

AD-A182 063

INTEGRATED INFORMATION SUPPORT SYSTEM (IISS) VOLUME 7
COMMUNICATIONS SUBS. (U) GENERAL ELECTRIC CO
SCHENECTADY NY PRODUCTION RESOURCES CONSU..

1/1

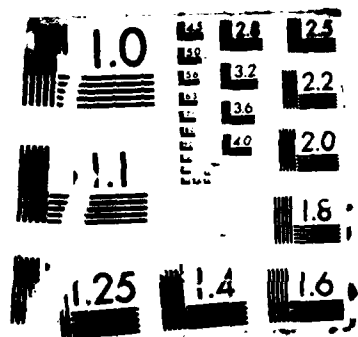
UNCLASSIFIED

K WELDON ET AL. 01 NOV 85 DS-620143000

F/G 12/5

NL

END
100
100
100



2

AD-A182 063

DTIC FILE COPY

**AFWAL-TR-86-4006
Volume VII
Part 1**



**INTEGRATED INFORMATION
SUPPORT SYSTEM (IISS)
Volume VII - Communications Subsystem
Part 1 - COMM Development Specification**

**General Electric Company
Production Resources Consulting
One River Road
Schenectady, New York 12345**

**DTIC
ELECTE
JUL 06 1987
S D D**

**Final Report for Period 22 September 1980 - 31 July 1985
November 1985**

Approved for public release; distribution is unlimited.

PREPARED FOR:

**MATERIALS LABORATORY
AIR FORCE WRIGHT AERONAUTICAL LABORATORIES
AIR FORCE SYSTEMS COMMAND
WRIGHT-PATTERSON AFB, OH 45433-6533**

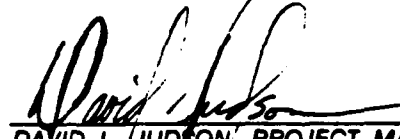
87 7 2 038

NOTICE

When Government drawings, specifications, or other data are used for any purpose other than in connection with a definitely related Government procurement operation, the United States Government thereby incurs no responsibility nor any obligation whatsoever; and the fact that the government may have formulated, furnished, or in any way supplied the said drawings, specifications, or other data, is not to be regarded by implication or otherwise as in any manner licensing the holder or any other person or corporation, or conveying any rights or permission to manufacture, use, or sell any patented invention that may in any way be related thereto.

This report has been reviewed by the Office of Public Affairs (ASD/PA) and is releasable to the National Technical Information Service (NTIS). At NTIS, it will be available to the general public, including foreign nations.

This technical report has been reviewed and is approved for publication.


DAVID L. JUDSON, PROJECT MANAGER
AFWAL/MLTC
WRIGHT PATTERSON AFB OH 45433

5 Aug 1986
DATE

FOR THE COMMANDER:


GERALD C. SHUMAKER, BRANCH CHIEF
AFWAL/MLTC
WRIGHT PATTERSON AFB OH 45433

7 Aug 86
DATE

"If your address has changed, if you wish to be removed from our mailing list, or if the addressee is no longer employed by your organization please notify AFWAL/MLTC, W-PAFB, OH 45433 to help us maintain a current mailing list."

Copies of this report should not be returned unless return is required by security considerations, contractual obligations, or notice on a specific document.

Unclassified

1 November 1985

SECURITY CLASSIFICATION OF THIS PAGE

REPORT DOCUMENTATION PAGE

1. REPORT SECURITY CLASSIFICATION Unclassified		10. RESTRICTIVE MARKINGS	
2a. SECURITY CLASSIFICATION AUTHORITY		3. DISTRIBUTION AVAILABILITY OF REPORT Approved for public release; distribution is unlimited.	
2b. DECLASSIFICATION/DOWNGRADING SCHEDULE			
4. PERFORMING ORGANIZATION REPORT NUMBER(S)		5. MONITORING ORGANIZATION REPORT NUMBER(S) AFVAL-TR-86-40006 Vol VII, Part 1	
6a. NAME OF PERFORMING ORGANIZATION General Electric Company Production Resources Consulting	6b. OFFICE SYMBOL (If applicable) AFVAL/MLTC	7a. NAME OF MONITORING ORGANIZATION AFVAL/MLTC	
6c. ADDRESS (City, State and ZIP Code) 1 River Road Schenectady, NY 12345		7b. ADDRESS (City, State and ZIP Code) WPAFB, OH 45433-6535	
8a. NAME OF FUNDING/SPONSORING ORGANIZATION Materials Laboratory Air Force Systems Command, USAF	8b. OFFICE SYMBOL (If applicable) AFVAL/MLTC	9. PROCUREMENT INSTRUMENT IDENTIFICATION NUMBER F33615-80-C-3155	
8c. ADDRESS (City, State and ZIP Code) Wright-Patterson AFB, Ohio 45433		10. SOURCE OF FUNDING NOS.	
		PROGRAM ELEMENT NO. 78011F	PROJECT NO. 7500
		TASK NO. 62	WORK UNIT NO. 01
11. TITLE (Include Security Classification) (See Reverse)			
12. PERSONAL AUTHOR(S): Veldon, Karen and Phillips, Donna			
13a. TYPE OF REPORT Final Technical Report	13b. TIME COVERED 22 Sept 1980 - 31 July 1985	13c. DATE OF REPORT (Yr., Mo., Day) 1985 November	13d. PAGE COUNT 64
14. SUPPLEMENTARY NOTATION ICAM Project Priority 6201			
The computer software contained herein are theoretical and/or references that in no way reflect Air Force-owned or -developed computer software.			
17. COSATI CODES		18. SUBJECT TERMS (Continue on reverse if necessary and identify by block number)	
FIELD	GROUP	SUB GR	
1308	0905		
19. ABSTRACT (Continue on reverse if necessary and identify by block number) This document defines the Development Specification for the Communications Subsystem (COMM) configuration item of the Integrated Information Support System (IISS) Test Bed. The Test Bed is an integrated hardware/software environment in which various Integrated Computer Aided Manufacturing (ICAM) Program applications will function cooperatively.			
20. DISTRIBUTION AVAILABILITY OF ABSTRACT UNCLASSIFIED/UNLIMITED ☒ SAME AS RPT. ☐ DTIC USERS ☐		21. ABSTRACT SECURITY CLASSIFICATION Unclassified	
22a. NAME OF RESPONSIBLE INDIVIDUAL David L. Jenson		22b. TELEPHONE NUMBER (Include Area Code) 815-255-6976	22c. OFFICE SYMBOL AFVAL/MLTC

DD FORM 1473, 83 APR

EDITION OF 1 JAN 79 IS OBSOLETE.

Unclassified
SECURITY CLASSIFICATION OF THIS PAGE

11. Title

Integrated Information Support System (IISS)
Vol VII - Communications Subsystem
Part 1 - COMM Development Specification

Accession For	
NTIS CRA&I	☒
DTIC TAB	☐
Unannounced	☐
Justification	
By	
Distribution/	
Availability Codes	
Dist	Avail. & d/or Special
A-1	



PREFACE

This development specification covers the work performed under Air Force Contract F33615-80-C-5155 (ICAM Project 6201). This contract is sponsored by the Materials Laboratory, Air Force Systems Command, Wright-Patterson Air Force Base, Ohio. It was administered under the technical direction of Mr. Gerald C. Shumaker, ICAM Program Manager, Manufacturing Technology Division, through Project Manager, Mr. David Judson. The Prime Contractor was Production Resources Consulting of the General Electric Company, Schenectady, New York, under the direction of Mr. Allan Rubenstein. The General Electric Project Manager was Mr. Myron Hurlbut of Industrial Automation Systems Department, Albany, New York.

Certain work aimed at improving Test Bed Technology has been performed by other contracts with Project 6201 performing integrating functions. This work consisted of enhancements to Test Bed software and establishment and operation of Test Bed hardware and communications for developers and other users. Documentation relating to the Test Bed from all of these contractors and projects have been integrated under Project 6201 for publication and treatment as an integrated set of documents. The particular contributors to each document are noted on the Report Documentation Page (DD1473). A listing and description of the entire project documentation system and how they are related is contained in document FTR620100001, Project Overview.

The subcontractors and their contributing activities were as follows:

TASK 4.2

<u>Subcontractors</u>	<u>Role</u>
Boeing Military Aircraft Company (BMAC)	Reviewer
D. Appleton Company (DACOM)	Responsible for IDEF support, state-of-the-art literature search
General Dynamics/ Ft. Worth	Responsible for factory view function and information models

DS 620143000
1 November 1985

Subcontractors

Role

Illinois Institute of
Technology

Responsible for factory view
function research (IITRI)
and information models of
small and medium-size business

North American Rockwell

Reviewer

Northrop Corporation

Responsible for factory view
function and information
models

Pritsker and Associates

Responsible for IDEF2 support

SofTech

Responsible for IDEF0 support

TASKS 4.3 - 4.9 (TEST BED)

Subcontractors

Role

Boeing Military Aircraft
Company (BMAC)

Responsible for consultation on
applications of the technology
and on IBM computer technology.

Computer Technology
Associates (CTA)

Assisted in the areas of
communications systems, system
design and integration
methodology, and design of the
Network Transaction Manager.

Control Data Corporation
(CDC)

Responsible for the Common Data
Model (CDM) implementation and
part of the CDM design (shared
with DACOM).

D. Appleton Company
(DACOM)

Responsible for the overall CDM
Subsystem design integration and
test plan, as well as part of
the design of the CDM (shared
with CDC). DACOM also
developed the Integration
Methodology and did the schema
mappings for the Application
Subsystems.

DS 620143000
1 November 1985

Subcontractors

Role

Digital Equipment
Corporation (DEC)

Consulting and support of the
performance testing and on DEC
software and computer systems
operation.

McDonnell Douglas
Automation Company
(McAuto)

Responsible for the support and
enhancements to the Network
Transaction Manager Subsystem
during 1984/1985 period.

On-Line Software
International (OSI)

Responsible for programming the
Communications Subsystem on the
IBM and for consulting on the
IBM.

Rath and Strong Systems
Products (RSSP) (In 1985
became McCormack & Dodge)

Responsible for assistance in
the implementation and use of
the MRP II package (PIOS) that
they supplied.

SofTech, Inc.

Responsible for the design and
implementation of the Network
Transaction Manager (NTM) in
1981/1984 period.

Software Performance
Engineering (SPE)

Responsible for directing the
work on performance evaluation
and analysis.

Structural Dynamics
Research Corporation
(SDRC)

Responsible for the User
Interface and Virtual Terminal
Interface Subsystems.

Subcontractors and other prime contractors under other
projects who have contributed to Test Bed Technology, their
contributing activities and responsible projects are as follows:

Subcontractors

Role

General Dynamics/
Ft. Worth

Responsible for
factory view

DS 620143000
1 November 1985

<u>Contractors</u>	<u>ICAM Project</u>	<u>Contributing Activities</u>
Boeing Military Aircraft Company (BMAC)	1701, 2201, 2202	Enhancements for IBM node use. Technology Transfer to Integrated Sheet Metal Center (ISMC)
Control Data Corporation (CDC)	1502, 1701	IISS enhancements to Common Data Model Processor (CDMP)
D. Appleton Company (DACOM)	1502	IISS enhancements to Integration Methodology
General Electric	1502	Operation of the Test Bed and communications equipment.
Hughes Aircraft Company (HAC)	1701	Test Bed enhancements
Structural Dynamics Research Corporation (SDRC)	1502, 1701, 1703	IISS enhancements to User Interface/Virtual Terminal Interface (UI/VTI)
Systran	1502	Test Bed enhancements. Operation of Test Bed.

TABLE OF CONTENTS

		<u>Page</u>
SECTION 1.0	SCOPE	1-1
1.1	Identification	1-1
1.2	Functional Summary	1-1
SECTION 2.0	DOCUMENTS	2-1
2.1	Applicable Documents	2-1
2.1.1	Specifications	2-1
2.1.2	Standards	2-1
2.1.3	Other	2-2
SECTION 3.0	REQUIREMENTS	3-1
3.1	Communications Subsystem	3-1
3.1.1	Communication Subsystem Constraints	3-1
3.1.2	Interface Requirements	3-6
3.1.2.1	Interface Block Diagram	3-7
3.1.2.2	Detailed Interface Definition of COMM	3-7
3.1.2.3	Detailed Interface Definition of Primitives	3-7
3.2	Detailed Functional Requirements	3-8
3.2.1	Communication Protocol - Characteristics and Description	3-8
3.2.1.1	Characteristics of the Communication Protocol	3-8
3.2.1.2	Sequence Number Handling	3-11
3.2.1.2.1	Send Sequence Number Checking	3-11
3.2.1.2.2	Receive Sequence Number Checking	3-11
3.2.1.3	Use of Protocol for Binary Data Transmission	3-11
3.2.1.4	Interface Between IBM and non-IBM Computers	3-12
3.2.1.4.1	Non-IBM Computers	3-12
3.2.1.4.2	IBM Computers	3-13
3.2.2	Message Format	3-13
3.2.3	Timing	3-17
3.2.3.1	Line Idle	3-17
3.2.3.2	Master	3-17
3.2.3.3	Slave	3-17

TABLE OF CONTENTS (Continued)

	<u>Page</u>
3.2.3.4 NTM Input Mailbox Full	3-17
3.2.4 Block Check Computation	3-17
3.2.5 Control Characters in Data	3-18
3.2.6 Communication Subsystem Primitives	3-18
3.2.6.1 Initialize Communication Port Interface	3-18
3.2.6.2 Send a Message to a Terminal Port	3-19
3.2.6.3 Receive a Message from a Terminal Port	3-20
3.2.6.4 Get a Message from a Terminal Port	3-21
3.2.6.5 Cancel a Receive Request from a Terminal Port	3-23
3.2.6.6 Terminate Communication Port Interface	3-24
3.2.7 Interprocess Communication Primitives	3-24
3.2.7.1 Create a Mailbox	3-24
3.2.7.2 Send a Message to Another Program	3-27
3.2.7.3 Receive a Message from Another Program	3-28
3.2.7.4 Get a Message from Another Program	3-29
3.2.7.5 Delete a Mailbox	3-29
3.2.7.6 Release an Event Block	3-30
3.2.7.7 Start a Timer	3-31
3.2.7.8 Stop a Timer	3-32
3.2.7.9 Wait for an Event to Occur	3-33
3.2.7.10 Terminate a Program	3-35
3.2.7.11 Save an Event Indicator	3-36
3.2.7.12 Request an Error Be Logged	3-36
3.2.7.13 Log an Error	3-37
3.3 Special Requirements	3-39
3.3.1 Programming Methods	3-39
3.3.2 Expandability	3-39
3.4 Human Performance	3-40
3.5 Data Description of Arguments Used with the Primitives	3-40
3.5.1 Data Descriptions of Arguments for Communication Subsystem Primitives	3-40

TABLE OF CONTENTS (Continued)

	<u>Page</u>
3.5.2 Data Description of Arguments for IPC Primitives	3-41
3.5.3 Data Description for the Event Block	3-45
3.6 Adaptation Requirements	3-46
SECTION 4.0 QUALITY ASSURANCE PROVISIONS	4-1
4.1 Introduction and Definition	4-1
SECTION 5.0 PREPARATION FOR DELIVERY	5-1

LIST OF ILLUSTRATIONS

<u>Figure</u>	<u>Title</u>	<u>Page</u>
3-1	IISS Test Bed Communications Subsystem (Structural Schematics)	3-2
3-2	IISS Hardware Configuration	3-3
3-3	Interprocess Communication Architecture ...	3-5
3-4	COMM Configuration Tree	3-8
3-5	Simple Scenario for Sending a Message Between NTM's	3-10
3-6	Communication Block Diagram	3-16

SECTION 1

SCOPE

1.1 Identification

This specification establishes the performance, development, test, and qualification requirements of the computer program identified as the Communications Subsystem, hereinafter referred to as COMM. COMM is one configuration item of the Integrated Information Support System (IISS).

1.2 Functional Summary

The COMM Computer Program Configuration Item (CPCI) provides a mechanism for transferring messages (data and control) between two tasks. These tasks can be executing on the same computer or on different computers. For the latter, the two tasks are Network Transaction Managers (NTM's).

The major functions of the COMM are:

1. **Interhost Communications:** Interhost Communications is responsible for moving variable length messages without error between computers. This function receives messages from an NTM, passes the messages to a local area network (LAN), receives messages from the LAN, and passes the messages to an NTM.
2. **Interprocess Communications:** Interprocess Communications allows the NTM to receive messages from and send messages to programs on the same computer. The programs may be another portion on the NTM, COMM, UI/VTI, the precompiler, NDDL, or user application programs. This function also allows programs to use timers and to process fatal errors.

SECTION 2

DOCUMENTS

2.1 Applicable Documents

The following documents were used in the definition of the COMM specification.

2.1.1 Specifications

- [1] Control Data Corporation and D. Appleton Co., Inc.; IISS Test Bed CDM Needs Analysis, 7 June 1982; IISS Test Bed CDM Environment, 7 June 1982; IISS Test Bed CDM System Requirements, 7 June 1982.
- [2] General Electric Co., Test Bed System Requirement Document (Draft), Revised 23 August 1982.
- [3] ICAM Computer-Based Information System (CBIS) System Requirements Document (Draft), 10 September 1981, CI #SRD3101400000.
- [4] General Electric Co., Test Bed System Specification (Draft), 23 August 1982.
- [5] General Electric Co., Test Bed System Design Specification, 7 February 1983.

2.1.2 Standards

- [6] American National Standards Committee X3, American National Dictionary for Information Processing, X3/TR-1-77, September 1977.
- [7] ICAM Documentation Standards, 28 December 1981, IDS150120000A.
- [8] SofTech, Inc., ICAM Test Bed Interim Standards and Procedures, 31 May 1982; ISP620150000.
- [9] General Electric Co., IISS Software Development Guidelines/Conventions (Draft), 23 August 1982.

2.1.3 Other

- [10] ICAM Program Office, The Integrated Sheet Metal Center, 30 September 1981.
- [11] ICAM Program Office, The Role of the ICAM Test Bed and Integrated Information Support System (Draft), 18 May 1982.
- [12] SofTech, Inc., IISS Response to CBIS Requirements and 'Threads': SofTech Reactions, 18 March 1982.
- [13] Digital, VAX-11 Architecture Handbook, Digital Equipment Corp., Maynard, MA, 1979.
- [14] IBM, A Guide to the IBM 3031 Processor Complex and Attached Processor Complex of System/370, GC20-18 54-3; System/370 Principles of Operation, GA22-70000.
- [15] Honeywell, Level 6 Minicomputer Systems Handbook.
- [16] Digital, VAX/VMS System Services Reference Manual, AA-D018B-TE; VAX-11 Information Directory and Index, AA-D016D-TE.
- [17] Users' Manual, IDBMS (2.0) Users' Manual, June 1980.
- [18] Systems Users' Manual, IDBMS (2.0) System Users' Manual, June 1980; IBM, OS/VS2 MVS Supervisor Services and Macro Instructions; GC28-0683-2.
- [19] Honeywell, Level 6 GCOS MOD600 System Concepts, CB50-02.

SECTION 3

REQUIREMENTS

This section includes functional and performance requirements for the COMM. In addition, it defines the COMM interfaces to other IISS CPCI's.

3.1 Communications Subsystem

The Communications Subsystem is primarily used to transfer messages (as opposed to files) between Network Transaction Managers on two different host computers. (The transfer of messages between tasks on the same computer is accomplished through Interprocess Communication Primitives.) Two copies of the COMM (the name of a single occurrence of the Communication Subsystem program) reside on each host, and each copy communicates between the NTM and a copy of COMM residing on one of the other hosts (see Figure 3-1). COMM communicates with the NTM via the IISS Interprocess Communication Primitives.

The requirement to implement IISS as a system residing on three computers has caused many system dependent issues to surface. The system design goal is to isolate system dependent coding by creating primitives which hide the system dependent programming from the IISS subsystems. The primitives used by COMM are:

1. Interprocess Communication (IPC)
2. Interhost Communication (IHC)

3.1.1 Communication Subsystem Constraints

The IISS hardware consists of a Local Area Network (LAN) and its interface into each computer (see Figure 3-2). The terminal interface standard is the interface into the LAN. The interface to the IBM is handled through an IBM 3271 station emulator (cluster controller) that manages the asynchronous to synchronous, the EBCDIC to ASCII character set, and the RS232C to 3270 protocol conversions.

The IISS protocol has permanent virtual circuits established between each pair of computers by the LAN when it is started (see Figure 3-1). By placing this requirement on the LAN, the overhead of having the communication software establish

DS 620143000
1 November 1985

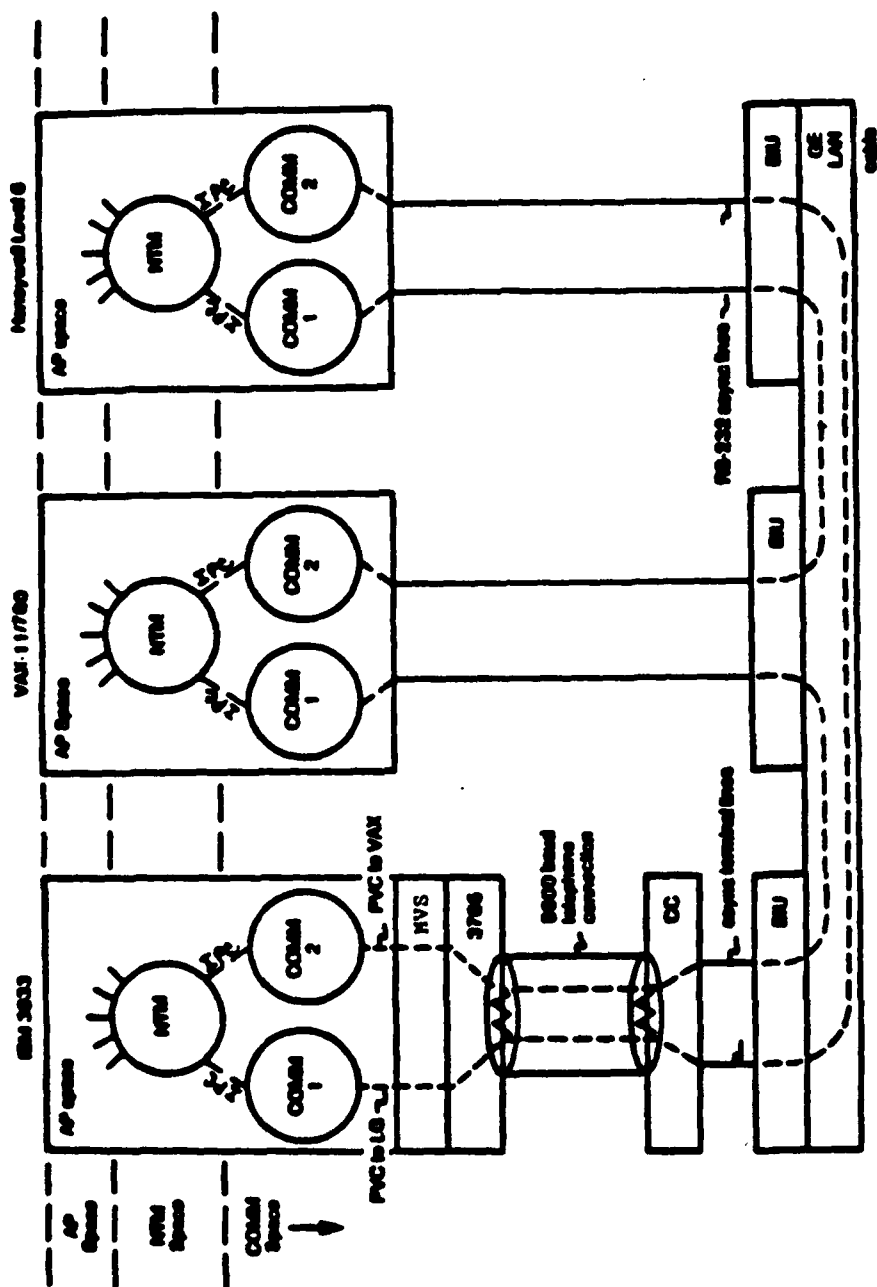


Figure 3-1. IISS Test Bed Communications Subsystem (Structural Schematic)



Figure 3-2. IISS Hardware Configuration

DS 620143000
1 November 1985

the circuit for each message of each session is removed. The COMM will be attached to each line when the computer is started or when IISS is initiated.

The communication subsystem is constrained in the following ways:

1. It must be implemented on each of the following systems with no modifications or additions to the operating systems or communication drivers.

VAX	-	VMS
Honeywell Level 6	-	MOD400
IBM 3084	-	MVS
2. A Local Area Network (LAN) will be used to interface these three systems. Terminal interfaces must be used since these three systems support no other way to uniformly interface to a LAN. The terminal interface to the IBM computer can be supported by using an IBM 3271 station emulator.
3. An MRP program has been chosen to run on IBM using the CICS sub-operating system. CICS supports only COBOL, PL-1, and assembly language. Of these three choices, COBOL has been selected as the implementation language. (This restriction was lifted in 1984.)
4. Programming in assembly language is to be avoided to the greatest extent possible.

These constraints place the following constraints on a data transmission protocol:

1. Full-duplex transmission is not supported.
2. Binary data transmission is supported with data translation.
3. Variable length data messages must be terminated by a carriage return.
4. Eight bit characters are not supported (seven bit characters with the eighth bit used for parity is supported).
5. Message headers, binary data translation, and the longitudinal redundancy checking technique have been chosen to allow for a COBOL implementation.

The communication protocol described in the detailed

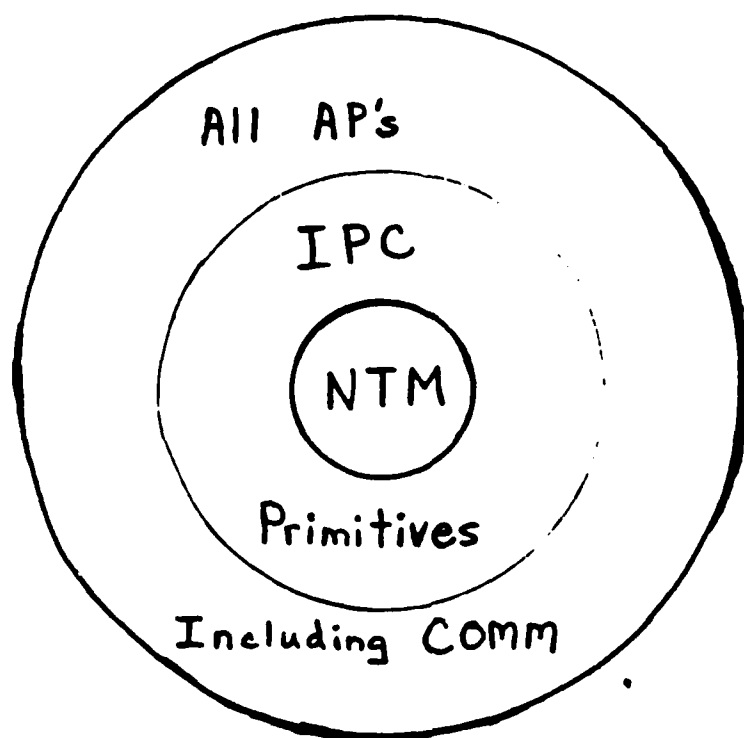


Figure 3-3. Interprocess Communication Architecture

DS 620143000
1 November 1985

functional requirements section satisfies these constraints and provides for point-to-point communication between any two of the three NTM's using three permanent virtual circuits on the Local Area Network.

The Interprocess Communication Primitives (IPC) are the mechanisms that the NTM will use to transfer data between itself and the application programs (AP) in the computer (see Figure 3-3). This approach removes the necessity of re-writing the NTM for each computer in the IISS configuration since the system-dependent software will be in the primitives.

Because of the highly system-dependent nature of the communication primitives, they must be implemented three times - once on each computer.

3.1.2 Interface Requirements

The Communication Subsystem consists of two copies of COMM, one resident in each host, enabling two NTM's to communicate with each other. In this sense the communication subsystem interfaces only with the NTM. This interface is accomplished with IPC's. To initiate communication with the NTM, COMM must use an NTM runtime service called INICOM. This service must supply to COMM its input communication mailbox name. COMM must call TRMNAT upon termination in order to stop all further communication with the NTM.

Two types of messages are received by COMM from the NTM:

1. Data messages (both binary and native character set)
2. Control messages

A field in the NTM message header indicates whether or not a message is a data message or a control message. If it is a data message, another field in the header determines whether the data message is "native" or "binary." If it is a control message, the type is included in the message header.

Two types of messages are sent to the NTM by COMM:

1. Data messages (both binary and native character set)
2. Status messages (includes statistics)

The communications subsystem interfaces with the LAN via the standard RS232C terminal interface.

3.1.2.1 Interface Block Diagram

The structural schematic for the communication subsystem is depicted in Figure 3-1. This shows the COMM interfaces to the NTM and the LAN.

3.1.2.2 Detailed Interface Definition of COMM

The Communication Subsystem receives messages from and sends messages to the NTM via the IPC's. Each message contains message data and a message header as described by the NTM. The COMM uses the following fields from the header:

- Destination AP Name
- Message Type
- Binary/Native Flag
- Priority Flag

Upon receiving messages from the NTM, the COMM checks for a control message by checking the destination AP name for COMM. If it is COMM, processing is determined by a message type of startup link (SL) or shutdown link (SD) or terminates (TR). If the destination is not COMM, COMM sends the message according to the binary/native flag.

Messages sent to the NTM from COMM are control messages or data messages. Control messages have the destination set equal to NTM, the priority flag equal to high, and are put in the high priority input queue for NTM. The message type may be link active (LA), link failure (LF) or recoverable error (RE). The data portion of the recoverable error message contains the error number. Data messages have the destination set equal to whatever it was on input to COMM. These messages are put in the correct NTM input queue according to the priority flag.

3.1.2.3 Detailed Interface Definition of Primitives

The VAX implementation of the primitives uses a combination of COBOL and FORTRAN. All system services are called using FORTRAN with most of the other codes such as error checking being done in COBOL.

The Level 6 implementation of the primitives uses a combination of COBOL and Assembly Language. All system services are called using Assembly Language because the Level 6 only supports an Assembly Language interface to system services.

In the IBM system, IISS is implemented in the Assembler and interfaces to the MVS operating system.

3.2 Detailed Functional Requirements

The node tree shown in Figure 3-4 illustrates the COMM functions currently defined.

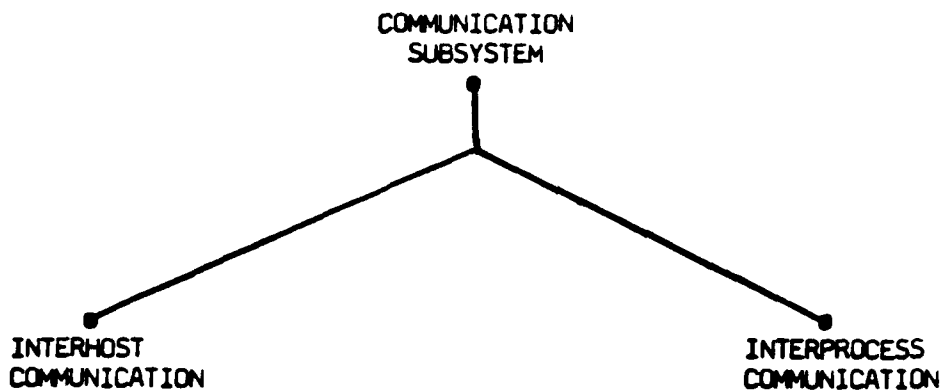


Figure 3-4. COMM Configuration Tree

Descriptions of these functions may be found in the following paragraphs.

3.2.1 Communication Protocol - Characteristics and Description

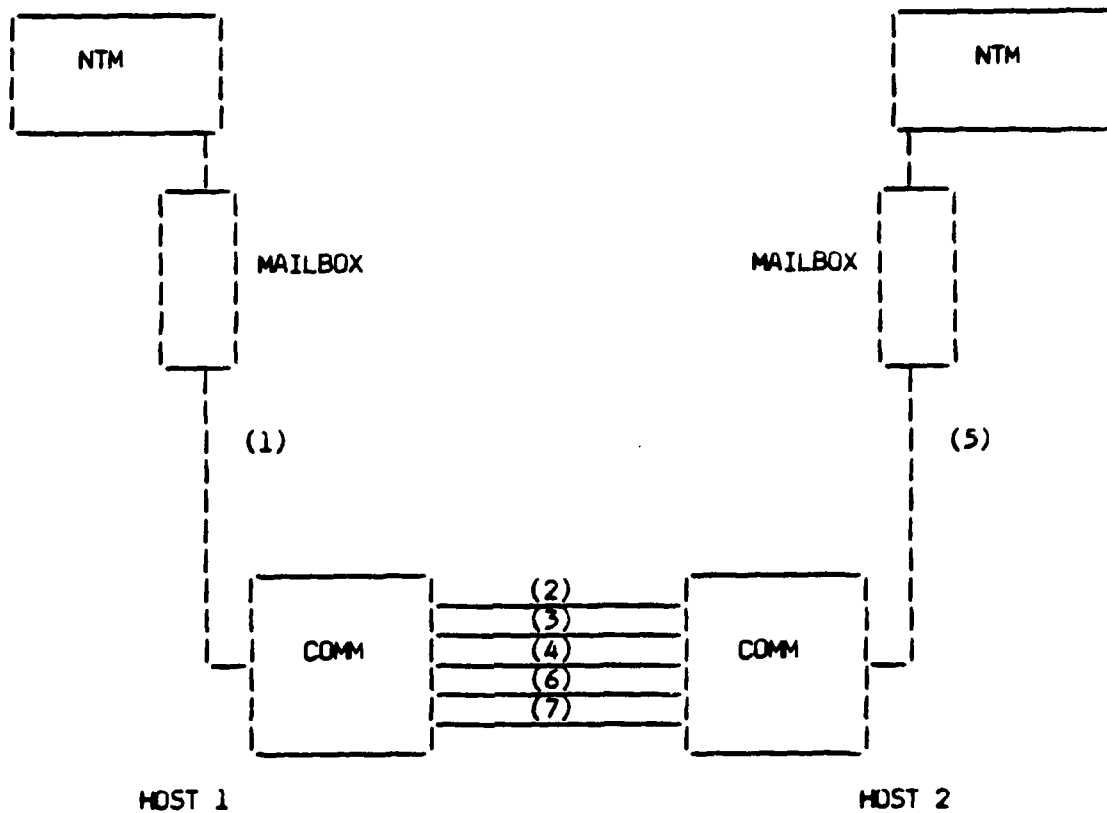
A simple scenario for sending a message between NTM's is depicted in Figure 3-5.

3.2.1.1 Characteristics of the Communication Protocol

1. Uses asynchronous communication lines.
2. Contention system with one end point designated as primary and the other as secondary.
3. Point-to-point.
4. Half-duplex (uses full-duplex communication lines).
5. Interleaved data transmission.

6. Uses the ASCII character set (excluding control characters 000-037 octal).
7. Error detection and correction by retransmission.
8. Byte stuffing is used to send the control characters. An exclamation mark (!) precedes a translated control character. The control character can then be reconstructed by the receiving COMM program.
9. Eighth bit used for even parity.
10. Accepts variable length NTM messages.
11. All messages terminated by carriage return.
12. Large messages are segmented into packets and reassembled by the receiving COMM.
13. Transmits variable length communication blocks.
14. Symmetric protocol with contention. The retry timing in case of simultaneous line bids is 1 second for the primary endpoint and 5 seconds for the secondary endpoint.
15. Retries a configurable number of times before reporting a link failure to the NTM. Currently, the number is three.
16. Master/Slave relationship determined by the endpoint that successfully bids for the line. Successful bidder becomes the master.
17. Data may be transmitted by either the master or the slave endpoint.
18. All timing is performed by the master. Time out interval is 3 seconds.
19. End of transmission bit may only be sent by the master and only when neither master nor slave has a

DS 620143000
1 November 1985



- (1) Message received
- (2) Line bid
- (3) Grant line
- (4) Message sent
- (5) Message delivered
- (6) ACK
- (7) EOT

Figure 3-5. Simple Scenario for Sending a Message Between NTM's

requirement to send data.

3.2.1.2 Sequence Number Handling

The send sequence number and the receive sequence number are used to prevent either duplication or loss of communication blocks. The first data block following a successful line bid is sent with a send sequence number of 1. Subsequent blocks are sent with the send sequence number cycling from 2 to 3 and back to 1. The received sequence number is used in reply to a sending COMM. A receiving COMM returns the send sequence number as the received sequence number after forwarding the data to the local NTM.

3.2.1.2.1 Send Sequence Number Checking

If the send sequence number is as expected, then the data is stored and that number is returned as the received sequence number. The expected send sequence number is incremented by one.

If the send sequence number is one less than expected, then that number is returned as the received sequence number. However, the data is not stored since it was stored on receipt of the prior message (which must have been repeated). The send sequence number should never be one more than expected.

3.2.1.2.2 Receive Sequence Number Checking

For a given copy of COMM the received sequence number from the correspondent COMM should be the same as its prior send sequence number indicating that the correspondent COMM received the communication block correctly. COMM then increments the send sequence number for the next block.

If the received sequence number is one less than the prior send sequence number, then COMM retransmits the data using the same send sequence number.

The received sequence number should never be one more than the prior send sequence number. If it is, COMM sends a negative acknowledgment indicating the error.

3.2.1.3 Protocol for Binary Data Transmission

When an NTM is required to send binary data using this transmission protocol, the binary information will be expanded

into an acceptable set of characters (0-9 and A-F), transmitted as characters, and transformed to binary upon receipt from the LAN. The following data translation algorithm will be applied to the data by the communication subsystem.

Each byte of data is transmitted as two bytes of ASCII data. The first byte represents the higher order 4 bits of data; the second, the last 4 bits.

This translation is reversed by the receiving station in order to rebuild the original NTM message.

3.2.1.4 Interface Between IBM and non-IBM Computers

When using standard terminal I/O drivers to transmit data between the IBM and non-IBM computers, there are four problems that must be considered:

1. Removing control characters from the message before sending it and inserting them back into the message after it has been received.
2. Using one character set to compute the longitudinal redundancy check on all computers.
3. Conversion or translation of certain characters because of unique problems. For example, spaces are converted to cursor control by protocol converters so spaces must be replaced by a special code before transmission.
4. Checking for characters in EBCDIC that do not have an equivalent character in ASCII.

These problems are solved in the three subroutines, KMINDA, EXOUDA, and KLCLRC, in the Communication Subsystem through the use of two tables. The tables are found in the include files CTLASC and ASCII on non-IBM computers and CTLEBC and EBCDIC on the IBM.

3.2.1.4.1 Non-IBM Computers

The CTLASC and ASCII include files are used in computers whose native character specification is ASCII. The table in the ASCII include file is used by the EXOUDA and KLCLRC subroutines to determine if there is an equivalent EBCDIC character, and, if it is a control character, what the substitution code is.

The table contains positive, negative, and zero values. The negative values indicate that the character being processed is a control character or a character that requires special conversion for transmission. The code to be used in the message is the negated negative value found in the table. If the value is positive, the original character is used in the message. If the value is zero, there is no equivalent character in the EBCDIC character set.

The table in the CTLASC include file is used by the KMINDA subroutine to restore the original message. Through the table, KMINDA determines the correct control character or special character to be inserted into the message in place of the code character found in the message. The code consists of alphanumeric characters that will not cause terminal I/O drivers to react in a special manner. The flag in the message that indicates a code character follows is the (!).

3.2.1.4.2 IBM Computers

The CTLEBC and EBCDIC include files are used in the IBM computers where the native character specification is EBCDIC. The table in the EBCDIC include file is used for two functions. In EXOUDA it is used to determine if there is a comparable ASCII character, and, if there is an ASCII equivalent, is it a control character or a special character. In KLCLRC the table in EBCDIC is used to convert all the characters in the message to ASCII in order to perform the sum for the longitudinal redundancy check. (Negative values in the table are never used in KLCLRC because conversion is done by EXOUDA before KLCLRC is called.)

The contents of the table in the EBCDIC include file on the IBM have the same definitions for positive, negative, and zero values as its counterpart file, ASCII, does in the non-IBM environment.

The table in the CTLEBC include file is used by the KMINDA subroutine in the Communication Subsystem on the IBM in exactly the same manner as KMINDA in the Communication Subsystem on the non-IBM computers.

3.2.2 Message Format

Each message contains seven characters in addition to the data characters. The first three characters are called the header, and the last four characters are called the trailer.

This is depicted in Figure 3-6.

Header Character #1 - Send Data Sequence Number Byte

ASCII Character

0	No send data message
1	Send Sequence Number 1
2	Send Sequence Number 2
3	Send Sequence Number 3

Source sequence number cycles from 1 to 2 to 3 and back to 1 when data is being transmitted.

Header Character #2 - Received Data Sequence Number Byte

ASCII Character

0	No received data message
1	Received Sequence Number 1
2	Received Sequence Number 2
3	Received Sequence Number 3

Header Character #3 - Control Byte

ASCII Character

The following 8 control chars do not accompany a data msg.

0	Positive Acknowledgment
1	Line Bid
2	On-Line
3	End of Transmission
4	NAK - Parity Error
5	NAK - No Buffer Space
6	NAK - Bad Receive Sequence Number
7	NAK - Bad Send Sequence Number

The following 4 control chars always accompany a data msg.

A	ACK - Sending Native Data
B	ACK - Sending Binary Data
C	ACK - Sending Continued Native Msg.
D	ACK - Sending Continued Binary Msg.

Trailer Characters

#1, 2, & 3 - Block Check

These three characters contain a message block check. The block check is computed by adding all bytes in the message (including

DS 620143000
1 November 1985

the header bytes). This block check could be up to 18 bits in size for a message block of size 2048 bytes. The 18 bit additive result is split into three 6 bit quantities. Each 6 bit quantity is represented by an ASCII character, with the three characters being stored as the last characters in the message (high order byte occurring first).

The block check characters occur only in messages containing data.

Trailer Character #4 -

Carriage Return

DS 620143000
1 November 1985

Send	Rcv.	Ctrl.						C
Seq.	Seq.	Byte	DATA	Bcc	Bcc	Bcc	R	
No.	No.			1	2	3		

or

Send	Rcv.	Ctrl.	C
Seq.	Seq.	Byte	R
No.	No.		

SEND SEQUENCE

- 0 - No send data
- 1 - Send sequence number 1
- 2 - Send sequence number 2
- 3 - Send sequence number 3

RECEIVE SEQUENCE NUMBER

- 0 - No received data
- 1 - Receive sequence number 1
- 2 - Receive sequence number 2
- 3 - Receive sequence number 3

CONTROL BYTE

THE FOLLOWING CHARACTERS DO NOT ACCOMPANY DATA

- 0 - Positive acknowledgment
- 1 - Line bid
- 2 - End of transmission
- 3 - NAK - Block check error
- 4 - NAK - NTM input mailbox full
- 5 - NAK - Bad send sequence number
- 6 - NAK - Bad receive sequence number

THE FOLLOWING CHARACTERS ALWAYS ACCOMPANY DATA

- A - ACK - Sending native data
- B - ACK - Sending binary data
- C - ACK - Sending continued native data
- D - ACK - Sending continued binary data

Figure 3-6. Communication Block Diagram

3.2.3 Timing

3.2.3.1 Line Idle

When the line is idle, the line bid message may be sent by either endpoint. If a collision of line bids occur or if there is no response, the primary waits for 1 second before retransmitting a line bid, and the secondary waits for 5 seconds.

3.2.3.2 Master

The endpoint successfully bidding for the line becomes the master. The master waits for 5 seconds for a response from a slave. A time-out causes retransmission of the prior communication block.

3.2.3.3 Slave

The slave waits for

(Maximum number of retries x 5 secs) + 5 seconds

for a response from the master. If data is not received within this time, the slave assumes an end of transmission was missed and returns to an idle state. If a partial message is being assembled, a link failure is reported and the partial message is discarded.

3.2.3.4 NTM Input Mailbox Full

When a copy of COMM discovers that the NTM input mailbox is full, it waits for one second and tries again to place the message in the mailbox. If the condition still exists, a NAK message is returned to the correspondent COMM. The correspondent COMM sends an error status message to its local NTM and retransmits the communication block.

3.2.4 Block Check Computation

The block check is computed by adding all bytes in the communication block excluding the block check itself and the carriage return. This could account for an 18 bit block check for a message block of size 2048 bytes. The 18 bit additive result is split into three 6 bit quantities. Each 6 bit quantity is represented by a valid ASCII character with the

three characters being stored as the last three characters in the message (high order byte occurring first).

3.2.5 Control Characters in Data

The 32 ASCII control characters (octal 0-37) affect the behavior of the I/O handlers and therefore cannot be allowed in the data stream. Byte stuffing will be used to convert each control character to a valid ASCII character preceded by a usable special character such as data link escape.

3.2.6 Communication Subsystem Primitives

The design of the COMM is made up of a large generic portion that will be the same for all computers, and a small host-specific part (primitives) that is specially developed for each computer. The host-specific part is a group of routines called the Interhost Communication Primitives (IHC's). The following is a description of these routines.

3.2.6.1 Initialize Communication Port Interface

Calling Sequence:

```
CALL INILAN USING PORT-NAME,  
                  RCV-BLOCK,  
                  XMIT-BLOCK,  
                  EVENT-BLOCK-nn,  
                  STATUS.
```

Description:

INILAN moves the PORT-NAME or some system dependent equivalent to the appropriate storage in the XMIT and RCV blocks. If initialization for some system services associated with the port is required, it is performed in this primitive. INILAN also initializes the XMIT and RCV blocks with character zeros.

Inputs:

```
PORT-NAME  
RCV-BLOCK  
XMIT-BLOCK  
EVENT-BLOCK-nn
```

Outputs:

STATUS

Possible Status Conditions:

**Successful
System dependent errors**

3.2.6.2 Send a Message to a Terminal Port

Calling Sequence:

**CALL "XMTLAN" USING XMIT-BLOCK,
EVENT-BLOCK-nn,
FLAGS,
STATUS.**

Description:

XMTLAN outputs a message to a given terminal port.

Inputs:

**XMIT-BLOCK
EVENT-BLOCK-nn
FLAGS**

Outputs:

STATUS

Possible Status Conditions:

**Successful
Number of bytes zero
Number of bytes greater than maximum
Receive LAN outstanding
System dependent errors**

Notes:

1. The XMIT-BLOCK contains the port name, the number of bytes to be transmitted, and the buffer with the message. (See the next section for a detailed description of the XMIT-BLOCK.)

DS 620145000
1 November 1985

2. A good status indicates that the message has been accepted for transmission. It does not mean a successful transmission.
3. The number of bytes to be transmitted must be at least one and no more than 1024. The maximum number of bytes is a communication variable that can be changed.
4. The event block that is passed to XMTLAN must be the same as the one that is passed to RCVLAN and GETLAN.

3.2.6.3 Receive a Message from a Terminal Port

Calling Sequence:

CALL "RCVLAN" USING RCV-BLOCK,
EVENT-NUMBER,
EVENT-BLOCK-nn,
FLAGS,
STATUS.

Description:

RCVLAN informs the operating system that the program will accept a message from the given terminal port.

Inputs:

RCV-BLOCK
EVENT-NUMBER
EVENT-BLOCK-nn

Outputs:

STATUS

Possible Status Conditions:

Successful
Only one receive outstanding permitted
Event number zero
Event number greater than maximum
System dependent errors

Notes:

1. Only one receive may be outstanding for a given terminal port.
2. The event number may have a value of 01 through 22. A value of zero or 23 through 99 causes an error status to be returned. The maximum number of events outstanding that can be waited on at any one time is 22 because COBOL on the Level 6 limits the number of arguments in a calling sequence to 25.
3. The event number must be unique.
4. The value of the event number is the priority of the receive message request in relation to the other outstanding requests. The lower the value, the higher the priority.
5. The RCV-BLOCK contains the port name, the buffer size, and the buffer into which the message will be stored. (See the next section for a detailed description of the RCV-BLOCK.)
6. The RCV-BLOCK that is passed to RCVLAN must be the same one that is passed to GETLAN when the program actually gets the message from the given terminal port.
7. The event block that is passed to RCVLAN must be the same one that is passed to GETLAN when the program actually gets the message from the given terminal port. The event block is also the same one that is passed to XMTLAN to send a message to the port.

3.2.6.4. Get a Message from a Terminal Port

Calling Sequence:

CALL "GETLAN" USING RCV-BLOCK,
EVENT-BLOCK-nn,
STATUS.

Description:

GETLAN accepts the message that was received from the given terminal port and moves it into the given buffer.

DS 620143000
1 November 1985

Inputs:

RVC-BLOCK
EVENT-BLOCK-nn

Outputs:

STATUS

Possible Status Conditions:

Successful
Receive not satisfied
Not a LAN event block
No receive outstanding
Buffer size zero
Buffer size greater than maximum
Buffer too small
System dependent errors

Notes:

1. The RCV-BLOCK contains the port name, the buffer size, the buffer, and a location into which the number of bytes in the message is stored by GETLAN.
2. The RVC-BLOCK that is passed to GETLAN must be the same one that was passed to RCVLAN for the given terminal port.
3. The event block that is passed to GETLAN must be the same one that was passed to RCVLAN for the given terminal port. The event block is also the same one that is passed to XMTLAN to send a message to the port.
4. If the buffer size is too small for the entire message, an error status is returned and the message is lost. There is no longer a receive outstanding.
5. If no receive is outstanding for the given terminal port, an error status is returned.

3.2.6.5 Cancel a Receive Request from a Terminal Port

Calling Sequence:

CALL "CNLLAN" USING RCV-BLOCK,
EVENT-BLOCK-nn,
STATUS.

Description:

CNLLAN removes the receive outstanding for a given terminal port.

Inputs:

RCV-BLOCK
EVENT-BLOCK-nn

Outputs:

STATUS

Possible Status Conditions:

Successful
No receive outstanding
Not a LAN event block
System dependent errors

Notes:

1. The RCV-BLOCK that is passed to CNLLAN must be the same one that was passed to RCVLAN when the receive was requested.
2. The event block that is passed to CNLLAN must be the same one that was passed to RCVLAN when the receive was requested.
3. The RCV-BLOCK contains the port name, the buffer size, the buffer into which the message will be stored, and a location for the actual number of bytes in the message. (See the next section for a detailed description of the RCV-BLOCK.)
4. If there was a message present, it is lost.

DS 620143000
1 November 1985

3.2.6.6 Terminate Communication Port Interface

Calling Sequence:

CALL "TRMLAN" USING PORT-NAME,
RCV-BLOCK,
XMIT-BLOCK,
EVENT-BLOCK-nn,
STATUS.

Description:

TRMLAN is needed for the IBM environment only. It issues VTAM calls to disconnect from the port. The implementation on the other computers is a stub.

Inputs:

PORT-NAME
RCV-BLOCK
XMIT-BLOCK
EVENT-BLOCK-nn

Outputs:

STATUS

3.2.7 Interprocess Communication Primitives

The IPC primitives are used to transfer messages between tasks on the same computer. They have been designed to localize all host dependent code in several small routines. These routines are referred to as the Interprocess Communication Primitives or IPC's.

3.2.7.1 Create a Mailbox

Calling Sequence:

CALL "CRTMBX" USING INPUT-MAILBOX-NAME,
MAILBOX-SIZE,
EVENT-BLOCK-nn,
STATUS.

Description:

CRTMBX creates a mailbox through which the program will

DS 620143000
1 November 1985

receive messages from another program running on the same computer.

Inputs:

INPUT-MAILBOX-NAME
MAILBOX-SIZE
EVENT-BLOCK-nn

Outputs:

STATUS

Possible Status Conditions:

Successful
Invalid mailbox name
Mailbox already exists
Mailbox size zero
Mailbox size greater than maximum
Event block not initialized
System dependent errors

Notes:

1. If the input mailbox has been created previously, an error status is returned.
2. The event block that is passed in the CRTMBX call is the same event block that is passed when the program issues RCVMSG and GETMSG for the input mailbox.
3. The mailbox size is only applicable on the VAX and IBM implementations of the CRTMBX primitive. The resource wait mode on the VAX must be disabled to permit the program to regain control immediately upon detection of mailbox full.

3.2.7.2 Send a Message to Another Program

Calling sequence:

CALL "SNDMSG" USING TARGET-MAILBOX-NAME,
BUFFER,
NUMBER-OF-BYTES,
EVENT-BLOCK-nn,
STATUS.

DS 620143000
1 November 1985

Description:

SNDMSG sends a message to another program running on the same computer through the input mailbox of the other program. The event block is system dependent storage that is required by SNDMSG. It is not associated with an event that can be waited on.

Inputs:

TARGET-MAILBOX-NAME
BUFFER
NUMBER-OF-BYTES
EVENT-BLOCK-nn

Outputs:

STATUS

Possible Status Conditions:

Successful
Mailbox not found
Mailbox full
Number of bytes zero
Number of bytes greater than maximum
System dependent errors

Notes:

1. The receiving program must have previously created its input mailbox. If no mailbox exists with the given name, an error status is returned.
2. A good status indicates that the message has been accepted by the operating system for transfer. It does not mean that the message has been received by the other program.
3. If the status of "mailbox full" is returned, the program must retry at a later time.
4. The number of bytes to be sent must be at least one and no more than 2000.
5. The event block must not be in use for an input

mailbox or a timer. An event block associated with a mailbox is in use if the mailbox has been created but not deleted. An event block associated with a timer is in use if the timer has been started but not cancelled or runout during a WAITnn.

6. A series of SNDMSG calls being directed to a single target mailbox should use the same event block. The VAX assigns a logical channel ID to a target mailbox at the time of the first SNDMSG call and uses that channel ID for subsequent SNDMSG calls. This event block cannot be used for other purposes until it is released (see RELEVB).

3.2.7.3 Receive a Message from Another Program

Calling sequence:

```
CALL "RCVMSG" USING INPUT-MAILBOX-NAME,  
                     EVENT-NUMBER,  
                     EVENT-BLOCK-nn,  
                     STATUS.
```

Description:

RCVMSG informs the operating system that the program will accept a message sent from another program to the given input mailbox. The program continues executing. The fact that a message has been sent by another program is obtained with the WAITnn or GETMSG primitive.

Inputs:

```
INPUT-MAILBOX-NAME  
EVENT-NUMBER  
EVENT-BLOCK-nn
```

Outputs:

```
STATUS
```

Possible Status Conditions:

```
Successful  
Invalid event block for mailbox named  
Not a receive event block  
Only one outstanding receive permitted
```

Event number zero
Event number greater than maximum
System dependent errors

Notes:

1. Only one receive may be outstanding for a given input mailbox.
2. The event block that is passed to RCVMSG must be the same event block that was passed to CRTMBX when the given input mailbox was created.
3. The event number must be unique. (See the description of the event number in Section 3.5.2)
4. The value of the event number is the priority of the receive message request in relation to the other outstanding requests. The lower the value, the higher the priority.
5. The event number may have a value of 01 through 22. A value of zero or 23 through 99 causes an error status to be returned. Because COBOL on the Level 6 limits the number of arguments in a calling sequence to 25, the maximum number of outstanding events that can be waited on at any one time is 22.

3.2.7.4 Get a Message from Another Program

Calling Sequence:

CALL "GETMSG" USING INPUT-MAILBOX-NAME,
BUFFER,
BUFFER-SIZE,
NUMBER-OF-BYTES,
EVENT-BLOCK-nn,
STATUS.

Description:

GETMSG accepts the message that was sent from another program running on the same computer and moves it into the given buffer.

Inputs:

INPUT-MAILBOX-NAME
BUFFER-SIZE
EVENT-BLOCK-nn

Outputs:

BUFFER
NUMBER-OF-BYTES
STATUS

Possible Status Conditions:

Successful
Invalid event block for mailbox name
Not a receive event block
No receive outstanding
Receive not satisfied
Buffer too small
Buffer size zero
Buffer size greater than maximum
System dependent errors

Notes:

1. If no receive is outstanding for the given input mailbox, an error status is returned.
2. If a receive is outstanding but no message has been delivered by the operating system, a status is returned indicating that no message has been received.
3. If the buffer size is too small for the entire message, an error status is returned and the message is lost. An outstanding receive for that mailbox no longer exists.
4. The event block that is passed to GETMSG must be the same one that was passed to RCVMSG for the given input mailbox.

3.2.7.5 Delete a Mailbox

Calling Sequence:

CALL "DELMBX" USING INPUT-MAILBOX-NAME,
EVENT-BLOCK-nn,
STATUS.

DS 620143000
1 November 1985

Description:

DELMBX removes the ability to receive messages from another program through the given input mailbox.

Inputs:

INPUT-MAILBOX-NAME
EVENT-BLOCK-nn

Outputs:

STATUS

Possible Status Conditions:

Successful
Invalid event block for mailbox named
Not a receive event block
System dependent errors

Notes:

1. The event block that is passed to DELMBX must be the same event block that was passed to CRTMBX when the given input mailbox was created.
2. If there are messages remaining in the mailbox, they are lost.
3. The given event block is re-initialized.

3.2.7.6 Release an Event Block

Calling Sequence:

Call "RELEVB" USING TARGET-MAILBOX-NAME,
EVENT-BLOCK-nn,
STATUS.

Description:

RELEVB releases an event block which had been used for sending messages to a target mailbox. The event block is cleared and may be used for another purpose.

DS 620143000
1 November 1985

Inputs:

TARGET-MAILBOX-NAME
EVENT-BLOCK-nn

Outputs:

STATUS

Possible Status Conditions:

Invalid event for mailbox named
System dependent errors (VAX only)

Notes:

1. The VAX implementation of RELEVb deassigns the logical channel previously assigned to the target mailbox by SNDMSG.
2. All logical channels assigned to target mailboxes by SNDMSG primitive calls on the VAX are deassigned when a process terminates. Therefore, for many applications, a call to RELEVb is not required.

3.2.7.7 Start a Timer

Calling Sequence:

CALL "SETTIM" USING TIME-INTERVAL,
EVENT-NUMBER,
EVENT-BLOCK-nn,
STATUS.

Description:

SETTIM requests the operating system start a timer with the given time interval. The program continues executing. The fact that the timer has elapsed is obtained from the WAITnn primitive. Only one timer may be active at a time.

Inputs:

TIME-INTERVAL
EVENT-NUMBER
EVENT-BLOCK-nn

Outputs:

STATUS

Possible Status Conditions:

Successful
Time interval zero
Time interval greater than maximum
Event number zero
Event number greater than maximum
Event block not initialized
System dependent errors
Timer already active

Notes:

1. The time interval is given as HHMMSS.
2. The minimum time interval is 000001. The maximum time interval is 24 hours, 59 minutes, 59 seconds.
3. The event number must be unique. (See the description of the event number in Section 3.5.2)
4. The value of the event number may be 01 through 22. A value of zero or 23 through 99 causes an error status to be returned. The maximum number of outstanding events that may be waited on at any one time is 22 because COBOL on the Level 6 limits the number of arguments in a calling sequence to 25.
5. The value of the event number in relation to the values of other event numbers is an indication of priority for the WAITnn primitive. The event number with the lowest value has the highest priority.
6. The event block that is passed to SETTIM must be the same event block that is passed to CNLTIM to cancel the timer.

3.2.7.8 Stop a Timer

Calling Sequence:

CALL "CNLTIM" USING EVENT-BLOCK-nn,
STATUS.

DS 620143000
1 November 1985

Description:

CNLTIM cancels the request made to the operating system by SETTIM to be notified when a given time interval has passed.

Inputs:

EVENT-BLOCK-nn

Outputs:

STATUS

Possible Status Conditions:

Successful
Not a timer event block
System dependent errors

Notes:

1. The event block that is passed to CNLTIM must be the same event block that was passed to SETTIM to start the timer.
2. The given event block is re-initialized.

3.2.7.9 Wait for an Event to Occur

Calling Sequence:

CALL "WAITnn" USING EVENT-NUMBER,
STATUS,
NUMBER-OF-EVENT-BLOCKS,
EVENT-BLOCK-01,
EVENT-BLOCK-02,
.
.
.
EVENT-BLOCK-xx.

where nn is a number from 01 through 22 that indicates the number of event blocks being passed

Description:

DS 62Q143000
1 November 1985

WAITnn waits for one of the outstanding requests that are associated with the list of event blocks to be satisfied. The event number associated with the completed request is returned in the EVENT-NUMBER variable.

Inputs:

NUMBER-OF-EVENT-BLOCKS
EVENT-BLOCK-01
EVENT-BLOCK-02
.
.
.
EVENT-BLOCK-xx

Outputs:

EVENT-NUMBER
STATUS

Possible Status Conditions:

Successful
Number of event blocks zero
Number of event blocks greater than
maximum
List of event blocks greater than
number of event blocks
Event numbers not unique
No requests outstanding
System dependent errors

Notes:

1. The order in which the testing for the completion of a request is performed is based upon the event number associated with each event block. The request with the lowest event number has the highest priority; therefore, the check for its completion is done first. The request with the next lowest event number is checked second, and so forth. This process continues until a request associated with one of the given event blocks has been satisfied.
2. If two requests have been completed, the one with the lowest event number is indicated to the calling program. The information that another request has been

satisfied is retained by the primitive until the WAITnn is called again.

3. The maximum number of event blocks that may be passed is 22. Because COBOL on the Level 6 limits the number of arguments in a calling sequence to 25, the maximum number of outstanding events that can be waited on at any one time is 22.
4. At least one event block must be passed.
5. Each event block must have a unique event number associated with it; otherwise, an error status is returned.
6. It is not required to have an outstanding request associated with each event block passed to WAITnn. Because the primitive is able to differentiate between event blocks that have outstanding requests and those that do not, the WAITnn primitive may be passed all the event blocks that would be needed for the worst case. If there are no requests outstanding, however, an error status is returned.
7. If an error status is returned, the event number is set to character zeros.
8. The 'list of event blocks greater than number of event blocks' error is checked only by the IBM.

3.2.7.10 Terminate a Program

Calling sequence:

CALL "ENDRUN".

Description:

ENDRUN calls the appropriate system routine to stop execution of the particular program without halting the IISS system.

Inputs:

None

Outputs:

None

3.2.7.11 Save an Event Indicator

Calling Sequence:
CALL "LOCKEF".

Description:

LOCKEF is applicable only to the VAX. Its purpose is to obtain event indicator 63 from the system so that none of the other primitives can use it. LOCKEF is needed to work around a bug in VAX-11 DBMS. It is called by the Request Processors that manipulate data store under the VAX-11 DBMS.

Inputs:

None

Outputs:

None

3.2.7.12 Request an Error Be Logged

ERRPRO -- Process Error

Description:

This module is used to process errors. It gets the current date and time, formats the message and writes the message to the mailbox ERRMBX. In case of an error, this module will print a message on the operator hardcopy console. It will not terminate the processing of the calling program, but will return control to the calling program regardless of whether the error message is being written to the mailbox or not. Another task, ERRLOG, will read the error message from the mailbox and log the message on a file. See module specification of ERRLOG for details of this error logging task. The format of an error message is as shown:

Byte	Data
1 - 2	function code (DA)
3 - 10	current date (YY/MM/DD)

DS 620143000
1 November 1985

11 - 18 current time (HH:MM:SS)
19 - 24 module name
25 - 29 return status
30 - 89 message description

INTERFACES:

ENTRY CONDITIONS:

RET-STATUS	PIC X(5).
MODULE-NAME	PIC X(6).
MESG-DESC	PIC X(60).

EXIT CONDITIONS:

(NONE)

GLOBAL BLOCKS:

(NONE)

DATA ORGANIZATION:
LOCAL VARIABLES:

ERR-STATUS	PIC X(5).	--ERROR STATUS
------------	-----------	----------------

DATABASE INTERACTION:

LIMITATIONS:

Depending on the operating system, we may have a problem trying to write a message to the operator console. If so, some other mechanism will be used to log the error status on a hardcopy device in case this module encounters an error while writing to the ERRMBX mailbox.

3.2.7.13 Log an Error

ERRLOG -- PERFORM ERROR LOGGING TO A FILE

Description:

This program reads a message from mailbox ERRMBX. The format of the message is as follows:

Byte	Data
------	------

DS 620143000
1 November 1985

1 - 2	Function code e.g., DA -- error message data to be logged to file ERRLOG.DAT CF -- close current version of file and open a new version of the file ST -- close current version of file and terminate processing
3 - 109	Data portion of message -- If the function code is DA, the data portion of the message is formatted as follows:
Byte	Data
3 - 10	current date (YY/MM/DD)
12 - 19	current time (YY/MM/DD)
21 - 35	process name
37 - 42	module name
44 - 48	return status
50 - 109	message description
11, 20, 36, 43, 49	(blank)

Depending on the function code, this routine will branch to the corresponding part of the program. If an error is detected in this module, it will try to display the error on the operator's hard copy device. This module is executed as soon as IISS is brought up. It will check to see if mailbox ERRMBX is already created or not. If not, it will create the mailbox. Otherwise it will set up the event block for a future receive. Mailbox ERRMBX is a temporary mailbox and there is no need to delete or clean up the mailbox when this program is brought up. If there is no message in the mailbox, this program will suspend processing and wait for a message to arrive at mailbox ERRMBX. If the message contains function code ST, this routine will close the current version of the file and terminate processing. If the message contains function code CF, this routine will close the current version of the file and open a new version of the file. Then it continues to read messages from the mailbox. If the message contains function code DA, it will write a record to the ERRLOG.DAT file and continue to read messages from the mailbox.

DATA ORGANIZATION:

LOCAL VARIABLES:

DATABASE INTERACTION:

ERRLOG.DAT -- this is an indexed file of record size (102) with the following format:

Byte	Data
1 - 8	current date (YY/MM/DD)
9 - 16	current time (YY/MM/DD)
17 - 31	process name
32 - 37	module name
38 - 42	return status
43 - 102	message description

LIMITATIONS:

LEVEL 6 MOD 400 operating system does not keep track of the version number of the file. In order to keep different versions of the file, we may have to rename the current version of the file and then create a new version using ERRPRO.DAT as the file name.

3.3 Special Requirements

3.3.1 Programming Methods

COMM programming methods shall conform to the standards set forth by General Electric in the IISS Software Development Guidelines/Conventions document. Principles of structured design and programming will be adhered to.

3.3.2 Expandability

The design constraints on the communications subsystem were to follow the ISO reference model, use a local area network, and, if possible, use standard, vendor-supported software drivers. The objective is to avoid writing new drivers with all of the maintenance problems that this entails, and to avoid developing any special hardware. Various communications packages and approaches were reviewed and evaluated with respect to their applicability to the local area network usage, the availability for the computers to be used on the Test Bed, and their adherence and extension to the ISO reference model. The final conclusion was that no communication protocol existed or was in the detailed specification stage, even for the lower

levels, to fully satisfy the needs and constraints of the Test Bed and at the same time not to unduly overburden the system with capabilities that are not needed for a LAN-based system.

Based on the evaluation of currently available packages, and in consultation with Computer Technology Associates (active on the ISO committee), it was decided to develop a protocol based on the standard "bisynch" protocol for level 2 that would serve in the interim period and be compatible with the substitution of new standard protocols when they become available. Also, local area network vendors will begin supplying compatible host protocols. In the meantime, the protocols developed for the project will use standard terminal drivers supplied by the computer vendors and will interface to the local area network as a standard terminal, making the interface completely standard for the LAN.

Since the system dependent software for each host computer on the IISS is implemented in the primitives, additional types of computers can be added to the IISS and only the primitives have to be reimplemented.

3.4 Human Performance

Not Applicable

3.5 Data Description of Arguments Used with the Primitives

3.5.1 Data Description of Arguments for Communication Subsystem Primitives

- Xmit-block

The argument XMIT-BLOCK is a variable name associated with a block of contiguous storage. It contains the port name, the number of bytes to be transmitted, and the buffer for the message. The message consists of the header, the data, and the longitudinal redundancy check (LRC). The location of the LRC is the three bytes immediately following the data. XMIT-BLOCK also contains the buffer size so that both blocks can be manipulated by the same subroutines. Its description in COBOL is

01 XMIT-BLOCK.

03 XMIT-PORT-NAME	PIC X(4).
03 XMIT-NO-OF-CHARS	PIC 9(4).
03 XMIT-BUFFER-SIZE	PIC 9(4).

DS 620143000
1 November 1985

```
03 XMIT-BUFFER.  
05 XMIT-HEADER1          PIC 9.  
05 XMIT-HEADER2          PIC 9.  
05 XMIT-HEADER3          PIC X.  
05 XMIT-DATA             PIC X(1021).  
03 XMIT-SEQUENCE-NO      PIC 9
```

- Rcv-block

The argument RCV-BLOCK is a variable name associated with a contiguous block of storage. It contains the port name, the number of bytes in the message just received, the maximum number of bytes the buffer can hold, and the buffer for the message. Its description in COBOL is

```
01 RCV-BLOCK.  
03 RCV-PORT-NAME          PIC X(4).  
03 RCV-NO-OF-CHARS        PIC 9(4).  
03 RCV-BUFFER-SIZE        PIC 9(4).  
03 RCV-BUFFER.  
05 RCV-HEADER1            PIC 9.  
05 RCV-HEADER2            PIC 9.  
05 RCV-HEADER3            PIC X.  
05 RCV-DATA               PIC X(1021).  
03 RCV-SEQUENCE-NO        PIC 9.
```

- Flags

The argument FLAGS is a variable name associated with a contiguous block of storage. It contains data that indicates whether this version of COMM is primary, what the state of COMM is, time settings, and the host and target computer indicators.

- Port-name

The argument PORT-NAME contains the alphanumeric characters required by the operating system to indicate the exact terminal port the COMM program will use.

```
01 PORT-NAME              PIC X(12)
```

3.5.2 Data Description of Arguments for IPC Primitives

- Input-mailbox-name

The argument INPUT-MAILBOX-NAME contains a 14-character alphanumeric (no embedded blanks) that is used by the program to indicate it will accept and process messages from other programs running on the computer if they are sent to a mailbox with this 14-character label. It is recommended that a different mailbox name be used to receive messages from different programs. Its description in COBOL is

01 input-mailbox-name-x PIC X(14).
- Target-mailbox-name

The argument TARGET-MAILBOX-NAME contains the 14-character alphanumeric that is the input mailbox name for the program to which the message is to be sent. Its description in COBOL is

01 target-mailbox-name-x PIC X(14).
- Mailbox-size

The argument MAILBOX-SIZE contains the the amount of storage in bytes the programmer wants allocated by the operating system for the given mailbox. Its description in COBOL is

01 mailbox-size-x PIC 9(5).
- Buffer

The argument BUFFER is a variable name that is associated with a given amount of contiguous memory. The amount of memory should be enough to contain the largest single message that will be sent or received. Its description in COBOL is

01 buffer-x PIC X(2000).
- Number-of-bytes

The argument NUMBER-OF-BYTES contains the actual number of characters that, in the case of SNDMSG, are to be sent. In the case of GETMSG it is the number of bytes that were moved into the buffer. Its description in COBOL is

01 number-of-bytes-x PIC 9(4).
- Buffer-size

The argument BUFFER-SIZE contains the maximum number of

DS 620143000
1 November 1985

bytes that can be stored in the buffer by GETMSG. Its description in COBOL is

01 buffer-size-x PIC 9(4).

- Number-of-event-blocks

The argument NUMBER-OF-EVENT-BLOCKS contains the number of event blocks that are being passed to the WAITnn block. Its description in COBOL is

01 number-of-event-blocks-x PIC 99.

- Time-interval

The argument TIME-INTERVAL contains the number of hours, minutes and seconds the timer is to count before it returns a request complete. Its description in COBOL is

01 time-interval-x.

03 time-in-hours PIC 99.

03 time-in-minutes PIC 99.

03 time-in-seconds PIC 99.

The maximum time interval is 24 hours, 59 minutes, 59 seconds.

- Status

The argument STATUS contains a code that indicates whether the primitive was successful or not; and, if it was not, what the problem was. Its description in COBOL is

01 status-x PIC X(5)

- Event-block

The argument EVENT-BLOCK is a variable name that is associated with a block of contiguous memory to be used by the primitive. It is, therefore, system dependent storage that may not be used by the program. Since it is system dependent, its size will vary from host to host. There are a set of include files that a programmer may use to define the event blocks. The names of the files are

EVBK01.system standard suffix
EVBK02. " " "

EVBK22. " " "

The programmer must decide the maximum number of event blocks that will be needed by the program at any one time. In the initialization section of the program, all the event blocks must be initialized to character zeros. The event block will be re-initialized by the primitives when the timer runs out or is cancelled and when a mailbox is deleted.

The event block differs from the other arguments passed to the primitives. For all the others, the primitives need the contents or value of the arguments. In the case of the event block, however, the primitives need its address. Therefore, a dummy argument for an event block will not work. The following code illustrates the incorrect passing of event blocks.

MOVE EVENT-BLOCK-01 TO EVENT-BLOCK.
PERFORM CREATE-MAILBOX.

MOVE EVENT-BLOCK-02 TO EVENT-BLOCK.
PERFORM CREATE-MAILBOX.

CREATE-MAILBOX.
CALL "CRTMBX" USING INPUT-MAILBOX-NAME,
MAILBOX-SIZE,
EVENT-BLOCK,
STATUS.

Programmers should use the CASE statement to invoke the primitives requiring an event block as an argument.

- Event-number

The argument EVENT-NUMBER is a code that is returned by the WAITnn primitive to indicate which of the outstanding requests was satisfied. The code is set by the programmer through the value in the EVENT-NUMBER passed to the RCVMSG and SETTIM primitives. The value must be in the range 01 through 22.

For example, a program could create MAILBOX-A, MAILBOX-B and MAILBOX-C and issue a receive on each of the mailboxes with EVENT-NUMBER set to 01, 02 and 03 respectively. If a message was sent by another program to MAILBOX-B, the WAITnn primitive would return 02. Thus, the programmer can use the EVENT-NUMBER as a condition data item to determine what program sent the message.

Its description in COBOL is

01 event-number PIC 99.

The event number is also used by the WAITnn primitive to determine the order in which the outstanding requests are to be tested. The request with the lowest event number is checked first. The request with the next lowest event number is checked second, and so forth.

3.5.3 Data Description for the Event Block

The contents of the event block are system dependent except for the first two fields. The first field is the event type with a COBOL description of PIC 99. The second field is the event outstanding flag with a COBOL description of PIC 9. Both fields are described in more detail below.

- Event-type

The possible values for event-types are

01 IPC Receive
02 Timer Runout
03 LAN Receive (Terminal input complete)

- Event-outstanding

The possible values of event-outstanding codes are

Zeros	No request outstanding
Character One	Request outstanding, wait not complete
Character Two	Request outstanding, wait complete. This state indicates that the request for a receive on either the LAN or a mailbox has been satisfied but that the user has not performed a get. The user is not required to perform a get after a wait or receive.

DS 620143000
1 November 1985

3.6 Adaptation Requirements

The COMM must be compiled, linked, and installed on each host (VAX/VMS, HL 6/GCOS MOD400, IBM/MVS) in the IISS test bed. These are located at the General Electric facility in Albany, New York. If a new host type is added to the configuration, the host specific primitives must be reimplemented for the new host and relinked with the recompiled nonhost-specific software.

SECTION 4

QUALITY ASSURANCE PROVISIONS

4.1 Introduction and Definition

"Testing" is a systematic process that may be preplanned and explicitly stated. Test techniques and procedures may be defined in advance and a sequence of test steps may be specified. "Debugging" is the process of isolation and correction of the cause of an error.

"Antibugging" is defined as the philosophy of writing programs in such a way as to make bugs less likely to occur and when they do occur, to make them more noticeable to the programmer and the users. That is, do as much error checking as is practical and possible in each routine. This approach will be followed in the COMM.

The quality assurance provisions for test will consist of the normal testing techniques that are accomplished during the construction process. They consist of design and code walk-throughs, unit testing, and integration testing. These tests will be performed by the design team. Structured design, design walk-throughs and the incorporation of "antibugging" facilitate this testing by exposing and addressing problem areas before they become coded "bugs." A detailed description of the unit tests for COMM is given in the Unit Test Plan for the Communication Subsystem, UTP620143000.

The integration testing will entail use of a test NTM implementation on each host. This test program will send messages to COMM, read messages from COMM, and print out results. COMM can also read messages from a terminal and output messages to a terminal in the absence of a LAN. A simple cable can also be used to connect the VAX and L6 instead of the LAN.

DS 620143000
1 November 1985

SECTION 5

PREPARATION FOR DELIVERY

The implementation site for the constructed software will be the ICAM Integrated Supost System (IISS) Test Bed site located at General Electric in Albany, NY. The software associated with each COMM CPCI release will be delivered on a media which is compatible with the IISS Test Bed. The release will be clearly identified and will include instructions on procedures to be followed for installation of the release.