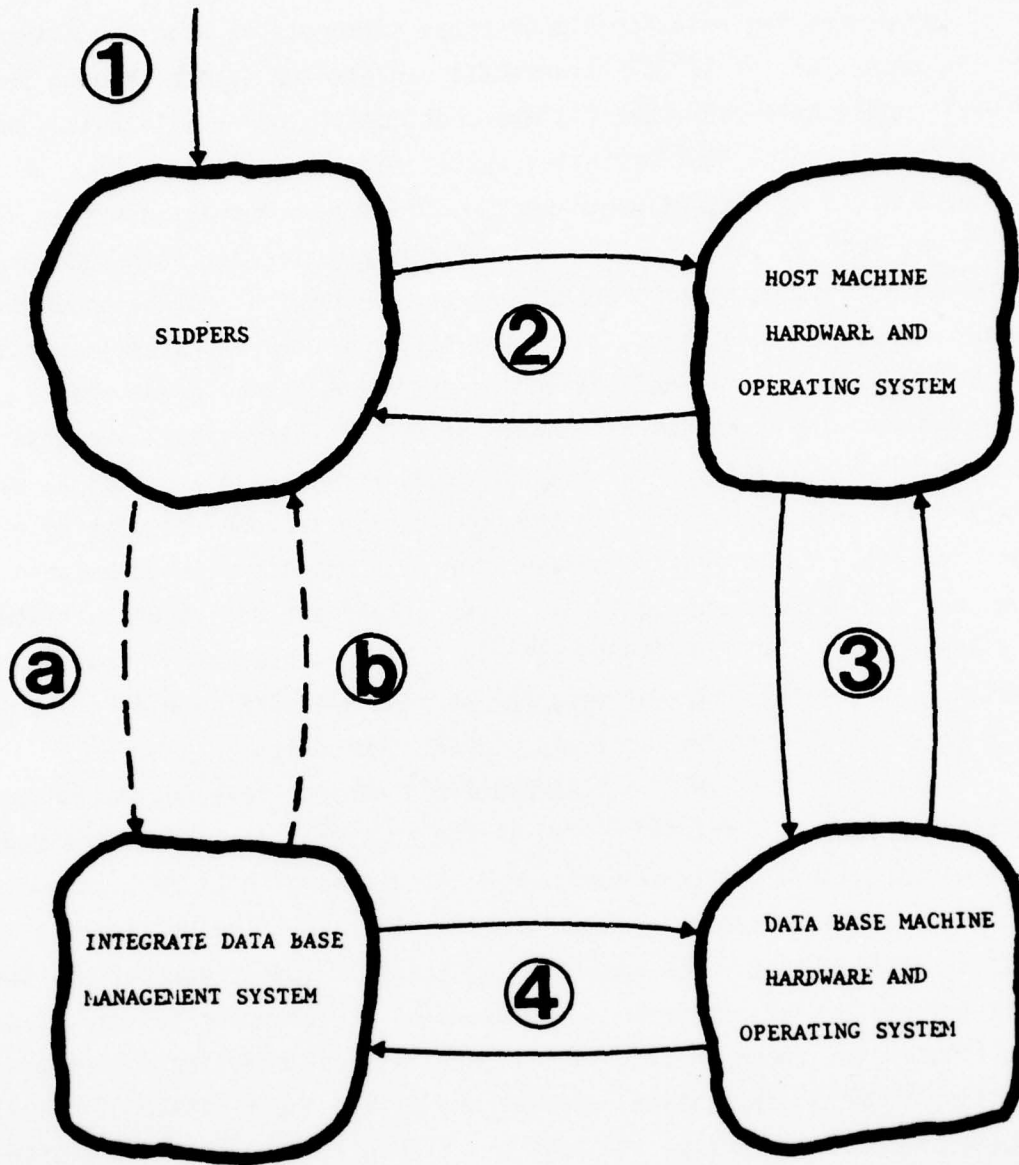


Figure A3.1-1 Relation of Submodel to Total Model

The processing relationship of these submodels is also demonstrated in Figure A2.1-1. A SIDPERS transaction enters the system through the SIDPERS application submodel (1) and continues to operate in this submodel until it requires a host operating system function. At that time it triggers off a host function to perform the needed operating system function, such as task, job or resource management (2). Both submodels continue processing other transactions and events that may be occurring simultaneously. In essence, then, the host submodel operates in parallel with the SIDPERS model, satisfying its operating system needs and coordinating the execution of concurrent SIDPERS processes. The host submodel is essentially in a reactive mode, residing in a continual wait state until awakened by the SIDPERS application submodel requesting a host service. Communication between the host is accomplished between the host and backend submodel (3). The backend submodel then initiates the appropriate IDMS service to satisfy the DML request sent from the host (4). From then on, however, the backend and IDMS submodels operate in a parallel as does the host and SIDPERS submodels.

Previous studies of the VIABLE/SIDPERS project have not designated a particular type of computer system as the host machine. Thus, the model that was developed has been designed to be a general host machine model, including those operating system functions that would be necessary on any machine that as host. The structure of the host model, however, is such that the particular software and hardware of any computer system can be substituted for those in the initial model without modification of the existing SIDPERS application submodel and little modification of the host submodel. Currently, the test bed SIDPERS/IDMS system resides on an IBM 370/165, both the host and the backend. Thus, although the structure of the host submodel in this project is general, the particular elements are specific to the USACSC HCC IBM 370/165. The backend machine, however, has been chosen to be a PDP-11/70 and the resultant backend model will reflect the operating system philosophies of this system.

In modeling these two types of environments, a similar approach was taken. It is assumed in the IPSS methodology that operating systems,

in general, perform certain types of services regardless of the particular machine. As Figure A3.1-2 illustrates, although an operating system may be servicing many different applications simultaneously, the coordinating and servicing of these applications is completed through a stage set of underlying operating system functions. The types of system functions performed are job scheduling, task management, resource allocation, I/O supervision and telecommunication control. Thus, both the host and the backend submodels were constructed by characterizing these functions in general and then particularizing them to the specific machine, IBM 370/165 and PDP-11/70. The structure of these submodels, however, is such that at any later time the hardware and/or software can be changed to represent another system.

It should be emphasized that these characterizations of the computer systems were not dependent on the applications to be run on them, i.e., SIDPERS and IDMS. A future application of these submodels is to change the loading on the computer submodels to include other, possible other SIDPERS, applications to determine the effect on SIDPER/IDMS processing. Thus, the scheduling and allocation policies that are coded for these submodels do not reflect the particulars of either SIDPER or IDMS and are, essentially, general task and resource managers.

A3.2 IPSS Model of Host Computer

The approach taken in using IPSS to model this computer system is to first model the static elements of the system, the hardware and the data and then model the dynamic software elements. The host computer that was modeled was the USACSC HCC IBM 370/165. To construct the hardware configuration of this machine, only those elements that were impacted by the SIDPERS application were modeled. Specifically, they were the secondary storage devices upon which the IDMS error message files were kept, the backend interfacing equipment, external terminals through which USACSC personnel invoke and interact with SIDPERS, and the connecting telecommunication equipment. Other hardware typical of an IBM

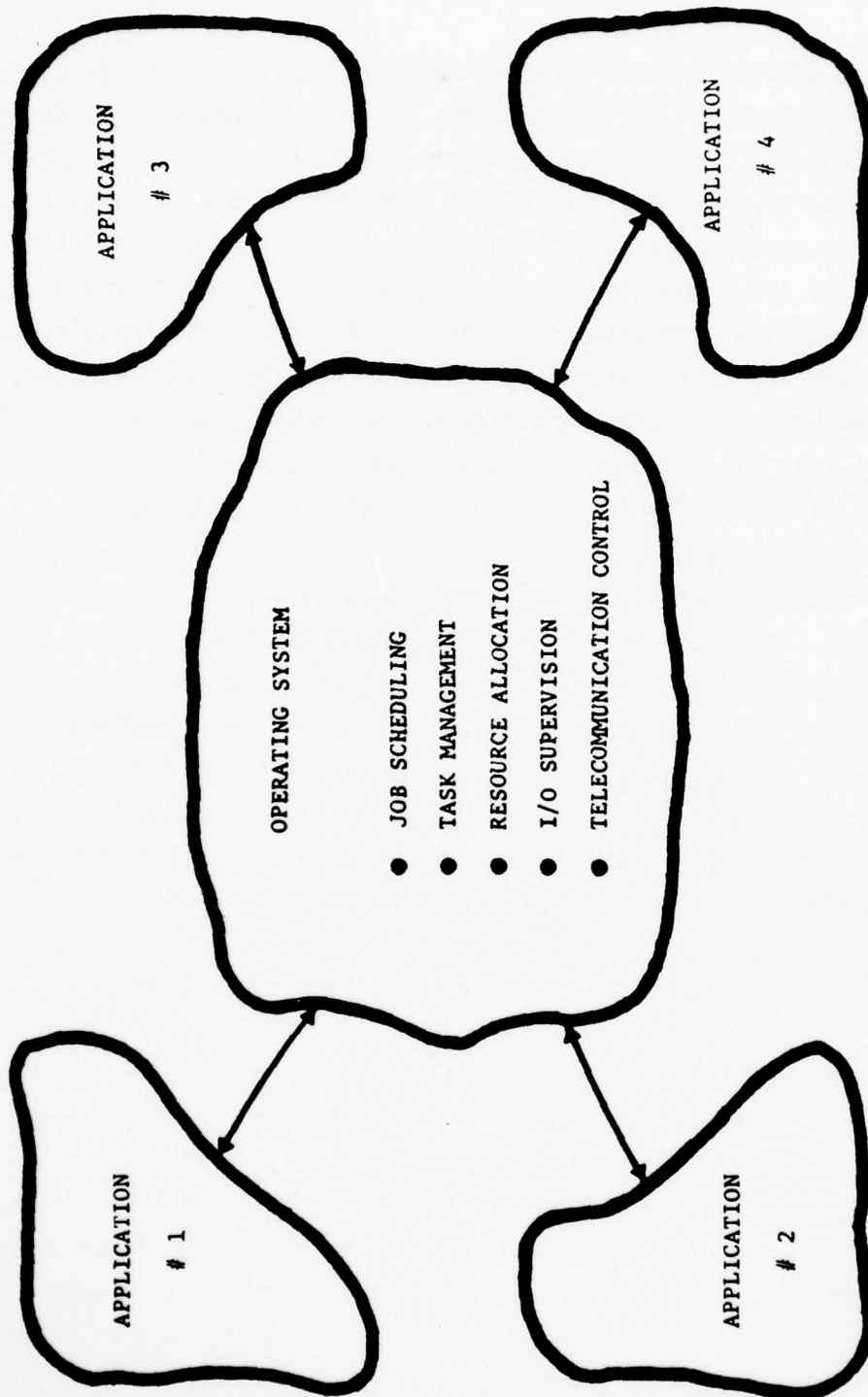


Figure A3.1-2 Modeled Operating System Functions

370/165 were not included as they would not be referenced in this model, but at a later point in time when other USACSC applications are added to the host model, additional hardware can be easily added.

Figure A3.2-1 illustrates the initial hardware configuration that was assumed for the host, IBM 370/165 computer. The particular secondary storage devices assumed were a bank of eight IBM 3330 DASDs, their associated control units, and single channel. A single CPU with 2 megabytes of main memory were assumed to characterize the host computer. The external terminals were assumed to be similar to IBM 2741 and the communication line between the host computer and these terminals was assumed to be a 2400 baud line. It should be noted that in an IPSS model the characteristics of anyone of these hardware devices can be modified without requiring the modification of the rest of the model. Thus, experiments evaluating different combinations of devices can be easily made.

The second step in constructing an IPSS model is to define the secondary storage data files. With respect to the SIDPERS application, the only data files associated with the host computer are the files of validity and compatibility error messages. The validity error file is accessed once for a compatibility error by the SIDPERS application software. Although there was little data on the size and internal characteristics of these files the following was assumed: logical record length of 80 bytes, physical record length of 4000, (to fit into the VM OS page size of 4096), 1000 validity records, 200 compatibility records and placement of the two files on the same volume.

The crucial aspect of this and any IPSS model is the characterization of software processing. It has been mentioned that the operating systems of both the host and the backend computers were assumed to be in a reactive mode, responding to the needs of their respective applications. Thus, in order to determine the operating system functions that needed to be modeled in the host (and backend) IPSS submodels, the processing requirements of the SIDPERS (and IDMS) applications must be examined. The requirements for the host can be seen by examining the functional

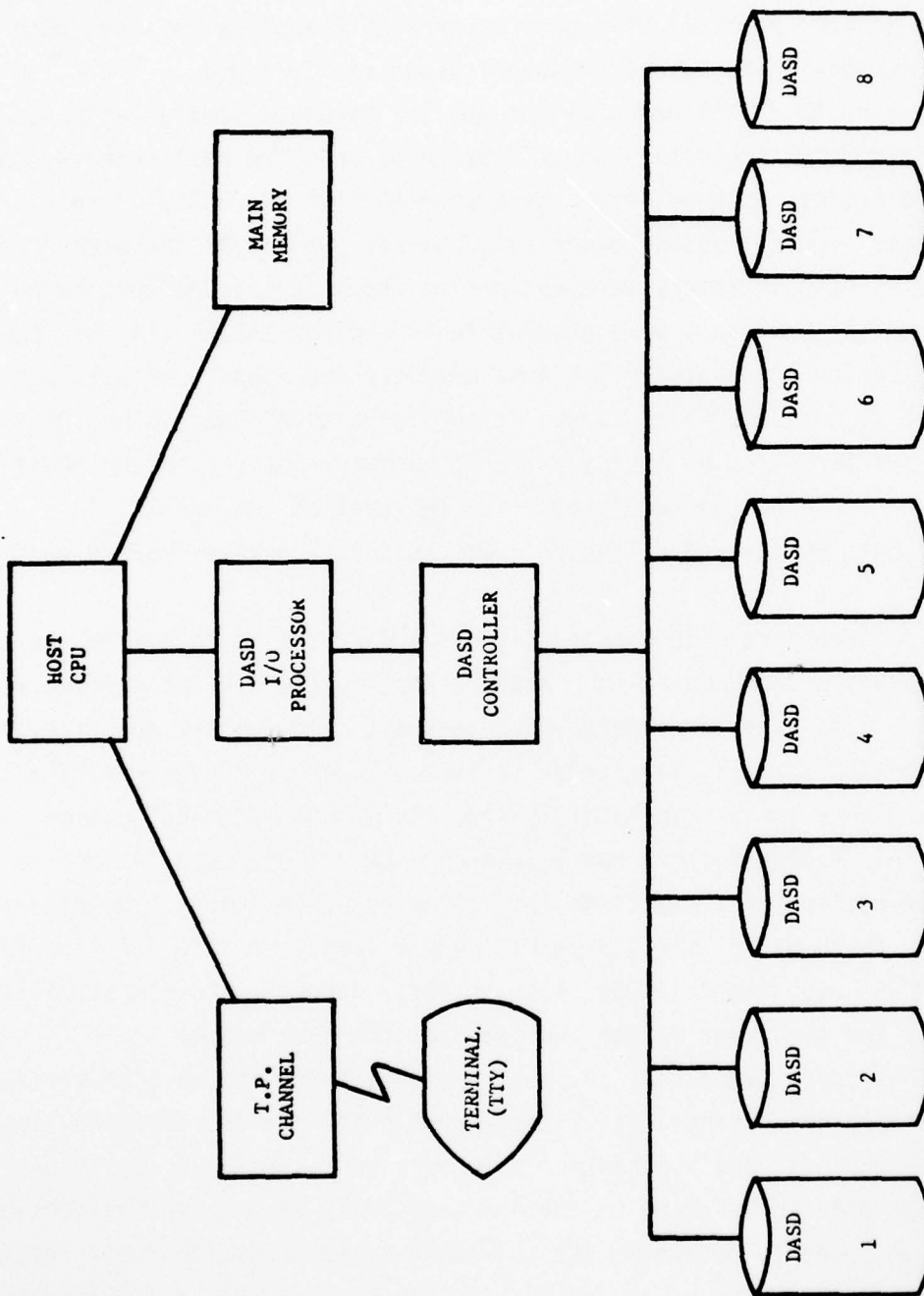


Figure A3.2-1 Assumed Host Hardware Configuration

steps that a SIDPERS transaction goes through that were detailed in Section A1. Summarizing, there were five basic steps as shown in Figure A3.2-2: log-on, SIDPERS scheduling, data input and validation, IDMS/backend interaction, and log-off. At each of these steps particular operating system functions are required at varying levels of complexity.

To satisfy these needs, certain salient operating system functions were identified. These were: the scheduling and transmission of data to/from the user, the scheduling of a SIDPERS application within the host, the communication requirements of interactive data input and the necessary resource allocation and I/O scheduling to access the error message files and the scheduling transmission and repetition of messages to/from the backend computer. Underlying all these processes is the synchronization and coordination of the SIDPERS application functions and the host operating system functions which may be executing in parallel, as provided by general task management and resource allocation.

The relationship between the IPSS services that characterized the SIDPERS application and the IPSS services that performed the host functions are shown in Figure A3.2-3. Being that SIDPERS is an interactive system, constant communication between the user and the SIDPERS system, through the host, must be maintained. Thus, each of the five SIDPERS application steps, in some way, require the scheduling and transmission of information to or from the user. In the resulting IPSS submodel this interface was handled by the service XTCAM, which acted as scheduler for TP to the user. It was awakened by a SIDPERS module and returned to the wait state. The requesting SIDPERS module was then in control of the use and disposition of the TP line. The second step was to specifically schedule the SIDPERS application for execution in the host system. This was performed by the IPSS service JOBSCH which controlled the allocation of main memory and central processor and placed the incoming SIDPERS application in the ready queue. It is general, though, so that other non-SIDPERS jobs could also be scheduled.

The third SIDPERS application step, data input and validation, was characterized in the SIDPERS application submodel by the services VALID,

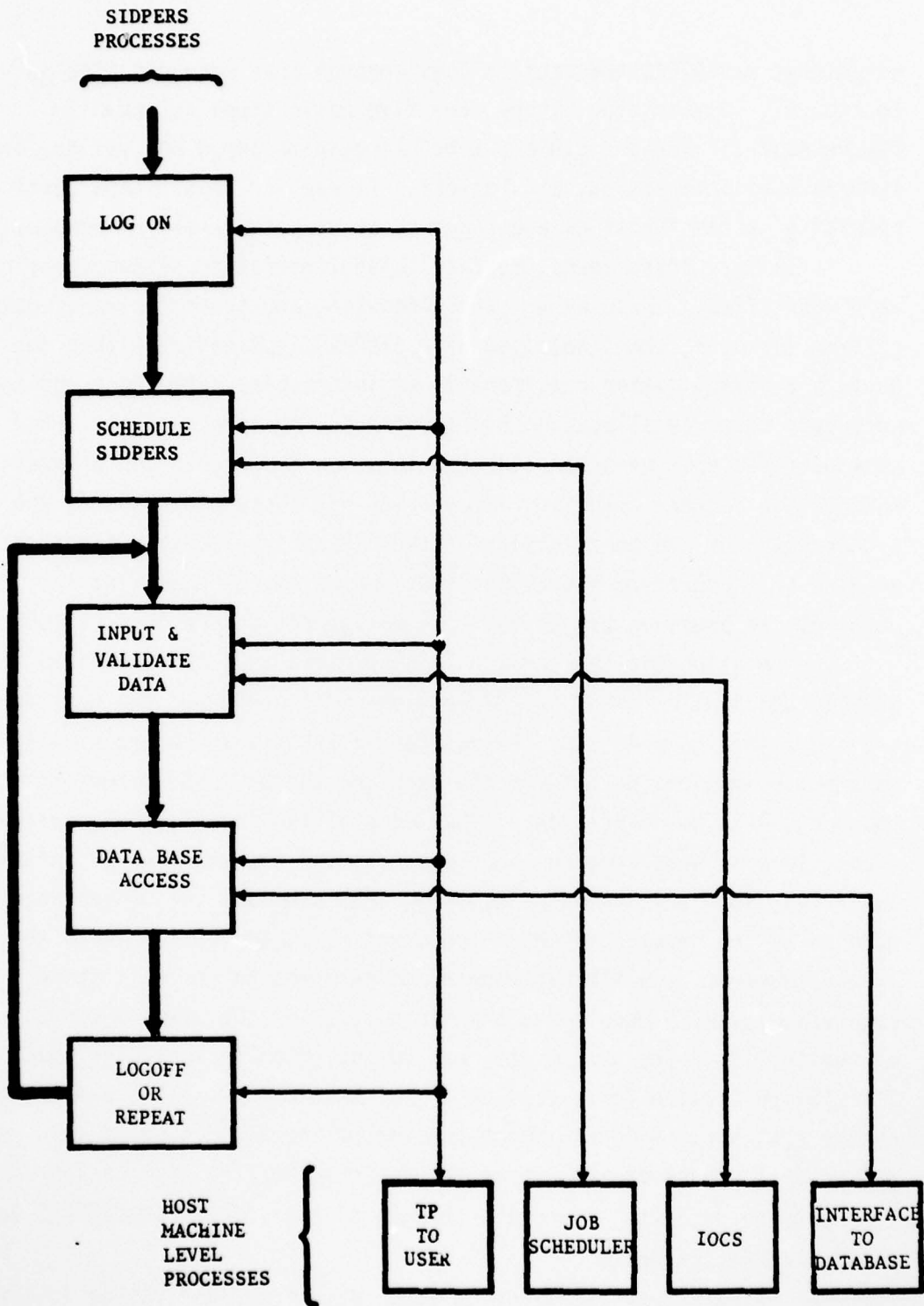


Figure A3.2-2 Relation of Host Machine Processes to SIUPERS Processes

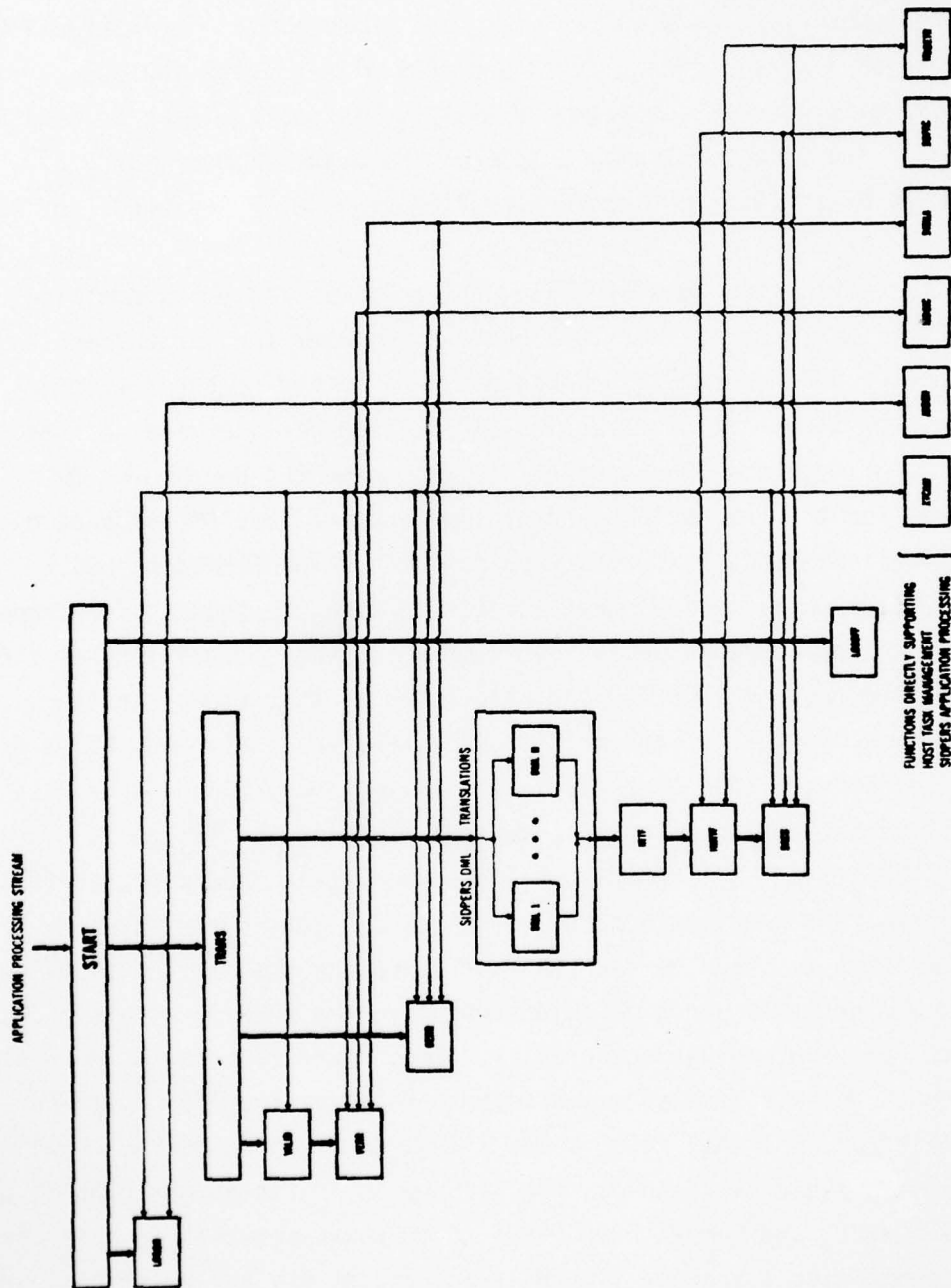


Figure A3.2-3 IPSS Realization of Host/SIDPERS Model

AD-A069 543

OHIO STATE UNIV RESEARCH FOUNDATION COLUMBUS
A SIMULATION LANGUAGE FOR EVALUATING INFORMATION PROCESSING SYS--ETC(U)
SEP 78 T 6 DELUTIS

F/G 9/4

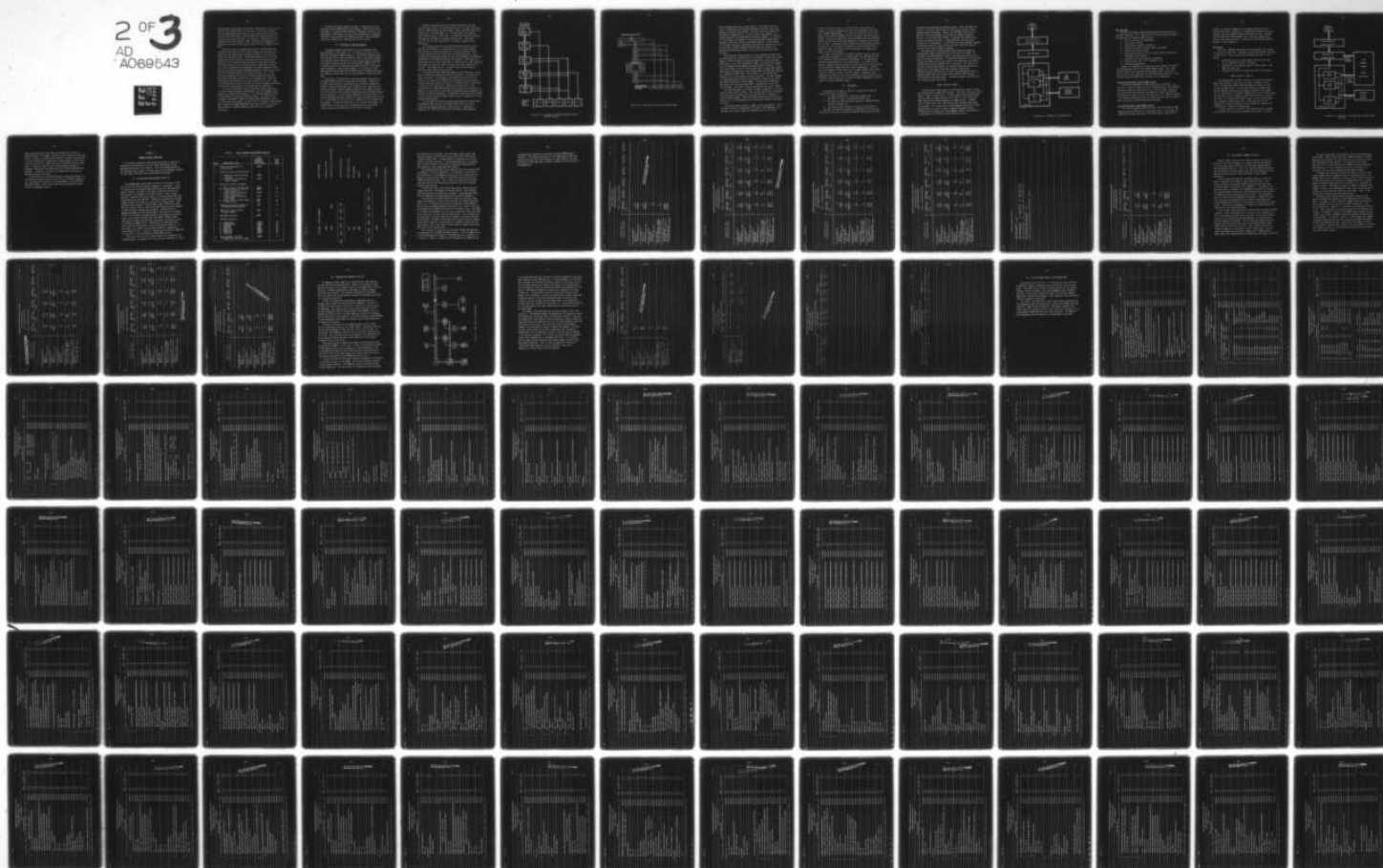
DAAG29-77-G-0203

UNCLASSIFIED

ARO-15356.1-A-M

NL

2 OF 3
AD
A069543



VERR and CERR, representing, respectively, input process, validity error accessing and compatibility error accessing. Each of these modules involves communication with the user, thus, necessitating interaction with the host service XCAM. VERR and CERR also require the access of the host disk storage subsystem. This process is regulated by a disk scheduler and device allocator, XHDSK. The actual I/O commands are performed by an IOCS routine labeled DISKA in the IPSS submodel of the host.

The fourth step in a SIDPERS application is the DML processing, including the interface with the backend computer for the purposes of sending the DML request to the backend, and receiving the retrieved record from the backend IDMS. The SIDPERS submodel services for each transaction type and the service INTF determine the particular DML required for the transaction and perform any pre/post DML processing. The actual interface to the backend computer is performed by HINTF. This service characterizes the packaging of the DML request for transmission to the backend, buffer management (BUFALC), scheduling of the TP line to the backend (HOBETR) and allocation of that line. In the preliminary submodel of the host, the backend processing was black boxed and service BACK characterized this implementation. BACK also controlled the deallocation of resources (buffers and TP line). The final SIDPERS step of log-off parallels the log-on procedure, performing allocation (now deallocation) of resources and job scheduling.

In order to characterize the host operating systems functions as they would normally operate, asynchronously, the IPSS submodels had to operate (or simulate) asynchronously. The interface between the host and the backend is also an asynchronous process. The IPSS built-in statements of POST/WAIT EVENT allow XTCAM, XHDSK and HOBETR to remain in a wait state until needed. The IPSS built-in statements FIND IN QUEUE, WAIT IN QUEUE, and REMOVE FROM QUEUE allow these services to maintain a queue containing suspended SIDPERS transactions via any queueing discipline.

Testing of this host submodel was done in conjunction with the previously tested SIDPERS application submodel. Since the host model was always in a reactive mode, completely independent testing was not appropriate. Additionally, the use of this submodel was to determine the effect of host support for SIDPERS processes to that the SIDPERS application submodel was a necessary part of this testing.

A3.3 IPSS Model of Backend Computer

The backend computer was modeled by first characterizing the great features of that system, i.e., the hardware and data, and then modeling the dynamic software features. Since previous USACSC studies had indicated that the PDP-11/70 would be used as the backend computer, a more specific model than that built for the host could be developed. The hardware configuration for the existing USACSC AERMICS PDP-11/70 at Georgia Institute of Technology was modeled as part of the Systems Reserve Component. Since backend computer is to be used for other applications besides supporting the IDMS system, the complete configuration needed to be modified, not only those elements directly affecting IDMS.

Figure A2.2-1 illustrates the idealized hardware configuration of the USACSC AIRMICS PDP-11/70 system used in the model. Although the physical system itself was unavailable, sufficient DEC documentation existed so that detailed modeling of individual device characteristics could be achieved. Several hardware elements of this model that have significant impact on the software processing and thus software modeling, of this system are: (1) the central UNIBUS data and control path; (2) the two dedicated high speed controllers to disk and tape drives; (3) the use of cache memory by the core and; (4) communication link between host and backend. It should be noted that in this and many IPSS models, the characteristics of any one of the hardware devices can be modified without requiring the modification of the rest of the model.

Defining the characteristics of the particular data sets that reside on this system is the second modeling step. For this application, only the IDMS test data base was modeled. Each of the IDMS records was, essentially, a 1024 byte page (i.e., physical record) with varying highly logical records in each page. The IDMS test data base contained 200 pages and the IDMS file organization was direct. Other data sets for this and other applications can be easily added to these definitions.

The characterization of the backend software processing follows much the same pattern as the characterization of the host software. Both operating systems are assumed to be in reactive models, waiting for and responding to the needs of the resident applications. Thus, to determine the operating system functions to be modeled, the processing requirements of the IDMS applications were examined. A DML request in the backend follows a five step process: request arrival, IDMS scheduling, IDMS processing, physical record retrieval, and record transmission to the host. (See Figure A3.3-1). At each of these steps particular operating system functions are required.

To satisfy functional IDMS requirements for DML processing, certain PDP-11/70 operating system functions were identified. As Figure A3.3-2 shows, these were: the scheduling and transmission of data to/from the host, the scheduling of IDMS within the backend, and the necessary resource allocation, buffer allocation, I/O scheduling and I/O execution to retrieve the desired record. Underlying processes such as the synchronization of IDMS and backend functions, which may be executing in parallel, is provided for in the model by generalized task management and resource allocations procedures. Although CPU scheduling was not included as part of these functions, UNIBUS scheduling was an integral part of the synchronization process.

The relationships between the general DML processing step and their corresponding IDMS services, and between these services and those characterizing backend operating system functions is depicted in Figure A3.3-2. Both the first and the last steps in DML processing

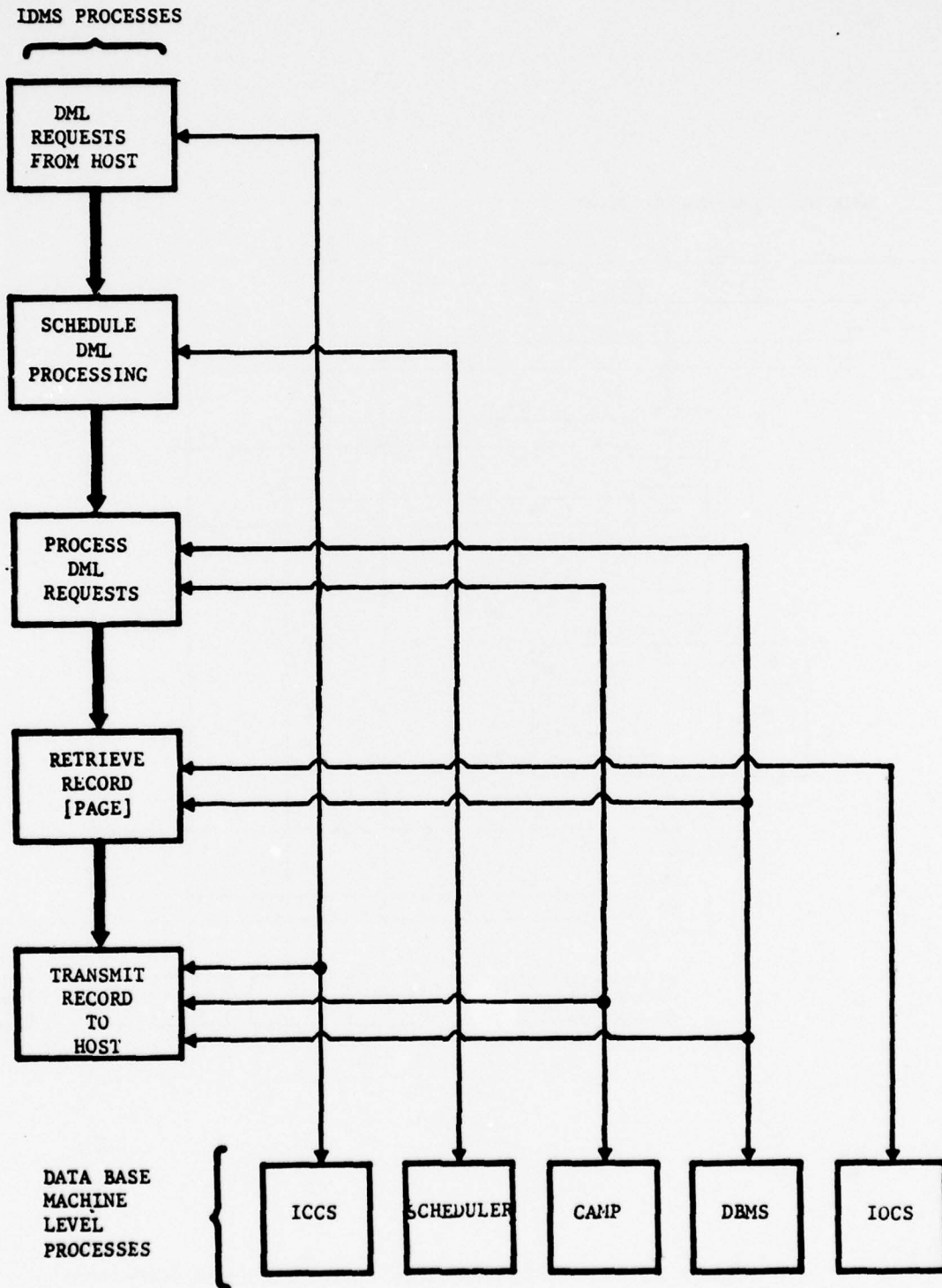


Figure A3.3-1 Relation of Data Base Machine Processes to IDMS Processes

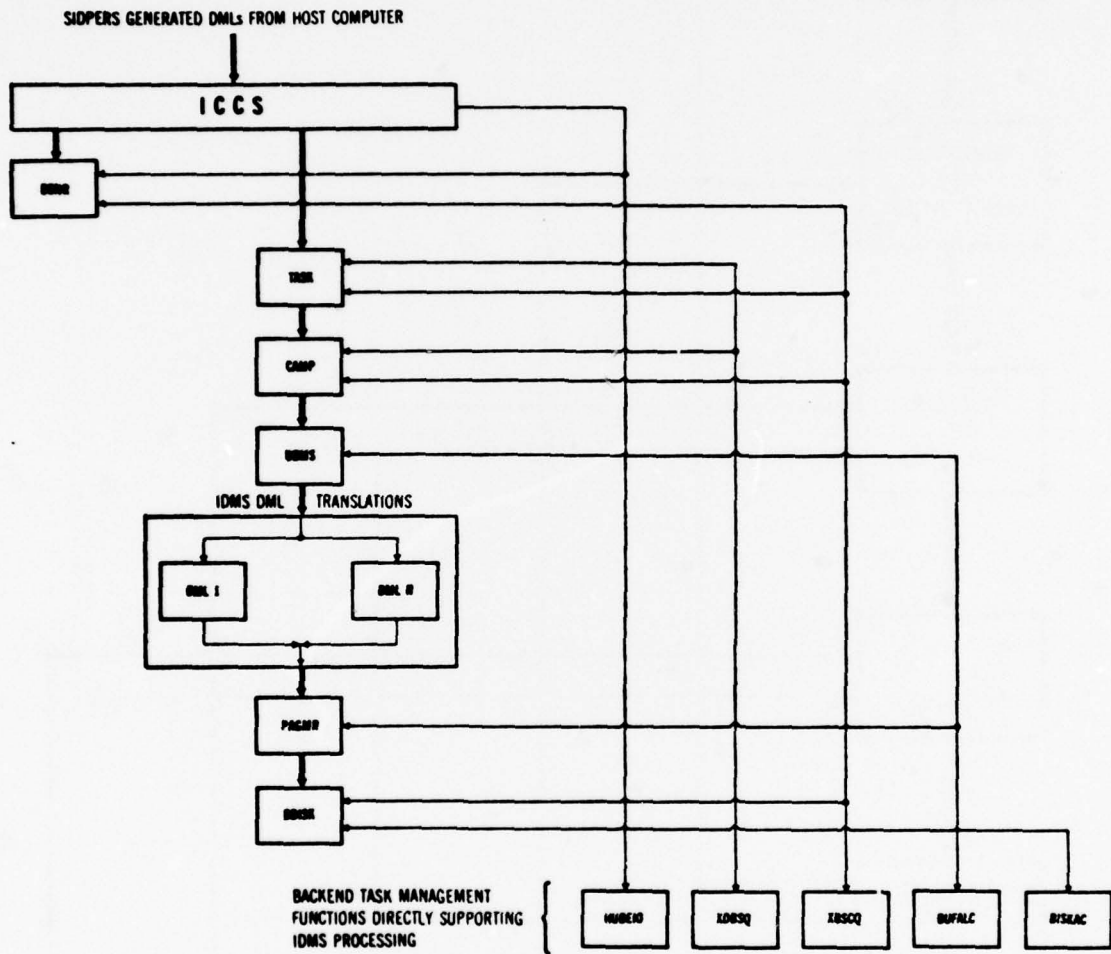


Figure A3.3-2 IPSS Realization of Back-end/IDMS Model

require communication with the host machine. This entails not only a service to characterize the backend interfacing functions such as the message printing and line protocol, but also a service to schedule the intercomputer communications. In the resulting IPSS submodel, this scheduling function was performed by the service, MUB910. It scheduled, in a FIFO manner, all requests to use the host/backend communication interface, from either the host or backend direction. It was assumed for this model that only one message could cross the interface at a time. Once the service HUBE10 scheduled and allocated the interface, the requesting service gained control of the use and disposition of the interface.

The second step in DML processing is the scheduling of IDMS itself. As mentioned before, the PDP-11/70 is not assumed to be dedicated to the support of IDMS, allowing the concurrent execution of other non-IDMS applications. Thus, each invocation of IDMS, via entering DML command, must be scheduled and the appropriate IDMS software allocated. In the structure of IDMS its twin in the host application signals CAMP (Central Access Monitor Program) that a DML requires IDMS, and CAMP then schedules the invocation of IDMS itself. Only one DML can use the data base functions at a time to prevent concurrent updatings. The IPSS submodel for IDMS contained IPSS services corresponding to CAMP and the actual data base manipulation routines (DBMS). The signaling to CAMP that a DML request has arrived and needs DBMS is accomplished via the IPSS service XDBSQ. The actual allocation and deallocation occurs in CAMP.

Once the DBMS module is involved, the DML is processed in the third step via IDMS internal routing. Each of these routines has an analogue IPSS service that accomposes the high level DML into a sequence of primitive DBMS commands. IDMS individually changes these primitives to requests to its virtual data store for particular record(s) which leads to step 4.

The actual record retrieval in IDMS is a two step process. First, the record location is determined with respect to an IDMS page. The currently core resident pages are examined to see if one of them

contains the desired record. If none do, then a new page must be accessed. This page management function is accomplished via the IPSS service PAGMR (page manager). If a page must be accessed, then buffer space must be allocated (BUFACE) and the I/O request must be scheduled (XDSCO). When the I/O operations can be performed, the actual hardware devices for the PDP-11/70 backend computer are referenced via the IPSS service DISKB. Furthermore, DISKB simulates the allocation/deallocation of resources and the actual I/O operations of each and data transfer. DISKB also takes into account the dedicated data lists between main memory and auxiliary storage.

In order to characterize the backend operating system's functions as they normally operate, i.e., asynchronously, the IPSS IDMS and PDP-11/70 submodel also had to execute asynchronously. The host/backend interface (through HUBE10) is also asynchronous. To accomplish this behavior, communication with operating system functions was characterized via the posting of common IPSS Event signals (semaphores). The IPSS built-in statements of WAIT EVENT cause the XBBSQ, HUBE10, XBXCQ, XDSKQ services to remain in a wait state until needed. Conversely, once awakened by the IDMS services, these schedulers utilize the IPSS built-in statements FIND IN QUEUE, WAIT IN QUEUE and REMOVE FROM QUEUE to maintain queues of suspended IDMS DML requests.

A3.4 IDMS Models

The purpose of the model of IDMS was to characterize IDMS DML processing in order to determine:

1. the IDMS contribution to transaction response time,
2. the effect of widely varying workload definitions,
3. the effect of processing DMLs concurrently (through multiple terminal loadings), and
4. the effect of data structure on processing efficiencies.

Two models of IDMS were prepared, one employing the current IPSS facilities and the other incorporating advanced IPSS features for data

structure characterization and data base access. These two models are similar in structure, they contain the same services and the calling sequences remains unchanged. However, the IPSS/BASIC model does not make use of the DML parameters passed from the SIDPERS transaction services. Virtual page identification is characterized through probability distribution functions. Schema definition and data access currency are not represented. It was found that these important features of data base systems are very difficult to model without the special simulation facilities provided by IPSS/DBS.

The second model of IDMS represents these features through the advanced data structure and data access facilities of IPSS/DBS. In this model, the DML parameters passed by the application services are used to traverse the data base. The data base structure is characterized through the SCHEMA facility and instances of the data base are created through special purpose IPSS/DBS built-in functions. This results in a much more accurate determination of data base processing activities in general and virtual page referencing sequences in particular. Secondary Storage functions were initially characterized by waiting an average elapsed time of 36 ms (the average time for PDP-11/70 head positioning and rotational delay). The detailed model of the PDP-11/70 was used to calculate these times.

Model Structure - Model I

The IPSS model of IDMS consists of services which interact in much the same way that the IDMS itself executes. The overall structure of the model is shown in Figure A3.4-1. The DML commands from the application services invoke the IDMS Interface service which in turn invokes the CAMP (Central Access Monitor Program) and DBMS services. The DBMS service invokes one of the 7 DML services which represent data access and retrieval. The interact with the Page Manager service and a single Schema Definition service in order to generate secondary storage references for the underlying file management system. The following is a description of the function of each of these services.

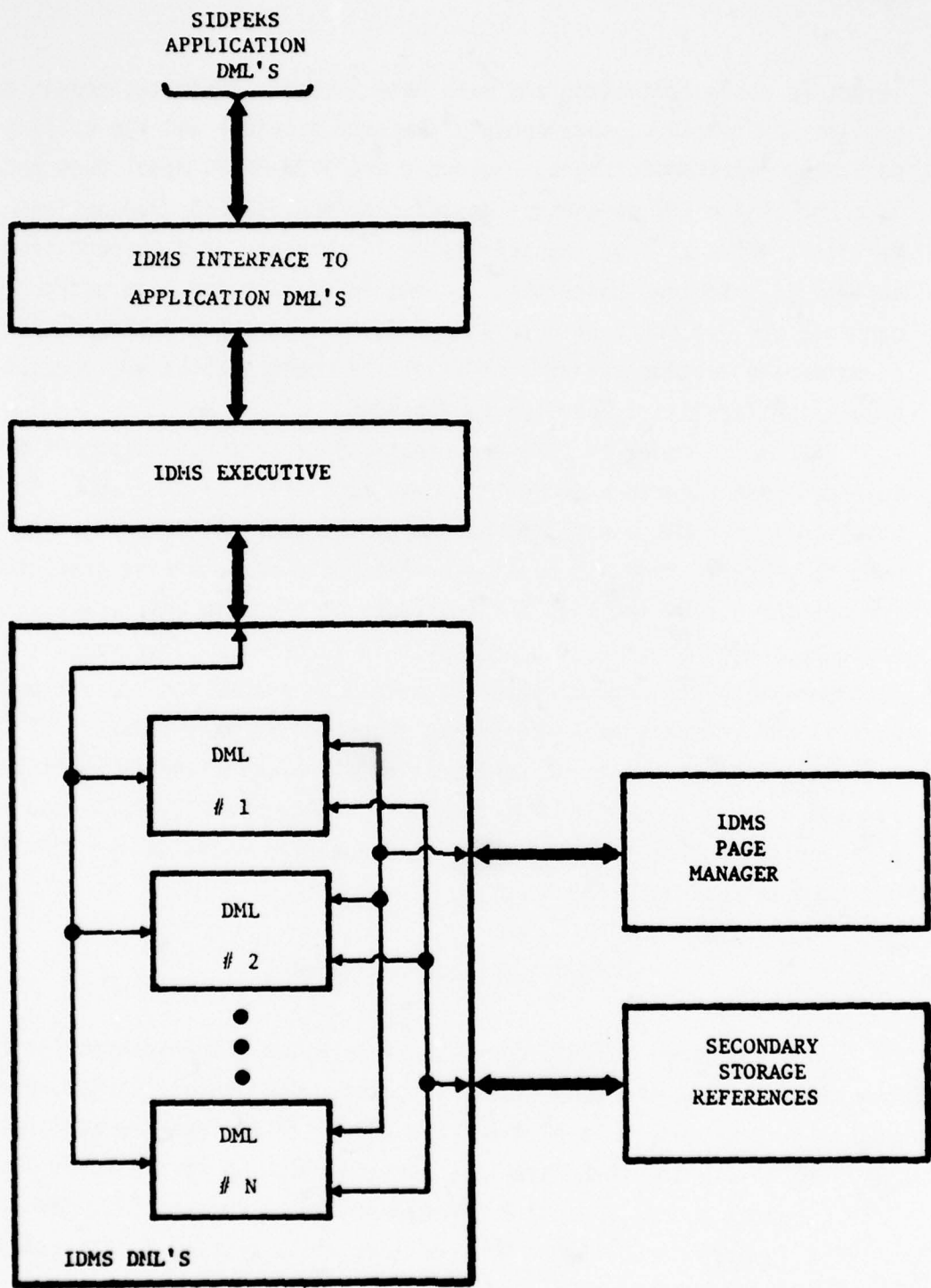


Figure A3.4-1 Schematic of the Modeled IDMS

IDMS Interface

This service represents IDMS processing activities that occur for every DML invocation. It is invoked by one of the application programs and represents the following IDMS functions:

1. validity checking and segmentation,
2. error status checking,
3. performing any 'before' procedures,
4. testing error status from the 'before' procedures,
5. testing the channel indicator
6. testing error status on return from performing the data base management system function,
7. executing any "on error during" procedures,
8. executing any 'after error' procedures, and
9. performing segmentation.

However, no information was available on the frequency or nature of these procedures within the test VIABLE/SIDPERS system. Thus, these functions were merely identified and documented within this service and can thus be easily expanded once more data becomes available. This service invokes the CAMP Service and waits for its completion.

Central Access Monitor Program (CAMP) Service

This service represents the IDMS program of the same name. Its function is simply to thread requests as the DBMS service can accommodate them thus holding the invoking services in the wait state. Because of the single thread of processing represented by a single user terminal, the queueing of requests at this level does not occur. However, queueing is an important performance factor for models with more than one terminal and/or concurrently executing application programs.

Data Base Management System (DBMS) Service

This service represents those functions at the very heart of IDMS, namely the materialization of an application record. This invokes the schema and subschema definitions, data base currency, the content of

buffers, and (usually) requests to the operating system for one or more data base pages. Because of the magnitude and complexity of these processes, they are modeled through several IPSS services. The DBMS service establishes the environment (e.g., page size, buffer size), collects cache paging statistics, invokes one of the IPSS DML services and waits for its completion.

DML Services

There are 7 IDMS DMLs represented in the IPSS model, each through its own service definition. These DMLs are the IDMS functions identified in Table A2.3-4. Each of these services performs the logic of the DML by:

1. determining the virtual page number from the record tape and set parameters as well as currency,
2. determining if the requested page is already present in the IDMS page cache, and
3. initiating a service to retrieve the record if not present.

Model Structure - Model II

The IPSS/DBS model of IDMS is similar in structure to the IPSS/BASIC model reported in the previous section. As shown in Figure A3.4-1, the application DMLs invoke the CAMP service which in turn invokes the DBMS service. In this model, however, the DBMS service has the additional functions of creating an instance of the schema. This is accomplished by the CREATE SCHEMA built-in facility which copies the Data Base Structure Component schema definition into an internal work area. This service then creates a set of arrays (called routes) which represent schema access paths. These routes are the basic mechanism for creating instances of sets, maintaining currency, and simulating the traversal of the data base.

A route consists of a series of set identifiers which includes the identification of the owner and member record types. A single

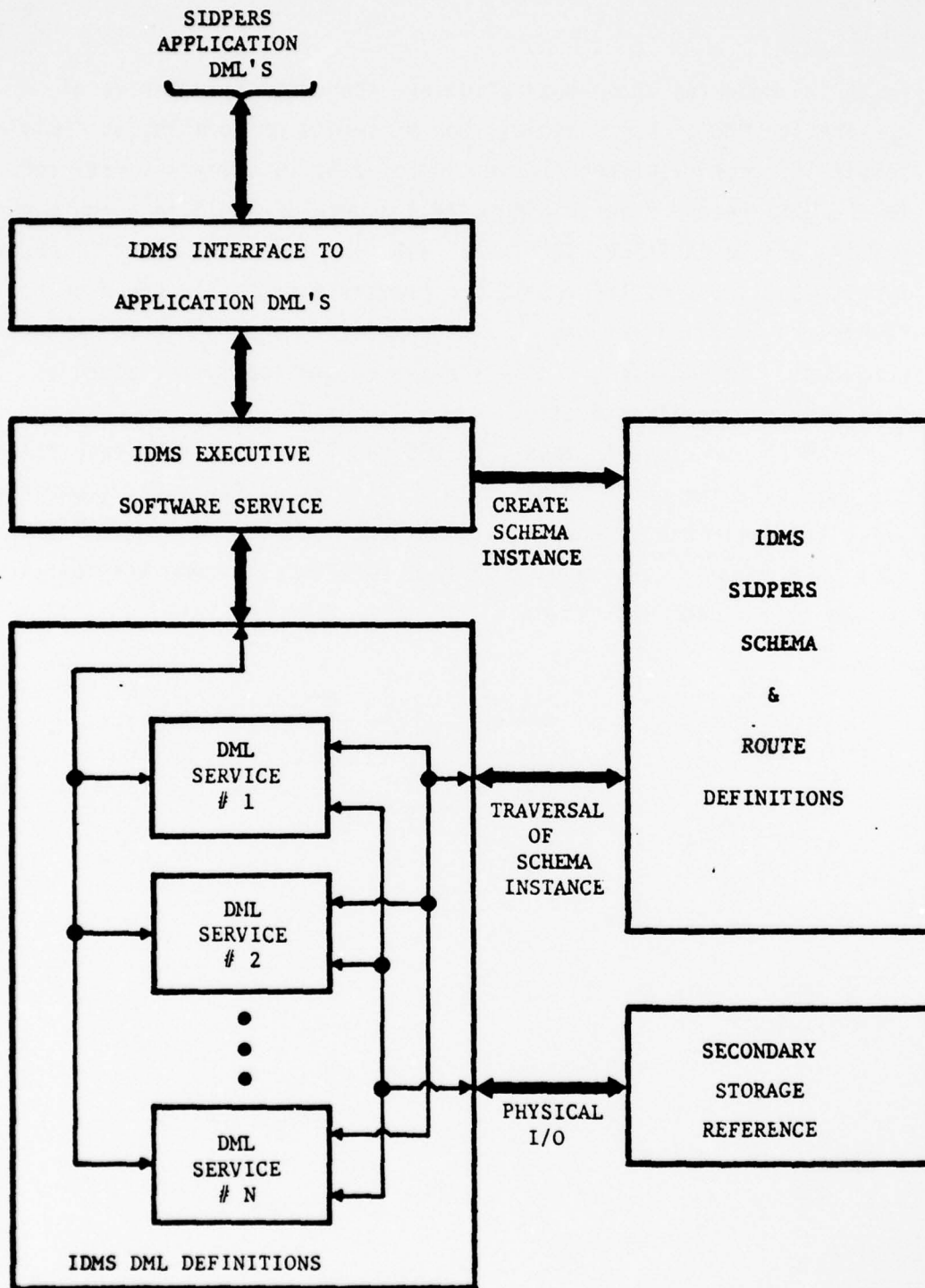


Figure A3.4-2 Schematic of IPSS/DBS Model of IDMS/SIDPERS Data Base

route incorporates those sets which are traversed by a series of application DMLs. For example, the DML series of locating an individual soldier's record via the CALC set followed by locating his data records in the UIST record type in Figure A2.3-2, would result in a route consisting of the CALC act, IDENT-UIST set, and the IDENT and UIST record types. Instances of these sets are created dynamically based on parameters supplied for the schema facility in the Data Base Structure component. In addition, statistics are automatically collected on data base traversal activities.

The DML services in this IPSS/DBS model simulate the traversal of the data base through such built-in facilities as FIND NEXT (member in set), and determine a record's virtual page address via the GET DATA BASE PAGE command. These virtual page references become the loading on the file management system.

APPENDIX B

SUMMARY OF MODEL STATISTICS

This appendix presents an overview and explanation of statistics generated by the execution of the IPSS models synthesized for the experimentation phase of the research. All of the defined models are not reproduced here, however, the final section of this Appendix contains the IPSS source listing for the SIDPERS/IDMS workload model.

B1. SIDPERS/IDMS WORKLOAD MODEL STATISTICS

The SIDPERS/IDMS workload model consists of a collection of IPSS service facilities whose invocation sequence is hierarchical. Figure B1-1 depicts the relationships among the services and indicates their generic function. As shown in the figure, transaction arrival at a work station is represented in the START service. Since a single terminal is modeled, the service queues all arriving transactions until the proceeding one has completed. Transaction parameters are assigned (TRANS) and the data fields are inputted and validated (VALID). Once the application programs is invoked (ADDSL, ARRIV, DEPTR, DTYC, GRCH, ING) which issues a DML by invoking the IDMS interface routine (modeled by INTF) and passing appropriate parameters. IDMS processing is sequenced by CAMP which then invokes DBMS. The function of the DBMS service is to establish the IDMS processing environment and to invoke a DML service (CONNECT, DISCON, ERASE, AND, MODIFY, OBTAIN, STORE). These services identify an IDMS page and request its presence in core by invoking the PAGMR service. PAGMR manages the IDMS page buffer using a modified least-recently-used page replacement algorithm. In this model, data transfer from secondary storage is simulated by a clock advance of 36 ms. Table B1-1 relates the IPSS services identified in Figure B1-1 to the corresponding SIDPERS/IDMS activity.

Statistics relative to this modeled process are contained in the following pages. An elapsed time of two hours was simulated; all times

Table B1-1. Index to Modeled SIDPERS/IDMS Processes

Level	SIDPERS/IDMS Process	IPSS Endogenous Exogenous Service Name	Source Code Page
1	Arrival of transaction at a work station	START	3
2	Transaction input and verification		
	a. Transaction input by terminal operator	TRANS	4
	b. Transaction verification by system software	VALID	10
3	Application software processing		
	a. Add a soldier to the data base	ADDSL	31
	b. Process soldier's arrival at the installation	ARRIV	23
	c. Process soldier's departure from installation	DEPTR	27
	d. Process change in soldier's duty status	DTYC	18
	e. Process soldier's grade change	CRCH	13
	f. Process inquiry	INQ	21
4	Software processing performed by the IDMS interface routine	INTF	35
5	IDMS Central Access Monitor Processor (CAMP)	CAMP	37
6	IDMS DBMS processing	DBMS	39
7	IDMS DML processing		
	a. CONNECT DML	CONECT	50
	b. DISCONNECT DML	DISCON	51
	c. ERASE DML	ERASE	52
	d. FIND DML	FIND	45
	e. MODIFY DML	MODIFY	46
	f. OBTAIN DML	OBTAIN	42
	g. STORE DML	STORE	47
8	Page management, retrieval of pages from auxiliary storage	PAGMR	54

IPSS MODEL - SERVICE HIERARCHY

SERVICE FUNCTION

START	VALID	transaction arrival					
TRANS		transaction input and verification					
<hr/>							
ADDSL	ARRIV	DEPTR	DTYC	GRCH	INQ	application processing	
<hr/>							
INTF						IDMS interface routine	
CAMP						DML sequencer	
DBMS						control	
<hr/>							
CONNECT	DISCON	ERASE	FIND	MODIFY	OBTAIN	STORE	DML's
<hr/>							
PAGMR							page manager

Figure B1-1. SIDPERS/IDMS Workload Model Service Invocation Sequence

are reported in milliseconds. The first page of model output shows that the exogenous service START was busy 62.4% of the time and was invoked nineteen times. The mean number of transactions at the work station at any one time was 1.1 (the maximum number was 3) indicating that a single terminal is sufficient to handle the processing load anticipated for a small data base (203 soldiers).

The average elapsed time from transaction arrival at a work station to complete is 5.64 min (338,430 ms). An examination of the transit time statistics for the endogenous service facilities reveals that the majority of this time was spent in transaction input and validation (TRANS and VALID services). This is due to the relatively long data input time required and the terminal operator error rates reflected through these services.

A significant drop in the percent busy and the busy period mean length statistics is evidenced for the INTF service. This indicates that more terminals could easily be supported. An appropriate number could be determined by experimentation with the model parameters.

An examination of the DML statistics reveals a relatively large number of executions of these services and that the average processing time is low. For example, there were 83 invocations of the FIND DML service and the average processing time was 37.3 ms. Most of this time was consumed by the PAGMR service which all the DML's invoke. The PAGMR is required to bring pages into core, and as reflected in the cache page fault statistics, the miss ratio was high (94.7%). This indicates that the data base is not well organized with respect to the DML processing. As inspection of the DML's in the application program services, however, reveals that a majority of the processing is random via CALC sets, and hence there is very little chance of more optimally organizing the data base given this processing load.

The last page of statistics for this model displays the queuing of transactions for the DBMS and TRANS services. Since the DBMS service is single-threaded (the function of the CAMP service), no queuing to it occurred. However, if CAMP was multi-threaded and the DBMS was not,

no change would be required in the model to collect DBMS queueing statistics. The queueing statistics for the TRANS service shows that the queue was never very long (maximum of 2, mean length of .6) again indicating a larger data base or heavier transaction rate could easily be accommodated.

MODEL SIMULATION PHASE

POST SIMULATION STATISTICS

PART C - UTILIZATION STATISTICS

FACILITY TYPE = EXS	FACILITY 1	FACILITY 2	FACILITY 3	FACILITY 4	FACILITY 5	FACILITY 6
STATISTICS DESCRIPTION	NAME	INDEX	NAME	INDEX	NAME	INDEX
	START	1				

BUSY PERIOD STATISTICS						
PER CENT BUSY	62.4					
NUMBER	8					
MEAN LENGTH	499438.9					
STND DEV OF MEAN LENGTH	736315.3					

IDLE PERIOD STATISTICS						
PER CENT IDLE	37.6					
NUMBER	10					
MEAN LENGTH	270504.5					
STND DEV OF MEAN LENGTH	203298.0					

CONCURRENCY STATISTICS						
CURRENT LEVEL	0					
MAXIMUM LEVEL	3					
MEAN LEVEL	7.1					
STND DEV OF MEAN LEVEL	0.8					

SEIZURE STATISTICS						
NUMBER OF COMPLETED SEIZURES	19					
NUMBER OF ZERO TIME SEIZURES	0					
PER CENT ZERO TIME SEIZURES	0.0					

TRANSITION TRANSIT TIME STATISTICS						
OVERALL MEAN TIME	338430.0					
OVERALL STND DEV OF MEAN TIME	252403.6					
MEAN TIME*	338430.0					
STND DEV OF MEAN TIME*	752403.6					
*EXCLUDING ZERO TIME SEIZURES						

THIS PAGE IS BEST QUALITY PRACTICES
FROM CASEY FURNISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

MODEL SIMULATION PHASE

POST SIMULATION STATISTICS

PART C - UTILIZATION STATISTICS

FACILITY TYPE = ENS

STATISTICS DESCRIPTION

FACILITY 1	FACILITY 2	FACILITY 3	FACILITY 4	FACILITY 5	FACILITY 6
NAME INDEX	NAME INDEX	NAME INDEX	NAME INDEX	NAME INDEX	NAME INDEX
ARRIV 1	CAMP 1	CONFCT 1	DEMS 1	DEPTR 1	DISCON 1

BUSY PERIOD STATISTICS

PER CENT BUSY	0.2	0.4	0.1	0.4	0.0	0.0
NUMBER	19	574	60	574	2	27
MEAN LENGTH	951.8	44.2	108.6	44.2	667.5	97.3
STND DEV OF MEAN LENGTH	243.6	36.6	58.7	36.6	25.5	54.5

IDLE PERIOD STATISTICS

PER CENT IDLE	99.8	99.6	99.9	99.6	100.0	100.0
NUMBER	19	575	61	575	3	28
MEAN LENGTH	37805.1	12477.6	117925.8	12477.6	2359254.0	257048.1
STND DEV OF MEAN LENGTH	67096.1	85898.4	318776.4	85898.4	3507636.0	474848.2

CONCURRENCY STATISTICS

CURRENT LEVEL	0	0	0	0	0	0
MAXIMUM LEVEL	1	1	1	1	1	1
MEAN LEVEL	0.5	0.5	0.5	0.5	0.5	0.5
STND DEV OF MEAN LEVEL	0.5	0.5	0.5	0.5	0.5	0.5

SEIZURE STATISTICS

NUMBER OF COMPLETED SEIZURES	18	574	60	574	2	27
NUMBER OF ZERO TIME SEIZURES	0	28	0	28	0	0
PER CENT ZERO TIME SEIZURES	0.0	4.9	0.0	4.9	0.0	0.0

TRANSACTION TRANSIT TIME STATISTICS

OVERALL MEAN TIME	951.8	44.2	108.6	44.2	667.5	97.3
OVERALL STND DEV OF MEAN TIME	243.6	36.6	58.7	36.6	25.5	54.5
MEAN TIME*	951.8	46.4	108.6	46.4	667.5	97.3
STND DEV OF MEAN TIME*	243.6	36.1	58.7	36.1	25.5	54.5

*EXCLUDING ZERO TIME SEIZURES

THIS PAGE IS BEST QUALITY FRAGMENT
FROM COPY FURNISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

MODEL SIMULATION PHASE

POST SIMULATION STATISTICS

PART C - UTILIZATION STATISTICS

FACILITY TYPE = EMS

STATISTICS DESCRIPTION

FACILITY 1 NAME INDEX	FACILITY 2 NAME INDEX	FACILITY 3 NAME INDEX	FACILITY 4 NAME INDEX	FACILITY 5 NAME INDEX	FACILITY 6 NAME INDEX
DTYC	1 ERASE	1 FIND	1 GRCH	1 INC	1 INTF
BUSY PERIOD STATISTICS					
PER CENT BUSY	0.0	0.0	0.0	0.1	0.4
NUMBER	4	83	3	4	574
MEAN LENGTH	413.1	36.0	37.3	1481.4	44.6
STD DEV OF MEAN LENGTH	54.0	0.0	8.8	20.8	36.6
IDLE PERIOD STATISTICS					
PER CENT IDLE	100.0	100.0	100.0	99.0	99.6
NUMBER	5	7	64	4	575
MEAN LENGTH	1439669.0	1028540.3	65677.1	179888.0	12477.6
STD DEV OF MEAN LENGTH	1308999.0	895815.3	212861.4	1613925.0	5593.4
CONCURRENCY STATISTICS					
CURRENT LEVEL	0	0	0	0	0
MAXIMUM LEVEL	1	1	1	1	1
MEAN LEVEL	0.5	0.5	0.5	0.5	0.5
STD DEV OF MEAN LEVEL	0.5	0.5	0.5	0.5	0.5
SEIZURE STATISTICS					
NUMBER OF COMPLETED SEIZURES	4	63	3	4	574
NUMBER OF ZERO TIME SEIZURES	0	1	0	0	2
PER CENT ZERO TIME SEIZURES	0.0	0.0	1.2	0.0	0.3
TRANSACTION TRANSIT TIME STATISTICS					
OVERALL MEAN TIME	413.1	36.0	37.3	1481.4	44.6
OVERALL STD DEV OF MEAN TIME	54.0	0.0	8.8	20.8	36.6
MEAN TIME*	413.1	36.0	37.8	1481.4	46.4
STD DEV OF MEAN TIME*	54.0	0.0	7.8	20.8	36.1

* (EXCLUDING ZERO TIME SEIZURES)

INFORMATION PROCESSING SYSTEM SIMULATOR

MODEL SIMULATION PHASE

POST SIMULATION STATISTICS

PART C - UTILIZATION STATISTICS

FACILITY TYPE = ENS	FACILITY 1		FACILITY 2		FACILITY 3		FACILITY 4		FACILITY 5		FACILITY 6	
	NAME	INDEX	NAME	INDEX	NAME	INDEX	NAME	INDEX	NAME	INDEX	NAME	INDEX
STATISTICS DESCRIPTION	MODIFY	1	OBTAIN	1	PAGMR	1	STORE	1	TRANS	1	VALID	1

EUSY PERIOD STATISTICS												
PER CENT BUSY	0.0		0.2		0.3		0.0		0.4		0.1	
NUMBER	62		319		699		17		19		34	
MEAN LENGTH	14.7		37.6		35.0		0.0		236576.1		131458.0	
STND DEV OF MEAN LENGTH	0.0		12.3		10.0		0.0		218316.4		138645.5	

IDLE PERIOD STATISTICS												
PER CENT IDLE	100.0		99.8		99.7		100.0		97.0		97.9	
NUMBER	63		320		700		18		40		35	
MEAN LENGTH	114270.9		22462.4		10250.7		305900.6		132222.3		78011.2	
STND DEV OF MEAN LENGTH	238647.5		114307.8		77980.0		622427.6		197001.9		92139.2	

CONCURRENCY STATISTICS												
CURRENT LEVEL	0		0		0		0		0		0	
MAXIMUM LEVEL	1		1		1		1		1		1	
MEAN LEVEL	0.5		0.5		0.5		0.5		0.5		0.5	
STND DEV OF MEAN LEVEL	0.5		0.5		0.5		0.5		0.5		0.5	

SEIZURE STATISTICS												
NUMBER OF COMPLETED SEIZURES	62		319		699		17		19		34	
NUMBER OF ZERO TIME SEIZURES	0		10		37		17		0		0	
PER CENT ZERO TIME SEIZURES	0.0		3.1		5.3		100.0		0.0		0.0	

TRANSACTION TRANSIT TIME STATISTICS												
OVERALL MEAN TIME	14.7		37.6		35.0		0.0		236576.1		131458.0	
OVERALL STND DEV OF MEAN TIME	0.0		12.3		10.0		0.0		218316.4		138645.5	
MEAN TIME*	14.7		38.8		32.9		0.0		236576.1		131458.0	
STND DEV OF MEAN TIME*	0.0		10.5		5.7		0.0		218316.4		138645.5	
* (EXCLUDING ZERO TIME SEIZURES)												

CACHE PAGE FAULT STATISTICS

CACHE SIZE (# PAGES): 5
PAGE SIZE (BYTES): 1024

PAGE WRITE REFERENCES: 17 CACHE MISSES: 0 MISS RATIO: 0.0
PAGE READ REFERENCES: 682 CACHE MISSES: 662 MISS RATIO: 0.9707

PAGE TOTAL REFERENCES: 699 CACHE MISSES: 662 MISS RATIO: 0.9471

PAGES REMITTEN BEFORE SPACE REUSED: 17
PAGES REPLACED WITHOUT REWRITING: 645

AVERAGE DISTANCE BETWEEN SUCCESSIVE PAGE REFERENCES: 37.1539

INFORMATION PROCESSING SYSTEM SIMULATOR

PAGE 6

MODEL SIMULATION PHASE

POST SIMULATION STATISTICS

PART D - QUEUING STATISTICS

FACILITY TYPE = ENS	FACILITY 1		FACILITY 2		FACILITY 3		FACILITY 4		FACILITY 5		FACILITY 6	
	NAME	INDEX	NAME	INDEX	NAME	INDEX	NAME	INDEX	NAME	INDEX	NAME	INDEX
STATISTICS DESCRIPTION	DBMS	1	TRANS	1								

BUSY PERIOD STATISTICS												
PER CENT BUSY	0.0		23.9									
NUMBER	574		16									
MEAN LENGTH	0.0		95422.4									
STND DEV OF MEAN LENGTH	0.0		144971.9									

IDLE PERIOD STATISTICS												
PER CENT IDLE	100.0		76.1									
NUMBER	575		19									
MEAN LENGTH	12521.6		288546.6									
STND DEV OF MEAN LENGTH	85898.6		195194.2									

QUEUE LENGTH STATISTICS												
CURRENT LENGTH	0		0									
MAXIMUM LENGTH	1		2									
MEAN LENGTH	0.5		0.6									
STND DEV OF MEAN LENGTH	0.5		0.6									

QUEUE ENTRY STATISTICS												
NUMBER OF COMPLETED ENTRIES	574		19									
NUMBER OF ZERO TIME ENTRIES	574		9									
PER CENT ZERO TIME ENTRIES	100.0		47.4									

TRANSACTION TRANSIT TIME STATISTICS												
OVERALL MEAN TIME	0.0		101853.7									
OVERALL STND DEV OF MEAN TIME	0.0		142386.9									
MEAN TIME*	0.0		193522.1									
STND DEV OF MEAN TIME*	0.0		144350.6									
*EXCLUDING ZERO TIME SEIZURES)												

B2. HOST/SIDPERS SUBMODEL STATISTICS

Previous USACSC studies did not identify a particular machine to be used for the host environment, so a generalized host was modeled. To do this, the software functions of the host were emphasized more than the hardware functions. Thus, only a minimum hardware configuration was described and more attention was paid towards determining and modeling the host operating system requirements of the SIDPERS application.

The host operating system was assumed to execute in a reactive manner, responding to the needs of the SIDPERS application. The particular steps in the modeled SIDPERS application were discussed in the main body of this report and in Appendix A. The needs of this application were, in general, task management and resource allocation. A crucial host/SIDPERS interface was the synchronization of the execution of SIDPERS services and operating system services. The asynchronous nature of their execution compounded the problem. To model this situation, IPSS services representing task schedulers communicated with existing SIDPERS IPSS services via the posting of and waiting for global events. These events corresponded to a request for scheduling and the completion of that scheduling process.

Four schedulers were identified for the modeling of the host/SIDPERS interface. These controlled access to the host's disk subsystem, the use of the teleprocessing (TP) line to the user, the use of the teleprocessing line to the back-end and the resumption of execution after a back-end request was completed. Each of these were characterized by a separate IPSS service and their use was controlled through a queue, exemplified by an IPSS chain facility. When a SIDPERS transaction required scheduling, it was placed in one of these chains. The four chains were HDSKQ (disk subsystem), TCAMQ (user TP communication), HOBEQ (back-end TP communication) and RSPONS (back-end request completion).

The use of these chains was measured in IPSS via the queueing statistics shown. Because of the particular SIDPERS workload, no SIDPERS applications overlapped, and, therefore, there was no queue in transit time associated with any scheduler. This condition also limited the possibility of having more than one SIDPERS transaction enqueued simultaneously. If such a situation occurred in the future, the concurrency statistics would demonstrate the degree of delay.

The queue entry statistics indicate the number of times that the current SIDPERS workload requested each of the scheduling tasks. The SIDPERS load for the run depicted is shown in the utilization statistics for the IPSS services that characterize SIDPERS processing. The total number of SIDPERS transaction for this run was 9, i.e., the number of requests to VALID. The number of requests to HDSKQ corresponds to the number of validity and/or compatability errors that occurred during the current run, since each error required a disk access. Thus, on the average, each incoming transaction had one error.

Statistics measuring the communication with the user and the back-end indicate the complexity of the SIDPERS application transactions. The number of user interactions (TCAMQ) reflects the amount of input required for the current set of SIDPERS transactions. Thus, for 9 transactions, a mean of 5.1 interactions were required per transaction. This included input of data and output of error messages and records. The number of entries in the RSPONS queue indicates that the current input stream of SIDPERS transactions generated 145 DML request to the back-end. This implies a mean of 16.1 DML requests per SIDPERS transaction types and does not take into account the DML requirements of particular transaction types. And finally, the number of entries in the back-end interface chain, HOBEQ, is twice the number of the actual DML requests since the TP interface must be acquired for both sending and receiving DML requests.

MODEL SIMULATION PHASE

POST SIMULATION STATISTICS

PART D - QUEUING STATISTICS

FACILITY TYPE = CHM

STATISTICS DESCRIPTION

FACILITY 1 FACILITY 2 FACILITY 3 FACILITY 4 FACILITY 5 FACILITY 6
 NAME INDEX NAME INDEX NAME INDEX NAME INDEX NAME INDEX NAME INDEX

MD5K0 1 HOBEO 1 RSPNS 1 TCAMQ 1

BUSY PERIOD STATISTICS

PER CENT BUSY 0.0 0.0 0.5 0.0 0.0
 NUMBER 10 290 145 46
 MEAN LENGTH 0.0 0.0 100.0 0.0
 STND DEV OF MEAN LENGTH 0.0 0.0 0.0 0.0

IDLE PERIOD STATISTICS

PER CENT IDLE 100.0 100.0 99.5 100.0
 NUMBER 11 291 146 57
 MEAN LENGTH 254545.1 9021.9 19078.7 59574.2
 STND DEV OF MEAN LENGTH 57145.4 69778.4 97742.3 140279.8

QUEUE LENGTH STATISTICS

CURRENT LENGTH 0 0 0 0
 MAXIMUM LENGTH 1 1 1 1
 MEAN LENGTH 0.5 0.5 0.5 0.5
 STND DEV OF MEAN LENGTH 0.5 0.2 0.5 0.5

QUEUE ENTRY STATISTICS

NUMBER OF COMPLETED ENTRIES 10 290 145 46
 NUMBER OF ZERO TIME ENTRIES 10 290 0 46
 PER CENT ZERO TIME ENTRIES 100.0 100.0 0.0 100.0

TRANSACTION TRANSIT TIME STATISTICS

OVERALL MEAN TIME 0.0 0.0 100.0 0.0
 OVERALL STND DEV OF MEAN TIME 0.0 0.0 0.0 0.0
 MEAN TIME* 0.0 0.0 100.0 0.0
 STND DEV OF MEAN TIME* 0.0 0.0 0.0 0.0

*EXCLUDING ZERO TIME SEIZURES

MODEL SIMULATION PHASE
PCST SIMULATION STATISTICS

PART C - UTILIZATION STATISTICS

FACILITY TYPE = ENS STATISTICS DESCRIPTION	FACILITY 1		FACILITY 2		FACILITY 3		FACILITY 4		FACILITY 5		FACILITY 6	
	NAME	INDEX	NAME	INDEX	NAME	INDEX	NAME	INDEX	NAME	INDEX	NAME	INDEX
ARRIV	1	BACK	1	DEPTR	1	DTYC	1	INO	1	INTF	1	
BUSY PERIOD STATISTICS												
PER CENT BUSY	0.3	0.5	0.1	0.1	0.1	0.1	0.0	0.0	0.0	0.5		
NUMBER	4	145	1	2	2	2	2	2	2	145		
MEAN LENGTH	2204.4	100.0	2104.2	1202.4	601.2	601.2	601.2	601.2	601.2	100.2		
STND DEV OF MEAN LENGTH	231.4	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.3		
IDLE PERIOD STATISTICS												
PER CENT IDLE	99.7	99.5	99.9	99.9	99.9	99.9	100.0	100.0	100.0	99.5		
NUMBER	3	146	2	3	3	3	3	3	3	146		
MEAN LENGTH	558236.2	19078.7	1390947.0	932531.0	932531.0	932531.0	932531.0	932531.0	932531.0	19078.7		
STND DEV OF MEAN LENGTH	814536.8	97742.3	251210.1	342562.8	342562.8	342562.8	342562.8	342562.8	342562.8	97742.3		
CONCURRENCY STATISTICS												
CURRENT LEVEL	0	0	0	0	0	0	0	0	0	0		
MAXIMUM LEVEL	1	1	1	1	1	1	1	1	1	1		
MEAN LEVEL	0.5	0.5	0.5	0.5	0.5	0.5	0.5	0.5	0.5	0.5		
STND DEV OF MEAN LEVEL	0.5	0.2	0.7	0.6	0.6	0.6	0.6	0.6	0.6	0.5		
SEIZURE STATISTICS												
NUMBER OF COMPLETED SEIZURES	4	145	1	2	2	2	2	2	2	145		
NUMBER OF ZERO TIME SEIZURES	0	0	0	0	0	0	0	0	0	0		
PER CENT ZERO TIME SEIZURES	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0		
TRANSACTION TRANSIT TIME STATISTICS												
OVERALL MEAN TIME	2204.4	100.0	2104.2	1202.4	601.2	601.2	601.2	601.2	601.2	100.2		
OVERALL STND DEV OF MEAN TIME	231.4	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.3		
MEAN TIME*	2204.4	100.0	2104.2	1202.4	601.2	601.2	601.2	601.2	601.2	100.2		
STND DEV OF MEAN TIME*	231.4	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.3		

THIS PAGE IS BEST QUALITY PRACTICES
FROM COPY FURNISHED TO DDG

MODEL SIMULATION PHASE

POST SIMULATION STATISTICS

PART C - UTILIZATION STATISTICS

FACILITY TYPE = EMS	FACILITY 1	FACILITY 2	FACILITY 3	FACILITY 4	FACILITY 5	FACILITY 6
STATISTICS DESCRIPTION	NAME INDEX	NAME INDEX	NAME INDEX	NAME INDEX	NAME INDEX	NAME INDEX

TRANS	1	VALID	1
-------	---	-------	---

BUSY PERIOD STATISTICS

PER CENT BUSY	23.4	22.9
NUMBER	7	10
MEAN LENGTH	93523.5	64008.0
STND DEV OF MEAN LENGTH	119212.9	106357.9

JULY PERIOD STATISTICS

PER CENT FULL	76.6	77.1
NUMBER	7	10
MEAN LENGTH	301476.1	215991.7
STND DEV OF MEAN LENGTH	229362.3	238402.2

CONCURRENCY STATISTICS

CURRENT LEVEL	1	1
MAXIMUM LEVEL	1	1
MEAN LEVEL	0.5	0.5
STND DEV OF MEAN LEVEL	0.5	0.5

SEIZURE STATISTICS

NUMBER OF COMPLETED SEIZURES	6	9
NUMBER OF ZERO TIME SEIZURES	0	0
PER CENT ZERO TIME SEIZURES	0.0	0.0

TRANSACTION TRANSIT TIME STATISTICS

OVERALL MEAN TIME	60278.0	71089.4
OVERALL STND DEV OF MEAN TIME	88146.7	110240.8
MEAN TIME*	60278.0	71089.4
STND DEV OF MEAN TIME*	88146.7	110280.8

*EXCLUDING ZERO TIME SEIZURES

THE INFORMATION CONTAINED HEREIN IS UNCLASSIFIED
 DATE 07-28-2010 BY 60322 UCBAW/STP

B3. IDMS/BACK-END SUBMODEL STATISTICS

The emphasis of this model was on interfacing the IDMS functions with the PDP-11/70 backend operations. Because of the wealth of documentation available on the hardware aspects of the USACSC PDP-11/70 (and the lack of software documentation), the impact of IDMS on hardware utilization was examined. The complete configuration of the USACSC PDP-11/70 installation at the Georgia Institute of Technology was modeled and is shown in Figure B3-1.

The PDP-11/70 operating system employs a dedicated data path between central memory and high speed disks (RWPO4AA in Figure B3-1). The IDMS data base resides on these disks, allowing for rapid retrieval of IDMS pages. DML requests entering the backend from the host were eventually mapped to specific data accesses to the RWPO4AA. Thus, to measure the direct effect of the IDMS loading on the PDP-11/70 backend, the IPSS statistics on the behavior of this disk subsystem (labeled RPO4AA in the corresponding IPSS model) are examined and are found immediately following this discussion.

Utilization statistics on the RPO4AA provide the best insight into its performance relative to the IDMS loading. Busy/idle statistics indicate the percentage of time that this facility was utilized. As can be seen, the RPO4AA was used only 23.1% of the time, remaining idle for 76.9%. Such an imbalance can be attributed to the characteristics of the device exceeding the needs of the application or the actual placement of data on the disks.

Concurrency statistics measure the degree of multiprocessing that occurred on the RPO4AA. In this case, in order to maintain the integrity of the IDMS data base, only one access at a time was permitted. This is evidenced by the value of the maximum level of concurrency being 1. The value of the current levels equalling 1 implies that at the end of the current simulation run, the RPO4AA device was in use.

The last two sets of utilization statistics characterize the behavior of individual accesses to the RPO4AA. Each access involves a seek and a data transfer operation. The RPO4AA was acquired (seized) throughout both operations. The total number of references(19) and the number

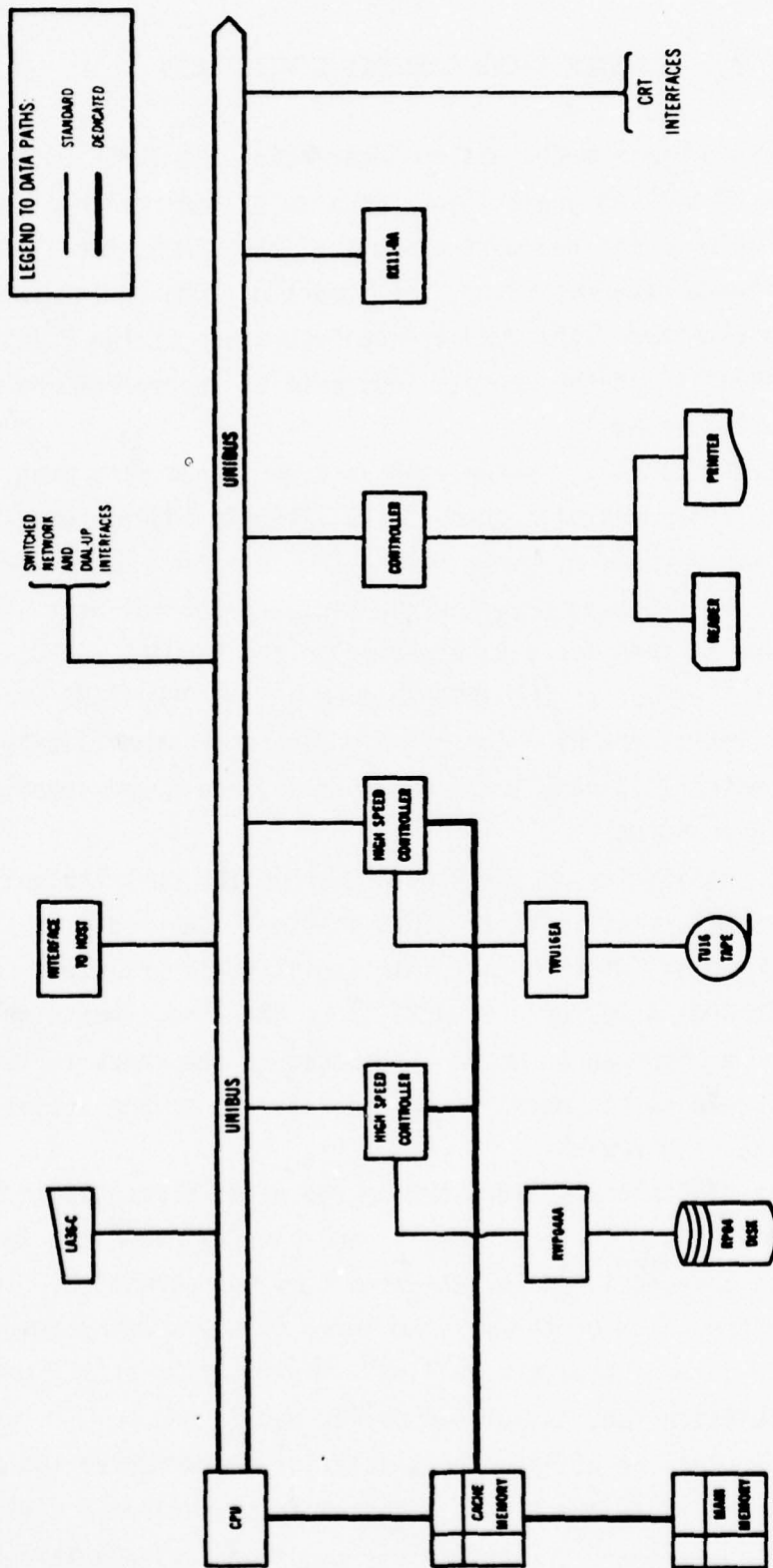


Figure B3-1 USACSC PDP 11/70 Configuration

of zero time references (0) indicate the relative closeness of successive page references to the IDMS data base. A low percentage, as is the case here, indicate that the data may be well spread out on the disk or that the individual accesses are randomly distributed. On the other hand, a high percentage of zero time seizures may indicate the compactness of the data or a deliberate ordering of accesses (e.g., sequential or within an IDMS set). The mean total time for an individual accesses to the RPO4AA, 23.2 milliseconds, includes the seek and data transfer operations. An additional mean is calculated removing the bias of the zero time seizures so that an average time per actual access can be determined. For each mean, the variability of that mean, i.e., the standard deviation, is also calculated to give some insight into the relative accuracy of the above means.

The impact of the IDMS functions can also be shown in other IPSS statistics. The IDMS DML requests are mapped into I/O requests for which IPSS has built-in primitives such as Seek and Data Transfer. Statistics are automatically calculated by IPSS on invocations of each of these I/O operations by each I/O operation (Part B1), by access mechanism (Part B2) and by data set (Part B3). The statistics by the access mechanism, RPO4AA, aid in analyzing the impact of IDMS on the backend. The sum of the mean routine times for the two I/O operations that referenced RPO4AA equals 23.2 milliseconds, the same utilization time for RPO4AA in the previous discussion. This allows us to break down the aggregate access time to RPO4AA into its component activities. It is evident that the Seek operation comprises the greatest amount of time in each DML request processing. Thus, in order to reduce this time factor affecting the Seek operation, such as placement of data files, compactness or device type, should be examined.

MODEL SIMULATION PHASE

POST SIMULATION STATISTICS

PART C - UTILIZATION STATISTICS

FACILITY TYPE = FA	FACILITY 1	FACILITY 2	FACILITY 3	FACILITY 4	FACILITY 5	FACILITY 6
STATISTICS DESCRIPTION	NAME INDEX	NAME INDEX	NAME INDEX	NAME INDEX	NAME INDEX	NAME INDEX
	RF04AA					
BUSY PERIOD STATISTICS						
PER CENT BUSY	24.1					
NUMBER	20					
MEAN LENGTH	23.4					
STND DEV OF MEAN LENGTH	17.4					
idle Period Statistics						
PER CENT Idle	76.9					
NUMBER	20					
MEAN LENGTH	78.0					
STND DEV OF MEAN LENGTH	17.4					
CONCURRENCY STATISTICS						
CURRENT LEVEL	1					
APPLIED LEVEL	1					
MEAN LEVEL	0.5					
STND DEV OF MEAN LEVEL	0.5					
SEIZURE STATISTICS						
NUMBER OF COMPLETED SEIZURES	19					
NUMBER OF ZERO TIME SEIZURES	0					
PER CENT ZERO TIME SEIZURES	0.0					
TRANSACTION TRAIL TIME STATISTICS						
OVERALL MEAN TIME	23.2					
OVERALL STND DEV OF MEAN TIME	17.7					
MEAN TIME*	23.2					
STND DEV OF MEAN TIME*	17.7					

THIS PAGE IS BEST QUALITY PRACTITIONER
FROM COPY FURNISHED TO DDC

* (EXCLUDING ZERO TIME SEIZURES)

INFORMATION PROCESSING SYSTEM SIMULATOR

MTEL SIMULATION PHASE

PST SIMULATION STATISTICS

PART 21 - I/O STATISTICS BY I/O ROUTINE

STATISTICS DESCRIPTION	TOTAL	ZERU	ZKEND	MEAN	STL DEV	*MEAN	*STU DEV
SEK STATISTICS	10	5	46.3				
NUMBER OF COMPLETES				20.6	12.7	28.0	6.0
TIME IN EXECUTION				2.2	2.2	3.1	1.9
NUMBER OF CYCLES CROSSED							
LATA TRANSFER STATISTICS	15						
DEVICE TYPE - LARD				2.0	0.0		
NUMBER OF READS				1024.0	0.0		
NUMBER OF WRITES							
TIME IN EXECUTION							
NUMBER OF CHARACTER TRANSMITTED							

THIS PAGE IS BEST QUALITY PRINTABLE
FROM JULY PUBLISHED TO DDC

MODEL SIMULATION PHASE

POST SIMULATION STATISTICS

PART 2 - I/O STATISTICS BY ACCESS MECHANISM

FACILITY IDENTIFICATION	I/O ROUTINE	NUMBER OF I/O CALLS				ROUTINE TIME			
		TOTAL	454C	824C	824D	MEAN	STU DEV	%MEAN	%STU DEV
040-4A	SEEK	19	5	20	3	20.6	12.7	28.0	0.0
	DATA TRANS	19	0	0	0	2.6	0.0	2.6	0.0

INFORMATION PROCESSING SYSTEM SIMULATOR

MODEL SIMULATION PHASE

POST SIMULATION STATISTICS

PART 13 - I/O STATISTICS BY DATA SET - FILE

FACILITY IDENTIFICATION	I/O CODE	NUMBER OF I/O CALLS				ROUTINE TIME		
		TOTAL	ZERO	2ZKRU	PLRN	STD DEV	MEAN	STD DEV
05	1	19	0	0.0	2.6	0.0	2.6	0.0

B4. IPSS FUNCTIONAL MODEL OF THE SIDPERS/IDMS

The purpose of this section is to illustrate the use of the IPSS to describe a complex information processing system at the functional level. This model does not contain characterizations for the computer hardware, executive software, data base management system software or data base. At this level they were treated as black boxes, however appropriate interfaces were defined so that respective characterizations could be included in later models.

The focus of this model was to model the processes (represented as IPSS endogenous services) encountered in updating the SIDPERS data base. Figure B1-1 identified the hierarchical relationship assumed for the model; Table B1-1 related the IPSS service identified in Figure B1-1 to a corresponding SIDPERS/IDMS data base update procedure. The IDMS data base management system backend was also represented in the model.

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT ALT INPUT
SEC-NO LEV SEQ-ID REF-AU

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00

DEFINE SYSTEM RESOURCES: NAME=CF, COMPILER=YES, DISPOSITION=KEEP; 0000010 1

PROCEDURE: NAME=COMENT, TYPE=SUBROUTINE; 0000020 2

C MODEL OF-- USACSC SIDPERS IDMS TEST DATA BASE --PERS3- 0000030 3

C FOR-- U.S. ARMY COMPUTER SYSTEM COMMAND 0000040 4

C APPLICATION-- SIX TYPES OF TRANSACTION PROCESSING - GRADE CHANGE, 0000050 5

C DUTY STATUS CHANGE, INQUIRY, ARRIVAL, DEPARTURE, ADD SOLDIER 0000060 6

C WRITTEN BY-- JOSEPH D. BROWNSMITH 0000070 7

C DATE-- JUNE-AUG 1977 0000080 8

C NOTE-- THIS MODEL IS WRITTEN IN JPSS (INFORMATION PROCESSING SYSTEM) 0000090 9

C SYSTEM SIMULATORY AND REPRESENTS THE USACSC'S SIDPERS IDMS 0000100 10

C TEST DATA BASE SYSTEM RUNNING ON THE CSC IMP370/165 COMPUTER 0000110 11

C (FRONT-END) AND THE PDPI/770 TRACK-END). 0000120 12

C THIS SIMULATION INCORPORATES THE FOLLOWING EXPERIMENT(S) - 0000130 13

C 1 0000140 14

C EXPERIMENT #1 0000150 15

C STANDARD SIDPERS TRANSACTION AND IDMS PARAMETERS 0000160 16

C EXPERIMENT #2 0000170 17

C PVERR = PCERR = 0 (PROB VALIDITY & COMPATIBILITY ERRORS) 0000180 18

C (SEE LINES 2540-2590) 0000190 19

C EXPERIMENT #3 0000200 20

C AVEDIS = 10 (DISTANCE BETWEEN SUCCESSIVE PAGE REFERENCES) 0000210 21

C (SEE LINE 9670) 0000220 22

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80
 INPUT CARD IMAGE
 SEC-NO ALT INPUT MFC-REF
 LEV SEC-NO MCF-MI

REF	SERVICE/PROCEDURE	APPROX	INVOICES/	DESCRIPTION	CALLS	PAGE	SEC-NO	ALT INPUT	MFC-REF
C							25	00000250	24
C							26	00000260	25
C							27	00000270	26
C							28	00000280	27
C							29	00000290	28
C							30	00000300	29
C	1	SYSTEM RESOURCES	(01)	-	COMPONENT DEFINITION		31	00000310	30
C	2	COMENT (PROC)	(01)	-	COMMENTS		32	00000320	31
C	3	START (EXS)	(03)	4	SCHEDULE TRANS (EVI)0000330		33	00000330	32
C	4	TRANS (ENS)	(04)	5-11	TRANS SELECTION		34	00000340	33
C	5	VALID (ENS)	(10)	-	TERMINAL INTERACTION0000350		35	00000350	34
C	6	GRCH (ENS)	(13)	12	GRADE CHANGE		36	00000360	35
C	7	DYTC (ENS)	(18)	12	DUTY STATUS		37	00000370	36
C	8	INO (ENS)	(21)	12	INQUIRY		38	00000380	37
C	9	ARRIV (ENS)	(23)	12	ARRIVAL		39	00000390	38
C	10	DEPTR (ENS)	(27)	12	DEPARTURE		40	00000400	39
C	11	ADDSL (ENS)	(31)	12	ADD SOLDIER		41	00000410	40
C	12	INTF (ENS)	(35)	73	HOST INTERFACE		42	00000420	41
C	13	CAMP (ENS)	(37)	14	ACCESS MONITOR		43	00000430	42
C	14	DBMS (ENS)	(39)	15-21	DBMS		44	00000440	43
C	15	CBTAN (ENS)	(42)	22-23	OBTAIN DML		45	00000450	44
C	16	FIND (ENS)	(45)	22-23	FIND DML		46	00000460	45
C	17	MODIFY (ENS)	(46)	22-23	MODIFY DML		47	00000470	46
C	18	STORE (ENS)	(47)	22-23	STORE DML		48	00000480	47

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE	INPUT SEQ-NO	ALT INPUT MEMBER	MEMBER REF-NO
.....10.....20.....30.....40.....50.....60.....70.....80			
C 19 CONECT (ENS) (50) 22-23 CONNECT DML C0000400	49		46
C 20 DISCON (ENS) (51) 22-23 DISCONNECT DML C0000500	50		49
C 21 FRASE (ENS) (52) 22-23 ERASE DML C0000510	51		50
C 22 PAGMR (ENS) (54) - PAGE MANAGER (CACHE)C0000520	52		51
C 23 GENRP (PKOC) (57) - PAGE # GENERATOR C0000530	53		52
C 24 BUFST (EXS) (60) - PRINT PAGE STATIEV2)C0000540	54		53
C 25 SNAPST (EXS) (61) - PRINT IPSS STATIEV3)C0000550	55		54
C 26 EXO EVENT STREAM (61) - COMPONENT DEFINITION C0000560	56		55
C 27 EVI (EXO EVENT) (62) 3 INVOKE PROCESSING C0000570	57		56
C 28 EV2 (EXO EVENT) (62) 24 DISPLAY PAGE STATS C0000580	58		57
C 29 EV3 (EXO EVENT) (62) 25 DISPLAY IPSS STATS C0000590	59		58
C 30 MODEL (65) - COMPONENT DEFINITION C0000600	60		59
C	61		60
C	62		61
C	63		62
C	64		63
C	65		64
C	66		65
C	67		66
C	68		67
C	69		68
C	70		69
C	71		70
C	72		71

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE
SFC-NO LFLV SEC-NL RFL-MC

.....10.....20.....30.....40.....50.....60.....70.....80

C U1SY U1ST 330 9. UMGR UNIT-MASTER-GRADE 73 72

C UMMA UNIT-MASTER 310 10. UMPOS UNIT-MASTER-POSNO 74 73

C U1. UMUI UNIT-MASTER-U1ST 75 74

RETURN 00000760 75 75

END: PROCEDURE: 00000770 77 76

00000780 78

00000790 79

END SERVICE: ID=START, NAME=START\$, 00000800 80 77

SAVE AREA SIZE = 10; 00000810 81 76

00000820 82 75

C THIS SERVICE REPRESENTS THE ARRIVAL OF TRANSACTIONS AT A SINGLE 00000830 83 80

WORK STATION. IT QUEUES THE TRANSACTION UNTIL THE PREVIOUS 00000840 84 81

C TRANSACTION IS COMPLETED. 00000850 85 82

END: DECLARATIONS: 00000860 86 83

END: INITIALIZATION: 00000870 87 84

SEIZE: FACILITY=START; 00000880 88 85

QUEUE: FACILITY=TRANS; 00000890 89 86

WAIT FACILITY: FACILITY=TRANS, STATUS=0, CONDITION=EQ; 00000900 90 87

DEPART QUEUE: FACILITY=TRANS; 00000910 91 88

SET STATUS: FACILITY=TRANS, STATUS=1; 00000920 92 89

INVOKE: SERVICE=TRANS; 00000930 93 90

WAIT RETURN: SERVICE=TRANS; 00000940 94 91

SET STATUS: FACILITY=TRANS, STATUS=0; 00000950 95 92

RELEASE: FACILITY=START; 00000960 96 93

.....10.....20.....30.....40.....50.....60.....70.....80.....90

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....	20.....	30.....	40.....	50.....	60.....	70.....	80.....	90.....	00.....	10.....	20.....	30.....	40.....	50.....	60.....	70.....	80.....	90.....	00.....	

END: EXO SERVICE:	00000970	97	94
	00000980	98	
	00000990	99	

END: SERVICE: ID=TRANS, NAME=TRANS\$,	00001000	100	95
SAVE AREA SIZE = 100;	00001010	101	96
	00001020	102	97

C THIS SERVICE REPRESENTS THE INTERACTION OF TERMINAL USER ENTERING 00001030	103	98
C TRANSACTIONS INTO THE SIDPERS SYSTEM. THIS TERMINAL CONTROL	104	99
C PROCESSING MODULE INCLUDES TERMINAL INTERACTION WITH THE DATA BASE 00001050	105	100
C 'VALID' RECORDS AND ALSO VALIDITY CHECKS. THE CHARACTERIZATION	106	101
C OF TRANSACTION PROCESSING REPRESENTED IN THIS SERVICE WAS	107	102
C DERIVED FROM "DESCRIPTION OF THE VIABLE TRANSACTIONS FOR THE	108	103
C "DINOT MODEL" BY HELMUT SCHAAFF, MARCH 1977.	109	104
C	110	105
C	111	106
C	112	107
C	113	108
C	114	109
C	115	110
C	116	111
C	117	112

C THE TYPES OF TRANSACTIONS REPRESENTED IN THE MODEL AND THEIR	110	105
C PERCENT OCCURRENCE ARE-	111	106
C	112	107
C	113	108
C	114	109
C	115	110
C	116	111
C	117	112

C DATA SERVICE/TRANS/	118	113
C	119	114
C	120	115

C INTEGER RETC, SEED, SUB, YLOC(17)	118	113
C	119	114
C	120	115

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT MEMLEF
SFC-NO LEV SFC-NU REF-NU

2	63., 8., 3.40, 6.40, 11.1, 26.4, 7.,	00001450	145	140
C	INQUIRY	00001460	146	141
3	82., 3., 1.75, 10.53, 18.0, 27.7, 0.,	00001470	147	142
C	ARRIVAL	00001480	148	143
4	5., 9., 2.00, 2.18, 16.9, 32.3, 16.,	00001490	149	144
C	DEPARTURE	00001500	150	145
5	57., 12., 2.60, 4.36, 13.7, 27.9, 18.,	00001510	151	146
C	ADD INDIVIDUAL	00001520	152	147
6	0., 70., 4.00, 1.43, 0.0, 20.5, 0. /	00001530	153	148
C	EQUIVALENCE (\$SAVE(15),IP), (\$SAVE(6),SEED), (\$SAVE(7),SUB)	00001540	154	149
C	END: DECLARATIONS:	00001550	155	150
C	SEED = 1	00001560	156	151
C	DU TO I = 1,10	00001570	157	152
C	TO YLOC(I) = 0	00001580	158	153
C	END: INITIALIZATION:	00001590	159	154
C	SEIZE: FACILITY=TRANS:	00001600	160	155
C	SELECT TYPE OF TRANSACTION	00001610	161	156
C	RN = SUNFM(0.0, 1.0, SEED)	00001620	162	157
C	RP = SPWLN(PUNTS,7,RN)	00001630	163	158
C		00001640	164	159
C		00001650	165	160
C		00001660	166	161
C		00001670	167	162
C		00001680	168	163

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

```

.....10.....20.....30.....40.....50.....60.....70.....80
INPUT CARD IMAGE
SEG-NO  LEV  SEQ-NO  REF-NO
      1P = RP                169                164
C
50 SUB = ((IP-1)*7)+1      170                165
      INVOKE: SERVICE = VALID$, PARAMETER LIST=(ISEED,SUB,TPARM);
      WAIT RETURN: SERVICE=VALID;
      IP(TLOC(8) .EQ. 1) GO TO 950
C
      TRANSACTION IS NOT CANCELLED
      GO TO (100,200,300,400,500,600),IP
C
      START GRADE CHANGE PROCESSING
      200 CONTINUE
      INVOKE: SERVICE=GRCH$, PARAMETER LIST=(TPARM(4),SEED);
      WAIT RETURN: SERVICE=GRCH;
      GO TO 900
C
      START DUTY STATUS CHANGE PROCESSING
      200 CONTINUE
      INVOKE: SERVICE=DITYC$, PARAMETER LIST=(TPARM(1),SEED);
      WAIT RETURN: SERVICE=DITYC;
      GO TO 900
C
      START INQUIRY PROCESSING
      300 CONTINUE

```

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

	INPUT CARD IMAGE	SEQ-NO	ALT INPUT MEMBER
.....10.....20.....30.....40.....50.....60.....70.....80.....	193	188
INVOKE: SERVICE=INQS, PARAMETER LIST=(TPARM(16),SEED);	00001920	194	189
WAIT RETURN: SERVICE=INQ;	00001940	195	190
GO TO 900	00001950	196	191
C	00001960	197	192
C START ARRIVAL PROCESSING	00001970	198	193
400 CONTINUE	00001980	199	194
INVOKE: SERVICE=ARRIVS, PARAMETER LIST=(TPARM(25),SEED);	00001990	200	195
WAIT RETURN: SERVICE=ARRIV;	00002000	201	196
GO TO 900	00002010	202	197
C	00002020	203	198
C START DEPARTURE PROCESSING	00002030	204	199
500 CONTINUE	00002040	205	200
INVOKE: SERVICE=DEPTRS, PARAMETER LIST=(TPARM(32),SEED);	00002050	206	201
WAIT RETURN: SERVICE=DEPTR;	00002060	207	202
GO TO 900	00002070	208	203
C	00002080	209	204
C START ADDITION PROCESSING	00002090	210	205
600 CONTINUE	00002100	211	206
INVOKE: SERVICE=ADDLS, PARAMETER LIST=(TPARM(39),SEED);	00002110	212	207
WAIT RETURN: SERVICE=ADDLS;	00002120	213	208
GO TO 900	00002130	214	209
C	00002140	215	210
900 CONTINUE	00002150	216	211
C CHECK FOR COMPATABILITY ERROR IN TRANSACTION	00002160		
.....10.....20.....30.....40.....50.....60.....70.....80.....		

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....

IF(TP .LT. 0) GO TO 950 00002170 217 212

IF(TP .GT. 10) GO TO 950 00002180 218 213

IF(VLOC(TP) .EQ. 0) GO TO 950 00002190 219 214

C COMPATABILITY ERROR 00002200 220 215

VLOC(TP) = 0 00002210 221 216

GO TO 50 00002220 222 217

950 CONTINUE 00002230 223 218

RELEASE: FACILITY=TRANS; 00002240 224 219

END: ENDO SERVICE; 00002250 225 220

00002260 226

00002270 227

ENDO SERVICE: ID=VALID, NAME=VALID, SAVE AREA SIZE=10, 00002280 228 221

PARAMETER LIST=(SEEDI,SUBI,TPARM): 00002290 229 222

C THIS SERVICE REPRESENTS THE TERMINAL INTERACTIONS INVOLVED 00002300 230 223

C IN OBTAINING A TRANSACTION WITH NO VALIDATION ERRORS. 00002310 231 224

COMMON /MK/ VLOC 00002320 232 225

INTEGER SEED,SUB,RETC,\$UNFMI,SEEDI,SUBI,VLOC(I0) 00002330 233 226

REAL DARR(42),DPA,NUMDI,PVERR,PCFERR,PTDM,PTTM,PCANCL,TPARM(42) 00002340 234 227

REAL SURFH 00002350 235 228

END: DECLARATIONS; 00002360 236 229

GO TO I = 1,42 00002370 237 230

I0 DARR(I) = TPARM(I) 00002380 238 231

SEED = SEEDI 00002390 239 232

SUB = SUBI 00002400 240 233

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....

THIS PAGE IS BEST QUALITY PRACTICES
FROM COPY FURNISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....	INPUT	ALT INPUT MEMBER
.....10.....20.....30.....40.....50.....60.....70.....80.....90.....	SFC=NO	LEV SEC=RU REF=NO
END: INITIALIZATION:	247	224
SEIZE: FACILITY=VALID:	242	235
C GET PROBABILITY OF DIRECT MODE OF OPERATION	243	236
DM = DARR(SUB+0)/100.	244	237
C GET NUMBER OF DATA ITEMS THIS TRANSACTION	245	238
NUMDI = DARR(SUB+1)	246	239
C GET PROBABILITY OF VALIDITY ERROR	247	240
PVERR = DARR(SUB+2)/100.	248	241
C PVERR = 0.	249	242
C GET PROBABILITY OF COMPATABILITY ERROR	250	243
PCERR = DARR(SUB+3)/100.	251	244
C PCERR = 0.	252	245
C GET DIRECT MODE PROCESS TIME IN MSEC	253	246
PTDM = DARR(SUB+4) * 1000.	254	247
C GET TUTORIAL MODE PROCESS TIME IN MSEC	255	248
PTTM = DARR(SUB+5) * 1000.	256	249
C GET PROBABILITY THAT THIS TRANSACTION WILL BE CANCELLED	257	250
PCANCL = DARR(SUB+6)/100.	258	251
C SELECTY INPUT MODE (EITHER DIRECTY OR TUTORIAL)	259	252
RM = SUNFM(0.0, 1.0, SEED)	260	253
IF(RM .GT. DM) GO TO 50	261	254
C DIRECT MODE OF OPERATION	262	255

THIS PAGE IS BEST QUALITY PRACTICE FROM COPY FURNISHED NO DDC

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80.....90

C WAIT FOR DATA TO BE ENTERED (SYSTEM HOOK) 265 258

PROCESS: TIME=PTDM; 266 259

C DETERMINE IF A VALIDITY ERROR HAS OCCURRED ON THE DATA ITEM 267 260

RN = \$URFM(0.0, 1.0, SEED) 268 261

IF(RN .GT. PVERR) GO TO 40 269 262

C VALIDITY ERROR - WAIT FOR CORRECTED INPUT (SYSTEM HOOK) 270 263

PROCESS: TIME=PTDM; 271 264

C NO VALIDITY ERROR 272 265

40 CONTINUE 273 266

GO TO 100 274 267

C 50 J = 1 275 268

C TUTORIAL MODE OF OPERATION 276 269

60 IF(J .GT. NUMDI) GO TO 100 277 270

C WAIT FOR INPUT FROM TERMINAL (SYSTEM HOOK) 278 271

PROCESS: TIME=PTDM; 279 272

C DETERMINE IF VALIDITY ERROR HAS OCCURRED FOR THIS DATA ITEM 280 273

RN = \$URFM(0.0, 1.0, SEED) 281 274

IF(RN .GT. PVERR) GO TO 70 282 275

C VALIDITY ERROR - WAIT FOR CORRECTED INPUT (SYSTEM HOOK) 283 276

PROCESS: TIME=PTDM; 284 277

70 J = J + 1 285 278

C GO TO 60 286 279

.....10.....20.....30.....40.....50.....60.....70.....80.....90

THIS PAGE IS BEST QUALITY PRACTICE FROM COPY FURNISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE
SEQ-NO ALT INPUT MEMBER
LEV SEQ-NL KIP-NC

.....10.....20.....30.....40.....50.....60.....70.....80

C 00002850 280 282

C 00002900 280 283

C 00002910 291 284

100 RW = SUNFM(0.0, I.0, SEED)

TLOC(R) = 0

IFIRM .LY. PCANCL) TLOC(8) = 1

RELEASE: FACILITY=VALID;

END: ENDD SERVICES;

00002960 286

00002970 297 285

00002980 298 286

00002990 299 287

00003000 300 288

00003010 301 289

00003020 302 290

00003030 303 291

00003040 304 292

00003050 305 293

00003060 306 294

00003070 307 295

00003080 308 296

00003090 309 297

00003100 310 298

00003110 311 299

00003120 312 300

.....10.....20.....30.....40.....50.....60.....70.....80

THIS PAGE IS BEST QUALITY PRINTING FROM COPY FURNISHED TO DOC

INFORMATION PROCESSING SYSTEM SIMULATOR
STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80
INPUT CARD IMAGE
SFC-NO ALT INPUT MEMOR. LEV SFC-NO RFT-NE

I / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10/
00003130 313 303

DATA CALC,GRUI,TDUI,INUM,PRID,PNUI,SMID,UMAM,UMGR,UMPOS,UMUI
00003140 314 304

I / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11/
00003150 315 305

END: DECLARATIONS;
00003160 316 306

PERR = PERRI/100.
00003170 317 307

SEED = SEEDI
00003180 318 308

END: INITIALIZATION;
00003190 319 309

SETZE: FACILITY=GRCH;
00003200 320 310

C RETURN CODE -TLOC(I) = 0 IF NO ERROR, 1 IF ERROR
00003210 321 311

C PARAMETER LIST CONTAINS-
00003220 322 312

C 1) DML TYPE - OBTAIN=1,FIND=2,MODIFY=3,STORE=4,CONNECT=5,
00003230 323 313

C DISCONNECT=6,ERASE=7
00003240 324 314

C 2) QUALIFIER - DMY=0,NEXT=1,PPTR=2,OWNER=3
00003250 325 315

C 2) RECORD TYPE - IDENTIFIER OF THE RECORD UPON WHICH THE DML IS
00003260 326 316

C TO OPERATE
00003270 327 317

C 4) SET IDENTIFIER - INDEX
00003280 328 318

C INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMY,IDEN,CALC);
00003290 329 319

WAIT RETURN: SERVICE=INTF;
00003300 330 320

WAIT RETURN: SERVICE=INTF;
00003310 331 321

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMY,ADDI,CALC);
00003320 332 322

WAIT RETURN: SERVICE=INTF;
00003330 333 323

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMY,MOSC,CALC);
00003340 334 324

WAIT RETURN: SERVICE=INTF;
00003350 335 325

INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,NEXT,UIST,TDUI);
00003360 336 326

.....10.....20.....30.....40.....50.....60.....70.....80

THIS PAGE IS BEST QUALITY PRACTICE
FROM COPY FURNISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

```

.....10.....20.....30.....40.....50.....60.....70.....80
INPUT CARD IMAGE
SEC-NO ALT INPUT MEM-LOC
LIV SIC-NO PER-NO
.....10.....20.....30.....40.....50.....60.....70.....80
WAIT RETURN: SERVICE=INTF; 337 00003370 327
INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,OMNER,DHY,UMUI); 338 00003380 326
WAIT RETURN: SERVICE=INTF; 339 00003390 324
INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DHY,POSN,CALC); 340 00003400 330
WAIT RETURN: SERVICE=INTF; 341 00003410 331
RN = $UMFH(0.0, 1.0, SEFD) 342 00003420 332
IFIPN .LY. PERRY GO TO 100 343 00003430 333
C 344 00003440 334
C 345 00003450 335
C 346 00003460 336
C 347 00003470 337
INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,DHY,ADDI,CALC); 348 00003480 338
WAIT RETURN: SERVICE=INTF; 349 00003490 339
INVOKE: SERVICE= INTFS, PARAMETER LIST=(MODIF,DHY,ADDI,CALC); 350 00003500 340
WAIT RETURN: SERVICE=INTF; 351 00003510 341
INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,DHY,IDENT,CALC); 352 00003520 342
WAIT RETURN: SERVICE=INTF; 353 00003530 343
INVOKE: SERVICE= INTFS, PARAMETER LIST=(MODIF,DHY,IDENT,DHY); 354 00003540 344
WAIT RETURN: SERVICE=INTF; 355 00003550 345
INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,DHY,IDENT,CALC); 356 00003560 346
WAIT RETURN: SERVICE=INTF; 357 00003570 347
INVOKE: SERVICE= INTFS, PARAMETER LIST=(DISCN,DHY,IDENT,PMID); 358 00003580 348
WAIT RETURN: SERVICE=INTF; 359 00003590 349
INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DHY,MOJSC,DHY); 360 00003600 350

```

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00
INPUT CARD IMAGE
SEC-ID LRV SEC-ID REP-ID

WAIT RETURN: SERVICE=INTF: 00003610 361

INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,DMY,IDEN,CALC): 00003620 362

WAIT RETURN: SERVICE=INTF: 00003630 363

INVOKE: SERVICE= INTFS, PARAMETER LIST=(CONCT,DMY,IDEN,PMIO): 00003640 364

WAIT RETURN: SERVICE=INTF: 00003650 365

WAIT RETURN: SERVICE=INTF: 00003660 366

C DISCONNECT THE PERSON'S UIST RECORD 0000367C 367

C 00003680 368

INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,NEXT,UIST,DUUI): 00003690 369

WAIT RETURN: SERVICE=INTF: 0000370C 370

INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,OWNER,DMY,GRUI): 00003710 371

WAIT RETURN: SERVICE=INTF: 00003720 372

INVOKE: SERVICE= INTFS, PARAMETER LIST=(DISCN,DMY,UIST,GRUI): 00003730 373

WAIT RETURN: SERVICE=INTF: 00003740 374

INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,NEXT,UIST,DUUI): 00003750 375

WAIT RETURN: SERVICE=INTF: 00003760 376

INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,OWNER,UIST,GRUI): 00003770 377

WAIT RETURN: SERVICE=INTF: 0000378C 378

INVOKE: SERVICE= INTFS, PARAMETER LIST=(DISCN,DMY,UIST,GRUI): 00003790 379

WAIT RETURN: SERVICE=INTF: 00003800 380

C 00003810 381

C CONNECT UIST RECORDS TO THE GRADE-UIST SET OF THE NEW GRADE 00003820 382

C 0000383C 383

INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,NEXT,UIST,DUUI): 0000384C 384

C 00003850 385

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00

THIS PAGE IS BEST QUALITY PRACTICE
REPRODUCED FROM THE ORIGINAL

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT ALT INPUT MEMBER
FROM LRU SIGNL FROM

.....10.....20.....30.....40.....50.....60.....70.....80

WAIT RETURN: SERVICE=INTF; 385 00003850 375

INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,DMV,IDEN,CALC); 386 00003860 376

WAIT RETURN: SERVICE=INTF; 387 00003870 377

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,NEXT,GRADE,UMGR); 388 00003880 378

WAIT RETURN: SERVICE=INTF; 389 00003890 379

INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,DMV,UIST,CALC); 390 00003900 380

WAIT RETURN: SERVICE=INTF; 391 00003910 381

INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,OMNER,DMV,GRUI); 392 00003920 382

WAIT RETURN: SERVICE=INTF; 393 00003930 383

INVOKE: SERVICE= INTFS, PARAMETER LIST=(CONCT,DMV,UIST,GRUI); 394 00003940 384

WAIT RETURN: SERVICE=INTF; 395 00003950 385

INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,NEXT,UIST,IDUI); 396 00003960 386

WAIT RETURN: SERVICE=INTF; 397 00003970 387

INVOKE: SERVICE= INTFS, PARAMETER LIST=(CONCT,DMV,UIST,GRUI); 398 00003980 388

WAIT RETURN: SERVICE=INTF; 399 00003990 389

C NORMAL RETURN 400 00004000 390

90 TLOC(1) = 0 401 00004010 391

GO TO 150 402 00004020 392

C COMPATABILITY ERROR 403 00004030 393

100 CONTINUE 404 00004040 394

TLOC(1) = 1 405 00004050 395

150 CONTINUE 406 00004060 396

RELEASE: FACILITY=GRCH; 407 00004070 397

END: ENDO SERVICE; 408 00004080 398

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/16/77

INPUT CARD IMAGE
SFC=IN LEV SEQ=LO PRT=NO

.....10.....20.....30.....40.....50.....60.....70.....80

00004090 409

00004100 410

00004110 411

00004120 412

00004130 413

00004140 414

00004150 415

00004160 416

00004170 417

00004180 418

00004190 419

00004200 420

00004210 421

00004220 422

00004230 423

00004240 424

00004250 425

00004260 426

00004270 427

00004280 428

00004290 429

00004300 430

00004310 431

00004320 432

THIS PAGE IS BEST QUALITY PRACTICABLE
FROM COPY FURNISHED TO DDC

ENDS SERVICE: ID=DTVC, NAME=DTVC, SAVE AREA SIZE=10,

PARAMETER LIST=(PERR,SEED);

COMMON /WK/ YLOC

INTEGER SUNFMT,SEED,SEED1,TLN(10)

INTEGER DBTAN,FINDR,MODIF,STOR,CONCT,DISCN,ERAS

INTEGER DMV,NEXT,PRIGR,OWNER

INTEGER ADDI,ADUN,ALOC,GRADE,IDFN,INST,MOSC,POSN,UIST,UNMA

INTEGER CALC,GRUI,IOUI,INUM,PMID,POUI,SMID,UMAM,UMGR,UMPOS,UMUI

EQUIVALENCE (\$SAVE(1),PERR), (\$SAVE(2),SEED)

REAL SUNFM,RN,PERR,PERRI

DATA DBTAN,FINDR,MODIF,STOR,CONCT,DISCN,ERAS/1,2,3,4,5,6,7/

DATA DMV,NEXT,PRIGR,OWNER/0,1,2,3/

DATA ADDI,ADUN,ALOC,GRADE,IDFN,INST,MOSC,POSN,UIST,UNMA

1 / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10/

DATA CALC,GRUI,IOUI,INUM,PMID,POUI,SMID,UMAM,UMGR,UMPOS,UMUI

1 / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11/

FND: DECLARATIONS:

PERR = PERR/100.

SEED = SEED1

ENDS INITIALIZATION:

SEIZE: FACILITY=DTVC:

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE

INPUT ALT INPUT PEREFF
SF(0-P) LFV SECU-LU KEF-LU

.....10.....20.....30.....40.....50.....60.....70.....80	00004330	433	426
C	00004340	434	427
C	00004350	435	428
C	00004360	436	429
C	00004370	437	430
C	00004380	438	431
C	00004390	439	432
C	00004400	440	433
C	00004410	441	434
C	00004420	442	435
C	00004430	443	436
C	00004440	444	437
C	00004450	445	438
C	00004460	446	439
C	00004470	447	440
C	00004480	448	441
C	00004490	449	442
C	00004500	450	443
C	00004510	451	444
C	00004520	452	445
C	00004530	453	446
C	00004540	454	447
C	00004550	455	448
C	00004560	456	449

THIS PAGE IS BEST QUALITY FRAGMENTED FROM COPY ANALYSIS TO DDC

RETURN CODE -TLOC(ZI) = 0 IF NO ERROR, I IF ERROR

PARAMETER LIST CONTAINS-

1) DML TYPE - OBTAIN=1, FIND=2, MODIFY=3, STORE=4, CONNECT=5,
DISCONNECT=6, ERASE=7

2) QUALIFIER - DMY=0, NEXT=1, PRIORITY=2, OWNER=3

3) RECORD TYPE - IDENTIFIER OF THE RECORD UPON WHICH THE DML IS TO OPERATE

4) SET IDENTIFIER - INDEX

COMPATABILITY CHECK

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMY,IOFN,CALC);
WAIT RETURN: SERVICE=INTF;

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMY,ADDI,CALC);
WAIT RETURN: SERVICE=INTF;

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMY,ADDI,CALC);
WAIT RETURN: SERVICE=INTF;

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,NEXT,UIST,IOUI);
WAIT RETURN: SERVICE=INTF;

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,NEXT,UIST,IOUI);
WAIT RETURN: SERVICE=INTF;

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,CWNER,DMY,UMUI);
WAIT RETURN: SERVICE=INTF;

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMY,ADUN,CALC);
WAIT RETURN: SERVICE=INTF;

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR
STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80
INPUT CARD IMAGE
SEQ-NO ALT INPUT MEMBER
LEV SEQ-NO REF-NO

WAIT RETURN: SERVICE=INTF; 457 444

C DETERMINE IF COMPATABILITY ERROR HAS OCCURRED 458 445

RN = SUNFM(0, 1, 0, SEED) 459 446

IF(RN .LT. FERR) GO TO 100 460 447

C NO COMPATABILITY ERROR 461 448

C 462 449

C MODIFY THE IDENT, UIST AND ADD-UNTY RECORDS 463 450

C 464 451

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(FINDR,DMY,IDEN,CALC); 465 452

WAIT RETURN: SERVICE=INTF; 466 453

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(MODIF,DMY,IDEN,DMY); 467 454

WAIT RETURN: SERVICE=INTF; 468 455

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(FINDR,DMY,UIST,CALC); 469 456

WAIT RETURN: SERVICE=INTF; 470 457

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(MODIF,DMY,UIST,DMY); 471 458

WAIT RETURN: SERVICE=INTF; 472 459

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(FINDR,DMY,ADUN,CALC); 473 460

WAIT RETURN: SERVICE=INTF; 474 461

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(MODIF,DMY,ADUN,DMY); 475 462

WAIT RETURN: SERVICE=INTF; 476 463

C NORMAL RETURN 477 464

90 TLOC(2) = 0 478 465

GO TO 150 479 466

C COMPATABILITY ERROR 480 467

THIS PAGE IS BEST QUALITY PRACTICE
FROM COPY FURNISHED TO DDC

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80
INPUT CARD IMAGE
SEC-NO ALT INPUT MEM-PH REF-MU

100 TLOC(2) = 1 468
150 CONTINUE 469
RELEASE: FACILITY=DIYC; 470
END: ENDD SERVICE; 471

THIS PAGE IS BEST QUALITY AVAILABLE
FROM COPY AVAILABLE TO YOU

ENDD SERVICE: ID=INQ, NAPE=IMOS, SAVF AREA SIZE=10, 472
PARAMETER LIST=(PERRI,SEEDI); 473
COMMON /NR/ TLOC 474
INTEGER SUNFM,SEED,SEEDI,TLOC(10) 475
INTEGER OBTAN,FINDR,MODIF,STOP,CONCT,DISCN,ERAS 476
INTEGER DMY,NEXT,PRIOR,OWNER 477

INTEGER ADDI,ADUN,ALOC,GRADE,JDEN,INST,MOSC,POSN,UIST,UNMA 478
INTEGER CALC,GRUI,IDUI,INUM,PMID,POUI,SMID,UMAM,UMGR,UMPOS,(IMU) 479
EQUIVALENCE ((SAVE(1),PERRI), ((SAVE(2),SEED) 480
REAL SUNFM,RN,PERR,PERRI 481
DATA OBTAN,FINDR,MODIF,STOP,CONCT,DISCN,ERAS/1,2,3,4,5,6,7/ 482
DATA DMY,NEXT,PRIOR,OWNER/0,1,2,3/ 483
DATA ADDI,ADUN,ALOC,GRADE,JDEN,INST,MOSC,POSN,UIST,UNMA 484
I / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10/ 485

DATA CALC,GRUI,IDUI,INUM,PMID,POUI,SMID,UMAM,UMGR,UMPOS,UMUI 486
I / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11/ 487
END: DECLARATIONS; 488

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE

.....10.....20.....30.....40.....50.....60.....70.....80

INPUT SEC-DIG LEV SEC-AC REF-CH MEMBER

PERR = PERR1/100.

505 489

SEED = SEEDI

506 490

END: INITIALIZATION;

507 491

SEIZE: FACILITY=INO;

508 492

C

509 493

C

510 494

C

511 495

PARAMETER LIST CONTAINS-

512 496

C 1) DML TYPE - OBTAIN=1, FIRM=2, MODIFY=3, STORE=4, CONNECT=5,

513 497

C DISCONNECT=6, PRAISE=7

514 498

C 2) QUALIFIER - DMY=0, NEXT=1, PRIOR=2, OWNER=3

515 499

C 3) RECORD TYPE - IDENTIFIER OF THE RECORD UPON WHICH THE DML IS

516 500

C TO OPERATE

517 501

C 4) SET IDENTIFIER - INDEX

518 502

C

519 503

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMY, IDEN,CALC);

520 504

WAIT RETURN: SERVICE=INTF;

521 505

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMY, INST,CALC);

522 506

WAIT RETURN: SERVICE=INTF;

523 507

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMY, ADDI,CALC);

524 508

WAIT RETURN: SERVICE=INTF;

525 509

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,NEXT,UIST,IOUI);

526 510

WAIT RETURN: SERVICE=INTF;

527 511

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMY, POSN,CALC);

528 512

.....10.....20.....30.....40.....50.....60.....70.....80

THIS PAGE IS BEST QUALITY PRACTICES FROM XEROX PUBLISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE
SERIAL NO. SEQ-NO. REF-NO.

.....10.....20.....30.....40.....50.....60.....70.....80

WAIT RETURN: SERVICE=INTF; 00005290 529 513

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMY,ALOC,CALC); 00005300 530 514

WAIT RETURN: SERVICE=INTF; 00005310 531 515

C DETERMINE IF COMPATABILITY ERROR HAS OCCURRED 00005320 532 516

RN = SUNFH(0.0, 1.0, SEED) 00005330 533 517

IFIPERR .LT. RN) GO TO 100 00005340 534 518

C NORMAL RETURN 00005350 535 519

90 TLOC(3) = 0 00005360 536 520

GO TO 150 00005370 537 521

C COMPATABILITY ERROR 00005380 538 522

100 TLOC(3) = 1 00005390 539 523

150 CONTINUE 00005400 540 524

RELEASE: FACILITY=INQ; 00005410 541 525

END: ENDO SERVICE; 00005420 542 526

00005430 543

00005440 544

00005450 545

00005460 546 527

00005470 547 528

00005480 548 529

00005490 549 530

00005500 550 531

00005510 551 532

00005520 552 533

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80 INPUT CARD IMAGE

THIS PAGE IS BEST QUALITY PRACTICES FROM COPY FURNISHED TO DDC

INPUT SEC-NO	ALY INPUT LFV SEC-NO	MEMO-CH REF-HE
553	00005530	524
554	00005540	535
555	00005550	536
556	00005560	537
557	00005570	538
558	00005580	539
559	00005590	540
560	00005600	541
561	00005610	542
562	00005620	543
563	00005630	544
564	00005640	545
565	00005650	546
566	00005660	547
567	00005670	548
568	00005680	549
569	00005690	550
570	00005700	551
571	00005710	552
572	00005720	553
573	00005730	554
574	00005740	555
575	00005750	556
576	00005760	557

FORM 8111

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT ALT INPUT MEMPHX
SEQ-NO LEV SIG-NU REF-NO

.....10.....20.....30.....40.....50.....60.....70.....80

577 558

00005770

578 559

00005780

579 560

00005790

580 561

00005800

581 562

00005810

582 563

00005820

583 564

00005830

584 565

00005840

585 566

00005850

586 567

00005860

587 568

00005870

588 569

00005880

589 570

00005890

590 571

00005900

591 572

00005910

592 573

00005920

593 574

00005930

594 575

00005940

595 576

00005950

596 577

00005960

597 578

00005970

598 579

00005980

599 580

00005990

600 581

00006000

THIS PAGE IS BEST QUALITY PRACTICES FROM COPY FACILISHED TO DDC

INPUT CARD IMAGE

.....10.....20.....30.....40.....50.....60.....70.....80

C

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMV,IDEN,CALC);

WAIT RETURN: SERVICE=INTF;

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMV,INST,CALC);

WAIT RETURN: SERVICE=INTF;

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMV,ADDI,CALC);

WAIT RETURN: SERVICE=INTF;

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMV,POSN,CALC);

WAIT RETURN: SERVICE=INTF;

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,NEXT,UIST,IOUI);

WAIT RETURN: SERVICE=INTF;

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,NEXT,UIST,IOUI);

WAIT RETURN: SERVICE=INTF;

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMV,UNMA,CALC);

WAIT RETURN: SERVICE=INTF;

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMV,ADUN,CALC);

WAIT RETURN: SERVICE=INTF;

INVOKE: SERVICE= INTFS, PARAMETER LIST=(OBTAN,DMV,ALOC,CALC);

WAIT RETURN: SERVICE=INTF;

RN = SUNFM(0.0, 1.0, SEED)

IFIRM -LY. PERR) GO TO 100

C

END COMPATABILITY CHECK - NO COMPATABILITY ERROR DETECTED

MODIFY IDENT AND ADD-IDENT RECORDS

C

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT ALT INPUT MEMBER
SEQ-NO LFV SEQ-NO REF-NO

.....10.....20.....30.....40.....50.....60.....70.....80
C 00006010 601 582

INVOKE: SERVICE= INTFS, PARAMETER LYST=(OBTAN,DMV,ALOC,CALC); 602 583

WAIT RETURN: SERVICE=INTF; 603 584

INVOKE: SERVICE= INTFS, PARAMETER LYST=(FINDR,DMV,IDEN,CALC); 604 585

WAIT RETURN: SERVICE=INTF; 605 585

INVOKE: SERVICE= INTFS, PARAMETER LYST=(MODIF,DMV,IDEN,DMV); 606 587

WAIT RETURN: SERVICE=INTF; 607 588

INVOKE: SERVICE= INTFS, PARAMETER LYST=(OBTAN,DMV,ADDI,CALC); 608 589

WAIT RETURN: SERVICE=INTF; 609 590

INVOKE: SERVICE= INTFS, PARAMETER LYST=(MODIF,DMV,ADDI,DMV); 610 591

WAIT RETURN: SERVICE=INTF; 611 592

C 00006120 612 593

C STORE THE NEW UIST RECORD AND ITS CONNECTION IN THE IDENT=UIST 613 594

C SET AND IN THE UNIT=UIST SET 614 595

C 00006150 615 596

INVOKE: SERVICE= INTFS, PARAMETER LYST=(STOR,DMV,UIST,DMV); 616 597

WAIT RETURN: SERVICE=INTF; 617 598

INVOKE: SERVICE= INTFS, PARAMETER LYST=(OBTAN,DMV,GRADE,UMGR); 618 599

WAIT RETURN: SERVICE=INTF; 619 600

INVOKE: SERVICE= INTFS, PARAMETER LYST=(CONCT,DMV,UIST,POUT); 620 601

WAIT RETURN: SERVICE=INTF; 621 602

INVOKE: SERVICE= INTFS, PARAMETER LYST=(OBTAN,DMV,UNMA,CALC); 622 603

WAIT RETURN: SERVICE=INTF; 623 604

INVOKE: SERVICE= INTFS, PARAMETER LYST=(COMCT,DMV,UIST,UMUT); 624 605

.....10.....20.....30.....40.....50.....60.....70.....80

THIS PAGE IS BEST QUALITY PRACTICE
FROM COPY FURNISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

PAGE 27

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

```

.....10.....20.....30.....40.....50.....60.....70.....80.....90
INPUT CARD IMAGE
INPUT SEC=0 ALT INPUT MEMBER
LFY SEC=00 KLF=100
00006250 625 606
WAIT RETURN: SERVICE=INTF;
00006260 626 607
INVOKE: SERVICE= INTFS, PARAMETER LIST=(CONCT,DMY,UJST,GRUI);
00006270 627 608
WAIT RETURN: SERVICE=INTF;
RN = SUNFR(0.0, 1.0, SEED)
00006280 628 609
IFRN .GT. 0.20) GO TO 100
00006290 629 610
INVOKE: SERVICE= INTFS, PARAMETER LIST=(DISCN,DMY,UJST,POUI);
00006300 630 611
WAIT RETURN: SERVICE=INTF;
00006310 631 612
INVOKE: SERVICE= INTFS, PARAMETER LIST=(DISCN,DMY,UJST,UMUI);
00006320 632 613
WAIT RETURN: SERVICE=INTF;
00006330 633 614
INVOKE: SERVICE= INTFS, PARAMETER LIST=(DISCN,DMY,UJST,GRUI);
00006340 634 615
WAIT RETURN: SERVICE=INTF;
00006350 635 616
INVOKE: SERVICE= INTFS, PARAMETER LIST=(ERAS,DMY,UJST,DMY);
00006360 636 617
WAIT RETURN: SERVICE=INTF;
00006370 637 618
C NORMAL RETURN
00006380 638 619
90 TLOC(4) = 0
00006390 639 620
GO TO 150
00006400 640 621
C COMPATABILITY ERROR
00006410 641 622
100 CONTINUE
00006420 642 623
TLOC(4) = 1
00006430 643 624
150 CONTINUE
00006440 644 625
RELEASE: FACILITY=ARRIV;
00006450 645 626
END: ENDU SERVICE;
00006460 646 627
00006470 647
00006480 648

```

THIS PAGE IS BEST QUALITY PRACTICE FROM COPY FURNISHED TO DDC

.....10.....20.....30.....40.....50.....60.....70.....80.....90

INFORMATION PROCESSING SYSTEM SIMULATOR
STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80
INPUT CARD IMAGE

ENDS SERVICE: ID=DEPTR, NAME=DEPTRS, SAVE AREA SIZE=10, INPUT SEC-NO 650 ALT INPUT MEMBER 624
C0006450 649

PARAMETER LIST=(PERRI,SEEDI);

COMMON /WK/ TLOC

INTEGER SUNFMI,SEED,SEEDI,TLOC(10) INPUT SEC-NO 651 ALT INPUT MEMBER 624

INTEGER OBTAN,FINDR,MODIF,STOR,CONCT,DISCN,ERAS INPUT SEC-NO 652 ALT INPUT MEMBER 630

INTEGER DRY,NEXT,PRIOR,OWNER INPUT SEC-NO 653 ALT INPUT MEMBER 631

INTEGER ADDI,ADUN,ALOC,GRADE,IDEN,INST,MOSC,POSN,UJST,UNMA INPUT SEC-NO 654 ALT INPUT MEMBER 632

INTEGER CALC,GRUI,IOUI,INUM,PHIC,POUI,SMID,UMAW,UMGR,UMPOS,UMUI INPUT SEC-NO 655 ALT INPUT MEMBER 633

EQUIVALENCE (SSAVET(1),PERR), (*SAVE(2),SEED)

REAL SUNFM,RN,PERR,PERRI INPUT SEC-NO 656 ALT INPUT MEMBER 634

DATA OBTAN,FINDR,MODIF,STOR,CONCT,DISCN,ERAS(1,2,3,4,5,6,7) INPUT SEC-NO 657 ALT INPUT MEMBER 635

DATA DRY,NEXT,PRIOR,OWNER(0,1,2,3) INPUT SEC-NO 658 ALT INPUT MEMBER 636

DATA ADDI,ADUN,ALOC,GRADE,IDEN,INST,MOSC,POSN,UJST,UNMA INPUT SEC-NO 659 ALT INPUT MEMBER 637

I / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10/ INPUT SEC-NO 660 ALT INPUT MEMBER 638

DATA CALC,GRUI,IOUI,INUM,PHIC,POUI,SMID,UMAW,UMGR,UMPOS,UMUI INPUT SEC-NO 661 ALT INPUT MEMBER 639

I / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11/ INPUT SEC-NO 662 ALT INPUT MEMBER 640

ENDS DECLARATIONS;

PERR = PERRI/100. INPUT SEC-NO 663 ALT INPUT MEMBER 641

SEED = SEEDI INPUT SEC-NO 664 ALT INPUT MEMBER 642

ENDS INITIALIZATION;

SETZE: FACILITY=DEPTR; INPUT SEC-NO 665 ALT INPUT MEMBER 643

C

RETURN CODE -TLOC(5) = 0 IF NO ERROR, 1 IF ERROR INPUT SEC-NO 666 ALT INPUT MEMBER 644

C

.....10.....20.....30.....40.....50.....60.....70.....80

THIS PAGE IS BEST QUALITY PRINTING

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT ALT INPUT NEW LI
SEQ-NO SEQ-NO KEYS

.....10.....20.....30.....40.....50.....60.....70.....80
C 00006730 673 651

PARAMETER LIST CONTAINS-

C 1) DML TYPE - OBTAIN=1, FIND=2, MODIFY=3, STORE=4, CONNECT=5, 674 652

C DISCONNECT=6, ERASE=7 675 653

C 2) QUALIFIER - DMY=0, NEXT=1, PRIOR=2, OWNER=3 676 654

C 2) RECORD TYPE - IDENTIFIER OF THE RECORD UPON WHICH THE DML IS 677 655

TO OPERATE 678 656

C 4) SET IDENTIFIER - INDEX 679 657

C 680 658

C 681 659

C 682 660

C 683 661

C 684 662

C 685 663

C 686 664

C 687 665

C 688 666

C 689 667

C 690 668

C 691 669

C 692 670

C 693 671

C 694 672

C 695 673

C 696 674

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE
SFC=FD LFLV SIC=FD PEF=FD

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....

WAIT RETURN: SERVICE=INTF; 697

C DETERMINE IF COMPATABILITY ERROR HAS OCCURRED 676

RN = SUNFH(0.0, 1.0, SEED) 699

IF(RN .LT. PERR) GO TO 100 700

C NO COMPATABILITY ERROR 674

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(OBTAN,NEXT,UIST,IOUI); 702 680

WAIT RETURN: SERVICE=INTF; 703 681

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(OBTAN,NEXT,UIST,IOUI); 704 682

WAIT RETURN: SERVICE=INTF; 705 683

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(FINDR,DHY,IDEN,CALC); 706 684

WAIT RETURN: SERVICE=INTF; 707 685

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(MODIF,DHY,IDEN,DHY); 708 686

WAIT RETURN: SERVICE=INTF; 709 687

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(FINDR,NEXT,UIST,IOUI); 710 688

WAIT RETURN: SERVICE=INTF; 711 689

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(FINDR,NEXT,UIST,IOUI); 712 690

WAIT RETURN: SERVICE=INTF; 713 691

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(MODIF,DHY,UIST,DHY); 714 692

WAIT RETURN: SERVICE=INTF; 715 693

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(MODIF,DHY,UIST,DHY); 716 694

WAIT RETURN: SERVICE=INTF; 717 695

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(FINDR,DHY,ADDI,CALC); 718 696

WAIT RETURN: SERVICE=INTF; 719 697

INVOKE: SERVICE= INTF\$, PARAMETER LYST=(MODIF,DHY,ADUN,DHY); 720 698

THIS PAGE IS BEST QUALITY PRACTICES FROM COPY FORWARDED TO DDC

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80
 IMPUT ALT INPUT MEMOIF
 SEQ-NO LEV SEQ-NU REF-NU

INTEGER DMY,NEXT,PRIOR,OWNER 745 00007450 740

INTEGER ADDI,ADUN,ALOC,GRADE,IDFN,INST,MOSC,POSN,UIST,URMA 746 00007460 721

INTEGER CALC,GRUI,IDUI,INUM,PHYD,PROY,SMIO,UMAM,UMGR,UMPOS,UMUI 747 00007470 722

EQUIVALENCE (SSAVE(1),PERR), (SSAVE(2),SEED) 748 00007480 723

REAL SUNFN,RN,PERR,PERRI 749 00007490 724

DATA OBTAN,FINDR,MODIF,STOR,CONCT,DISCN,ERAS/1,2,3,4,5,6,7/ 750 00007500 725

DATA DMY,NEXT,PRIOR,OWNER/0,1,2,3/ 751 00007510 726

DATA ADDI,ADUN,ALOC,GRADE,IDFN,INST,MOSC,POSN,UIST,UNMA 752 00007520 727

I / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10/ 753 00007530 728

DATA CALC,GRUI,IDUI,INUM,PHYD,PROY,SMIO,UMAM,UMGR,UMPOS,UMUI 754 00007540 729

I / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11/ 755 00007550 730

END: DECLARATIONS; 756 00007560 731

PERR = PERRI 757 00007570 732

SEED = SEEDI 758 00007580 733

END: INITIALIZATION; 759 00007590 734

SEIZE: FACILITY=ADSL; 760 00007600 735

C RETURN CODE -(LOC(6)) = 0 IF NO ERROR, 1 IF ERROR 761 00007610 736

C 762 00007620 737

C 763 00007630 738

C PARAMETER LIST CONTAINS- 764 00007640 739

C 1) DML TYPE - OBTAIN=1, FIND=2, MODIFY=3, STORE=4, CONNECT=5, 765 00007650 740

C DISCONNECT=6, EPASE=7 766 00007660 741

C 2) QUALIFIER - DMV=0, NEXT=1, PRIOR=2, OWNER=3 767 00007670 742

C 2) RECORD TYPE - IDENTIFIER OF THE RECORD UPON WHICH THE DML IS 00007680 768 00007680 743

.....10.....20.....30.....40.....50.....60.....70.....80

THIS PAGE IS BEST QUALITY PRACTICES
 DOCUMENT PUBLISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80
INPUT CARD IMAGE
SEQ-NO ALT INPUT PEMPLR
LEV SECT-PC KEY-PC

TO OPERATE 749 744

4) SET IDENTIFIER - INDEX 770 745

771 746

772 747

773 748

774 749

775 750

776 751

777 752

778 753

779 754

780 755

791 756

782 757

783 758

784 759

785 760

786 761

787 762

788 763

789 764

790 765

791 766

792 767

THIS PAGE IS A COPY OF THE ORIGINAL DOCUMENT
IT IS NOT A REPRODUCTION OF THE ORIGINAL DOCUMENT

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80
INPUT CARD IMAGE
SEQ-NUM ALT INPUT MEMPH
LEV SEQ-ID REF-NO

INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,NEXT,GRADE,UMGR); 00007930 793 765

WAIT RETURN: SERVICE=INTF; 00007940 794 765

INVOKE: SERVICE= INTFS, PARAMETER LIST=(CONCT,DHY,UIST,GRUI); 00007950 795 770

WAIT RETURN: SERVICE=INTF; 00007960 796 771

INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDK,DHY,MGSC,CALC); 00007970 797 772

WAIT RETURN: SERVICE=INTF; 00007980 798 772

INVOKE: SERVICE= INTFS, PARAMETER LIST=(CONCT,DHY,IDEN,PHID); 00007990 799 774

WAIT RETURN: SERVICE=INTF; 00008000 800 775

I = SUNFMT(1,2,SEED) 00008010 801 776

IF(1.GT. 1) GO TO 200 00008020 802 777

INVOKE: SERVICE= INTFS, PARAMETER LIST=(FINDR,DHY,MGSC,CALC); 00008030 803 778

WAIT RETURN: SERVICE=INTF; 00008040 804 779

INVOKE: SERVICE= INTFS, PARAMETER LIST=(CONCT,DHY,IDEN,SHID); 00008050 805 780

WAIT RETURN: SERVICE=INTF; 00008060 806 781

200 CONTINUE 00008070 807 782

C NORMAL RETURN 00008080 808 783

TLCC(1) = 0 00008090 809 784

GO TO 350 00008100 810 785

C COMPATABILITY ERROR 00008110 811 786

300 CONTINUE 00008120 812 787

TLCC(1) = 1 00008130 813 788

350 CONTINUE 00008140 814 789

RELEASE: FACILITY=ADDSL; 00008150 815 790

END: ENDO SERVICE; 00008160 816 791

.....10.....20.....30.....40.....50.....60.....70.....80

THIS PAGE IS BEST QUALITY PRACTICES FROM COPY FURNISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT ALT INPUT MEMBER
SEC-NO LEV S-UNLU MFP-NO

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00

INPUT CARD IMAGE

00008170 P17

00008180 P18

00008190 P19 792

00008200 P20 793

00008210 P21 794

00008220 P22 795

00008230 P23 796

00008240 P24 797

00008250 P25 798

00008260 P26 799

00008270 P27 800

00008280 P28 801

00008290 P29 802

00008300 P30 803

00008310 P31 804

00008320 P32 805

00008330 P33 806

00008340 P34 807

00008350 P35 808

00008360 P36 809

00008370 P37 810

00008380 P38 811

00008390 P39 812

00008400 P40 813

ENDS SERVICE: ID=INTF, NAME= INTFS, SAVE AREA SIZE=10,

PARAMETER LIST = (P1,P2,P3,P4);

C IOWS INTERFACE ROUTINE

COMMON /WK/ TLOC

INTEGER ERRST,CIND,VALERR

INTEGER DML,OPER,RT,SET,P1,P2,P3,P4

INTEGER SUNFH,SEED,SEEDI,TLOC(10)

INTEGER OBTAN,FINDR,MODIF,STOR,CONCT,DISCN,ERAS

INTEGER DMV,NEXT,PRIOR,OWNER

INTEGER ADDI,ADUN,ALOC,GRADE,IDEN,INST,MOSC,POSN,U1ST,UNMA

INTEGER CALC,GRUI,IDUI,INUM,PHID,POUI,SHID,UMAW,UMGR,UMPOS,UMUI

EQUIVALENCE

1 (SSAVE(6),OPER), (SSAVE(7),RT), (SSAVE(8),SET)

REAL SUNFH,RN

DATA OBTAN,FINDR,MODIF,STOR,CONCT,DISCN,ERAS/1,2,3,4,5,6,7/

DATA DMV,NEXT,PRIOR,OWNER/0,1,2,3/

DATA ADDI,ADUN,ALOC,GRADE,IDEN,INST,MOSC,POSN,U1ST,UNMA

1 / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10/

DATA CALC,GRUI,IDUI,INUM,PHID,POUI,SHID,UMAW,UMGR,UMPOS,UMUI

1 / 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11/

ENDS DECLARATIONS;

DML = P1

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80.....	INPUT	ALT INPUT	MEMBER
								SEC-NO	LFV SEC-NO	REF-NO
	OPFR = P2	00008410	P41							814
	RT = P3	00008420	P42							815
	SEY = P4	00008430	P43							816
	END: INITIALIZATION	00008440	P44							817
	SEIZE: FACILITY=INTF;	00008450	P45							818
C	STEP 1. PERFORM FUNCTION VALIDITY CHECKING AND SEGMENTATION	00008460	P46							819
C	(SYSTEM HOOK)	00008470	P47							820
	VALERR = 0	00008480	P48							821
C	STEP 2. TEST ERROR STATUS AFTER VALIDITY CHECKING	00008490	P49							822
	IF(VALERR .NE. 0) GO TO 80	00008500	P50							823
C	NO VALIDITY ERRORS OCCURRED	00008510	P51							824
C	STEP 3. PERFORM ANY 'BEFORE' PROCEDURES	00008520	P52							825
	(SYSTEM HOOK)	00008530	P53							826
C	NOTE - ERNST IS THE 'ERROR STATUS' INDICATOR	00008540	P54							827
C	CIND IS THE 'CANCEL' INDICATOR	00008550	P55							828
	ERRSY = 0	00008560	P56							829
	CIND = 0	00008570	P57							830
C	STEP 4. TEST THE ERROR STATUS FROM THE 'BEFORE' PROCEDURES	00008580	P58							831
	(SYSTEM HOOK)	00008590	P59							832
	IF(ERRSY .NE. 0) GO TO 80	00008600	P60							833
C	NO ERRORS OCCURRED	00008610	P61							834
C	STEP 5. TEST THE CANCEL INDICATOR	00008620	P62							835
	IF(CIND .NE. 0) GO TO 100	00008630	P63							836
C	DML NOT CANCELLED	00008640	P64							837

THIS PAGE IS BEST QUALITY PRACTICES FROM COPY FURNISHED TO DDC

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT ALT JUMP MEMSIT
SF0-9F LTV SEC-ED PLT-ED

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....

C STEP 6. PERFORM DATA BASE MANAGEMENT FUNCTION

C ALL DML'S ARE ASSUMED TO REQUIRE DRMS SERVICE

C (SYSTEM HOOK FOR FRONT-END PACK-END COMMUNICATION)

INVOKE: SERVICE=CAMPS, PARAMETER LIST=(DML,OPER,RT,SFT);

WAIT RETURN: SERVICE=CAMP;

C STEP 7. TEST ERROR STATUS

IF(ERST .EQ. 0) GO TO 100

C ERROR OCCURRED IN THE EXECUTION OF THE DML

C STEP 8. EXECUTE ANY 'ON ERROR NURSING' PROCEDURES

C (SYSTEM HOOK)

80 CONTINUE

C STEP 9. EXECUTE ANY 'AFTER ERROR' PROCEDURES

C (SYSTEM HOOK)

C STEP 10. PERFORM SEGMENTATION

C (SYSTEM HOOK)

100 CONTINUE

C STEP 11. RETURN CONTROL TO THE APPLICATION PROGRAM

RELEASE: FACILITY=INTF;

END: ENDU SERVICE;

ENDU SERVICE: ID=CAMP, NAME=CAMPS, SAVE AREA SIZE=10,

PARAMETER LIST = (P1,P2,P3,P4);

C ***** CENTRAL ACCESS MONITOR PROGRAM *****

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....

THIS PAGE
FROM COPY MADE BY DDD

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

```

.....10.....20.....30.....40.....50.....60.....70.....80
INPUT CARD IMAGE
SFC-TC LEV SFC-MC REF-NU
C THIS PROGRAM *THREADS REQUESTS AS (THE) DBMS CAN ACCOMMODATE 0C0R8R90 PFC EOC
C THEM - HOLDING THE NECESSARY USERS IN THE WAIT STATE. 00008900 F00 F01
C 0000801G P01 F02
COMMON 7MK/ TL0C 0000E920 P02 b62
INTEGER DML,OPER,RT,SET,PI,P2,P3,P4,TL0C(10) 0000R930 P03 c64
EQUIVALENCE ($SAVE(1),DML), ($SAVE(2),OPER), ($SAVE(3),RT), 00C98940 P04 b75
) ($SAVE(4),SET) 00008950 P05 E66
END: DECLARATIONS: 00008960 P06 c67
DML = PI 00008970 P07 F08
OPER = P2 00008980 P08 P09
RT = P3 00009000 P09 b71
SET = P4 00009010 P10 P72
END: INITIALIZATION: 00009020 P11 E73
SET: FACILITY=CAMP: 00C09030 P12 b74
QUEUE: FACILITY=DBMS:
WAIT FACILITY: FACILITY=DBMS, STATUS=0, CONDITION=EC; 00009040 P13 b75
DEPART QUEUE: FACILITY=DBMS; 00009050 P14 F76
SET STATUS: FACILITY=DBMS, STATUS=1; 00009060 P15 F77
TWORKE: SERVICE=DBMS$, PARAMETER: LYST=(DML,OPER,RT,SET); 00009070 P16 F78
WAIT RETURN: SERVICE=DBMS; 00009080 P17 R79
SET STATUS: FACILITY=DBMS, STATUS=0; 00009090 P18 c80
RELEASE: FACILITY=CAMP; 00009100 P19 b81
END: ENDJ SERVICE: 00009110 P20 F82
00009120 P21
.....10.....20.....30.....40.....50.....60.....70.....80

```

THIS PAGE IS BEST QUALITY PRACTICE
FROM MICROFILM ANALYSIS TO DOC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT MEMETH
FCO-NO (EV SEQ-NU REF-DC

.....10.....20.....30.....40.....50.....60.....70.....80

CC009130 013

014

015

016

017

018

019

020

021

022

023

024

025

026

027

028

029

030

031

032

033

034

035

036

037

038

INPUT CARD IMAGE

ENDJ SERVICE: IO=DBMS, NAME= DBMS, SAVE AREA SIZE=10.

PARAMETER LIST = (PI,P2,P3,P4);

***** DATA BASE MANAGEMENT SYSTEM FUNCTIONS*****

COMMON /WK/ TLDC

COMMON /PE/ PAGFT

COMMON /PAGAR/ FSTPN,LSTPN,AVDIS,RNSAV,ADREC,BUF,RNRET,INDX,NPAG

INTEGER FSTPN,LSTPN,AVDIS,RNSAV,ADREC,BUF(10,3)

INTEGER DML,OPER,RT,SET,PI,P2,P3,P4,TLDC(10),INDX

INTEGER I,J,RN,SUNFMI,RETP,TOTY,TOTZ,PAGFT(10)

EQUIVALENCE (SAVE(5),DML), (SAVE(6),OPER), (SAVE(7),RT),

I (SAVE(8),SET)

BUFFER - BUF(10,3) - FIRST SUBSCRIPT IS BUFFER NUMBER

- SECOND SUBSCRIPT

(1) PAGE NUMBER 10 TO NUMBER OF PAGES IN DB)

(2) LEAST RECENTLY USED (LRU) INDICATOR (0-NPAG)

(3) MUST REWRITE INDICATOR (0-1)

REWRITE PAGE WHEN = 1

PAGE FAULT STATISTICS AREA - PAGFT(10) -

(1) NUMBER OF PAGES CONCURRENTLY IN THE PAGE BUFFER

(2) PAGE SIZE (BYTES)

(3) PAGE WRITES

(4) CACHE MISSES ON WRITE

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....

NOT REPRODUCIBLE TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR
STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80
INPUT CARD IMAGE
SEG-NO LEV SEQ-NO RPT-NO

C (5) PAGE READS 937 506

C (6) CACHE MISSES ON READ 938 507

C (7) NUMBER OF PAGES REWRITTEN BEFORE PAGE SPACE REUSED 939 508

C (8) NUMBER OF PAGE REPLACEMENTS WITHOUT REWRITING 940 509

C (9) NUMBER OF PAGE NUMBERS GENERATED (BY GENER) 941 510

C (10) TOTAL DISTANCE BETWEEN PAGE REFERENCES 942 511

END: DECLARATIONS:

DML = P1

OPER = P2

RY = P3

SET = P4

IF(PAGFY(2) .EQ. 1024) GO TO 1R

DO I=1,1,10

DO J=1,1,3

IO BUF(I,J) = 0

DO I=1,1,10

15 PAGFY(I) = 0

C PAGE SIZE IS 1024 BYTES ON THE PDP1170 AND 3156 ON THE IEM370/16500009540

PAGFY(2) = 1024

FSTPH = 1000

LSTPN = 1200

AVEDIS = 45

RNRFT = 1000

HPAG = 5

THIS PAGE IS BEST QUALITY PRACTICES
FROM COPY FURNISHED TO DDC

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE

.....10.....20.....30.....40.....50.....60.....70.....80

PAGE(1) = MPAG

INDEX = 0

18 CONTINUE

END: INITIALIZATION;

SETZE: FACILITY=DBMS;

C INVOKE THE REQUESTED DML SERVICE

GO TO (100,200,300,400,500,600,700),DML

C OBTAIN DML

100 CONTINUE

INVOKE: SERVICE=OBTAIN\$, PARAMETER LYST=(DML,OPER,RT,SET);

WAIT RETURN: SERVICE=OBTAIN;

GO TO 960

C FIND DML

200 CONTINUE

INVOKE: SERVICE=FIND\$, PARAMETER LYST=(DML,OPER,RT,SET);

WAIT RETURN: SERVICE=FIND;

GO TO 960

C MODIFY DML

300 CONTINUE

INVOKE: SERVICE=MODIFY\$, PARAMETER LYST=(DML,OPER,RT,SET);

WAIT RETURN: SERVICE=MODIFY;

GO TO 900

C STORE DML

400 CONTINUE

.....10.....20.....30.....40.....50.....60.....70.....80

THIS PAGE IS BEST QUALITY PRACTICES FROM COPY FURNISHED TO DOC

THIS PAGE IS BEST QUALITY PRACTICES FROM COPY FURNISHED TO DOC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....

INVOKE: SERVICE=STORES, PARAMETER LYST=(DML,OPER,RT,SET); 00009850 985 554

WAIT RETURN: SERVICE=STORE; 00009860 986 555

GO TO 900 00009870 987 556

C CONNECT DML 00009880 988 557

500 CONTINUE 00009890 989 558

INVOKE: SERVICE=CONNECT, PARAMETER LYST=(DML,OPER,RT,SET); 00009900 990 559

WAIT RETURN: SERVICE=CONNECT; 00009910 991 560

GO TO 900 00009920 992 561

C DISCONNECT DML 00009930 993 562

600 CONTINUE 00009940 994 563

INVOKE: SERVICE=DISCONNECT, PARAMETER LYST=(DML,OPER,RT,SET); 00009950 995 564

WAIT RETURN: SERVICE=DISCONNECT; 00009960 996 565

GO TO 900 00009970 997 566

C ERASE DML 00009980 998 567

700 CONTINUE 00009990 999 568

INVOKE: SERVICE=ERASES, PARAMETER LYST=(DML,OPER,RT,SET); 00010000 1000 569

WAIT RETURN: SERVICE=ERASES; 00010010 1001 570

C 00010020 1002 571

900 CONTINUE 00010030 1003 572

RELEASE: FACILITY=DBMS; 00010040 1004 573

END: END: SERVICE; 00010050 1005 574

00010060 1006 575

00010070 1007 576

00010080 1008 577

00010090 1009 578

00010100 1010 579

THIS PAGE IS BEST QUALITY PRACTICE FROM COPY PARALISED TO DOC

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....

INFORMATION PROCESSING SYSTEM SIMULATOR
STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

```

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....
      INPUT CARD IMAGE
      ***** DBMS OBTAIN DML PROCESSING FUNCTIONS*****
      COMMON /WK/ TLGC
      COMMON /PE/ PAGFY
      COMMON /PAGAR/ FSTPN,LSTPN,AVENIS,RNSAV,POROC,BUF,RNRET,INDX,NPAG
      INTEGER FSTPN,LSTPN,AVENIS,RNSAV,RDRFC,BUF(10,3)
      INTEGER DML,OPER,RT,SEY,PI,P2,P3,P4,TLOC(10),INDX
      INTEGER I,J,RN,SUMFHI,RETP,TOT1,TOT2,PAGFT(10)
      EQUIVALENCE (SSAVE(1),DML), (SSAVE(2),OPER), (SSAVE(3),RT),
      I (SSAVE(4),SEY)
      END: DECLARATIONS:
      DML = P1
      OPER = P2
      RT = P3
      SEY = P4
      END: INITIALIZATION:
      SEIZE: FACILITY=OBTAIN:
      C GENERATE A RANDOM PAGE NUMBER 1001-1200 - RETURNED VALUE IS KNRET
      20 CALL GENRP
      C GET THIS PAGE IF NOT ALREADY PRESENT IN THE CACHE
      INVOKE: SERVICE=PAGMR, PARAMETER LIST=(1);
      WAIT RETURN: SERVICE=PAGMR;
      C DETERMINE IF THE NEXT PAGE IN THE CALC SET NEEDS TO BE ACCESSED
  
```

THIS PAGE IS BEST QUALITY PRACTICES
FROM COPY MAINTAINED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80	INPUT CARD IMAGE	INPUT SEQ-NO	ALT INPUT MEM-NO	MEM-NO
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1033	1002	1002
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1034	1001	1001
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1035	1002	1002
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1026	1003	1003
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1037	1004	1004
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1038	1005	1005
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1039	1006	1006
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1040	1007	1007
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1041	1006	1006
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1042	1006	1006
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1043	1009	1009
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1044	1010	1010
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1045	1011	1011
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1046	1012	1012
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1047	1015	1015
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1048	1016	1016
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1049	1014	1014
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1050	1015	1015
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1051	1016	1016
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1052	1017	1017
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1053	1018	1018
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1054	1019	1019
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1055	1020	1020
.....10.....20.....30.....40.....50.....60.....70.....8010.....20.....30.....40.....50.....60.....70.....80	1056	1021	1021

C THIS IS THE CASE ON 'FIND' AND/OR 'ACQUIRE' FOR CALC SETS

C WHEN THE RECORD COULD NOT BE PLACED ON THE PAGE DETERMINED BY

C THE CALC ALGORITHM

RN = SUNFMI(1,1000,2)

IFIRM .GT. 900) GO TO Z0

C NEXT PAGE IN THE CALC SET NEED NOT BE ACCESSED

C PROCESSING OF 'OBTAIN' FINISHED

RELEASE: FACILITY=OBTAIN;

END: ENDO SERVICE;

END: ENDO SERVICE: ID=FIND, NAME=FIND\$, SAVE AREA SIZE=10,

PARAMETER LST = (PI,P2,P3,P4);

***** DBMS OBTAIN DML PROCESSING FUNCTIONS*****

COMMON /MK/ TLOC

COMMON /PE/ PAGFT

COMMON /PAGAR/ FSTPN,LSTPN,AVEDIS,RNSAV,RDREC,BUF,RNRET,INDX,RPAG

INTEGER FSTPN,LSTPN,AVEDIS,RNSAV,RDREC,BUF(10,3)

INTEGER DML,OPER,RT,SET,PI,P2,P3,P4,VLLOC(10),INDX

INTEGER I,J,RN,SUNFMI,RETP,TOT1,TOT2,PAGFT(10)

EQUIVALENCE ((SAVE(1),DML), ((SAVE(2),OPER), ((SAVE(3),RT),

I ((SAVE(4),SET)

END: DECLARATIONS;

DML = PI

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00

INPUT ALT INPUT MEMBER
SFC-NO LFN SFC-NO LFN-MC

OPER = P2 1027 00010570 1022 1022

RY = P3 1028 00010580 1023 1023

SLY = P4 1029 00010590 1024 1024

END: INITIALIZATION; 1030 00010600 1025 1025

SEIZE: FACILITY=FIND; 1031 00010610 1026 1026

GENERATE A RANDOM PAGE NUMBER 1001-1200 1032 00010620 1027 1027

20 CALL GENRP 1033 00010630 1028 1028

GET THIS PAGE IF NOT ALREADY PRESENT IN THE CACHE 1034 00010640 1029 1029

INVOKE SERVICE=PAGMR\$, PARAMETER LYST=(1); 1035 00010650 1030 1030

WAIT RETURN: SERVICE=PAGMR; 1036 00010660 1031 1031

DETERMINE IF THE NEXT PAGE IN THE CALC SET NEEDS TO BE ACCESSED 1037 00010670 1032 1032

THIS IS THE CASE ON 'FIND' AND/OR 'OBTAIN' FOR CALC SETS 1038 00010680 1033 1033

WHEN THE RECORD COULD NOT BE PLACED ON THE PAGE DETERMINED BY 1039 00010690 1034 1034

THE CALC ALGORITHM 1040 00010700 1035 1035

RM = \$UNFHI(1,1000,3) 1041 00010710 1036 1036

IFRN .GT. 900) GO TO 20 1042 00010720 1037 1037

NEXT PAGE IN THE CALC SET NEED NOT BE ACCESSED 1043 00010730 1038 1038

PROCESSING OF 'FIND' FINISHED 1044 00010740 1039 1039

500 CONTINUE 1045 00010750 1040 1040

RELEASE: FACILITY=FIND; 1046 00010760 1041 1041

END: ENDO SERVICES; 1047 00010770 1042 1042

1048 00010780 1043 1043

1049 00010790 1044 1044

END: ENDO SERVICE: IO=MODIFY, NAME= MODIFS, SAVE AREA SIZE=10; 1050 00010800 1045 1045

1051 00010810 1046 1046

1052 00010820 1047 1047

1053 00010830 1048 1048

1054 00010840 1049 1049

1055 00010850 1050 1050

QUALITY PRACTICES

INFORMATION PROCESSING SYSTEM SIMULATOR
STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....
INPUT CARD IMAGE
MEMBER REF-NO

C ***** DBMS MODIFY DML PROCESSING FUNCTIONS*****
COMMON /WK/ YLOC

PARAMETER LIST = (PI,P2,P3,P4);

COMMON /PE/ PAGFY

COMMON /PAGAR/ FSTPN,LSTPN,AVEDTS,RNSAV,ROREC,BUF,RNREY,INDX,NPAG

INTEGER FSTPN,LSTPN,AVEDIS,RNSAV,ROREC,BUFF(10,3)

INTEGER DML,OPER,RT,SET,PI,P2,P3,P4,YLOC(10),INDX

INTEGER I,J,RN,UNFHI,RETP,TOTI,TOTZ,PAGFY(10)

EQUIVALENCE (\$SAVE(1),DML), (\$SAVE(2),OPER), (\$SAVE(3),RT),

I (\$SAVE(4),SET)

END: DECLARATIONS;

DML = P1

OPER = P2

RT = P3

SET = P4

END: INITIALIZATION;

SEIZE: FACILITY=MODIFY;

C WALGREEN DATA - 50% READ & 100% WRITE MISSES

C (SYSTEM HOOK)

PROCESS: TIME=14.7;

RELEASE: FACILITY=MODIFY;

END: ENDO SERVICE;

END: SERVICE: ID=STORE, NAME= STORE\$, SAVE AREA SIZE=10,

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....

THIS PAGE IS BEST QUALITY PRACTICE
FROM COPY FURNISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

```

.....10.....20.....30.....40.....50.....60.....70.....80
INPUT CARD IMAGE
SEC-NO  ALT INPUT MEMCHK
LEV  SLU-NC  K1P-M

PARAMETER LIST = (P1,P2,P3,P4);
***** DBMS STORE DML PROCESSING FUNCTIONS*****
COMMON /WK/ TLDC
COMMON /PE/ PAGFT
COMMON /PAGAR/ FSTPN,LSTPN,AVEDIS,RNSAV,RDREC,BUF,RNRET,INDX,NPAG
INTEGER FSTPN,LSTPN,AVEDIS,RNSAV,RDREC,BUF(10,3)
INTEGER DML,OPEK,RT,SET,P1,P2,P3,P4,TLOC(10),JMXD
INTEGER I,J,RN,SUNFMI,RETP,TOT1,TOT2,PAGFT(10)
EQUIVALENCE (SSAVE(1),DML), (SSAVE(2),OPER), (SSAVE(3),RT),
              (SSAVE(4),SET)
END: DECLARATIONS;
DML = P1
OPER = P2
RT = P3
SET = P4
END: INITIALIZATION;
SETZE: FACILITY=STORE;
C DETERMINE IF THE *SPACE MANAGEMENT PAGE* NEEDS TO BE ACCESSED
C (ON *STORE* WHEN NOT ENOUGH SPACE AVAILABLE IN PRESENT PAGE)
RN = SUNFMI(1,1000,4)
IFRN .LT. 990) GO TO 200
C MUST ACCESS THE SPACE MANAGEMENT PAGE
C GET THIS PAGE IF NOT ALREADY IN THE CACHE
RNRET = FSTPN

```

THIS PAGE IS BEST QUALITY REPRODUCTION FROM COPY PARALLEL TO DDC

.....10.....20.....30.....40.....50.....60.....70.....80

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80
INPUT CARD IMAGE

INVOKE: SERVICE=PAGRRS, PARAMETER LIST=(1); 1091

WAIT RETURN: SERVICE=PAGRR; 1092

C ACCESS THE PAGE ON WHICH SPACE IS AVAILABLE(POINTED TO BY THE 1092

C SPACE MANAGEMENT PAGE) 1094

C GENERATE A RANDOM PAGE NUMBER 1095

CALL GENRP 1096

C GET SPACE-AVAILABLE PAGE IF NOT ALREADY IN THE CACHE 1097

INVOKE: SERVICE=PAGRRS, PARAMETER LIST=(2); 1098

WAIT RETURN: SERVICE=PAGRR; 1099

C DETERMINE IF THE *SPACE MANAGEMENT PAGE* MUST BE REWRITTEN 1100

C (ONLY WHEN PAGE IS > 70% FULL) 1101

RN = SUNFMI(1,1000,4) 1102

IFRN .LY. 990) GO TO 200 1103

C REWRITE NO-SPACE-AVAILABLE PAGE TO UPDATE POINTER TO NEXT IN CALC 1104

C SET 1105

C DETERMINE IF NO-SPACE-AVAILABLE PAGE IS IN THE CACHE 1106

200 CONTINUE 1107

INVOKE: SERVICE=PAGRRS, PARAMETER LIST=(2); 1108

WAIT RETURN: SERVICE=PAGRR; 1109

RELEASE: FACILITY=STORE; 1110

END: ENDO SERVICE; 1111

.....10.....20.....30.....40.....50.....60.....70.....80

END SERVICE: ID=CONNECT, NAME=CONNECTS, SAVE AREA SIZE=10, 1112

.....10.....20.....30.....40.....50.....60.....70.....80

THIS PAGE IS BEST QUALITY PRACTICES FROM COPY FURNISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00

INPUT CARD IMAGE
SEC=0 ALT INPUT MEMBER
LEV=00-TC EXT=0

PARAMETER LIST = (P1,P2,P3,P4);

C ***** DBMS CONNECT DML PROCESSING FUNCTIONS*****

COMMON /WK/ TLOC

COMMON /PE/ PAGFY

COMMON /PAGAR/ FSTPN,LSTPN,AVFOIS,RNSAV,RDREC,BUF,RNRET,INDX,MPAG 00011570

INTEGR FSTPN,LSTPN,AVEOIS,RNSAV,RDREC,BUF(10,3)

INTEGR DML,OPER,RT,SET,PI,P2,P3,P4,TLOC(1G),INDX 00011590

INTEGR I,J,RN,SUNFHI,KEIP,TPY,INT2,PAGFY(10) 00011600

EQUIVALENCE (\$SAVE(1),DML), (\$SAVE(2),OPER), (\$SAVE(3),RT), 00011610

(\$SAVE(4),SET), (\$SAVE(7),RN), (\$SAVE(8),I) 00011620

END: DECLARATIONS:

DML = P1 00011630

OPER = P2 00011640

RT = P3 00011650

SET = P4 00011660

END: INITIALIZATION:

SETZF: FACILITY=CONNECT; 00011670

C ASSUME UNIFORM DISTRIBUTION OF 1-5 MEMBER RECORDS TO BE CONNECTED 00011700

RN = \$UNFHI(1,5,5) 00011710

I = 0 00011720

150 I = I + 1 00011730

IF (I .GT. RN) GO TO 500 00011740

CALL GENRP 00011750

INVOKE: \$SERVICE=PAGARS, PARAMETER LIST=(1); 00011760

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00

THIS PAGE IS BEST QUALITY PRACTICE
FROM COPY PARALISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT	ALT INPUT	MEMBER	
SFC-P.O	LEV	SEC-P.O	REF-NU
00011770	1177	1127	
00011780	1178	1138	
00011790	1179	1139	
00011800	1180	1140	
00011810	1181	1141	
00011820	1182		
00011830	1183		
00011840	1184	1142	
00011850	1185	1143	
00011860	1186	1144	
00011870	1187	1145	
00011880	1188	1146	
00011890	1189	1147	
00011900	1190	1148	
00011910	1191	1149	
00011920	1192	1150	
00011930	1193	1151	
00011940	1194	1152	
00011950	1195	1153	
00011960	1196	1154	
00011970	1197	1155	
00011980	1198	1156	
00011990	1199	1157	
00012000	1200	1158	

WAIT RETURN: SERVICE=PAGMP;

GO TO 150

500 CONTINUE

RELEASE: FACILITY=CONNECT;

END: ENDO SERVICE;

ENDO SERVICE: ID=DISCON, NAME= DYSO\$, SAVE AREA SIZE=10,

PARAMETER LIST = (P1,P2,P3,P4);

C ***** DBMS DISCONNECT PROCESSING FUNCTIONS*****

COMMON /HK/ TLOC

COMMON /PE/ PAGFT

COMMON /PAGAR/ FSTPN,LSTPN,AVFOIS,RNSAV,RDREC,BUF,RNRET,INDX,MPAG 00011600

INTEGER FSTPN,LSTPN,AVEDIS,RNSAV,RDREC,BUF(10,3)

INTEGER DML,OPER,RT,SET,P1,P2,P3,P4,TLOC(10),INDX

INTEGER I,J,RN,SUNFH,RETP,TOY,TOY2,PAGFT(10)

EQUIVALENCE (\$SAVE(1),DML), (\$SAVE(2),OPER), (\$SAVE(3),RT),

I (\$SAVE(4),SET), (\$SAVE(7),RN), (\$SAVE(8),I)

END: DECLARATIONS;

DML = P1

OPER = P2

RT = P3

SET = P4

END: INITIALIZATION;

.....10.....20.....30.....40.....50.....60.....70.....80

THIS PAGE IS BEST QUALITY PRACTICABLE FROM COPY FINALISED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE
MEMBER
MEMBER
MEMBER

.....10.....20.....30.....40.....50.....60.....70.....80.....
INPUT CARD IMAGE
MEMBER
MEMBER
MEMBER

SETZE: FACILITY=DISCON; 1201 1159

C ASSUME UNIFORM DISTRIBUTION OF 1-5 RECORDS TO BE DISCONNECTED 1202 1160

RN = \$UNFMI(1,5,6) 1203 1161

I = 0 1204 1162

150 I = I + 1 1205 1163

IF (I .GT. RN) GO TO 500 1206 1164

CALL GENRP 1207 1165

INVOKE: SERVICE=PAGMR, PARAMETER LIST=(1); 1208 1166

WAIT RETURN: SERVICE=PAGMR; 1209 1167

GO TO 150 1210 1168

500 CONTINUE 1211 1169

RELEASE: FACILITY=DISCON; 1212 1170

END: ENDD SERVICE; 1213 1171

ENDG SERVICE: IO=ERASE, NAME= ERASES, SAVE AREA SIZE=10, 1214 1172

PARAMETER LIST = (PI,P2,P3,P4); 1215 1173

***** DBMS ERASE DML PROCESSING FUNCTIONS***** 1216 1174

COMMON /MK/ TLUC 1217 1175

COMMON /PE/ PAGFT 1218 1176

COMMON /PAGAR/ FSTPN,LSTPN,AVEDIS,RNSAV,RDREC,BUF,RMRET,INDX,MPAG 00012219 1177

INTEGR FSTPN,LSTPN,AVEDIS,RNSAV,RDREC,BUF(10,3) 1220 1178

INTEGER DML,OPER,KT,SET,PI,P2,P3,P4,TLOC(10),INDX 1221 1179

INTEGER I,J,K,M,SUNFMI,RETP,TOT1,TOT2,PAGFT(10) 1224 1180

.....10.....20.....30.....40.....50.....60.....70.....80.....

THIS PAGE IS BEST QUALITY PRACTICES
FROM COPY AVAILABLE TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR
STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT METHOD
REC-NO SEQ-NR

REC-NO	SEQ-NR	INPUT METHOD	MEMBER
1	1225	00012250	1181
1	1226	00012260	1182
	1227	00012270	1183
	1228	00012280	1184
	1229	00012290	1185
	1230	00012300	1186
	1231	00012310	1187
	1232	00012320	1188
	1233	00012330	1189
	1234	00012340	1190
	1235	00012350	1191
	1236	00012360	1192
	1237	00012370	1193
	1238	00012380	1194
	1239	00012390	1195
	1240	00012400	1196
	1241	00012410	1197
	1242	00012420	1198
	1243	00012430	1199
	1244	00012440	1200
	1245	00012450	1201
	1246	00012460	1202
	1247	00012470	1203
	1248	00012480	1204

EQUIVALENCE (\$SAVE(1),DML), (\$SAVE(2),OPER), (\$SAVE(3),RT),
 (\$SAVE(4),SET), (\$SAVE(5),RN), (\$SAVE(6),I)
 END: DECLARATIONS:
 DML = P1
 OPER = P2
 RY = P3
 SEY = P4
 END: INITIALIZATION:
 SEIZF: FACILITY=ERASE:
 WALGREEN DATA - 50% READ & 100% WRITE MISS RATIO
 %ERASE ALL MEMBERS - .3% (101,360 MSEC AVE)
 %ERASE <RECORD NAME> - 99.7% (49 MSEC AVE)
 RN = \$UNFH(1,1000,7)
 IF(RN .GT. 997) GO TO 100
 PROCESS %ERASE <RECORD NAME> DML
 CALL GENRP
 INVOKE: SERVICE=PAGMR\$, PARAMETER LTST=(1);
 WAIT RETURN: SERVICE=PAGMR\$
 GO TO 500
 PROCESS %ERASE ALL MEMBERS' STATEMENT
 100 CONTINUE
 ASSUME UNIFORM DISTRIBUTION OF 1 TO 5 MEMBER RECORDS
 RN = \$UNFH(1,5,7)
 I = 0

THIS PAGE IS BEST QUALITY PRACTICES
FROM COPY PARALLEL TO DDC

.....1C.....20.....30.....40.....50.....60.....70.....80.....90.....

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE
.....10.....20.....30.....40.....50.....60.....70.....80

150 Y = I + J 1249 00012490 1205

IF (I .GT. RN) GO TO 500 1250 00012500 1206

CALL GENRP 1251 00012510 1207

INVOKE: SERVICE=PAGMR\$, PARAMETER LIST=(I); 1252 00012520 1208

WAIT RETURN: SERVICE=PAGMR; 1253 00012530 1209

GO TO 150 1254 00012540 1210

500 CONTINUE 1255 00012550 1211

RELEASE: FACILITY=ERASE; 1256 00012560 1212

END: ENDO SERVICE; 1257 00012570 1213

END SERVICE: ID=PAGMR, NAME= PAGMR\$, PARAMETER LIST=(RW), 1258 00012580 1214

SAVE AREA SIZE=10; 1259 00012590 1215

***** BUFFER MANAGEMENT ROUTINE***** 1260 00012600 1216

(IDMS' MODIFIED LEAST-RECENTLY-USED ALGORITHM) 1261 00012610 1217

THIS ROUTINE IS RESPONSIBLE FOR BRINGING THE REQUESTED PAGE 1262 00012620 1218

(I.E., RNRET) INTO THE PAGE BUFFER. THIS MAY INVOLVE REWRITING 1263 00012630 1219

ONE OF THE PAGES IN THE BUFFER. 1264 00012640 1220

INPUT PARAMETER - 1=READ, 2=WRITE 1265 00012650 1221

COMMON /WK/ TLDC 1266 00012660 1222

COMMON /PE/ PAGFT 1267 00012670 1223

COMMON /PAGAR/ FSTPN, LSTPN, AVEPTS, RN SAV, RUREC, RUF, RNRET, INDX, NPAG 00012700 1224

INTEGER FSTPN, LSTPN, AVEPTS, RN SAV, RUREC, RUF(10,3) 1271 00012710 1225

INTEGER DML-OPER, RT, SET, PI, P2, P3, P4, TLOC(10), INDX, RW, RDWT 1272 00012720 1226

.....10.....20.....30.....40.....50.....60.....70.....80

THIS PAGE IS THE BEST QUALITY AVAILABLE FROM OUR PUBLISHING SERVICE

THIS PAGE IS BEST QUALITY PRACTICES
FROM COPY FURNISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR
STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80.....
INPUT CARD IMAGE
SF0=NO LIV SEC=LU PFF=LU

INTEGER I,J,RN,SUNFMI,RETP,YOY1,YOY2,PAGFT(10),IX 1273 1271

EQUVALENCE (SSAVE(1),DML), (SSAVE(2),OPER), (SSAVE(3),RT), 1274 1226

I (SSAVE(4),SET), (SSAVE(5),RN), (SSAVE(9),IX), (SSAVE(10),RDWT) 1275 1229

C PAGE FAULT STATISTICS AREA - PAGFT(10) - 1276 1250

C (1) NUMBER OF PAGES CONCURRENTLY IN THE PAGE BUFFER 1277 1251

C (2) PAGE SIZE (BYTES) 1278 1232

C (3) PAGE WRITES 1279 1231

C (4) CACHE MISSES ON WRITE 1280 1234

C (5) PAGE READS 1281 1235

C (6) CACHE MISSES ON READ 1282 1236

C (7) NUMBER OF PAGES REWRITTEN BEFORE PAGE SPACE REUSED 1283 1237

C (8) NUMBER OF PAGE REPLACEMENTS WITHOUT REWRITING 1284 1238

C (9) NUMBER OF PAGE NUMBERS GENERATED (BY GENPR) 1285 1239

C (10) TOTAL DISTANCE BETWEEN PAGE REFERENCES 1286 1240

END: DECLARATIONS: 1287 1241

RN = RWREY 1288 1242

RDWT = RW 1289 1243

END: INITIALIZATION: 1290 1244

SEIZE: FACILITY= PAGMR; 1291 1245

IFIRDWT .EQ. 1) RDWT = 1 1292 1246

IFIRDWT .GT. 2) RDWT = 2 1293 1247

C UPDATE NUMBER OF READS/WRITES 1294 1248

IFIRDWT .EQ. 1) PAGFT(5) = PAGFT(5) + 1 1295 1249

IFIRDWT .EQ. 2) PAGFT(3) = PAGFT(3) + 1 1296 1250

.....10.....20.....30.....40.....50.....60.....70.....80.....

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00

C DETERMINE IF PAGE IS IN BUFFER POOL

DO IS I = 1, NPAG

IX = 1

IF(BUFTX,1) .EQ. RN) GO TO 70

TO CONTINUE

C PAGE WAS NOT IN THE BUFFER POOL

IF(RDWT .EQ. 1) PAGFT(6) = PAGFT(6) + 1

IF(RDWT .EQ. 2) PAGFT(4) = PAGFT(4) + 1

C SIMULATE RECORD RETRIEVAL FROM DATA BASE - ELAPSED TIME = 36 MS.

C (SYSTEM HOOK)

PROCESS: TIME=36.0;

C STORE THE PAGE IN THE BUFFER POOL

IX = 0

30 IX = IX + 1

IF(IX .GT. NPAG) GO TO 70

IF(BUFTX,2) .NE. 0) GO TO 30

C DETERMINE IF THE PAGE MUST BE WRITTEN BEFORE IT IS REPLACED

IF(BUFTX,3) .EQ. 0) GO TO 40

C PAGE TO BE REPLACED HAS BEEN MODIFIED - REWRITE

PAGFT(7) = PAGFT(7) + 1

C SIMULATE RECORD STORAGE TO DATA BASE - ELAPSED TIME = 36 MS.

C (SYSTEM HOOK)

PROCESS: TIME=36.0;

GO TO 45

THIS PAGE IS BEST QUALITY PRACTICES FROM COPY FURNISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT SEQ-NO	ALT LFLV	INPUT SEQ-NO	MEMBER REF-NO
1321		1321	1275
1322		1322	1276
1323		1323	1277
1324		1324	1278
1325		1325	1279
1326		1326	1280
1327		1327	1281
1328		1328	1282
1329		1329	1283
1330		1330	1284
1331		1331	1285
1332		1332	1286
1333		1333	1287
1334		1334	1288
1335		1335	1289
1336		1336	1290
1337		1337	1291
1338		1338	1292
1339		1339	1293
1340		1340	1294
1341		1341	1295
1342		1342	1296
1343		1343	1297
1344		1344	1298

40 PAGFT(8) = PAGFT(8) + 1

C AGE ALL EXISTING PAGES IN THE BUFFER POOL

45 DO 50 J = 1, NPAG

IF(BUF(J,2) .GT. 0) BUF(J,2) = BUF(J,2) - 1

50 CONTINUE

C STORE THE NEW PAGE IN THE BUFFER POOL

BUF(IX,1) = RN

BUF(IX,2) = NPAG - 1

BUF(IX,3) = 0

C

C

C

70 INDX = IX

IF(RDWT .EQ. 1) GO TO 80

C TURN MUST-WRITE SWITCH 'ON'

BUF(INDX,3) = 1

80 CONTINUE

RELEASE FACILITY=PAGMR;

END: ENDD SERVICE;

C

C

C

C

C

C

C

C

C

C

THIS PAGE IS BEST QUALITY PRACTICES FROM COPY FURNISHED TO DDC

FORM 5411 1

INFORMATION PROCESSING SYSTEM SIMULATOR

PAGE 27

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE

INPUT ALT INPUT PERM-BK
SEC-NR LEV SEU-FC REF-FC

.....10.....20.....30.....40.....50.....60.....70.....80

C DISTANCE BETWEEN SUCCESSIVE PAGES. USING THIS ALGORITHM, 00013450 1345 1300

C 'CLUSTERED' PAGE REFERENCES CAN BE GENERATED REGARDLESS OF THE 00013460 1346 1298

C SIZE OF THE DATA BASE. 00013470 1347 1299

C 00013480 1348 1300

C FSTPN - LOW ORDER BOUND ON GENERATED PAGE NUMBERS 00013490 1349 1301

C LSTPN - HIGH ORDER BOUND ON GENERATED PAGE NUMBERS 00013500 1350 1302

C AVEDIS - THE DESIRED AVERAGE DISTANCE BETWEEN SUCCESSIVE PAGE 00013510 1351 1303

C REFERENCES 00013520 1352 1304

C RNRET - THE COMPUTED PAGE NUMBER 00013530 1353 1305

C COMMON 7NK/ TLOC 00013540 1354 1306

C COMMON 7PE/ PAGFT 00013550 1355 1307

C COMMON 7PAGAR/ FSTPN,LSTPN,AVEDIS,RNSAV,KDREC,BUF,RNRET,INDX,NPAC 00013560 1356 1308

C INTEGFA FSTPN,LSTPN,AVEDIS,RNSAV,KDREC,BUF(10,3) 00013570 1357 1309

C INTEGER OML,OPER,RT,SET,P1,P2,P3,P4,TLOC(10),INDX 00013580 1358 1310

C INTEGER I,J,RN,SUNFMI,RETP,TOT1,TOT2,PAGFT(10),K,L 00013590 1359 1311

C INTEGER KNI,KN2,VARI,VAR2 00013600 1360 1312

C REAL SUNFM,Y,RN3 00013610 1361 1313

C INITIALIZE 00013620 1362 1314

C MAXPN = LSTPN - FSTPN 00013630 1363 1315

C IF(MAXPN .LE. 0) GO TO 100 00013640 1364 1316

C RNI = RNRET 00013650 1365 1317

C COMPUTE KN2 USING THE FULL RANGE OF PAGES 00013660 1366 1318

C TO RN2 = SUNFMI(FSTPN,LSTPN,8) 00013670 1367 1319

C VARI = RNI + AVEDIS - FSTPN 00013680 1368 1320

THIS PAGE IS BEST QUALITY PRINTING FROM COPY FURNISHED TO DDC

.....10.....20.....30.....40.....50.....60.....70.....80

THIS PAGE IS BEST QUALITY PRACTICES
FROM COPY FURNISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR
STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

```

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00
INPUT CARD IMAGE
.....40.....50.....60.....70.....80.....90.....00
IF(VARI .GT. MAXPN) VARI = MAXPN
VAR2 = RNI - AVEDIS - FSTPN
IF(VAR2 .LT. 0) VAR2 = 0
C DON'T RECOMPUTE RN2 IF NOT OUTSIDE RANGE RNI+-AVEDIS
IF(RN2 .GT. RNI+AVEDIS) GO TO 20
IF(RN2 .LT. RNI-AVEDIS) GO TO 20
C GO TO 50
20 CONTINUE
C Y = PROBRN2 > RNI+VAR) OR PROBRN2 < RNI-VAR)
IF(RN2 .GT. RNI) Y = 1. - (VARI/(MAXPN*1.0))
IF(RN2 .LT. RNI) Y = (VAR2/(MAXPN*1.0))
C TEST BOUNDS ON Y - ALGORITHM COULD LOOP - EXIT IF OUT OF RANGE
IF(Y .LT. 0.) GO TO 50
IF(Y .GT. 1.) GO TO 50
C GENERATE RN3 TO TEST AGAINST Y
RN3 = $UNFMT(0.,1.,0.,8)
30 CONTINUE
IF(RN3 .GE. Y) GO TO 50
C RECOMPUTE RN2 - RANGE CRITERIA NOT MET (ON AVERAGE)
K = RNI - AVEDIS
IF(K .LT. FSTPN) K = FSTPN
L = RNI + AVEDIS
IF(L .GT. LSTPN) L = LSTPN
RN2 = $UNFMT(K.,L.,8)
.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00

```

INPUT SEC-NO	MEMBER NO
1360	1321
1370	1322
1371	1323
1372	1324
1373	1325
1374	1326
1375	1327
1376	1328
1377	1329
1378	1330
1379	1331
1380	1332
1381	1333
1382	1334
1383	1335
1384	1336
1385	1337
1386	1338
1387	1339
1388	1340
1389	1341
1390	1342
1391	1343
1392	1344

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 05/10/77

INPUT CARD IMAGE

INPUT ALT INPUT PIMLFR
SEQ-NO LEV SEQ-NO REF-AL

10.....20.....30.....40.....50.....60.....70.....80
IDIF = (RN2-RN1) 00013930 1343 1345

C AVERAGE DISTANCE CRITERIA IS SATISFIED - UPDATE VARIABLES & RETURN00013940 1394 1346

50 RNRET = RN2 00013950 1395 1347

PAGFT(9) = PAGFT(9) + 1 00013960 1396 1348

PAGFT(10) = PAGFT(10) + IARS(RN1-RN2) 00013970 1397 1349

IDIF = (RN2-RN1) 00013980 1398 1350

C BYPASS SECOND WRITE OF GENRP DATA 00013990 1399 1351

IFIRN3 -LT. Y) GG TO 500 00014000 1400 1352

GG TO 500 00014010 1401 1353

100 CONTINUE 00014020 1402 1354

C ERROR RETURN ON FSTPN LSTPN VALUES - FSTPN > LSTPN 00014030 1403 1355

RNRET = \$UNFRI(LSTPN,FSTPN,8) 00014040 1404 1356

500 CONTINUE 00014050 1405 1357

RETURN 00014060 1406 1358

END: PROCEDURE: 00014070 1407 1359

00014080 1408

00014090 1409

00014100 1410

00014110 1411

C PRINT BUFFER PAGE FAULT STATISTICS 00014120 1412 1362

COMMON /PE/ PAGFT 00014130 1413 1363

INTEGER PAGFT(10),TOT1,TOT2 00014140 1414 1364

REAL MISR 00014150 1415 1365

END: DECLARATIONS: 00014160 1416 1366

10.....20.....30.....40.....50.....60.....70.....80

THIS PAGE IS BEST QUALITY PRINTING
FROM COPY PARALSHED TO DDC

AD-A069 543

OHIO STATE UNIV RESEARCH FOUNDATION COLUMBUS
A SIMULATION LANGUAGE FOR EVALUATING INFORMATION PROCESSING SYS--ETC(U)
SEP 78 T 6 DELUTIS

F/G 9/4

DAAG29-77-G-0203

UNCLASSIFIED

ARO-15356.1-A-M

NL

3 OF 3
AD
A069543



END
DATE
FILMED
7-79
DDC

INFORMATION PROCESSING SYSTEM SIMULATOR
STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE
SFCTO ALT INPUT MMLER
LIV SEC-TO PFF-NU

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....	INPUT	ALT INPUT	MMLER
.....10.....20.....30.....40.....50.....60.....70.....80.....90.....	SFCTO	LIV SEC-TO	PFF-NU
END: INITIALIZATION;	00014170	1417	1467
WRITE(6,10)	000141F0	141F	1468
10 FORMAT('1',10X,'CACHE PAGE FAULT STATISTICS//')	00014190	1419	1469
WRITE(6,20) PAGFT(1)	00014200	1420	1470
20 FORMAT('1',10X,'CACHE SIZE (# PAGES): ',12)	00014210	1421	1471
WRITE(6,30) PAGFT(2)	00014220	1422	1472
30 FORMAT('1',10X,'PAGE SIZE (BYTES): ',14)	00014230	1423	1473
MISR = PAGFT(4)/(PAGFT(3)*1.0)	00014240	1424	1474
WRITE(6,40) PAGFT(3),PAGFT(4),MISR	00014250	1425	1475
40 FORMAT('0',10X,'PAGE WRITE REFERENCES: ',16,	00014260	1426	1476
1 3X,'CACHE MISSES:',16,3X,'MISS RATIO: ',F6.4)	00014270	1427	1477
MISR = PAGFT(6)/(PAGFT(5)*1.0)	00014280	1428	1478
WRITE(6,50) PAGFT(5),PAGFT(6),MISR	00014290	1429	1479
50 FORMAT('1',10X,'PAGE READ REFERENCES: ',16,	00014300	1430	1480
1 3X,'CACHE MISSES:',16,3X,'MISS RATIO: ',F6.4)	00014310	1431	1481
TOT1 = PAGFT(3) + PAGFT(5)	00014320	1432	1482
TOT2 = PAGFT(4) + PAGFT(6)	00014330	1433	1483
MISR = TOT2/(TOT1*1.0)	00014340	1434	1484
WRITE(6,60) TOT1,TOT2,MISR	00014350	1435	1485
60 FORMAT('0',10X,'PAGE TOTAL REFERENCES: ',16,	00014360	1436	1486
1 3X,'CACHE MISSES:',16,3X,'MISS RATIO: ',F6.4)	00014370	1437	1487
WRITE(6,70) PAGFT(7)	00014380	1438	1488
70 FORMAT('0',10X,'PAGES REWRITEN BEFORE SPACE REUSED: ',16)	00014390	1439	1489
WRITE(6,80) PAGFT(8)	00014400	1440	1490

THIS PAGE IS BEST QUALITY FRAGMENTED FROM COPY FURNISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE
SFC=MM ALT INPUT MEM=K
LEV=50-60 REF=PC

80 FORMAT(' ,10X,'PAGES REPLACED WITHOUT REWRITING: ',I6) 00014410 1447 1392

MISR = PAGET(10)/(PAGET(9)+1.0) 00014420 1442 1392

WRITE(6,90) MISR 00014430 1443 1393

90 FORMAT(' ,10X,'AVERAGE DISTANCE BETWEEN SUCCESSIVE PAGES', 00014440 1444 1394

I * REFERENCES: ' ,F9.4//') 00014450 1445 1395

END: EXO SERVICE: 00014460 1446 1396

00014470 1447

00014480 1448

00014490 1449 1397

END SERVICE: ID=SNAPST, NAME=SNAP\$, 00014500 1450 1398

SAVE AREA SIZE = 10; 00014510 1451 1399

C PRINT STANDARD IPSS STATISTICS AT THE DESIGNATED INTERVALS 00014520 1452 1400

END: DECLARATIONS: 00014530 1453 1401

END: INITIALIZATION: 00014540 1454 1402

CALL \$SNAP 00014550 1455 1403

END: EXO SERVICE: 00014560 1456

00014570 1457

00014580 1458 1404

END: SYSTEM RESOURCES: 00014590 1459

00014600 1460

00014610 1461

DEFINE EXOGENOUS EVENT STREAM: NAME=FX, COMPILE=YES, 00014620 1462

DISPOSITION=KEEP; 00014630 1463

00014640 1464

EXOGENOUS EVENT: IO=EVI, TIME=PROCTIME); 00014650 1465

00014660 1466

00014670 1467

00014680 1468

00014690 1469

00014700 1470

THIS PAGE IS BEST QUALITY PRINTED COPY AVAILABLE TO YOU

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00.....
			INPUT CARD IMAGE						

EXOGENOUS EVENT: ID=EV2, TIME=PPROC(TIME2);

EXOGENOUS EVENT: ID=EV3, TIME=PPROC(TIME3);

PROCEDURE: NAME=TIME1, TYPE=SUBROUTINE;

THIS MODEL ASSUMES-

AN AVERAGE OF 8 TRANSACTIONS PER MONTH PER SOLDIER AND A DATA

BASE OF 203 SOLDIERS. THE CHARACTERIZATION OF EACH

TRANSACTION IS GIVEN IN THE 'TARM' ARRAY. THE AVERAGE STAY

PER SOLDIER IS ASSUMED TO BE 6 MONTHS

GIVEN THESE ASSUMPTIONS, THIS EXO SERVICE-

IS INVOKED ONCE EVERY 5.89 MINUTES AND STARTS A RANDOMLY

SELECTED TRANSACTION ACCORDING TO THE PROBABILITY DISTRIBUTION

FUNCTION GIVEN IN THE 'POINTS' ARRAY IN 'START' EXO SERVICE.

NOTE- THE DATA USED IN CALCULATING THE TRANSACTION ARRIVAL TIMES

WAS BASED ON 20 WORKING DAYS PER MONTH AND 8 HOURS OF TERMINAL

INPUT PER DAY.

COMMON /SYSTEM/ %CLOCK,%CLRI,%CLR?

COMMON /SYSINF/ SYSINF

INTEGER %UNFM,1,SEED,%SYSINF(40)

REAL %RN,%SPMLT,%SIRFH,%POINTS(6,2)

REAL %R TIME,RP

REAL %B %CLOCK,%CLRI,%CLR?

.....10.....

.....20.....

.....30.....

.....40.....

.....50.....

.....60.....

.....70.....

THIS PAGE IS BEST QUALITY PRACTICE FROM COPY FURNISHED TO DDC

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

```

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00
INPUT CARD IMAGE
.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00
DATA ITIME/0.00+00/
DATA SFED/1/
DATA POINTS/ 0.0, 1.0, 2.0, 3.0, 4.0, 5.0,
I 0.0, 0.2, 0.5, 0.7, 0.8, 1.0/
C IFITIME .NE. 0.) GO TO 20
DO TO I=1,33
IO SYSCOM(I)=0
20 CALL $SREAL(ITIME,SYSCOM,34)
C ONE TRANSACTION ARRIVES EVERY 5.89 MINUTES +-5 MIN
C (= I TRANS EVERY 353635 MSEC)
I = SUNFRI(I,2,SEED)
C CALCULATE THE DEVIATION (MIN IS -5 MIN, MAX IS +5 MIN)
RN = SUNFH(0.0, 1.0, I)
RP = $PHL(INPOINTS,0,RN)
C CONVERT DEVIATION MINUTES TO MSEC
RP = RP * 60000.
C NEXT ARRIVAL TIME = CURRENT TIME + 5 MIN +- THE DEVIATION
IF(I .EQ. 1) ITIME = 353.6350+03 + RP
IF(I .EQ. 2) ITIME = 353.6350+03 - RP
RETURN
END: PROCEDURE:

```

THIS PAGE IS BEST QUALITY PRACTICES FROM COPY PARALSHED TO DDO

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....00

INFORMATION PROCESSING SYSTEM SIMULATOR
STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE
SEQ-NO LCV SEC-NO KLF-NO

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....	00015130	1513	40
PROCEDURE: NAME=TIME2, TYPE=SUBROUTINE;			
C THIS PROCEDURE ESTABLISHES THE TIME AT WHICH THE PAGE	00015140	1514	44
C CACHE STATISTICS ARE TO BE PRINTED	00015150	1515	50
REAL*8 ITIME	00015160	1516	51
DATA ITIME/7.20000ID*067	00015170	1517	52
CALL \$SREAL(ITIME,SYSCOM,34)	00015180	1518	53
RETURN	00015190	1519	54
END: PROCEDURE;	00015200	1520	55
PROCEDURE: NAME=TIME3, TYPE=SUBROUTINE;	00015210	1521	56
C THIS PROCEDURE ESTABLISHES THE TIME AT WHICH A SNAPSHOT	00015220	1522	57
C OF THE STANDARD IPSS STATISTICS ARE TO BE PRINTED	00015230	1523	58
REAL*8 ITIME	00015240	1524	59
DATA ITIME/7.20*067	00015250	1525	60
CALL \$SREAL(ITIME,SYSCOM,34)	00015260	1526	61
RETURN	00015270	1527	62
END: PROCEDURE;	00015280	1528	63
PROCEDURE: NAME=TIME4, TYPE=SUBROUTINE;	00015290	1529	64
C THIS PROCEDURE ESTABLISHES THE TIME AT WHICH ANOTHER	00015300	1530	65
C SNAPSHOT OF THE STANDARD IPSS STATISTICS ARE TO BE PRINTED	00015310	1531	66
REAL*8 ITIME	00015320	1532	67
DATA ITIME/7.20*067	00015330	1533	68
CALL \$SREAL(ITIME,SYSCOM,34)	00015340	1534	69
RETURN	00015350	1535	70
END: PROCEDURE;	00015360	1536	71
PROCEDURE: NAME=TIME5, TYPE=SUBROUTINE;	00015370	1537	72
C THIS PROCEDURE ESTABLISHES THE TIME AT WHICH ANOTHER	00015380	1538	73
C SNAPSHOT OF THE STANDARD IPSS STATISTICS ARE TO BE PRINTED	00015390	1539	74
REAL*8 ITIME	00015400	1540	75
DATA ITIME/7.20*067	00015410	1541	76
CALL \$SREAL(ITIME,SYSCOM,34)	00015420	1542	77
RETURN	00015430	1543	78
END: PROCEDURE;	00015440	1544	79
PROCEDURE: NAME=TIME6, TYPE=SUBROUTINE;	00015450	1545	80
C THIS PROCEDURE ESTABLISHES THE TIME AT WHICH ANOTHER	00015460	1546	81
C SNAPSHOT OF THE STANDARD IPSS STATISTICS ARE TO BE PRINTED	00015470	1547	82
REAL*8 ITIME	00015480	1548	83
DATA ITIME/7.20*067	00015490	1549	84
CALL \$SREAL(ITIME,SYSCOM,34)	00015500	1550	85
RETURN	00015510	1551	86
END: PROCEDURE;	00015520	1552	87
PROCEDURE: NAME=TIME7, TYPE=SUBROUTINE;	00015530	1553	88
C THIS PROCEDURE ESTABLISHES THE TIME AT WHICH ANOTHER	00015540	1554	89
C SNAPSHOT OF THE STANDARD IPSS STATISTICS ARE TO BE PRINTED	00015550	1555	90
REAL*8 ITIME	00015560	1556	91
DATA ITIME/7.20*067	00015570	1557	92
CALL \$SREAL(ITIME,SYSCOM,34)	00015580	1558	93
RETURN	00015590	1559	94
END: PROCEDURE;	00015600	1560	95
PROCEDURE: NAME=TIME8, TYPE=SUBROUTINE;	00015610	1561	96
C THIS PROCEDURE ESTABLISHES THE TIME AT WHICH ANOTHER	00015620	1562	97
C SNAPSHOT OF THE STANDARD IPSS STATISTICS ARE TO BE PRINTED	00015630	1563	98
REAL*8 ITIME	00015640	1564	99
DATA ITIME/7.20*067	00015650	1565	100
CALL \$SREAL(ITIME,SYSCOM,34)	00015660	1566	101
RETURN	00015670	1567	102
END: PROCEDURE;	00015680	1568	103
PROCEDURE: NAME=TIME9, TYPE=SUBROUTINE;	00015690	1569	104
C THIS PROCEDURE ESTABLISHES THE TIME AT WHICH ANOTHER	00015700	1570	105
C SNAPSHOT OF THE STANDARD IPSS STATISTICS ARE TO BE PRINTED	00015710	1571	106
REAL*8 ITIME	00015720	1572	107
DATA ITIME/7.20*067	00015730	1573	108
CALL \$SREAL(ITIME,SYSCOM,34)	00015740	1574	109
RETURN	00015750	1575	110
END: PROCEDURE;	00015760	1576	111
PROCEDURE: NAME=TIME10, TYPE=SUBROUTINE;	00015770	1577	112
C THIS PROCEDURE ESTABLISHES THE TIME AT WHICH ANOTHER	00015780	1578	113
C SNAPSHOT OF THE STANDARD IPSS STATISTICS ARE TO BE PRINTED	00015790	1579	114
REAL*8 ITIME	00015800	1580	115
DATA ITIME/7.20*067	00015810	1581	116
CALL \$SREAL(ITIME,SYSCOM,34)	00015820	1582	117
RETURN	00015830	1583	118
END: PROCEDURE;	00015840	1584	119
PROCEDURE: NAME=TIME11, TYPE=SUBROUTINE;	00015850	1585	120
C THIS PROCEDURE ESTABLISHES THE TIME AT WHICH ANOTHER	00015860	1586	121
C SNAPSHOT OF THE STANDARD IPSS STATISTICS ARE TO BE PRINTED	00015870	1587	122
REAL*8 ITIME	00015880	1588	123
DATA ITIME/7.20*067	00015890	1589	124
CALL \$SREAL(ITIME,SYSCOM,34)	00015900	1590	125
RETURN	00015910	1591	126
END: PROCEDURE;	00015920	1592	127
PROCEDURE: NAME=TIME12, TYPE=SUBROUTINE;	00015930	1593	128
C THIS PROCEDURE ESTABLISHES THE TIME AT WHICH ANOTHER	00015940	1594	129
C SNAPSHOT OF THE STANDARD IPSS STATISTICS ARE TO BE PRINTED	00015950	1595	130
REAL*8 ITIME	00015960	1596	131
DATA ITIME/7.20*067	00015970	1597	132
CALL \$SREAL(ITIME,SYSCOM,34)	00015980	1598	133
RETURN	00015990	1599	134
END: PROCEDURE;	00016000	1600	135

THIS PAGE IS BEST QUALITY PRACTICABLE
FROM COPY FORWARDED TO DDC

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....

INFORMATION PROCESSING SYSTEM SIMULATOR

STANDARD INPUT STREAM LISTING

DATE ... 09/10/77

INPUT CARD IMAGE

.....10.....20.....30.....40.....50.....60.....70.....80.....90.....

INPUT ALT INPUT MEMPER
SFC-IP LEV SLU-LO PFF-10

EQUATE: CF=START, EX=EV1;

EQUATE: CF=8UFST, EX=EV2;

EQUATE: CF=SNAPST, EX=EV3;

00015370 1537 2

00015380 1538 3

00015390 1539 4

00015400 1540

SIMULATE: STREAM=START, TRACE=OFF, STOP=(TIME=2800010.);

00015410 1541 6

ERROR=(TERMINATE=10, STATISTICS=NO);

00015420 1542 7

00015430 1543

00015440 1544 9

END: MODEL;

THIS PAGE IS BEST QUALITY PRACTICES FROM ANY MATERIAL TO DO

APPENDIX C

RESULTS OF THE RESEARCH ACTIVITIES

This appendix contains the conclusions drawn from the research activities. These conclusions are based in a large part from the experiences gained with the IPSS SIDPERS/IDMS modeling activities. Recall that the purpose of the modeling activities was to demonstrate the feasibility and suitability of the Information Processing System Simulator (IPSS) with respect to modeling large complex information systems. The system modeled to demonstrate these capabilities was a host/backend processor configuration which supported interactive application processing and a data base management system. The modeled software processes included a detailed characterization of application loading, DBMS processing, and the operating system functions of task management, job scheduling and resource allocation.

The research proposal identified 164 different extensions to the then existing IPSS. As a result of this research, 139 new language and statistics gathering features have been incorporated into the IPSS. Other extensions not in the original proposal were also identified as useful features and are being added to the IPSS. The results are reported in the following three sections of this appendix:

<u>Section</u>	<u>Area Covered</u>
C1	Data Base Systems - specific extensions
C2	Information System - related performance statistics
C3	Computer Networks - specific extensions

The vehicle used for assessing the usefulness of the suggested IPSS extensions was to model the complex systems architecture described in previously in this report. For evaluative purposes, the proposed extensions to the IPSS are each assigned one of the following assessment dispositions:

<u>Assessment Conclusion</u>	<u>Category/Disposition</u>
Not Examined	NE
Examined and Rejected	ER
Examined and found to be useful	
- implemented as proposed	EU-I
- implemented but merged into another facility.	EU-IM
- defined and not implemented	EU-D
- defined and not implemented but used in paper model	EU-DP

Dispositions D and DP are intended to show the depth to which an IPSS language statement under consideration has been examined. Paper models demonstrate effectively the definitional capability of IPSS to characterize information system activities of interest, and as such are a firm indication of the desire of implementing the statement into IPSS.

C1. IPSS LANGUAGE STATEMENTS FOR CHARACTERIZING DATA BASE SYSTEMS

The IPSS system was designed to analyze information processing systems in their totality. To date, IPSS mechanisms for the definition and evaluation of hardware, computer system software and storage level data structures have been extensively developed and tested. The next phase in the continued development of IPSS was the incorporation of features for the description, manipulation and evaluation of logical data structures of the kind normally encountered in modern data base systems.

The goal for this aspect of the research was to identify the usefulness of incorporating a set of generalized data base-related characterization facilities into the IPSS. Data Base Systems (DBS) characterization is provided through two components: one definitional in nature which is specifically designed to represent logical data structures, and the other procedural to represent application and DBS

software processing. These components have been named the Data Base Structure component and the Data Base Access component, respectively. The relation of these two components to the existing model architecture is shown in Figure C1-1. For purposes of identification they have been titled the IPSS/DBS submodel architecture. The proposed language for characterizing DBS facilities for the IPSS/DBS are identified in Table C1-1. Also shown are the results of their evaluation as part of this research effort.

C1.1 The Data Base Structure Component

The purpose of the Data Structure component is to provide the modeler with a set of facilities which allows one to define logical data structures and to characterize the relationships among them. This component is defined independently of the data base access component (See Section C1.2) and thus these facilities can be applied to a variety of DBS architectures and application environments.

The Data Base Structure component facilities serve the following functions:

1. They define templates for schemas consisting of record types and sets, instances of which can be transient entity definitions in the DBA component;
2. They allow data base structures to be viewed as a collection of contiguous logical areas and to map these onto non-contiguous domains of a DBS's storage structure and also onto other logical data structures;
3. They specify allocation policies of record type records to the sets in which they participate; and
4. They allow the characterization of access paths between record types of a schema.

The Data Base Structure component permits the modeler to investigate the effects on system behavior caused by alternate set, record type, and access path definitions. The definitional facilities provided allow the modeler to investigate a wide spectrum of logical data structure organizations and allocation policies.

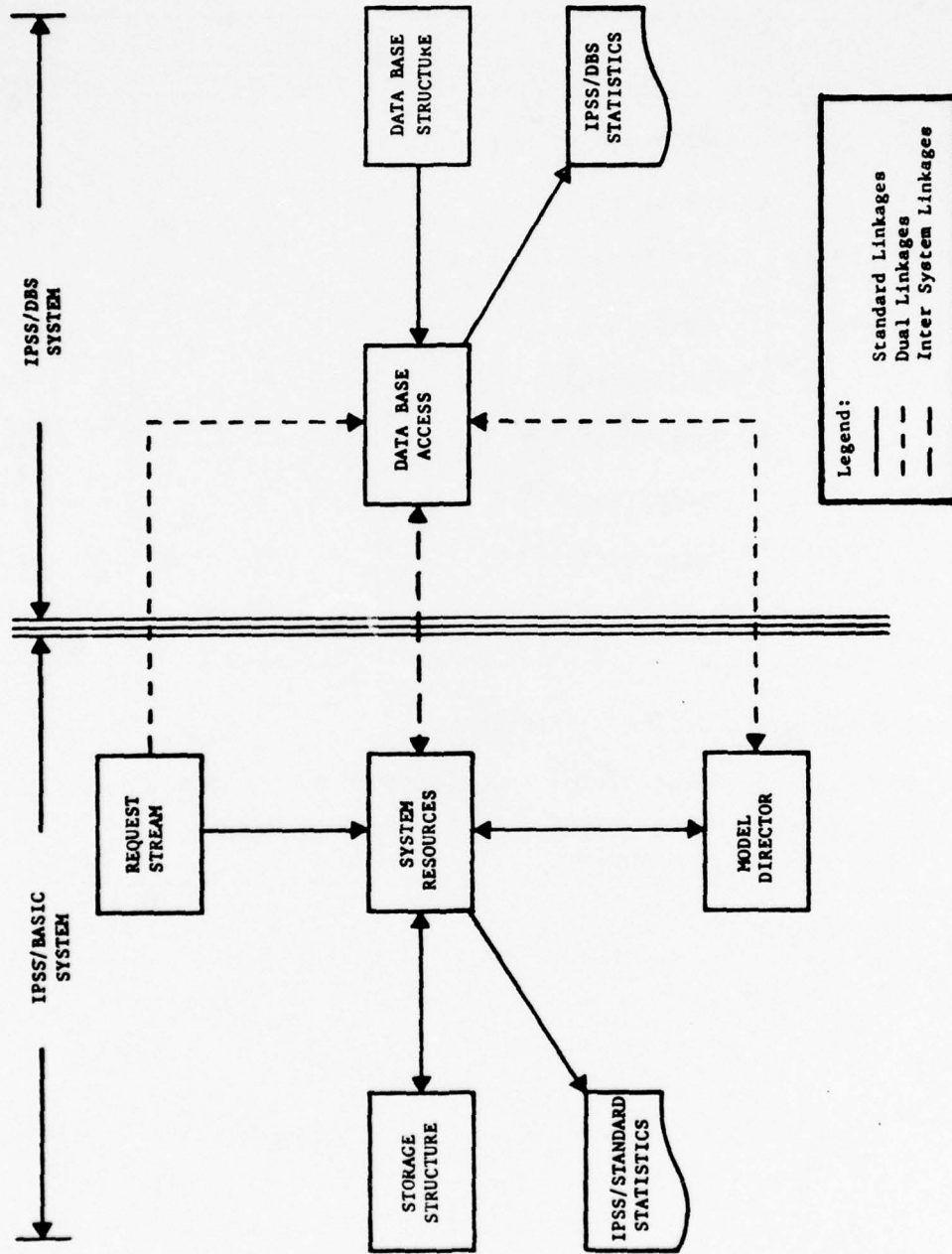


Figure C1-1 IPSS System Architecture with DBMS Extensions

Table C1-1. IPSS/DBS Facilities

Statement	Type	Status in IPSS		Disposition*
		Current	Additional	
BUFFER POOL	Definitional	X		
DML SERVICE	Executorial		X	EU-I
ENDOGENOUS SERVICE	Executorial	X		
EXOGENOUS SERVICE	Executorial	X		
EXTENT	Definitional		X	EU-IM
REALM	Definitional		X	EU-I
RECORD TYPE	Definitional		X	EU-I
SCHEMA	Definitional		X	EU-I
SET TYPE	Definitional		X	EU-I

*Disposition codes are defined as follows:

EU-I : Examined, found Useful - Implemented as proposed;

EU-IM: Examined, found Useful - Implemented by Merging
into another facility.

Schema Definition

The DEFINE SCHEMA - END SCHEMA statement sequence defines a logical data structure of inter-related elements. Within this sequence are the definitions of sets, record types and extents.

Set Definition

The SET TYPE statement characterizes connections between record types in a schema. This statement identifies a collection of record types, names the collection, specifies the type of linkage between the record types, identifies the type of participation (either owner or member) of the record types in the set, and characterizes the relative frequency of set occurrences.

Record Type Definition

The RECORD TYPE statement characterizes those properties of a record type which are independent of secondary storage. The physical attributes are expressed in terms of the record format, the number of record occurrences, and record origin. An association with an extent is also provided for linearization. These physical characteristics are independent of the record location in secondary storage. (The IPSS Storage Structure component performs the required mappings from logical records to storage structure representation.) This Data Base Structure component facility completes the record type definition of the Data Base Access component whose major purpose is to define the source of the records.

Extent Definition

The EXTENT statement defines one or more extent facilities. The extent provides the interface between a record type's logical address space (defined via the RECORD TYPE statement) and the Realm specification defined via the REALM statement. Any number of extents may be defined.

Realm Definition

The REALM facility is used to linearize the address space of a schema. The REALM statement provides the external identifier to the REALM facility and provides the space management area associated with the REALM. Record types may be assigned to realms through the EXTENT statement. As many realms as desired may be defined for each schema definition. IPSS views a realm to be the equivalent of a virtual data set.

C1.2 The Data Base Access Component

The Data Base Access component is the central component in an IPSS/DBS defined model. Within it are contained the definitions of all the required DBS and application level software. Hardware resources are limited to main memory buffers used by the DBS software and user work areas. All DBS-related entity type facilities are defined within this component. This subsection describes those facility definition statements which can appear within a Data Base Access component definition. These represent extensions to the basic IPSS language.

In addition to the facilities identified below, the modeler can include any number of standard IPSS procedure facilities and Fortran Subprogram definitions. These definitions can appear anywhere in a Data Base Access component definition with the exception that they cannot be imbedded within any other facility definition. This implies that a Procedure or Subprogram facility definition cannot appear within:

1. Another Procedure,
2. Another Subprogram,
3. An Endogenous Service, or
4. An Exogenous Service.

They may, however, be referenced from within any of the above.

Procedure and Subprogram syntax and semantics are discussed in Volume I of the IPSS Syntax & Semantics document (Section 3.6).

DML Service Definition

The DML SERVICE - END DML SERVICE statement sequence corresponds in function to the standard IPSS Endogenous Service definition. It has been named DML service in the DBA component to more accurately reflect DBS service request processing.

Record Type Definition

The RECORD TYPE statement defines a record type facility within the Data Base Access Component and identifies the source of the records which are created for the specified record type facility.

Schema Definition

The SCHEMA statement identifies a schema facility to the Data Base Access component. The named Schema is associated with a schema definition in the Data Structure component through EQUATE statements in the Model Stream Component.

C1.3 IPSS/DBS Oriented Built-In Facilities

IPSS/DBS Built-in Procedure Facilities provide the modeler with a large number of predefined capabilities whose purpose is to reduce the time needed to develop data base management system models. These procedural facilities free the modeler from the detailed mechanics of the simulation process. The built-in procedures have been grouped into three categories based on features or functions common to each group they are:

1. I/O-oriented procedures,
2. Schema-oriented procedures, and
3. Data structure traversal-oriented procedures.

Table C1-2 contains an alphabetized listing of all the IPSS/DBS Built-in procedures and denotes the dispositions of the statements. All Built-in procedure statements are converted to Fortran subroutine calls with the parameters associated with that statement transposed to subprogram parameters. Formal descriptions for these IPSS facilities

Table C1-2 IPSS/DBS Built-in Procedures

Statement	Type	Status in IPSS		Disposition
		Current	Additional	
ACQUIRE RECORD	Executorial		X	EU-DP
ALLOCATE SCHEMA EXTENT	Executorial		X	EU-IM
ALTER SET	Executorial		X	NE
CONFIGURE CEQ	Executorial	X		
COPY SET	Executorial		X	NE
COPY ROUTE	Executorial		X	EU-DP
CREATE ROUTE	Executorial		X	EU-I
CREATE SCHEMA	Executorial		X	EU-I
CREATE SCHEMA EXTENT	Executorial		X	EU-IM
CREATE TRANSACTION	Executorial	X		
DESTROY ROUTE	Executorial		X	EU-I
DESTROY SCHEMA	Executorial		X	EU-I
DESTROY SCHEMA EXTENT	Executorial		X	EU-IM
DISABLE EVENT	Executorial	X		
DO TO LABEL	Executorial	X		
ENABLE EVENT	Executorial	X		
END DO	Executorial	X		
FIND IN QUEUE	Executorial	X		
FIND MEMBER	Executorial		X	EU-I
FIND NEXT BUFFER	Executorial	X		
FIND ROUTE	Executorial		X	EU-DP
FREE BUFFER	Executorial	X		
GET BUFFER	Executorial	X		
GET SAVE AREAS	Executorial	X		
IDENTIFY OWNER	Executorial		X	EU-DP
IDENTIFY SET OCCURRENCE	Executorial		X	EU-DP
INITIATE TASK	Executorial	X		
MERGE QUES	Executorial	X		
MODIFY BUFFER POOL	Executorial	X		
MODIFY QUEUE	Executorial	X		
MODIFY REALM	Executorial		X	EU-DP
MODIFY ROUTE	Executorial		X	EU-DP
MODIFY SCHEMA EXTENT	Executorial		X	EU-IM
MODIFY SCHEMA RECORD TYPE	Executorial		X	EU-DP
MODIFY SCHEMA SET	Executorial		X	EU-DP
PLACE IN QUEUE	Executorial	X		
POST SEMAPHORE	Executorial	X		
POST SIGNAL	Executorial	X		
PROCESS TIME	Executorial	X		
PUT SAVE AREAS	Executorial	X		
RELEASE SCHEMA EXTENT	Executorial		X	EU-IM
REQUEST DML SERVICE	Executorial		X	NE
REQUEST SERVICE	Executorial	X		
RESUME SCAN	Executorial	X		
SET FACILITY STATUS	Executorial	X		
SET INDIRECT REFERENCE	Executorial	X		

Table C1-2 IPSS/DBS Built-in Procedures (Continued)

Statement	Type	Status in IPSS		Disposition
		Current	Additional	
SET PRIORITY	Executorial	X		
SET SYSTEM PARAMETER	Executorial	X		
START QUEUE STATISTICS	Executorial	X		
START USAGE STATISTICS	Executorial	X		
STOP QUEUE STATISTICS	Executorial	X		
STOP SIMULATION	Executorial	X		
STOP USAGE STATISTICS	Executorial	X		
STORE RECORD	Executorial		X	NE
TERMINATE SERVICE	Executorial	X		
TERMINATE TASK	Executorial	X		
WAIT DML SERVICE COMPLETE	Executorial		X	NE
WAIT FACILITY STATUS	Executorial	X		
WAIT RECORD	Executorial		X	NE
WAIT SEMAPHORE	Executorial	X		
WAIT SERVICE COMPLETE	Executorial	X		
WAIT SIGNAL	Executorial	X		

* EU-I : Examined, found Useful - Implemented as proposed
 EU-IM: Examined, found Useful - Implemented by Merging into
 another facility
 EU-DP: Examined, found Useful - defined and not implemented but
 used in paper model
 NE : Not examined.

are not presented in this report. IPSS Syntax and Semantics, Vols. I and II presents the following items of information for each Built-in Procedure:

1. Statement syntax,
2. The function of the statement,
3. Semantics for the statement's parameters,
4. A narrative description of the execution mechanism,
5. IPSS generated Fortran source code which replaces IPSS statements and
6. Examples of the use of the statement.

C1.3.1 IPSS/DBS I/O Related Built-in Procedures

The following Built-in Procedures are used to initiate I/O activity and to wait for the I/O to be completed. The ACQUIRE and STORE RECORD procedures determine the source of the record from the DBA component record type facility definition. If a service is specified as the source, then the service is automatically invoked. WAIT RECORD statements are defined to enable service processing to be suspended pending the I/O operation. This section also contains the definition of the statement for initiating and waiting for a DML Service. These procedures can be referenced from within either Procedure facilities or Exogenous and Endogenous Service facilities.

Acquire and Store Record

The ACQUIRE and STORE RECORD Built-in Procedures invoke a service which simulates the acquisition and storage of the specified record to and from the service in which the ACQUIRE or STORE statement appears. The identity of service to be invoked is contained in the definition of the specified record type facility in the DBA component.

Wait Record

The function of the WAIT RECORD Built-in Procedure is to place the specified transaction into the wait state until the completion of the Endogenous service which the specified transaction invoked via an ACQUIRE or STORE RECORD statement. Service for the suspended transaction resumes after the specified wait criteria (count and request number) are satisfied. Each completed Endogenous Service routine satisfying the request number criterion decrements the suspended transaction's Wait Count by one and the delayed transaction is reactivated when the count reaches zero and is returned to the active state.

Request DML Service

The REQUEST DML SERVICE Built-in Procedure is used to initiate the execution of a DML Endogenous Service.

Wait DML Service Complete

The function of the WAIT DML SERVICE COMPLETE Built-in Procedure is to place the named transaction into the wait state until the completion of the specified Endogenous DML Service which it invoked. Service for suspended transaction resumes after the specified wait criteria (count and request number) are satisfied. Each completed Endogenous DML Service routine satisfying the request number criterion decrements the suspended transaction's wait count by one and the delayed transaction is reactivated when the count reaches zero and is returned to the active state.

C1.3.2 IPSS/DBS Schema Related Built-in Procedures

The following Built-in procedures are used to change the current state of facilities defined in the DATA Structure component of an IPSS/DBS model. These procedures can be referenced from within either Procedure facilities or Exogenous and Endogenous Service facilities.

Allocate and Release Schema Extent

The effect of executing the ALLOCATE SCHEMA EXTENT Built-in Procedure is to simulate the allocation of a logical address space to the Extent subfacility. The allocation causes the extent to take on a displacement attribute which specifies an extent's displacement relative to a Realm facility. This is done by placing the calculated displacement of the extent start into the Schema facility's extent table. The displacement is a real number whose dimension is in terms of the reference units implied for the Realm. This information is used by the Acquire and Store Record Built-in procedures (when the desired record type has an origin of Realm) to determine the logical address of a record type record. If a Schema Extent is not allocated, it is assumed to have zero capacity relative to its ability to hold schema records. The RELEASE SCHEMA EXTENT Built-in Procedure returns allocated address space.

Create and Destroy Schema

The purpose of the CREATE SCHEMA Built-in Procedure is to create for the named Schema facility a copy of its associated Data Structure component definition. The copy is placed in a pre-specified area of the Transaction Area. Once copied, simulation functions involving the Schema facility can occur. The definition remains valid until the issuing of a DESTROY SCHEMA statement for the same Schema facility.

Create and Modify Schema Extent

The function of the CREATE AND MODIFY SCHEMA EXTENT Built-in Procedure is to alter the current attributes associated with a Schema Extent facility. The attributes are first assigned via the extent definition in the Schema facility definition of the Data Structure component. A Schema Extent once created can only be modified or destroyed and can only be "created" when it is not in a "created" state. The procedure can only be invoked after a Schema has been created.

Destroy Schema Extent

The DESTROY SCHEMA EXTENT Built-in Procedure deletes the specified schema extent definition from the named Schema facility. The destroy procedure releases the logical address space associated with the extent by setting its size attribute to zero and removing its Realm facility association.

Modify Realm

The MODIFY REALM Built-in Procedure enables the modeler to change the attributes associated with a Realm facility definition. The procedure changes only those attributes explicitly identified in the Modify Realm Statement.

Modify Schema Record Type

The MODIFY SCHEMA RECORD TYPE Built-in Procedure changes the parameter values of the specified record type facility definition. The original parameters are specified in the RECORD TYPE statements in both the Data Structure and DBS components. The procedure changes only those attributes explicitly identified in the MODIFY RECORD TYPE statement.

Modify Schema Set

The function of the MODIFY SCHEMA SET Built-in Procedure is to modify parameters associated with a Schema's Set subfacility. The original parameters are specified via the Data Structure component's Set facility definition. The procedure changes only those attributes explicitly identified in the MODIFY SCHEMA SET Statement.

C1.3.3 IPSS/DBS Data Structure Traversal Oriented Built-in Procedures

The following Built-in Procedures are used a) to create routes or individual sets based upon the schema definitions in the Data Structure component, b) to modify routes or sets by specifying the elements to be changed, and c) to destroy an existing route. The procedure statements are

divided into two groups, those which operate on sets, and those which operate on routes. These procedures can be referenced from within either Procedure facilities or Endogenous and Exogenous Service facilities.

Copy Set

The COPY SET Built-in Procedure creates a set occurrence from the specified set definition. The set occurrence is placed into an output array which the modeler designates. The set may either be copied from the corresponding set definition in the Data Structure component or from an existing route. In either case the modeler may change the set traversal currency indicator during the copy process.

Alter Set

The ALTER SET Built-in Procedure modifies a previously created set occurrence. The specified set occurrence may either be the output area of a previously executed Copy Set procedure or may exist within a route. This statement can be used to change the owner and member record types of the set as well as the record occurrences for these record types.

Find Member

The FIND MEMBER Built-in Procedure identifies a member record occurrence of a set relative to the specified set traversal currency. The output of this procedure is an updated member record occurrence and/or member record type of the set's traversal currency. A NO MEMBER EXIT parameter is provided for those cases in which it is not possible to update this currency; for example, FIND NEXT MEMBER when the number of occurrences associated with the member record type is zero. The FIND MEMBER examines the set and the record type facility definitions in the Data Structure component in order to determine the existence of the desired record. The statement types are as follows: a) FIND FIRST MEMBER specifies that the first member record occurrence for the designated set type is to be located and its identity placed in the set's traversal currency. By "identity" we mean the member record type, the

record occurrence within the set and the record occurrence within the record type. The success of this operation depends on a data structure which directly links the set traversal currency before execution of this procedure to the first record in the set; b) FIND PRIOR MEMBER specifies that the record to be identified is the record occurrence which immediately precedes the one specified by the set traversal currency. The success of this option likewise depends upon a direct link between the record in the set traversal currency and the prior record occurrence; c) FIND NEXT MEMBER specifies that the record to be identified is the record occurrence which immediately follows the record occurrence of the set traversal currency. The successful completion of this procedure depends upon the existence of a direct link between the record specified by the set traversal currency and the next record occurrence in the set; d) FIND LAST MEMBER specifies that the last record occurrence for the specified set type is to be located and its identity placed in the set's traversal currency. The success of this option depends upon a data structure which directly links this record to the set traversal currency before execution of this procedure.

Identify Owner

The IDENTIFY OWNER Built-in Procedure identifies the record type which is the owner of the specified record type and identifies the set in which they both participate. This procedure examines the set and record type facility definitions in the Data Structure component of the current schema to determine this relationship. All the sets which own a given record type can be obtained by specifying IDENTIFY OWNER followed by successive IDENTIFY NEXT OWNER statements.

Identify Set Occurrences

The IDENTIFY SET OCCURRENCE Built-in Procedure identifies the sets which the current member record occurrence owns. By current member record occurrence we mean the traversal currency as reflected in the

specified set. This procedure examines the set and record type facility definitions in the Data Structure component of the current schema to determine ownership of a non-empty set of records. All of the sets owned by a record can be obtained by specifying IDENTIFY SET OCCURRENCE followed by successive IDENTIFY NEXT SET OCCURRENCE statements.

Create Route

The CREATE ROUTE Built-in Procedure determines the existence of a path between two record types of a schema. The beginning and ending nodes of the path are specified in terms of record type identifiers. The function of this procedure is to determine if these two record types are connected and if the connections meet all the set constraints. The set declarations of the identified schema are examined in the order of their declaration until either a path is found or all the declarations for the schema have been explained. All the paths between two record types can be located by specifying PATH = FIRST followed by successive PATH = NEXT statements. When a path is found, its sets are placed in the \$TRANS array beginning with the location returned as the value of the route identifier. If no path is found, the value of the route identifier will be a zero.

Destroy Route

The DESTROY ROUTE Built-in Procedure deletes the specified route from the \$TRANS array. This allows the re-use of the space occupied by the route for other modeling activities. The execution of this procedure does not affect schema definitions in the Data Structure component.

Modify Route

The MODIFY ROUTE Built-in Procedure changes the route definition by either: a) reestablishing the current set, b) deleting or inserting sets, or c) concatenating another route to the existing one. A route consists of a sequence of set traversal tables. This procedure can change this

sequence by deleting, inserting or concatenating set traversal tables relative to any of the sets in the route. Only one type of change can be specified for each Modify Route Statement.

Copy Route

The COPY ROUTE Built-in Procedure creates a route and copies the specified route data into it. The sets in the output route can be modified during this process. These modifications consist of: a) re-establishing the current set, b) inserting or deleting sets, or c) concatenating another route to the output route. These modifications can be specified either by naming specific sets or by referencing the sets in the route relative to the set currency.

Find Route

The FIND ROUTE Built-in Procedure uses a previously created route and the current schema definition to establish specific record occurrences within the sets of the route. A set other than the first set of the route may be designated as the beginning set for the procedure to establish record occurrences. A set other than the last set in the route may also be specified.

C2. IPSS PERFORMANCE MEASURES

The proposal identified 120 extensions to the currently available model statistics. The purpose of this section of Appendix C is to present the results of the research in the area of added statistical outputs from IPSS models.

The Information Processing System Simulator provides a modeler with a number of statistics concerning the behavior of modeler defined entities and IPSS supplied built-in services. These statistics fall into eight general categories listed below. The evaluative results of the research is found in Subsections C2.2 through C2.8 which correspond to these statistics categories:

1. Operational Statistics,
2. Request Stream Statistics,
3. I/O Activity,
4. Queueing Statistics,
5. Utilization Statistics,
6. Wait Statistics,
7. Service Statistics, and
8. Task/Activity Statistics.

In addition to these "automatic" statistics, the modeler can employ the complete facilities of the Fortran language to develop his own statistics. Statistics are printed automatically at the conclusion of each model simulation unless explicitly inhibited.

Time constraints prohibited the implementation of any new statistics. However, definitions for a number of the proposed statistics were completed and are reported in the appropriate subsection. These statistics employ both IPSS facility specific variables and those global to an entire model. For each defined statistic the variables required to support their collection and calculation are also discussed along with the method for data collection including the specific IPSS statements responsible for the activity. Note: all variable names used in these definitions are for publication purposes only.

C2.1 IPSS Operational Statistics

These statistics provide an overview of the operational characteristics of the current simulation execution. The Clock and Termination Statistics involve the length of the run, in terms of both simulation time and CPU time, and the reason for termination. The Transaction Statistics provide insight into the degree of transaction activity generated by the current simulation run. Statistics on the number of transactions generated and the average life span of a transaction are collected by transaction type. The results of their evaluation are shown in Table C2.1-1. Table C2.1-2 identifies the required data to meet the definitional requirements for these statistical extensions to the IPSS. The explicit statistic definitions are shown in Table C2.1-3.

Table C2.1-1 IPSS Operational Statistics

Statistic Type	Specific Statistic	Currently Available	Proposed Extension	Disposition*
Clock	Absolute elapsed simulation time	X		
	Elapsed simulation time since last reset	X		
	Elapsed simulation time of current model execution	X		
	Elapsed CPU time (IPSS run time)		X	EU-DP
	Scaling factor		X	EU-DP
Termination	Number of errors	X		
	Termination status		X	EU-DP
Transaction (By Transaction type)	Total number generated		X	EU-DP
	Current number in system	X		
	Maximum number in system	X		
	Starting number in system	X		
	Total number terminated		X	EU-DP
	Mean time of existence		X	EU-DP

*Distribution codes used in this table are defined as follows:

DU-DP: Examined, found useful - Defined, not implemented but used in paper models

Table C2.1-2 IPSS Operational Statistics - Required Model Statements

Variable Name	Description of the Variable	Method of data Collection
VARIABLES GLOBAL TO AN IPSS MODEL		
CLOCK1	Absolute elapsed simulation time	Automatically by simulation driver
CLOCK2	Elapsed simulation time since last reset	Automatically by simulation driver
CLOCK3	Elapsed simulation time of current model execution	Automatically by simulation driver
CPUTIM	Beginning CPU time of current IPSS run	SVC call by simulation driver
SCALEF	Scaling factor	Specified by user
ERRM	Total number of errors in current model execution	Updated by internal IPSS error routine
TERMST	Termination status	Specified by user
VARIABLES WHICH ARE TRANSACTION SPECIFIC		
TNGEN _A	Total number generated	Updated by simulation driver upon generation of new transaction
TNCUR _A	Current number in system	Maintained by driver in SYSINF
TNMAX _A	Maximum number in system	Either user specified or system default
TNSTRT _A	Starting number in system	Either user specified or system default
TNTRM _A	Total number terminated	Updated by driver upon termination

Table C2.1-3 IPSS Operational Statistics - Definition

Stastic Type	Specific Statistic Name	Definition
Clock	Absolute elapsed simulation time	CLOCK1
	Elapsed simulation time since last reset	CLOCK2
	Elapsed simulation time of current model execution	CLOCK3
	*Elapsed CPU time (IPSS run time)	CPUTIM
	*Scaling Factor	SCALEF
Termination	Number of errors	ERRM
	*Termination Status	TERMST
Transaction (by trans- action type)	*Total number generated	TNGEN _A
	Current number in system	TNCUR _A
	Maximum number in system	TNMAX _A
	Starting number in system	TNSTRT _A
	*Total number terminated	TNTRM _A
	*Mean time of existence	(CLOCK3/TNGEN _A)
*Statistic currently exists in IPSS models.		

C2.2 IPSS Request Stream Statistics

These statistics describe the behavioral characteristics of the Request Stream facilities. For each Exogenous Event a profile of the interarrival time mechanism is presented. The event of the IPSS built-in Endogenous Services, ENABLE/DISABLE Event, on the arrival of events is also summarized. The collection of these statistics is an automatic by-product of the simulation drive mechanism. Table C2.2-1 presents the evaluation results. Table C2.2-2 identifies the required model data and Table C2.2-3 the actual statistic definition.

Table C2.2-1 IPSS Request Stream Statistics

Statistic Type	Specific Statistic	Currently Available	Proposed Extension	Disposition*
General	Number generated	X		EU-DP
	Number terminated		X	
Arrival Time	First generated arrival time		X	EU-DP
	Average interarrival time		X	EU-DP
	Last arrival time generated		X	EU-DP
Enable/Disable	Number of enables		X	EU-DP
	Total time enabled		X	EU-DP
	Average enable time		X	EU-DP
	Number of disables		X	EU-DP
	Total Time Disabled		X	EU-DP
	Average disable time		X	EU-DP

*Disposition Codes used in this table are defined as follows:

EU-DP: Examined, found useful - Defined, not implemented but useful in paper models

Table C2.2-2
IPSS Request Stream Statistics - Required Model Statements

Variable Name	Description of the Variable	Method of Data Collection
EXGEN	Number generated	Updated by driver
EXTRM	Number terminated	Updated by driver
EFARR	First generated arrival time	Specified by user (either directly or via Procedure)
ELARR	Last (latest) arrival time generated	Updated by driver
ENEN	Number of enables	Updated by ENABLE statement
ETLEN	Time of last enable/disable	Updated by both the ENABLE and DISABLE statements
ETTEN	Total time enabled	Calculated by DISABLE statement as $(\text{CLOCK3} - \text{ETLEN}) + \text{ETTEN}$
ENDSA	Number of disables	Updated by DISABLE statement
ETTDSA	Total time disables	Calculated by ENABLE statement as $(\text{CLOCK3} - \text{ETLEN}) + \text{ETTDSA}$

Table C.2.2-3
IPSS Request Stream Statistics - Definitions

Statistic Type	Specific Statistic Name	Definition
General	Number Exogenous requests generated	EXGEN
	*Number Exogenous requests terminated	EXTRM
Arrival time	*First generated Exogenous event arrival time	EFARR
	*Mean interarrival time	(ELARR - EFARR)/EXGEN
	*Last arrival time	ELARR
Enable/Disable	*Number of enables	ENEN
	*Total time enabled	ETTEN
	*Mean enable time	(ETTEN/EXEN)
	*Number of disables	ENDSA
	*Total time disabled	ETTDSA
	*Mean disable time	(ETTDSA/ENDSA)
*Statistic currently exists in IPSS models		

C.2.3 IPSS I/O Activity Statistics

All of the proposed new statistics in the area of I/O activity were concerned with the concept of locality. Statistics were proposed at the Data Set level and within the Data Set, by File, Extent, Volume, Cylinder and Track, and at the Area level, and within Area by Segment and by Volume. All of these statistics are to be gathered by the GET ADDRESS statement. Major modifications however, will be required to the existing IPSS code supporting the IPSS GET ADDRESS Built-in Procedure in order to gather the required data. The purpose of these statistics is to measure the degree of hardware and file management level I/O activity that occurred during the current model. Measures of I/O routine performance are collected by I/O routine type, by associated access mechanism, and by referenced Data Set and File. Statistics showing the degree of locality of consecutive logical and physical references by Data Set and Area were also proposed.

A second set of statistics of similar intent was proposed for the IPSS/DBS subsystem. Since these statistics are closely associated with other I/O activities, the decision was made to delay the examination of this class of statistics until the IPSS/DBS subsystem was operational. However, the I/O-related statistics are defined in Table C2.3 for report completeness.

Table C.2:3 IPSS I/O Activity Statistics

Statistics Type	Specific Statistic	Currently Available	Proposed Extension	Disposition
I/O Invocations (By I/O Statement)	Number of calls	X		
	Percent zero-time calls	X		
	Mean time per call	X		
	Mean time for non-zero calls	X		
	Standard deviation of time per call	X		
	Standard deviation of time per call (excluding zero-time calls)	X		
	(Mean and standard deviation of various physical attributes, i.e., cylinders crossed, characters read, etc.)	X		
I/O Invocations (By Access Mechanism)	By I/O type for each Access Mechanism:			
	Number of calls	X		
	Number of zero-time calls	X		
	Percent zero-time calls	X		
	Mean routine time	X		
	Mean routine time (excluding zero-time calls)	X		
	Standard deviation of mean routine time	X		
	Standard deviation of mean routine time (excluding zero time calls)	X		
I/O Invocations (By Data Set and File)	By I/O type for each Data Set and File:			
	Number of calls	X		
	Number of zero-time calls	X		
	Percent of zero-time calls	X		
	Mean routine time	X		
	Mean routine time (excluding zero-time calls)	X		
	Standard deviation of mean routine time	X		
	Standard deviation of mean routine time (excluding zero-time calls)	X		

Table C.2.3 IPSS I/O Activity Statistics (Continued)

Statistics Type	Specific Statistic	Currently Available	Proposed Extension	Disposition	
Locality (Data Set) (Conditional)	By Data Set: Number of consecutive references for each file		X	NE	
	By File: Number of consecutive references for each Extent		X	NE	
	By Extent: Number of consecutive references to the same Volume		X	NE	
	By Volume: Number of consecutive references to the same cylinder		X	NE	
	Average number of cylinders crossed		X	NE	
	By Cylinder: Number of consecutive references to the same track		X	NE	
	Average number of tracks crossed		X	NF	
	By Track: Number of consecutive references to the same block		X	NE	
	Average number of blocks crossed		X	NE	
	Locality (Area) (Conditional)	By Area: Total number of references to area		X	NE
		Number of consecutive references to same area		X	NE
		Average number of consecutive references		X	NE
		Average length between consecutive references in terms of reference units		X	NE
		By Segment: Total number of references to segment		X	NF

Table C.2.3 IPSS I/O Activity Statistics (Continued)

<u>Statistics Type</u>	<u>Specific Statistic</u>	<u>Currently Available</u>	<u>Proposed Extension</u>	<u>Disposition</u>
	Number of consecutive references to same segment		X	NE
	Average number of consecutive references		X	NE
	Average length between consecutive references in terms of reference units		X	NE
	By Volume:			
	Number of consecutive references to same volume		X	NE
	Average number of consecutive references		X	NE
	Average length between consecutive references in terms		X	NE

*Disposition codes used in this table are:

NE: Not evaluated

C.2.4 IPSS Queueing Statistics

An extensive number of queueing statistics are currently gathered during an IPSS model's execution. Several additional extensions to these statistics were proposed and evaluated. These are now being incorporated into the IPSS statistics generation procedures. Table C.2.4-1 identifies both the existing and proposed statistics and the result of the evaluative effort. The purpose of Table C.2.4-2 is to identify the specific data for these statistics. Table C.2.4-3 provides definitions for the statistics.

Table C.2.4-1 IPSS Queueing Statistics

<u>Statistics Type</u>	<u>Specific Statistic</u>	<u>Currently Available</u>	<u>Proposed Extension</u>	<u>Disposition</u>	
Busy Period	Total number of busy periods	X			
	Number of zero time busy periods	X			
	Percent zero time busy periods	X			
	Total busy time	X			
	Percent busy	X			
	Mean length of busy period	X			
	Mean length of busy period (excluding zero time periods)		X	EU-DP	
	Standard deviation for non-zero time busy periods		X	EU-DP	
	Maximum length of a busy period		X	EU-DP	
	Minimum (non-zero) length of a busy period		X	EU-DP	
	Length of the current busy period		X	EU-DP	
	Idle Period	Total number of idle periods	X		
		Number of zero-time idle periods	X		
		Percent zero-time idle periods	X		
Total idle time		X			
Percent idle		X			
Mean length of an idle period		X			
Mean length of an idle period (excluding zero-time periods)			X	EU-DP	
Standard deviation for non-zero time idle periods			X	EU-DP	
Maximum length of an idle period			X	EU-DP	
Minimum (non-zero) length of an idle period			X	EU-DP	

Table C.2.4-1 IPSS Queueing Statistics (Continued)

Statistics Type	Specific Statistic	Currently Available	Proposed Extension	Disposition
Queue Length (During a Busy Period)	Mean length	X		
	Maximum length	X		
	Current length	X		
	Standard deviation of mean length	X		
Queue	Total number of queue entries	X		
	Total number of completed queue entries	X		
	Number of zero-time queue entries	X		
	Percent of zero-time queue entries	X		
	Mean queue transit time	X		
	Mean queue transit (excluding zero- time transit)	X		
	Standard deviation of non-zero time queue transit	X		
	Maximum transit time		X	EU-DP
	Minimum (non-zero) transit time		X	EU-DP

Disposition codes used for this Table are:

EU-DP: Examined, found useful - Defined not implemented but useful in paper model

Table C.2.4-2 IPSS Queueing Statistics - Required Model Statements

<u>Variable Name*</u>	<u>Description of the Variable</u>	<u>Method of Data Collection</u>
QCNCE	Current number of concurrent elements in queue	Update by Queue and Depart Queue Statement
QMXNCE	Maximum number of current concurrent elements in queue	Update by Queue statement Queue
QSNCE	Sum of number of concurrent elements in queue	Updated by Depart Queue statement
QSNCES	Sum of (number of concurrent elements in queue)**2	Updated by Queue and Depart Queue statement
QSQE	Total number of queue elements	Update by Queue statement
QSZQE	Total number of zero-time queue elements	Updated by Depart Queue statements
QSTT	Sum of time per transaction	Updated by Depart Queue statement
QSTTS	Sum of (time per trans)**2	Updated by Depart Queue statement
QCTI	Cumulative Time integral	Updated by Queue/Depart Queue statement
QTQDQ	Time of last occurrence of Queue or Depart Queue invocation	Updated by Queue and Depart Queue statement
QSBP	Total number of busy periods	Updated by Queue statement
QSBPL	Sum of length of busy periods	Updated by Depart Queue statement
QSBPLS	Sum of (length of busy periods)**2	Updated by Depart Queue statement
QSIP	Number of idle periods	Updated by Depart Queue statement
QSIPL	Sum of length of time for each idle period	Update by Queue statement

Table C.2.4-2 IPSS Queueing Statistics - Required Model Statements (Continued)

Variable Name*	Description of the Variable	Method of Data Collection
QSIPLS	Sum of (length of idle periods)**2	Updated by Queue statement
QTBIP	Time of last Busy or Idle period start	Updated by Queue statement and Depart Queue
QSZBP	Number of zerotime busy periods	Updated by Queue statement
QSZIP	Number of zerotime idle periods	Updated by Depart Queue statement
QMNBP	Minimum non-zero busy period length	Updated by Depart Queue statement
QMXBP	Maximum busy period length	Updated by Depart Queue statement
QCBP	Length of current busy period	Updated by Depart Queue statement
QMNIP	Minimum non-zero idle period length	Updated by Queue statement
QMXIP	Maximum idle period length	Updated by Queue statement
QCIPL	Length of current idle period	Updated by Queue statement
QMXTT	Maximum transit time per transaction	Updated by Depart Queue statement
QMNTT	Minimum non-zero transit time per transaction	Updated by Depart Queue statement

*All data are stored in the statistics area associated with the designated facility.

Table C.2.4-3 IPSS Queueing Statistics - Definitions

Statistic Type	Specific Statistic	Definition
Busy Period	Total number of busy periods	QSBP
	Number of zerotime busy periods	QSZBP
	%zero time busy periods	$(QSZBP/QSBP)*100$
	Total Busy time	QSBPL
	%busy	$(QSBPL/CLOCK3)*100$
	Mean length of busy period	$(QSBPL/QSBP)$
	Standard deviation of mean length	$((QSBP*QSBPLS-QSBPL**2)/(QSBP*(QSBP-1)))**0.5$
	Mean length of nonzero time busy periods	$(QSBPL/(QSBP-QSZBP))$
	Standard deviation of nonzero time busy periods	$((((QSBP-QSZBP)*QSBPLS-QSBPL**2)/((QSBP-QSZBP)*(QSBP-QSZBP-1))))**0.5$
	Maximum length of a busy period	QMXBPL
	Minimum non-zero length	QMNBP
	Length of current busy period	QCBPL
	Idle Period	Total number of idle periods
Number of zerotime idle periods		QSZIP
%zero time idle periods		$(QSZIP/QSIP)*100$
Total idle time		QSIPL
%idle		$(QSIPL/CLOCK3)*100$
Mean length of idle period		$(QSIPL/QSIP)$
Standard deviation of mean length		$((((QSIP-QSIPLS)-QSIPL**2)/(QSIP*(QSIP-1)))**0.5$

Table C.2.4-3 IPSS Queueing Statistics - Definitions (Continued)

Statistic Type	Specific Statistic	Definition
	Mean length of nonzero time idle periods	$(QSIPL / (QSIP - QSZIP))$
	Standard deviation of nonzero time idle periods	$((((QSIP - QSZIP) * QSIPLS - QSIPL ** 2) / ((QSIP - QSZIP) * (QSIP - QSZIP - 1))) ** 0.5)$
	Maximum length of an idle period	QMXIPL
	Minimum nonzero length	QMNIPL
	Length of current idle period	QCIPL
Queue Length	Mean length	$(QSNCE / QSQE)$
	Maximum length	QMXNCE
	Current length	QCNCCE
	Standard deviation of mean length	$((QSQE * QSNCE - QSNCE ** 2) / (QSQE * (QSQE - 1))) ** 0.5$
Queue	Total number of queue entries	QSQE
	Total number of completed queue entries	$QSQE - QCNCCE = QSCQE$
	Number of zerotime queue entries	QSZQE
	%zerotime queue entries	$(QSZQE / QSQE) * 100$
Transit Time	Mean transit time per transaction	$(QSTT / (QSQE - QCNCCE))$
	Mean nonzero time per transaction	$(QSTT / (QSCQE - QSZQE))$
	Standard deviation of mean transaction time	$((QSCQE * QSTTS - QSTT ** 2) / (QSCQE * (QSCQE - 1))) ** 0.5$

Table C.2.4-3 IPSS Queueing Statistics - Definitions (Continued)

<u>Statistic Type</u>	<u>Specific Statistic</u>	<u>Definition</u>
	Standard deviation of mean non-zero time transactions	$\frac{(((QSCQE-QSZQE)*QSTTS-QSTT**2)}{((QSCQE-QSZQE)*(QSCQE-QSZQE-1))**0.5}$
	Maximum transit time	QMXTT
	Minimum nonzero transit time	QMNTT

C.2.5 IPSS Utilization Statistics

An IPSS model generates utilization statistics on user defined modeled facilities. Several extensions to these statistics were proposed and evaluated in the research. Table C.2.5-1 identifies all IPSS utilization statistics, whether current or proposed and their desposition after evaluation. Table C.2.5-2 identifies the means for their generation and Table C.2.5-3 contains their definitions.

Table C.2.5-1 IPSS Utilization Statistics

<u>Statistic Type</u>	<u>Specific Statistics</u>	<u>Currently Available</u>	<u>Proposed Extension</u>	<u>Disposition</u>	
Busy Period	Total number of busy periods	X			
	Number of zero time busy periods	X			
	Percent zero time busy periods	X			
	Total busy time	X			
	Percent busy	X			
	Mean length of busy period	X			
	Mean length of busy period (excluding zero time periods)		X	EU-DP	
	Standard deviation for nonzero time busy periods		X	EU-DP	
	Maximum length of a busy period		X	EU-DP	
	Minimum (non-zero) length of a busy period		X	EU-DP	
	Length of the current busy period		X	EU-DP	
	Idle Period	Total number of idle periods	X		
		Number of zero time idle periods	X		
Percent zero-time idle periods		X			
Total idle time		X			
Percent idle		X			
Mean length of an idle period		X			
Mean length of an idle period (excluding zero-time periods)			X	EU-DP	
Standard deviation for nonzero time idle periods			X	EU-DP	
Maximum length of an idle period			X	EU-DP	
Minimum (non-zero) length of an idle period			X	EU-DP	
Length of the current idle period			X	EU-DP	
Concurrency		Mean level of concurrency	X		
		Maximum level of concurrency	X		
	Current level of concurrency	X			
	Standard deviation of mean level of concurrency	X			

Table C.2.5-1 IPSS Utilization Statistics (Continued)

<u>Statistic Type</u>	<u>Specific Statistics</u>	<u>Currently Available</u>	<u>Proposed Extension</u>	<u>Disposition</u>
Seizure	Total number of seizures	X		
	Total number of completed seizures	X		
	Number of zero-time seizures	X		
	Percent zero-time seizures	X		
	Mean seizure time	X		
	Mean seizure time (excluding zero-time seizures)	X		
	Standard deviation of non-zero time seizures	X		
	Maximum length of a seizure		X	EU-DP
	Minimum (non-zero) length of a seizure		X	EU-DP

Disposition codes used in this table are:

EU-DP: Examined, found useful - defined, not implemented but useful in paper model.

Table C.2.5-2 IPSS Utilization Statistics - Required Model Statement

<u>Variable Name*</u>	<u>Description of the Variable</u>	<u>Method of Data Collection</u>
UCNCS	Current number of Concurrent seizures	Updated by Seize/Release statements
UMXNCS	Maximum number of current concurrent seizures	Updated by Seize statement
USNCS	Sum of number of concurrent seizures	Updated by Seize/Release statement
USNCSS	Sum of (number of concurrent seizures)**2	Updated by Seize/Release statement
USSEZ	Total number of seizures	Updated by Seize statement
USZSEZ	Total number of zero-time seizures	Updated by Release statement
USTT	Sum of time per transaction	Updated by Release statement
USTTS	Sum of (time per trans)**2	Updated by Release statement
UCTI	Cumulative Time Integral	Update by Seize and Release statements
UTSR	Time of last seize or release	Updated by Seize and Release statements
USBP	Total number of busy periods	Update by Seize statement
USBPL	Sum of length of busy periods	Update by Release statement
USBPLS	Sum of (length of busy periods)**2	Updated by Release statement
USIP	Number of idle periods	Updated by Release statement
USIPL	Sum of length of idle periods	Updated by Seize statement
USIPLS	Sum of (length of idle periods)**2	Updated by Seize statement
UTBIP	Time of last Busy or idle period start	Updated by Seize and Release statement

Table C.2.5-2 IPSS Utilization Statistics - Required Model Statement (Continued)

<u>Variable Name*</u>	<u>Description of the Variable</u>	<u>Method of Data Collection</u>
USZBP	Number of zerotime busy periods	Updated by Release statement
USZIP	Number of zerotime idle periods	Updated by Seize statement
UMNBPL	Minimum non-zero busy period length	Updated by Release statement
UMXBPL	Maximum busy period length	Updated by Release statement
UCBPL	Length of current busy period	Updated by Release statement
UMNIPL	Minimum non-zero idle period length	Updated by Seize statement
UMXIPL	Maximum idle period length	Updated by Seize statement
UCIPL	Length of current idle period	Updated by Seize statement
UMXTT	Maximum transit time per transaction	Updated by Release statement
UMNTT	Minimum non-zero transit time per transaction	Updated by release statement

*All data are stored in the statistics area associated with the designated facility.

Table C.2.5-3 IPSS Utilization Statistics - Definitions

<u>Statistic Type</u>	<u>Specific Statistic</u>	<u>Definition</u>
Busy Period	Total number of busy periods	USBP
	Number of zero time busy periods	USZBP
	%zero time busy periods	$(USZBP/USBP)*100$
	Total busy time	USBPL
	%busy	$(USBPL/CLOCK3)*100$
	Mean length of busy period	$(USBPL/USBP)$
	Standard deviation of mean length	$((USBP-USBPLS)-USBPL**2) / (USBP*(USBP-1))**0.5$
	Mean length of non-zero time busy periods	$(USBPL/(USBP-USZBP))$
	Standard deviation of non-zero time busy periods	$((USBP-USZBP)*USBPLS-USBPL**2) / ((USBP-USZBP)*(USBP-USZBP-1))**0.5$
	Maximum length of a busy period	UMXBPL
	Minimum non-zero length	UMNBPL
	Length of current busy period	UCBPL
	Idle Period	Total number of idle periods
Number of zero time idle periods		USZIP
%zero time idle periods		$(USZIP/USIP)*100$
Total idle time		USIPL
%idle		$(USIPL/CLOCK3)*100$
Mean length of idle period		$(USIPL/USIP)$
Standard deviation of mean length		$((USIP*USIPLS-USIPL**2) / (USIP*(USIP-1))**0.5$

Table C.2.5-3 IPSS Utilization Statistics - Definitions (Continued)

<u>Statistic Type</u>	<u>Specific Statistic</u>	<u>Definition</u>
	Mean length of non-zero time idle periods	$(USIPL / (USIP - USZIP))$
	Standard deviation of non-zero time idle periods	$(((USIP - USZIP) * USIPLS - USIPL ** 2) / ((USIP - USZIP) * (USIP - USZIP - 1))) ** 0.5$
	Maximum length of an idle period	UMXIPL
	Minimum non-zero length	UMNIPL
	Length of current Idle period	UCIPL
Concurrency	Mean level	$(USNCS / USSEZ)$
	Maximum level	UMXNCS
	Current level	UCNCS
	Standard deviation of mean level	$((USSEZ * USNCS - USNCS ** 2) / (USSEZ * (USSEZ - 1))) ** 0.5$
Seizures	Total number of seizures	USSEZ
	Total number of completed seizures	USCSEZ = USSEZ - UCNCS
	Number of zero-time seizures	USZSEZ
	%zero-time seizures	$(USZSEZ / USSEZ) * 100$
Transit Time	Mean transit time per transaction	$(USTT) / USCSEZ$
	Mean non-zero time per transaction	$(USTT) / (USCSEZ - USZSEZ)$
	Standard deviation of mean transaction time	$((USCSEZ * USTTS - USTT ** 2) / (USCSEZ * (USCSEZ - 1))) ** 0.5$
	Standard deviation of mean non-zero-time transactions	$(((USCSEZ - USZSEZ) * USTTS - USTT ** 2) / ((USCSEZ - USZSEZ) * (USCSEZ - USZSEZ - 1))) ** 0.5$

Table C.2.5-3 IPSS Utilization Statistics - Definitions (Continued)

<u>Statistic Type</u>	<u>Specific Statistic</u>	<u>Definition</u>
	Maximum transit time	UMXTT
	Minimum non-zero transit time	UMNTT

C.2.6 IPSS Wait Statistics

These statistics measure the degree to which individual transactions were delayed until specific system resources attained appropriate status. They are similar to types of statistics gathered for Queueing Statistics. IPSS automatically collects these statistics as a by-product of the following IPSS Built-in procedures:

1. Wait Facility Status,
2. Wait Service Complete,
3. Wait I/O Complete,
4. Wait Semaphore,
5. Wait Access Path, and
6. Wait Signal.

The effect of these built-in procedures is to suspend the execution of a service until the identified condition is met. The statistics for each are the same. Only the WAIT FACILITY STATUS definitions are described in this section. Table C2.6-1 identifies all the IPSS Wait Statistics, Table C2.6-2 the means for their operation and Table C2.6-3 their definitions.

Table C.2.6-1 IPSS Wait Statistics

Statistic Type	Specific Statistics	Currently Available	Proposed Extension	Disposition	
Busy Period	Total number of busy periods	X			
	Number of zero-time busy periods	X			
	Percent zero-time busy periods	X			
	Total busy time	X			
	Percent busy	X			
	Mean length of a busy period	X			
	Mean length of a busy period (excluding zero-time periods)		X	EU-DP	
	Standard deviation for non-zero time busy periods		X	EU-DP	
	Maximum length of a busy period		X	EU-DP	
	Minimum (non-zero) length of a busy period		X	EU-DP	
	Length of the current busy period		X	EU-DP	
	Idle Period	Total number of idle periods	X		
		Number of zero-time idle periods	X		
Percent zero-time idle periods		X			
Total idle time		X			
Percent idle		X			
Mean length of an idle period		X			
Mean length of an idle period (excluding zero-time periods)			X	EU-DP	
Standard deviation for non-zero time idle periods			X	EU-DP	
Maximum length of an idle period			X	EU-DP	
Minimum (non-zero) length of an idle period			X	EU-DP	
Length of the current idle period			X	EU-DP	
Wait Chain Length Characteristics		Mean length	X		
		Maximum length	X		
	Current length	X			
	Standard deviation of mean length	X			

Table C.2.6-1 IPSS Wait Statistics (Continued)

<u>Statistic Type</u>	<u>Specific Statistics</u>	<u>Currently Available</u>	<u>Proposed Extension</u>	<u>Disposition</u>
Wait Chain Entry Characteristics	Total number of entries	X		
	Total number of completed entries	X		
	Number of zero-time entries	X		
	Percent zero-time entries	X		
	Mean time in chain	X		
	Mean time in chain (excluding zero-time entries)	X		
	Standard deviation of mean time non-zero-time entries	X		
	Maximum time in chain		X	EU-DP
	Minimum (non-zero) time in chain		X	EU-DP

Disposition codes used in this table are:

EU-DP: Examined, found useful - defined, not implemented but useful in paper model.

Table C.2.6-2 IPSS Wait Statistics - Required Model Statements

<u>Variable Type*</u>	<u>Description of the Variable</u>	<u>Method of Data Collection</u>
WCNE	Current number of Concurrent wait chain elements	
WMXNCE	Maximum number of current concurrent wait chain elements	
WSNCE	Sum of number of concurrent wait chain elements	
WSNCES	Sum of (number of concurrent wait chain elements)**2	
WSWCE	Total number of wait chain elements	
WSZWCE	Total number of zero-time wait chain elements	
WSTT	Sum of time per transaction	All statistics are maintained by the Wait statements
WSTTS	Sum of (time per trans)**2	
WCTI	Cumulative Time integral	
WTWFSS	Time of last wait start or wait end	
WSBP	Total number of busy periods	
WSBPL	Sum of length of busy periods	
WSBPLS	Sum of (length of busy periods)**2	
WSIP	Number of idle periods	
WSIPL	Sum of length of idle periods	
WSIPLS	Sum of (length of idle periods)**2	
WTBIP	Time of last Busy or Idle period start	

Table C.2.6-2 IPSS Wait Statistics - Required Model Statements (Continued)

<u>Variable Type*</u>	<u>Description of the Variable</u>	<u>Method of Data Collection</u>
WSZBP	Number of zerotime busy periods	
WSZDP	Number of zerotime idle periods	
WMNBPL	Minimum non-zero busy period length	
WMXBPL	Maximum busy period length	
WCBPL	Length of current busy period	All statistics are maintained by the Wait Statement
WMNIPL	Minimum non-zero idle period length	
WMXIPL	Maximum idle period length	
WCIPL	Length of current idle period	
WMXTT	Maximum transit time per transaction	
WMNTT	Minimum non-zero transit time per transaction	

*All are stored in the statistics area associated with the designated facility.

Table C.2.6-3 IPSS Wait Statistics - Definitions

<u>Statistics Type</u>	<u>Specific Statistic</u>	<u>Definitions</u>
Busy Period	Total number of busy periods	WSBP
	Number of zerotime busy periods	WSZBP
	%zero time busy periods	$(WSZBP/WSBP)*100$
	Total busy time	WSBPL
	%busy	$(WSBPL/CLOCK3)*100$
	Mean length of busy period	$(WSBPL/WSBP)$
	Standard deviation of mean length	$((WSBP*WSBPLS-WSBPL**2)/(WSBP*(WSBP-1)))**0.5$
	Mean length of non-zero time busy periods	$(WSBPL/(WSBP-WSZBP))$
	Standard deviation of non-zero time busy periods	$((((WSBP-WSZBP)*WSBPLS-WSBPL**2)/((WSBP-WSZBP)*(WSBP-WSZBP-1)))**0.5$
	Maximum length of a busy period	WMXBPL
	Minimum non-zero length	WMNBPL
	Length of current busy period	WCBPL
Idle Period	Total number of idle periods	WSIP
	Number of zerotime idle periods	WSZIP
	%Zerotime idle periods	$(WSZIP/WSIP)*100$
	Total idle time	WSIPL
	%Idle	$(WSIPL/CLOCK3)*100$
	Mean length of idle period	$(WSIPL/WSIP)$
	Standard deviation of mean length	$((WSIP*WSIPLS-WSIPL**2)/(WSIP*(WSIP-1)))**0.5$
	Mean length of non-zero time idle periods	$(WSIPL/(WSIP-WSZIP))$

Table C.2.6-3 IPSS Wait Statistics - Definitions (Continued)

Statistics Type	Specific Statistic	Definitions
	Standard deviation of non-zero time idle periods	$((\text{WSIP}-\text{WSZIP}) * \text{WSIPLS}-\text{WSIPL}^{**2}) / ((\text{WSIP}-\text{WSZIP}) * (\text{WSIP}-\text{WSZIP}-1))^{**0.5}$
	Maximum length of an idle period	WMXIPL
	Minimum non-zero length	WXNIPL
	Length of current Idle period	WC IPL
Chain Length	Mean length	(WSNCE/WSWCE)
	Maximum length	WMXNCE
	Current length	WCNCE
	Standard deviation of mean length	$((\text{WSWCE} * \text{WSNCEs}-\text{WSNCE}^{**2}) / (\text{WSWCE} * (\text{WSWCE}-1)))^{**0.5}$
Chain	Total number of wait chain entries	WSWCE
	Total number of completed chain entries	WSCCE = WSWCE - WCNCE
	Number of zero-time entries	WSZWCE
	%zero-time chain entries	(WSZWCE/WSWCE)*100
Transit Time	Mean transit time per transaction	(WSTT/WSCCE)*100
	Mean non-zero time per transaction	(WSTT/(WSCCE-WSZWCE))
	Standard deviation of mean transaction time	$((\text{WSCCE} * \text{WSTTS}-\text{WSTT}^{**2}) / (\text{WSCCE} * (\text{WSCCE}-1)))^{**0.5}$
	Standard deviation of mean non-zero-time transactions	$((\text{WSCCE}-\text{WSZWCE}) * \text{WSTTS}-\text{WSTT}^{**2}) / ((\text{WSCCE}-\text{WSZWCE}) * (\text{WSCCE}-\text{WSZWCE}-1))^{**0.5}$
	Maximum transit time	QMXTT
	Minimum non-zero transit time	WMNTT

C.2.7 IPSS Service Related Statistics

IPSS Exogenous and Endogenous services are modeler-defined model components representing information system processes. These processes could be manual or automated. IPSS provides the modeler with over 100 statements in addition to the use of all of ANSI Fortran to define these services. In the current IPSS, services are treated the same as other acquirable resources for statistic gathering purposes.

In the proposal a number of additional statistics specific to services were proposed. The purpose of these new statistics is to characterize service performance in a more meaningful manner. This is accomplished in the following two ways:

A. Modeler Explicit Statistics

Three new sets of IPSS statements permit the modeler to divide overall service time into three segments;

1. Queue Time - time which a request for service spends in a service maintained queue waiting for its service to commence,
2. Acquire Time - time spent by the service acquiring the resources needed to service a request, and
3. Service Time - time spent by the service servicing the request.

These statistics are generated only when explicitly requested by the modeler through the employment respectively of the following IPSS statement pairs:

1. MARK QUEUE START - MARK QUEUE END
2. MARK ACQUIRE START - MARK ACQUIRE END
3. MARK SERVICE START - MARK SERVICE END

B. Modeler Implicit Statistics

The purpose of these statistics is to allocate a service's time to subordinate-services. The IPSS translator will identify

all Endogenous services (modeler defined or IPSS Built-in) requested by a modeler defined service. Statistics on subservice performance will be gathered and displayed with the calling service. These statistics show total service time for the invoked service with respect to the invocations by the requestor service and total time. The requestor service was delayed wait for the requested service to complete processing. These statistics are automatically generated by the REQUEST SERVICE - WAIT SERVICE statements.

Table C.2.7-1 identifies the specific statistics generated. Tables C.2.7-2 and C.2.7-3 specify the IPSS mechanisms used in their generation.

Table C.2.7-1 IPSS Service Related Statistics

<u>Statistic Type</u>	<u>Specific Statistics</u>	<u>Currently Available</u>	<u>Proposed Extension</u>	<u>Disposition*</u>
Execution	Total number of executions		X	EU-DP
	Percent zero-time		X	EU-DP
	Mean execution time		X	EU-DP
	Mean non-zero execution time		X	EU-DP
Queue Time	Total number of periods		X	EU-DP
	Percent zero-time		X	EU-DP
	Mean period length		X	EU-DP
	Mean non-zero period length		X	EU-DP
	Mean percent of execution time		X	EU-DP
	Mean non-zero percent of execution time		X	EU-DP
Acquire Time	Total number of periods		X	EU-DP
	Percent zero-time		X	EU-DP
	Mean period length		X	EU-DP
	Mean non-zero period length		X	EU-DP
	Mean percent of execution time		X	EU-DP
	Mean non-zero percent of execution time		X	EU-DP
Service Time	Total number of periods		X	EU-DP
	Percent zero-time		X	EU-DP
	Mean period length		X	EU-DP
	Mean non-zero period length		X	EU-DP
	Mean percent of execution time		X	EU-DP
	Mean non-zero percent of execution time		X	EU-DP
Service by Service (For each Service invoked by a parent service)	Total number of invocations		X	EU-DP
	Percent zero-time invocations		X	EU-DP
	Mean time per invocation		X	EU-DP
	Mean time per invocation (excluding zero-time invocations)		X	EU-DP
	Total number of Wait Service Completes		X	EU-DP

Table C.2.7-1 IPSS Service Related Statistics (Continued)

<u>Statistic Type</u>	<u>Specific Statistics</u>	<u>Currently Available</u>	<u>Proposed Extension</u>	<u>Disposition*</u>
	Percent zero-time Waits		X	EU-DP
	Mean time per Wait		X	EU-DP
	Mean time per Wait (excluding zero-time Waits)		X	EU-DP

*Disposition codes used in this table are:

EU-DP: Examined, found useful - defined, but not implemented

Table C.2.7-2 IPSS Service Related Statistics - Required Model Statistics

Variable Type	Description of the Variable	Method of Data Collection
<p>The following variables identify storage locations in the fixed statistics area associated with each modeler defined Exogenous and Endogenous Service</p>		
SNT	Total number of executions	Updated by REQUEST SERVICE Statement
SNZT	Total number of zero-time executions	Updated by WAIT RETURN* statement
SCUMT	Cumulative execution time	Update by TERMINATE* statement
SNQ	Number of queue periods	Updated by MARK QUEUE START statement
SNZQ	Number of zero-time queue periods	Updated by MARK QUEUE END, MARK ACQUIRE START, MARK SERVICE START or TERMINATE statements
SCUMQ	Cumulative queue period time	
SNAC	Number of acquire periods	Updated by MARK ACQUIRE START statement
SNZAC	Number of zero time acquire periods	Updated by MARK ACQUIRE END, MARK SERVICE START or TERMINATE statements
SCUMAC	Cumulative acquire period time	
SNS	Number of service periods	Updated by MARK SERVICE START statement
SNZS	Number of zero time periods	Updated by MARK SERVICE END or TERMINATE statements
SCUMS	Cumulative service period time	

The following variables identify storage locations in the variable part of the statistics area. These are replaced for each subordinate service requested by a service. (Note: intermediate statistics are kept in the transactions created for a requested service until the transaction is terminated.)

SSNI _A	Total number of invocations of Service A	Updated by REQUEST SERVICE statement
-------------------	--	--------------------------------------

Table C.2.7-2 IPSS Service Related Statistics - Required Model Statistics (Continued)

<u>Variable Type</u>	<u>Description of the Variable</u>	<u>Method of Data Collection</u>
SSNZI _A	Total number of zero-time invocations of Service A	Update by TERMINATE statement
SSTMI _A	Total execution time	Update by TERMINATE statement
SSNW _A	Total number of Wait Service	Update by WAIT SERVICE Complete statement
SSNZW _A	Total number of zerotime	Updated by WAIT RETURN statement
SSTWS _A	Total Wait service time	Updated by WAIT RETURN statement

*These statements are internal to IPSS and not available to the modeler.

Table C.2.7-3 IPSS Service Related Statistics - Definitions

<u>Statistics Type</u>	<u>Specific Statistic</u>	<u>Definitions</u>
Execution	Total number of executions	SNT
	Percent zero-time	$(SNZT/SNT)*100$
	Mean execution time	$SAVT = (SCUMT/SNT)$
	Mean non-zero time execution time	$SAVTX = (SCUMT/(SNT - SNZT))$
Queue Time	Total number of queue periods	SNQ
	Percent zero-time	$(SNZQ/SNQ)*100$
	Mean period length	$SAVQ = (SCUMQ/SNQ)$
	Mean non-zero time period length	$SAVZQ = (SCUMQ/(SNQ - SNZQ))$
	Mean percent of execution time	$SAVQ/SAVT$
	Mean non-zero % of non-zero execution time	$SAVQZ/SAVTX$
Acquire Time	Total number	SNAC
	Percent zero-time	$(SNZAC/SNAC)*100$
	Mean length	$SAVAC = (SCNMAC/SNAC)$
	Mean non-zero length	$SAVACX = (SCUMAC/(SNAC - SNZAC))$
	Mean % of execution time	$(SAVAC/SAVT)*100$
	Mean non-zero % of non-zero execution time	$(SAVACX/SAVTX)*100$
Service Time	Total number	SNS
	Percent zero-time	$(SNZS/SNS)*100$
	Mean length	$SAVS = (SCUMS/SNS)$

Table C.2.7-3 IPSS Service Related Statistics - Definitions (Continued)

Statistics Type	Specific Statistic	Definitions
	Mean non-zero length	$SAVSX = (SCUMS / (SMS - SNZS))$
	Mean % execution time	$(SAVS / SAVT) * 100$
	Mean non-zero % of non-zero execution time	$(SAVSX / SAVTX) * 100$
Subservice Behavior for Requestor Service	Total number of invocations	$SSNI_A$
	Percent zero-time	$(SSNZI_A / SSNI_A) * 100$
	Mean time per invocation	$(SSTMI_A / SSNI_A)$
	Mean time per non-zero time execution	$(SSTMI_A / (SSNI_A - SSNZI_A))$
	Total number of WAIT SERVICE COMPLETES	$SSNW_A$
	Percent zero time Waits	$(SSNZW_A / SSNW_A) * 100$
	Mean time per Wait	$(SSTW_A / SSNW_A)$
	Mean time per non-zero time Wait	$(SSTW_A / (SSNW_A - SSNZW_A))$

C.2.8 IPSS Task and Activity Statistics

These statistics provide insight into the service activity subordinate to a specific task, where a Task is a modeler defined facility activated as part of the execution of a Exogenous and Endogenous service. A Task is assumed to be comprised of a linear combination of one or more modeler defined activities which are identified as part of the Task declaration. A modeler defines an activity by associating an activity name with an endogenous service. Statistics are gathered for a Task by gathering statistics for its activities. This occurs in the following two step process:

1. The Task is explicitly activated via the INITIATE TASK statement.
2. Statistics are automatically gathered for each Endogenous service which
 - a. is subordinate to the inservice containing the INITIATE TASK statement, and
 - b. has been identified as an activity associated with the Task.

The statistics gathering continues until the TERMINATE TASK statement is executed.

Table C.2.8-1 identifies the statistics to be gathered. However, the actual collection mechanism to be used within an IPSS model has not yet been defined.

Table C.2.8-1 IPSS Task and Activity Statistics

Statistic Type	Specific Statistics	Currently Available	Proposed Extension	Disposition*	
By Task -- General	Total number of tasks initiated		X	EU-IM	
	Total number of tasks completed		X	EU-IM	
	Total number of tasks still in execution		X	EU-IM	
	Mean time per task initiation		X	EU-IM	
By Activity -- Within Task -- Task Initiation Statistics	Total number of invocations by task		X	EU-IM	
	Mean total activity processing		X	EU-IM	
	Mean total queue time per task initiation		X	EU-DP	
	Percent queue time of processing time		X	EU-DP	
	Mean total activity acquire time per task initiation		X	EU-DP	
	Percent acquire time of processing time		X	EU-DP	
	Mean total activity service time per task initiation		X	EU-DP	
	Percent service time of processing time		X	EU-DP	
	Activity Invoke Statistics	Mean processing time per invocation by task		X	EU-IM
		Mean activity queue time per invocation by task		X	EU-DP
Percent queue time of processing time per invocation			X	EU-DP	
Mean activity acquire time per invocation by task			X	EU-DP	
Percent acquire time of processing time per invocation			X	EU-DP	
Mean activity service time per invocation by task			X	EU-DP	
Percent service time of processing time by task			X	EU-DP	

*Disposition codes used in this table have the following meanings:

EU-IM:

EU-DP:

C3. IPSS EXTENSIONS TO CHARACTERIZE COMPUTER NETWORKS

The standard IPSS and the IPSS/DBS have been designed to operate as independent simulators. However, an integrated model capability is required in order to characterize a complete system. In the IPSS this capability was achieved by permitting the concurrent execution of two asynchronous simulation processes which have the ability to communicate. Once this capability was defined, it was possible to extend the mechanism to accommodate any number of asynchronous processes.

The purpose of this section is to describe the mechanism to be used which will permit the modeling of information processing system networks. For this discussion a mode is assumed to be an individual model containing IPSS components:

<u>Component</u>	<u>Number Permitted</u>
Request Stream	0 or 1
System Resources	0 or 1*
Storage Structure	0 or 1**
Data Base Access	0 or 1*
Data Structure	0 or 1***

*Either the System Resources or Data Base Access Component must be present. Both can also be present at each mode.

**Storage Structure Component can be present only when the System Resources Component is also present.

***Data Structure Component can be present only when the Data Base Access Component is also present.

A network is comprised of multiple nodes which communicate by initiating requests for services to another node. These services requests are treated the same as exogenous requests generated by a node's request stream. To establish a network the modeler defines a multiple facility pool - which is described in the following subsection.

C3.1 Multiple Facility Pools

Currently, in order to simulate a computer network, all nodes of the network have to be combined together to form a single facility pool consisting of at most one of each of the following components: system resource component, storage structure component, request stream component and model director component. Since the above approach does not facilitate independent testing of each node or evaluating the interaction between nodes, a modeling structure which allows multiple facility pools is needed. The basic structure of the IPSS nucleus which directs model execution is given in Figure C3.1-1a. In order to consider multiple facility pools, this structure will have to be modified to that shown in Figure C3.1-1b. In the multiple facility pool configuration, the nucleus consists of three drivers:

1. Model Driver

The function of the model driver is the same as in the case of a single facility pool. It's basic function is to start and terminate the simulation drivers.

2. Global Simulation Driver

The global simulation driver controls the FEQ (Future Events Queue) and the simulation clock. All future events generated by the different facility pools are combined together into a single FEQ so as to provide correct simulated time synchronization between the facility pools. The global simulation driver has a priority list that is used to determine which local simulation driver has control.

3. Local Simulation Driver

Each facility pool has a local simulation driver which controls its corresponding CEQ (Current Event Queue). Once a local simulation driver is given control, it retains control until either (i) a transaction has to be placed in a different CEQ or (ii) a transaction has to be placed in the FEQ. In these cases, the local driver

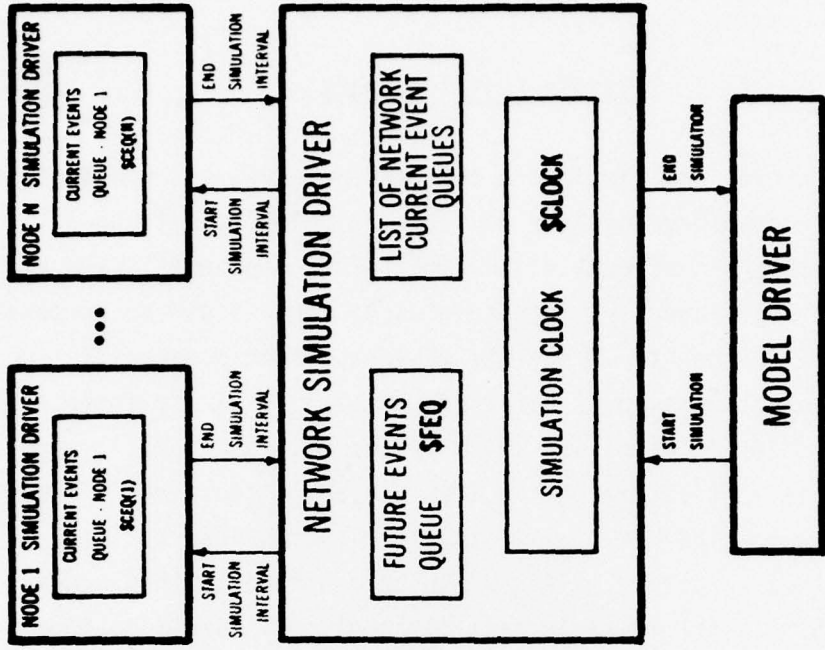


Figure C3.1-1b Structure of Model Nucleus to Support Networking

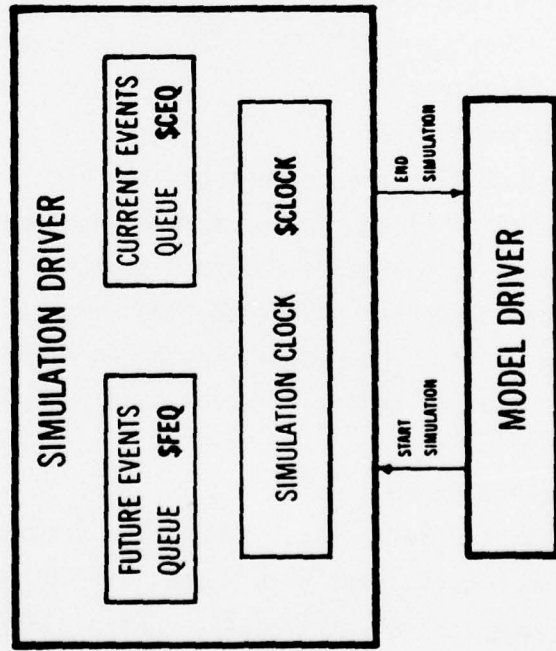


Figure C3.1-1a Basic Structure of IPSS Simulation Nucleus

returns control to the global driver which has the responsibility for placing the transaction into the FEQ or calling the appropriate local nucleus to insert a transaction into its CEQ. If placement is in a CEQ with higher priority than the currently active one, the interrupt flag is set. The currently active nucleus again resumes processing its active transaction until its processing is suspended. If the interrupt flag is set, control is given to the local driver with the higher priority via the global nucleus, otherwise the current driver continues with its next transaction. The process continues until all the CEQ's are empty at which time the clock is advanced to the next most imminent departure time in the FEQ. One or more transactions with this time are then moved to their corresponding CEQ and control is given to the nucleus with highest priority and the process begins a new.

C3.2 Changes to IPSS Simulation Nucleus

In order to accommodate the multiple facility pool structure, a number of changes had to be made to the IPSS nucleus and system subroutines.

1. The IPSS definitional tables (e.g., \$TOC, \$LBL, \$DIR, \$DEF) of the individual facility pools have to be combined together to form a single set of tables.
2. To avoid naming conflicts between facility pools, a global interface routine is needed to invoke the currently implemented local interface routines.
3. The system subroutines (e.g., EQUATE, INVOKE) that handle intra-facility pool communications have to be modified to handle inter-facility pool communications.

Detailed programming changes to IPSS are given below.

Changes to the IPSS Fortran System Subroutines to Handle
Multiple Facility Pools

(A) The following procedures are to be modified:

1. \$FDEF
This procedure has to be modified so that it will accept base pointers to the definitional tables as parameters. A procedure for searching the System Resources and Data Base Access \$DEF tables through \$TOC, \$LBL and \$DIR has to be added to this procedure.
2. \$EQUAT
This procedure has to be modified so that it will accept definitional table addresses and bases for the equated components as input parameters.
3. \$CREAT
This procedure has to be modified so that both the Facility's Pool ID and Driver ID are stored in the created transaction.
4. \$GENER, \$GADR, \$SYSEX2, \$SEEK, \$EXINT
All calls to the interface routine are to be modified.
5. \$LINK, \$UNLINK, \$INSRT, \$REMOV
All these procedures have to be modified to accept multiple CEQ's.
6. \$DRIVR
This procedure has to be modified to handle multiple CEQ's. The driver should automatically interchange base for the definitional table when a new CEQ is selected. All calls to interface routines should be modified. Also, all calls to initialization routines have to be modified so that all Facility Pool/Driver-relative components are initialized.
7. \$DVINT, \$ABORT
Parameter lists are to be deleted from these procedures and appropriate COMMON's should be added.

8. All statistical routines are to be modified so that separate statistics will be printed for each facility pool.
9. All procedures that use the Create Data Set (\$CDS) arrays have to be changed.
10. All procedures that use absolute addresses for accessing the definitional table have to be changed so that they use relative addresses in the base table.

(B) The following procedures to be added.

1. \$BLOC
When the Facility Pool ID and Driver ID are given, this procedure will search through the base table to identify the bases to be used for the local simulation driver.
2. \$SUPEQ
This procedure handles all inter-facility pool equates and makes use of \$EQUAT to resolve all intra-facility pool equates.
3. \$INIT
This procedure initializes all endo and exo services with the facility pool and driver ID's. The endo and exo service definitions are found by searching through \$TOC, \$LBL and \$DIR.
4. \$INVEX
This procedure handles the invoking of an Exo service in another facility pool. Basically, this procedure facilitates inter-facility pool communications.

(C) The following Fortran tables and statements are to be generated by the language parsers.

1. CEQ Table
This table will be used by the flobal simulation driver to determine the priority of the CEQ's.

\$PRTY	\$FPID	\$DRIVR	\$CBEG	\$CEND
:	:	:	:	:
:	:	:	:	:
:	:	:	:	:

\$PRTY: Priority of the CEQ

\$DPIF: Facility Pool ID

\$DRIVR: type of driver

1. if SR, RS or SS component driver
2. if DBA or DT component driver

\$CBEG: beginning addresses of CEQ's (initialized to zero's)

\$CEND: ending addresses of CEQ's (initialized to zero's)

2. BASE Table

This table stores the bases used for addressing the definitional tables.

\$FPID	\$CMPNT	\$STOC	\$LBL	\$DIR	\$DEF	\$CDS	\$NEVNT
:	:	:	:	:	:	:	:
:	:	:	:	:	:	:	:
:	:	:	:	:	:	:	:

\$FPID: Facility Pool ID

\$CMPNT: Component ID

- type = 1 if SR component
- type = 2 if RS component
- type = 3 if SS component
- type = 4 if DBA component
- type = 5 if DT component

\$STOC: list of bases for \$TØC
\$\$ LBL: list of bases for \$LBL
\$\$DIR: list of bases for \$dir
\$\$DEF: list of bases for \$DEF
\$CDS: list of bases for create data set array
\$NEVNT: number of events that has occurred in the associated
facility pool.

3. The following interface routine will be generated for each
(facility pool, driver) pair.

```
SUBROUTINE $xxxxx (TYPE, SYSIND, INDX1, INDX2,*)  
INTEGER TYPE, SYSIND, INDX1, INDX2  
GO TO (100, 200, 300, 400), TYPE  
50 RETURN 1  
100 CALL $CFPR (SYSIND, &50)  
RETURN  
200 CALL $YSER (INDX2, INDX1, &50)  
RETURN  
300 CALL $DBPR (SYSIND, &50)  
RETURN  
400 CALL $EXPR (SYSIND, &50)  
RETURN  
END
```

Where xxxxx is a unique name generated for each (facility
pool, driver) pair.

4. The following global interface routine is generated for the
model pool.

```
SUBROUTINE $SUPER (FPID, DRIV, TYPE, SYSIND, INDX1, XINDX2,*)  
INTEGER FPID, DRIV, TYPE, SYSIND  
INTEGER INDX1, INDX2, I  
GO TO 9999  
1 CALL $xxxxx (TYPE, SYSIND INDX1, INDX2 $100000)  
RETURN  
.  
.  
.
```

```
9999 I = FPID *2 + DRV-2  
      GO TO (1, 2, ... n), I  
10000 RETURN 1  
      END
```

Where n is equal to the number of (facility pool, driver) pairs.