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SOLUTION SETS FOR GAMES ON THE SQUARE
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Summary: Some necessary and sufficient* conditions that a pair of non-void weak closed convex sets of strategies form the solution set of a game with continuous payoff on the square are given.

SOLUTION SETS FOR GAMES ON THE SQUARE

I. Glilksberg and O. Gross

Let K denote the set of all optimal strategies for one player, L the corresponding set for his opponent in a game. We shall refer to $K \times L$, the set of all pairs (f, g) , $f \in K$, $g \in L$, as the solution set of the game. Any non-void weak* closed convex set K is the set of all optimal strategies for one player in some game with continuous payoff, as was shown in [1], but of course not all pairs K, L of such sets will yield solution sets. By means of constructions similar to those used in [1] we shall determine which pairs do occur in terms of the spectra,¹⁾ $\sigma K, \sigma L$ of these sets and the number of independent containing hyperplanes.

1. Preliminaries. As was shown in [1], any non-void weak* (w^*) closed convex set K of strategies is the intersection

1) $\sigma K = \bigcup_{f \in K} \sigma(f)$, which is easily seen to be a closed set.

of a sequence of half spaces, which we may express by

$$(1) \quad K = \left\{ f \mid (\varphi_n, f) = \int \varphi_n(x) df(x) \geq 0, n = 1, \dots \right\}$$

where $\{\varphi_n\}$ is a sequence of continuous functions and we may assume, for each n , $(\varphi_n, f) = 0$ for some f in K . Certain of these φ 's will yield $(\varphi, f) = 0$ for all f in K , and these we shall denote by p 's. Thus we shall write

$$K = S(\varphi_m; p_n)^{2)}$$

to express the fact that $K = \{f \mid (\varphi_m, f) \geq 0 = (p_n, f)\}$ as well as the fact that (φ_m, K) is a non-degenerate interval. The functions p_m thus define hyperplanes containing K while the φ_m do not. If the set K is the intersection of a set of hyperplanes, one may show exactly as in the proof of (1) that it is the intersection of a sequence of these and one may write $K = S(p_m)$.

What we shall be concerned with in large part in the following constructions will be the hyperplanes containing K . It is immediately evident that if we select from the functions $\{p_n\}$ a maximal subsequence $\{p'_n\}$ which is linearly independent on σK then the relations $(p_n, f) = 0$ are consequences of the relations $(p'_n, f) = 0$ for f for which

2) For the opponent we shall write $L = S(\psi_m; q_n)$ where we take $(\psi_m, g) \leq 0$.

$\sigma(f) \subset \sigma K$. Consequently if we set $p^*(x) = \text{dist}(x, \sigma K)$, then $(p^*, f) = 0$, $(p'_n, f) = 0$ all $n \iff (p_n, f) = 0$ all n ; thus in most of what follows we shall assume the $\{p_n\}$ to be linearly independent,³⁾ and actually orthonormal:

Suppose we define a measure on σK in the following way: select a sequence $\{x_n\}$ dense in σK and place weight 2^{-n} at x_n . Then clearly we may apply the Gram-Schmidt process to the $\{p_n\}$ to obtain an orthonormal sequence $\{p'_n\}$ of the same length (we take $\{x_n\}$ dense to insure that only the function 0 has the integral of its square zero). Just as clear is the fact that $(p_n, f) = 0$ for all n is equivalent to $(p'_n, f) = 0$ for all n .

2. Constructions. We shall now construct payoffs which will have three types of solution sets. That these are the only types which occur will be shown later.

Case I: Suppose $\sigma K = [0, 1] = \sigma L$ and K and L are the intersections of the same number of independent hyperplanes. The orthonormal sequences $\{p_n\}$ and $\{q_n\}$ defining K and L are thus of the same length, and if we set

$$M(x, y) = \sum a_n p_n(x) q_n(y),$$

3) We shall say that the hyperplanes H_n defined by $H_n = \{f \mid (p_n, f) = 0\}$ are independent hyperplanes containing K if the p_n are linearly independent on σK , $K \subset H_n$.

where the a_n are chosen to insure uniform convergence of the series, then for f in K and g in L ,

$$\int Mdf = \sum a_n (p_n, f) q_n(y) = 0 = \sum a_n p_n(x) (q_n, g) = \int Mdg;$$

on the other hand, if f is optimal

$$\int Mdf = \sum a_n (p_n, f) q_n(y) = 0,$$

and in view of the orthogonality of the q_n , f is in K .

Similarly every optimal g is in L , and $K \times L$ is the solution set.

Case II: Suppose $\sigma K = [0, 1] \neq \sigma L$, and $K = S(\varphi_m; p_n)$, $L = S(q_n)$ where there are at least as many independent hyperplanes containing K as there are containing L (thus we may assume that a maximal linearly independent set of p_n 's is at least as long as the set of q_n 's linearly independent on σL). Since σL is not the full unit interval we may select an open interval I which has one end point y_0 in σL . Select a disjoint sequence $\{I_n\}$ of open subintervals of I for which $\text{dist}(y_0, I_n) \rightarrow 0$, and an open subinterval I_n^* of each I_n whose closure lies entirely in I_n . Let k_n be a continuous non-negative function which vanishes outside I_n but is non-zero inside I_n , and which assumes the value 1 at a point y_n of I_n^* . Define a continuous function m_n which vanishes at y_n and outside I_n^* , but takes on the values ± 1 .

If we then set $q(y) = \text{dist}(y, \sigma L \cup I_n^*)$, then for every y not in σL one of the non-negative functions q, k_n is non-zero at y .

We now define our payoff as follows: we divide the sequence $\{p_n\}$ into $\{p_n\}$, orthonormal and of the same length as the $\{q_n\}$, and $\{p'_n\}$. If either of the sequences $\{p'_n\}$ or $\{\varphi_n\}$ are finite we use repetitions to form a sequence, and if there are no φ_n 's say, we take $\varphi_n = 1$ for all n . We set (for $b_n > 0$, chosen to insure uniform convergence)

$$M(x, y) = \sum a_n p_n(x) q_n(y) + \sum b_n [k_n(y) \varphi_{N_n}(x) + n m_n(y) p'_{N_n}(x)] + q(y)$$

where $\{N_n\}$ is an enumeration of the integers in which each integer occurs infinitely often. For f in K and g in L

$$\int Mdf = \sum b_n k_n(y) (\varphi_{N_n}, f) + q(y) \geq 0 = \int Mdg,$$

so that both are optimal.

Suppose f is optimal; then for y in σL ,

$$\sum a_n (p_n, f) q_n(y) = 0$$

whence $(p_n, f) = 0$, and thus

$$0 \leq \sum b_n [k_n(y) (\varphi_{N_n}, f) + n m_n(y) (p'_{N_n}, f)] + q(y),$$

and at setting $y = y_n$, $b_n (\varphi_{N_n}, f) \geq 0$, so that $(\varphi_n, f) \geq 0$ for all n . For y in I_n^* we have

$$0 \leq b_n [k_n(y) (\varphi_{N_n}, f) + n m_n(y) (p'_{N_n}, f)]$$

whence $0 \leq (\varphi_{N_n}, f) + nm_n(y)(p'_{N_n}, f)$, and since m_n assumes the values ± 1 ,

$$(\varphi_{N_n}, f) \geq \pm n(p'_{N_n}, f),$$

hence

$$(\varphi_{N_n}, f) \geq n|(p'_{N_n}, f)|.$$

Since N_n takes on the value n_0 infinitely often, $(\varphi_{n_0}, f) \geq n|(p'_{n_0}, f)|$ for arbitrarily large n , and $(p'_{n_0}, f) = 0$ for each n_0 . Thus f is in K .

If g is optimal, then for any f in K ,

$$0 = \iint Mdfdg = \sum b_n(k_n, g)(\varphi_{N_n}, f) + (q, g).$$

But each term of this sum is non-negative ($(k_n, g) \geq 0$ since $k_n \geq 0$) so that surely $(q, g) = 0$. If $(k_n, g) > 0$ for some n then since there is an f in K for which $(\varphi_n, f) > 0$, we would have a contradiction. Thus

$$(q, g) = 0, (k_n, g) = 0, \text{ and } (m_n, g) = 0$$

since $(k_n, g) = 0$ implies g places no weight on I_n . Thus $\sigma(g) \subset \sigma L$, and since we now may write

$$0 = \sum a_n p_n(x)(q_n, g), \quad x \text{ in } \sigma K,$$

and $(q_n, g) = 0$, g is in L .

Case III: $\sigma K \neq [0,1] \neq \sigma L$. Here we may take any K and L without further restriction, so that $K = S(\psi_m; p_n)$ and $L = S(\psi_m; q_n)$ ($(\psi_m, g) \leq 0$ here, however, in our definitions). We construct functions h_n similar to the k_n of case II, and l_n similar to the m_n , on an interval abutting σK . We set

$$M(x,y) = \sum a_n [h_n(x)\psi_{N_n}(y) + n l_n(x)q_{N_n}(y) + k_n(y)\psi_{N_n}(x) + n m_n(y)p_{N_n}(x)] .$$

Arguments entirely similar to those used in case II show $K \times L$ to be the solution set.

3. Generality. In case I ($\sigma K = [0,1] = \sigma L$) we restricted our attention to the case in which K and L were intersections of the same number of independent hyperplanes. Suppose now that a game with payoff M has as its solution set $K \times L$ where σK and σL are the full intervals. K is determined as the set of all f for which

$$\int M(x,y)df(x) = 0$$

(for convenience we take the value to be zero), and thus is the intersection of hyperplanes given by the functions $\{M(\cdot, y)\}$, and similarly L is the intersection of the hyperplanes determined by the functions $\{M(x, \cdot)\}$.

If a maximal linearly independent set $\{M(x_i, \cdot)\}$ of the first set, say, is finite, $i = 1, \dots, n$, then the same is true of the second, indeed there are just as many. For, as is

well known, n functions F_1, \dots, F_n are linearly independent on a set X if and only if there exist x_1, \dots, x_n in X for which

$$\det (F_i(x_j)) \neq 0;$$

consequently we have y_1, \dots, y_n for which

$$(2) \quad \det (M(x_i, y_j)) \neq 0,$$

so that the functions $\{M(\cdot, y_j)\}_{j=1, \dots, n}$ are linearly independent.

Of course if $\{M(\cdot, y_j)\}_{j=1, \dots, n+1}$ were linearly independent by the same argument we should have an x_{n+1} for which

$\{M(x_i, \cdot)\}_{i=1, \dots, n+1}$ were, which contradicts our assumption,

and there are exactly n . Thus the type of solution sets considered in case I are the only type which can occur. (One might note that here finite set of independent containing hyperplanes can only occur in a polynomial-like game, since for every x we have coefficients $a_i(x)$ for which

$$M(x, y) = \sum a_i(x) M(x_i, y),$$

and (2) shows the functions a_i to be continuous.)

In case II, ($\sigma K = [0, 1] \neq \sigma L$) we considered only those K and L for which we had as many independent hyperplanes containing K as there are containing L . But if M is the payoff of a game with solution set $K \times L$, $\sigma K = [0, 1] \neq \sigma L$, then as before since L is determined by

$$\int M(x, y) dg(y) = 0, \quad \text{all } x,$$

L is just the intersection of hyperplanes. If there are only n independent containing hyperplanes, then, as we shall see in a moment, these must be given by the functions

$\{M(x_i, \cdot)\}_{i=1, \dots, n}$ linearly independent on σL , for some set x_1, \dots, x_n ; consequently there exist y_1, \dots, y_n in σL for which (2) holds, and $\{M(\cdot, y_j)\}_{j=1, \dots, n}$ are linearly independent.

Since $\int M(x, y) df(x) = 0$ for y in σL , these functions define n independent hyperplanes containing K .

To see that the n independent hyperplanes containing L arise from functions $M(x_i, \cdot)$ we note that for each x , $M(x, \cdot)$ defines a containing hyperplane since x is in $\sigma K = [0, 1]$. Consequently there can be only m points, $m \leq n$, x_1, \dots, x_m for which $\{M(x_i, \cdot)\}$ are linearly independent, so that clearly $L = \{g \mid (M(x_i, \cdot), g) = 0, i = 1, \dots, m\}$.

If $m < n$, we can find a function q_0 for which, denoting $M(x_i, \cdot)$ by q_i , the set q_0, \dots, q_m is linearly independent on σL and $(q_0, g) = 0$ for all g in L . But then the mapping

$$T: g \rightarrow ((q_0, g), \dots, (q_m, g)),$$

of the set S of all strategies into $m + 1$ space, takes S into a convex subset containing $(0, \dots, 0)$ (since L is non-void). But $T(S)$ intersects the line $(t, 0, \dots, 0)$ in only one point (since $(q_0, g) = 0$ for g in L)—thus $(0, 0, \dots, 0)$ is a boundary point and we have a supporting hyperplane at this point given by constants (not all zero) a_0, \dots, a_m . Thus $\sum_{i=0}^m a_i (q_i, g) \geq 0$

for all g in S , hence $\sum a_i q_i(y) \geq 0$ for y in σL . If inequality holds for any y it holds in some neighborhood, and this is, of course, of positive measure with respect to some g in L (from the definition of σL), whence $\sum a_i (q_i, g) > 0$ for some g in L - a contradiction. Thus $\sum a_i q_i = 0$ on σL , which contradicts the linear independence on σL , and we must have $m = n$.

Thus the theme of things is as follows: The necessary and sufficient condition that $K \times L$ be the solution set for a game with continuous payoff on the square (where K and L are non-void ω^* closed convex sets of strategies) is that one of the following hold:

- (a) $\sigma K = [0, 1] = \sigma L$ and K and L are the intersection of the same number (finite if and only if the game is polynomial-like) of independent containing hyperplanes
- (b) $\sigma K = [0, 1] \neq \sigma L$, L is the intersection of hyperplanes and K has as many independent containing hyperplanes as L
- (c) $\sigma K \neq [0, 1] \neq \sigma L$.

The constructions we have used can be duplicated in the case of a game with continuous payoff played on a pair of infinite compact metric spaces; the character of solution sets, however, involves slightly different conditions:

$\sigma K = [0,1] = \sigma L$ must be replaced by $\sigma K, \sigma L$ open,
 $\sigma K = [0,1] \neq \sigma L$ by σK open, σL not open,
 $\sigma K \neq [0,1] \neq \sigma L$ by σK and σL not open. In the case of
a unique optimal strategy forming K and another forming L
we are thus guaranteed a game having $K \times L$ as the solution
set, which generalizes the result of [2].

As a final remark, we note that solution sets for
symmetric games on the square (where $M(x,y) = -M(y,x)$) can be
easily described. For such games the value is always zero
and any optimal strategy for one player is optimal for his
opponent, so that a solution set is of the form $K \times K$. The
necessary and sufficient condition that $K \times K$ be the solution
set of a symmetric game is that either

- (a) $\sigma K = [0,1]$ and K is the intersection of an even
(we take ∞ as even) number of independent hyper-
planes, or
- (b) $\sigma K \neq [0,1]$.

For if $\sigma K = [0,1]$ and K is the intersection of an
even number of independent hyperplanes given by functions
 $\{p_n\}$ (which we may take orthonormal), then, dividing these
into two sets $\{p_n\}, \{p'_n\}$ of equal cardinality, we may set

$$M(x,y) = \sum a_n [p_n(x)p'_n(y) - p'_n(x)p_n(y)],$$

which is easily seen to have $K \times K$ as its solution, and is
symmetric. On the other hand, if $K \times K$ is the solution set

of a game with payoff M and $\sigma K = [0,1]$, then K is, of course, the intersection of a set of hyperplanes. If only a finite number of these are independent, then, as before, M is polynomial-like, that is,

$$M(x,y) = \sum_{n=1}^k \varphi_n(x) \psi_n(y),$$

where $\{\varphi_n\}$ and $\{\psi_n\}$ are linearly independent sets of functions. Since M is symmetric

$$M(x,y) = -M(y,x) = -\sum_{n=1}^k \varphi_n(y) \psi_n(x),$$

so

$$M(x,y) = \frac{1}{2} \sum_{n=1}^k [\varphi_n(x) \psi_n(y) - \varphi_n(y) \psi_n(x)].$$

If the functions $\{\varphi_n, \psi_n\}$ are not a linearly independent set, replacement of a dependent φ or ψ again yields a sum of the same type, and we finally obtain a similar expression for M in which the set $\{\varphi_n, \psi_n\}$ is linearly independent; however, there are an even number of terms in the resulting sums, and thus there must be an even number of independent hyperplanes determining K .

In case (b), $K = S(\varphi_m; p_n)$, and we may set

$$M(x,y) = \sum a_n [k_n(y)\varphi_{N_n}(x) + nm_n(y)p_{N_n}(x) - k_n(x)\varphi_{N_n}(y) - nm_n(x)p_{N_n}(y)]$$

to obtain a symmetric game in which $K \times K$ is the solution set.

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2. I. Glicksberg and O. Gross, Continuous Games with Given Unique Solutions, RM-620.