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# **Empirically Based Analysis: The DDoS Case**

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# Report Documentation Page

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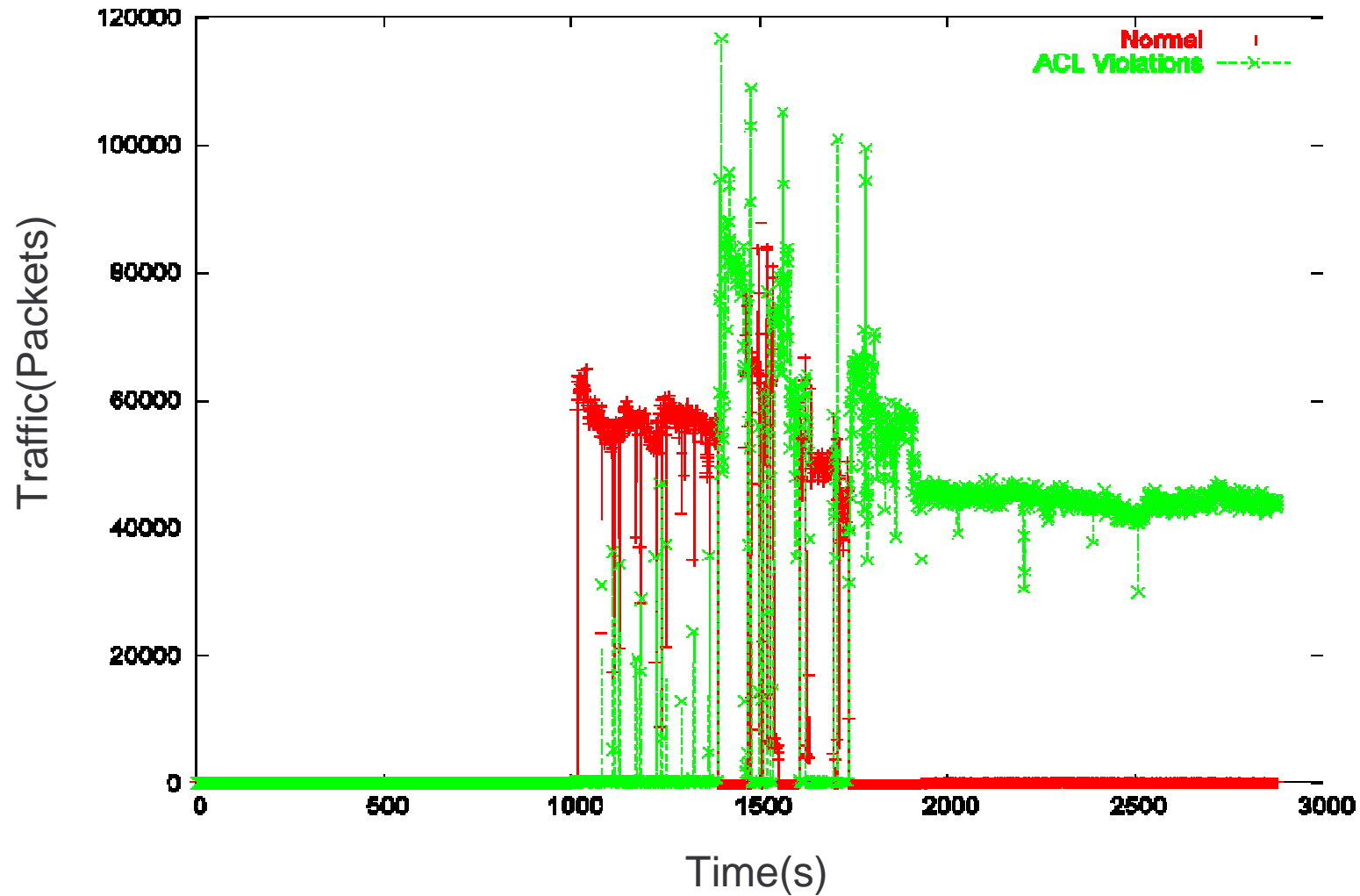


# Introduction

- Ø Access to the dataset gives us a large enough record of traffic to test hypotheses in network security.
- Ø Given this, we select and evaluate various security measures against real traffic
  - Or a reasonable facsimile thereof
- Ø One example: target resident DDoS Filters
  - Heavily constrain the problem— not considering SYN floods, smurfing, reflection attacks...



# Attacks like this





# How Do We Test?

Ø Any analysis opens a can of worms...err,  
“assumptions”

- The network constantly changes
- What is a representative host?

Ø Rerunning attacks is of debatable value

- Most of the legitimate traffic is dropped, that's what a DoS is *for*

Ø We want our results to be representative

- Test and summarize over multiple machines

Ø We want our results to be reproducible

- Depend heavily on SiLK structures and tools

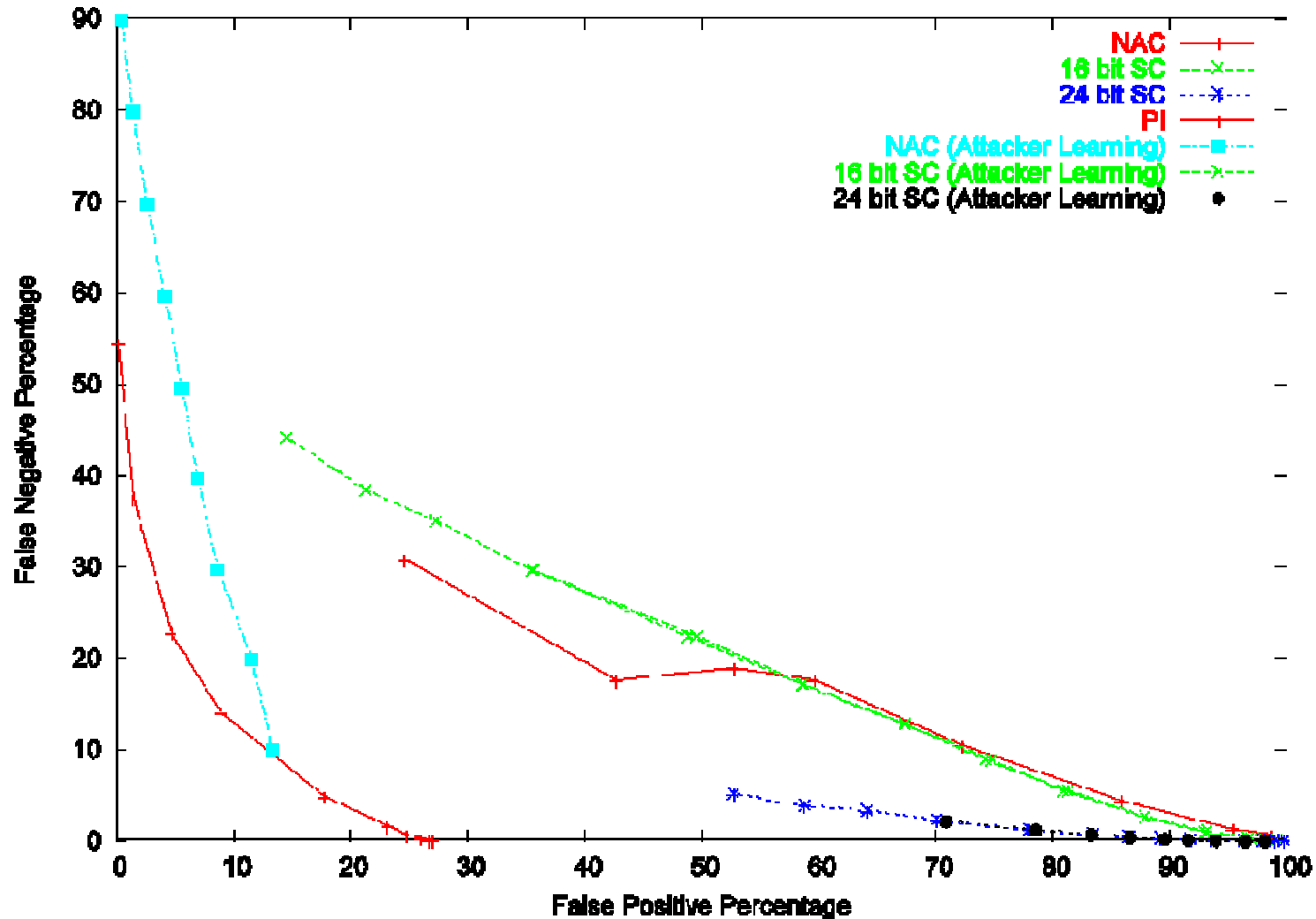


# Evaluation

- Ø Trained filters on 15 days of legitimate traffic
  - Built a representation of IP address: volume relationship (via `rwaddrcount`)
- Ø Then generated a simulated DoS
  - Botnet IPs collected with `rwset`
  - Normal traffic selected from another day
- Ø Resulting traffic was then evaluated for failure rates
- Ø Tested 2 types of filters:
  - Clustering – groups of adjacent IP addresses
  - PI – path marking approach



# DoS Filters





# Initial Observations

## Ø Two groups

- One group assumes a magic DoS Detection Oracle
  - That's the group with better results

## Ø In general, the filters don't do well

- Should we compare IP addresses, or packets?
- Is traffic different for different servers?

## Ø Let's look at one result in more depth





# Observations

## ∅ Normal traffic varies extensively

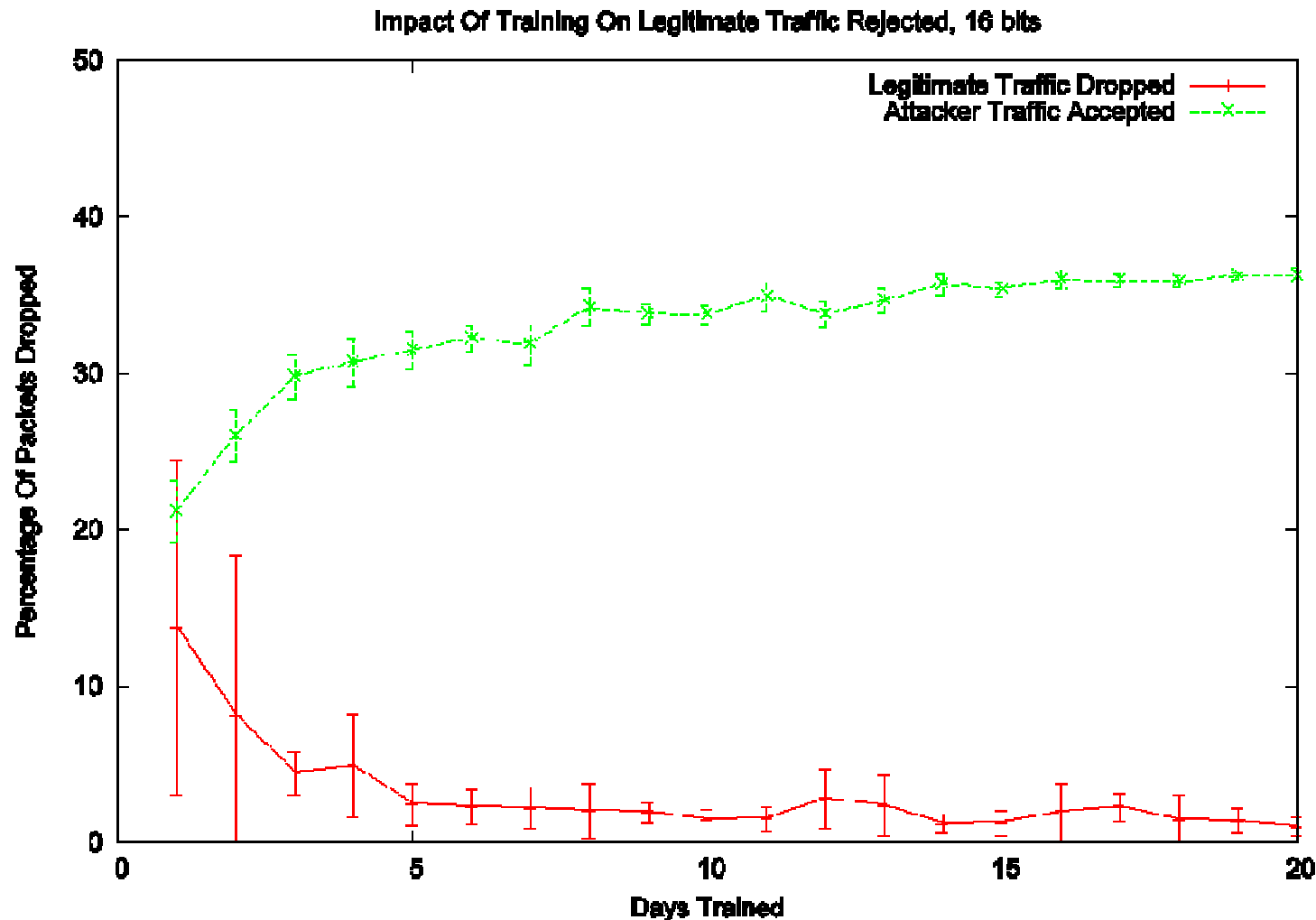
- Although it seems to vary more with “smaller” servers
- And it’s better when you look at packet counts
  - Which makes sense, given the absurd number of scanners we see.

## ∅ False negative rate (attackers accepted) seems to be related to server activity – the busier the higher.

- Attackers don’t vary as much



# Learning Curves – 95% threshold





# Other Observations

Ø In the majority of cases, packets are dropped because they've never been seen before

- Short learning curves – effectively no change in false positive rate after a week of learning.
- Especially true for spoofed traffic

Ø Entropy is lower than expected

- Filters that rely on spoof defense (HCF, PI) drop less than 10% of their packets because they detect a spoof



# Further Work

## Ø Exploiting our DoS attack traffic records further

- We know how the network reacts
- We know how the attack starts and ends
  - Which impacts learning curve for defenses that *only* profile the attack

## Ø Further use of other network maps

- Skitter (used for PI), &c.

## Ø Formalization of the techniques used

- Developed a matrix based approach for the final iteration
- Tools are going to be available publicly



# A Final Note

ØURL for the SiLK tools:

<http://silktools.sourceforge.net>